# Linear programming

- Example Numpy: PageRank
- scipy.optimize.linprog
- Example linear programming: Maximum flow

# PageRank

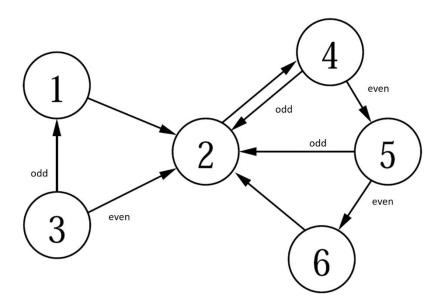
# PageRank - A NumPy / Jupyter / matplotlib example

- Google's original search engine ranked webpages using PageRank
- View the internet as a graph where nodes correspond to webpages and directed edges to links from one webpage to another webpage
- Google's PageRank algorithm was described in (ilpubs.stanford.edu:8090/361/, 1998)

### The Anatomy of a Large-Scale Hypertextual Web Search Engine

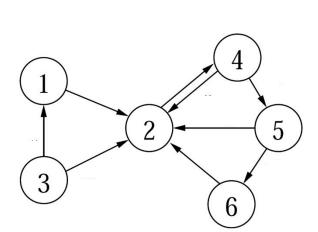
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# Five different ways to compute PageRank probabilities

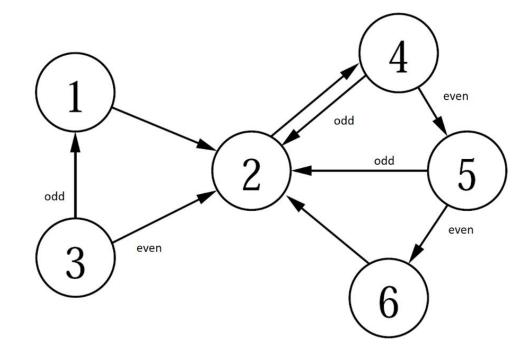
- 1) Simulate random process manually by rolling dices
- 2) Simulate random process in Python
- 3) Computing probabilities using matrix multiplication
- 4) Repeated matrix squaring
- 5) Eigenvector for  $\lambda = 1$



### Random surfer model (simplified)

The PageRank of a node (web page) is the fraction of the time one visits a node by performing an *infinite random traversal* of the graph starting at node 1, and in each step

- with probability 1/6 jumps to a random page (probability 1/6 for each node)
- with probability 5/6 follows an outgoing edge to an adjacent node (selected uniformly)



The above can be simulated by using a dice: Roll a *dice*. If it shows 6, jump to a random page by rolling the dice again to figure out which node to jump to. If the dice shows 1-5, follow an outgoing edge - if two outgoing edges roll the dice again and go to the lower number neighbor if it is odd.

### Adjacency matrix and degree vector

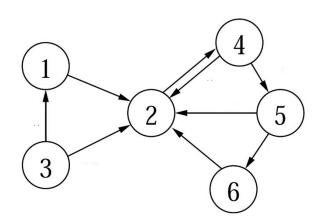
```
pagerank.ipynb
import numpy as np
# Adjacency matrix of the directed graph in the figure
# (note that the rows/column are 0-indexed, whereas in the figure the nodes are 1-indexed)
G = np.array([[0, 1, 0, 0, 0, 0],
              [0, 0, 0, 1, 0, 0],
              [1, 1, 0, 0, 0, 0],
              [0, 1, 0, 0, 1, 0],
              [0, 1, 0, 0, 0, 1],
              [0, 1, 0, 0, 0, 0]]
n = G.shape[0] # number of rows in G
degree = np.sum(G, axis=1, keepdims=True) # column vector with row sums = out-degrees
# The below code handles sinks, i.e. nodes with outdegree zero (no effect on the graph above)
G = G + (degree == 0) # add edges from sinks to all nodes
degree = np.sum(G, axis=1, keepdims=True)
```

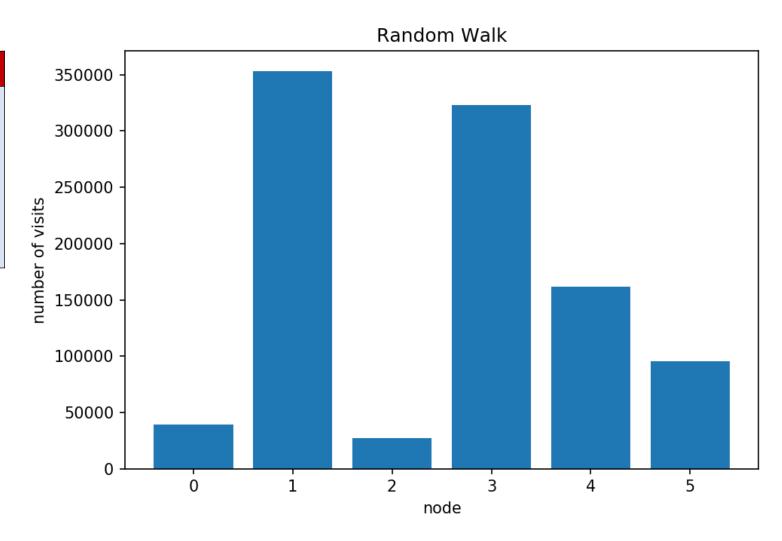
### Simulate random walk (random surfer model)

```
pagerank.ipynb
from random import randint
STEPS = 1000000
# adjacency list[i] is a list of all j where (i, j) is an edge of the graph.
adjacency list = [[j for j, e in enumerate(row) if e] for row in G]
count = np.zeros(n) # histogram over number of node visits
state = 0
                         # start at node with index 0
for in range(STEPS):
   count[state] += 1  # increment count for state
   if randint(1, 6) == 6: # original paper uses 15% instead of 1/6
       state = randint(0, 5)
   else:
       state = adjacency list[state][randint(0, degree[state] - 1)]
print(adjacency list, count / STEPS, sep='\n')
Python shell
  [[1], [3], [0, 1], [1, 4], [1, 5], [1]]
  [0.039365 0.353211 0.02751 0.322593 0.1623
                                              0.0950211
```

### Simulate random walk (random surfer model)

# pagerank.ipynb import matplotlib.pyplot as plt plt.bar(range(6), count) plt.title("Random Walk") plt.xlabel("node") plt.ylabel("number of visits") plt.show()





### Transition matrix A

```
pagerank.ipynb
A = G / degree # Normalize row sums to one. Note that 'degree'
               # is an n x 1 matrix, whereas G is an n x n matrix.
               # The elementwise division is repeated for each column of G
print(A)
Python shell
 [[0. 1. 0. 0. 0. 0.]
   [0. 0. 0. 1. 0. 0.]
   [0.5 0.5 0. 0. 0. 0.]
   [0. 0.5 0. 0. 0.5 0.]
   [0. 0.5 0. 0. 0. 0.5]
   [0. 1. 0. 0. 0. 0.]]
```

### Repeated matrix multiplication

We now want to compute the probability  $p^{(i)}_{j}$  to be in vertex j after i steps. Let  $p^{(i)} = (p^{(i)}_{0}, \dots, p^{(i)}_{n-1})$ .

Initially we have  $p^{(0)} = (1, 0, ..., 0)$ .

We compute a matrix M, such that  $p^{(i)} = M^i \cdot p^{(0)}$  (assuming  $p^{(0)}$  is a column vector).

If we let  $\mathbf{1}_n$  denote the  $n \times n$  matrix with 1 in each entry, then M can be computed as:

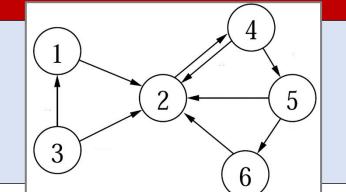
$$p_j^{(i+1)} = \frac{1}{6} \cdot \frac{1}{n} + \frac{5}{6} \sum_{k} p_k^{(i)} \cdot A_{k,j}$$

$$p^{(i+1)} = \left(\frac{1}{6} \cdot \frac{1}{n} \mathbf{1}_n + \frac{5}{6} A^{\mathsf{T}}\right) \cdot p^{(i)}$$

### pagerank.ipynb

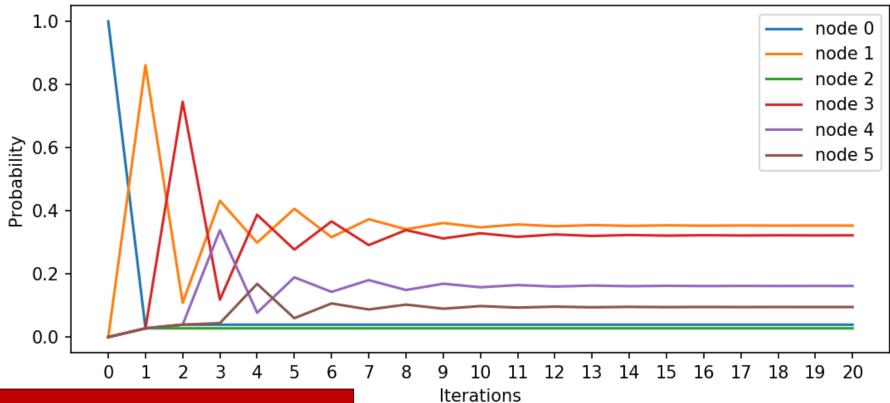
### Python shell

[[0.03935185] [0.35326184] [0.02777778] [0.32230071] [0.16198059] [0.09532722]]



### Random Surfer Probabilities

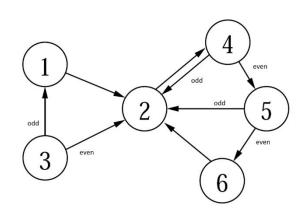
# Rate of convergence



```
pagerank.ipynb

x = range(ITERATIONS + 1)
for node in range(n):
    plt.plot(x, prob[node], label="node %s" % node)

plt.xticks(x)
plt.title("Random Surfer Probabilities")
plt.xlabel("Iterations")
plt.ylabel("Probability")
plt.legend()
plt.show()
```



### Repeated squaring

 $\mathcal{M} \cdot (\cdots (\mathcal{M} \cdot (\mathcal{M} \cdot p^{(0)})) \cdots) = \mathcal{M}^k \cdot p^{(0)} = \mathcal{M}^{2 \log k} \cdot p^{(0)} = (\cdots (\mathcal{M}^2)^2)^2 \cdots)^2 \cdot p^{(0)}$ 

log k

k multiplications, k power of 2

```
pagerank.ipynb

from math import log
MP = M

for _ in range(1 + int(log(ITERATIONS, 2))):
        MP = MP @ MP

p = MP @ p_0
print(p)

Python shell
```

### [[0.03935185]

[0.03933183] [0.35332637] [0.02777778] [0.32221711] [0.16203446]

[0.09529243]]

# PageRank: Computing eigenvector for $\lambda = 1$

• We want to find a vector p, with |p| = 1, where Mp = p, i.e. an *eigenvector* p for the eigenvalue  $\lambda = 1$ 

```
pagerank.ipynb
eigenvalues, eigenvectors = np.linalg.eig(M)
idx = eigenvalues.argmax()  # find the largest eigenvalue (= 1)
p = np.real(eigenvectors[:, idx]) # .real returns the real part of complex numbers
p /= p.sum()  # normalize p to have sum 1
print(p)
Python shell
```

0.02777778 0.32221669 0.16203473 0.095292251

[0.03935185 0.3533267

### PageRank: Note on practicality

 In practice an explicit matrix for billions of nodes is infeasible, since the number of entries would be order of 10<sup>18</sup>

 Instead use sparse matrices (in Python modul scipy.sparse) and stay with repeated multiplication

# Linear programming

### scipy.optimize.linprog

 scipy.optimize.linprog can solve linear programs of the following form, where one wants to find an  $n \times 1$  vector x satisfying:

**Subject to**:  $A_{ub} \cdot x \le b_{ub}$  $A_{eq} \cdot x = b_{eq}$ 

dimension

 $c: n \times 1$ 

 $A_{\text{ub}}: m \times n$   $b_{\text{ub}}: m \times 1$   $A_{\text{eq}}: k \times n$   $b_{\text{eq}}: k \times 1$ 

### Linear programming example

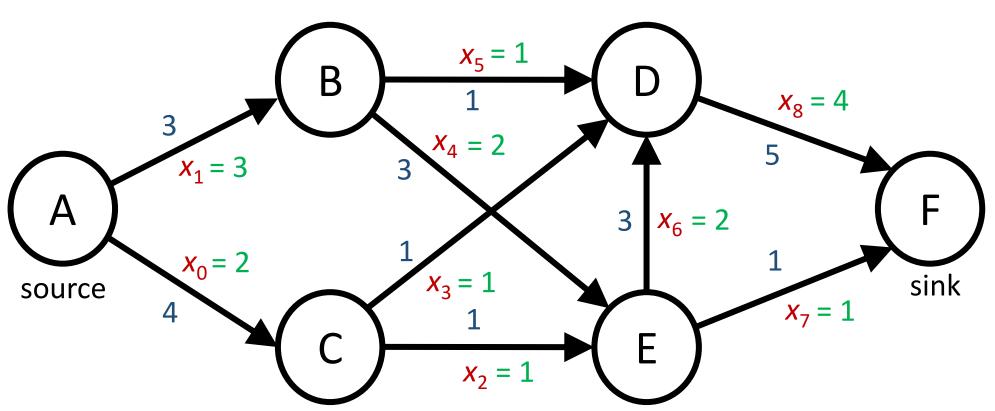
### Maximize $3x_1 + 2x_2$ $3 \cdot x_1 + 2 \cdot x_2$ $2x_1 + x_2 \le 10$ $5x_1 + 6x_2 \ge 4$ **Subject to** 10 $-3x_1 + 7x_2 = 8$ $2 \cdot x_1 + 1 \cdot x_2 \le 10$ (3.65, 2.71) $5 \cdot x_1 + 6 \cdot x_2 \ge 4$ $-3 \cdot x_1 + 7 \cdot x_2 = 8$ - 20 $X_2$ - 10 **Minimize** $-(3\cdot x_1 + 2\cdot x_2)$ **Subject to -**5 · $2 \cdot x_1 + 1 \cdot x_2 \le 10$ $-5 \cdot x_1 + -6 \cdot x_2 \le -4$ -10 $-3 \cdot x_1 + 7 \cdot x_2 = 8$ 2.5 5.0 7.5 0.0 10.0

```
linear programming.py
import numpy as np
from scipy.optimize import linprog
c = np.array([3, 2])
A ub = np.array([[2, 1],
                 [-5, -6]) # multiplied by -1
b ub = np.array([10, -4])
A eq = np.array([-3, 7])
b eq = np.array([8])
res = linprog(-c, # maximize = minimize the negated
              A ub=A ub,
              b ub=b ub,
              A eq=A eq,
              b eq=b eq)
print(res) # res.x is the optimal vector
```

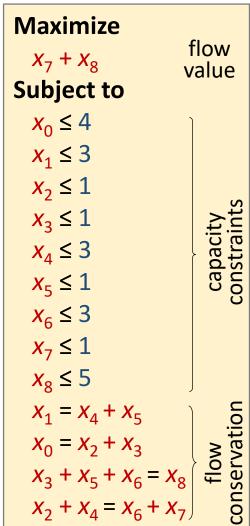
### Python shell

### Maxmium flow

# Solving maximum flow using linear programming



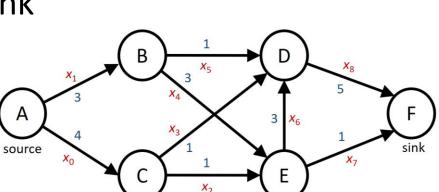
We will use the <u>scipy.optimize.linprog</u> function to solve the *maximum flow* problem on the above directed graph. We want to send as much *flow* from node A to node F. Edges are <u>numbered 0..8</u> and each edge has a maximum *capacity*.

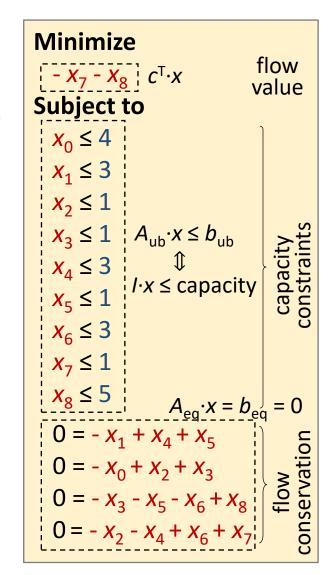


Note: solution not unique

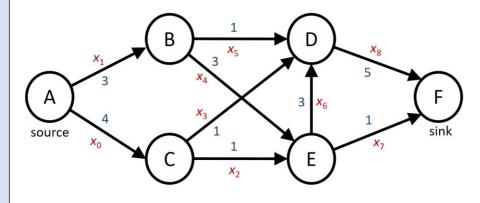
# Solving maximum flow using linear programming

- x is a vector describing the flow along each edge
- c is a vector that to add the flow along the edges (7 and 8) to the sink (F), i.e. a function computing the flow value
- $A_{\text{ub}}$  and  $b_{\text{ub}}$  is a set of *capacity constraints*, for each edge flow ≤ capacity
- $A_{eq}$  and  $b_{eq}$  is a set of flow conservation constraints, for each non-source and non-sink node (B, C, D, E), requiring that the flow into equals the flow out of a node





```
maximum-flow.py
import numpy as np
from scipy.optimize import linprog
#
conservation = np.array([[0,-1, 0, 0, 1, 1, 0, 0, 0], #B])
                        [-1, 0, 1, 1, 0, 0, 0, 0, 0], \# C
                        [0, 0, 0, -1, 0, -1, -1, 0, 1], # D
                        [0, 0, -1, 0, -1, 0, 1, 1, 0]]) \# E
                 0 1 2 3 4 5 6 7 8
sinks = np.array([0, 0, 0, 0, 0, 0, 1, 1])
#
                    0 1 2 3 4 5 6 7 8
capacity = np.array([4, 3, 1, 1, 3, 1, 3, 1, 5])
res = linprog(-sinks,
             A eq=conservation,
             b eq=np.zeros(conservation.shape[0]),
             A ub=np.eye(capacity.size),
             b ub=capacity)
print(res)
```



#### fun: -5.0 message: 'Optimization terminated successfully.' nit: 9 slack: array([2., 0., 0., 0., 1., 0., 1., 0., 1.]) status: 0 success: True $\rightarrow$ x: array([2., 3., 1., 1., 2., 1., 2., 1., 4.])

Python shell

the solution found varies with the scipy version