

# Recursion

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# ❖ Recursion

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Recursion is a powerful problem-solving strategy

- employing a variant of divide-and-conquer
- leading to simple, elegant solutions

It is related to *induction* in mathematics, which has

- a base case (a problem instance where the solution is trivial)
- an inductive step (build solution from a simpler version of the problem)

## ❖ Example #1: factorial

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A simple example: computing factorial ( $n!$ )

- base case:  $n$  is 1  $\Rightarrow n!$  is 1
- for larger values:
  - I can't solve the whole problem directly
  - but I do know the value of  $n$
  - I could compute  $(n-1)!$  (easier than  $n!$  ?)
- multiply  $n$  by  $(n-1)!$ , giving  $n!$

E.g.

$$\begin{aligned}\text{factorial}(3) &= 3 * \text{factorial}(2) \\ &= 3 * (2 * \text{factorial}(1)) \\ &= 3 * (2 * 1) = 6\end{aligned}$$

## ❖ ... Example #1: factorial

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Expressing this as a C function:

```
int factorial(int n) {  
    if (n == 1)    // base case  
        return 1;  
    else           // recursive case  
        return n * factorial(n-1);  
}
```

compared to iterative version

```
int factorial(int n) {  
    int fac = 1;  
    for (int i = 1; i <= n; i++)  
        fac = fac * i;  
    return fac;  
}
```

## ❖ Example #2: Summing values in a list

Another simple example: summing integer values in a list

- base case: empty list  $\Rightarrow$  sum is zero
- for larger lists:
  - I can't solve the whole problem directly
  - but I do know the first value in the list
  - sum the rest of the list (smaller than whole list, easier?)
- add first value to sum-of-rest, giving sum of whole

E.g.

$$\begin{aligned}\text{sum } [1,2,3] &= 1 + \text{sum } [2,3] \\ &= 1 + (2 + \text{sum } [3]) \\ &= 1 + (2 + (3 + \text{sum } [])) \\ &= 1 + (2 + (3 + 0)) = 6\end{aligned}$$

## ❖ ... Example #2: Summing values in a list

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Expressing previous method as an (abstract) function

```
int sum(List L) {  
    if (empty(L))  
        return 0;  
    else {  
        int first, sumRest;  
        first = head(L);  
        sumRest = sum(tail(L));  
        return first + sumRest;  
    }  
}
```

## ❖ ... Example #2: Summing values in a list

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And then expressing using typical list data structure:

```
struct Node { int val; struct Node *next; };
```

```
int sum(struct Node *L) {  
    if (L == NULL)  
        return 0;  
    else {  
        int first, sumRest;  
        first = L->val;  
        sumRest = sum(L->next);  
        return first + sumRest;  
    }  
}
```

or

```
int sum(struct Node *L) {  
    if (L == NULL)  
        return 0;
```

```
    else  
        return L->val + sum(L->next);  
}
```



## ❖ How it works

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Recursion is a function calling itself

Won't the system get confused?

No, because each call to the function is a separate **instance**

- each function call creates a new mini-environment
- this holds all of the data needed by the function

The "mini-environments" are called **stack frames**

- they are created as part of the function call
- they are removed when the function **returns**



## ❖ Using Recursion

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While it is useful to know how it works ...

Sometimes it is confusing to think about stacks, etc.

When designing (or reading) recursive functions

- return to recursion basics
- identify the base case
- see how the problem can be reduced
- see how results can be built from base + recursive case

COMP2521 has many examples of recursively-defined algorithms

## ❖ Postscript

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Generally, recursive solutions are efficient (except stack space)

Sometimes, they can be very inefficient, e.g.

```
// returns n'th fibonacci number
int fibonacci(int n) {
    if (n == 1)
        return 1;
    else if (n == 2)
        return 1;
    else
        return fibonacci(n-1) + fibonacci(n-2);
}
```

Trace the recursive calls for **fibonacci(5)** to see the problem

