# **SQL:** Queries on Multiple Tables

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# Queries on Multiple Tables

Queries involving a single table are useful.

Exploiting all data in the DB requires

- combining data from multiple tables
- typically involving primary/foreign key matching

**Example:** Which brewers makes beers that John likes?

```
select b.brewer
from Beers b join Likes L on (b.name = L.beer)
where L.drinker = 'John';
```

Info on brewers is in **Beers**; info on who likes what in **Likes**.

Need to combine info from both tables using "common" attributes

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# Queries on Multiple Tables (cont)

### Example **Beers** and **Likes** tuples:

```
Beers(80/-, Caledonian, Scotch Ale)
Beers(New, Toohey's, Lager)
Beers(Red Nut, Bentspoke, Red IPA)
Beers(Sculpin, Ballast Point, IPA)
Likes(John, Sculpin)
Likes(John, New)
Likes(John, 80/-)
```

"Merged" tuples resulting from

• Beers b join Likes L on (b.name = L.beer)

```
Joined(80/-, Caledonian, Scotch Ale, John, 80/-)
Joined(New, Toohey's, Lager, Adam, New)
Joined(Red Nut, Bentspoke, Red IPA, John, Red Nut)
Joined(Sculpin, Ballast Point, IPA, John, Sculpin)
```

In the query, the where clause removes all tuples not related to John

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Join is the SQL operator that combines tuples from tables.

Such an important operation that several variations exist

- natural join matches tuples via equality on common attributes
- equijoin matches tuples via equality on specified attributes
- theta-join matches tuples via a boolean expression
- outer join like theta-join, but includes non-matching tuples

We focus on theta-join and outer join in this course

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Join fits into **SELECT** queries as follows:

```
SELECT Attributes

FROM R1

JOIN R2 ON (JoinCondition<sub>1</sub>)

JOIN R3 ON (JoinCondition<sub>2</sub>)

...

WHERE Condition
```

Can include an arbitrary number of joins.

WHERE clause typically filters out some of the joined tuples.

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Alternative syntax for joins:

```
SELECT brewer

FROM Likes L, Beers b

WHERE L.beer = b.name

AND L.drinker = 'John';
```

Join condition(s) are specified in the WHERE clause

We prefer the explicit JOIN syntax, but this is sometimes more compact

Note: duplicates could be eliminated by using **distinct** 

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## Operational semantics of R1 JOIN R2 ON (Condition):

```
FOR EACH tuple t1 in R1 D0

FOR EACH tuple t2 in R2 D0

check Condition for current

t1, t2 attribute values

IF Condition holds THEN

add (t1,t2) to result

END

END

END
```

Easy to generalise: add more relations, include WHERE condition

Requires one tuple variable for each relation, and nested loops over relations. But this is not how it's actually computed!

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## Name Clashes in Conditions

#### If a **SELECT** statement

- refers to multiple tables
- some tables have attributes with the same name

use the table name to disambiguate.

**Example:** Which hotels have the same name as a beer?

```
SELECT Bars.name

FROM Bars, Beers

WHERE Bars.name = Beers.name;

-- or, using table aliases ...

SELECT r.name

FROM Bars r, Beers b

WHERE r.name = b.name
```

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# Explicit Tuple Variables

Table-dot-attribute doesn't help if we use same table twice in **SELECT**.

To handle this, define new names for each "instance" of the table

```
SELECT r1.a, r2.b FROM R r1, R r2 WHERE r1.a = r2.a
```

**Example:** Find pairs of beers by the same manufacturer.

```
SELECT b1.name, b2.name

FROM Beers b1 JOIN Beers b2 ON (b1.brewer = b2.brewer)

WHERE b1.name < b2.name;
```

The WHERE condition is used to avoid:

- pairing a beer with itself e.g. (New, New)
- same pairs with different order e.g. (New, Old) (Old, New)

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## Outer Join

Join only produces a result tuple from  $t_R$  and  $t_S$  where ...

- there are appropriate values in both tuples
- so that the join condition is satisfied

SELECT \* FROM R JOIN S WHERE (Condition)

Sometimes, we want a result for every R tuple

even if some R tuples have no matching S tuple

These kinds of requests often include "for each" or "for every"

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**Example**: for each suburb with a bar, find out who drinks there.

Theta-join only gives results for suburbs where people drink.

```
SELECT B.addr, F.drinker
FROM
      Bars B
       JOIN Frequents F ON (F.bar = B.name)
ORDER BY addr;
   addr
            drinker
 Coogee
            Adam
 Coogee
        John
 Kingsford | Justin
 Sydney
         Justin
 The Rocks |
            John
```

But what if we want all suburbs, even if some have are no drinkers?

This is from an older and simpler instance of the beers database.



An outer join solves this problem.

For R OUTER JOIN S ON (Condition)

- all "tuples" in *R* have an entry in the result
- if a tuple from R matches tuples in S, we get the normal join result tuples
- if a tuple from R has no matches in S, the attributes supplied by S are **NULL**

This outer join variant is called **LEFT OUTER JOIN**.

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Solving the example query with an outer join:

```
SELECT B.addr, F.drinker
FROM
      Bars B
      LEFT OUTER JOIN Frequents F on (F.bar = B.name)
ORDER BY B.addr;
   addr
            | drinker
Coogee | Adam
Coogee
        John
          | Justin
Kingsford
Randwick
Sydney
           | Justin
The Rocks
             John
```

Note that Randwick is now mentioned (because of the Royal Hotel).

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Operational semantics of R1 LEFT OUTER JOIN R2 ON (Cond):

```
FOR EACH tuple t1 in R1 D0
    nmatches = 0
    FOR EACH tuple t2 in R2 D0
        check Cond for current
           t1, t2 attribute values
        IF Cond holds THEN
            nmatches++
            add (t1,t2) to result
        END
    FND
    IF nmatches == 0 THEN
        t2 = (null, null, null,...)
        add (t1,t2) to result
    END
END
```

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Many RDBMSs provide three variants of outer join:

- R LEFT OUTER JOIN S
  - behaves as described above
- R RIGHT OUTER JOIN S
  - includes all tuples from *S* in the result
  - **NULL**-fills any *S* tuples with no matches in *R*
- R FULL OUTER JOIN S
  - includes all tuples from *R* and *S* in the result
  - those without matches in other relation are NULL-filled

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# Subqueries

The result of a query can be used in the WHERE clause of another query.

Case 1: Subquery returns a single, unary tuple

```
SELECT * FROM R WHERE R.a = (SELECT S.x FROM S WHERE Cond<sub>1</sub>)
```

Case 2: Subquery returns multiple values

```
SELECT * FROM R WHERE R.a IN (SELECT S.x FROM S WHERE Cond<sub>2</sub>)
```

This approach is often used in the initial discussion of SQL in some textbooks.

These kinds of queries can generally be solved *more efficiently* using a join

```
SELECT * FROM R JOIN S ON (R.a = S.x) WHERE Cond
```

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