

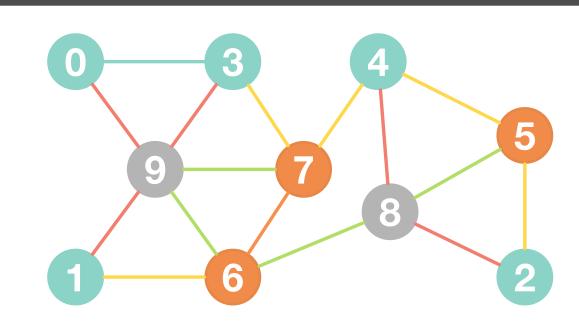
Parallel Graph Algorithms by Blocks

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Introduction



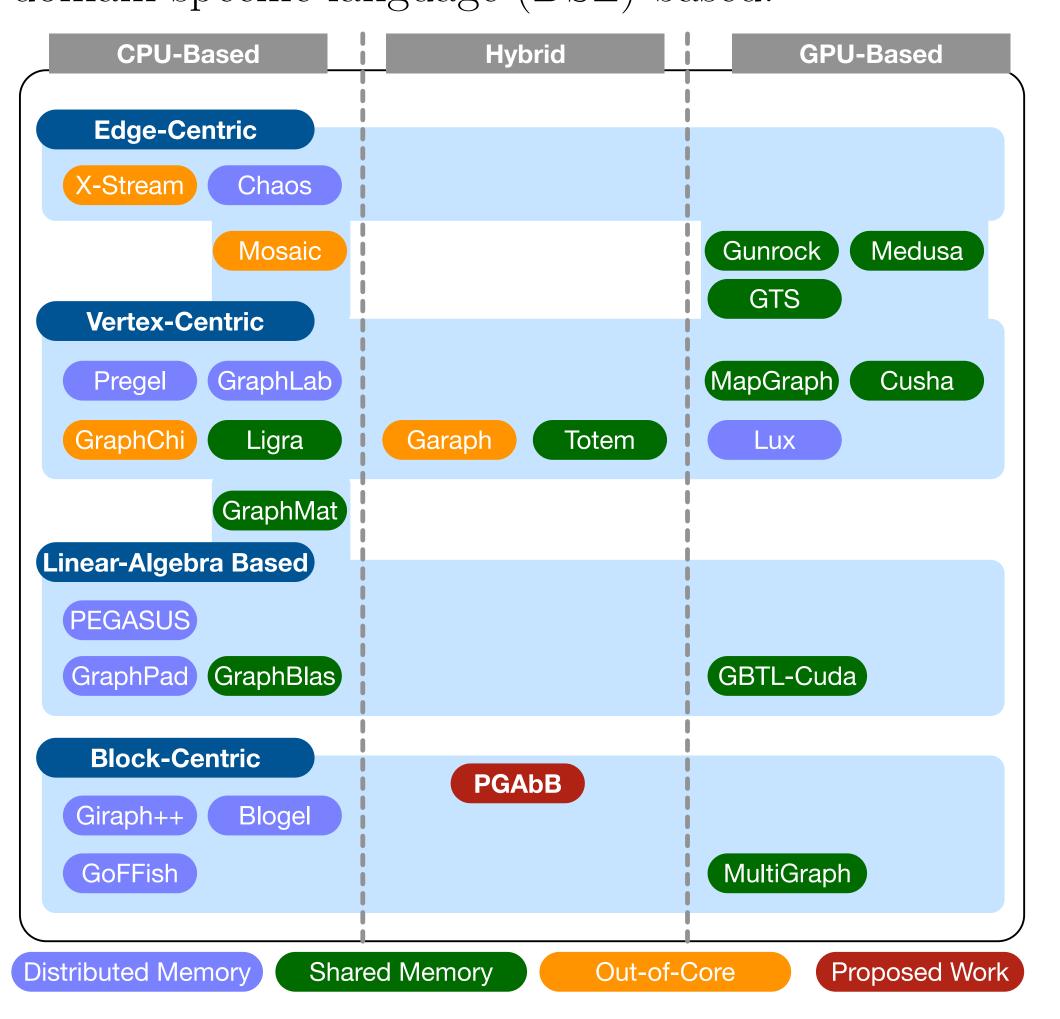
Graphs are very useful data-structures. They can represent different applications, such as social network, biological networks, and scientific simulations.

- Graphs are growing to billions of vertices and edges
- Fast, efficient analysis is crucial
- Many graph processing frameworks have been proposed

Goal: Developing an algorithmic framework (PGAbB) in which people can implement architecture agnostic graph algorithms that perform well on different architectures. Here we show our approach for the triangle counting problem as a use case.

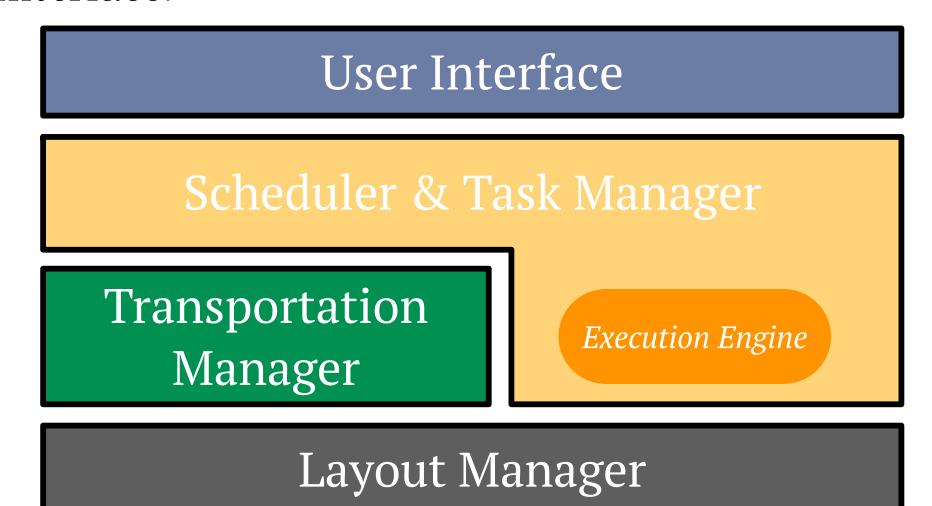
Graph Processing Systems

Programming abstractions of current systems can be broadly divided into five categories: vertex-centric; edge-centric; block-centric; linear-algebra based and domain-specific-language (DSL) based.



Parallel Graph Algorithm by Blocks (PGAbB)

PGAbB targets multi-core shared memory machines (host) with GPU accelerators (devices) and is composed by four components; Layout Manager, Scheduler, Transportation Manager Scheduler&Task Manager and User Interface.



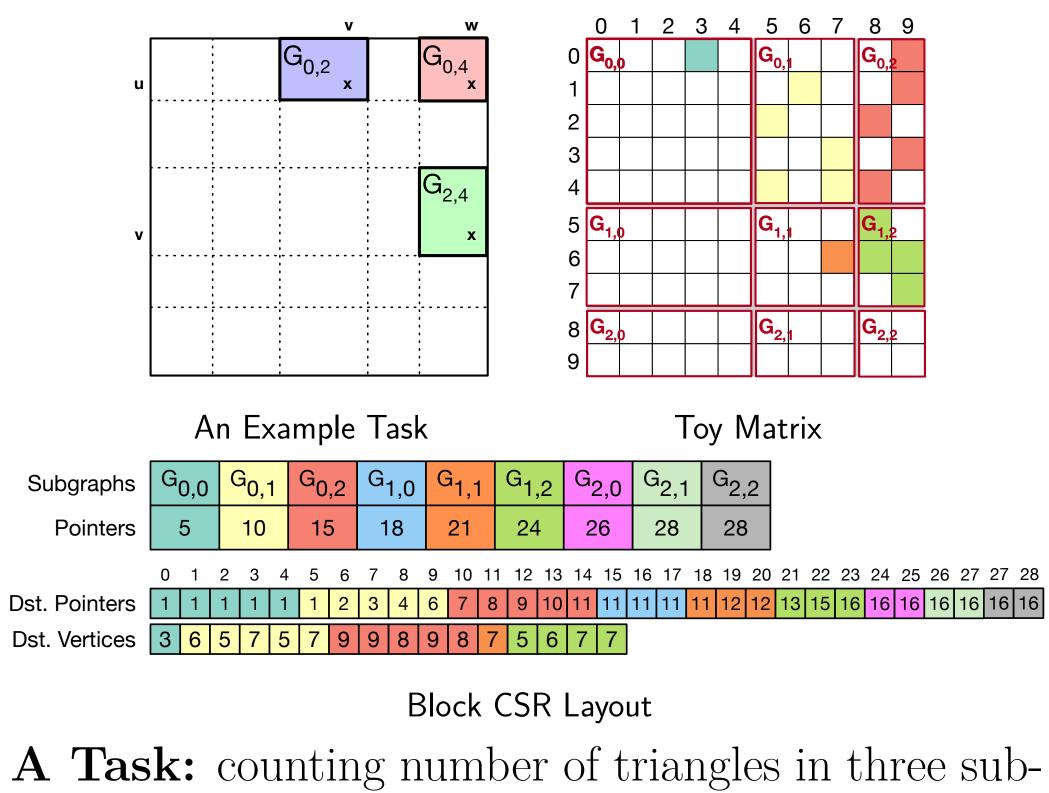
- Layout Manager: (1) Row-major, and (2) Hilbert order.
- Partitioner: Symmetric Rectilinear Partitioning.
- Blocks: (1) CSR and (2) COO style storage.
- Transporter: Communication overlapping.
- Scheduler&Task Manager: (1) locality based, and (2) workload estimation based.
- User Interface: Block-centric API and tools.

Use Case: Triangle Counting

Given a graph $G = \{V, E\}$, the triangle counting problem is to find the number (T) of all sets of three vertices, $u, v, w \in V$, such that:

 $T = |\{u, v, w \mid (u, v), (v, w), (w, u) \in E\}|.$

A Block-Based Triangle Counting Algorithm (bbTC)



A Task: counting number of triangles in three subgraphs; $G_{i,j}$, $G_{j,k}$ and $G_{i,k}$ where $i \leq j \leq k$.

Algorithm 1: $bbTC(G_{i,j}, G_{j,k}, G_{i,k})$

 $\tau \leftarrow 0$ > Initialize number of triangles to 0

for each $u \in V_S(G_{i,j})$ do for each $v \in N(G_{i,j}, u)$ do

Heterogeneous Execution

To utilize the massive parallelism on GPUs:

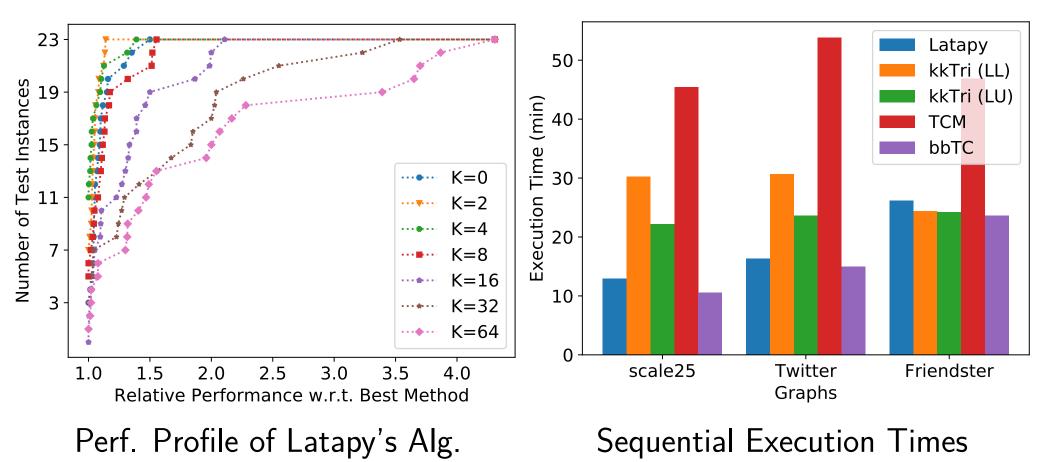
- Tasks are ordered based on their estimations.
- GPUs: starts with the heaviest tasks, and CPUs start with lighter ones
- © CUDA streams: simultaneous task execution and communication.

Triangle Counting Algorithms

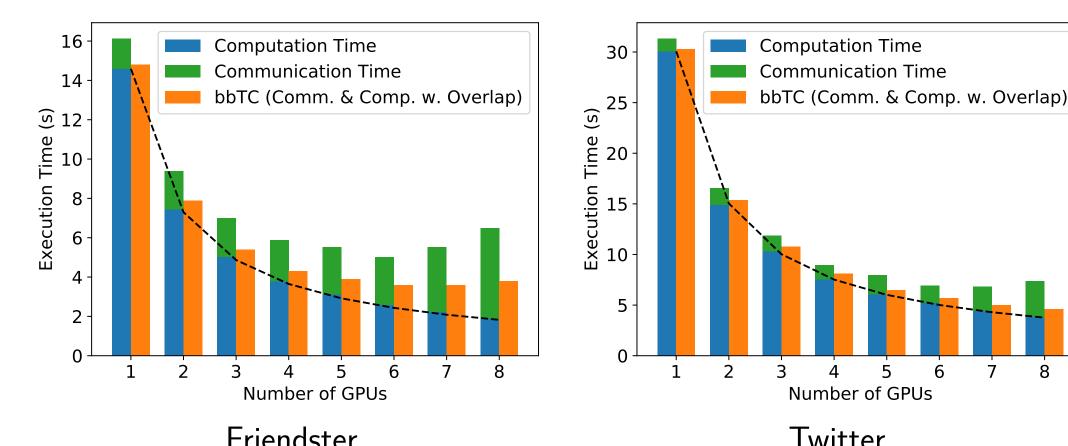
	TCM	kkTri	TriCore	bbTC	
Algorithmic properties					
List Intersect.	√	X	X	√	
Map Intersect.	X	\checkmark	X	\checkmark	
Search Intersect.	X	X		X	
Multi-Core	√	\checkmark	X	√	
Multi-GPU	X	X		\checkmark	
Complexity	$O(m^{rac{3}{2}})$	$O(m^{rac{3}{2}})$	$O(m^{\frac{3}{2}}\log\sqrt{m})$	$O(rac{\lambda p}{c}m^{rac{3}{2}})$	
Parallelization strategy					
One-dimension	√	√	√	X	
Two-dimension	X	X		\checkmark	
Used	runtime	es for im	plementation		
Pthread based	√	X	X	X	
OpenMP	√	\checkmark	X	\checkmark	
Cilk	√	\checkmark	X	\checkmark	
TBB	X	X	X	\checkmark	
CUDA	X	X		\checkmark	

Experiments

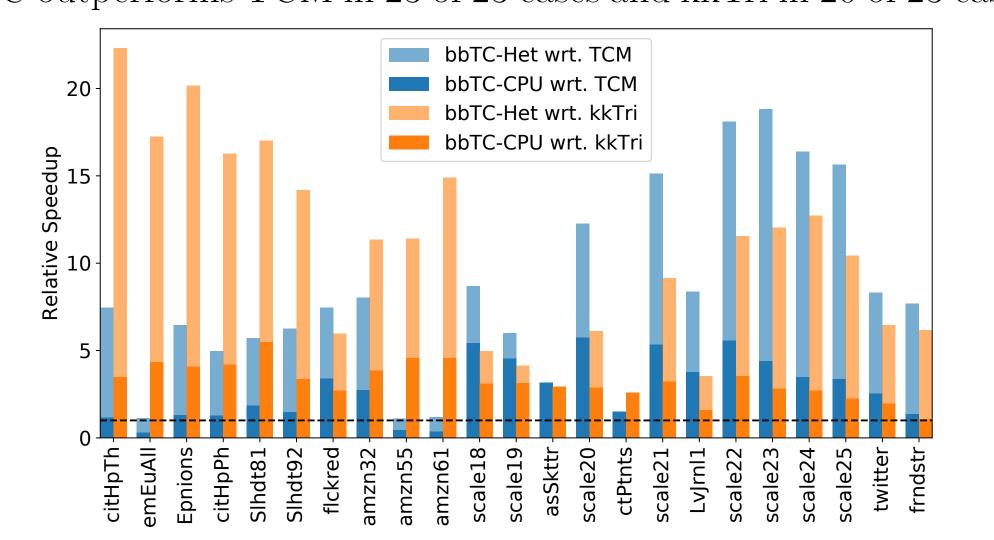
bbTC has a better memory utilization and outperforms the best sequential algorithm.



Overlap communication with computation with 10% overhead (5 GPUs).



bbTC outperforms TCM in 23 of 23 cases and kkTri in 20 of 23 cases.



bbTC outperforms kkTri and TCM 1.6× and 1.8 times, and processes the WDC-2014 graph in 116 seconds.

	Haswell	\mathbf{DGX}	Newell
kkTri	423	Out of memory	Out of memory
\mathbf{TCM}	476	Out of memory	Out of memory
$\mathbf{TriCore}$	N/A	_	_
\mathbf{bbTC}	265	116	1120

Table: Execution times (seconds) on a WDC graph.

References

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- [4] M. M. Wolf, M. Deveci, J. W. Berry, S. D. Hammond, and S. Rajamanickam. Fast linear algebra-based triangle counting with kokkoskernels.

^{*}This work was partially supported by the NSF grant CCF- 1919021.