Mapping ER to SQL

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# Mapping ER to SQL

We have explored mapping ER designs to relational schemas

SQL schemas are essentially more detailed versions of relational schemas

The mapping is much the same, except that

- you need to provide more details on allowed values
- you can map some ideas from ER that are not in relational schemas

There are also some ideas from ER than do not map to an SQL schema

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## Reminder: SQL/Relational Model vs ER Model

#### Correspondences between SQL/relational and ER data models:

- attribute(ER) ≅ attribute(Rel), entity(ER) ≅ row/tuple(Rel)

#### Differences between SQL and ER models:

- SQL uses tables to model entities *and* relationships
- SQL has no composite or multi-valued attributes (only atomic)
- SQL has no object-oriented notions (e.g. subclasses, inheritance)

#### Note that ...

- not all aspects of ER can be represented exactly in an SQL schema
- some aspects of SQL schemas (e.g. domains) do not appear in ER

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# Mapping ER to SQL

Some conventions that we use in mapping ER to SQL

- stop using upper-case for SQL keywords (use table vs TABLE)
- all tables based on entities are given plural names
- attributes in entities are given the same name in ER and SQL
- attributes in relationships are given the same name in ER and SQL
- ER key attributes are defined using **primary key**
- text-based attributes are defined with type **text**, unless there is a size which is obvious from the context
- attribute domains can be PostgreSQL-specific types where useful
- foreign keys within entity tables are named after the relationship
- foreign keys in relationship tables are named table\_id

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## Mapping Strong Entities

An entity set E with atomic attributes  $a_1, a_2, ... a_n$ 

maps to

A table R with attributes (columns)  $a_1, a_2, ... a_n$ 

### Example:

# StulD name Student degree

ER Model

#### SQL Version

```
create table Students (
stuID integer primary key,
name text not null,
degree char(4)
):
```

Note: the key is preserved in the mapping.

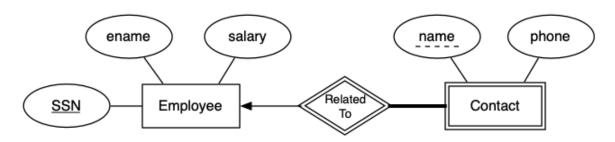
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## Mapping Weak Entities

#### Example:

#### ER Model

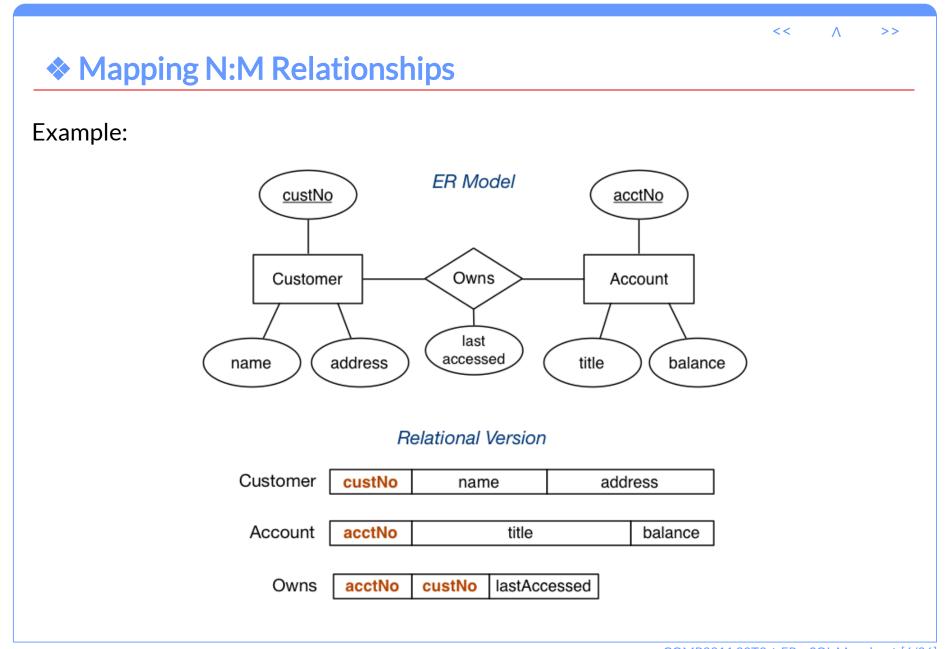


#### SQL Version

```
create table Employees (
    SSN text primary key,
    ename text,
    salary currency
);
```

```
create table Contacts (
    relatedTo text not null, -- total participation
    name text, -- not null implied by PK
    phone text not null,
    primary key (relatedTo, name),
    foreign key (relatedTo) references Employees (ssn)
);
```

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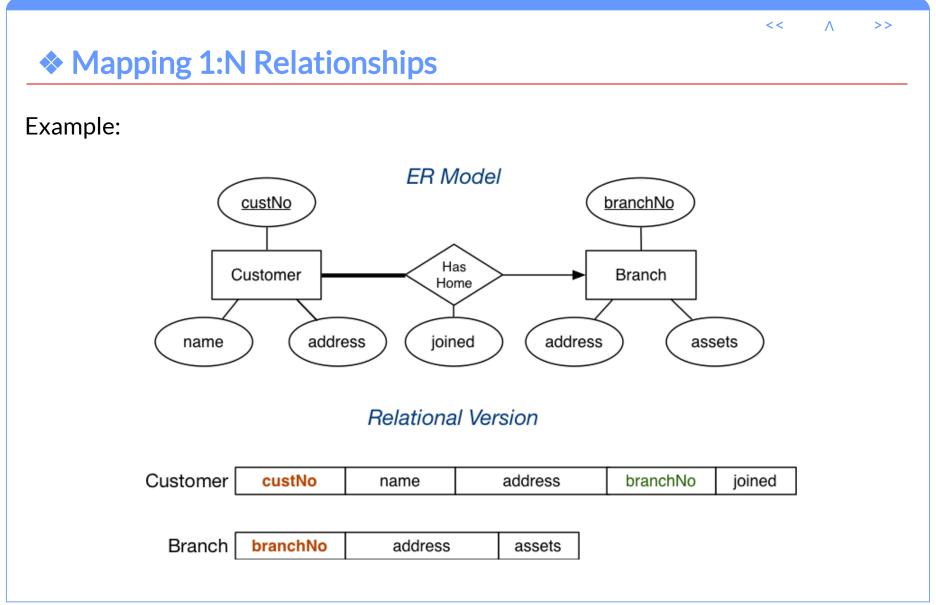


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## Mapping N:M Relationships (cont)

```
create table Customers (
    custNo serial primary key,
    name text not null,
    address text -- don't need to know customer's address
);
create table Accounts (
    acctNo char(5) check (acctNo \sim '[A-Z]-[0-9]{3}'),
    title text not null, -- acctNos are like 'A-123'
    balance float default 0.0,
    primary key (acctNo)
);
create table Owns (
    customer id integer references Customers(custNo),
    account_id char(5) references Accounts(acctNo),
    last_accessed timestamp,
    primary key (customer id, account id)
);
```

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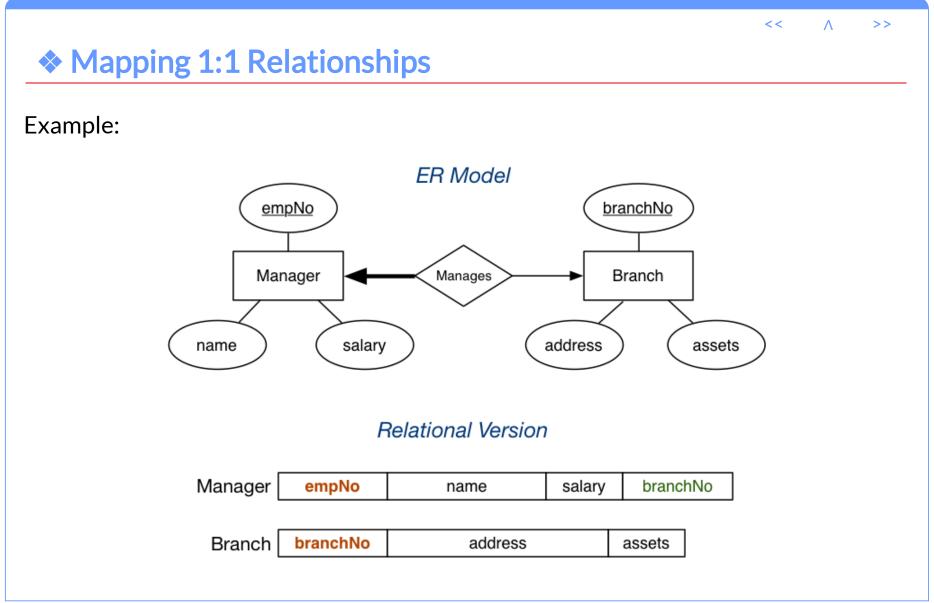
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## Mapping 1:N Relationships (cont)

```
create table Branches (
    branchNo serial primary key,
    address text not null,
    assets currency
);
create table Customers (
    custNo serial primary key,
    name text not null,
    address text,
    hasHome integer not null, -- total participation
    joined date not null,
    foreign key (hasHome) references Branches(branchNo)
);
```

hasHome implements the 1:n relationship; not null implements total participation

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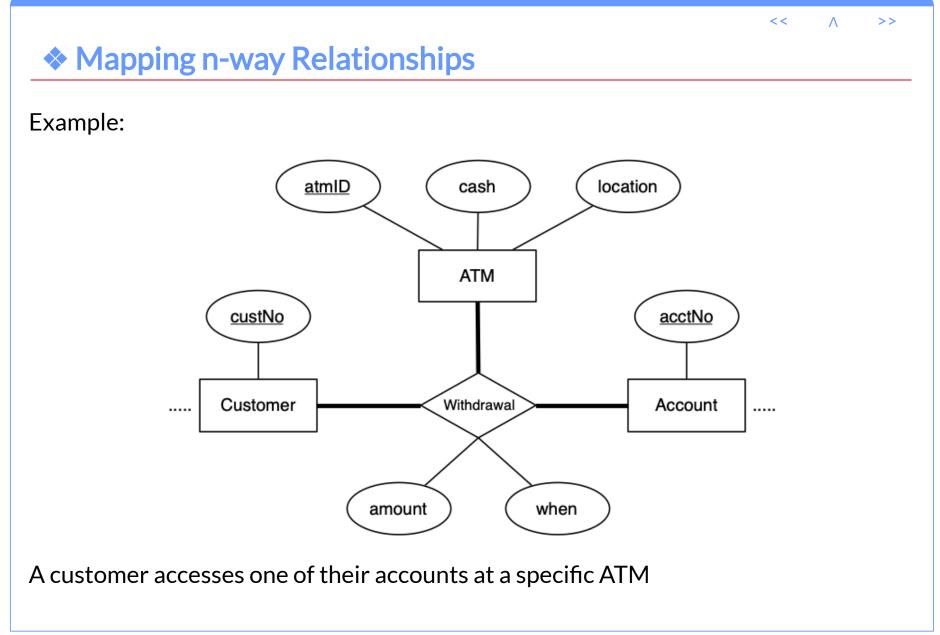
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## **♦ Mapping 1:1 Relationships** (cont)

If both entities have total participation, cannot express this in SQL except by putting a (redundant) **not null** foreign key in one table

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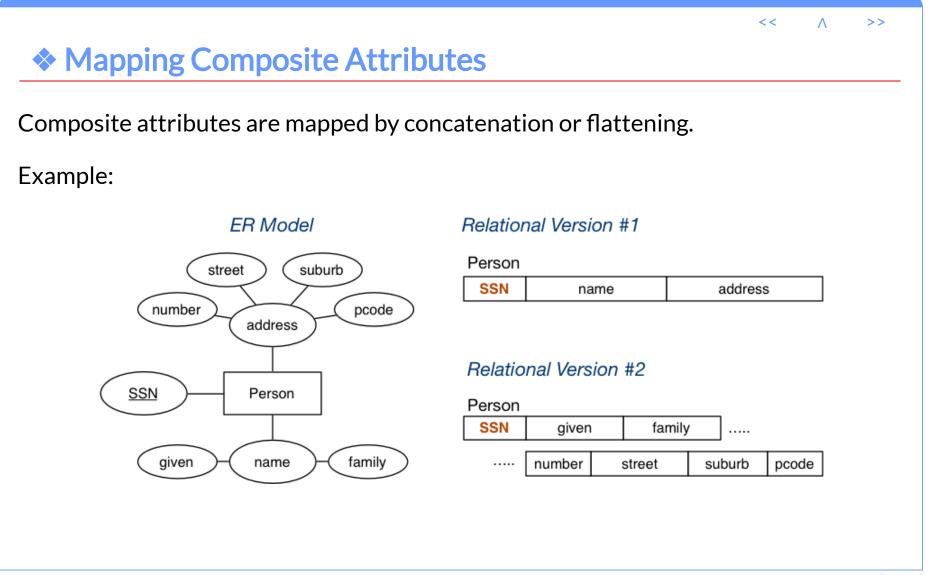


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## Mapping n-way Relationships (cont)

```
create table Customers (
    custNo serial primary key, ...
);
create table Accounts (
    acctNo char(5) ... primary key, ...
);
create table ATMs (
    atmID serial primary key,
    cash currency check (cash >= 0),
    location text not null
);
create table Withdrawal (
    customer id integer references Customers(custNo),
    account id char(5) references Accounts(acctNo),
                integer references ATMs(atmID),
    atm id
    amount currency not null,
    when timestamp default now(),
    primary key (customer_id,account_id,atm_id)
```

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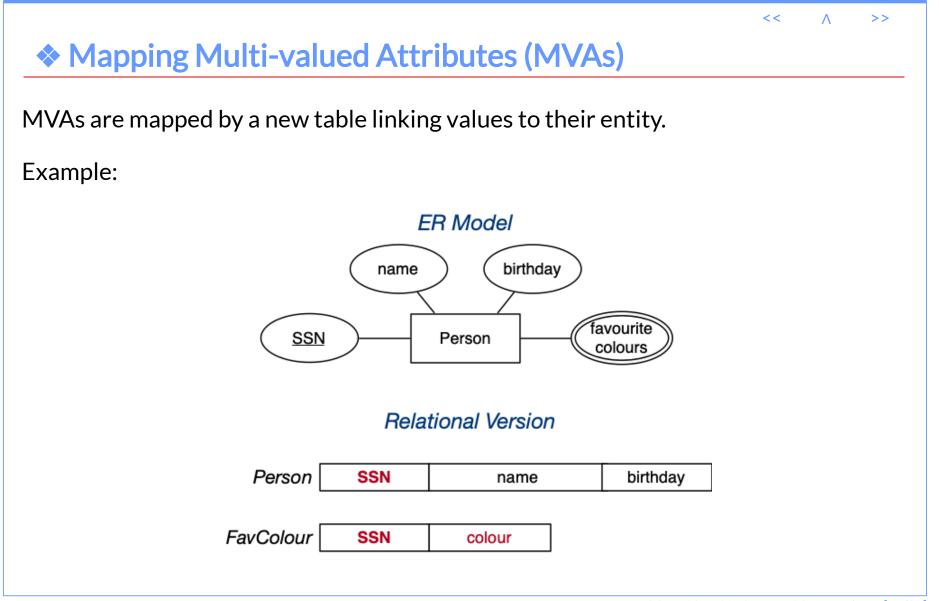
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## Mapping Composite Attributes (cont)

```
-- Version 1: concatenated
 create table People (
              integer primary key,
      ssn
      name text not null.
      address text not null
 );
 -- Version 2: flattened
 create table People (
              integer primary key,
      ssn
     given text not null,
     family text,
      number integer not null,
      street text not null,
      suburb text not null,
             char(4) not null check (pcode \sim '[0-9]{4}')
      pcode
 );
address = (number::text||' '||street||', '||suburb||' '||pcode)
Searching: suburb = 'Coogee' vs address like '%Coogee%'
Sorting: order by family vs can't be done (easily)
```

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## Mapping Multi-valued Attributes (MVAs) (cont)

```
create table People (
    ssn integer primary key,
    name text not null,
    birthday date
);
create table FavColour (
    person_id integer references People(ssn),
    colour text,
    primary key (person_id,colour)
);
```

Note that **colour** is implicitly **not null** because it is part of the primary key

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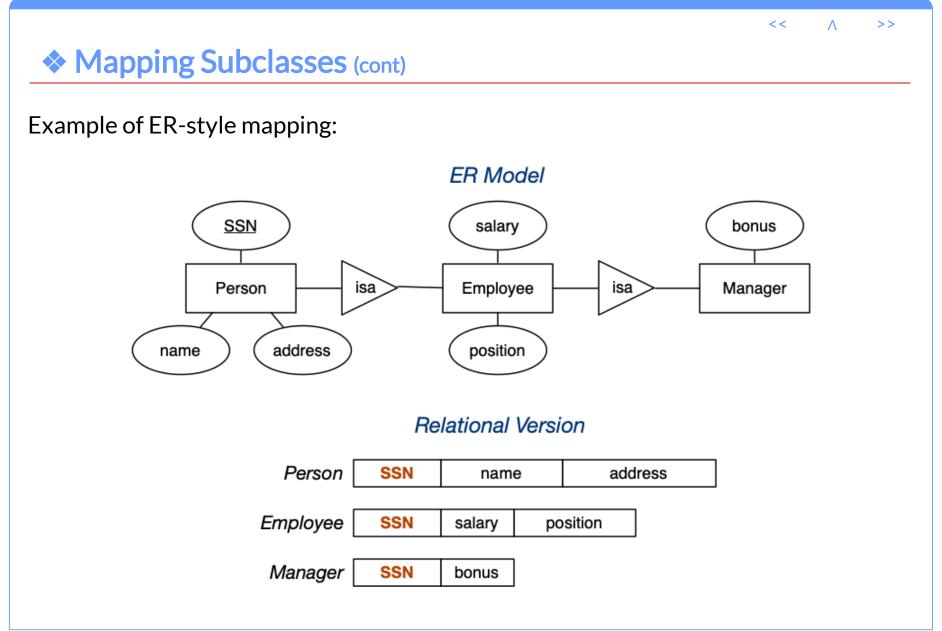
# Mapping Subclasses

Three different approaches to mapping subclasses to tables:

- ER style
  - each entity becomes a separate table,
  - containing attributes of subclass + FK to superclass table
- object-oriented
  - each entity becomes a separate table,
  - inheriting all attributes from all superclasses
- single table with nulls
  - whole class hierarchy becomes one table,
  - containing all attributes of all subclasses (null, if unused)

Which mapping is best depends on how data is to be used.

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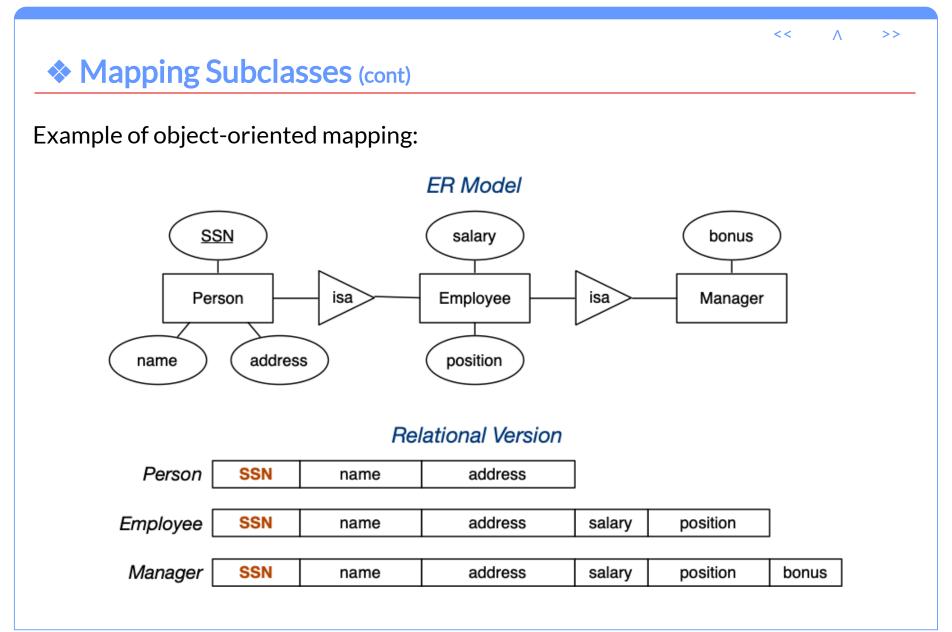
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## Mapping Subclasses (cont)

```
create table People (
           integer primary key,
    ssn
    name text not null,
    address text
);
create table Employees (
    person_id integer primary key,
    salary currency not null,
    position text not null,
    foreign key (person_id) references People(ssn)
create table Managers (
    employee_id integer primary key,
    bonus
                currency,
    foreign key (employee_id)
                references Employees(person id)
);
```

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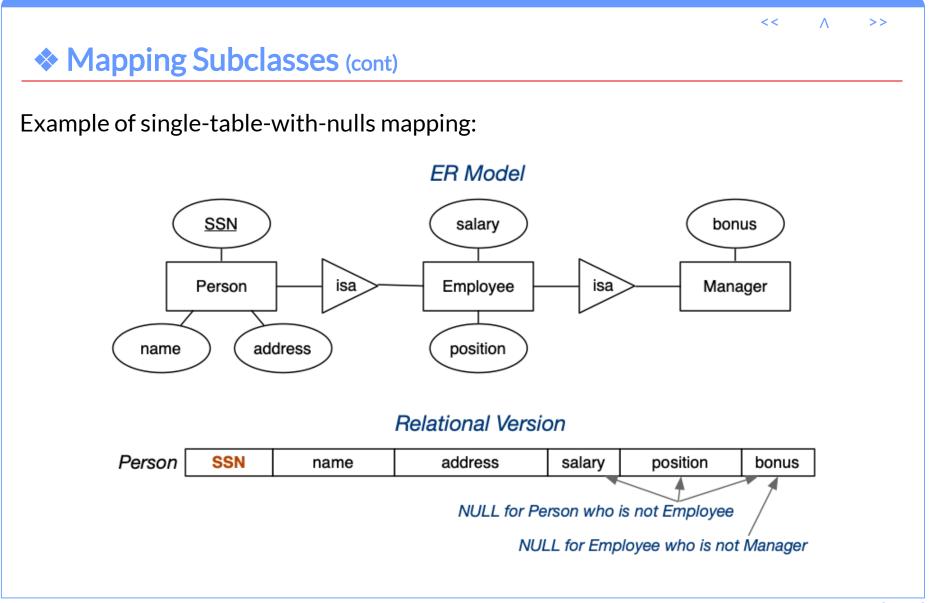


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# Mapping Subclasses (cont)

```
create table People (
           integer primary key,
   ssn
   name text not null,
   address text
);
create table Employees (
             integer primary key,
   ssn
   name text not null,
   address text
   salary currency not null,
   position text not null,
   foreign key (snn) references People(ssn)
);
create table Managers (
             integer primary key,
   ssn
         text not null,
   name
   address text
   salary currency not null,
   position text not null,
   bonus
             currency,
   foreign key (snn) references People(ssn)
);
```

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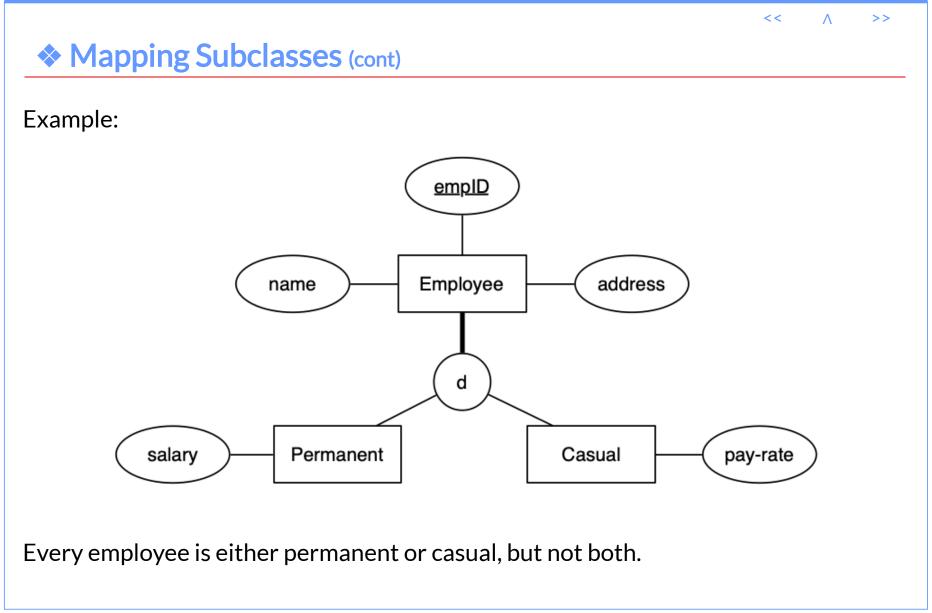


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## Mapping Subclasses (cont)

```
create table People (
              integer primary key,
    ssn
             char(1) not null
    ptype
                       check (ptype in ('P', 'E', 'M')),
              text not null,
    name
    address
              text
    salary currency,
    position text,
    bonus
              currency,
    constraint subclasses check
               ((ptype = 'P' and salary is null
                and position is null and bonus is null)
               or
                (ptype = 'E' and salary is not null
                 and position is not null and bonus is null)
               or
                (ptype = 'M' and salary is not null
                 and position is not null and bonus is not null))
);
```

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## Mapping Subclasses (cont)

ER-style mapping to SQL schema:

```
create table Employees (
    empID serial primary key,
    name text not null,
    address text not null
);
create table Permanents (
    employee_id integer primary key,
    salary currency not null,
    foreign key (employee_id) references Employees(empID)
);
create table Casuals (
    employee_id integer primary key,
    pay_rate currency not null,
    foreign key (employee_id) references Employees(empID)
);
```

Does not capture either participation or disjoint-ness constraints!

Would need to program a solution to this e.g web-form that requires user to enter both Employee and subclass info

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