

Introduction to Dependent Types

Eagan Technology Unconference

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Section Outline

1 Preface

Quick Question

How many are familiar with this topic?

A Joke

This is not a $\mathsf{m-}$ tutorial.

A Joke

This is not a `m-` tutorial.
Nor is it a `lens` tutorial

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This is not a `m-` tutorial.

Nor is it a `lens` tutorial (aka the new new `m-` tutorial...)

A Joke

This is not a `m-` tutorial.

Nor is it a `lens` tutorial (aka the new new `m-` tutorial...

...because arrows *were* the new `m-` tutorials).

About This Talk

Agda, Idris, Coq and co^* have full support for dependent types.

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Because of that, it's harder to see the build up, so we won't be directly using them in this talk.

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Agda, Idris, Coq and co^* have full support for dependent types.

Because of that, it's harder to see the build up, so we won't be directly using them in this talk.

Honestly though, it's because they're way over my head :(

() There was another mini joke here. . .*

About This Talk

But we will be using Haskell though :)

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It's not truly dependent, but we can do more and more with each language extension that comes along.

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But we will be using Haskell though :)

It's not truly dependent, but we can do more and more with each language extension that comes along.

For the examples, there will also be *very* loose translation to imperative/OOP. Though please keep in mind that these are merely syntax translations, the actual concepts can differ vastly.

Section Outline

- 2 Review of Basics
 - Test
 - Values and Types
 - Defining Data Types
 - Functions

Test

Syntax highlighting test reference, to be removed later.

```
-- Comment
data Maybe a = Nothing | Just a
               deriving (Show, Eq)

fmap :: Functor f => (a -> b) -> f a -> f b
map _ []          = []
map f (x:xs) = f x : map f xs

type family TF a :: *
type instance TF Int = Bool
```

Test

Couldn't quite yet get listing to work with overlay yet.

```
{- block comment -}  
foo :: Bool -> Int -> String  
foo False 0 = "Bad"  
foo True 0 = "Questionable"  
foo False n = "Fake"  
foo True n = "Read"
```

Test

Pausing within listing is ok?

```
{-# LANGUAGE KitchenSink #-}  
zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]
```

Test

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Test

Pausing within listing is ok?

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{-# LANGUAGE KitchenSink #-}  
zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]  
zipWith _ [] _ = []  
zipWith _ _ [] = []
```


Test

Pausing within listing is ok?

```
{-# LANGUAGE KitchenSink #-}
zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]
zipWith _ [] _ = []
zipWith _ _ [] = []
zipWith f (x:xs) (y:ys) = f x y : zipWith f xs ys
```

better yet

```
{-# LANGUAGE KitchenSink #-}
zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]
zipWith f (x:xs) (y:ys) = f x y : zipWith f xs ys
zipWith _ _ _ = []
```

Values and Types

Values has Types, or Values are classified by Types.

```
..., -1, 0, 1, 2, 3, ... :: Int
```

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True, False :: Bool
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```
'a', 'b', 'c' :: Char
```

Values and Types

Values has Types, or Values are classified by Types.

```
..., -1, 0, 1, 2, 3, ... :: Int
```

```
True, False :: Bool
```

```
'a', 'b', 'c' :: Char
```

```
"abc" :: String ~ [Char]
```

Values are also called Terms

About Types

How are data types defined?

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- Some are built in magic: `Int`, `Char`, functions

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 - We can define equivalent non-sugar-ed version ourselves

About Types

How are data types defined?

- Some are built in magic: `Int`, `Char`, functions
- Some are built in sugar: `list`, `tuples`
 - We can define equivalent non-sugar-ed version ourselves
- Rest can be user defined: `Bool`, `String`, `Maybe`

About Types

What are the data types like?

About Types

What are the data types like?

- Multiple **Value** constructors

About Types

What are the data types like?

- Multiple **Value** constructors
- Parametrize over another type

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- Recursive definition

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- Synonyms of other types

About Types

What are the data types like?

- Multiple **Value** constructors
- Parametrize over another type
- Recursive definition
- Synonyms of other types
- A combination of the above

Defining Data Types

Define new data type with `data`.

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- Left hand side (`LHS`) - `Type` constructor
- Right hand side (`RHS`) - `Value` constructor

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- Left hand side (`LHS`) - `Type` constructor
- Right hand side (`RHS`) - `Value` constructor

`Type` and `Value` constructors are capticalized.

Our First Example!

Define a person:

```
-- | params for firstname , lastname , age respectively  
data Person = Person String String Int
```

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A loose translation:

```
enum Person {  
  Person(String firstname, String lastname, Int age)  
}
```

Our First Example!

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data Person = Person String String Int
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A loose translation:

```
enum Person {  
  Person(String firstname, String lastname, Int age)  
}
```

In this example, the **Type** and **Value** constructor have the same name. The **Type** of the **Person** constructor:

```
Person :: String -> String -> Int -> Person
```

```
bobby :: Person
```

```
bobby = Person "Bobby" "Smith" 23
```

Our First Example!

Define a person:

```
-- | params for firstname, lastname, age respectively
data Person = Person String String Int
```

A loose translation:

```
enum Person {
  Person(String firstname, String lastname, Int age)
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In this example, the **Type** and **Value** constructor have the same name. The **Type** of the **Person** constructor:

```
Person :: String -> String -> Int -> Person
```

```
bobby :: Person
```

```
bobby = Person "Bobby" "Smith" 23
```

```
-- a loose translation:
```

```
Person bobby = new Person("Bobby", "Smith", 23)
```

Multiple Value Constructors

Data can have multiple **Value** constructors:

```
data Bool = False | True
```

```
data Weekdays = Sunday | Monday | Tuesday | Wednesday  
               | Thursday | Friday | Saturday
```

Does this remind you of anything?

Multiple Value Constructors

Data can have multiple **Value** constructors:

```
data Bool = False | True
```

```
data Weekdays = Sunday | Monday | Tuesday | Wednesday  
               | Thursday | Friday | Saturday
```

Does this remind you of anything?

A loose translation:

```
enum Bool { False, True }
```

```
enum Weekdays {  
    Sunday, Monday, Tuesday, Wednesday, Thursday, Friday,  
    Saturday  
}
```


Multiple Value Constructor

You can do type aliasing with `type`:

```
type Side = Double  
type Radius = Double
```

Multiple Value Constructor

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For example:

```
data Shape = Triangle Side Side Side  
          | Rectangle Side Side  
          | Circle Radius
```

Multiple Value Constructor

You can do type aliasing with `type`:

```
type Side = Double
type Radius = Double
```

For example:

```
data Shape = Triangle Side Side Side
           | Rectangle Side Side
           | Circle Radius
```

A loose translation:

```
enum Shape {
  Triangle(Double side1, Double side2, Double side3),
  Rectangle(Double length, Double width),
  Circle(Double radius)
}
```

Multiple Value Constructor

Recall $\text{Side} \sim \text{Radius} \sim \text{Double}$:

```
data Shape = Triangle Side Side Side
           | Rectangle Side Side
           | Circle Radius
```

Multiple Value Constructor

Recall $\text{Side} \sim \text{Radius} \sim \text{Double}$:

```
data Shape = Triangle Side Side Side
           | Rectangle Side Side
           | Circle Radius
```

Types of the **Value** constructors:

```
Triangle  :: Side -> Side -> Side -> Shape
Rectangle :: Side -> Side -> Shape
Circle    :: Radius -> Shape
```

Multiple Value Constructor

Recall $\text{Side} \sim \text{Radius} \sim \text{Double}$:

```
data Shape = Triangle Side Side Side
           | Rectangle Side Side
           | Circle Radius
```

Types of the **Value** constructors:

```
Triangle  :: Side -> Side -> Side -> Shape
Rectangle :: Side -> Side -> Shape
Circle    :: Radius -> Shape
```

Example **Shapes**:

```
myTri, myRect, myCir :: Shape
myTri  = Triangle 2.1 3.2 5
myRect = Rectangle 4 4
myCir  = Circle 7.2
```

Parametrization

Types can parametrize over another type:

```
data Identity a = Identity a
```

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data Identity a = Identity a
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A loose translation:

```
enum Identity<A> {  
  Identity(A a)  
}
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Types can parametrize over another type:

```
data Identity a = Identity a
```

A loose translation:

```
enum Identity<A> {  
  Identity(A a)  
}
```

The `Type` of the `Identity` constructor:

```
Identity :: a -> Identity a
```

```
intIdwrtSum :: Identity Int  
intIdwrtSum = Identity 0
```

Tuple

Parametrize over 2 types - 2-tuple!

```
data Tuple a b = Tuple a b
```

Tuple

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```
data Tuple a b = Tuple a b
```

A loose translation:

```
enum Tuple<A, B> {  
  Tuple(A a, B b)  
}
```

Tuple

Parametrize over 2 types - 2-tuple!

```
data Tuple a b = Tuple a b
```

A loose translation:

```
enum Tuple<A, B> {  
  Tuple(A a, B b)  
}
```

With:

```
Tuple :: a -> b -> Tuple a b
```

Tuple

Actual built-in sugar:

```
data Tuple a b = Tuple a b
=> data (,) a b = (,) a b
=> data (a, b) = (a, b)
```

Tuple

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```
data Tuple a b = Tuple a b
=> data (,) a b = (,) a b
=> data (a, b) = (a, b)
```

An example:

```
type Employed = Bool

barbara, chet, luffy :: (Person, Employed)
barbara = (Person "Barbara" "Sakura" 30, True)
chet    = (Person "Chet" "Awesome-Laser" 2, False)
luffy   = (Person "Luff D." "Monkey" 19, False)
```

Maybe

Like `Bool`, but parametrize an `a` over the `True` part:

```
data Maybe a = Nothing | Just a
```

Maybe

Like `Bool`, but parametrize an `a` over the `True` part:

```
data Maybe a = Nothing | Just a
```

A loose translation:

```
enum Maybe<A> {  
  Nothing,  
  Just(A a)  
}
```


Maybe

Like `Bool`, but parametrize an `a` over the `True` part:

```
data Maybe a = Nothing | Just a
```

A loose translation:

```
enum Maybe<A> {  
  Nothing,  
  Just(A a)  
}
```

The `Types` of the two `Value` constructors:

```
Nothing :: Maybe a  
Just    :: a -> Maybe a
```

Maybe

From previous slide:

```
data Maybe a = Nothing | Just a
```

Maybe

From previous slide:

```
data Maybe a = Nothing | Just a
```

Say more with the occupation:

```
type Occupation = Maybe String
```

```
barbara, chet, luffy :: (Person, Occupation)  
barbara = (Person "Barbara" "Sakura" 30, Just "dancer")  
chet    = (Person "Chet" "Awesome-Laser" 2, Nothing)  
luffy   = (Person "Luff D." "Monkey" 19, Just "pirate")
```

Either

Like `Bool`, but parametrize over both `True` and `False`:

```
data Either a b = Left a | Right b
```

Either

Like `Bool`, but parametrize over both `True` and `False`:

```
data Either a b = Left a | Right b
```

A loose translation:

```
enum Either<A, B> {  
  Left(A a),  
  Right(B b)  
}
```

Either

Like `Bool`, but parametrize over both `True` and `False`:

```
data Either a b = Left a | Right b
```

A loose translation:

```
enum Either<A, B> {  
  Left(A a),  
  Right(B b)  
}
```

The two `Value` constructors have `Types`:

```
Left  :: a -> Either a b  
Right :: b -> Either a b
```

Either

From previous slide:

```
data Either a b = Left a | Right b
```

Either

From previous slide:

```
data Either a b = Left a | Right b
```

Refine with more details:

```
type Earning = Either String Int
```

```
barbara, chet, luffy :: (Person, Earning)
barbara = (Person "Barbara" "Sakura" 30,
           Right 100000)
chet    = (Person "Chet" "Awesome-Laser" 2,
           Left "Is a baby")
luffy   = (Person "Luff D." "Monkey" 19,
           Right 2000000)
```


Types with Recursion

Natural number:

```
data Nat = Z | S Nat
```

```
Z :: Nat
```

```
S :: Nat -> Nat
```

Types with Recursion

Natural number:

```
data Nat = Z | S Nat
```

```
Z :: Nat
```

```
S :: Nat -> Nat
```

A loose translation:

```
enum Nat {  
  Z,  
  S(Nat n)  
}
```

Types with Recursion

Natural number:

```
data Nat = Z | S Nat
```

```
Z :: Nat
```

```
S :: Nat -> Nat
```

```
0 ~ Z
```

```
1 ~ S Z
```

```
2 ~ S (S Z)
```

```
3 ~ S (S (S Z))
```

Types with Recursion

List - recursive type while parametrize over another type:

```
data List a = Nil | Cons a (List a)
```

```
Nil :: List a
```

```
Cons :: a -> List a -> List a
```

Types with Recursion

List - recursive type while parametrize over another type:

```
data List a = Nil | Cons a (List a)
```

```
Nil :: List a
```

```
Cons :: a -> List a -> List a
```

A loose translation:

```
enum List<A> {  
  Nil,  
  Cons(A a, List<A> as)  
}
```

Types with Recursion

Actual built-in sugared version is something like:

```
data List a = Nil | Cons a (List a)
=> data [] a = [] | (:) a ([] a)
=> data [a] = [] | a : [a]
```

Types with Recursion

Actual built-in sugared version is something like:

```
data List a = Nil | Cons a (List a)
=> data [] a = [] | (:) a ([] a)
=> data [a] = [] | a : [a]
```

De-sugar that list:

```
ints :: List Int
ints = Cons 1 (Cons 2 (Cons 3 (Cons 4 Nil)))

-- built-in sugar
ints :: [] Int
ints = 1 : 2 : 3 : 4 : []

-- 2x the sugar!
ints :: [Int]
ints = [1, 2, 3, 4]
```

Functions

Maps **Values** of a **Type** to another **Type**:

```
even :: Int -> Bool
even 0 = True
even n = if rem n 2 == 0
         then True
         else False
```


Functions

Maps **Values** of a **Type** to another **Type**:

```
even :: Int -> Bool
even 0 = True
even n = if rem n 2 == 0
         then True
         else False
```

Not as loose translation:

```
static Bool even (Int n) {
  switch n:
    case n == 0:
      return True;
  default:
    if rem@(n, 2) == 0
      return True;
    else
      return False;
}
```

Functions with Recursion

Use recursion for recursive types:

```
toInt :: Nat -> Int
toInt Z    = 0
toInt (S n) = 1 + toInt n
```

Functions with Recursion

Use recursion for recursive types:

```
toInt :: Nat -> Int
toInt Z   = 0
toInt (S n) = 1 + toInt n
```

Not as loose translation:

```
static Int toInt (Nat n) {
  switch n:
    case Z:
      return 0;
    case (S m): -- n ~ (S m)
      return 1 + toInt(m);
}
```

Functions with Parametric Polymorphism

Functions can be parametric:

```
id :: a -> a  
id a = a
```

Functions with Parametric Polymorphism

Functions can be parametric:

```
id :: a -> a  
id a = a
```

Not as loose translation:

```
static A id<A>(A a) {  
    return a;  
}
```

Functions with Parametric Polymorphism

Functions can be parametric:

```
append :: [a] -> [a] -> [a]
append []      ys = ys
append (x:xs) ys = x : append xs ys
```

Functions with Parametric Polymorphism

Functions can be parametric:

```
append :: [a] -> [a] -> [a]
append []      ys = ys
append (x:xs)  ys = x : append xs ys
```

A translation:

```
static List<A> append(List<A> l1, List<A> l2) {
  switch l1:
    case Nil:
      return l2;
    case Cons(x, xs):
      List<A> rest = append(xs, l2);
      return Cons(x, rest);
}
```

Higher-order Functions

Functions that take functions as params:

```
( $\$$ ) :: (a -> b) -> a -> b  
f $ x = f x
```

```
(.) :: (b -> c) -> (a -> b) -> (a -> c)  
f . g = \x -> f (g x)
```


Higher-order Functions

Functions that take functions as params:

```
($) :: (a -> b) -> a -> b  
f $ x = f x
```

```
(.) :: (b -> c) -> (a -> b) -> (a -> c)  
f . g = \x -> f (g x)
```

Yay translations:

```
static B apply(Func<A, B> f, A a) {  
    return f(a);  
}
```

```
static Func<A,C> compose(Func<B,C> f, Func<A,B> g) {  
    return x => f(g(x));  
}
```

More Functions Examples

map:

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (x:xs)      = f x : map f xs
```

More Functions Examples

map:

```
map :: (a -> b) -> [a] -> [b]
map f []      = []
map f (x:xs) = f x : map f xs
```

A translation:

```
static List<B> map(Func<A,B> f, List<A> la) {
  switch la:
  case Nil:
    return Nil;
  case Cons(a, as):
    B b = f(a)
    List<B> rest = map(f, as);
    return Cons(b, rest);
}
```

More Functions Examples

zip:

```
zip :: [a] -> [b] -> [(a,b)]
zip []      ys      = []
zip xs      []      = []
zip (x:xs) (y:ys) = (x,y) : zip xs ys
```

More Functions Examples

zip:

```
zip :: [a] -> [b] -> [(a,b)]
zip []      ys      = []
zip xs      []      = []
zip (x:xs) (y:ys) = (x,y) : zip xs ys
```

A translation:

```
static List<Tuple<A,B>> zip(List<A> l1, List<A> l2) {
  switch l1:
    case Nil:
      return Nil;
    case Cons(a, as):
      switch l2:
        case Nil:
          return Nil;
        case Cons(b, bs):
          Tuple<A,B> front = Tuple(a, b);
          List<Tuple<A,B>> rest = zip(as, bs);
          return Cons(front, rest);
}
```

Section Outline

3 What is Dependent Type

Lambda Calculus

So far, we have seen:

- function application

Lambda Calculus

So far, we have seen:

- function application
- function abstraction (aka higher-order functions)

Lambda Calculus

So far, we have seen:

- function application
- function abstraction (aka higher-order functions)
- variable binding

Lambda Calculus

So far, we have seen:

- function application
- function abstraction (aka higher-order functions)
- variable binding
- substitution

Lambda Calculus

So far, we have seen:

- function application
- function abstraction (aka higher-order functions)
- variable binding
- substitution

\Rightarrow basis for λ – *calculus*.

Lambda Cube

Section Outline

- 4 Steps toward Dependent Types
 - Kinds
 - Language Extensions

Kinds

GADTs

KindSignatures

ConstraintKinds

Type Operators

DataKinds

Type Families

Section Outline

5 Questions

Questions?