ECON 6100

Problem Set 8

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1. The integer program for this problem is described as follows. Say that we have N villagers, where $x_{nf} = 1$ if n is assigned to fish, and $x_{nh} = 1$ if n is assigned to hunt.

$$\max_{x} \sum_{n=1}^{N} \alpha_n \cdot x_{nh} + \beta_n \cdot x_{nf}$$

subject to

s.t.
$$x_{nh} + x_{nf} \le 1 \ \forall \ n \in \{1, ..., N\}$$

 $x_{nh}, x_{nf} \in \{0, 1\} \ \forall \ n \in \{1, ..., N\}$

2. The linear relaxation is

$$\max_{x} \sum_{n=1}^{N} \alpha_n \cdot x_{nh} + \beta_n \cdot x_{nf}$$

subject to

s.t.
$$x_{nh} + x_{nf} \le 1 \ \forall \ n \in \{1, ..., N\}$$

 $x_{nh}, x_{nf} \ge 0 \ \forall \ n \in \{1, ..., N\}$

- 3. Observe that if, for some n, $\alpha_n > \beta_n$, a matching is only efficient if n is matched to hunt, because if assigned to fish (by any proportion greater than 0), they could strictly increase their total surplus by choosing to hunt. The only efficient match is if all villagers with $\alpha_n > \beta_n$ hunt and if all villagers with $\beta_n > \alpha_n$ fish. The stable payoffs are anything greater than their lower characteristic.
- 4. The linear relaxation is

$$\max_{x} \sum_{n=1}^{N} \alpha_n \cdot x_{nh} + \beta_n \cdot x_{nf}$$
s.t.
$$x_{nh} + x_{nf} \le 1 \ \forall \ n \in \{1, \dots, N\}$$

$$\sum_{n} x_{nh} \le V_h$$

$$x_{nh}, x_{nf} \ge 0 \ \forall \ n \in \{1, \dots, N\}$$

The dual of this relaxation is

$$\min_{q} \sum_{n=1}^{N} q_n + q_0 \cdot V_n$$

s.t.
$$q_n + q_0 \ge \alpha_n \ \forall \ n \in \{1, \dots, N\}$$
$$q_n \ge \beta_n \ \forall \ n \in \{1, \dots, N\}$$
$$q \ge 0$$

5. No. Consider the case where $V_n = 1$, N = 3, and the three villagers have:

$$\{(\alpha_n, \beta_n)\} = \{(10, 10), (9, 1), (8.9, 1)\}$$

The efficient match is to assign 1 and 3 to fish, and 2 to hunt, regardless of the fact that 1 is the best hunter.

- 6. Order the villagers in order of their $\alpha_n \beta_n$. If there are V_n or fewer villagers with $\alpha_n \beta_n \geq 0$, assign all of those to hunt and the rest to fish. If there are more than V_n villagers with $\alpha_n \beta_n > 0$, assign the V_n villagers with the largest gap to hunt and the rest to fish. This is feasible as long as the payment to each hunter is less than α_n and the payment to each fisher is less than β_n , and it is stable as long as the hunter with the lowest $\alpha_n \beta_n$ is paid at least $\alpha_n \beta_n$ more to hunt than to fish, and as long as the fisher with the highest $\alpha_n \beta_n$ is paid at most $\alpha_n \beta_n$ more to hunt than to fish.
- 7. We now have two capacity constraints. The primal is

$$\max_{x} \sum_{n} \alpha_{n} \cdot x_{nh} + \beta_{n} \cdot x_{nf}$$
 s.t.
$$x_{nh} + x_{nf} \leq 1 \,\forall \, n$$

$$\sum_{n} x_{nh} \leq V_{h}$$

$$\sum_{n} x_{nf} \leq V_{f}$$

$$x \geq 0 \,\forall \, n$$

The dual is

$$\min_{q} \sum_{n} q_{n} + q^{h} \cdot V_{h} + q^{f} \cdot V_{f}$$
 s.t.
$$q_{n} + q^{h} \ge \alpha_{n} \ \forall \ n$$

$$q_{n} + q^{f} \ge \beta_{n} \ \forall \ n$$

$$q \ge 0$$

- 8. The efficient assignment will be if we again sort villagers by their $\alpha_n \beta_n$, and we assign the top V_h to hunt and the bottom V_f to fish. Some villagers will be assigned to both, and will go to wherever they produce higher. One of the constraints will hold, and the other will be slack. WLOG, assume V_h binds. Then $q^h > 0$ and $q^f = 0$, and each hunter is paid β_n of the lowest hunter, each fisher is paid nothing. If V_f instead binds, the alternative holds.
- 9. We again order villagers by their $\alpha_n \beta_n$. The top V_h will hunt, the bottom V_f will fish, and the remaining will be idle. Note that this does not necessarily maximize the total caloric payoffs, but it ensures that the allocation will be stable.