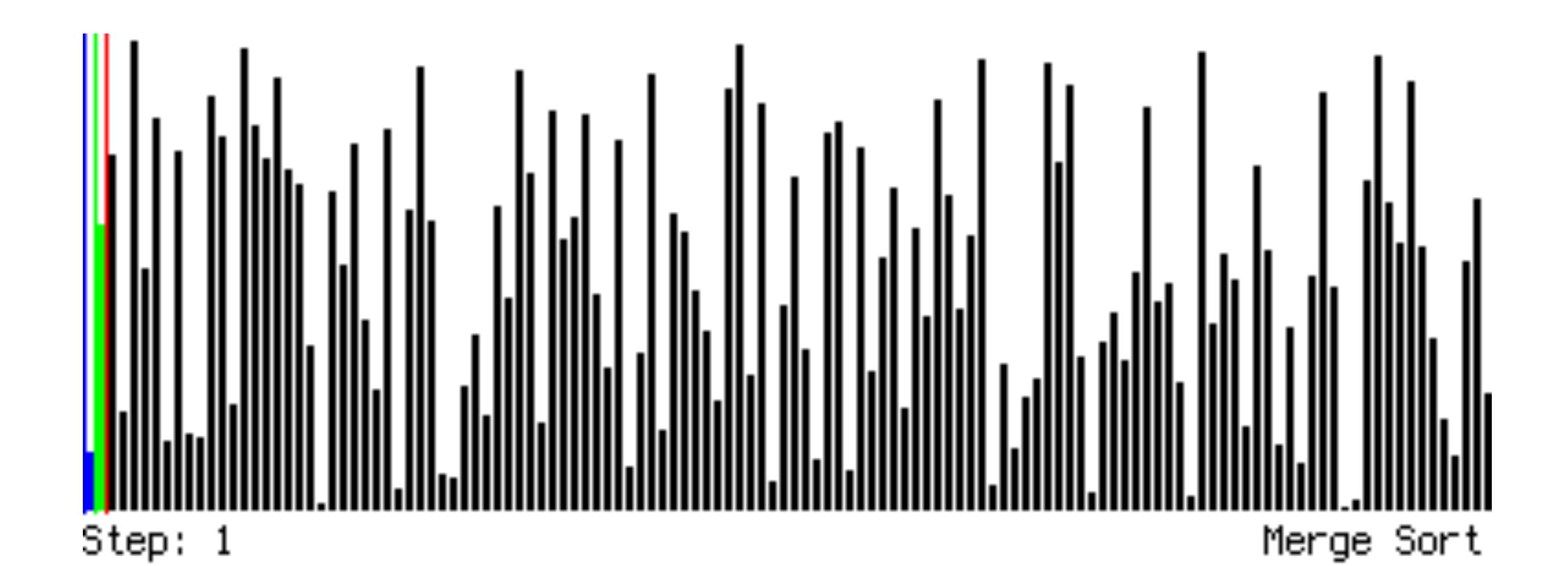
Parallel Programming DS295

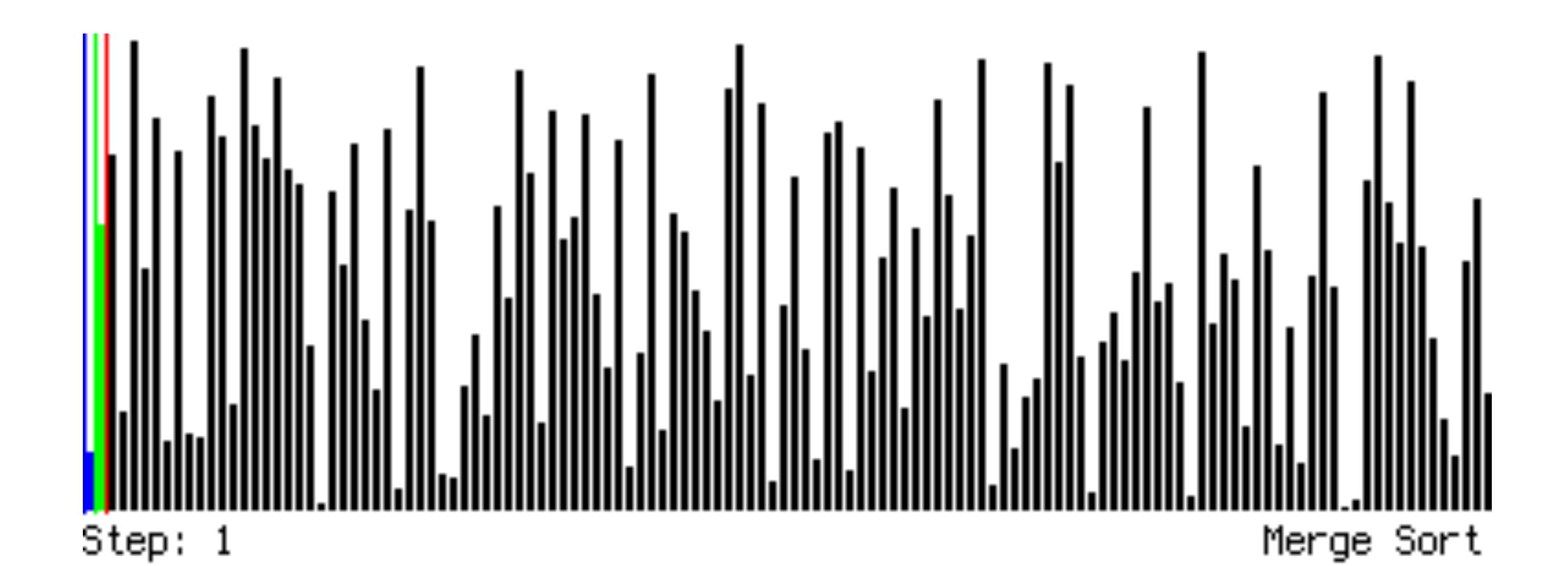
Agenda

- Sorting Algorithms (Complexity Overview)
- Merge-Sort Algorithm
- Parallel Merge-Sort Algorithm
- Distributed Implementation of Merge-Sort Algorithm (MPI)
- Theoretical Analysis

Sorting Animation



Sorting Animation

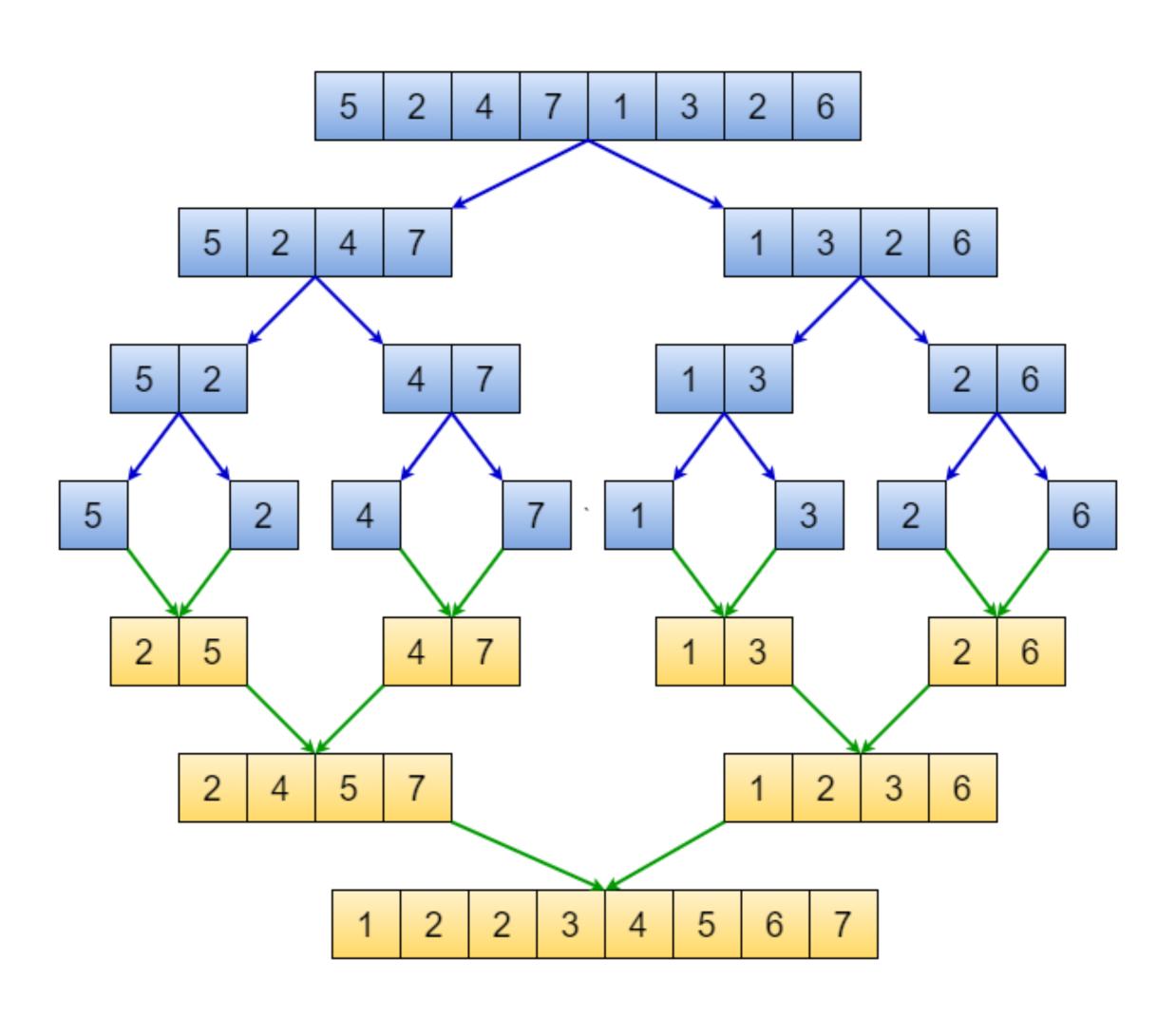


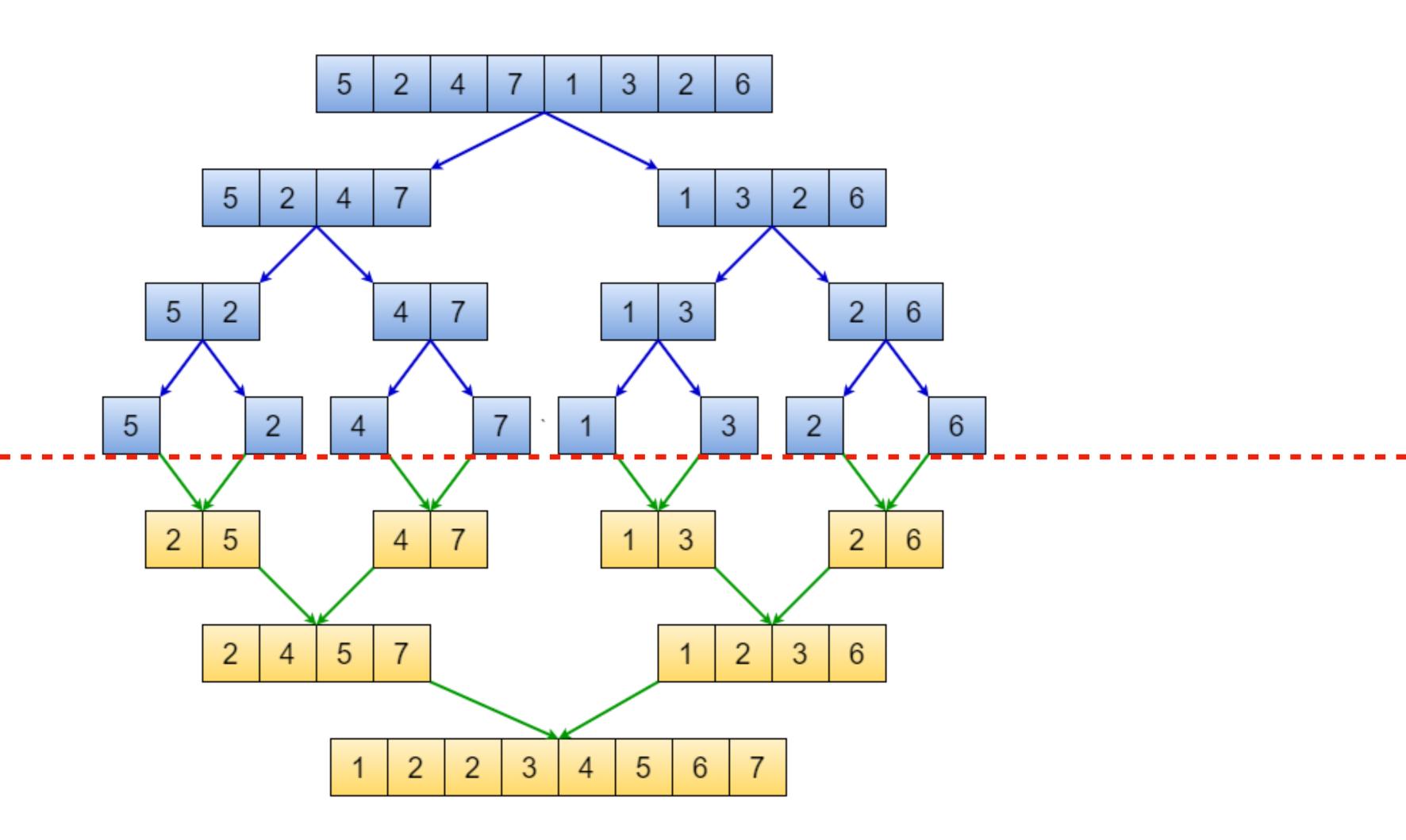
Sorting Algorithms

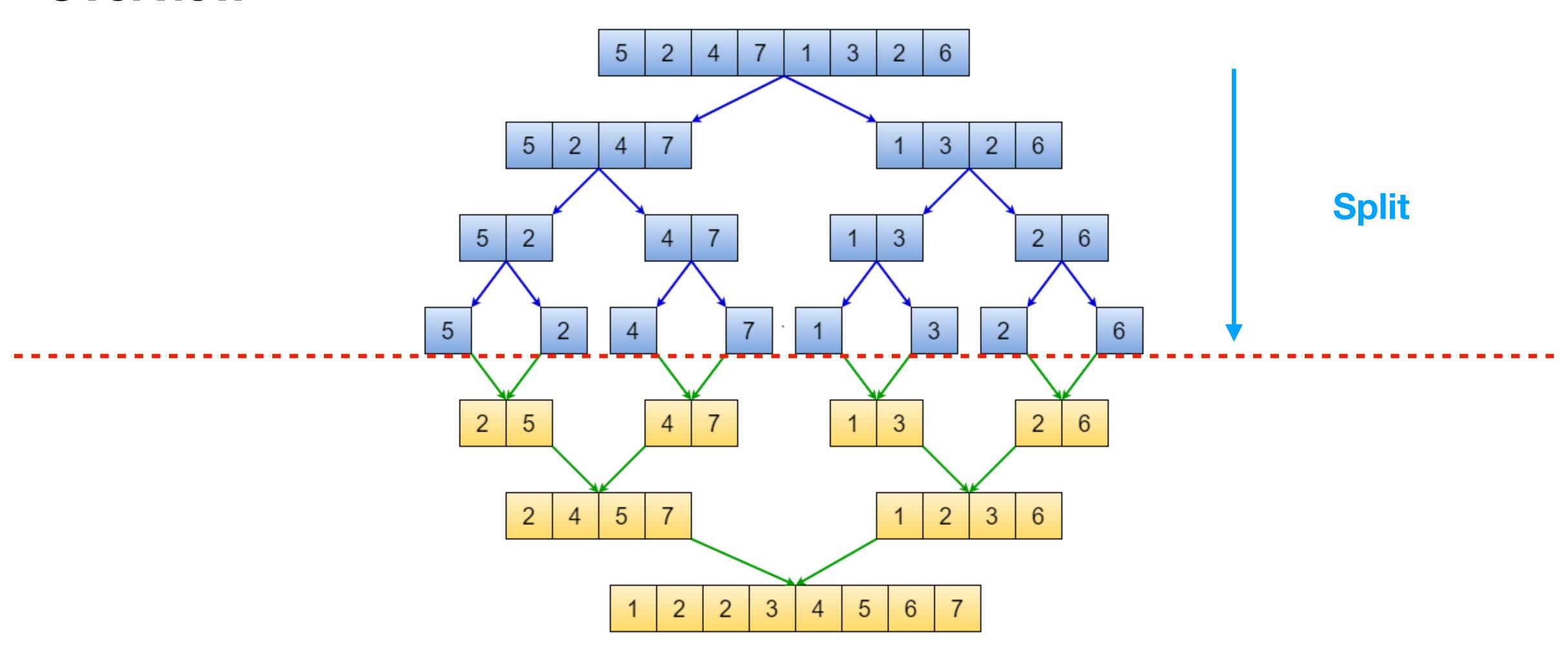
(Complexity Overview)

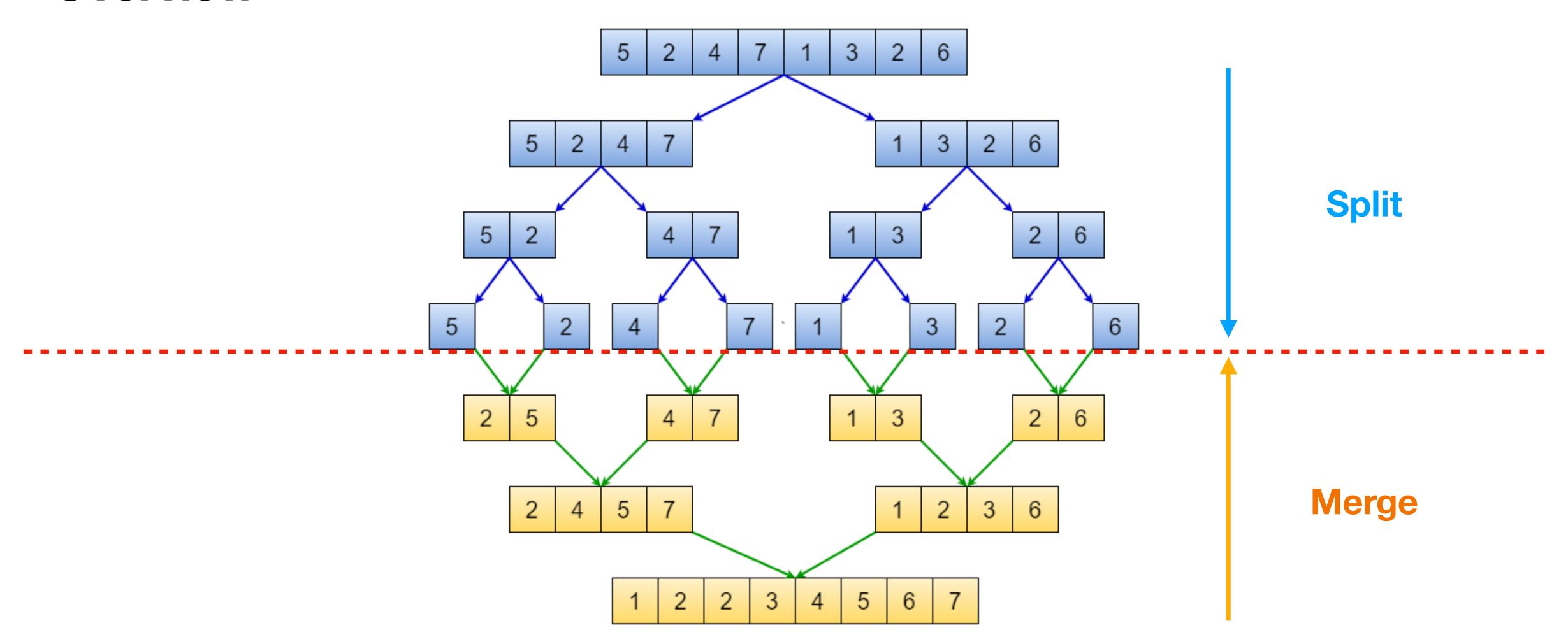
N: Input Size

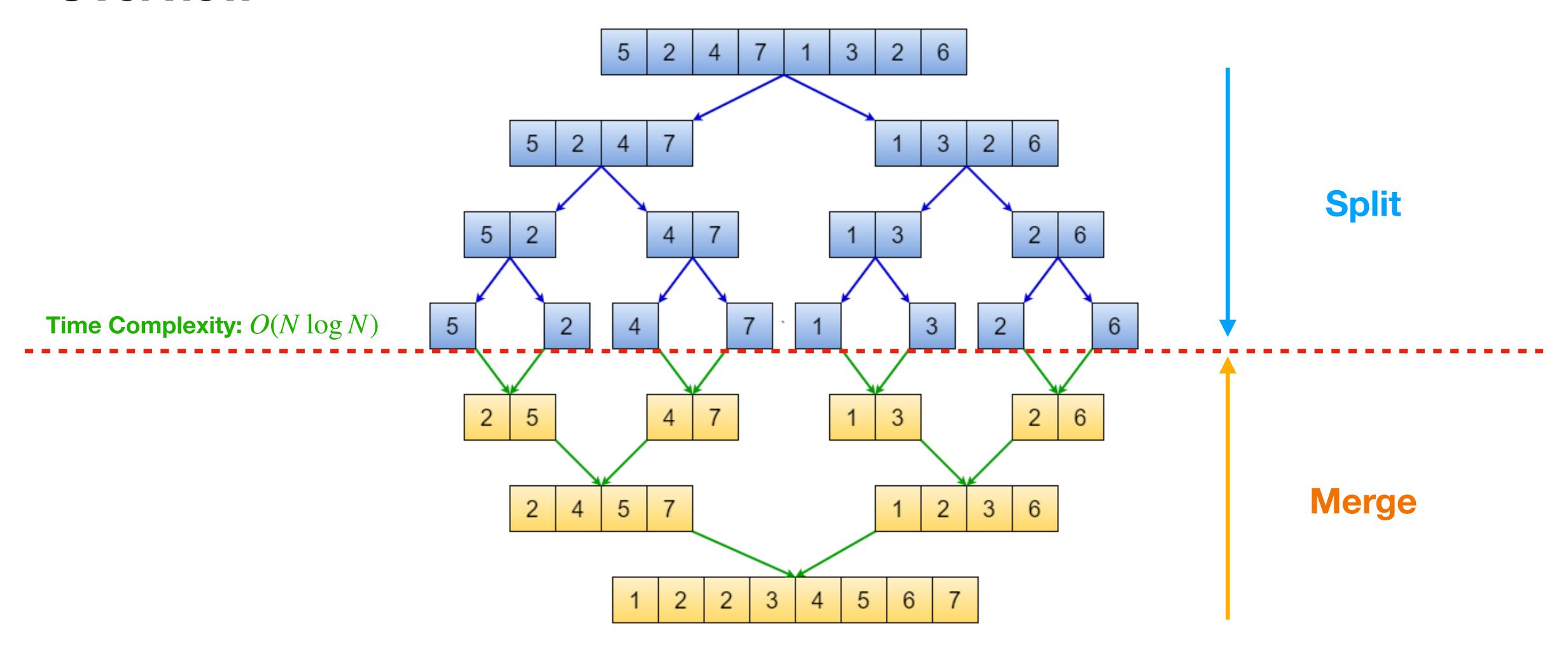
Algorithm	Best	Average	Worst	Memory	Stable	Method	Notes
Quick-Sort	N log N	N log N	N^2	log N	No	Partitioning	Easy to Implement
Heap-Sort	$N \log N$	$N \log N$	$N \log N$	In-place	No	Selection	
Intro-sort	N log N	N log N	N log N	log N	No	Partitioning & Selection	Used in C++ STL Containers
Merge-Sort	N log N	N log N	N log N	N	Yes	Merging	Highly Parallelizable

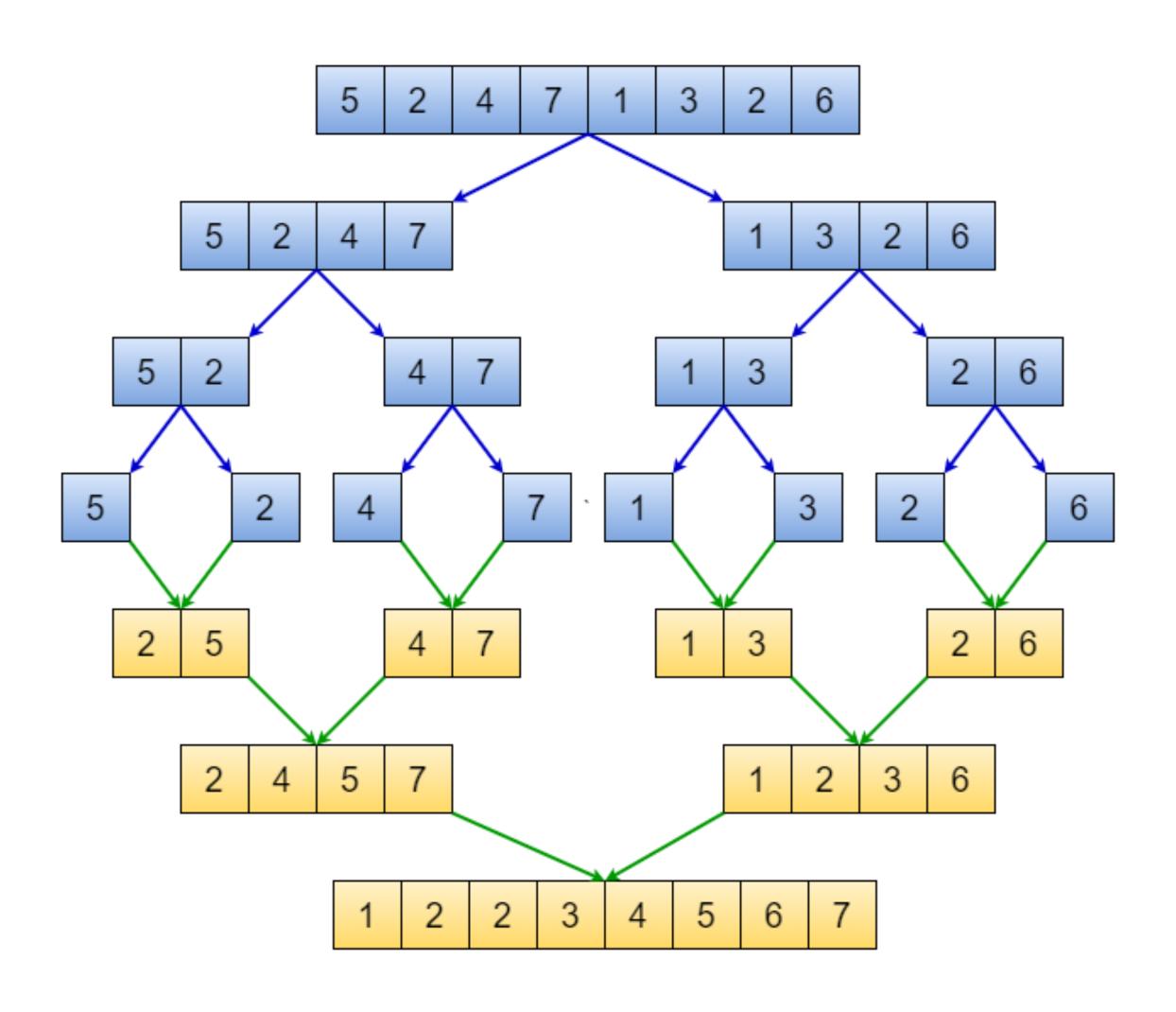


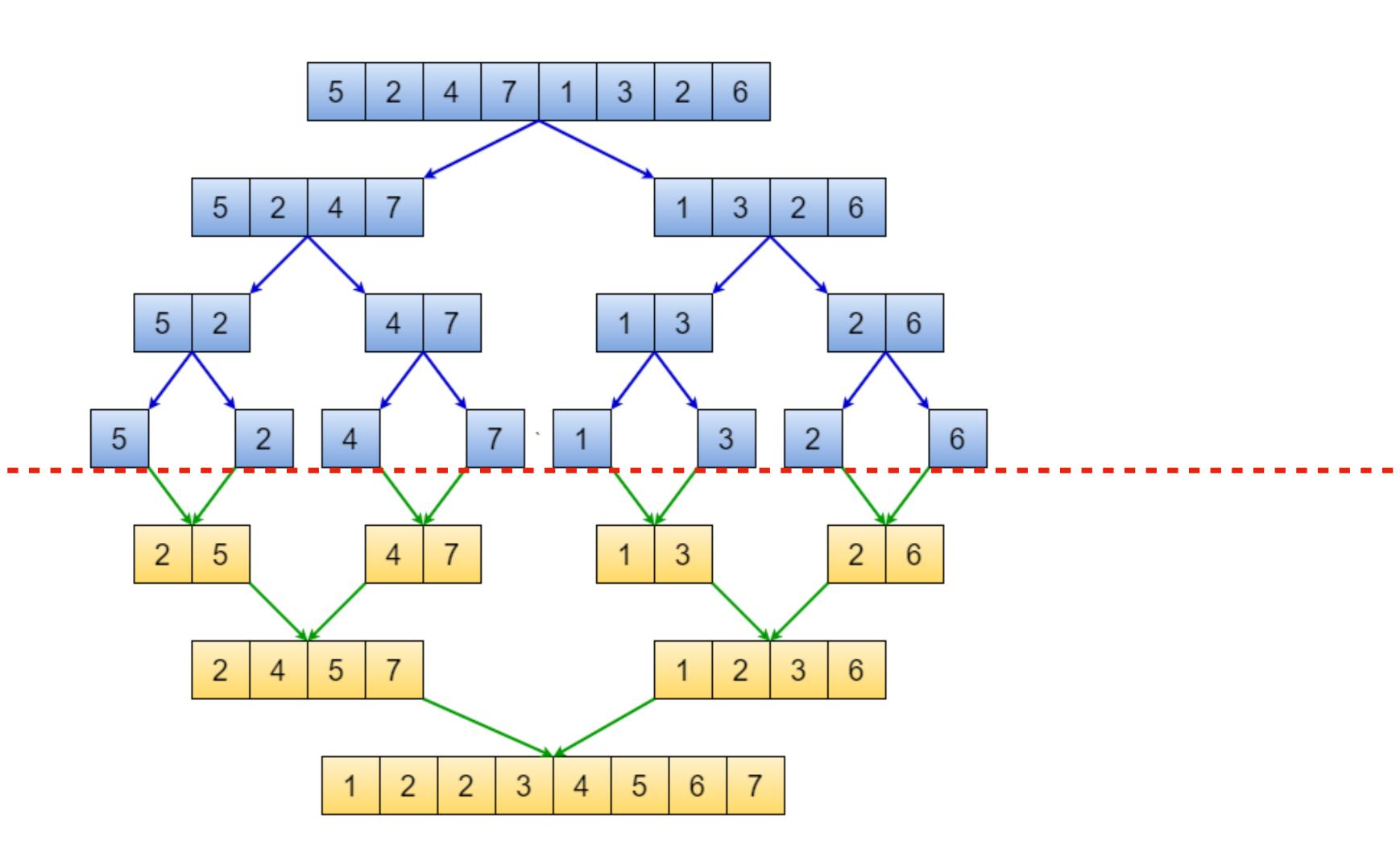


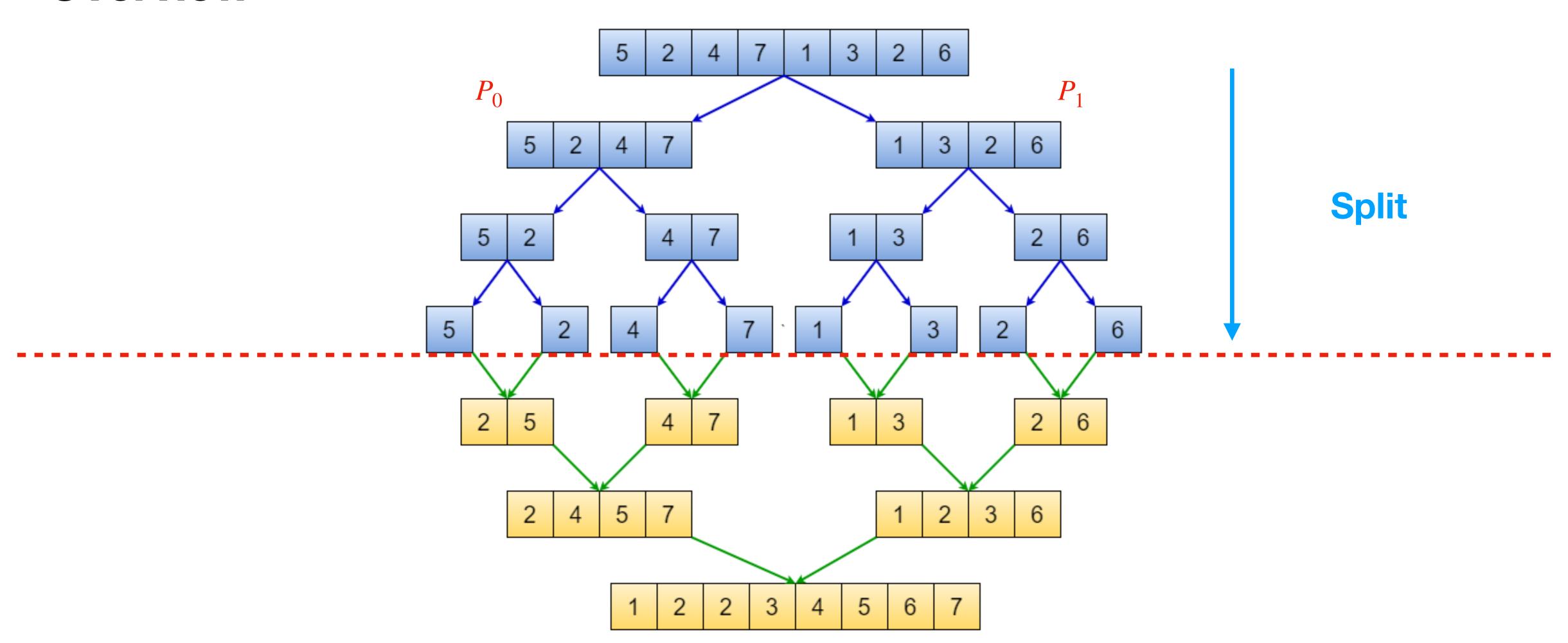


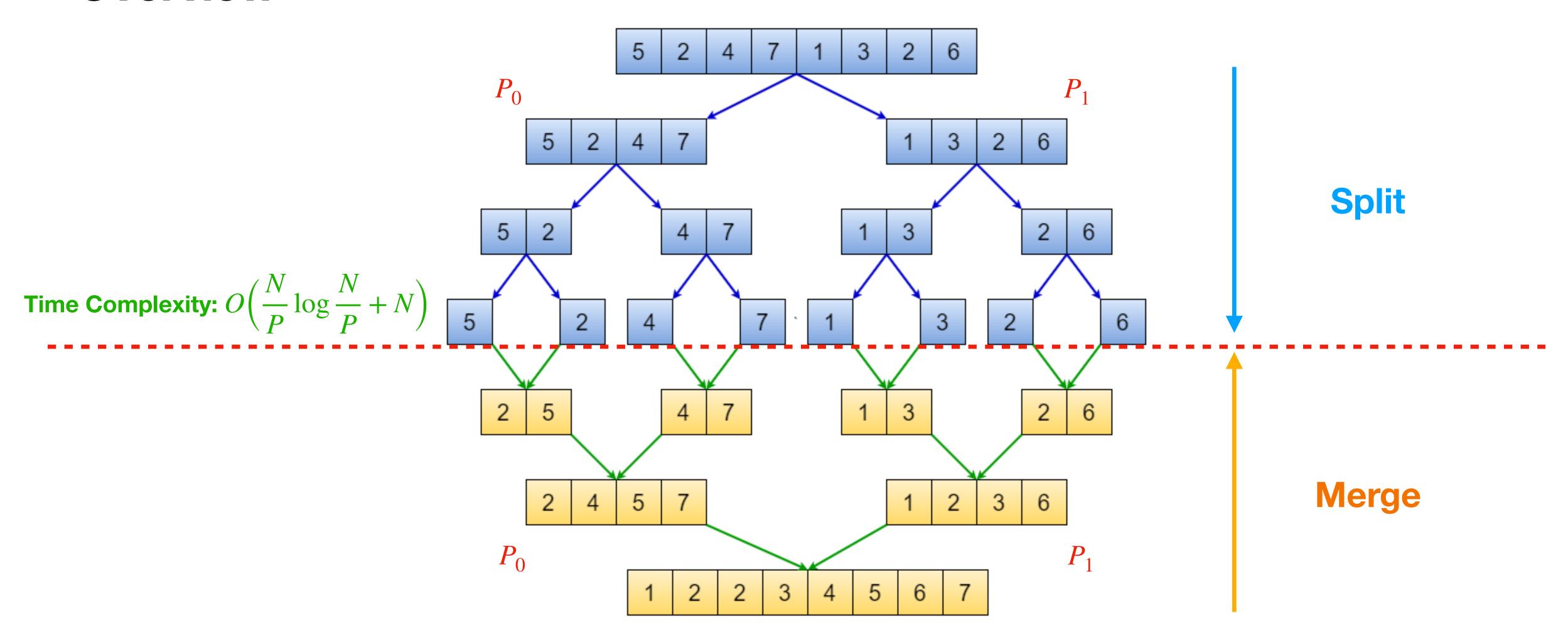








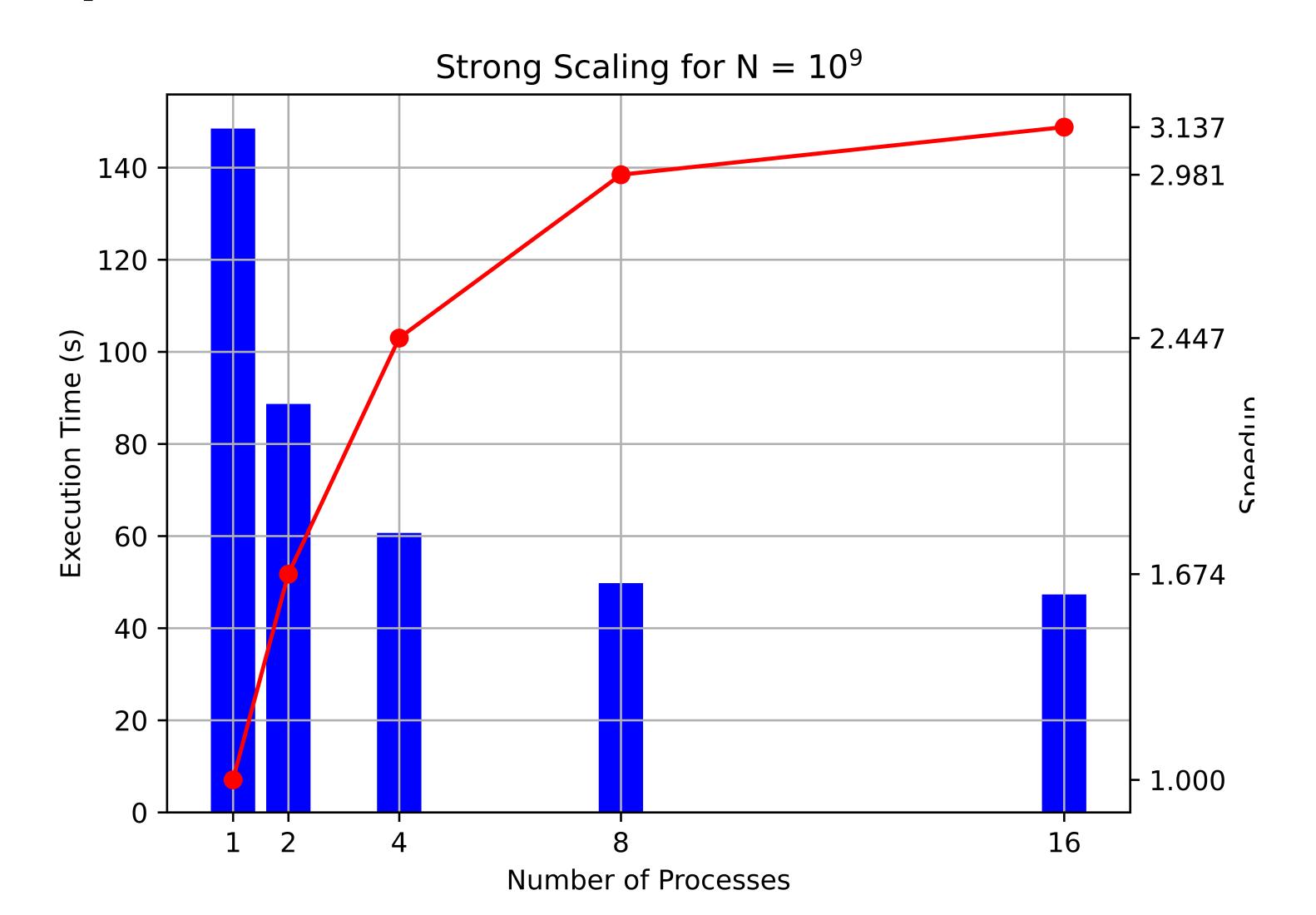




Merge-Sort Algorithm Distributed Implementation

Live Demo!

Distributed Implementation



Theoretical Analysis

Sequential Runtime:

$$T(n) = 2T(\frac{n}{2}) + O(n) = O(n \log n)$$

Parallel Runtime:

(with sequential merge)
$$T(n,p) = 2T(\frac{n}{2},\frac{p}{2}) + O(n) = O(\frac{n}{p}\log\frac{n}{p} + (n + \frac{n}{2} + \frac{n}{4}\cdots\frac{n}{p}))$$
$$\Longrightarrow T(n,p) = O(\frac{n}{p}\log\frac{n}{p} + n)$$

Upper bound on S (speedup):

$$\lim_{p \to \infty} \frac{O(n \log n)}{O(\frac{n}{p} \log \frac{n}{p}) + O(n)} = O(\log n)$$

Thank you! धन्यवाद!

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