Functional Programming Excercise Sheet 3

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Excercise 1

a)

```
collatz :: Int -> [Int]
collatz x = iterate (\y -> if mod y 2 == 0 then div y 2 else 3*y+1) x

total_stopping_time :: Int -> Int
total_stopping_time 1 = 3
total_stopping_time x = length (takeWhile (/= 1) (collatz x))
```

b)

```
check_collatz :: Int -> Bool
check_collatz 1 = False
check_collatz n = elem 1 (take n (collatz n))
```

Excercise 2

a)

b)

```
range :: [a] \rightarrow Int \rightarrow Int \rightarrow [a] range xs m n = [ x | (i, x) <- zip [0 ..] xs, i >= m, i <= n b]
```

Excercise 3

a)

```
main :: IO ()
main = do
    -- task a)
    putStrLn "Welcome to your library"
    library []
    putStrLn "Bye!"
    return ()
```

b)

```
getInput :: IO LibraryInput
getInput = do
   -- task b)
putStrLn "Would you like to put back or take a book?\n Enter Book:
    Title's name; Author's name \nAre you looking for an author?\n
    Enter Author: Author's name \nAre you looking for a special book
    ?\n Enter Title: Title's name."
putChar '>'
input <- getLine
return (parseLibraryInput input)
return Exit</pre>
```

c)

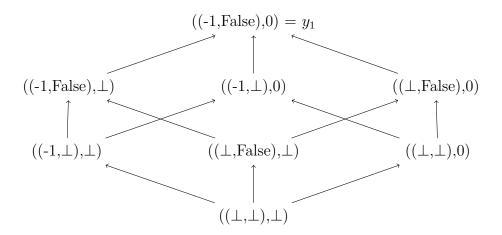
```
library ::[(String,String)] -> IO ()
library books = do
   -- task c)
input <- getInput
getLibraryAction books input
return ()</pre>
```

```
getLibraryAction :: [(String, String)] -> LibraryInput -> IO()
getLibraryAction books Exit = do return ()
getLibraryAction books (Error e) = do putStrLn (show (Error e))
                library books
getLibraryAction books (Book (t,a)) = do putStrLn (show (Book (t,a)))
                putStrLn "Do you want to (p)ut the book back or do you
                   want to (t)ake the book?"
                input <- getLine</pre>
                evaluateAction input books (t,a)
getLibraryAction books (Author a) = do putStrLn (show (Author a))
                putStrLn ("You have the following books from " ++ a)
                displayBooks_A a books
                library books
getLibraryAction books (Title t) = do putStrLn (show (Title t))
                putStrLn ("You have the following books with the title:
                    " ++ t)
                displayBooks_T t books
                library books
```

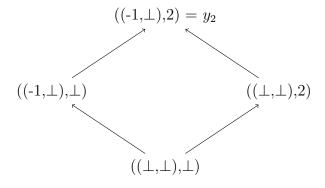
d) No, it is not possible to write a function main :: Int that behaves similar to the function main :: IO () of the previous excercises. The number of books in the library is a value, which is encapsulated in an IO action. This value cannot be taken out of this action, since that would be against referential transparency. Furthermore, having the function main be of type Int would also consequently change the behaviour of the function as IO actions such as putStrLn :: String -> IO () or let alone >> :: IO () -> IO () -> IO() would cause a type matching error.

Excercise 4

a)



All elements $x \in ((\mathbb{Z}_{\perp} \times \mathbb{B}_{\perp}) \times \mathbb{Z}_{\perp}$ that are less defined than y_1 are given by $((-1, \perp), 0), ((\perp, False), 0), ((-1, False), \perp), ((\perp, False), \perp), ((-1, \perp), \perp), ((\perp, \perp), 0)$ and the smallest Element $((\perp, \perp), \perp)$.



All elements $x \in ((\mathbb{Z}_{\perp} \times \mathbb{B}_{\perp}) \times \mathbb{Z}_{\perp}$ that are less defined than y_2 are given by $((\perp, \perp), 2), ((-1, \perp), \perp)$ and the smallest Element $((\perp, \perp), \perp)$.

b) Given a domain $D = \underbrace{\mathbb{Z}_{\perp} \times ... \times \mathbb{Z}_{\perp}}_{n \text{ times}}$ for $0 < n \in \mathbb{N}$ and a chain $S \subseteq D$ then $\sup\{|S| \mid S \subseteq D, S \text{ is a chain}\} = n + 1.$

We use induction on n to show that this holds for all $n \in \mathbb{N}, n > 0$.

Base case n=1

Let $D = \mathbb{Z}_{\perp}$, then $\perp \in D$ and for all $x \in \mathbb{Z} : x \in \mathbb{Z}_{\perp}$. For $x, y \in S$ with $x, y \in \mathbb{Z}$ it follows from the definition of S that either $x \sqsubseteq y$ or $y \sqsubseteq x$. By the definition of \sqsubseteq and x and y, this is only true iff x = y. As a result, $|\{x \in \mathbb{Z} \mid x \in S \text{ with } S \subseteq D, S \text{ is a chain}\}| \leq 1$.

However, $\bot \sqsubseteq x$ for all $x \in \mathbb{Z}$ hence $\sup\{|S| \mid S \subseteq D, S \text{ is a chain}\} = 1 + 1 = 2 = n + 1.$

Induction hypothesis

sup{
$$|S| \mid S \subseteq D$$
, S is a chain} = $n+1$ holds for $n \in \mathbb{N}$, $0 < n$ with $D = \underbrace{\mathbb{Z}_{\perp} \times ... \times \mathbb{Z}_{\perp}}_{n \text{ times}}$.

Induction step
$$n \to n+1$$

Let
$$D = \underbrace{\mathbb{Z}_{\perp} \times ... \times \mathbb{Z}_{\perp}}_{n+1 \text{ times}}$$
.

It follows from the hypothesis that if $D_n = \underbrace{\mathbb{Z}_{\perp} \times ... \times \mathbb{Z}_{\perp}}_{n \text{ times}}$ then $\sup\{|S_n| \mid S_n \subseteq$

 D_n, S_n is a chain $\} = n + 1$.

Now we extend each n-tuple $(n_1, n_2, ..., n_n)$ in the chain S_n $(n_1, n_2, ..., n_n \in \mathbb{Z}_{\perp})$ so that we get (n+1)-tuples $(n_1, n_2, ..., n_n, x) \in D, x \in \mathbb{Z}_{\perp}$.

Case 1: $x \in \mathbb{Z}$

For two tuples $t_i = (n_{i_1}, n_{i_2}, ..., n_{i_n}), t_j = (n_{j_1}, n_{j_2}, ..., n_{j_n}) \in S_n$ either $t_i \sqsubseteq t_j$ or $t_j \sqsubseteq t_i$ applies. Then also $(n_{i_1}, n_{i_2}, ..., n_{i_n}, x) \sqsubseteq (n_{j_1}, n_{j_2}, ..., n_{j_n}, x)$ or $(n_{j_1}, n_{j_2}, ..., n_{j_n}, x) \sqsubseteq (n_{i_1}, n_{i_2}, ..., n_{i_n}, x)$ since in particular $x \sqsubseteq x$ and consequently $(n_{j_1}, n_{j_2}, ..., n_{j_n}, x), (n_{i_1}, n_{i_2}, ..., n_{i_n}, x) \in S$ with $\sup\{|S| \mid S \subseteq D, S \text{ is a chain}\} \ge \sup\{|S_n| \mid S_n \subseteq D_n, S_n \text{ is a chain}\} = n + 1.$

If $(n_1, n_2, ..., n_n, x) \in S$ and $(n_1, n_2, ..., n_n, y) \in S$ then x = y because neither $x \sqsubseteq y$ nor $y \sqsubseteq x$ applies. Similarly, If $t_i = (n_1, n_2, ..., n_n, x) \in S$ and $t_j = (n_1, n_2, ..., x, ..., n_n) \in S$ then neither $t_i \sqsubseteq t_j$ nor $t_j \sqsubseteq t_i$ applies, because $t_i \neq t_j$ and both are equally defined. This leads on to our second case.

Case 2:
$$(n_1, n_2, ..., n_n, x), n_1 = n_2 = ... = n_n = x = \bot$$

Then clearly $(n_1, n_2, ..., n_n, x) \in S$ since it is the smallest element of D and therefore $(n_1, n_2, ..., n_n, x) \sqsubseteq y$ for all $y \in D$. Resulting from this, $\sup\{|S| \mid S \subseteq D, S \text{ is a chain}\} = (n+1)+1=n+2$. The equation holds.