

SEP: A Stable Election Protocol for clustered heterogeneous wireless sensor networks

Ibrahim Matta

Joint work with

Georgios Smaragdakis and Azer Bestavros



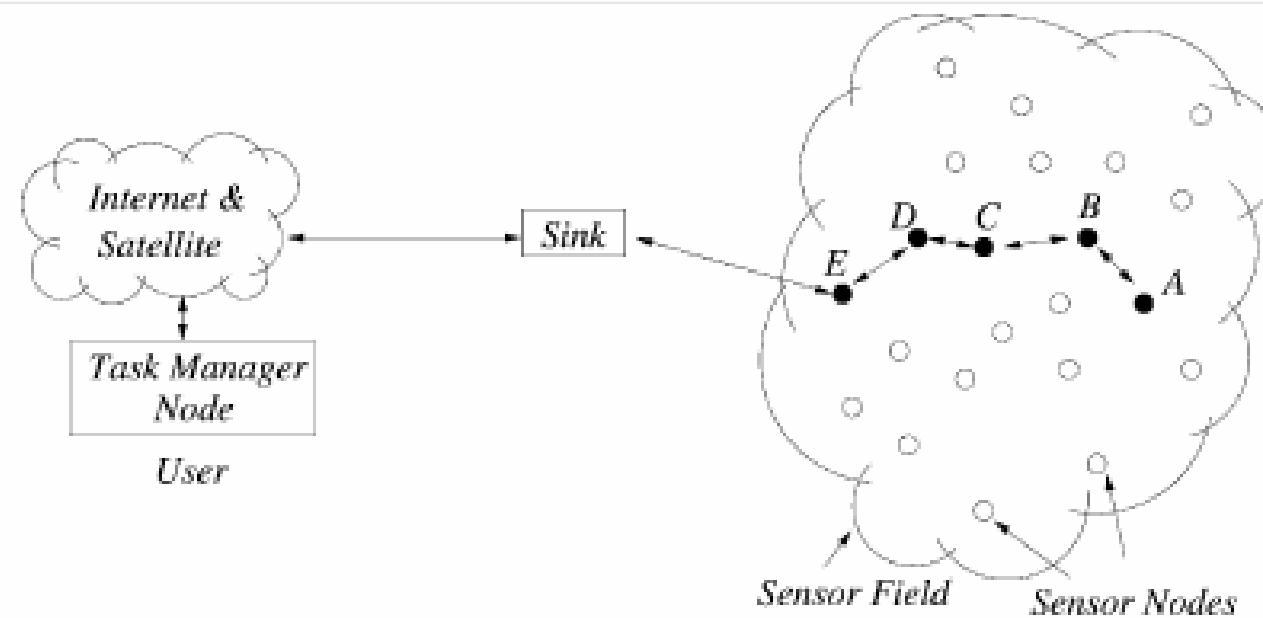
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Heterogeneous WSN

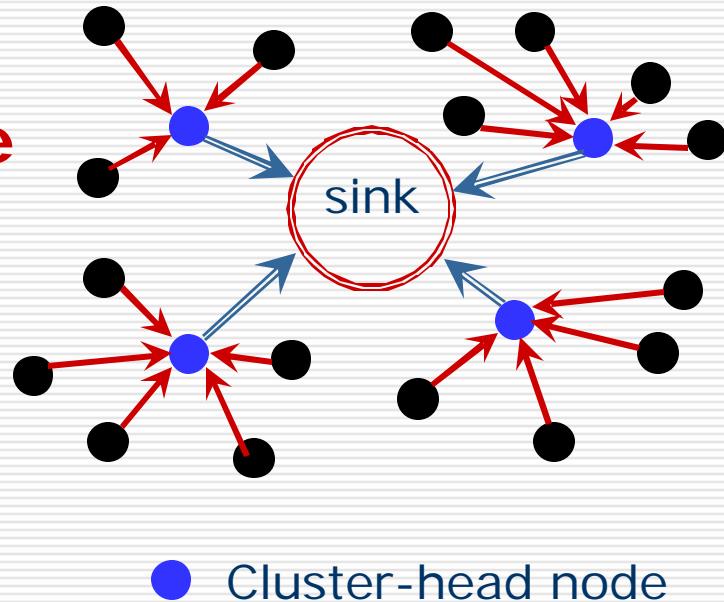
- Nodes have different energy levels
- Initial setting, after some operation time, or after re-energizing some sensors



Routing

- **Goal:**
prolong network lifetime/coverage

- Direct transmission to sink
- Min-energy routing
- **Sensing process can become biased**
- LEACH – Low Energy Adaptive Clustering Hierarchy
[Heinzelman *et al.*, 2000]



Selecting Cluster-heads in LEACH

- Node i chooses random number, s , between 0 and 1
- If $s < T(i)$, node i becomes a cluster head in current round where:

$$T(i) = \begin{cases} \frac{P}{1 - P \times (r \bmod \frac{1}{P})} & \text{if } i \in G \\ 0 & \text{otherwise} \end{cases}$$

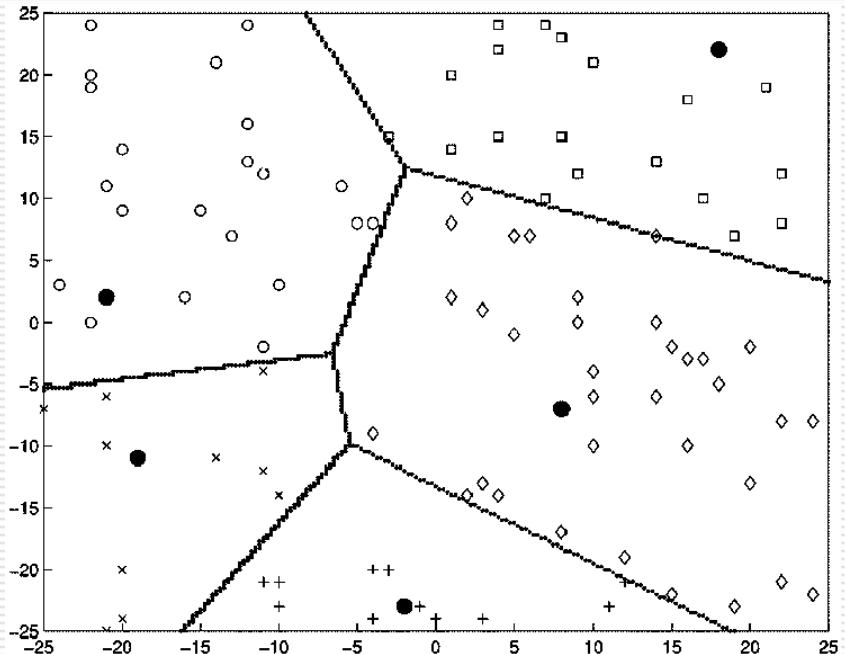
where:

P = desired percentage of cluster heads

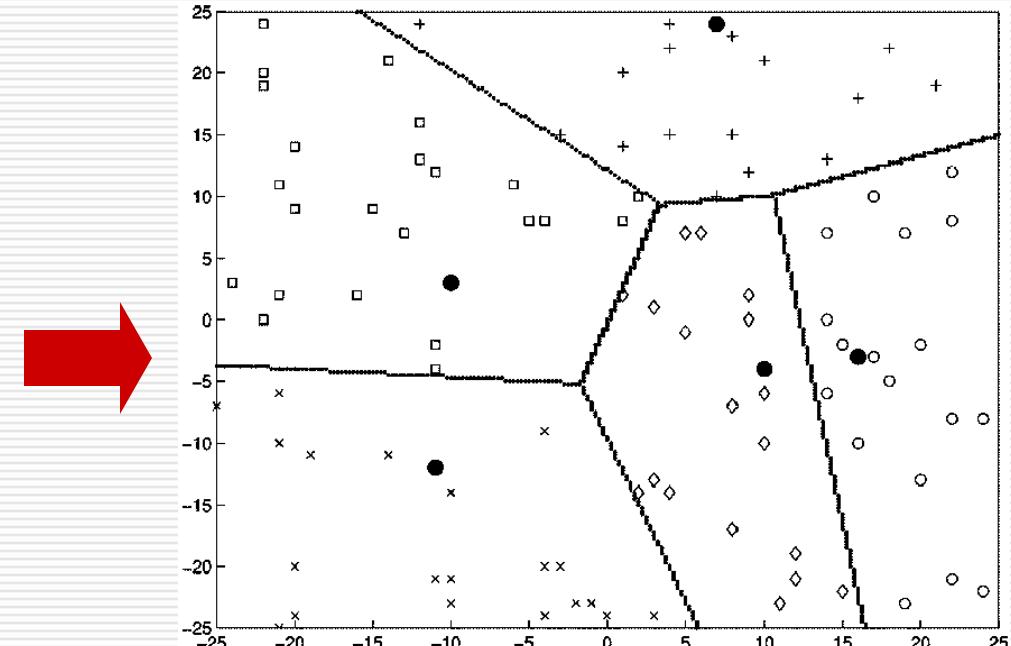
G = set of nodes that have not been a cluster head in the last $1/P$ rounds

- Each node is elected cluster-head once every $1/P$ rounds (epoch length)
- On average, $n \times P$ nodes elected per round
 - n = total number of nodes

Rotating Cluster-heads



Round r

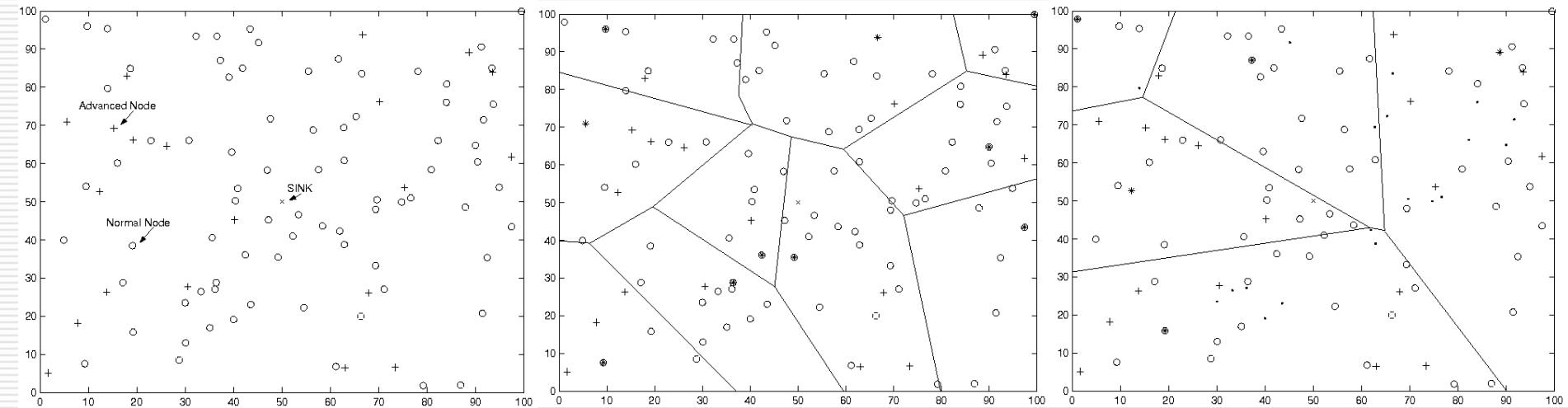


Round $r + 1$

When all nodes start with the same energy level (i.e., homogeneous setting), nodes “randomly” die and within a short period

But in a heterogeneous environment...

- The result of initial setting or evolution of the sensor network's operation
 - “advanced” node has α times more energy than “normal” node
- Once the first node dies, feedback and cluster-head election stays **unreliable** for a long time
 - advanced nodes don't get elected as often as they should

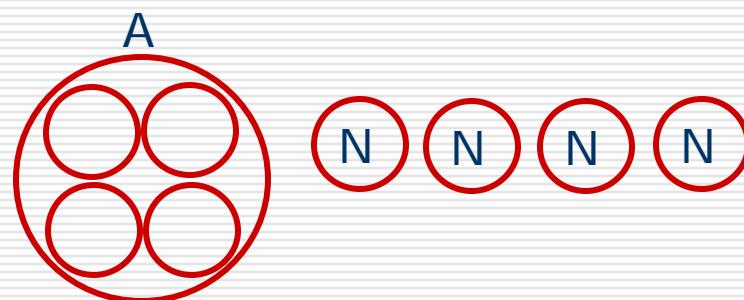


Naïve Modification to LEACH

- View network as homogeneous
 - Every “virtual” node is “normal” with unit energy
 - We have $n + \alpha \times m \times n = n(1 + \alpha \times m)$ “virtual” nodes
 - m = fraction of nodes that are advanced
- Extend epoch length to $(1 + \alpha \times m)/P$ rounds
- **Hope is:**
 - Each normal node is elected cluster-head once every epoch
 - Each advanced node is elected cluster-head $(1 + \alpha)$ times every epoch
 - This ensures well balanced energy consumption

Problem with Naïve Solution

- Probability of electing an advanced node is not weighted by its additional energy
 - advanced nodes may be underutilized
 - Lifetime is prolonged, but time until the first node dies can still be short!



$$n = 5, m = \frac{1}{5}, a = 3, P = \frac{1}{5}$$

"A" node elected with probability $T(A) = \frac{P}{1 - P \times 4}$





Our SEP (Stable Election Protocol)

- **Idea:** force each advanced node to be elected every sub-epoch of length $(1+\alpha \times m)/P / (1+\alpha)$ rounds
- Probability of a normal node getting elected as cluster-head is P_{normal}
- Probability of an advanced node getting elected as cluster-head is $P_{advanced}$
- Average number of nodes elected per round = $n \times P$

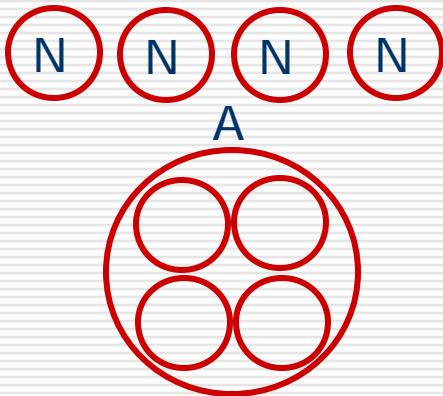
$$P_{normal} = \frac{P}{1 + a \times m}$$

$$T(i) = \begin{cases} \frac{P_{normal}}{1 - P_{normal} \times (r \bmod \frac{1}{P_{normal}})} & \text{if } i \in G_{normal} \\ 0 & \text{otherwise} \end{cases}$$

$$P_{advanced} = \frac{P}{1 + a \times m} (1 + a)$$

$$T(i) = \begin{cases} \frac{P_{advanced}}{1 - P_{advanced} \times (r \bmod \frac{1}{P_{advanced}})} & \text{if } i \in G_{advanced} \\ 0 & \text{otherwise} \end{cases}$$

Numerical Example



$$n = 5, m = 1/5, a = 3, P = 1/5$$



Our SEP Scheme

$$P_{normal} = \frac{P}{1 + a \times m} = \frac{1/5}{1 + 3 \times 1/5} = \frac{1}{8}$$

$$T(i) = \begin{cases} \frac{P_{normal}}{1 - P_{normal} \times (r \bmod \frac{1}{P_{normal}})} & \text{if } i \in G_{normal} \\ 0 & \text{otherwise} \end{cases}$$

$$P_{advanced} = \frac{P}{1 + a \times m} (1 + a) = \frac{1}{8} (1 + 3) = \frac{4}{8}$$

$$T(i) = \begin{cases} \frac{P_{advanced}}{1 - P_{advanced} \times (r \bmod \frac{1}{P_{advanced}})} & \text{if } i \in G_{advanced} \\ 0 & \text{otherwise} \end{cases}$$

Performance Measures

□ Stability Period

- Time until death of the first node

□ Network lifetime

- Time until death of the last alive node

□ Number of cluster-heads per round

- Nodes which will directly send aggregated information to the sink

□ Number of alive nodes of each type per round

□ Throughput

- Rate of data reporting to cluster-heads and to sink



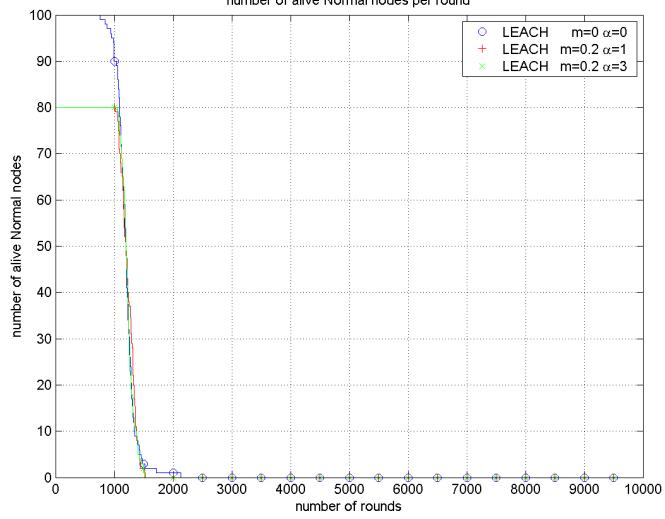
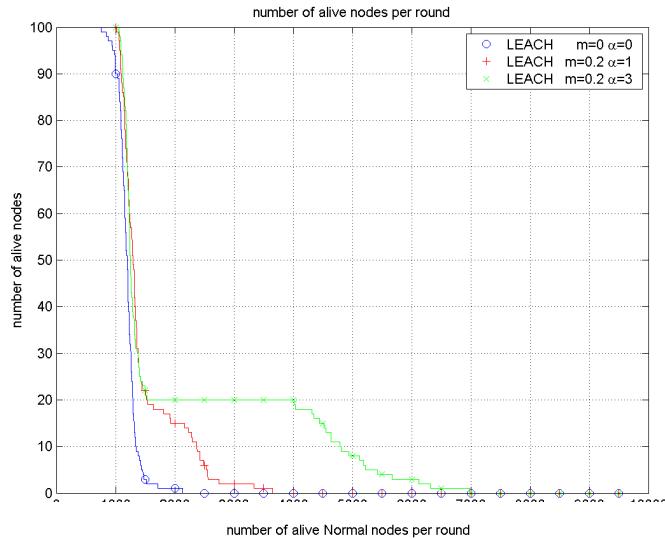
Simulation Model

- Two networks
 - 100 nodes uniformly distributed over 100mx100m
 - 900 nodes uniformly distributed over 300mx300m
- Sink is in the center of the sensor field
- Message size = 4000 bits
- Initial energy for a normal node = 0.5 Joules
- Energy consumed depends on radio characteristics

Results for LEACH

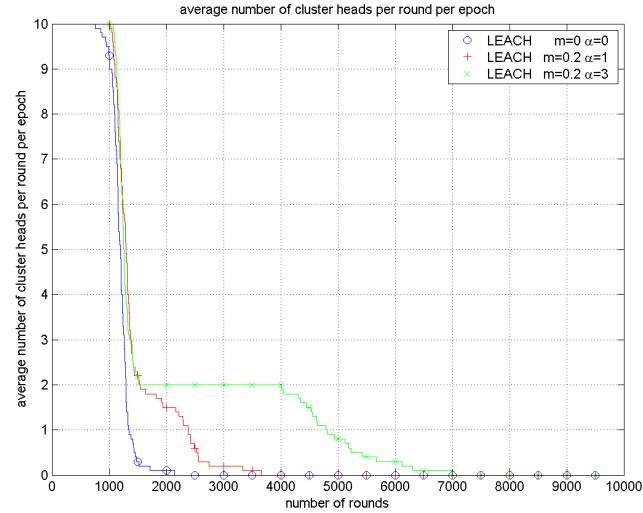
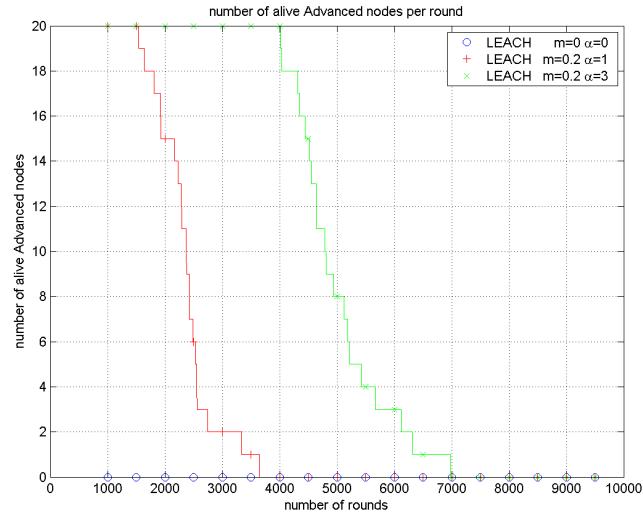
- LEACH does not take full advantage of the extra energy of advanced nodes

- Normal nodes die very fast
 - Sensing field becomes sparse very fast
 - Election process becomes unstable fast
 - Throughput is low



Results for LEACH (cont'd)

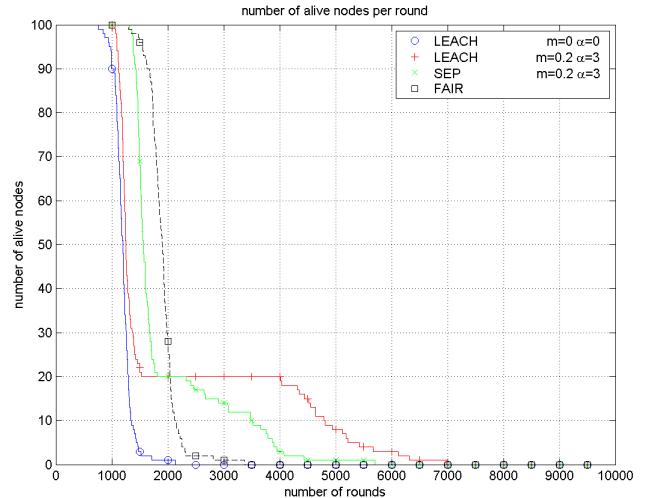
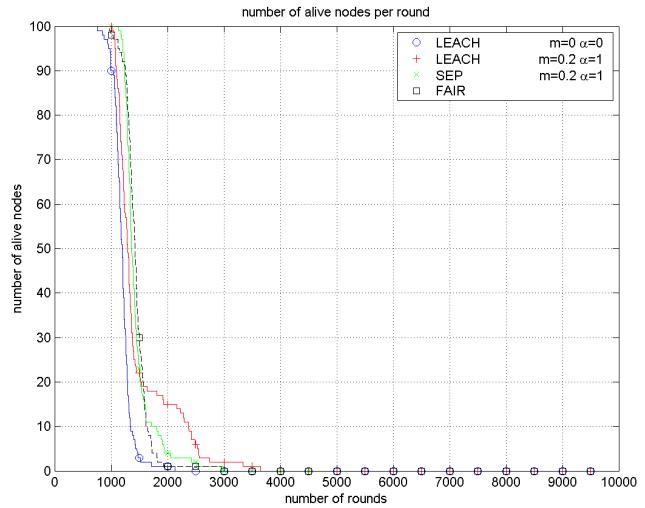
- Advanced nodes die very slowly
 - They are not elected as often as they should
 - Throughput is low



Results for SEP

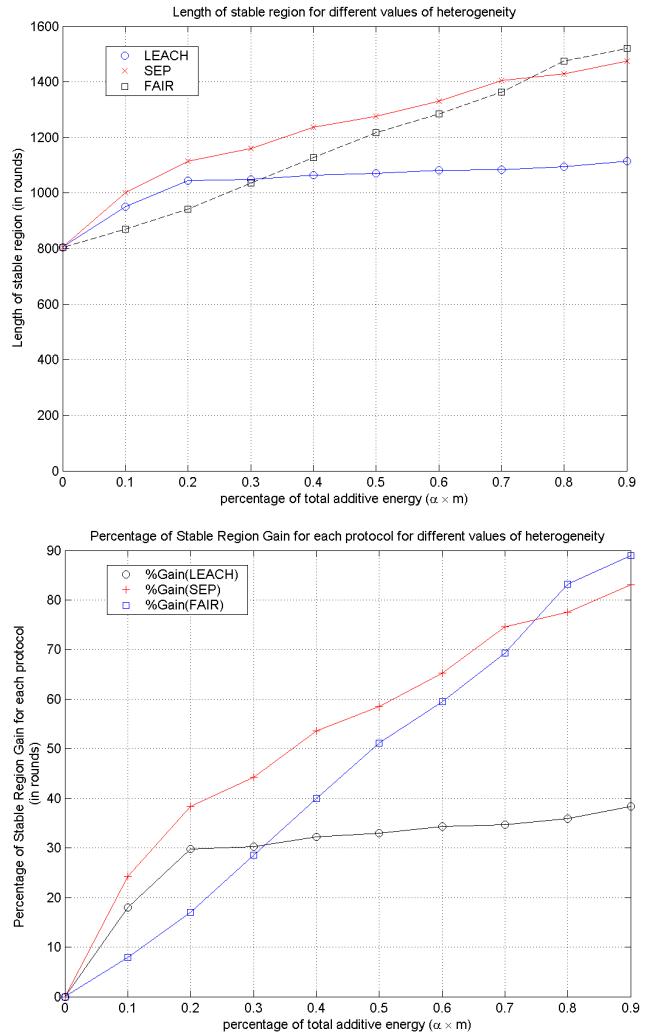
- SEP takes full advantage of the extra energy of advanced nodes
 - Stable region increases by up to 26% over LEACH
 - Higher throughput

- Advanced nodes follow the death process of normal nodes
 - Smaller unstable region



Sensitivity of SEP

- Performance depends on the product $m \times a$
- FAIR: ideal distribution of extra energy
- Gain in stability period (over system with no extra energy, i.e., $m=0$ & $\alpha=0$) is maximized under SEP
 - SEP > LEACH
 - SEP > FAIR for up to 75% relative extra energy





Conclusion

- SEP is heterogeneous-aware
 - Cluster-heads elected based on node's energy relative to that of others
- SEP is scalable and dynamic
 - Even normal node can be elected
 - No global knowledge required at every round
 - No prior distribution of energy levels assumed
- Study SEP for more than 2 hierarchical levels and more than 2 types of nodes
- Implementation in Berkeley/Crossbow motes
- Deployment issues
 - Dynamic updates of weighted election probabilities
 - Integration with power-aware MAC protocols
 - ...

**For more information, please check
<http://csr.bu.edu/sep/>**