

# Formal Languages — CSE 4083 & CSE 5210

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March 5, 2022 (11:59 P.M.)

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| Answer the questions in the space provided. |
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Your Name, Class, & Major: \_\_\_\_\_

You are not to consult other people when answering these questions.

You may consult sources (books, web pages, journals, *etc*, but cite your sources and don't simply copy and paste answers.

# 1 Deterministic Finite Automata

1. (10 points) Consider a DFA  $M = (Q, \Sigma, \delta, s, f)$  with States  $Q = \{s, q_1, q_2, f\}$ , where  $s$  is the *start* and  $f$  is the *final* state; Alphabet  $\Sigma = \{0, 1\}$ ; and transition function  $\delta$ .

Construct a state transition table for  $\delta$  (or you can draw a state transition diagram) that recognizes regular expressions that are binary strings and multiples of 3, for example, the strings

0, 11, 110, 1001, 1100, ...

would be accepted strings, but

1, 10, 100, 101, ...

would not be accepted.

(Hint: Think, if  $n = 3k$  is a multiple of 3, then the next multiple of 3 is  $3k + 3$ . this could be accomplished by a transition from the current state to a next state by scanning 3 ones.)

## 2 Nondeterministic Finite Automata

2. (15 points) Explain how any NFA (with  $\lambda$  (or  $\epsilon$ ) transitions) can be converted into a DFA that accepts the same language as that accepted by the NFA. That is, the expressive power of NFAs and DFAs are equivalent. This is known as the *Rabin-Scott Theorem*.

### 3 Regular Expressions in the Programming World

3. (10 points) Consider a programming language that has identifiers that start with a lowercase ASCII letter

$$A = \{a...z\}$$

followed by a string of 1 or more digits

$$D = \{0...9\}$$

or 1 or more lowercase ASCII letters. Show how to write this specification as a regular expression.

## 4 Closure Properties of Languages

Answer these **True** (T) or **False** (F) questions. Give a brief explanation of your answer (for example, explain how to construct a machine that implements the property.)

4. (5 points) \_\_\_\_ Regular languages are closed under *intersection*.

5. (5 points) \_\_\_\_ Regular languages are closed under *Kleene-star*.

## 5 Decision Properties of Languages

6. (5 points) What does it mean to say that a “yes” or “no” question is *undecidable*?

7. (5 points) \_\_\_\_ Answer **True** (T) or **False** (F): It is decidable whether or not the language of a DFA is empty or non-empty. Give an explanation of your answer.

8. (5 points) \_\_\_\_ It is decidable whether or not the language of a DFA is finite or infinite.

9. (5 points) \_\_\_\_ It is undecidable whether or not the a string  $s$  is accepted by a DFA.  
Answer **True** (T) or **False** (F): Give an explanation of your answer.

10. (5 points) \_\_\_\_ Answer **True** (T) or **False** (F): It is decidable whether or not two regular languages  $L_1$  and  $L_2$  are equal.

Give an explanation of your answer.

## 6 Equivalence Relations

11. (10 points) On the set  $\mathbb{N}$  of natural numbers define an equivalence relation  $n \equiv m$  if and only if

$$n \bmod 3 = m \bmod 3$$

(Hint: Recall any natural number  $n$  can be written as  $n = 3q + r$  with quotient  $q$  and remainder  $r$ . And  $n \bmod 3 = \{kr : k \in \mathbb{N}\}$  The set of all natural numbers that have a remainder of  $r$  when divided by 3.

The slick way of saying this is:  $n \equiv m$  if and only if they both have the same remainder when divided by 3.

Prove that  $\equiv$  is an equivalence relation on the set of natural numbers.



## 7 The Pumping Lemma for Regular Languages

12. (10 points) DFAs can't count to an arbitrary natural number! Use the pumping lemma for regular languages to show that language

$$\text{EQ} = \{w \in \{a, b\}^* : w = a^i b^i\}$$

is not regular. Here the number of  $a$ 's in the prefix of  $w$  equals the number of  $b$ 's in the suffix of  $w$ .

## 8 Context Free Languages

13. (5 points) Consider the CFG  $G$  defined by the productions:

$$S \rightarrow aS|Sb|a|b$$

Prove by induction that no string in  $L(G)$  has  $ba$  as a sub-string.

Hint: To show this do induction on the length of the strings.

14. (5 points) Give simple English language descriptions for the strings generated by the productions following four grammars ( $G = (V, T, P, S)$ ):

$$1. G1 \rightarrow S|aS|a$$

$$2. G2 : S!aSa|aa|a$$

$$3. G3 : S|SaS|a$$

| Question | Points | Score |
|----------|--------|-------|
| 1        | 10     |       |
| 2        | 15     |       |
| 3        | 10     |       |
| 4        | 5      |       |
| 5        | 5      |       |
| 6        | 5      |       |
| 7        | 5      |       |
| 8        | 5      |       |
| 9        | 5      |       |
| 10       | 5      |       |
| 11       | 10     |       |
| 12       | 10     |       |
| 13       | 5      |       |
| 14       | 5      |       |
| Total:   | 100    |       |