

Kinetic assumptions in Boolean networks: a case for buffering

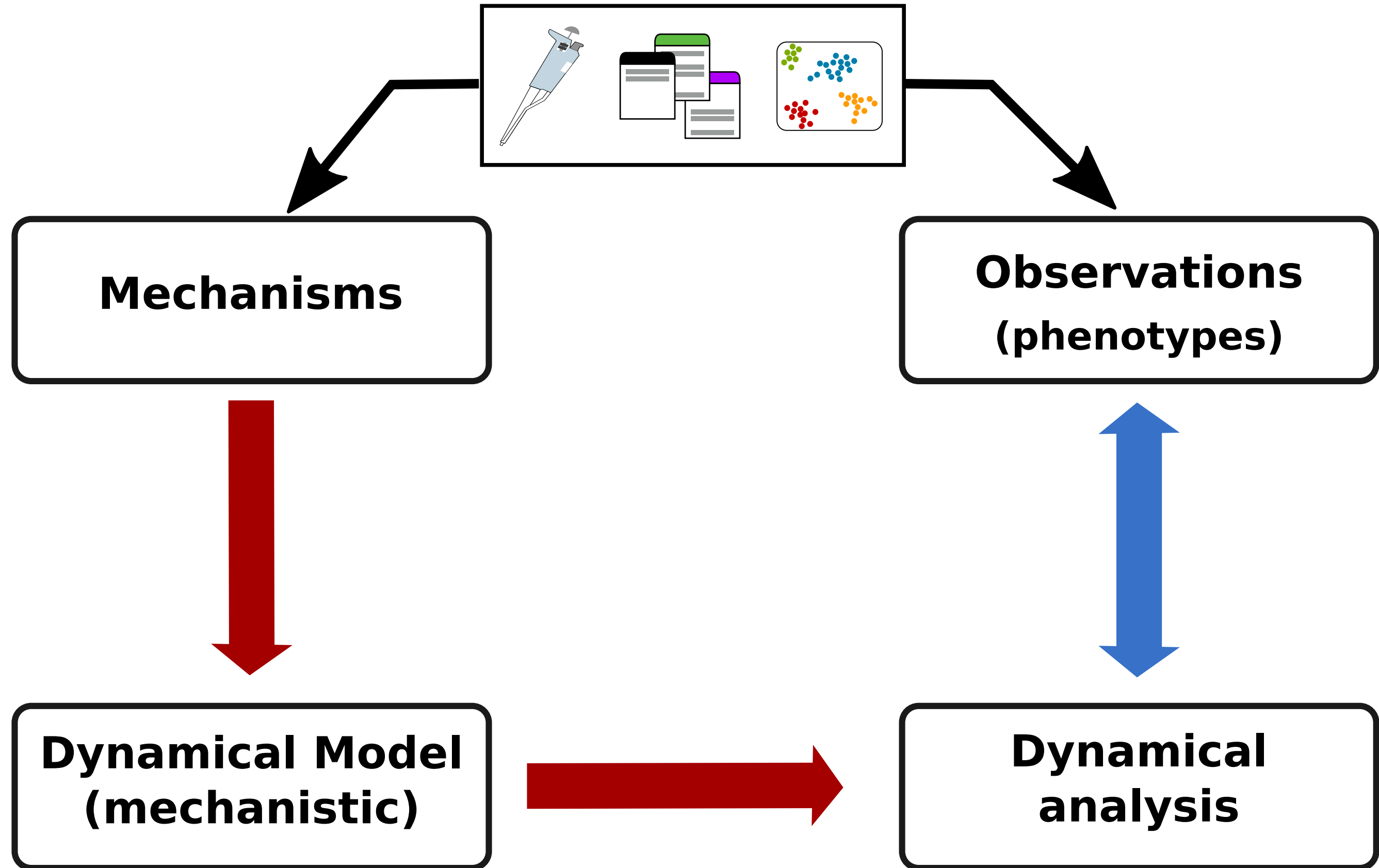
Aurélien Naldi

Inria

Elisa Tonello
Heike Siebert



Qualitative Dynamical Models in Biology



Logical modelling framework

Logical rules

$$f_1(x) = \neg x_2$$

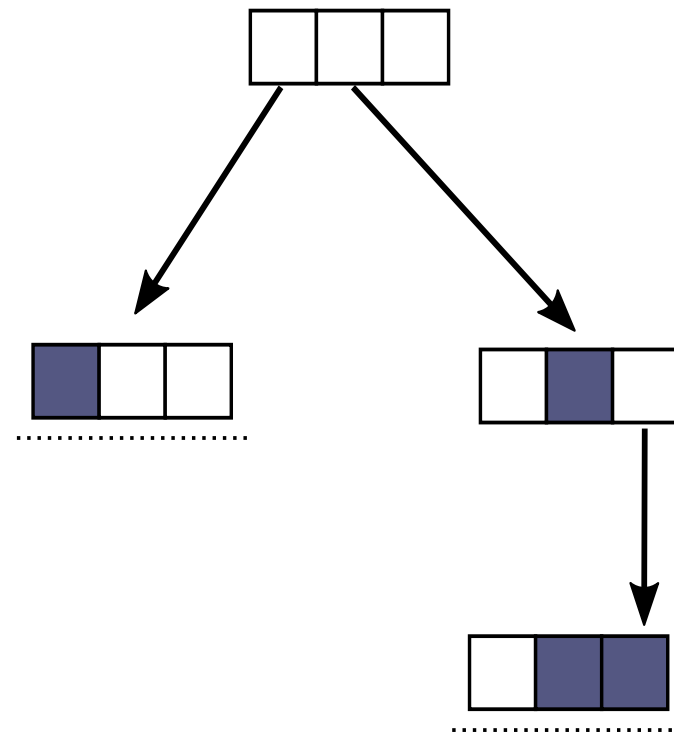
$$f_2(x) = \neg x_1$$

$$f_3(x) = \neg x_1 \wedge x_2$$

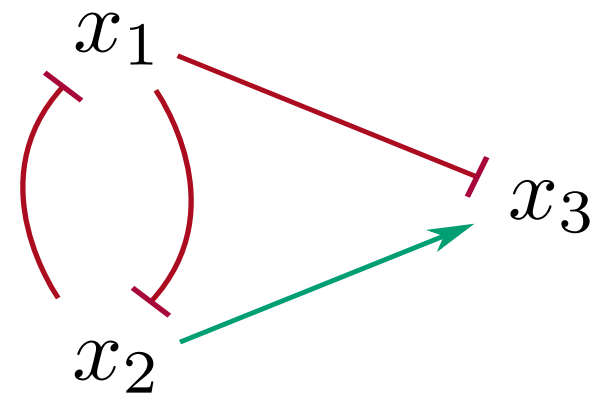


Dynamics as State Transition Graph

$$x = (x_1, x_2, x_3)$$



Regulatory graph



No quantitative parameters

Formal description of competing effects

Updating semantics

x : state

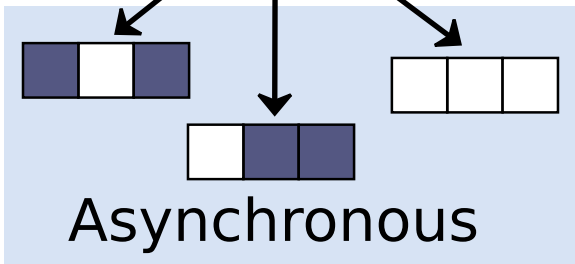
\bar{x}^i : flip i

$f(x) = \bar{x}^U$: image

$U = \{i : f_i(x) \neq x_i\}$

$x \rightarrow f(x)$

Synchronous



Asynchronous

$\forall i \in U: x \rightarrow \bar{x}^i$

Updating semantics

x : state

\bar{x}^i : flip i

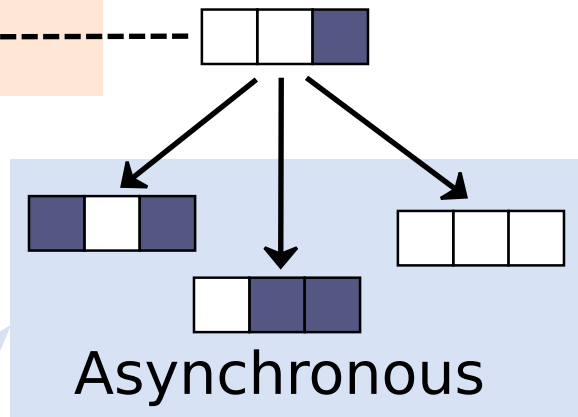
$f(x) = \bar{x}^U$: image

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Everything happens
at the same time

$x \rightarrow f(x)$

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Asynchronous

$\forall i \in U: x \rightarrow \bar{x}^i$

Nothing ever happens
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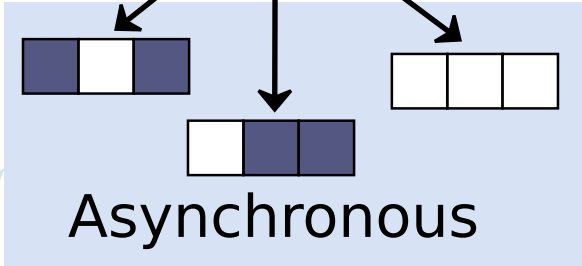
Everything happens at the same time

$$\forall C \subseteq U: x \rightarrow \bar{x}^C$$

Generalized Asynchronous

$$x \rightarrow f(x)$$

Synchronous



$$\forall i \in U: x \rightarrow \bar{x}^i$$

Nothing ever happens at the same time

Updating semantics

x : state

\bar{x}^i : flip i

$f(x) = \bar{x}^U$: image

$U = \{i : f_i(x) \neq x_i\}$

Everything happens at the same time

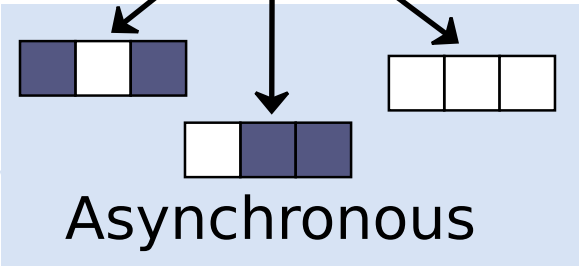
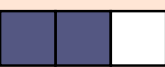
ON/OFF transitions are atomic

$$\forall C \subseteq U: x \rightarrow \bar{x}^C$$

Generalized Asynchronous

$$x \rightarrow f(x)$$

Synchronous



$$\forall i \in U: x \rightarrow \bar{x}^i$$

Nothing ever happens at the same time

Updating semantics

x : state

\bar{x}^i : flip i

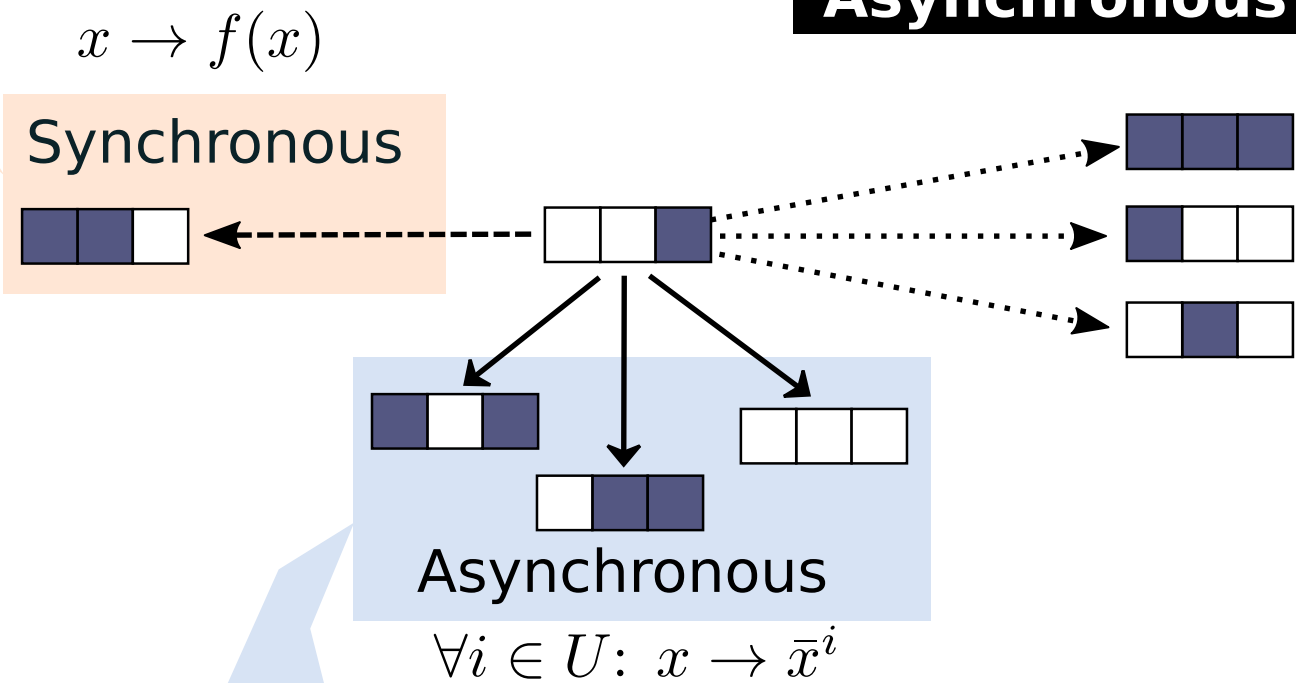
$f(x) = \bar{x}^U$: image

$U = \{i : f_i(x) \neq x_i\}$

Everything happens at the same time

$$\forall C \subseteq U: x \rightarrow \bar{x}^C$$

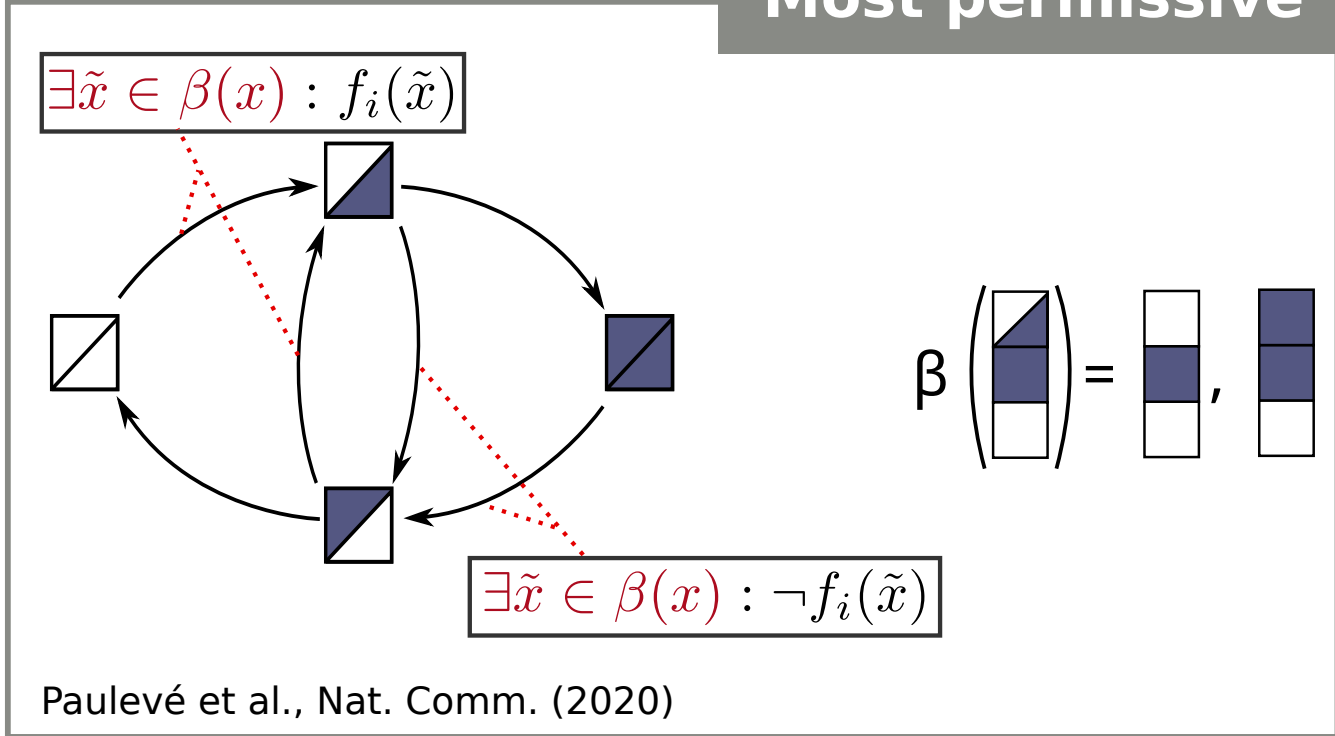
Generalized Asynchronous



Nothing ever happens at the same time

ON/OFF transitions are atomic

Most permissive



Updating semantics

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Everything happens at the same time

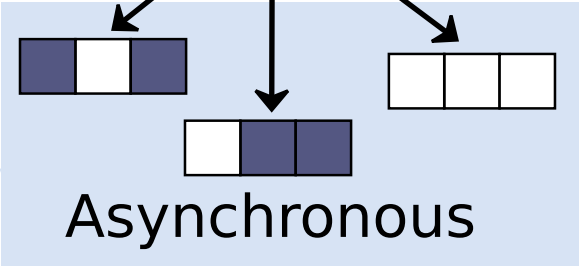
ON/OFF transitions are atomic

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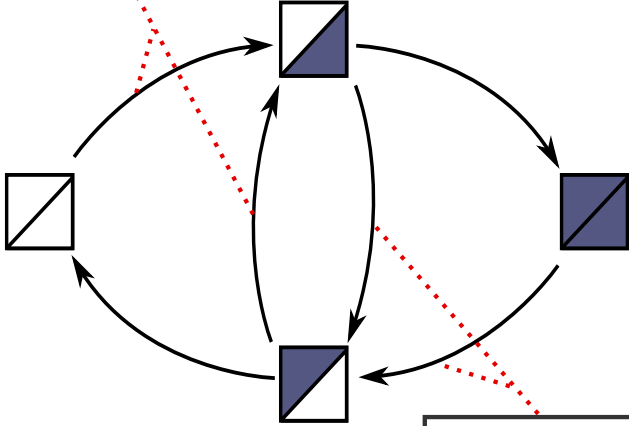
Asynchronous

$$\forall i \in U: x \rightarrow \bar{x}^i$$

Nothing ever happens at the same time

Most permissive

$$\exists \tilde{x} \in \beta(x) : f_i(\tilde{x})$$



$$\exists \tilde{x} \in \beta(x) : \neg f_i(\tilde{x})$$

$$\beta \left(\begin{bmatrix} \text{diagonal} \\ \text{diagonal} \\ \text{diagonal} \end{bmatrix} \right) = \begin{bmatrix} \text{diagonal} \\ \text{diagonal} \\ \text{diagonal} \end{bmatrix}, \begin{bmatrix} \text{diagonal} \\ \text{diagonal} \\ \text{diagonal} \end{bmatrix}$$

Paulevé et al., Nat. Comm. (2020)

Kinetic assumptions

Alternative trajectories

Updating semantics

Huge number of trajectories

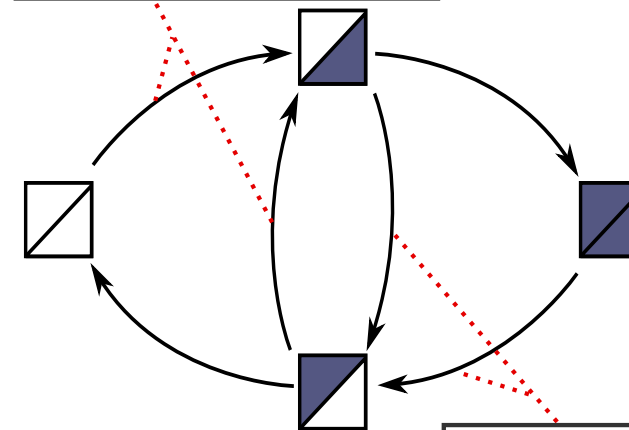
Low computational complexity

- attractors
- reachability

Multivalued refinements

Most permissive

$$\exists \tilde{x} \in \beta(x) : f_i(\tilde{x})$$



$$\exists \tilde{x} \in \beta(x) : \neg f_i(\tilde{x})$$

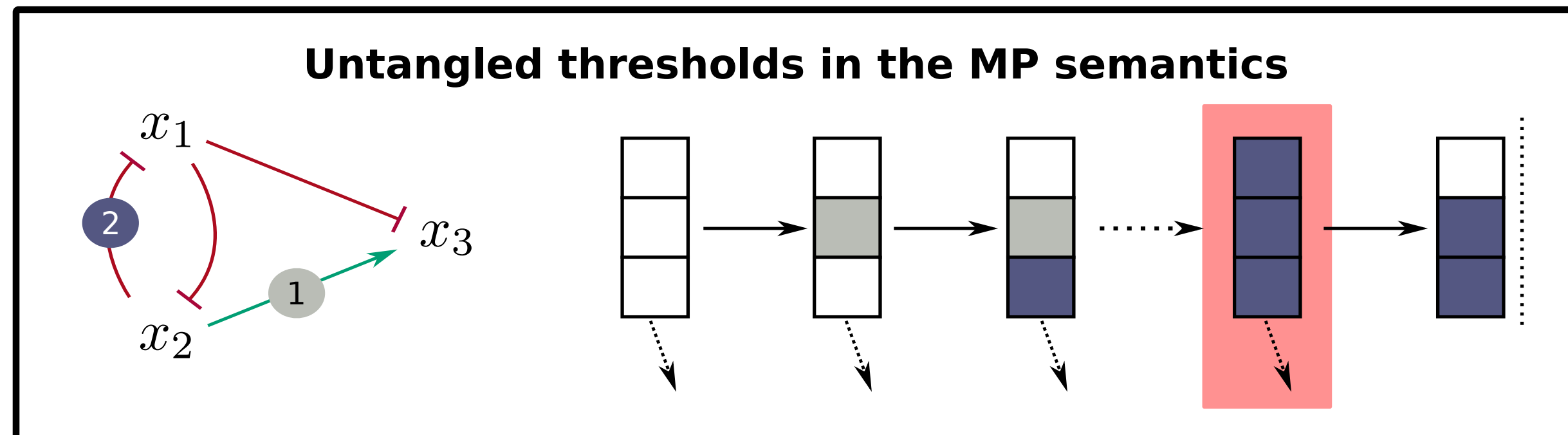
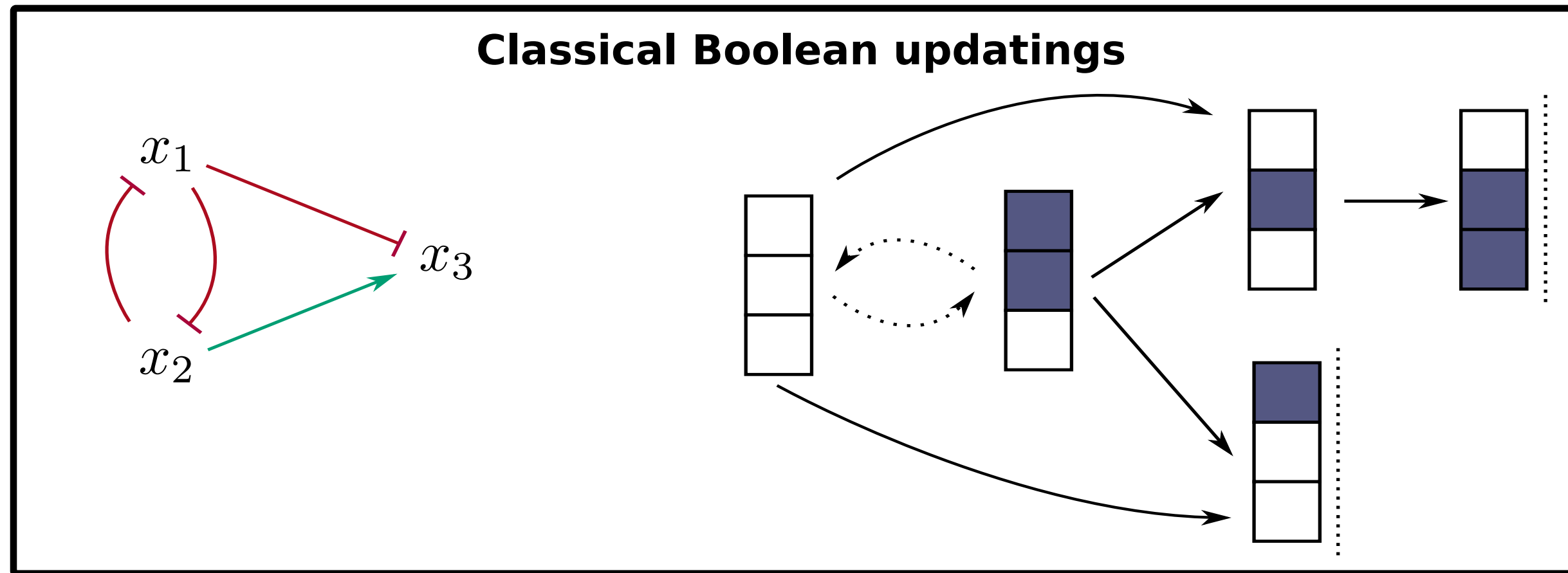
$$\beta \left(\begin{array}{|c|} \hline \text{diagonal} \\ \hline \end{array} \right) = \begin{array}{|c|} \hline \text{white} \\ \hline \end{array}, \begin{array}{|c|} \hline \text{blue} \\ \hline \end{array}$$

Paulevé et al., Nat. Comm. (2020)

Alternative trajectories

Kinetic assumptions

Trajectories with untangled thresholds

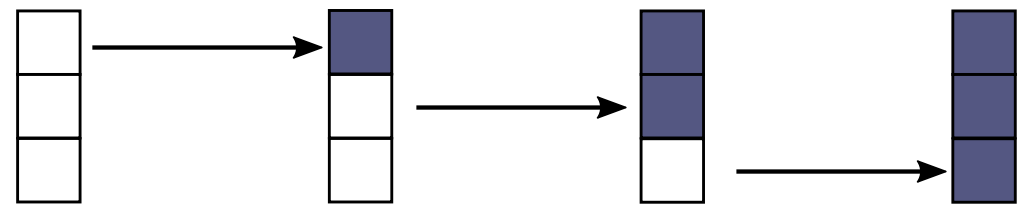


Refinements reveal effect of separate thresholds

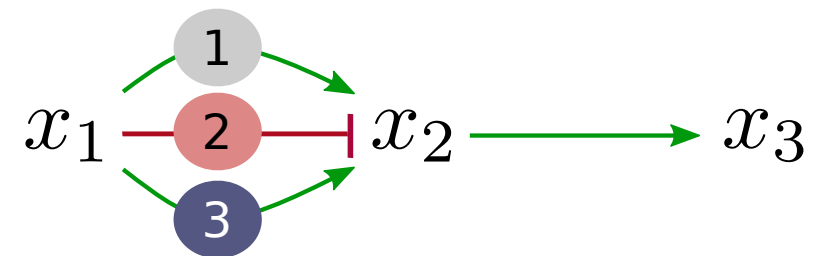
Trajectories with dual interactions

Classical updatings

$$x_1 \longrightarrow x_2 \longrightarrow x_3$$



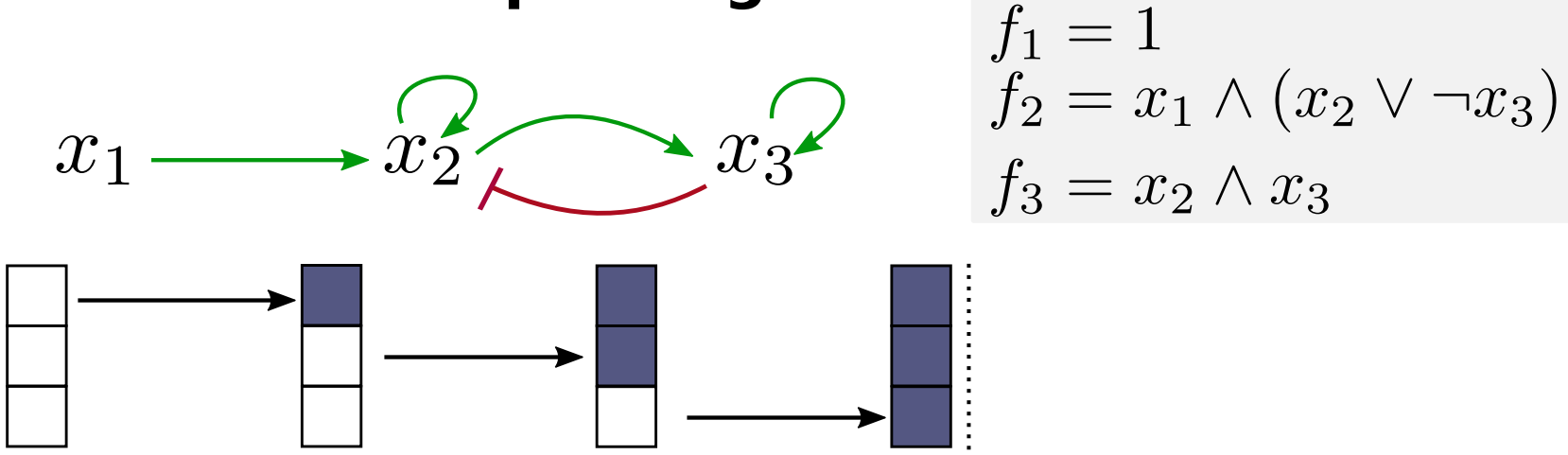
Dual interaction in the MP semantics



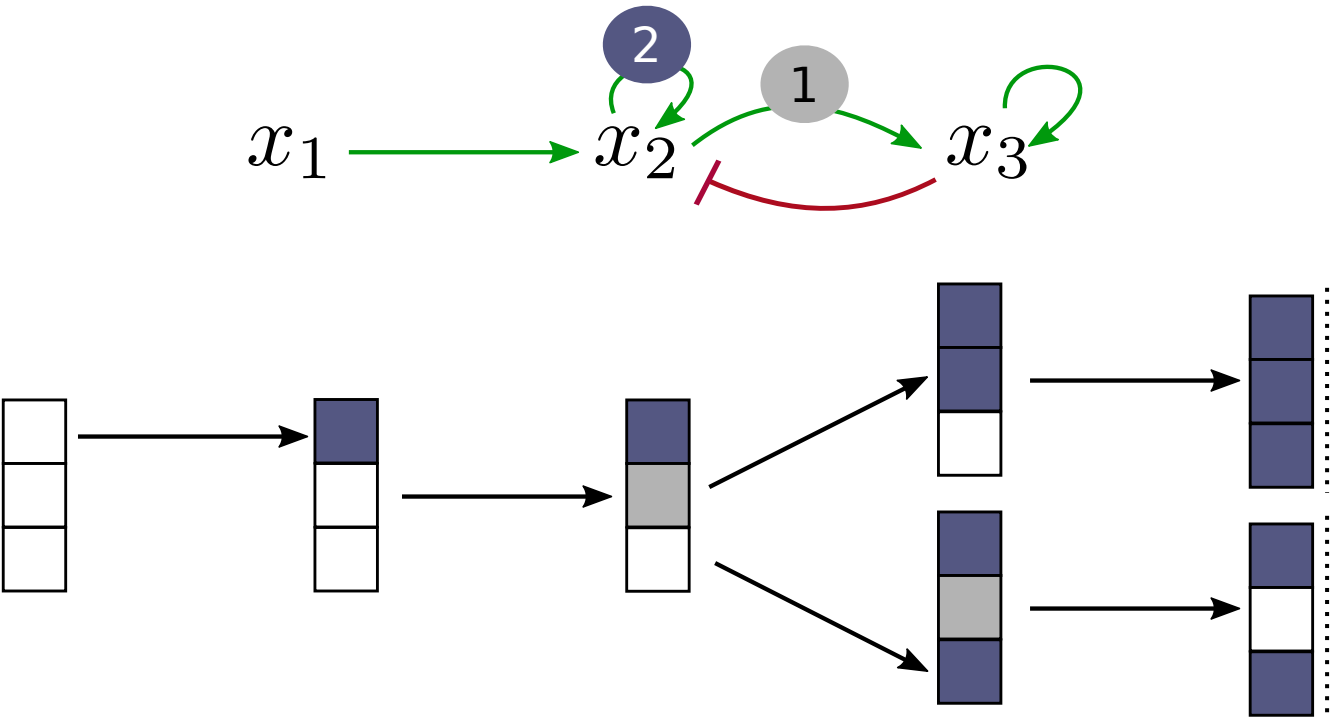
Refinements can introduce dual interactions

Stable states/patterns seem special

Classical updatings

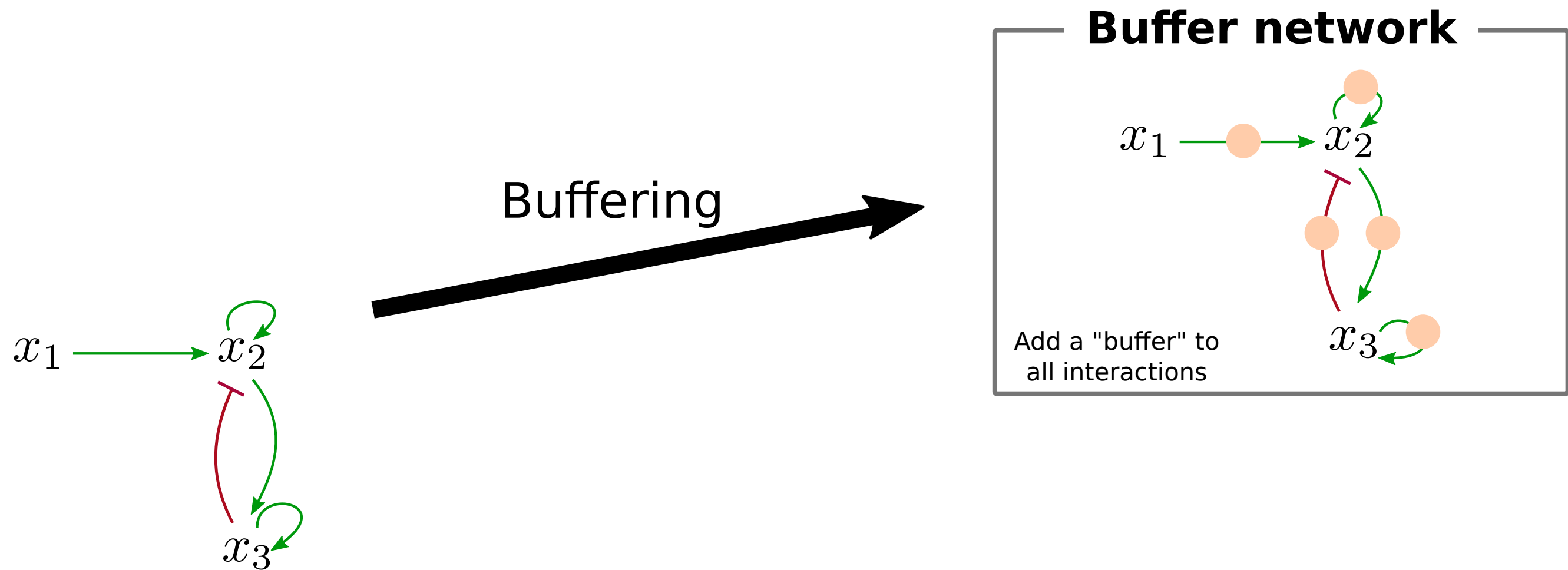


Reachable without dual interaction



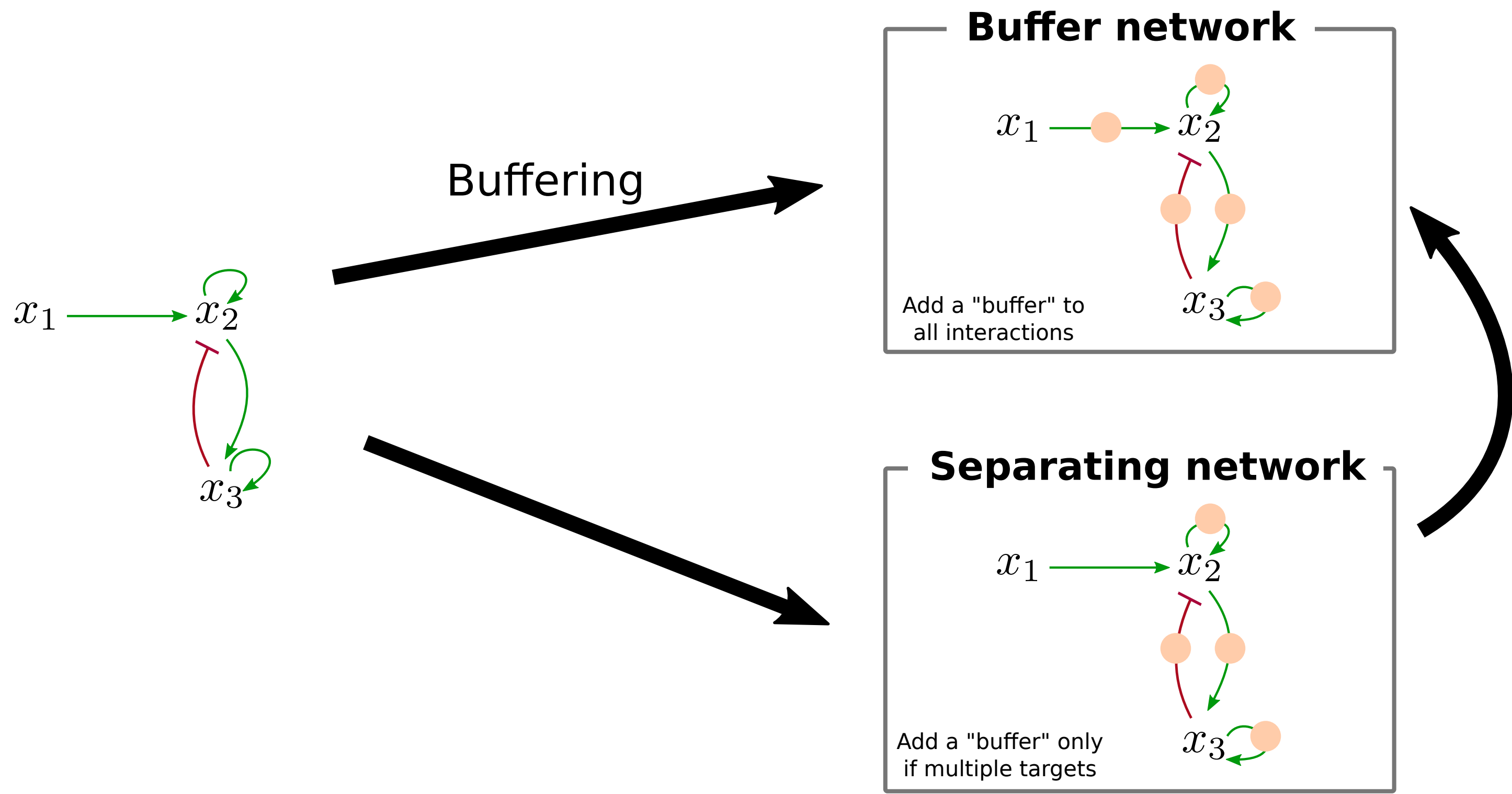
Are stable patterns always easier to reach?

Buffer networks



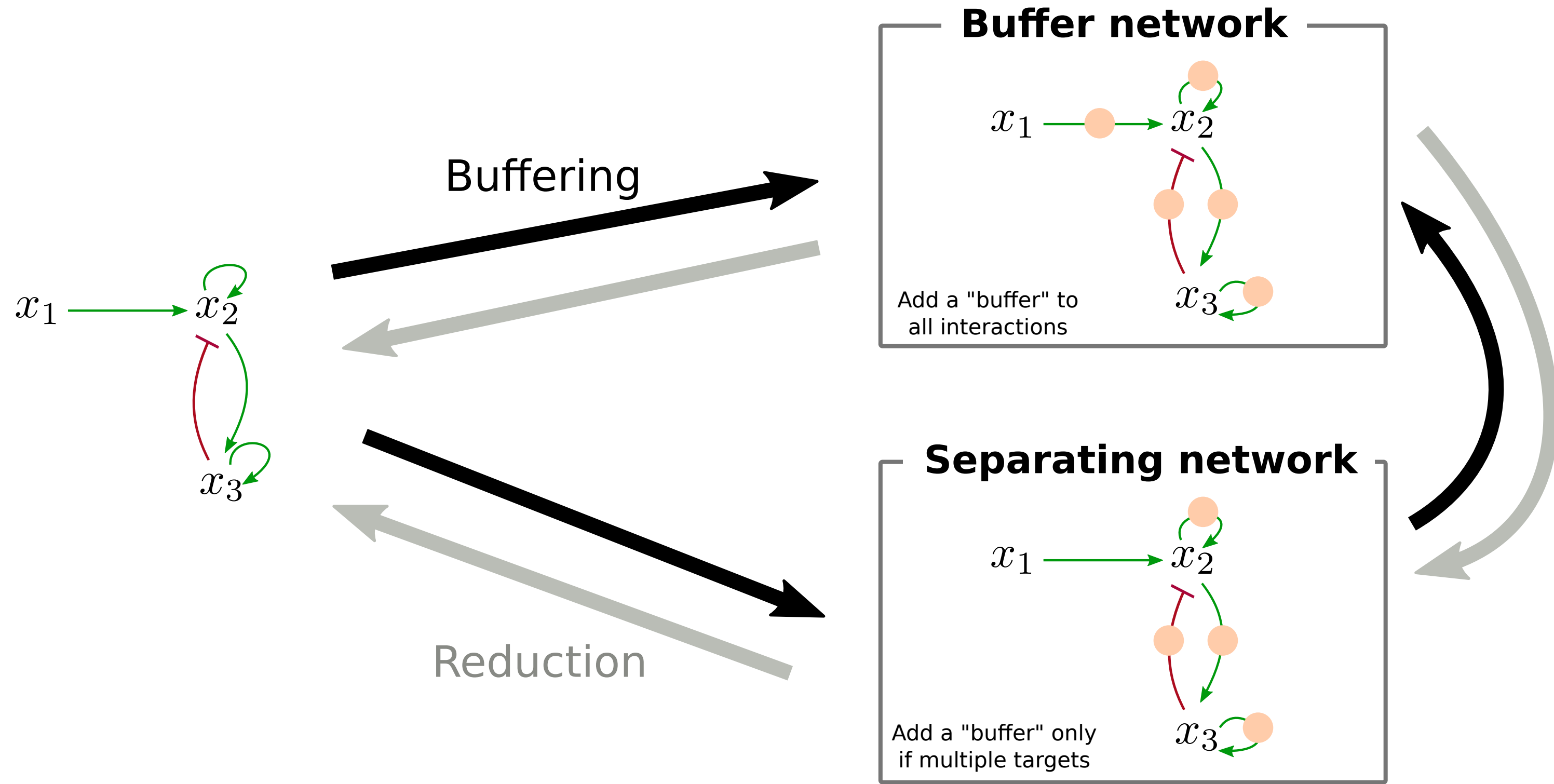
Asynchronous updating of the extended model

Buffer networks



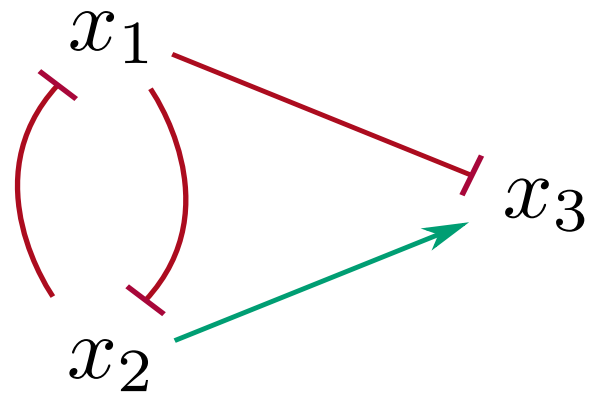
Asynchronous updating of the extended model

Buffer networks



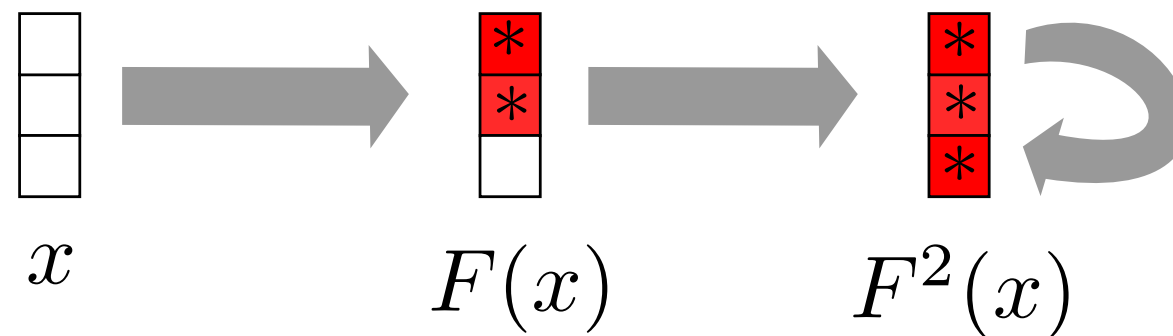
Asynchronous updating of the extended model

Enclosing stable pattern



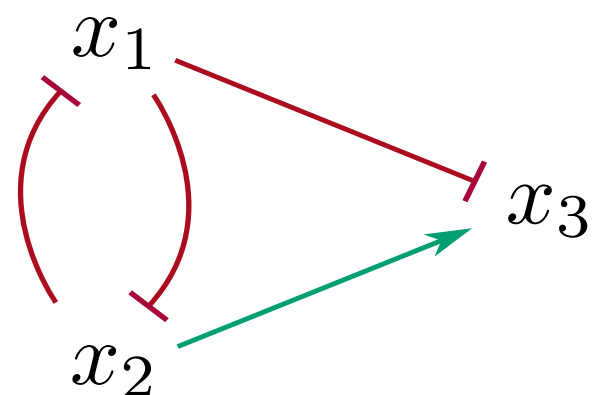
Extend f to patterns

$$F_i(p) = \begin{cases} * & \text{if } p_i = *, \\ * & \text{if } \exists x \in p : f_i(x) \neq p_i, \\ p_i & \text{otherwise.} \end{cases}$$



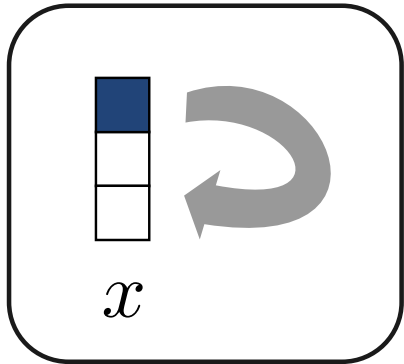
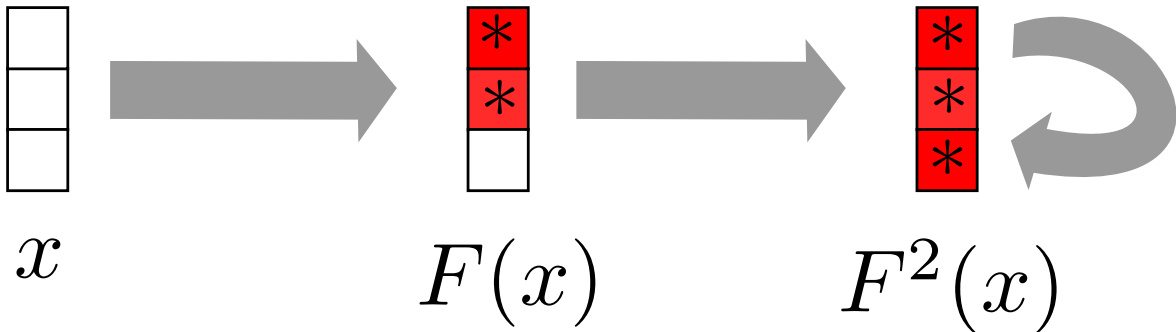
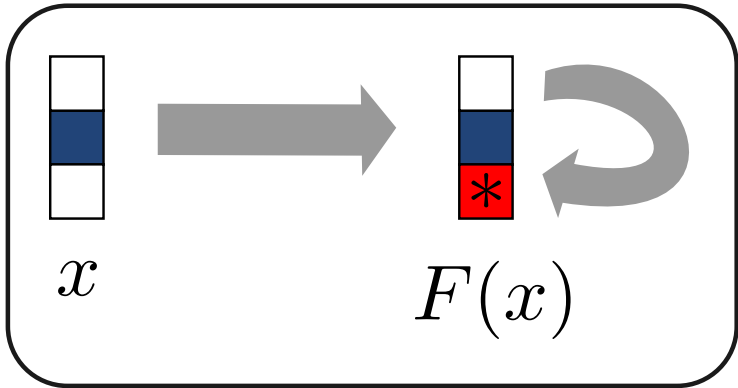
$F^n(x)$ is the smallest stable pattern
containing the state x

Enclosing stable pattern



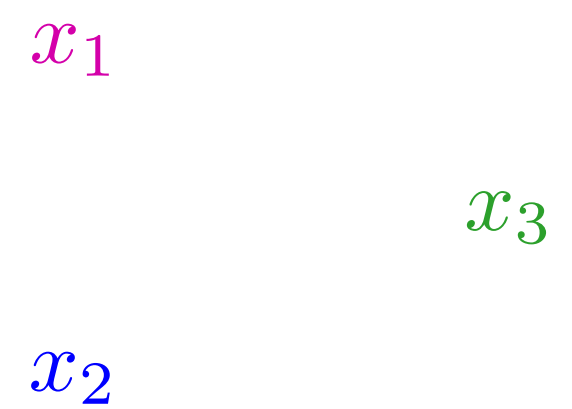
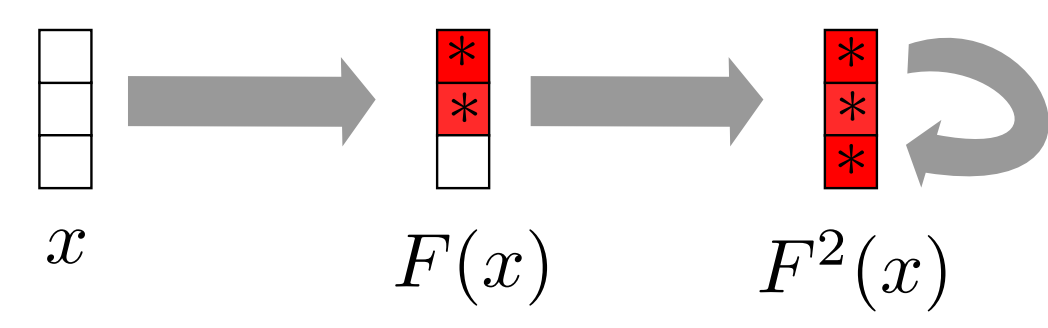
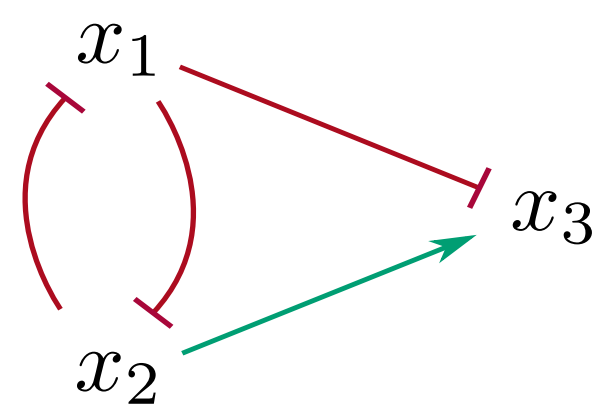
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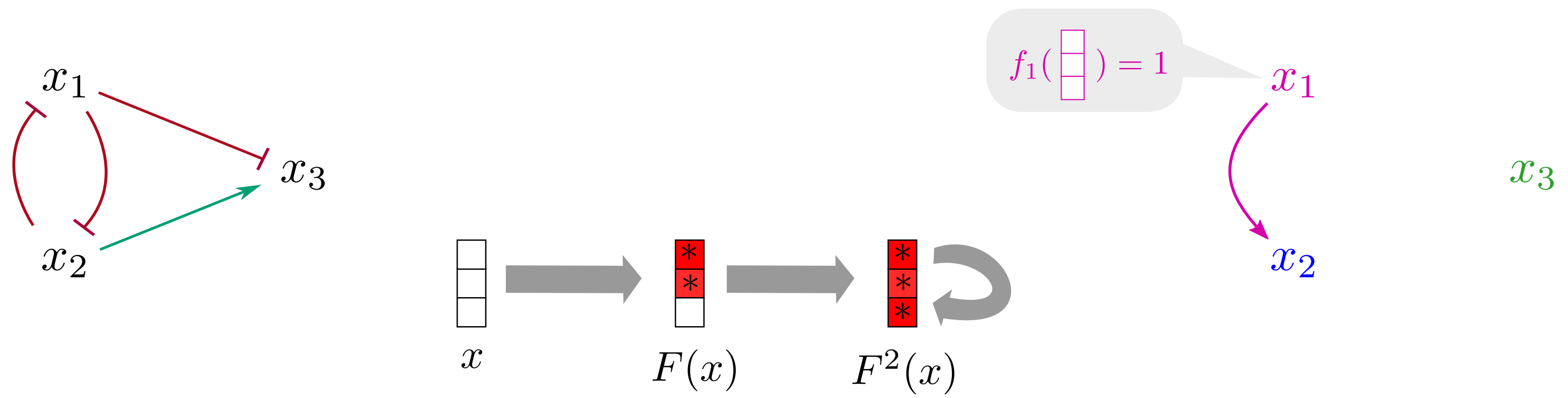


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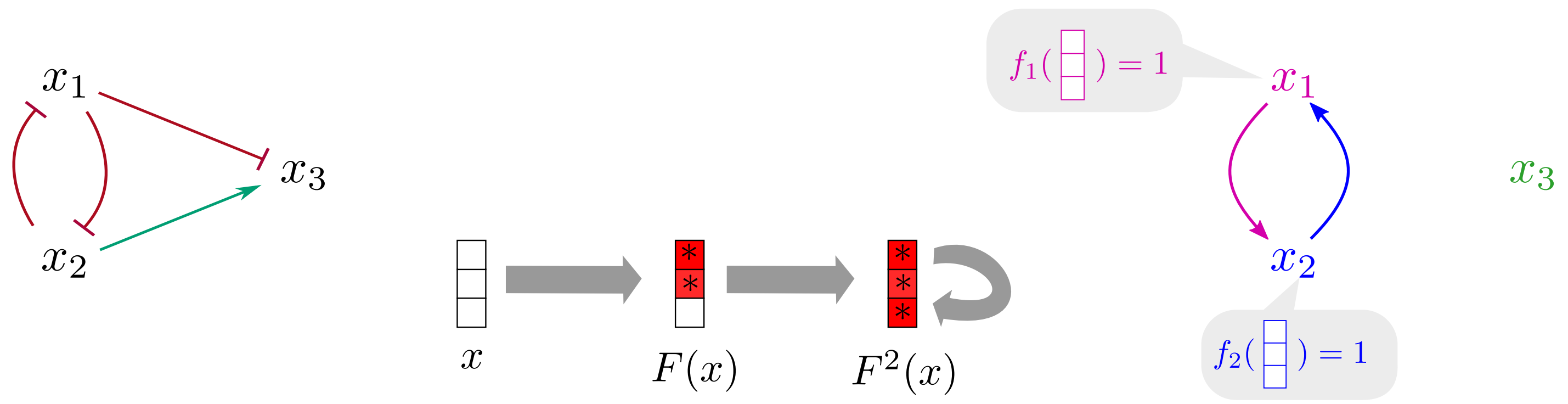
Partial order on transitions



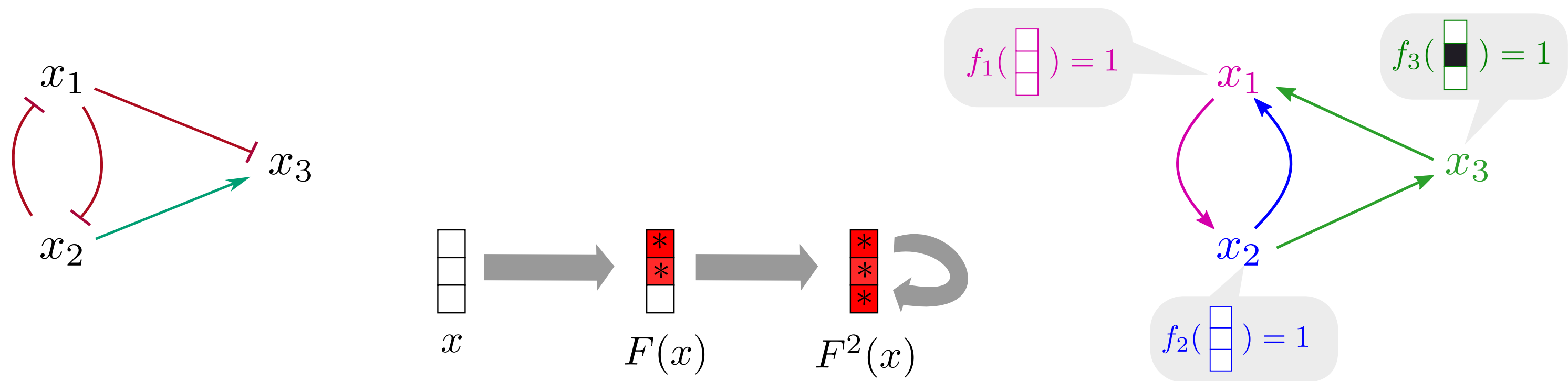
Partial order on transitions



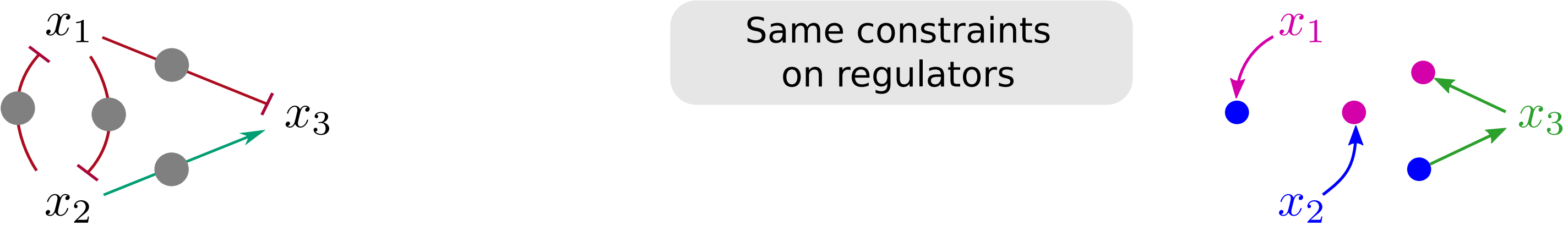
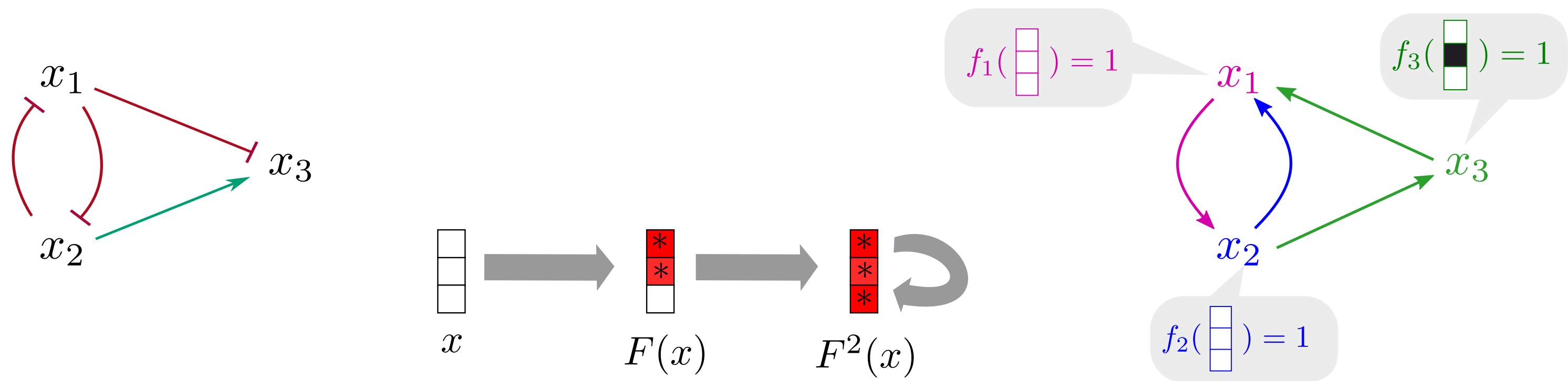
Partial order on transitions



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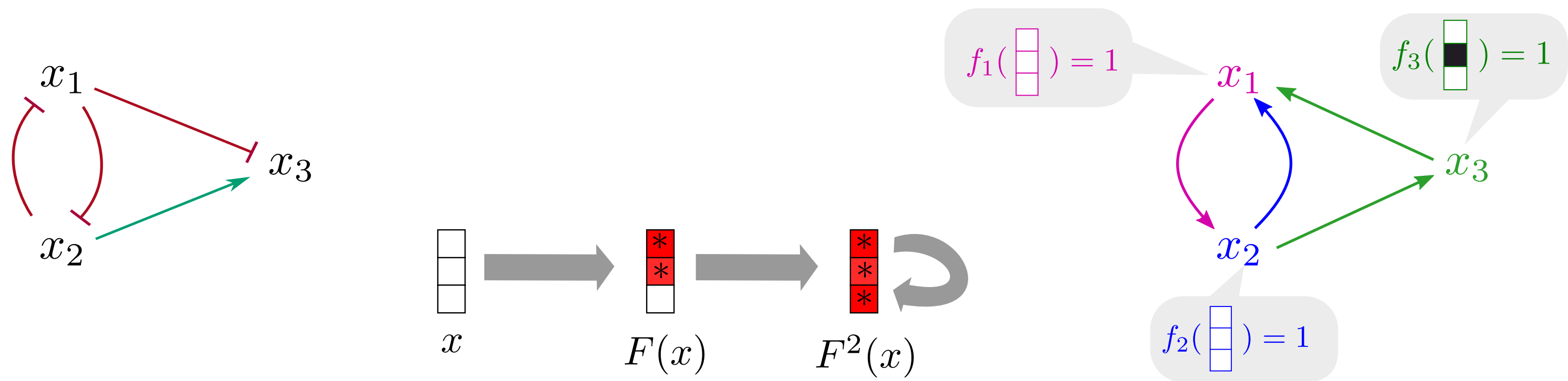


Partial order on transitions

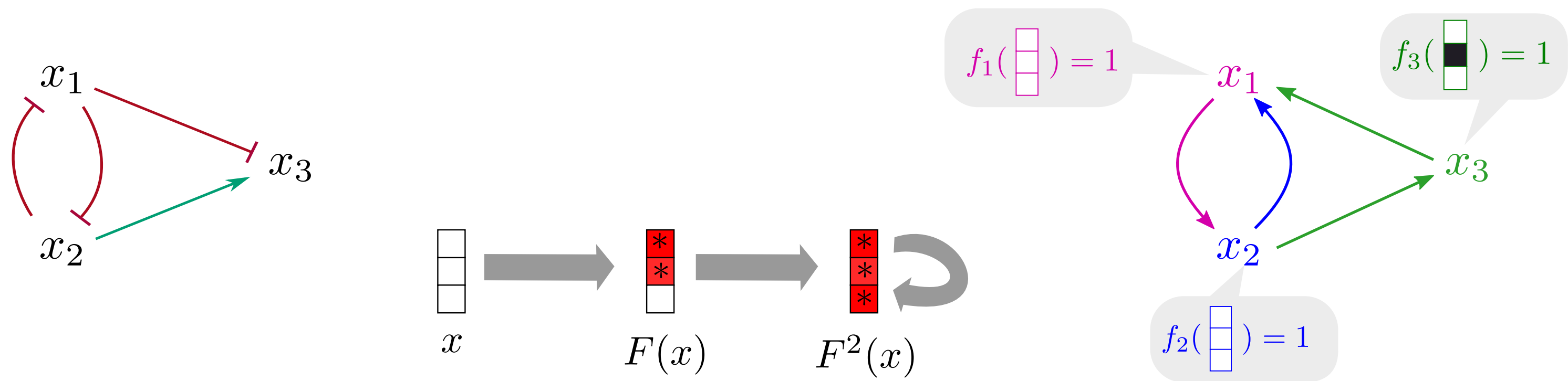


Same constraints on regulators

Partial order on transitions



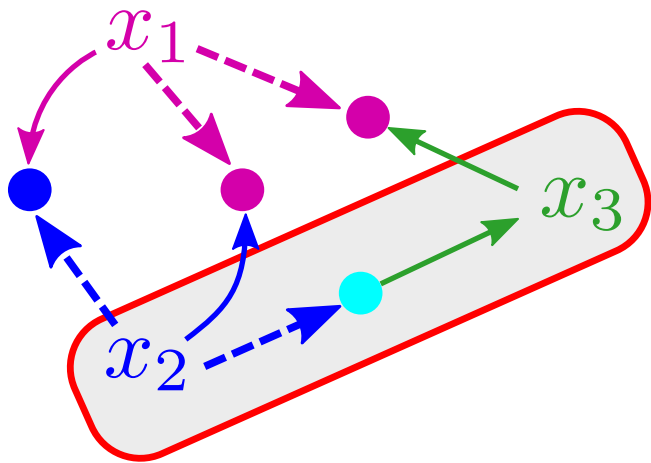
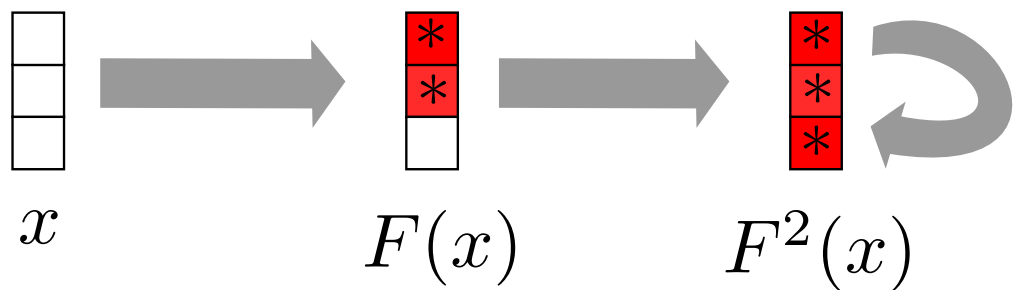
Partial order on transitions



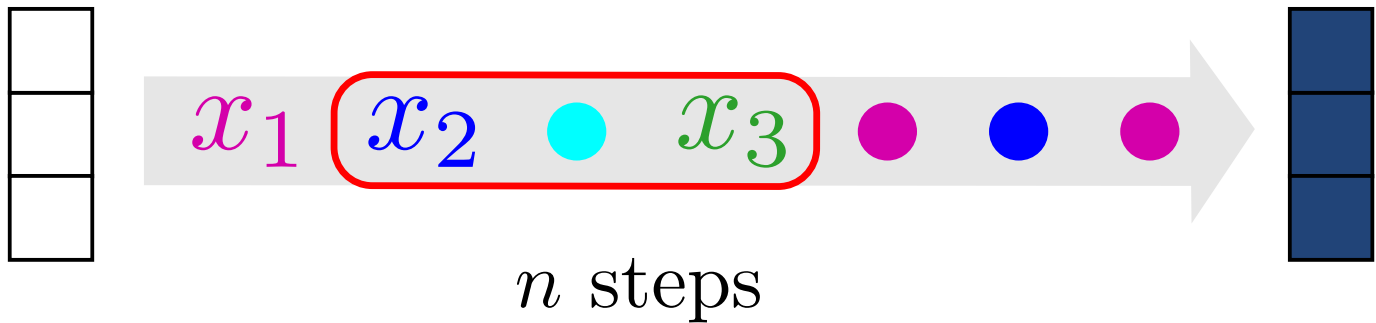
Buffers eliminate cycles in the partial order graph

Reaching the opposing corner

Acyclic partial order

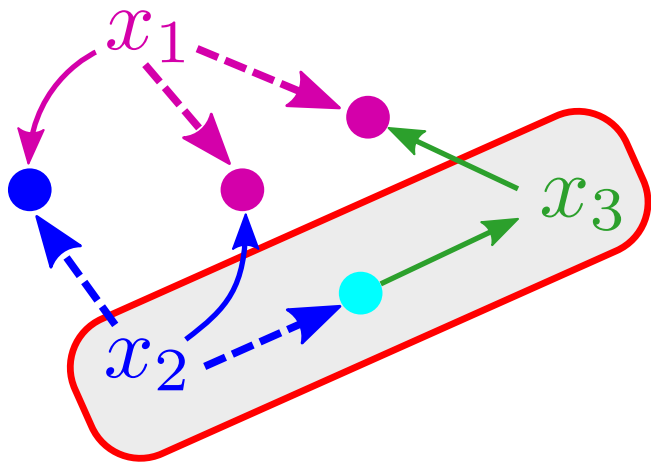
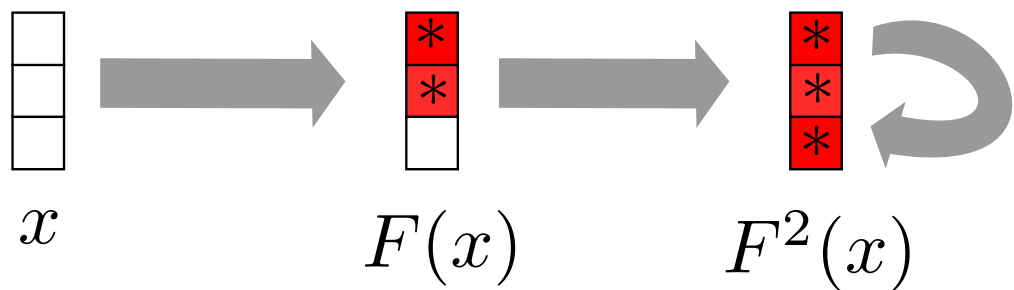


Trajectory to the opposing corner

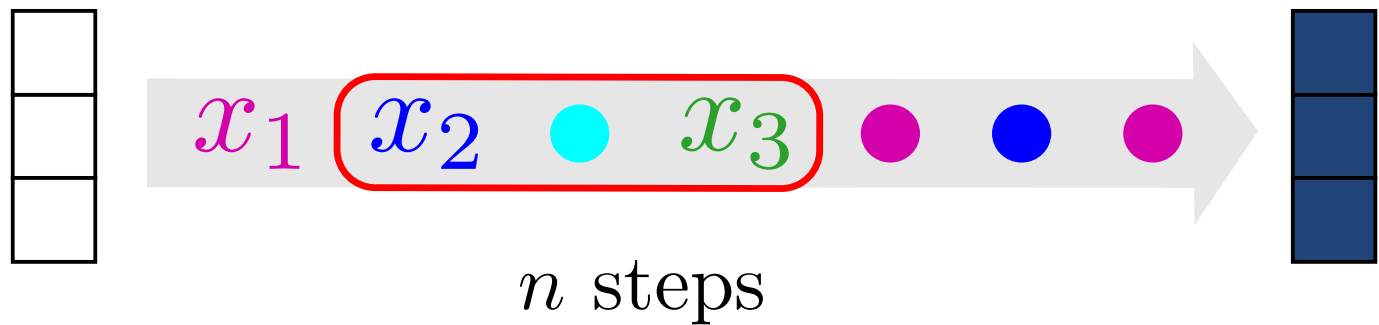


Reaching the opposing corner

Acyclic partial order

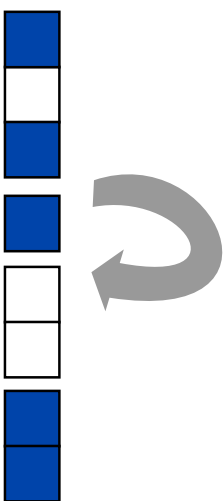
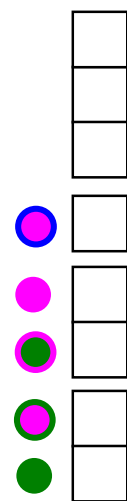
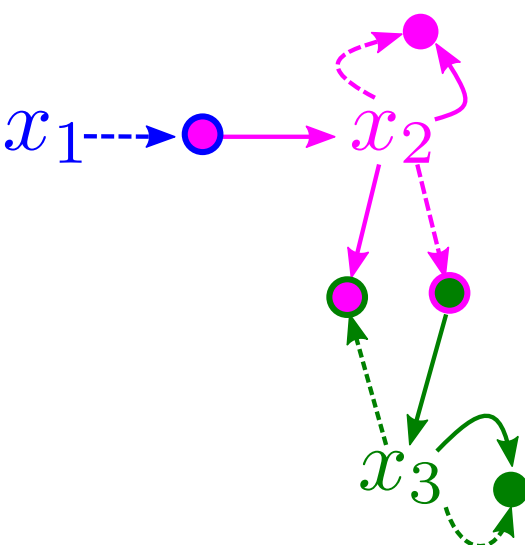
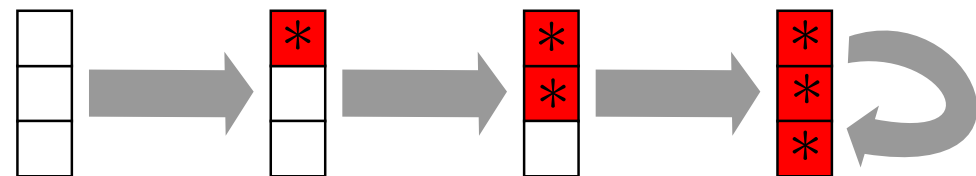
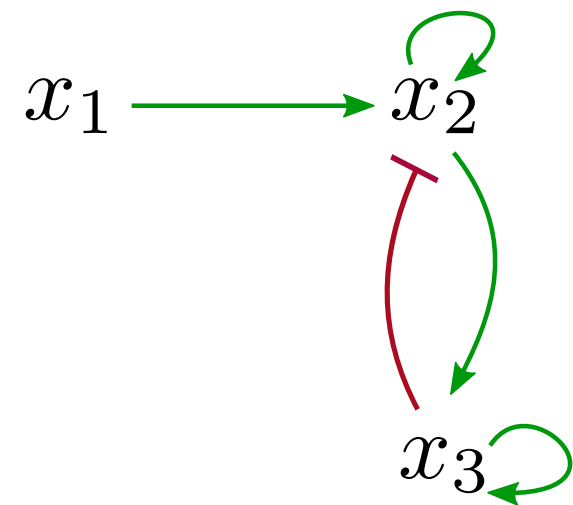


Trajectory to the opposing corner



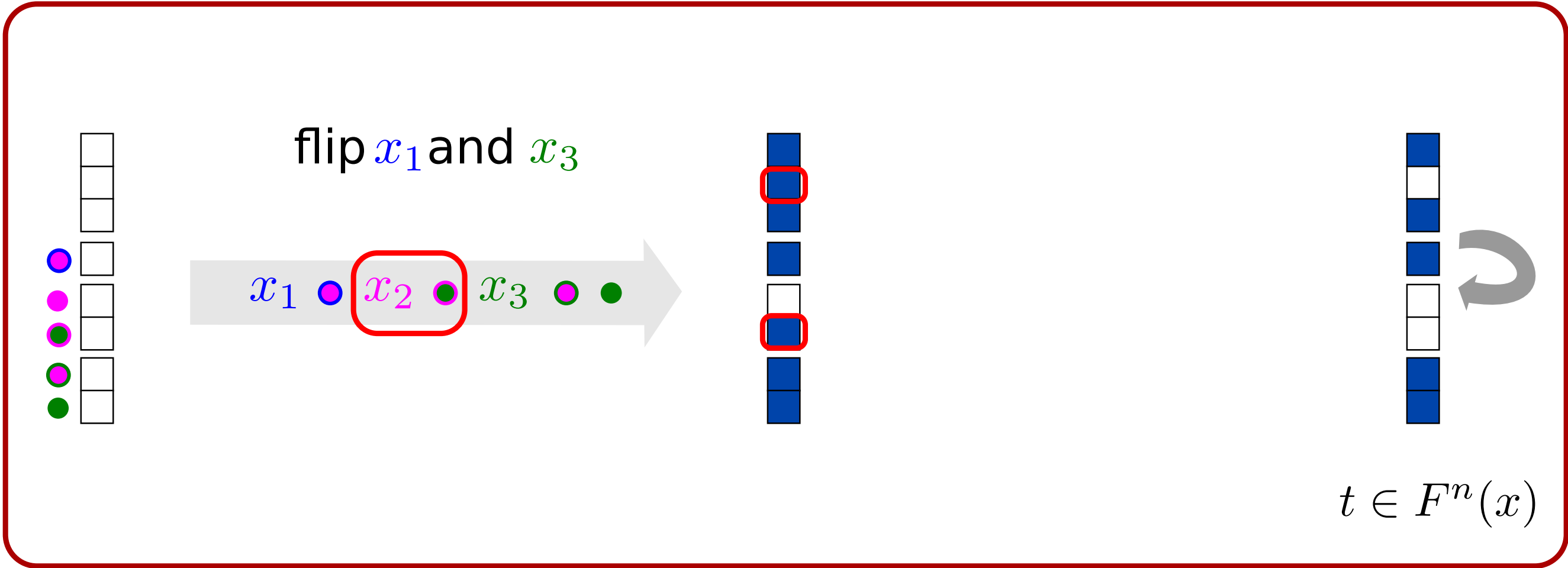
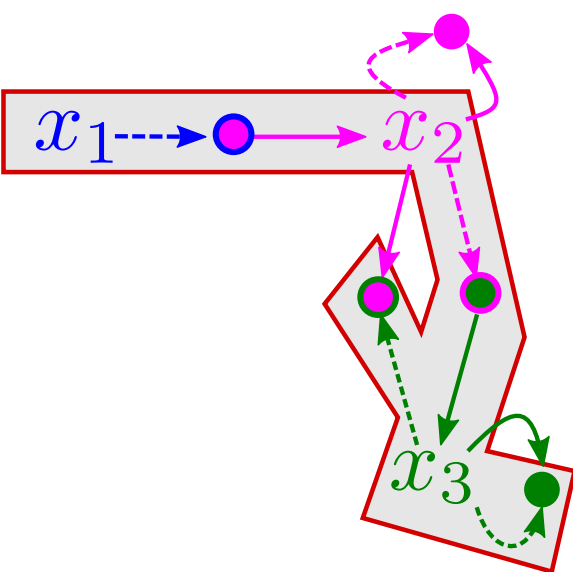
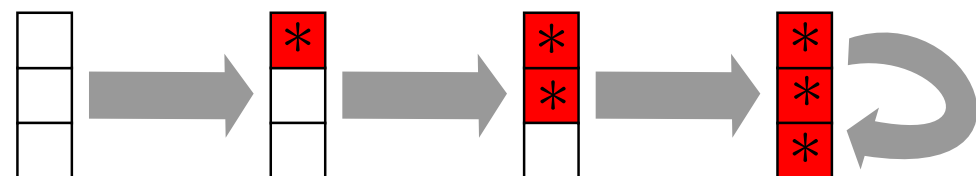
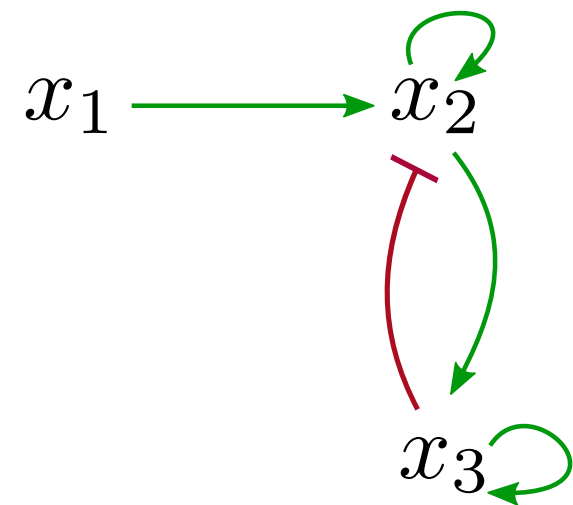
The minimal hypercube enclosing an attractor is a stable pattern

Reaching stable patterns

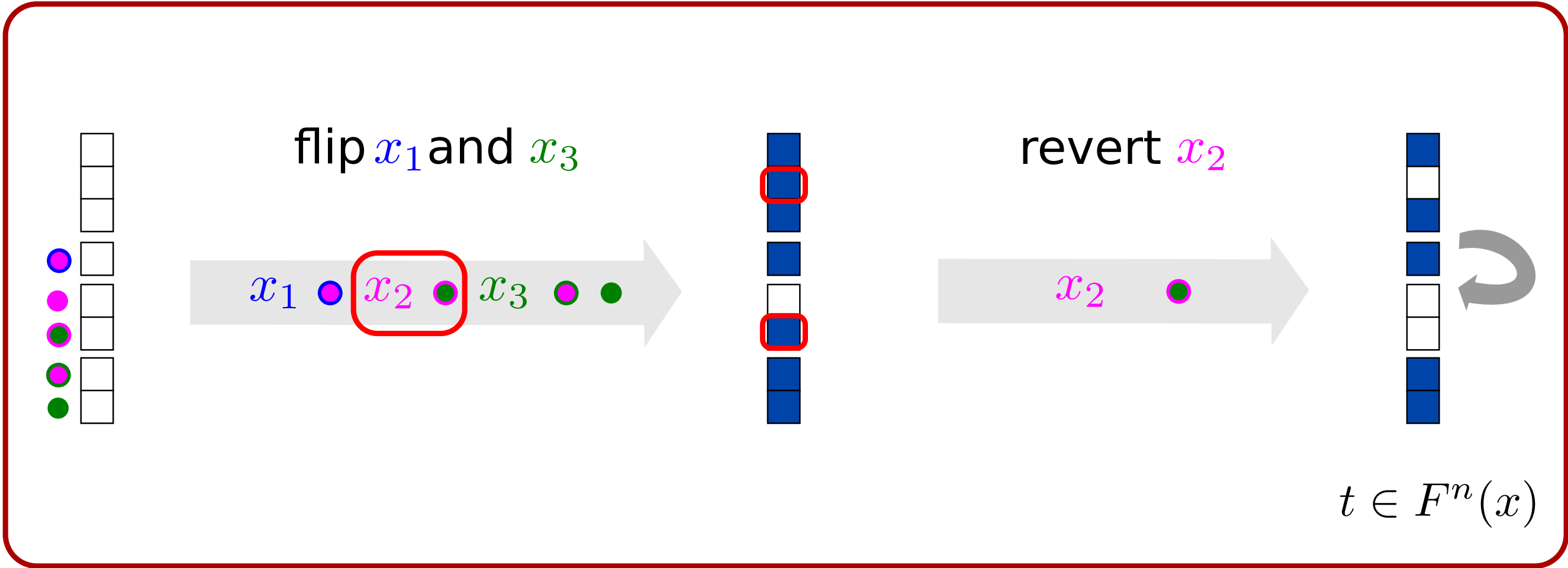
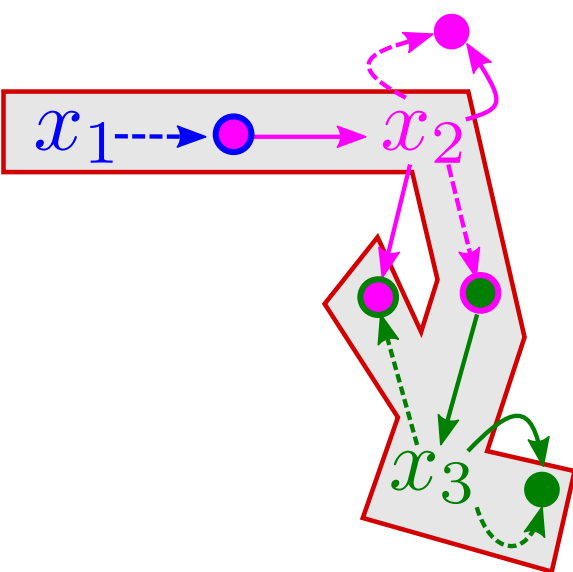
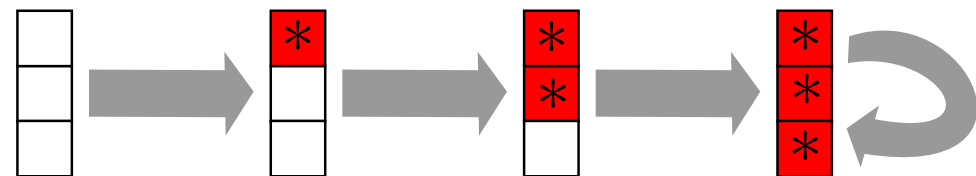
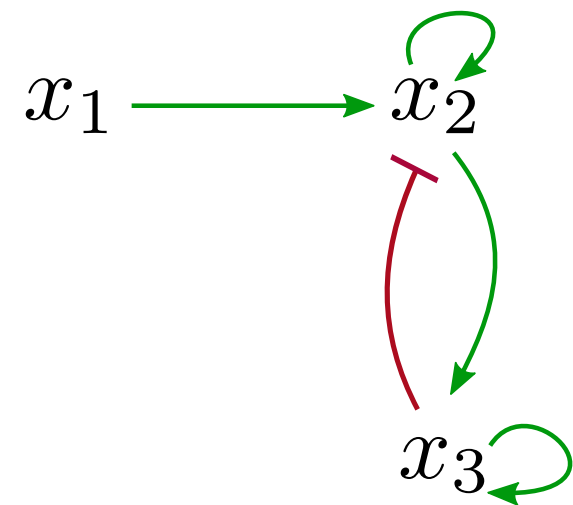


$t \in F^n(x)$

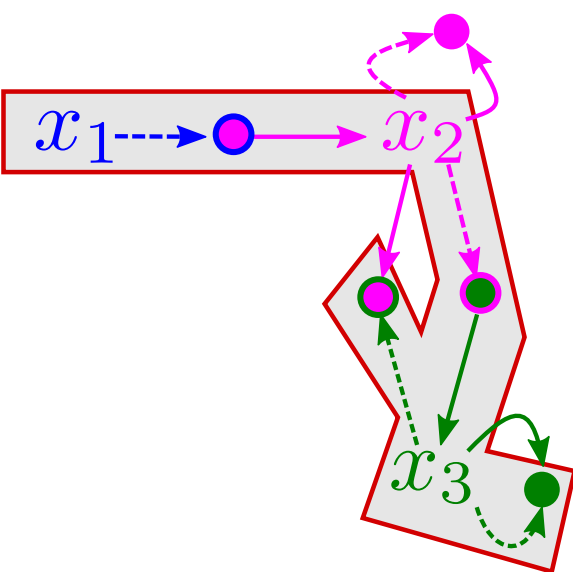
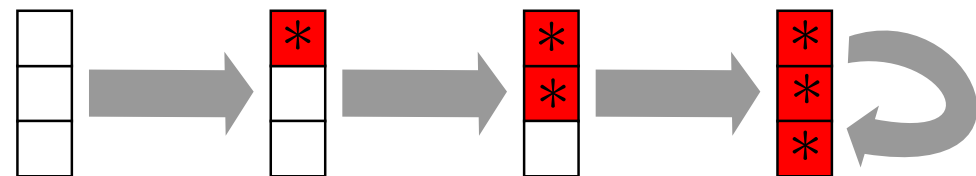
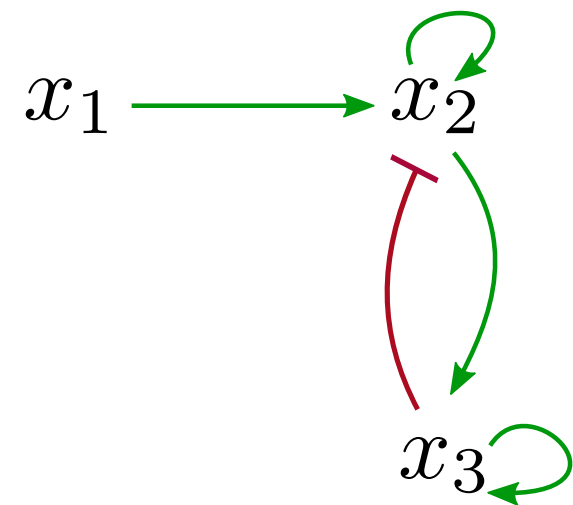
Reaching stable patterns



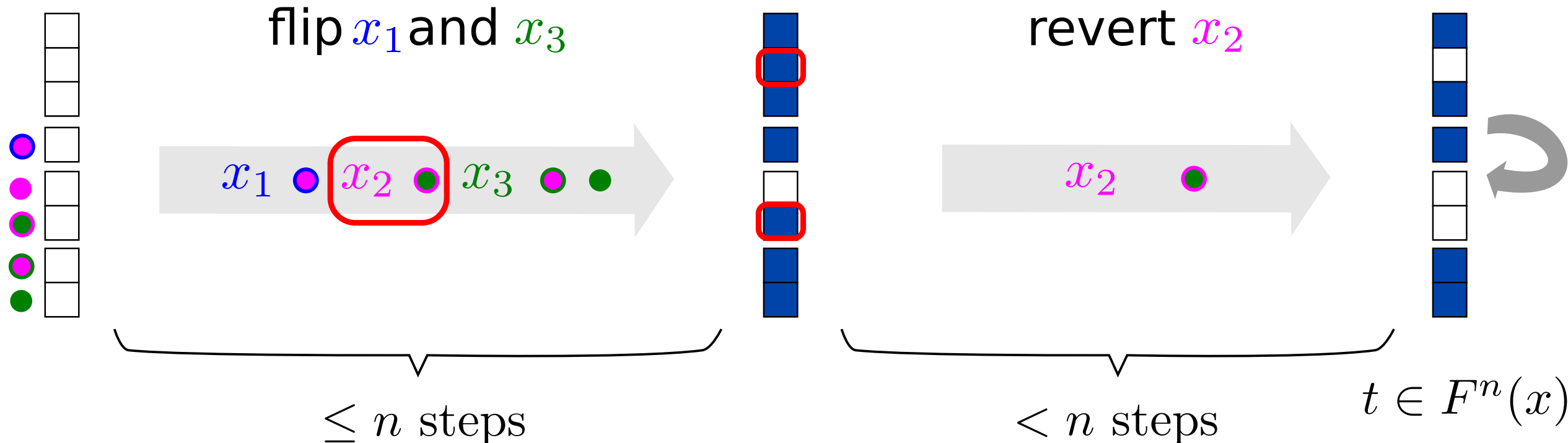
Reaching stable patterns



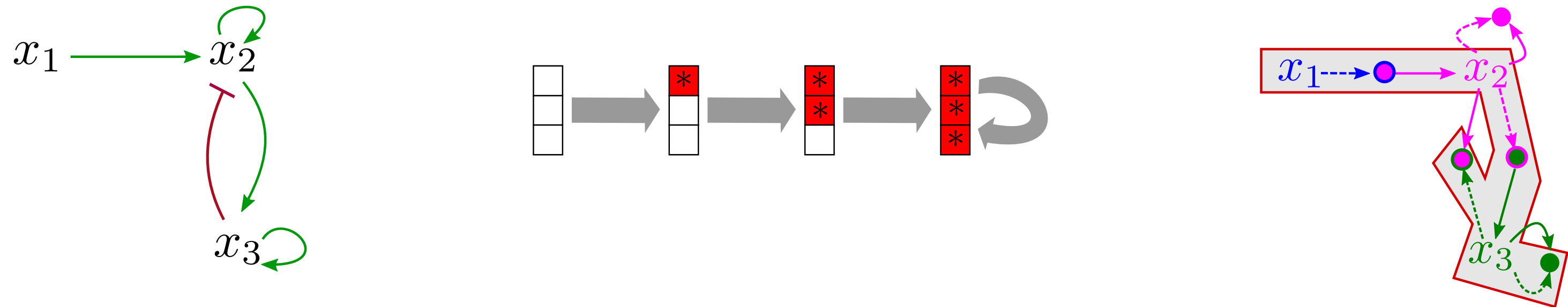
Reaching stable patterns



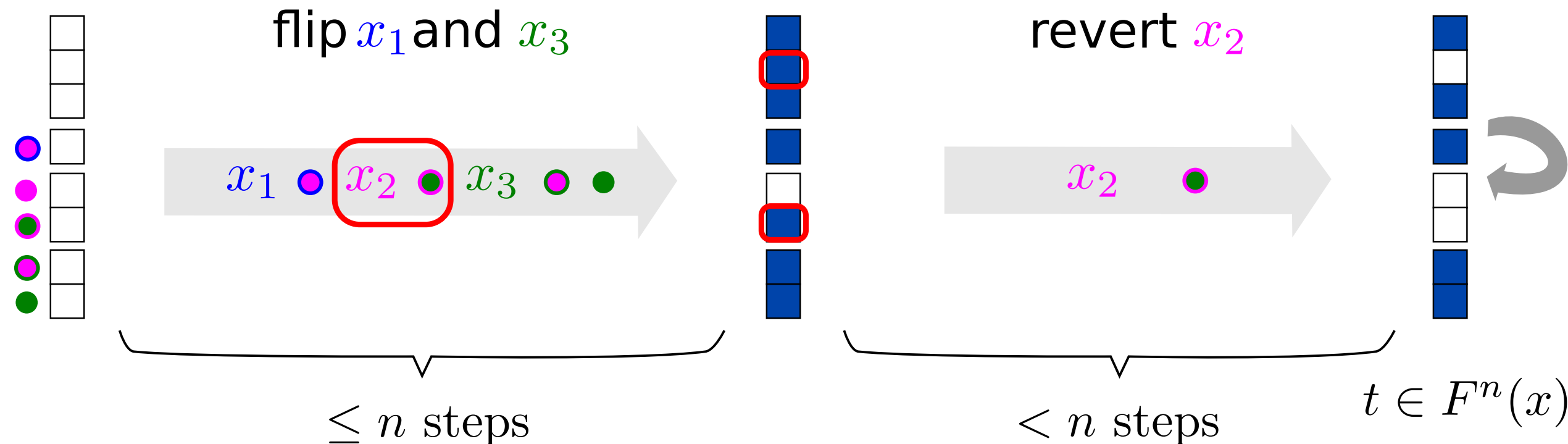
All enclosed stable patterns can be reached



Reaching stable patterns



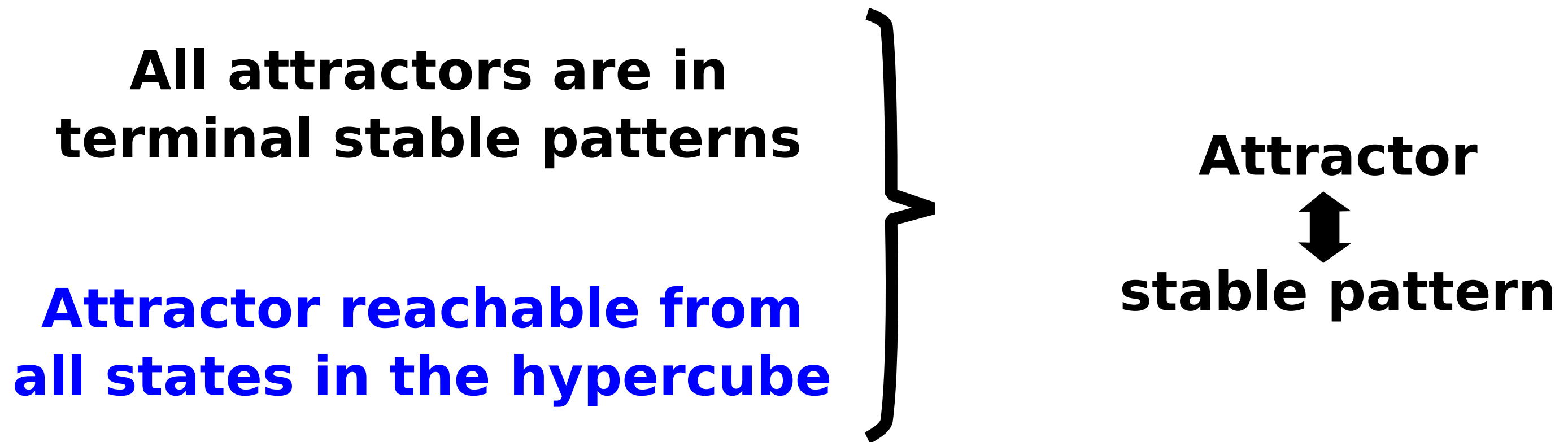
All enclosed stable patterns can be reached



All attractors are in terminal stable patterns

1 attractor \Leftrightarrow 1 trap space

The minimal enclosing hypercube is a stable pattern



Stable patterns are good estimators for attractors

Summary

	Async	Buffer	MP
TrapSPACE vs attractor	lower bound	1:1	Exact
Reach. states	2^n	???	$3n$
Reach. trapSPACE	2^n	$2(n + m)$	$3n$

**Stable patterns as robust abstraction
for attractors and reachability**

Analysis directly on un-buffered model

Discussion

Limits

Assumption: variable threshold order

↗ kinetic information → ↘ alternative trajectories

Perspectives

From partial order to constraints on kinetic parameters

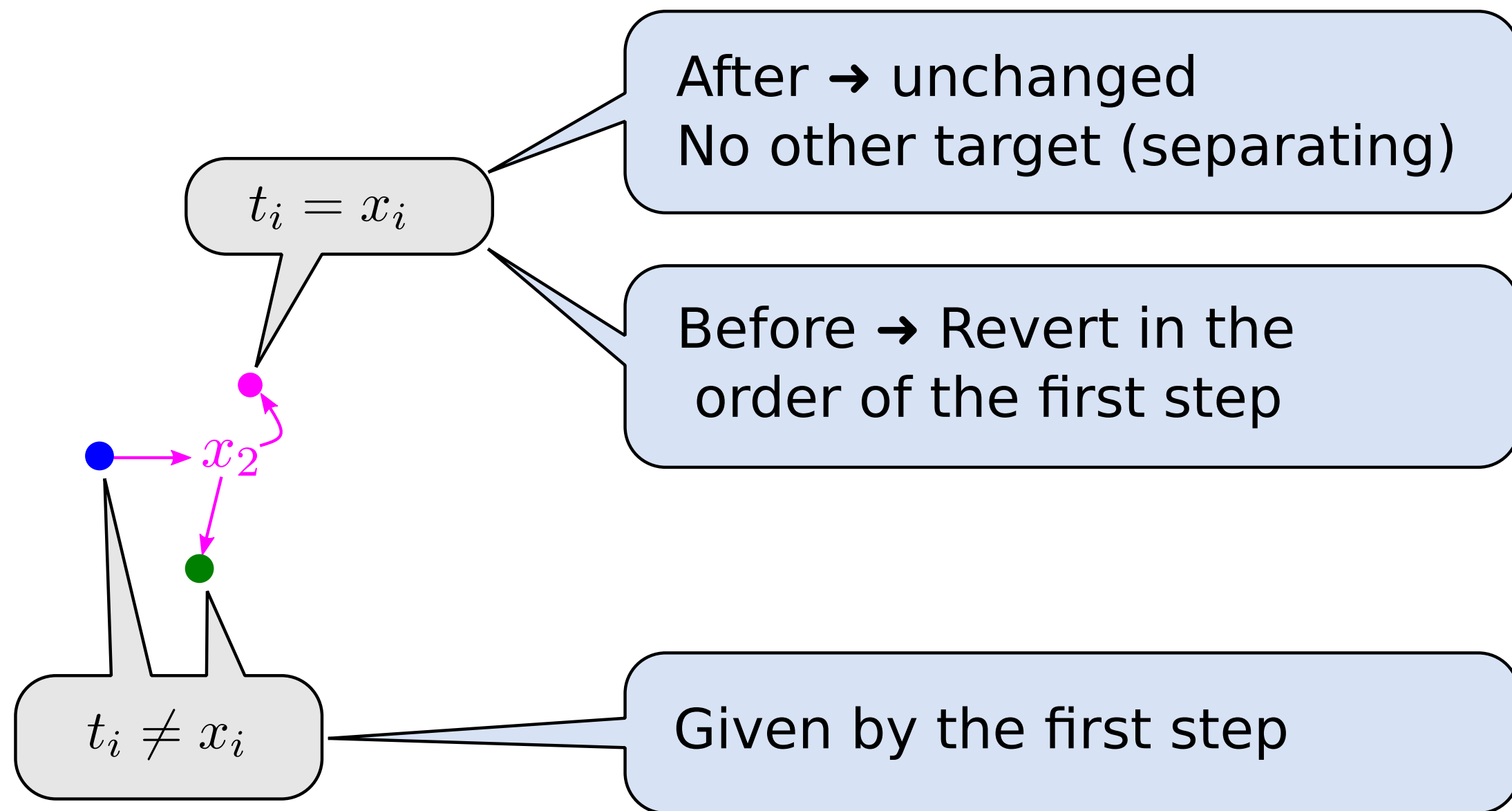
Prime implicant graph to improve partial order

Can use existing stochastic tools (MaBoSS)

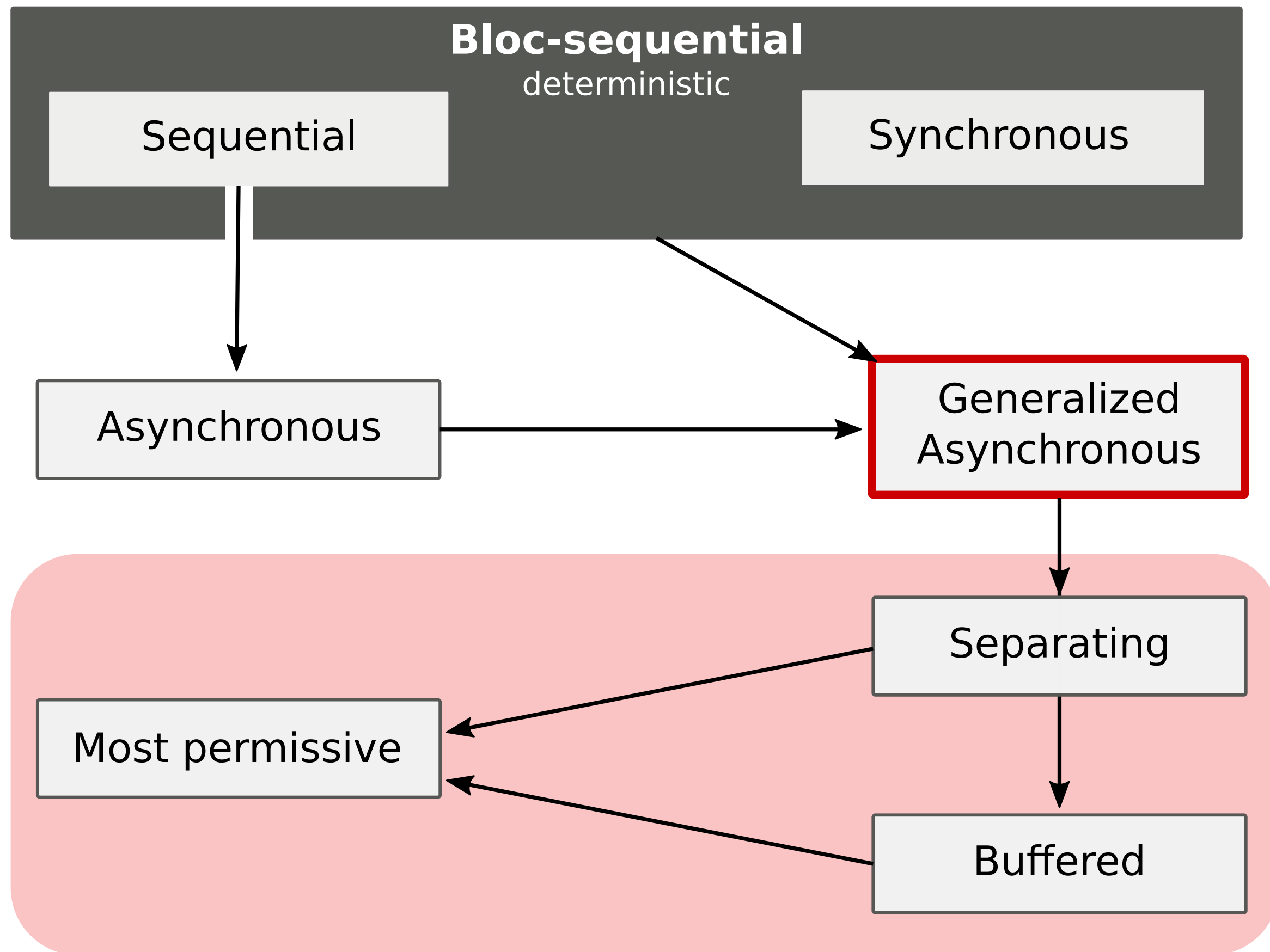
Derive analytic estimation of probabilities

Reverting to reach a trapspace

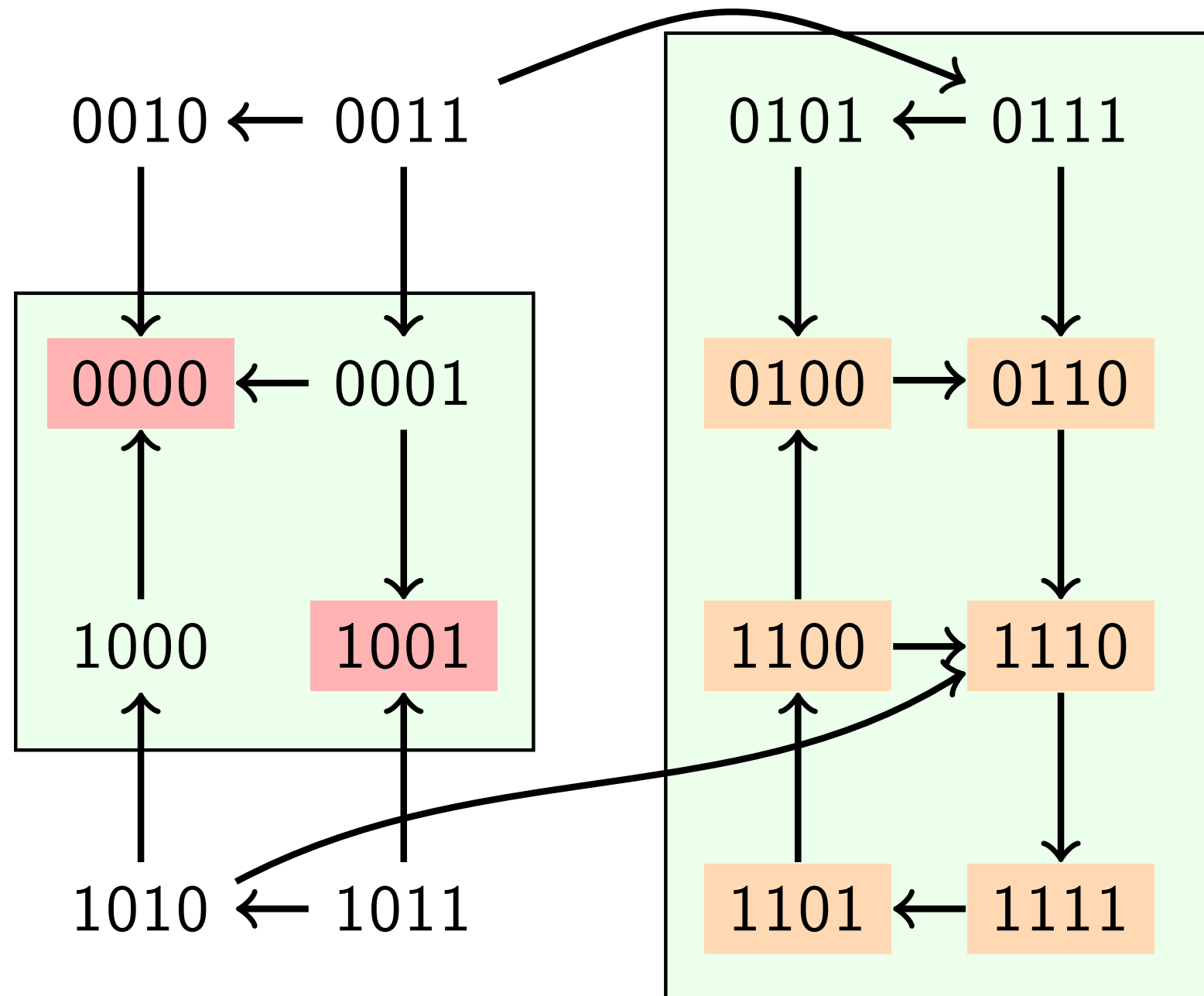
For each core: fix regulators to the trapspace value



Reachability across updating semantics



Attractors = stable phenotypes



Trap spaces

(stable patterns/hypercubes)

hypercube h :

$\forall x \in h : f(x) \in h$

Constraint solving

Zanudo et al., 2013
Klarner et al., 2014

Stable states (fixed points)

state x : $f(x) = x$

Constraint solving

Naldi et al., 2007

Complex attractors

states C : $\{x : F(x) = C\}$

Symbolic exploration

Garg et al., 2008