

# Mergesort

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# ❖ Mergesort

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## Mergesort: basic idea

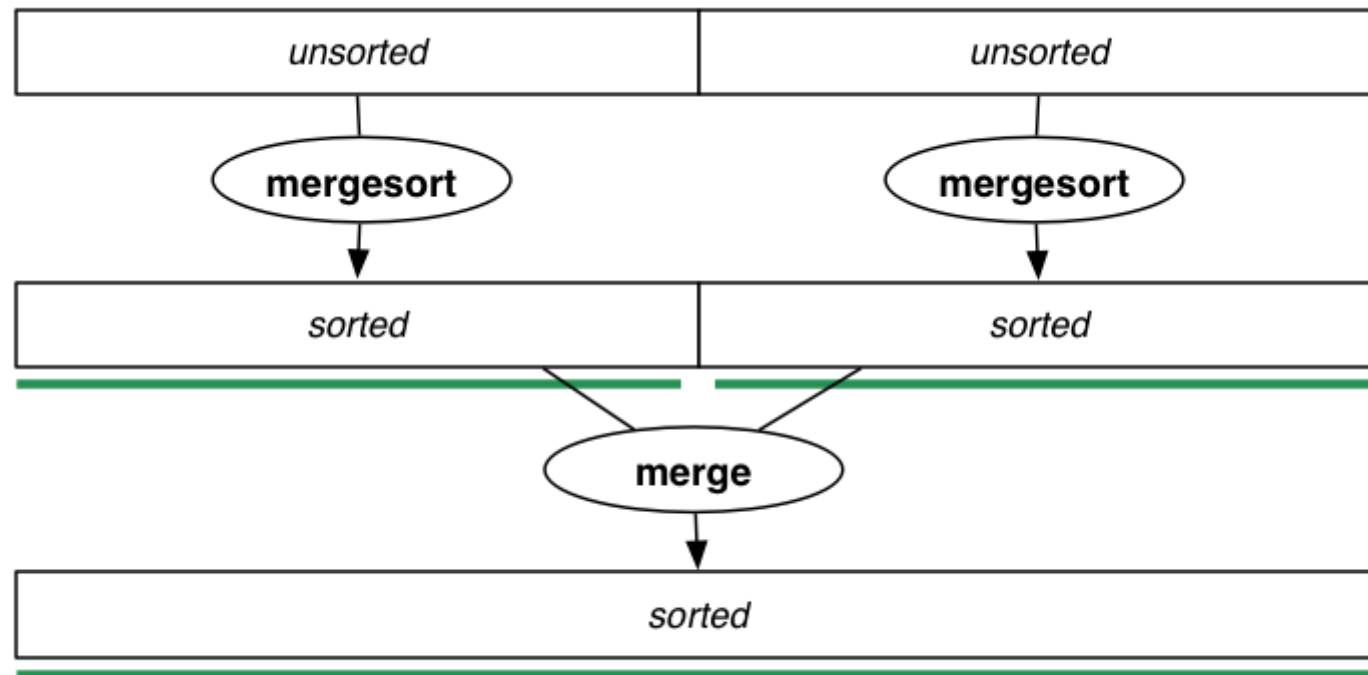
- split the array into two equal-sized partitions
- (recursively) sort each of the partitions
- merge the two partitions into a new sorted array
- copy back to original array

## Merging: basic idea

- copy elements from the inputs one at a time
- give preference to the smaller of the two
- when one exhausted, copy the rest of the other

## ❖ ... Mergesort

### Phases of mergesort



## ❖ Mergesort Implementation

Mergesort function:

```
void mergesort(Item a[], int lo, int hi)
{
    int mid = (lo+hi)/2; // mid point
    if (hi <= lo) return;
    mergesort(a, lo, mid);
    mergesort(a, mid+1, hi);
    merge(a, lo, mid, hi);
}
```

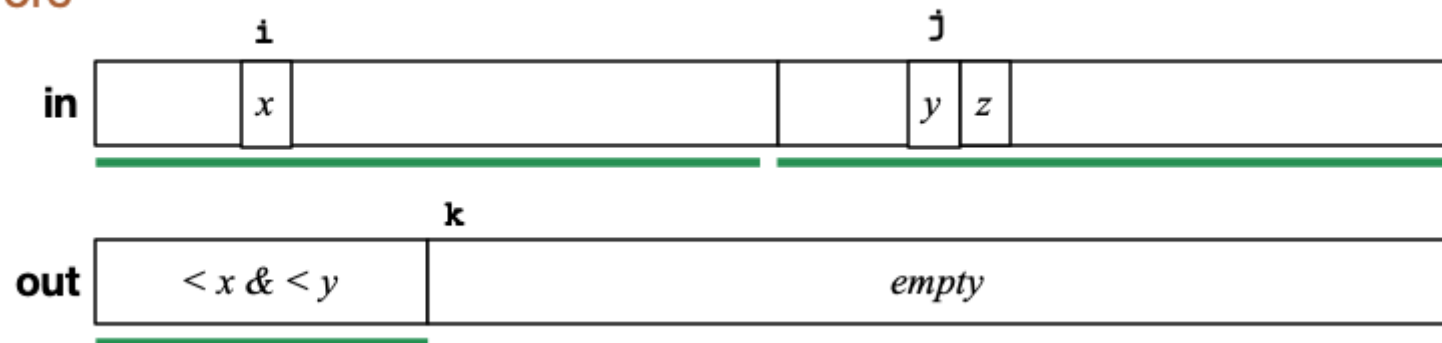
Example of use (**typedef int Item**):

```
int nums[10] = {32,45,17,22,94,78,64,25,55,42};
mergesort(nums, 0, 9);
```

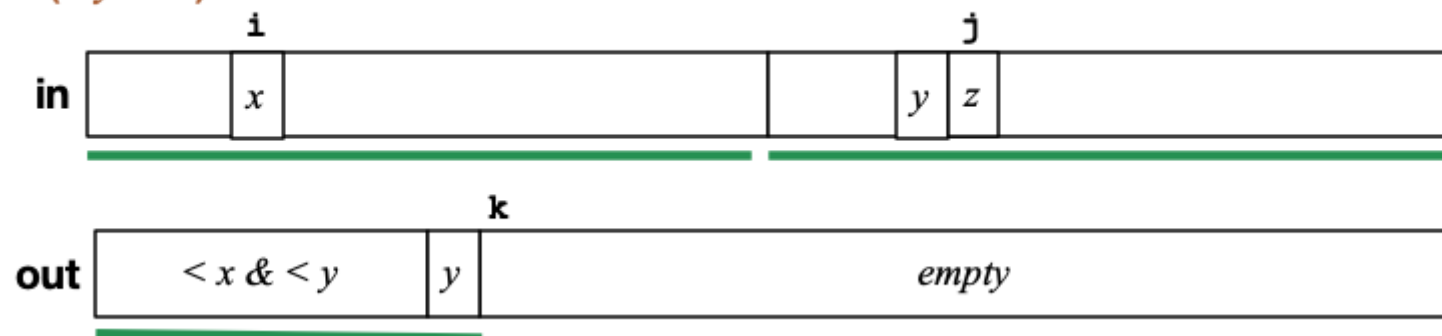
## ❖ ... Mergesort Implementation

One step in the merging process:

Before



After (if  $y < x$ )



## ❖ ... Mergesort Implementation

Implementation of merge:

```
void merge(Item a[], int lo, int mid, int hi)
{
    int i, j, k, nitems = hi-lo+1;
    Item *tmp = malloc(nitems*sizeof(Item));

    i = lo; j = mid+1; k = 0;
    // scan both segments, copying to tmp
    while (i <= mid && j <= hi) {
        if (less(a[i],a[j]))
            tmp[k++] = a[i++];
        else
            tmp[k++] = a[j++];
    }
    // copy items from unfinished segment
    while (i <= mid) tmp[k++] = a[i++];
    while (j <= hi) tmp[k++] = a[j++];

    //copy tmp back to main array
    for (i = lo, k = 0; i <= hi; i++, k++)
        a[i] = tmp[k];
    free(tmp);
}
```



## ❖ Mergesort Performance

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Best case:  $O(N \log N)$  comparisons

- split array into equal-sized partitions
- same happens at every recursive level
- each "level" requires  $\leq N$  comparisons
- halving at each level  $\Rightarrow \log_2 N$  levels

Worst case:  $O(N \log N)$  comparisons

- partitions are exactly interleaved
- need to compare all the way to end of partitions

Disadvantage over quicksort: need extra storage  $O(N)$



## ❖ Bottom-up Mergesort

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Non-recursive mergesort does not require a stack

- partition boundaries can be computed iteratively

Bottom-up mergesort:

- on each pass, array contains sorted **runs** of length  $m$
- at start, treat as  $N$  sorted runs of length 1
- 1st pass merges adjacent elements into runs of length 2
- 2nd pass merges adjacent 2-runs into runs of length 4
- continue until a single sorted run of length  $N$

This approach can be used for "in-place" mergesort.

## ❖ ... Bottom-up Mergesort

Original

[0]	[1]	[2]													[15]
A	S	O	R	T	I	N	G	E	X	E	M	P	L	A	R

After 1st pass  
sorted slices of length 2

A	S	O	R	I	T	G	N	E	X	E	M	L	P	A	R
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

After 2nd pass  
sorted slices of length 4

A	O	R	S	G	I	N	T	E	E	M	X	A	L	P	R
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

After 3rd pass  
sorted slices of length 8

A	G	I	N	O	R	S	T	A	E	E	L	M	P	R	X
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

After 4th pass  
sorted slice of length 16

A	A	E	E	G	I	L	M	N	O	P	R	R	S	T	X
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

## ❖ ... Bottom-up Mergesort

Bottom-up mergesort for arrays:

```
#define min(A,B) ((A < B) ? A : B)

void mergesort(Item a[], int lo, int hi)
{
    int m;      // m = length of runs
    int len;    // end of 2nd run
    Item c[];   // array to merge into
    for (m = 1; m <= lo-hi; m = 2*m) {
        for (int i = lo; i <= hi-m; i += 2*m) {
            len = min(m, hi-(i+m)+1);
            merge(&a[i], m, &a[i+m], len, c);
            copy(&a[i], c, m+len);
        }
    }
}
```

## ❖ ... Bottom-up Mergesort

```
// merge arrays a[] and b[] into c[]
// aN = size of a[], bN = size of b[]
void merge(Item a[], int aN, Item b[], int bN, Item c[])
{
    int i; // index into a[]
    int j; // index into b[]
    int k; // index into c[]
    for (i = j = k = 0; k < aN+bN; k++) {
        if (i == aN)
            c[k] = b[j++];
        else if (j == bN)
            c[k] = a[i++];
        else if (less(a[i], b[j]))
            c[k] = a[i++];
        else
            c[k] = b[j++];
    }
}
```

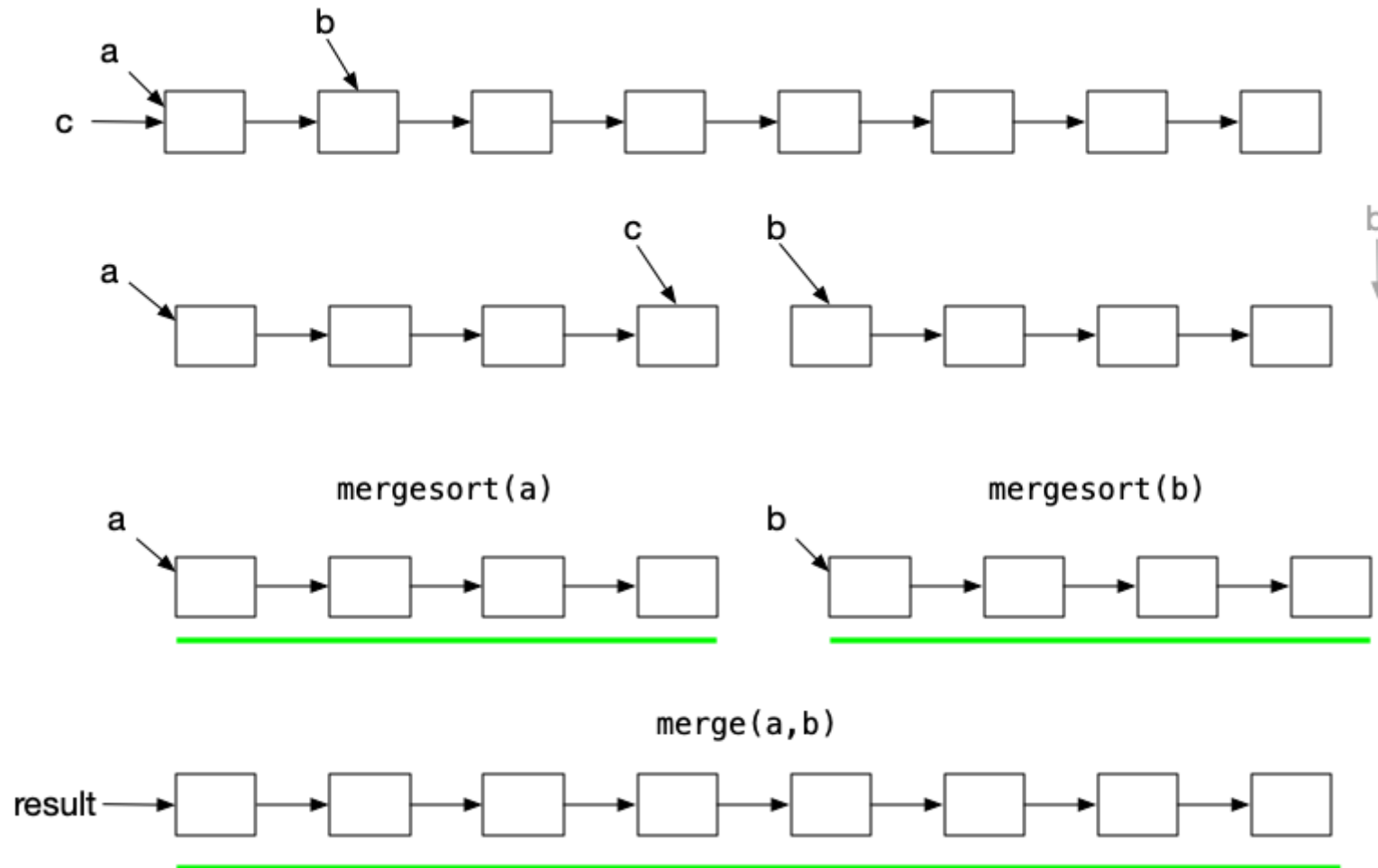
## ❖ Mergesort and Linked Lists

Merging linked lists is relatively straightforward:

```
List merge(List a, List b)
{
    List new = newList();
    while (a != NULL && b != NULL) {
        if (less(a->item, b->item))
            { new = ListAppend(new, a->item); a = a->next; }
        else
            { new = ListAppend(new, b->item); b = b->next; }
    }
    while (a != NULL)
        { new = ListAppend(new, a->item); a = a->next; }
    while (b != NULL)
        { new = ListAppend(new, b->item); b = b->next; }
    return new;
}
```

## ❖ ... Mergesort and Linked Lists

Mergesort method using linked lists



## ❖ ... Mergesort and Linked Lists

Recursive linked list mergesort, built with list merge:

```
List mergesort(List c)
{
    List a, b;
    if (c == NULL || c->next == NULL) return c;
    a = c; b = c->next;
    while (b != NULL && b->next != NULL)
        { c = c->next; b = b->next->next; }
    b = c->next; c->next = NULL; // split list
    return merge(mergesort(a), mergesort(b));
}
```

