

Introduction to Dependent Types

Eagan Technology Unconference

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Agenda

- 1 Preface
- 2 Review of Basics
- 3 What is Dependent Type
- 4 Steps toward Dependent Types
- 5 Closing

Section Outline

1 Preface

Quick Question

How many are familiar with this topic?

A Joke

This is not a \mathbb{M} -tutorial, and nothing here will involve burritos.

Disclaimer

There will be many code examples with *very* loose translations to imperative/OOP as we go along. Though please keep in mind that these are merely made up syntactical translations, the actual concepts may differ vastly.

About This Talk

Example languages with dependent types:

- Idris
- Epigram
- Agda
- Coq

About This Talk

Example languages with dependent types:

- Idris
- Epigram
- Agda
- Coq

But we will be using Haskell though.

Honestly, it's because they're way over his head...

Dependent Types

Why do we want them?

- more expressive type system
- encode stronger invariants
- proving correctness of code

Teaser

Example please:

```
infixr 5 :>
data Vect (n :: Nat) a where
  VNil :: Vect 0 a
  (:>) :: a -> Vect n a -> Vect (n + 1) a

vs :: Vect 6 Int
vs = 4 :> 8 :> 15 :> 16 :> 23 :> 42 :> VNil
```

Translation* please:

```
enum Vect<Nat n, A> {
  Vect<0, A> VNil,
  Vect<n '+ 1, A> VCons(A a, Vect<n, A> va)
}

Vect<6, Int> vs = VCons(4, VCons(8, VCons(15, VCons(16,
  VCons(23, VCons(42, VNil))))));
```

(*) *supreme looseness and totally made-up syntax!!!*

Section Outline

- 2 Review of Basics
 - Values and Types
 - Defining Data Types
 - Functions

Values and Types

Values has Types, or Values are classified by Types.

```
..., -1, 0, 1, 2, 3, ... :: Int
```

```
True, False :: Bool
```

```
'a', 'b', 'c' :: Char
```

```
"abc" :: String ~ [Char]
```

About Data Types

How are data types defined?

- Some are built in magic: `Int`, `Char`, function arrow
- Some are built in sugar: list, tuples
 - We can define equivalent non-sugar version ourselves
- Rest can be user defined: `Bool`, `String`, `Maybe`

About Data Types

What are the data types like?

- Multiple **Value** constructors
- Parametrize over another **Type**
- Recursive definition
- Synonyms of other **Types**
- Combination of the above

Defining Data Types

Define new data type with `data`.

- Left hand side (LHS) - `Type` constructor
- Right hand side (RHS) - `Value` constructor

`Type` and `Value` constructors are capticalized.

Our First Example!

Define a person:

```
-- | params for firstname, lastname, age respectively
data Person = Person String String Int

barbara :: Person
barbara = Person "Barbara" "Smith" 30
```

The **Type** of the **Person Value** constructor:

```
Person :: String -> String -> Int -> Person
```

A loose translation:

```
enum Person {
  Person(String firstname, String lastname, Int age)
}

Person barbara = new Person("Barbara", "Smith", 30)
```


Multiple Value Constructors

Data can have multiple **Value** constructors:

```
data Bool = False | True
```

```
data Weekdays = Sunday | Monday | Tuesday | Wednesday  
               | Thursday | Friday | Saturday
```

A loose translation:

```
enum Bool { False, True }
```

```
enum Weekdays {  
    Sunday, Monday, Tuesday, Wednesday, Thursday, Friday,  
    Saturday  
}
```

Multiple Value Constructor

You can do type aliasing with `type`:

```
type Side = Double
type Radius = Double
```

For example:

```
data Shape = Triangle Side Side Side
           | Rectangle Side Side
           | Circle Radius
```

A loose translation:

```
enum Shape {
  Triangle(Double side1, Double side2, Double side3),
  Rectangle(Double length, Double width),
  Circle(Double radius)
}
```

Multiple Value Constructor

Recall `Side ~ Radius ~ Double`:

```
data Shape = Triangle Side Side Side
           | Rectangle Side Side
           | Circle Radius
```

Types of the 3 `Value` constructors:

```
Triangle  :: Side -> Side -> Side -> Shape
Rectangle :: Side -> Side -> Shape
Circle    :: Radius -> Shape
```

Example `Shapes`:

```
myTri, myRect, myCir :: Shape
myTri  = Triangle 2.1 3.2 5
myRect = Rectangle 4 4
myCir  = Circle 7.2
```

Parametrization

Types can parametrize over another type:

```
data Tuple a b = Tuple a b
```

With:

```
Tuple :: a -> b -> Tuple a b
```

A loose translation:

```
enum Tuple<A, B> {  
  Tuple(A a, B b)  
}
```

Tuple

Actual built-in sugar:

```
data Tuple a b = Tuple a b
=> data (,) a b = (,) a b
=> data (a, b) = (a, b)
```

An example:

```
type Employed = Bool

barbara, chet, luffy :: (Person, Employed)
barbara = (Person "Barbara" "Smith" 30, True)
chet    = (Person "Chet" "Awesome-Laser" 2, False)
luffy   = (Person "Luffy D." "Monkey" 19, False)
```

Maybe

Like `Bool`, but parametrizes a `Type a` over the `True` part:

```
data Maybe a = Nothing | Just a
```

With:

```
Nothing :: Maybe a  
Just    :: a -> Maybe a
```

A loose translation:

```
enum Maybe<A> {  
  Nothing,  
  Just(A a)  
}
```

Maybe

From previous slide:

```
data Maybe a = Nothing | Just a
```

Say more with **Occupation**:

```
type Occupation = Maybe String
```

```
barbara2, chet2, luffy2 :: (Person, Occupation)  
barbara2 = (Person "Barbara" "Smith" 30, Just "dancer")  
chet2    = (Person "Chet" "Awesome-Laser" 2, Nothing)  
luffy2   = (Person "Luffy D." "Monkey" 19, Just "pirate"  
            )
```

Types with Recursion

Natural number:

```
data Nat = Z | S Nat
```

With:

```
Z :: Nat  
S :: Nat -> Nat
```

A loose translation:

```
enum Nat {  
  Z,  
  S(Nat n)  
}
```


Types with Recursion

Natural number:

```
data Nat = Z | S Nat
```

```
Z :: Nat
```

```
S :: Nat -> Nat
```

```
0 ~ Z
```

```
1 ~ S Z
```

```
2 ~ S (S Z)
```

```
3 ~ S (S (S Z))
```

Types with Recursion

List - recursive **Type** that parametrizes over another **Type**:

```
data List a = Nil | Cons a (List a)
```

With:

```
Nil    :: List a  
Cons   :: a -> List a -> List a
```

A loose translation:

```
enum List<A> {  
  Nil,  
  Cons(A a, List<A> as)  
}
```

Types with Recursion

Actual built-in sugar is something like:

```
data List a = Nil | Cons a (List a)
=> data [] a = [] | (:) a ([] a)
=> data [a] = [] | (:) a [a]
```

Sugar that `List`:

```
ints :: List Int
ints = Cons 1 (Cons 2 (Cons 3 (Cons 4 Nil)))

-- built-in sugar
ints :: [] Int
ints = 1 : 2 : 3 : 4 : []

-- 2x the sugar!
ints :: [Int]
ints = [1, 2, 3, 4]
```

Functions

Maps **Values** of a **Type** to **Values** of another **Type**:

```
even :: Int -> Bool
even 0 = True
even n = if rem n 2 == 0
         then True
         else False
```

Not as loose translation:

```
Bool even(Int n) {
  switch n:
    case n == 0:
      return True;
  default:
    if rem(n, 2) == 0
      return True;
    else
      return False;
}
```

Functions with Recursion

Use recursion for recursive **Types**:

```
data Nat = Z | S Nat
```

```
toInt :: Nat -> Int
```

```
toInt Z = 0
```

```
toInt (S n) = 1 + toInt n
```

Not as loose translation:

```
Int toInt(Nat n) {  
  switch n:  
  case Z:  
    return 0;  
  case (S m): -- n ~ (S m)  
    return 1 + toInt(m);  
}
```

Functions with Recursion

Use recursion for recursive **Types**:

```
data Nat = Z | S Nat
```

```
toInt :: Nat -> Int
```

```
toInt Z = 0
```

```
toInt (S n) = 1 + toInt n
```

Evaluation is a series of substitutions:

```
three = S (S (S Z)) :: Nat
```

```
toInt three :: Int
```

```
= toInt (S (S (S Z)))
```

```
= 1 + toInt (S (S Z))
```

```
= 1 + 1 + toInt (S Z)
```

```
= 1 + 1 + 1 + toInt Z
```

```
= 1 + 1 + 1 + 0
```

```
= 1 + 1 + 1
```

```
= 1 + 2
```

```
= 3
```

Functions with Parametric Polymorphism

Functions can be parametric:

```
data [a] = [] | (:) a [a]

length :: [a] -> Int
length []      = 0
length (x:xs) = 1 + length xs
```

A translation:

```
Int length(List<A> ls) {
  switch l1:
  case Nil:
    return 0;
  case Cons(x, xs):
    return 1 + length(xs);
}
```

Functions with Parametric Polymorphism

Functions can be parametric:

```
data [a] = [] | (:) a [a]

length :: [a] -> Int
length []      = 0
length (x:xs) = 1 + length xs
```

Evaluation is a series of substitutions:

```
xs = [4, 8, 15] = 4 : 8 : 15 : [] :: [Int]

length xs :: Int
= length [4, 8, 15]
= 1 + length [8, 15]
= 1 + 1 + length [15]
= 1 + 1 + 1 + length []
= 1 + 1 + 1 + 0
= 1 + 1 + 1
= 1 + 2
= 3
```


Higher-order Functions

Functions that take functions as params:

```
-- actual name is ($)
apply :: (a -> b) -> a -> b
apply f x = f x

-- actual name is (.)
compose :: (b -> c) -> (a -> b) -> (a -> c)
compose f g = \x -> f (g x)
```

Yay translations:

```
B apply(Func<A, B> f, A a) {
  return f(a);
}

Func<A, C> compose(Func<B, C> f, Func<A, B> g) {
  return x => f(g(x));
}
```

Section Outline

- 3 What is Dependent Type
 - λ -Calculus
 - Extensions on λ -calculus

λ -Calculus

So far, we have seen:

- function application
- function abstraction (aka higher-order functions)
- variable binding
- substitution

\Rightarrow basis for simply typed λ -calculus.

λ -Calculus

You: Sure... Me: Ah, yes, we want to extend λ -calculus so we can have more forms of abstractions!

λ -Calculus

You: Sure... Me: Ah, yes, we want to extend λ -calculus so we can have more forms of abstractions!

Q: But how?

A: What if I tell you...

...you should already be familiar with 2 axes of extension :)

Subtype Polymorphism

Given data types T and P , if there is a relation between T and P by some notion of substitutability with T in place of P , then we say T is a subtype of the supertype P , denoted by $T <: P$.

There is an extension on λ -calculus with subtype polymorphism and is denoted by $\lambda_{<:}$.

=> Object Oriented Programming

Though this is not an axis that we are interested in.

Parametric Polymorphism

Introduce a mechanism of universal quantification over **Types**: **Types** can abstract over **Types**, allows for generic data types and generic functions.

=> Generic Programming

Recall:

```
data Maybe a = Nothing | Just a
data List a = Nil | Cons a (List a)

(.) :: (b -> c) -> (a -> b) -> (a -> c)
map :: (a -> b) -> [a] -> [b]
```

The name for this extension is formally second order λ -calculus, aka System F, denoted by $\lambda 2$,

Value and Type Interdependency

Re-thinking functions:

```
even :: Int -> Bool
even 0 = True
even n = if rem n 2 == 0
         then True
         else False
```

even maps **Ints** to **True** and **False**.

=> **Values** on RHS depends on the **Values** on LHS

=> **Values** depending on **Values** => Ordinary λ -calculus

Value and Type Interdependency

Re-thinking parametrized data types:

```
data Maybe a = Nothing | Just a
```

```
data List a = Nil | Cons a (List a)
```

`Maybe` and `List` take a `Type` and return `Value` constructors

=> `Values` on RHS depends on the `Type` on LHS

=> `Values` depending on `Types`

=> Parametric polymorphism of $\lambda 2$

Value and Type Interdependency

Then what about the other cases of dependencies?

- **Values** depending on **Values**: λ -calculus
- **Values** depending on **Types**: $\lambda 2$, System F

Value and Type Interdependency

Then what about the other cases of dependencies?

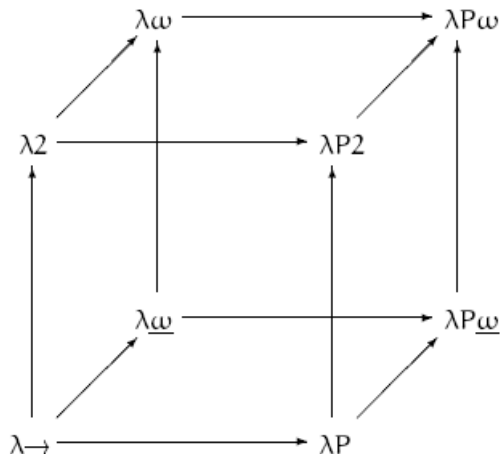
- **Values** depending on **Values**: λ -calculus
- **Values** depending on **Types**: $\lambda 2$, System F
- **Types** depending on **Types**: λ_{ω}
=> Type-level programming via type operators

Value and Type Interdependency

Then what about the other cases of dependencies?

- **Values** depending on **Values**: λ -calculus
- **Values** depending on **Types**: $\lambda 2$, System F
- **Types** depending on **Types**: λ_{ω}
=> Type-level programming via type operators
- **Types** depending on **Values**: $\lambda\Pi$
=> Dependent types

Lambda Cube



System F_c

Currently, Haskell as of GHC 7.10.2

- no true type operators
- type-level programming through:
 - type families
 - equalities and coercions on [Types](#)

This axis of extension on $\lambda 2$ is termed System F_c .

System F_c

Currently, Haskell as of GHC 7.10.2

- not truly dependent either:
 - strong distinction between **Values** and **Types**
- emulate dependent types with:
 - handful of language extensions
 - **Kind** system

Section Outline

- 4 Steps toward Dependent Types
 - Kinds
 - Language Extensions
 - Dependent Type Programming with Vectors
 - Pi and Sigma Types
 - Heterogeneous Collections

Kinds

Q: **Types** classify **Values**, but what classifies **Types**?

A: **Kinds**

Introducing ★

```
-- built-in magic: infinitely many value constructors
data Int = ... | -1 | 0 | 1 | 2 | ...
data Bool = False | True
data [a] = Nil | (:) a [a]
data Maybe a = Nothing | Just a
data (a, b) = (a, b)
```

```
Int :: *
Bool :: *
[Int] :: *
[] :: * -> *
Maybe Person :: *
Maybe :: * -> *
(Person, Bool) :: *
(,) Person :: * -> *
(,) :: * -> * -> *
```

Other Kinds

There are other **Kinds** aside from *****:

```
import GHC.Prim

(*)          -- kind of fully realized type
(#)         -- kind of unboxed stuff used internally
Constraint  -- kind of constraints and type equality
OpenKind    -- superkind of (*) and (#)
AnyK        -- polymorphic kind for flexible arity
```

Other Kinds

There are other **Kinds** aside from *****:

```
import GHC.Prim

(*)          -- kind of fully realized type
(#)         -- kind of unboxed stuff used internally
Constraint  -- kind of constraints and type equality
OpenKind    -- superkind of (*) and (#)
AnyK        -- polymorphic kind for flexible arity

-- the only sort, sorts classify kinds
(*), (#), Constraint, OpenKind, AnyK :: BOX
BOX :: BOX
```

All these **Kinds** are built-in and inferred as of GHC 7.10.2.

All of this will be changed with the next GHC 8.0 release.

Language Extensions

Compiler extensions that enable a variety of new functionalities:

- Syntax extension
- Type-level programming
- Generic deriving
- FFI
- Type disambiguation
- Typeclass extension

Each extension has a name, and is enabled with the LANGUAGE pragma.

GADTs

Define data and explicit give type signatures to the **Value** constructors.

```
data Bool = False | True
data Maybe a = Nothing | Just a
data List a = Nil | Cons a (List a)
```

Becomes:

```
{-# LANGUAGE GADTs #-}
data Bool where
  False :: Bool
  True  :: Bool

data Maybe a where
  Nothing :: Maybe a
  Just    :: a -> Maybe a

data List a where
  Nil    :: List a
  Cons   :: a -> List a -> List a
```

GADTs

Define data and explicit give type signatures to the **Value** constructors.

```
data Bool = False | True
data Maybe a = Nothing | Just a
data List a = Nil | Cons a (List a)
```

Loose translations:

```
enum Bool {
  Bool False,
  Bool True
}

enum Maybe<A> {
  Maybe<A> Nothing,
  Maybe<A> Just(A a)
}

enum List<A> {
  List<A> Nil,
  List<A> Cons(A a, List<A> as)
}
```

KindSignatures

Specify the **Kind** of the **Type** variables:

```
{-# LANGUAGE GADTs #-}  
{-# LANGUAGE KindSignatures #-}  
data Bool :: * where  
  False :: Bool  
  True  :: Bool  
  
data Maybe :: * -> * where  
  Nothing :: Maybe a  
  Just    :: a -> Maybe a  
  
data List :: * -> * where  
  Nil    :: List a  
  Cons   :: a -> List a -> List a
```


DataKinds

Kinds are built-in; no user defined Kinds.

Want Values at the Type level though!

=> Data kind promotion :)

DataKinds

Example:

```
data Bool = False | True
```

With DataKinds, we get something like:

```
{- # LANGUAGE DataKinds #-}
```

Kind			Bool
Type	Bool		'True 'False
Value	True	False	

DataKinds

Example:

```
data Nat = Z | S Nat
```

With DataKinds, we get something like:

```
{- # LANGUAGE DataKinds #-}
```

Kind		Nat
Type	Nat	'Z 'S Nat
Value	Z S Nat	

Type Families

Type families - type level functions, computed and checked at compile time.

Comes in 2 flavors:

- type synonym families
- data families

and have a few options:

- associated vs. unassociated
- open vs. closed¹
- injectivity²

Type Families

At **Value** level:

```
data Nat = Z | S Nat
```

```
add :: Nat -> Nat -> Nat
```

```
add Z      m = m
```

```
add (S n) m = S (add n m)
```

```
    add (S (S Z)) (S Z)
```

```
=> S (add (S Z) (S Z))
```

```
=> S (S (add Z (S Z)))
```

```
=> S (S (S Z))
```

Type Families

At **Type** level:

```
{-# LANGUAGE DataKinds #-}
{-# LANGUAGE TypeFamilies #-}
```

```
data Nat = Z | S Nat
```

```
type family Add (n :: Nat) (m :: Nat) :: Nat where
  Add 'Z      m = m
  Add ('S n) m = 'S (Add n m)
```

```
    Add ('S ('S 'Z)) ('S 'Z)
=> 'S (Add ('S 'Z) ('S 'Z))
=> 'S ('S (Add 'Z ('S 'Z)))
=> 'S ('S ('S 'Z))
```

Type Operators

Allows usage of symbols in place of **Type** constructors and **Type** families.

```
{-# LANGUAGE DataKinds #-}
{-# LANGUAGE TypeFamilies #-}
{-# LANGUAGE TypeOperators #-}
```

```
data Nat = Z | S Nat
```

```
type family ('+) n m where
  'Z      '+' m = m
  ('S n) '+' m = 'S (n '+' m)
```

```
      ('S ('S 'Z)) '+' ('S 'Z)
=> 'S (('S 'Z) '+' ('S 'Z))
=> 'S ('S ('Z '+' ('S 'Z)))
=> 'S ('S ('S 'Z))
```

Vectors

Like `List`, but also *indexed* by `Nat` to indicate length.

`List`:

```
data List a where
  Nil  :: List a
  Cons :: a -> List a -> List a
```

`Vector`:

```
-- 'Z ~ 0
-- 'S n ~ n '+ 1
data Vect (n :: Nat) a where
  VNil  :: Vect Z a
  (:>) :: a -> Vect n a -> Vect (S n) a

type Six = S (S (S (S (S (S Z))))))
vs :: Vect Six Int
vs = 4 :> 8 :> 15 :> 16 :> 23 :> 42 :> VNil
```


Vectors

Like `List`, but also indexed by `Nat` to indicate length.

`List`:

```
data List a where
  Nil    :: List a
  Cons   :: a -> List a -> List a
```

Module `GHC.TypeLits` provide type-level literals:

```
-- 'Z ~ 0
-- 'S n ~ n '+ 1
data Vect (n :: Nat) a where
  VNil    :: Vect 0 a
  (:>)    :: a -> Vect n a -> Vect (n '+ 1) a

vs :: Vect 6 Int
vs = 4 :> 8 :> 15 :> 16 :> 23 :> 42 :> VNil
```

Head

head returns the first element of the [List](#):

```
-- from standard library
-- useless unless we know the list is non-empty
head :: [a] -> a
head []      = error "empty list"
head (x:xs) = x
```

Head

head returns the first element of the [List](#):

```
-- from standard library
-- useless unless we know the list is non-empty
head :: [a] -> a
head []      = error "empty list"
head (x:xs) = x
```

Elm now uses [Maybe](#):

```
mhead :: [a] -> Maybe a
mhead []      = Nothing
mhead (x:xs) = Just x
```

Head

head returns the first element of the [List](#):

```
-- from standard library
-- useless unless we know the list is non-empty
head :: [a] -> a
head []      = error "empty list"
head (x:xs) = x
```

Elm now uses [Maybe](#):

```
mhead :: [a] -> Maybe a
mhead []      = Nothing
mhead (x:xs) = Just x
```

With [Vector](#):

```
vhead :: Vect (S n) a -> a
vhead (x:>xs) = x
```

Append

append concatenates 2 Lists:

```
data [a] = [] | (:) a [a]
```

```
append :: [a] -> [a] -> [a]
```

```
append []      ys = ys
```

```
append (x:xs)  ys = x : append xs ys
```

Evaluation is a series of substitutions:

```
xs = [4, 8] = 4 : 8 : [] :: [Int]
```

```
ys = [15, 16, 23, 42] = 15 : 16 : 23 : 42 : [] :: [Int]
```

```
append xs ys :: [Int]
= append [4, 8] [15, 16, 23, 42]
= 4 : append [8] [15, 16, 23, 42]
= 4 : 8 : append [] [15, 16, 23, 42]
= 4 : 8 : [15, 16, 23, 42]
= 4 : [8, 15, 16, 23, 42]
= [4, 8, 15, 16, 23, 42]
```

Append

append concatenates 2 Lists:

```
append :: [a] -> [a] -> [a]
append []    ys = ys
append (x:xs) ys = x : append xs ys
```

With Vector:

```
vappend :: Vect n a -> Vect m a -> Vect (n + m) a
vappend VNil    ys = ys
vappend (x:>xs) ys = x :> vappend xs ys
```

Map

map:

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (x:xs) = f x : map f xs
```

Evaluation is a series of substitutions:

```
xs = [4, 8, 15, 16, 23, 42] :: [Int]
even :: Int -> Bool
```

```
map even xs :: [Bool]
= map even [4, 8, 15, 16]
= even 4 : map even [8, 15, 16]
= True : even 8 : map even [15, 16]
= True : True : even 15 : map even [16]
= True : True : False : even 16 : map even []
= True : True : False : True : []
= [True, True, False, True]
```

Map

map maps a function over a List:

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (x:xs)      = f x : map f xs
```

With Vector:

```
vmap :: (a -> b) -> Vect n a -> Vect n b
vmap f VNil       = VNil
vmap f (x:>xs)    = f x :> vmap f xs
```


Zip

zip:

```
zip :: [a] -> [b] -> [(a,b)]
zip []      ys      = []
zip xs      []      = []
zip (x:xs) (y:ys) = (x,y) : zip xs ys
```

Evaluation is a series of substitutions:

```
xs = ['a', 'b', 'c'] :: [Char]
ys = [1, 2, 3, 4] :: [Int]

zip xs ys :: [(Char, Int)]
= zip ['a', 'b', 'c'] [1, 2, 3, 4]
= ('a', 1) : zip ['b', 'c'] [2, 3, 4]
= ('a', 1) : ('b', 2) : zip ['c'] [3, 4]
= ('a', 1) : ('b', 2) : ('c', 3) : zip [] [4]
= ('a', 1) : ('b', 2) : ('c', 3) : []
= [('a', 1), ('b', 2), ('c', 3)]
```

Zip

zip creates pair-wise tuples:

```
zip :: [a] -> [b] -> [(a,b)]  
zip (x:xs) (y:ys) = (x,y) : zip xs ys  
zip xs         ys         = []
```

With **Vector**:

```
vzip :: Vect n a -> Vect n b -> Vect n (a, b)  
vzip (x:>xs) (y:>ys) = (x,y) :> vzip xs ys  
vzip VNil     VNil     = VNil
```

Zip

zip2 with **Min** type family:

```
type family Min n m where
  Min Z      m      = Z
  Min n      Z      = Z
  Min (S n) (S m) = S (Min n m)

vzip2 :: Vect n a -> Vect m b -> Vect (Min n m) (a, b)
vzip2 (x:>xs) (y:>ys) = (x,y) :> vzip2 xs ys
vzip2 xs      VNil    = VNil
vzip2 VNil    ys      = VNil
```

Replicate

replicate repeats an element n times:

```
replicate :: Int -> a -> [a]
replicate 0 x = []
replicate n x = x : replicate (n - 1) x
```

For `Vector`, we want its `Type Vector n a` to reflect the `Value` of the first parameter `Int`.

Pi Types

Π -types - **Values** in **Type** signatures. Fake by deriving singleton instances of **Sing** data family to reflect values to the type level.

```
-- fake with singleton types
data instance Sing (n :: Nat) where
  SZ :: Sing Z
  SS :: Sing n -> Sing (S n)

-- so we have
SZ :: Sing Z
SS SZ :: Sing (S Z)
SS (SS SZ) :: Sing (S (S Z))
```

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-- so we have
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SS SZ :: Sing (S Z)
SS (SS SZ) :: Sing (S (S Z))

-- finally:
vreplicate :: Sing (n :: Nat) -> a -> Vect n a
vreplicate SZ      a = VNil
vreplicate (SS n) a = a :> vreplicate n a
```

Filter

filter selects elements from a list for given predicate:

```
filter :: (a -> Bool) -> [a] -> [a]
filter f []      = []
filter f (x:xs) = if f x
                  then x : filter f xs
                  else filter f xs
```

We do not the number of elements for which the predicate will return **True**, but **Vector** depends on this number in its **Type**.

Sigma Types

Σ -types - tuple where 2nd value depends on 1st:

```
-- using Idris 's ** dependent pair syntax
(3 ** 'a' :> 'b' :> 'c' :> VNil)
:: (n :: Nat ** Vect n Char)

vfilter :: (a -> Bool) -> Vect n a
        -> (p :: Nat ** Vect p a)
```


Sigma Types

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        -> (p :: Nat ** Vect p a)
```

Credit to Ertugrul Söylemez:

```
-- Sugar ** for Sigma type constructor and Exists value
  constructor
data Sigma :: KProxy a -> (a -> *) -> * where
  (Exists :: Sing (x :: a) -> b x
   -> Sigma ('KProxy :: KProxy a) b
```

Heterogeneous List

Heterogeneous List indexed by List of Types:

```
infixr 5 ::>
data HList (t :: [*]) where
  HNil    :: HList '[]
  (::>) :: t -> HList ts -> HList (t ': ts)

defaults :: HList '[Int, Bool, Maybe a]
defaults = 0 ::> False ::> Nothing ::> HNil
```

Heterogeneous Vector

Heterogeneous **Vector** indexed by a **List** of **Types**:

```
infixr 5 :>>
data HVect (n :: Nat) (t :: [*]) where
  HVNil    :: HVect Z '[]
  (:>>)    :: t -> HVect n ts -> HVect (S n) (t ': ts)

vdefaults :: HVect 3 '[Int, Bool, Maybe a]
vdefaults = 0 :>> False :>> Nothing :>> HVNil
```

Section Outline

- 5 Closing
 - Beyond
 - Questions

Beyond Dependent Types

- Total functional languages
 - termination and totality check
 - disallow partial functions
 - distinction between **data** and **codata**
- Proof assistant languages
 - Ph.D. first please

Questions?