

Understanding Evolution:

A Methodology for Evaluating the Extensibility of Boolean Networks' Structure and Function

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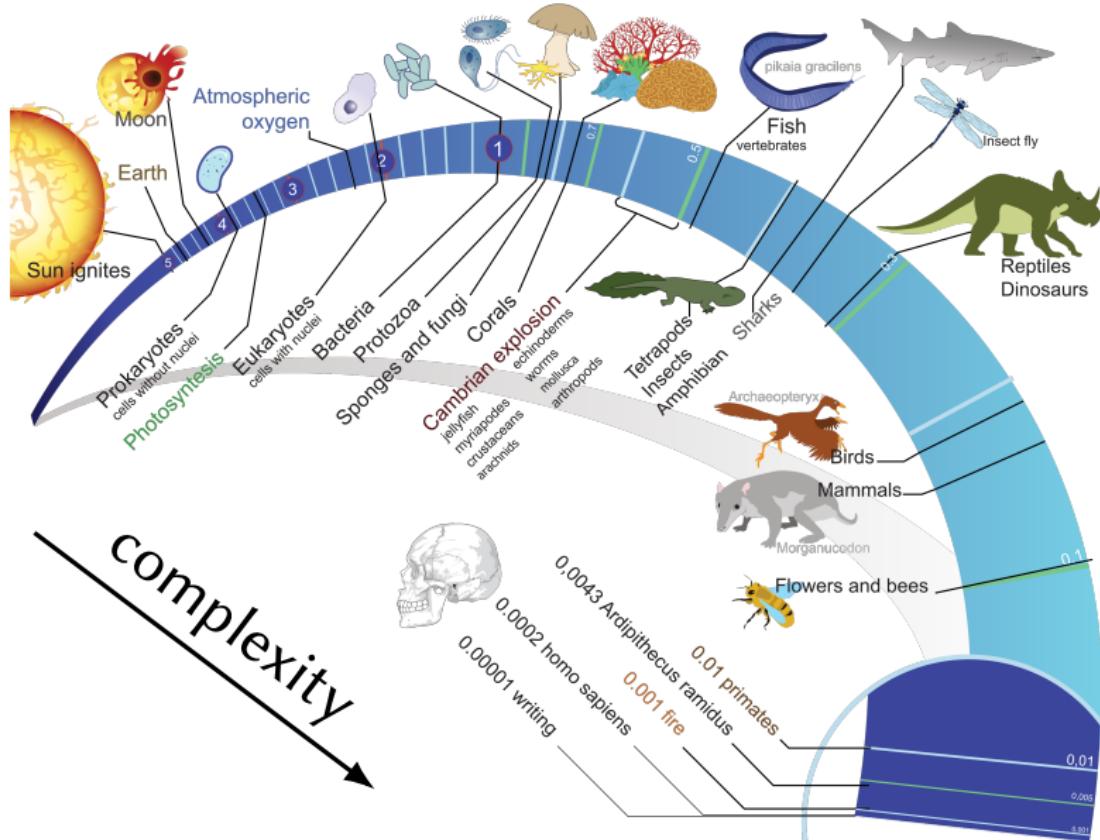
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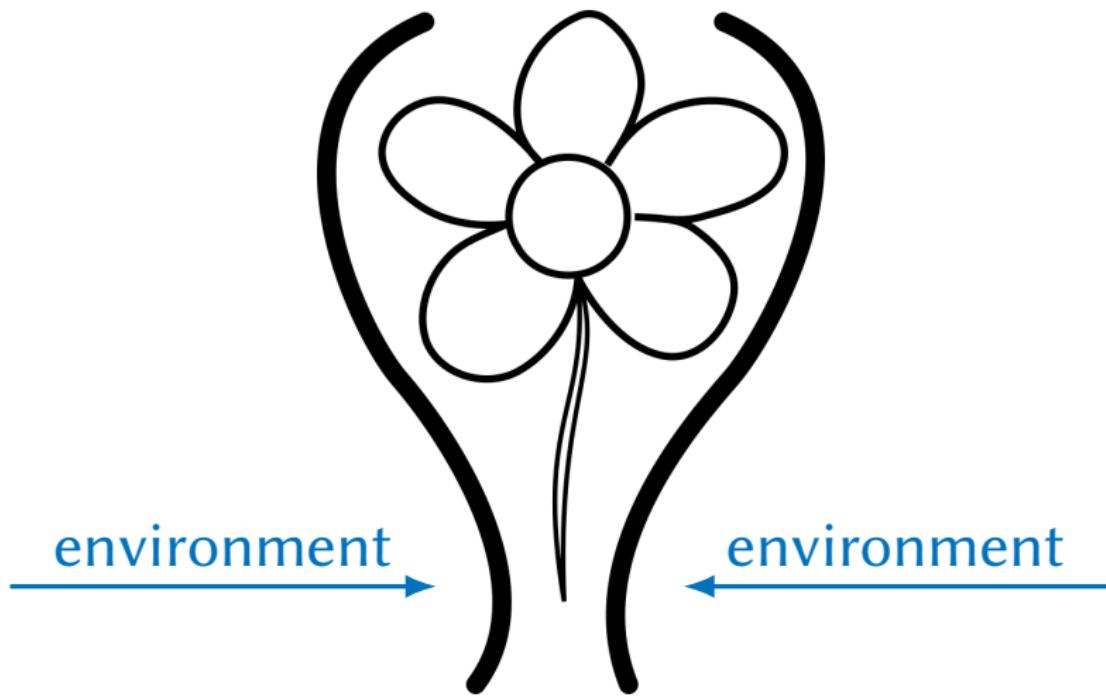
Evolution



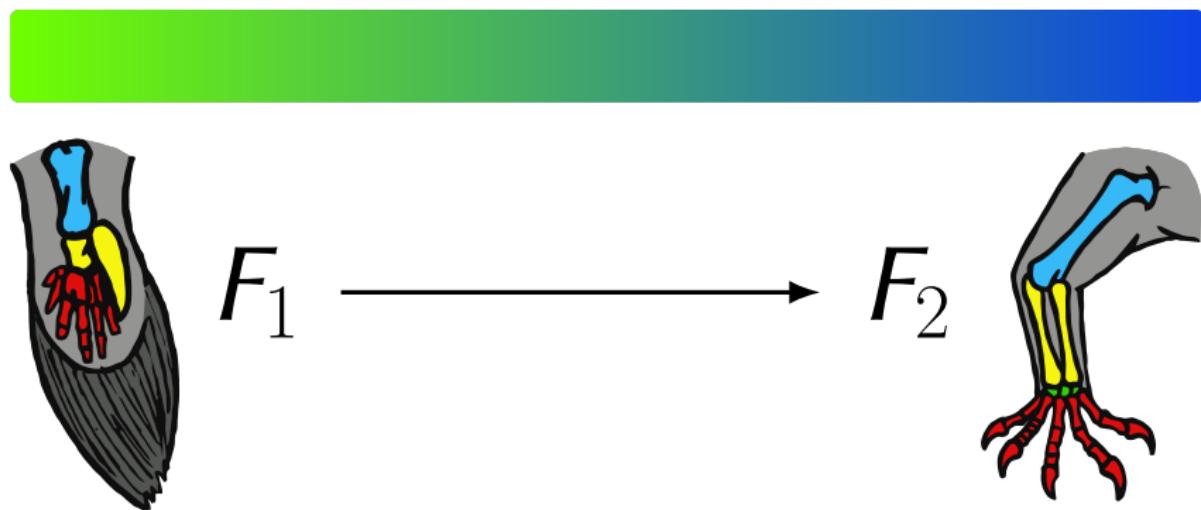
<https://openclipart.org/download/320330/timelineevolutionoflife-pd-ladyofhats-wikimedia.svg>

Natural selection

reinforces the features best suited to the environment.



Evolution is gradual

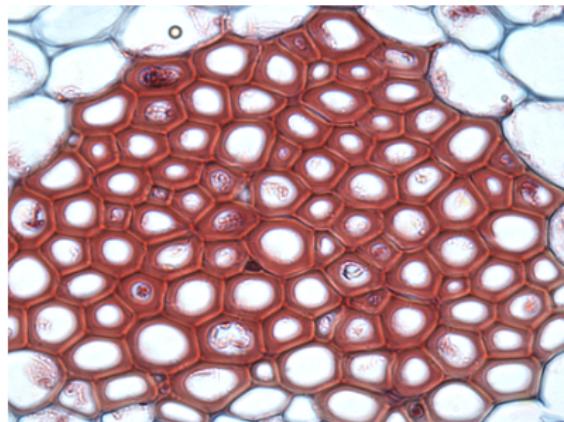


The cards are rarely entirely reshuffled: maintaining the functioning of the whole system is essential for survival.

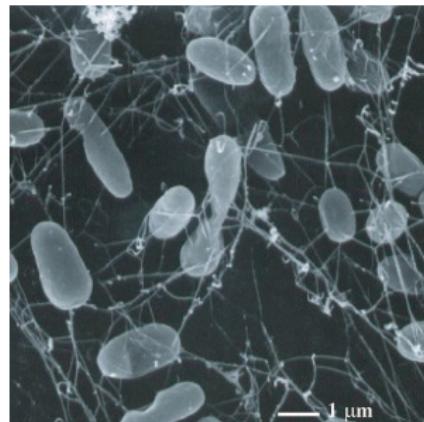
https://en.wikipedia.org/wiki/File:Crossopterygii_fins_tetrapod_legs.svg

Evolution and cancer

Atavistic theory of cancer: cancerous cells devolve to primitive unicellular mode.



1



2

¹ https://en.wikipedia.org/wiki/File:Plant_cell_type_sclerenchyma_fibers.png

² https://en.wikipedia.org/wiki/File:Thermophile_bacteria2.jpg

How does an organism grow more complex, while maintaining original structures and functions?

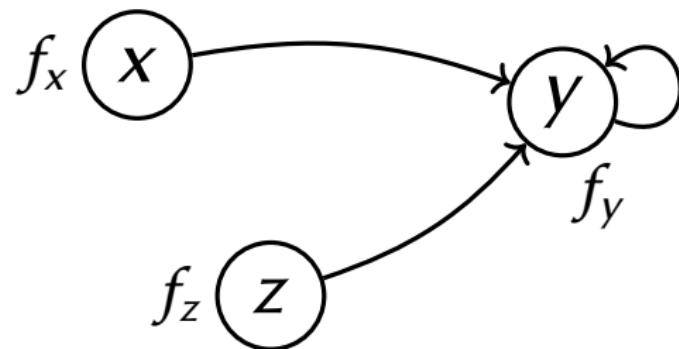
Which structures are the most extensible?

What language to answer this question?

- 1 Threshold/Sign Boolean networks (BN)
- 2 Extensibility in Sign BN
- 3 Complexity of Sign BN
- 4 Preliminary results and analysis

Boolean networks

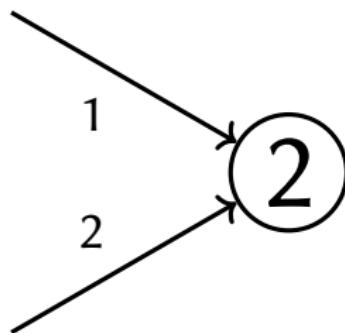
Boolean nodes with Boolean update functions.



$x, y, z \in \{0, 1\}$, a state: $\{x : 0, y : 1, z : 0\} \equiv 010$

$$f_x, f_y, f_z : \{0, 1\}^3 \rightarrow \{0, 1\}$$

TBF: Threshold Boolean Functions



$$f(x_1, x_2) \equiv w_1 x_1 + w_2 x_2 > \theta$$

$$f = (w_1, w_2, \theta)$$

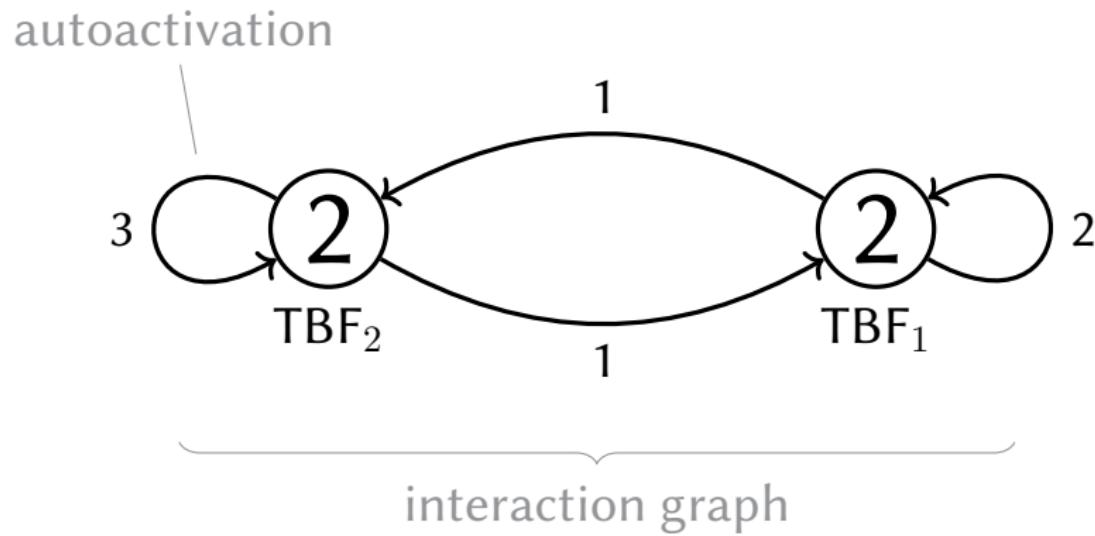
$$w_1 = 1$$

$$w_2 = 2$$

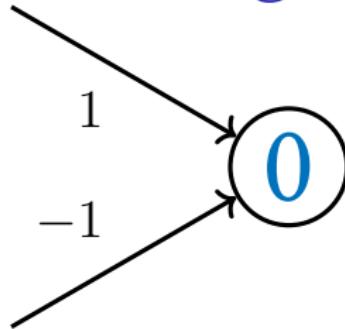
$$\theta = 2$$

w_1	w_2	f
x_1	x_2	
0	0	0
0	1	0
1	0	0
1	1	1

TBN: Threshold Boolean networks



SBF/N: Sign Boolean functions/networks



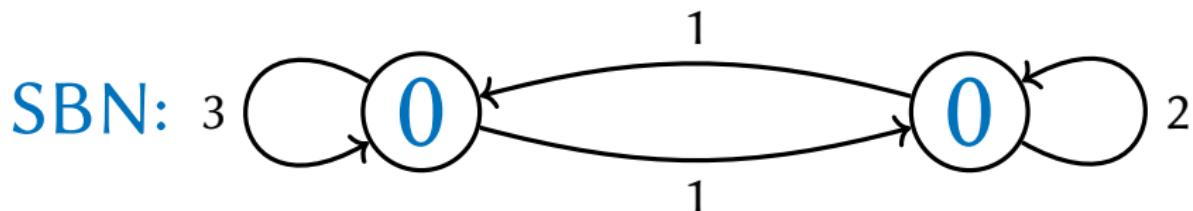
$$f(x_1, x_2) \equiv w_1 x_1 + w_2 x_2 > 0$$

$$f = (w_1, w_2)$$

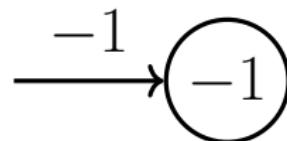
$$w_1 = 1$$

$$w_2 = -1$$

x_1	x_2	f
0	0	0
0	1	0
1	0	1
1	1	0



SBF < TBF

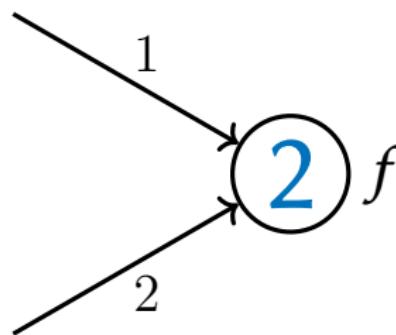


$$\begin{aligned} f &= (w_1 = -1, \theta = -1) \\ &\equiv \neg x_1 \end{aligned}$$

x_1	f
0	1
1	0

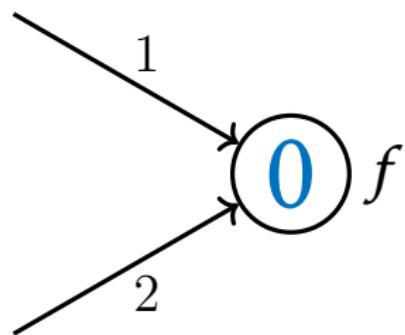
$$\forall f \in \text{SBF} : f(\mathbf{0}) = 0$$

SBN vs. TBF



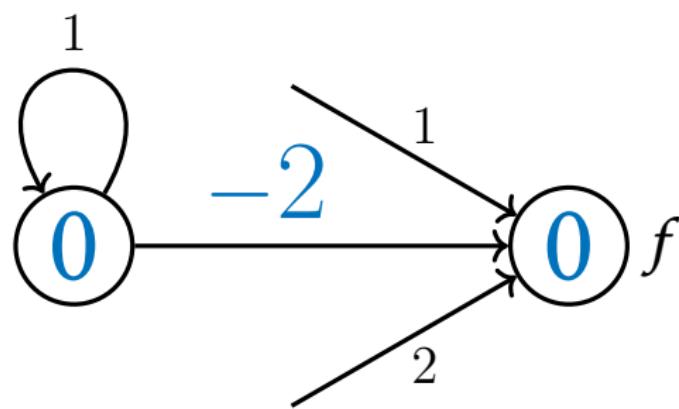
x_1	x_2	f
0	0	0
0	1	0
1	0	0
1	1	1

SBN vs. TBF



x_1	x_2	f
0	0	0
0	1	1
1	0	1
1	1	1

SBN vs. TBF



x_1	x_2	f
0	0	0
0	1	0
1	0	0
1	1	1

TBF = SBF + a hidden SBF

TBF vs. propositional BF

		1	1	0		1	1	1
		x_1	x_2	$x_1 \vee x_2$		x_1	x_2	$x_1 \wedge x_2$
w_1	θ							
-1	-1	0	0	0		0	0	0
x_1	$\neg x_1$	0	1	1		0	1	0
0	1	1	0	1		1	0	0
1	0	1	1	1		1	1	1

$\{\neg, \vee, \wedge\}$ is a basis for BF

\implies any BF = a combination of TBF (a TBN)

SBN vs. TBN vs. propositional BN

Propositional BN: much power hidden in the formula

TBN: the power is explicit in the connections

→ the interaction graph completely specifies the SBN

SBN: a more symmetric and explicit model

→ all numerical parameters have the same role

syntactic simplicity

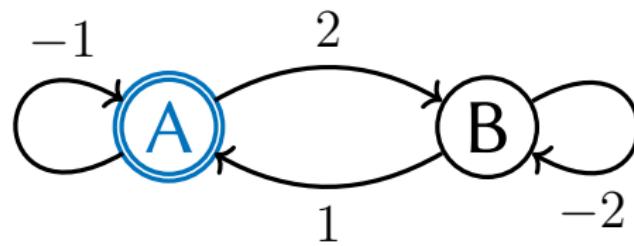
$$\text{power}(\text{SBN}) = \text{power}(\text{TBN}) = \text{power}(\text{propositional BN})$$

Update modes

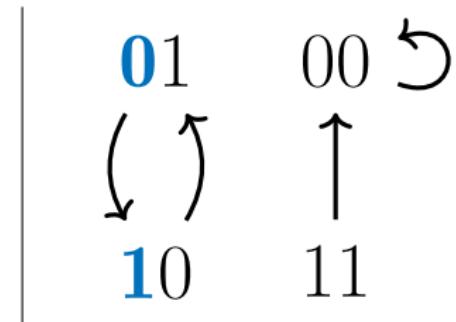
We typically consider the **synchronous** mode
→ all variables are updated at every step.

Other modes can be considered, too.

SBN: Transition graphs and outputs



interaction graph



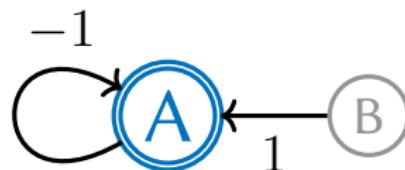
transition graph

Output = the binary sequence observed on (A)

Synchronous mode \implies outputs $\subseteq uv^*$

$u \in \{0, 1\}^k$ – the pre-period $v \in \{0, 1\}^m$ – the period

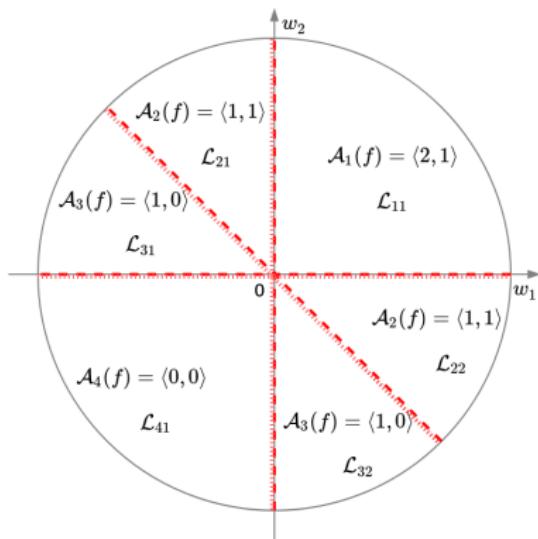
Equivalence by truth tables



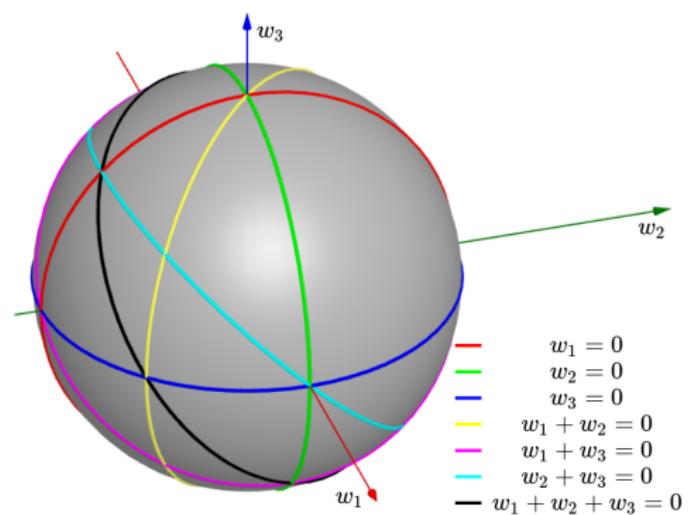
A	B	A	
0	0	0	$0 \leq 0$
0	1	1	$w_B > 0$
1	0	0	$w_A \leq 0$
1	1	0	$w_A + w_B \leq 0$

The inequalities determine subspaces of the parameter space = equivalence classes

Equivalence classes



$$d = 2$$



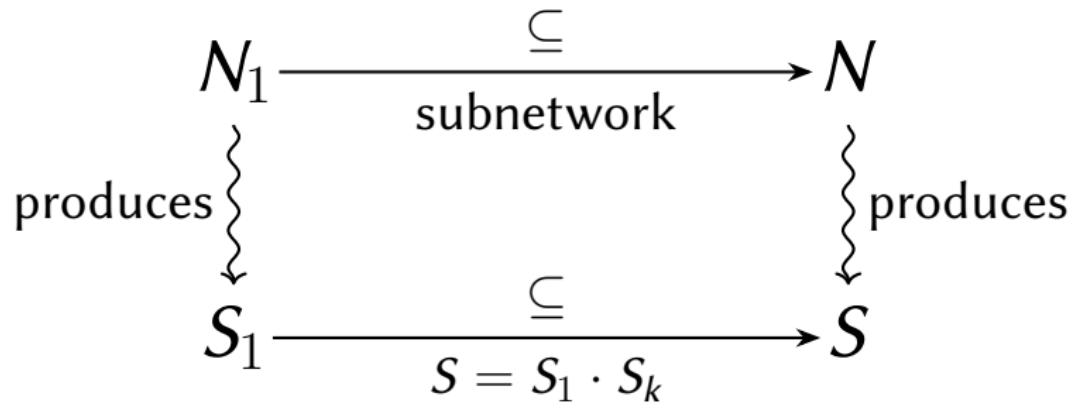
$$d = 3$$

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Definition of extensibility in SBN

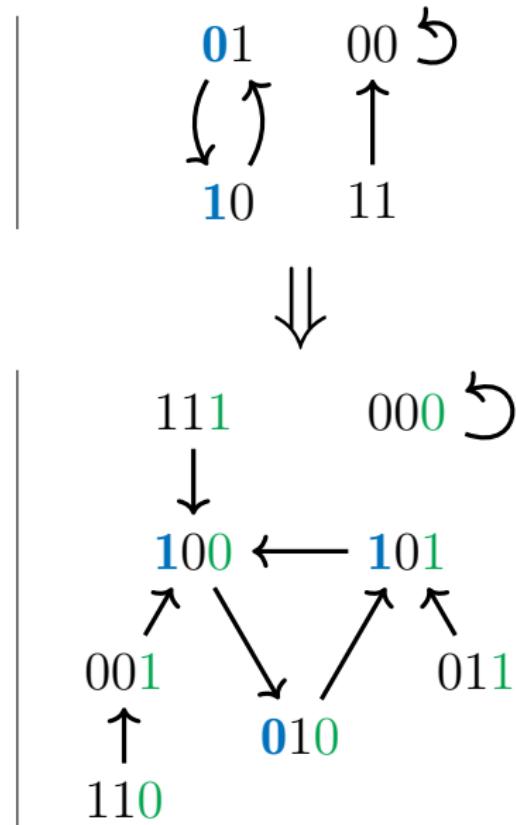
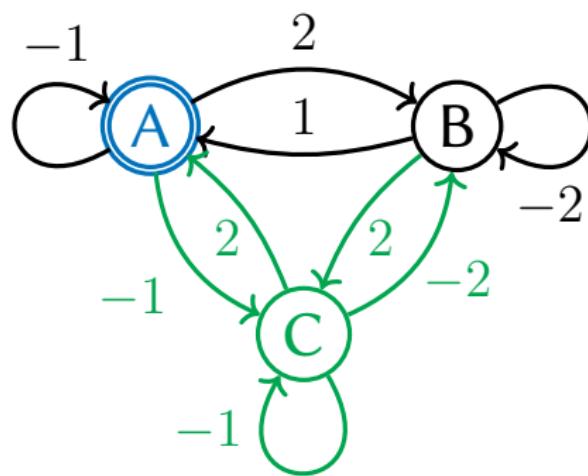
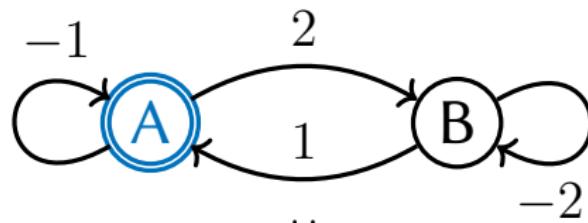
How can an SBN gain a new function, maintaining

- the prior structures? (interaction graph)
- the prior functions?

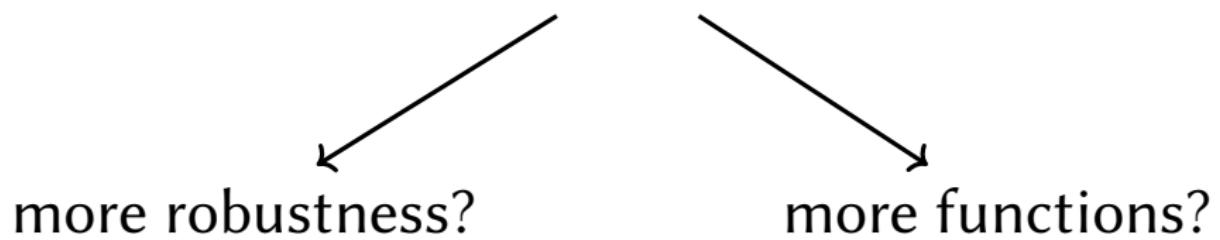


commutative diagram

Extensibility in SBN: an example



Growth of a network

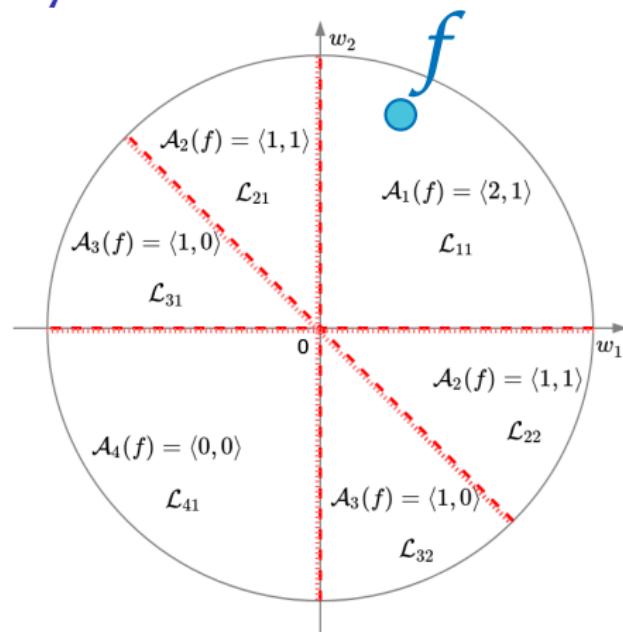


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Which structures are the most extensible?

Complexity is a good way to discriminate between SBN.

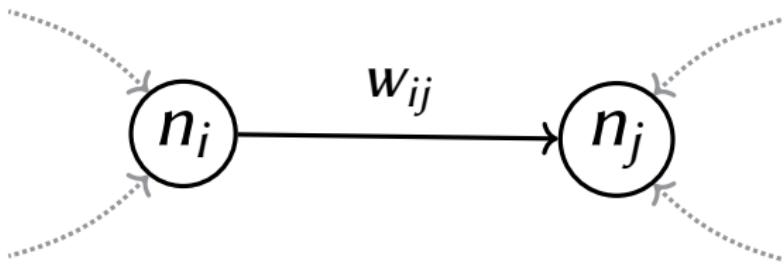
Complexity of an SBF



$\mathcal{C}^f = \text{complexity}(f) = \text{probability of not picking } f$
 \hookrightarrow uniform distribution over the unit ball

Complex function \implies low probability

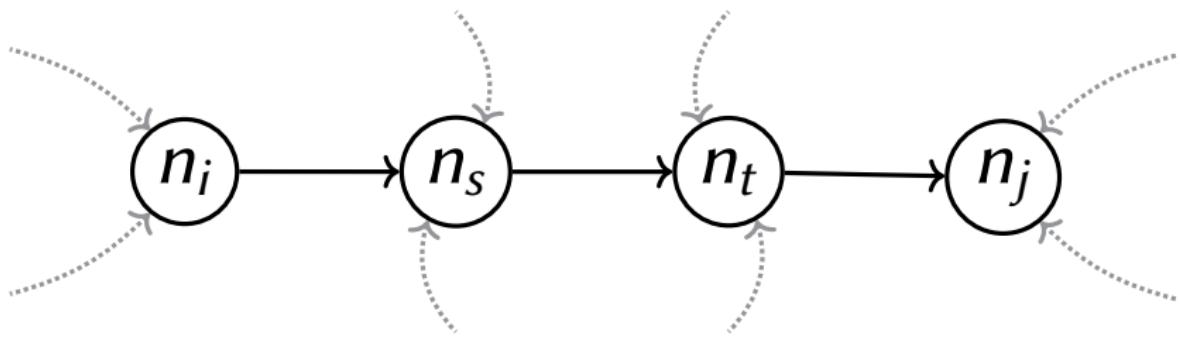
Complexity of an edge E_{ij} in SBN



The probability that n_i influences n_j via E_{ij} :

$$P^I(E_{ij}) = \frac{|w_{ij}|}{\sum_{k=1}^d |w_{kj}|}$$

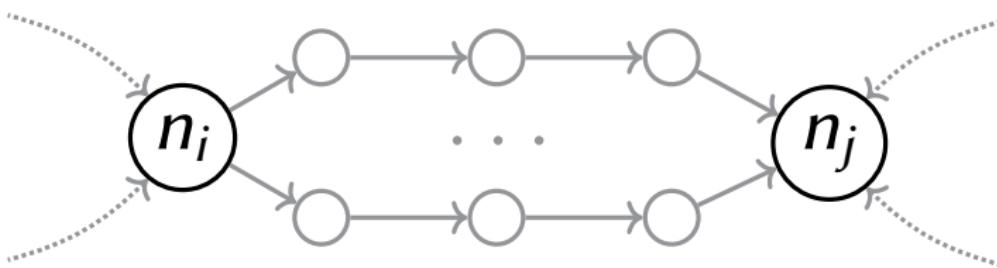
Complexity of a path in SBN



normalize by the length of the path

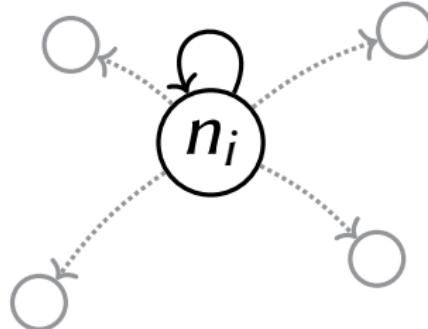
$$P^I(Path_{ab}) = \sqrt[3]{P^I(E_{is}) \cdot P^I(E_{st}) \cdot P^I(E_{tj})}$$

Probability that n_i influences n_j



$$P_{ab}^I = P \left(\bigcup_{Path \in Paths_{ij}} Path \right)$$

A centrality of a node



centrality(n_i) = probability that n_i influences at least one other node, including itself

$$C_i^s = P \left(\bigcup_{k=1}^d \left(\bigcup_{Path \in Paths_{ik}} Path \right) \right)$$

Complexity of an SBN

Product of complexities of connected nodes, modulated by their centralities:

$$\mathcal{C}(SBN) = \prod_{i=1}^d \begin{cases} \mathcal{C}_i^f \times \mathcal{C}_i^s & \text{if } \mathcal{C}_i^s > 0 \\ 1 & \text{otherwise} \end{cases}$$

1

Threshold/Sign Boolean networks (BN)

2

Extensibility in Sign BN

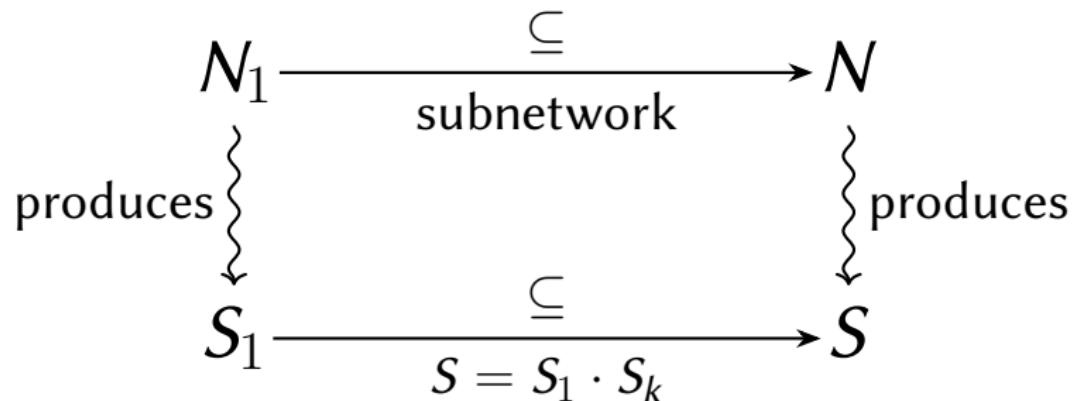
3

Complexity of Sign BN

4

Preliminary results and analysis

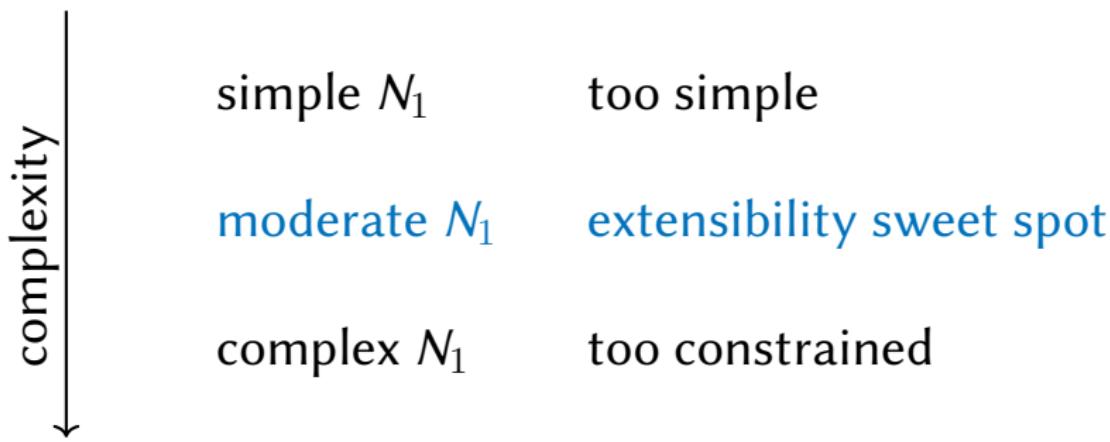
Extensibility = function(complexity)



Which N_1 yield the most extensions?

Hypothesis

Which N_1 yield the most extensions?



Testing the hypothesis for $d = 2$

- ① Enumerate all N_1 .
- ② For each N_1 , fix an output node, S_1 , and S .
- ③ Enumerate all $\textcolor{blue}{N}$ for which the diagram commutes:

$$\begin{array}{ccc} N_1 & \xrightarrow{\subseteq} & \textcolor{blue}{N} \\ \downarrow & & \downarrow \\ S_1 & \xrightarrow{\subseteq} & S \end{array}$$

- ④ Compute the complexities of N_1 , $\textcolor{blue}{N}$, S_1 , and S .

a version of Kolmogorov complexity
adapted to short strings

Soler-Toscano F., Zenil H., Delahaye J.-P., Gauvrit N.: Calculating Kolmogorov complexity from the output frequency distributions of small Turing machines; PLoS ONE 9(5): e96223 (2014)

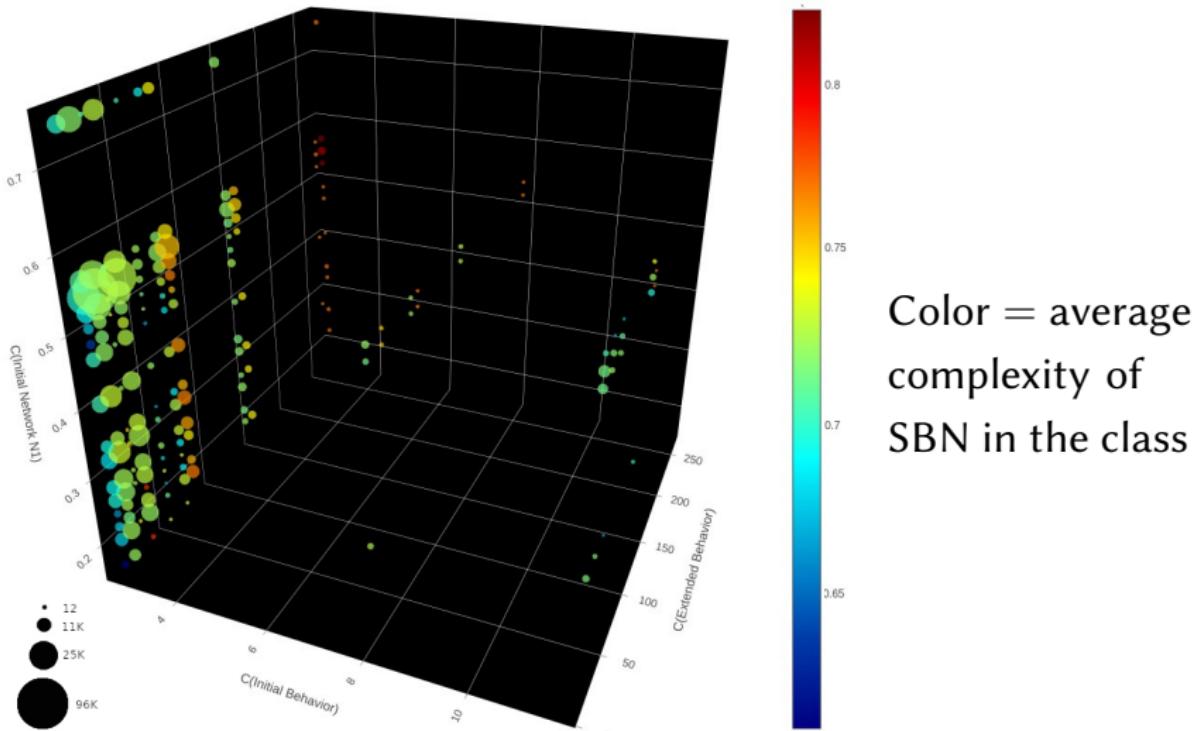
Combinatorics of the method

d	SBF	SBN	$\{N_1, S_1, S\}$	N
2	7	101	$96 \cdot 10^3$	777216
3	17	$206 \cdot 10^3$	$[180 \cdot 10^6 - 32 \cdot 10^9]$	$[1.5 \cdot 10^9 - 280 \cdot 10^9]$
4	47	$76 \cdot 10^9$		

estimations



First results



Discussion

- Behaviours of most $d = 2$ networks are not complex.
- Most extensions have moderate structural and behavioural complexity.
- Complex extended behaviors and networks are mostly obtained from networks of moderate complexity.

Improvement directions

- Refine the analysis.
- Look into higher dimensions.
- Consider other modes of evolution:
 - ▶ increase in the number of attractors,
 - ▶ increase in the size of basins of attraction.

simple $\xrightarrow[\text{??}]{\nabla}$ complex

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Laurent Trilling



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Grenoble Alpes



Île de France

