

ELEC 405 - Homework 2

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1. Nearest Neighbor partitioning and halfspaces

$$\{x | \|x - a\| \leq \|x - b\|\} = \{x | c^T x \leq d\}, \quad c \in \mathcal{R}^n, d \in \mathcal{R}^n \quad (1)$$

Take the square of 1, as $\|\cdot\| \geq 0$, then

$$\begin{aligned} \|x - a\|^2 &\leq \|x - b\|^2 \\ \langle x - a, x - a \rangle &\leq \langle x - b, x - b \rangle \\ \|x\|^2 - 2a^T x + \|a\|^2 &\leq \|x\|^2 - 2b^T x + \|b\|^2 \\ 2(b - a)^T x &\leq \|b\|^2 - \|a\|^2 \end{aligned} \quad (2)$$

Then $c = 2(b - a)$ and $d = \|b\|^2 - \|a\|^2$.

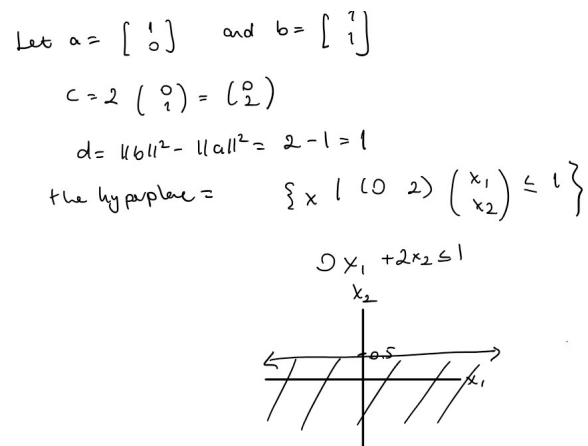


Figure 1: Hyperplane

2. Rank and matrix products

For the sake of this question let $A \in \mathcal{R}^{m \times n}$ and $B \in \mathcal{R}^{n \times p}$

1. If AB is full rank then A and B are full rank.
Not necessarily. Let $A = \begin{bmatrix} 1 & 1 \end{bmatrix}$ and $B = \begin{bmatrix} 1 & 1 \\ 0 & 0 \end{bmatrix}$. $AB = \begin{bmatrix} 1 & 1 \end{bmatrix}$. AB is full rank but B is not.
2. If A and B are full rank then AB is full rank.
Not necessarily. Let $A = \begin{bmatrix} 0 & 1 \end{bmatrix}$ and $B = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$. $AB = \begin{bmatrix} 0 \end{bmatrix}$. A and B are full rank but AB is not.
3. If A and B have zero nullspace, then so does AB .
Assume $\mathcal{N}(AB) \neq \{0\}$ and let $x \in \mathcal{R}^p$ such that $ABx = 0$ then $A(Bx) = 0$ so either $Bx \in \mathcal{R}^n$ is in null space of A but null space of A is trivial. So $Bx = 0$, then $x \in \mathcal{N}(B)$ but since null space of B is also trivial, x must be 0. Hence AB has zero nullspace.
4. If A and B are onto, then so is AB .
We need to show that $\forall y \in \mathcal{R}^m \exists x \in \mathcal{R}^p$ such that $ABx = y$. Since A is onto, $\exists z \in \mathcal{R}^n$ such that $Az = y$. Similarly, since B is onto, $\exists x \in \mathcal{R}^p$ such that $Bx = z$. Hence $A(Bx) = Az = y$.

3. Fundamental Subspaces and Determinant

Let $A \in \mathcal{R}^{m \times n}$

1. If $\mathcal{R}(A) = \mathcal{R}^m$, then AA^T is full rank.
If $\mathcal{R}(A) = \mathcal{R}^m$ then A^T is full rank and its null space is trivial ($\{0\}$). Let $x \in \mathcal{R}^m$ such that $(AA^T)x = 0$, multiply each side with $x^T \Rightarrow x^T AA^T x = (A^T x)^T (A^T x) = \langle A^T x, A^T x \rangle = \|A^T x\|^2$, since A^T is full rank we know that $x = 0$ which proves that the null space of AA^T is trivial and the matrix is full rank.
2. If $\mathcal{N}(A) = \{0\}$ then $A^T A$ is full rank.
For $A^T A$ to be full rank, we need to show that $\dim(\mathcal{R}(A^T A)) = n$. In other words, if we show that the null space of $A^T A$ is trivial then we prove that it is full rank. Let $x \in \mathcal{R}^n$, such that $(A^T A)x = 0$ is we multiply both sides with x^T then we get $x^T A^T A x = (Ax)^T (Ax) = \langle Ax, Ax \rangle = \|Ax\|^2$ and we know that $\|y\| = 0$ if and only if $y = 0$ since A is full rank therefore $A^T A$ is also full rank.

4. Characteristic Polynomial of a Square Matrix

Let $A \in \mathcal{R}^{n \times n}$ and $X(s) = \det(sI - A)$ be the characteristic polynomial of A .

1. When we think about how we take the determinant, we can construct the result of the determinant as follows:
First take the first row and column into account so that A' is A without the first row and column. Then

$$\det(sI - A) = (s - a_{11})\det A' + C$$

where C is a constant. Iteratively, this will yield a similar result so that we can write $\det(sI - A)$ of the form:

$$\det(sI - A) = \prod_{i=1}^n (s - a_{ii}) + C^* \quad (3)$$

Therefore, the coefficient of s^n is 1.

2. Show that s^{n-1} coefficient of $X(s)$ is given by $-TrA$.
Note that in (3) C is a polynomial of degree at most $n-1$ because A' is a $(n-1) \times (n-1)$ matrix. Similarly, A' is the only matrix with $n-1 \times n-1$ elements, so the coefficient of s^{n-1} comes from $\Pi_{i=1}^n (s - a_{ii})$ which is $\sum_{i=1}^n -a_{ii} = TrA$.
3. Show that the constant coefficient of $X(s)$ is given by $det(-A)$.
Like in any polynomial, we can take $X(0)$ to find the constant coefficient and find $det(0I - A) = det(-A)$.
4. Consider $n = 1$, then $X(s) = det(sI - A) = s - \lambda_1$, and the relationship holds. Now assume that the relationship holds for some positive integer k such that $a_{k-1} = -\sum_{i=1}^k \lambda_i$ and $a_0 = \prod_{i=1}^k -\lambda_i$. Multiply $X(s)$ with $(s - \lambda_{k+1})$ which will give us $sX(s) - \lambda_{k+1}X(s)$. By part (1) $sX(s)$ is monic and of degree $k+1$. The coefficient of s^k is $\sum_{i=1}^k (-\lambda_i) - \lambda_{k+1}$ and the constant coefficient is $\prod_{i=1}^{k+1} -\lambda_i$, note that the constant coefficient of $sX(s)$ is 0. Therefore, the relationship holds.

5. Schur's Theorem and Cayley-Hamilton Theorem

1. By Schur's Theorem, we can construct T such that $T = U^*AU$ where U is described in the theorem. $p(t) = \prod_{i=1}^n (t - \lambda_i)$ is the characteristic polynomial of T . Since λ_i are at the diagonals of T , note that λ_i 's don't need to be unique, $P(T) = \prod_{i=1}^n (T - \lambda_i I) = 0$. Evaluate

$$\begin{aligned}
 U^*P(A)U &= U^*(A - \lambda_1 I) \dots (A - \lambda_n I)U \\
 &= U^*(A - \lambda_1 I)UU^* \dots UU^*(A - \lambda_n I)U \\
 &= (U^*AI - \lambda_1 U^*U) \dots (U^*AI - \lambda_n U^*U) \text{ note } U^*U = I \\
 &= (T - \lambda_1 I) \dots (T - \lambda_n I) \\
 &= P(T)
 \end{aligned} \tag{4}$$

Reorganizing the equality in (4) yields

$$P(A) = UP(T)U^*(A) = 0$$

2. Evaluating $p_A(A) = A^n + a_{n-1}A^{n-1} + \dots + a_1A + a_0 = 0$. Then

$$A^n = -a_{n-1}A^{n-1} - \dots - a_1A - a_0 \tag{5}$$

$$\begin{aligned}
 A^{n+1} &= A(-a_{n-1}A^{n-1} - \dots - a_1A - a_0) \\
 &= -a_{n-1}A^n - \dots - a_1A^2 - a_0A
 \end{aligned} \tag{6}$$

$$= (a_{n-1}^2 + a_{n-1}a_{n-2})A^{n-1} + \dots + (a_2a_{n-1} - a_1)A^2 + (a_1a_{n-1} - a_0)A + a_{n-1}a_0 \tag{7}$$

$$\begin{aligned}
 A^{n+2} &= A(a_{n-1}^2A^{n-1} + \dots + (a_2a_{n-1} - a_1)A^2 + (a_1a_{n-1} - a_0)A + a_{n-1}a_0) \\
 &= a_{n-1}^2A^n + \dots + (a_2a_{n-1} - a_1)A^3 + (a_1a_{n-1} - a_0)A^2 + (a_{n-1}a_0A) \\
 &= -a_{n-1}^3A^{n-1} + \dots + (-a_2a_{n-1}^2 + a_1a_{n-1} - a_0)A^2 + (-a_{n-1}^2a_1 + a_{n-1}a_0)A - a_{n-1}^2a_0
 \end{aligned} \tag{8}$$

(7) follows from plugging (5) into (6). Formally in the expansion of A^k where $k \geq n$ the coefficient of A^{n-1} is $(-a_{n-1})^{k-n+1}$ and the coefficient of A^i is $\sum_{j=1}^{i+1} (-1)^{j-1} a_{n-1}^{j-1} a_{j-1}$

6. Theory in Action for Big Data!: Eigenvectors as community detectors

1. $E(C) \in \mathcal{R}^{2N \times 2N}$ where E is of the form:

$$E(C)_{ij} = \begin{cases} p & \text{if both } i, j \in \mathcal{I}_A \text{ or } \mathcal{I}_B \\ q & \text{if } i \in \mathcal{I}_A \text{ and } j \in \mathcal{I}_B \text{ or vice versa} \\ 0, & \text{if } C_{ij} = 0 \end{cases} \tag{9}$$

2. $2 \times \text{rank } C$