Coordinating Heterogenous Agents for Fast Package Delivery — The Package Handoff Problem

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Abstract

How do you get a package from an initial location S to a destination point T using a fleet of "heterogenous" carrier agents (e.g. drones, taxis). By "heterogenous" we mean the two agents can have different capabilities like different maximum speed or different fuel capacity.

In the simplest version of the category of problems, we are given as input the initial locations of n agents in \mathbb{R}^2 each capable of a maximum speed $u_i > 0$ (where u_i need not be equal to v_j for $i \neq j$). Each agent can pick up the package and move to another point to rendezvous with and hand off the package to another agent. This other agent then either proceeds to T or decides to meet with and hand off the package to another agent, until the last agent decides to head directly to T.

The objective is to get the agents to cooperate to send the package from S to T in the least possible time. We call this the $Package\ Handoff\ Problem$.

To solve this problem and its various avatars we need to

- 1. Figure out which subset $S = \{i_1, i_2, \dots i_k\}$ of the drones are used in the optimal schedule.
- 2. Find the order in which the handoffs happend between the drones used in a schedule.
- 3. Calculate the "handoff" points where drone i_m hands over the package to drone i_{m+1} , for $1 \le m \le k-1$

This report is an algorithmic study of various heuristics developed to solve different variants of Package Handoff.

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Chapter 1

Single Package Handoff

1.1 Introduction

How do you get a package from an initial location S to a destination point T using a fleet of "heterogenous" carrier agents (e.g. drones, taxis). By "heterogenous" we mean the two agents can have different capabilities like different maximum speed or different fuel capacity.

In the simplest version of the category of problems, we are given as input the initial locations of n agents in \mathbb{R}^2 each capable of a maximum speed $u_i > 0$ (where u_i need not be equal to v_j for $i \neq j$). Each agent can pick up the package and move to another point to rendezvous with and hand off the package to another agent. This other agent then either proceeds to T or decides to meet with and hand off the package to another agent 1 and so on and so forth.

The objective is to get the agents to cooperate to send the package from S to T in the least possible time. We call this the $Package\ Handoff\ Problem$.

To solve this problem and its various a vatars $^{\mathbf{2}}$ we need to

- 1. Figure out which subset $S = \{i_1, i_2, \dots i_k\}$ of the drones are used in the optimal schedule.
- 2. Find the order in which the handoffs happend between the drones used in a schedule.
- 3. Calculate the "handoff" points where drone i_m hands over the package to drone i_{m+1} 3

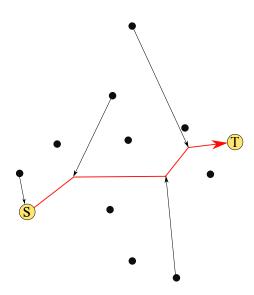


Figure 1.1: An instance of the Package Handoff problem for a single package being transported from S to T. Agents are located at the dots marked in black. The package travels along the red path. The agents all have different velocities, and in this example, assumed to have infinite battery capacity.

A real world instance of the basic Package Handoff problem, as described in the abstract, is when a ride hailing service must co-ordinate its fleet of taxis to transport a passenger from a given location in the quickest possible time to the target destination on the map. In this model, a passenger "hops rides" when two taxis meet: a taxi first gets to the passenger and takes him/her to a point where it rendezvous with another taxi, at which point the passenger swaps taxis. This process continues until the passenger hops onto a taxi that goes straight to the target.

 $^{^{1}}$ If it makes the package get to T faster

²Say when there are multiple packages to be delivered or a bound on fuel

 $^{^{3}}$ The last drone in the computed schedule, of course, flies directly to T

This package handoff process is depicted in Figure 1.1.

1.2 A note on source code

Many of the heuristics algorithms described here are implemented as literate programs [6] in Python 2.7.12 using the NuWeb tool [2] available from http://nuweb.sourceforge.net/ alongside associated theoretical and empirical analysis. All the algorithmic code goes into the file pholib.py, and any associated helper codes go into the files which are named as utils_*.py. The code for these utility files has been given in the appendices.

The pholib.py file looks like

```
"src/pholib.py" 3≡
```

```
from colorama import Fore, Style from matplotlib import rc import matplotlib as mpl import matplotlib.pyplot as plt from scipy.optimize import minimize from sklearn.cluster import KMeans import numpy as np import argparse, inspect, itertools, logging import os, time, sys import pprint as pp, randomcolor import utils_algo, utils_graphics

\( \langle Algorithms 6, \ldots \rangle \langle Run Handlers 26 \rangle \langle Plotting 33 \rangle \langle Plotting 33 \rangle \langle \langle Plotting 35 \rangle Plotting 35 \r
```

The chapters are devoted to fleshing out the chunks $\langle \texttt{Algorithms} \rangle$ and $\langle \texttt{Experiments} \rangle$. The $\langle \texttt{Run Handlers} \rangle$ and $\langle \texttt{Plotting} \rangle$ chunks is mainly to deal with interactive matplotlib input, and as such are boring and banished to the Appendix \odot .

All source code files are tangled to the **src** directory. The point of entry for the code are **main*.py** which are implemented separately in the **src** directory, since their contents can change based on what library code is being called for during development and testing. Since these files are very short and the mechanics clear, they are implemented as standalone files (i.e. not inside this document) but directly in the **src** folder itself. To run the code in interactive mode run the code as **python src/main_interactive.py** on a Unix / Windows terminal in the root folder of the project ⁴.

For a short overview of previous work on this problem see Appendix.

The README file containing instructions for running the source code and experiments is listed in the appendix (also available on the Github repository)

Each of the following sections correspond to a fixed variant of the package handoff problem and describe algorithms for that specific variant. Enough talk! Onto algorithms!

1.3 Problem: Unlimited Fuel, Different Drone Speeds

Much of the machinery developed in solving this basic basic basic question will be generalized and extended to other variants of the package handoff problem.

 $^{^4}$ This code has been tested on an 64 bit machine running Linux Mint 18.3 (Sylvia) running the Linux Kernel version 4.10.0-38-generic with an Intel(R) Core(TM) i7 CPU 960 @ 3.20GHz CPU

We repeat the problem definition and fix some notation that will be used for the remainder of the section

We are given as input the initial locations P_i of n agents in \mathbb{R}^2 each capable of a maximum speed $u_i > 0$ (where u_i need not be equal to u_j for $i \neq j$). Each agent can pick up the package and move to another point to rendezvous with and hand off the package to another agent. This other agent then either proceeds to T or decides to meet with and hand off the package to another agent and so on. The objective is to get the agents to cooperate to send the package from S to T in the least possible time.

We represent the handoff points as follows $H_{i_1} \dots H_{i_k}$ for $0 \le i_0, \dots i_k \le n$ stand for points where the drones with labels $i_0, \dots i_k$ hand the package off in that order. More precisely H_{i_j} is the point where drone i_{j-1} hands off the package to drone i_j for $1 \le j \le k$.

A solution to the package handoff problem is completely specified by computing the handoff points and the drone ids involved in the exchange at each handoff point.

The optimal schedule is denoted OPT. It is easy to see the statements of the following structural lemma always hold for OPT.

Lemma 1. In OPT

- A. A package is always transferred to a faster drone at a handoff point.
- **B.** A drone handles a package at most once i.e. if a drone hands off the package, it will never be involved in a handling that package again.
- C. All drones involved in the handoff start moving simultaneously at time t=0
- D. No two drones wait at a rendezvous point before rendezvous happens. ^a
- **E.** The path of the package is a radially monotone piecewise straight polygonal curve with respect to the direction ST no matter what the initial positions P_i or speeds u_i of the drones.
- **F.** $\frac{|ST|}{v_{max}}$ is a (trivial) lower bound for OPT, where v_{max} denotes the speed of the fastest drone.

^awaiting can happen in other problem variants say when there is limited fuel or only a finite set of allowed rendezvous points

Proof. TODO!

1.3.1 Handoff in a fixed order

If we know the drones involved in the handoff along with the order of handoff then we can compute the handoff points — and hence the path of the package — exactly via convex optimization as outlined in Lemma 2. This fact will be exploited in many heuristics: such methods will compute a subset of drones involved in the handoff (alongwith the handoff order) followed by a call to the convex program.

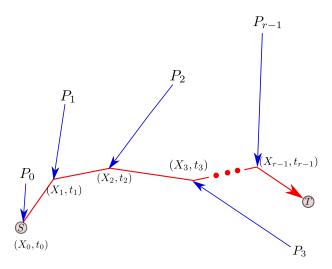


Figure 1.2: The path of the package is shown in red. The drones involved in the handoff are labelled P_i in the prespecified handoff order.

Lemma 2. Given as input are drones with initial positions $P_i \in \mathbb{R}^2$, with speeds $u_i > 0$ for $1 \le i \le r - 1$, the initial position S and final destination T for the package. The drones are expected to transport the package by handing of the package in the order $1, 2, \ldots, r$. Let t_i denote the departure time on a global clock from the i'th handoff point X_i .

Then the minimum time and handoff points for transporting the package and the handoff points can be calculated by the following convex program

$$\min_{t_i, X_i} \quad t_{r-1} + \frac{||T - X_{r-1}||}{u_{r-1}}$$

subject to the constraints

$$X_0 = S$$

$$t_i \ge \frac{||P_i - X_i||}{u_i} \qquad 0 \le i \le r - 1$$

$$t_i + \frac{||X_{i+1} - X_i||}{u_i} \le t_{i+1} \qquad 0 \le i \le r - 2$$

^aSee Figure 1.3.1 for an illustration of the notation used in this lemma

The following function is just an implentation of the convex program just described. Here **drone_info** is a list of tuples, where each tuple consits of the initial position and speed of the drone. The order of the drones is assumed to be that in which the list of drones is provided. **source** and **target** are just coordinate locations of S and T respectively. We use the CVXPY [3] library as a black-box convex optimization solver.

```
\langle Algorithms 6 \rangle \equiv
     def algo_pho_exact_given_order_of_drones ( drone_info, source, target ):
         import cvxpy as cp
         source = np.asarray(source)
         target = np.asarray(target)
         r = len(drone_info)
         source = np.asarray(source)
         target = np.asarray(target)
         # Variables for rendezvous points of drone with package
         X, t = [], []
         for i in range(r):
            X.append(cp.Variable(2)) # vector variable
            t.append(cp.Variable( )) # scalar variable
         # Constraints
         constraints_S = [ X[0] == source ]
         constraints_I = []
         for i in range(r):
           constraints_I.append(0.0 <= t[i])</pre>
           constraints_I.append(t[i] >= cp.norm(np.asarray(drone_info[i][0])-X[i])/drone_info[i][1])
         constraints_L = []
         for i in range(r-1):
           constraints\_L.append(t[i] + cp.norm(X[i+1] - X[i])/drone\_info[i][1] <= t[i+1])
         objective = cp.Minimize(t[r-1]+cp.norm(target-X[r-1])/drone_info[r-1][1])
         prob = cp.Problem(objective, constraints_S + constraints_I + constraints_L)
         print Fore.CYAN
         prob.solve(solver=cp.SCS,verbose=True)
         print Style.RESET_ALL
         package_trail = [ np.asarray(X[i].value) for i in range(r) ] + [ target ]
         return package_trail
Fragment defined by 6, 7, 10a, 12a, 14a.
Fragment referenced in 3.
```

We next describe a heuristic that use Continuous Dijkstra [11] type approach in computing approximate solutions to OPT.

1.3.2 Algorithm: One Dimensional Greedy Wavefront

In this heuristic we first constrain the package to travel along the line \widetilde{ST} , then compute the subset of the drones involved in the schedule, and finally pass of the list of drones involved to the convex program given in Lemma 2 to calculate the rendezvous points.

Here is a sketch of the implementation of algo_odw. Use Figure 1.3.2 as a reference while reading the description below.

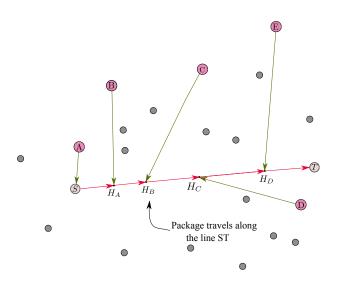


Figure 1.3: first figure

Figure 1.4: The package travels along the straight line \vec{ST} . The point where drone A hands off the package to drone B depicted as H_B , and similarly for other drones. Drones involved in the handoff are marked in pink. Those not involved are marked in gray. Two drones may have different speed.

Defines: algo_odw 26.

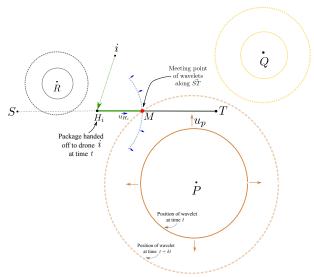


Figure 1.5: second figure

Figure 1.6: Intersection of two expanding wavelets along H_iT . The figure shows an overlay of snapshots of two times; one at time t (when the package has just been handed off to drone i at H_i) and another at time $t + \delta t$ (when a wavelet corresponding a drone faster than drone i meets it at M)

```
\langle Algorithms 7 \rangle \equiv
      def algo_odw(drone_info, source, target, plot_tour_p = False):
          from scipy.optimize import minimize
          source
                       = np.asarray(source)
          target
                       = np.asarray(target)
          sthat
                       = (target-source)/np.linalg.norm(target-source)
          numdrones = len(drone_info)
          clock_time = 0.0 # time on the global clock
          ⟨ Find the drone which can get to the source the quickest 8a⟩
          package_reached_p
                                     = False
          while not(package_reached_p):
                 \langle Find \ a \ faster \ wavelet \ that \ meets \ up \ with \ the \ package \ wavelet \ along \ line \ \vec{ST} \ at \ the \ earliest \ 8b \rangle
                (Check if package wavelet reaches target before meeting wavelet computed above. Update states accordingly 11a)
           ⟨ Run the convex optimization solver to retrieve the exact tour package_trail_cvx for given drone order 11c⟩
           ⟨ Plot tour if plot_tour_p == True 12b⟩
          return used_drones, package_trail_straight, mspan_straight, package_trail_cvx, mspan_cvx
Fragment defined by 6, 7, 10a, 12a, 14a.
Fragment referenced in 3.
```

At the start of the algorithm, the package is stationary. We find the drone which can reach it at the earliest, and update the time on the global clock. It is easy to see that this drone must be involved in the optimal handoff schedule.

```
\langle Find the drone which can get to the source the quickest 8a \rangle \equiv
     tmin = np.inf
     imin = None
     for idx in range(numdrones):
               initdroneposn = drone_info[idx][0]
               dronespeed
                              = drone_info[idx][1]
               tmin_idx = time_of_travel(initdroneposn, source, dronespeed)
               if tmin_idx < tmin:</pre>
                   tmin = tmin_idx
                   imin = idx
     clock_time = tmin
     current_package_handler_idx = imin
     current_package_position
                                   = source
     drone_pool = range(numdrones)
     drone_pool.remove(imin)
     used_drones = [imin]
     package_trail_straight = [current_package_position]
Fragment referenced in 7.
```

Uses: time_of_travel 12a.

 $Uses: \ \mbox{\tt get_interception_time} \ \ \mbox{\tt 10a}.$

In the optimal handoff order, as we have already noted, the handoff happens from a slower to a faster robot. In the next chunk, we calculate the wavelet that meets the package wavelet at the earliest, subject to the constraint that the wavelet should be faster than the package wavelet.

 \langle Find a faster wavelet that meets up with the package wavelet along line \vec{ST} at the earliest $8b\,\rangle \equiv$

```
tI_min
                = np.inf
     idx_tI_min = None
     for idx in drone_pool:
           us = drone_info[current_package_handler_idx][1]
           up = drone_info[idx][1]
           if up <= us: # slower drones are useless, so skip rest of the iteration
               continue
           else:
             s = current_package_position
             p = np.asarray(drone_info[idx][0])
             tI = get_interception_time(s, us, p, up, target, clock_time)
             if tI < tI_min:</pre>
                tI_min
                         = tI
                idx_tI_min = idx
Fragment referenced in 7.
```

We now implement get_interception_time, the time on the global clock when two wavelets meet. More precisely, the function computes the time on the global clock at which a wavelet that started expanding from $P = (\alpha, \beta)$ at time 0, meets a wavelet that started expanding from S at time S at time S along the half-line S. We make a change of coordinages such that S at time S at

is horizontal and pointing to the right, and S = (0,0) as in Figure 1.3.2.

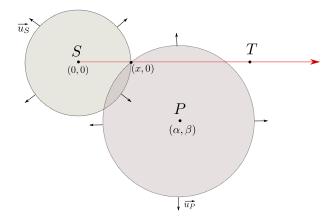


Figure 1.7: Reference figure for the function $get_interception_time$. Without loss of generality (by changing the coordinate system) we can assume the line \vec{ST} to be horizontal.

$$\frac{x-0}{u_S} = \frac{\sqrt{(x-\alpha)^2 + \beta^2}}{u_P} - t_0$$
$$\left(\frac{x}{u_S} + t_0\right)^2 = \frac{(x-\alpha)^2 + \beta^2}{u_P^2}$$

Rearranging the terms, we get a quadratic equation in x that we can solve solve using standard non-linear solvers 5 .

$$x^{2} \left(\frac{1}{u_{S}^{2}} - \frac{1}{u_{P}^{2}} \right) + x \left(\frac{2t_{0}}{u_{S}} + \frac{2\alpha}{u_{P}^{2}} \right) + \left(t_{0}^{2} - \frac{\alpha^{2}}{u_{P}^{2}} - \frac{\beta^{2}}{u_{P}^{2}} \right) = 0$$

$$(1.1)$$

Once we obtain x getting the interception time t_I (i.e. the time on the global clock when the two wavelets meet) as $t_I = \frac{x}{u_S} + t_0$.

⁵Using the exact formula for solving the quadratic equation is unstable in computer arithmetic as is well known from numerical analysis. To avoid such tricky issues, it is best to use the standard roots solver for polynomials available from numpy https://docs.scipy.org/doc/numpy/reference/generated/numpy.roots.html

```
\langle Algorithms 10a \rangle \equiv
      def get_interception_time(s, us, p, up, t, t0) :
          \langle Change\ coordinates\ to\ make\ s=(0,0)\ and\ t\ to\ lie\ along\ X-axis\ as\ in\ Figure\ 1.3.2\ 10b \rangle
          # Solve quadratic equation as documented in main text
          qroots = np.roots([ (1.0/us**2 - 1.0/up**2),
                                2*t0/us + 2*alpha/up**2,
                                t0**2 - alpha**2/up**2 - beta**2/up**2])
          # The quadratic should always have a root.
          qroots = np.real(qroots) # in case the imaginary parts
                                     # of the roots are really small
          qroots.sort()
          x = None
          for root in groots:
              if root > 0.0:
                  x = root
                 break
          assert abs(x/us+t0 - np.sqrt((x-alpha)**2 + beta**2)/up) <= 1e-6,
                  "Quadratic not solved perfectly"
          tI = x/us + t0
          return tI
Fragment defined by 6, 7, 10a, 12a, 14a.
Fragment referenced in 3.
Defines: get_interception_time 8b, 17a.
    Massage the input into such that \vec{ST} lies along the positive X-axis as in Figure 1.3.2.
\langle Change\ coordinates\ to\ make\ s=(0,0)\ and\ t\ to\ lie\ along\ X-axis\ as\ in\ Figure\ 1.3.2\ 10b \rangle \equiv
      t_m = t - s # the _m subscript stands for modify
      t_m = t_m / np.linalg.norm(t_m) # normalize to unit
      # For rotating a vector clockwise by theta,
      # to get the vector t_m into alignment with (1,0)
      costh = t_m[0]/np.sqrt(t_m[0]**2 + t_m[1]**2)
      sinth = t_m[1]/np.sqrt(t_m[0]**2 + t_m[1]**2)
     rotmat = np.asarray([[costh, sinth],
                            [-sinth, costh]])
      assert np.linalg.norm((rotmat.dot(t_m) - np.asarray([1,0]))) \leq 1e-6,\
             "Rotation matrix did not work properly. t_m should get rotated\
              onto [1,0] after this transformation"
     p_shift = p - s
               = rotmat.dot(p_shift)
     p_rot
      [alpha, beta] = p_rot
Fragment referenced in 10a.
```

If the package reaches the target before it meets any other wavelet along the line \vec{ST} , then there is no point in handing off the package to some other drone. Just terminate the handoff! Otherwise handoff the package and update variables

```
accordingly.
```

```
\langle Check if package wavelet reaches target before meeting wavelet computed above. Update states accordingly 11a\rangle \equiv
      time_to_target_without_handoff = np.linalg.norm((target-current_package_position))/ \
                                           drone_info[current_package_handler_idx][1]
      if time_to_target_without_handoff < tI_min :</pre>
           package_reached_p = True
           package_trail_straight.append(target)
      else:
           ⟨ Update package information (current speed, position etc.) and drone information (available and used drones) 11b⟩
Fragment referenced in 7.
\langle Update\ package\ information\ (current\ speed,\ position\ etc.)\ and\ drone\ information\ (available\ and\ used\ drones)\ 11b 
angle \equiv
      package_handler_speed
                                  = drone_info[current_package_handler_idx][1]
      current_package_position = current_package_position + \
                                     package_handler_speed * (tI_min - clock_time) * sthat
     package_trail_straight.append(current_package_position)
                                     = tI_min
      clock time
      current_package_handler_idx = idx_tI_min
     drone_pool.remove(idx_tI_min)
     used_drones.append(idx_tI_min)
Fragment referenced in 11a.
```

Now that we have a list of drones involved in the handoff, (along with an approximate trail for the package) we use the convex optimization solver to extract the exact tour for the given set of drones. We plot both tours for a visual comparison, if plot_tour_p is set to True.

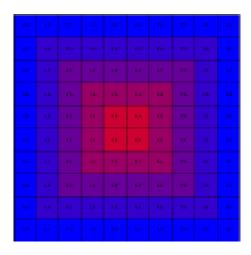
The next chunk implements a function that computes the makespan of the delivery process, i.e. the time it takes for the package to get from the source to the destination, given the points on the trajectory of the package and the drones involved alongwith the handoff order. The variable **drone_info** is list of tuples, where the **i**th tuple gives the initial position and speed of the drone in the zeroth and first position respectively. The function **time_of_travel** is simply a function used to compute the time it takes for an agent with uniform speed to travel between a given **source** and **target**. Both **source** and **target** are numpy arrays of size two (if not, they are converted into numpy arrays at the start of the function).

```
\langle Algorithms 12a \rangle \equiv
     def time_of_travel(start, stop, speed):
           start = np.asarray(start)
           stop = np.asarray(stop)
           return np.linalg.norm(stop-start)/speed
     def extract_coordinates(points):
          xs, ys = [], []
          for pt in points:
              xs.append(pt[0])
              ys.append(pt[1])
          return np.asarray(xs), np.asarray(ys)
     def makespan(drone_info, used_drones, package_trail):
          assert len(package_trail) == len(used_drones)+1, ""
          makespan = 0.0
          counter = 0
          for idx in used_drones:
               dronespeed
                              = drone_info[idx][1]
               makespan += time_of_travel(package_trail[counter],\
                                             package_trail[counter+1],
                                             dronespeed)
               counter += 1
          return makespan
Fragment defined by 6, 7, 10a, 12a, 14a.
Fragment referenced in 3.
Defines: extract_coordinates 33, makespan, Never used, time_of_travel 8a.
\langle \mathit{Plot\ tour\ if}\ \mathtt{plot\_tour\_p} == \mathsf{True}\ 12b\, \rangle \equiv
      if plot_tour_p:
           fig0, ax0 = plt.subplots()
           plot_tour(fig0, ax0, "ODW: Straight Line", source, target, \
                      drone_info, used_drones, package_trail_straight)
           fig1, ax1 = plt.subplots()
           plot_tour(fig1, ax1, "ODW: Straight Line, Post Convex Optimization", source, target, \
                      drone_info, used_drones, package_trail_cvx)
           plt.show()
Fragment referenced in 7.
```

Chapter 2

Multiple Package Handoff

What's more fun than delivering a single package? Delivering multiple packages! The moment, we generalize from a single to multiple packages the problem becomes enormously more interesting along with several possible generalizations to the statement of the problem. Each section is dedicated to one such generalization along with a description of algorithms and heuristics to solve them.



2.1 Problem: Drone Assignment To Packages

We are given as input the initial locations P_i of n agents in \mathbb{R}^2 each capable of a maximum speed $u_i > 0$. Also given are $k \leq n$ source target pairs S_j, T_j where S_j for $j \leq k$ denotes the beginning position of a package that has to get to target point T_j . Each agent is allowed to be involved in the transport of at most one package. More than one agent is allowed to be assigned to a package. Agents assigned to a package can coordinate to rendezvous and relay the package from its source to its destination (exactly as in the single package handoff case)

The objective is to perform an assignment of agents to the packages so that the time taken to deliver the last package to its target is minimized, i.e. the makespan of the deliveries of the packages from sources to their destinations is minimized. See for an example instance (alongside tours computed for the packages and drones by the "match and move" algorithm to be described next.)

2.1.1 Algorithm: Match-and-Move

The approach taken here is again based on Continuous Dijkstra. At the start of the algorithm, we assume each of the packages are constrained to travel along the the straight line segment joining its source and target. We then imagine wavelets expanding at speed u_i from each of the drone positions and perform an incremental assignment of the drones to the packages as the wavelet expansion proceeds. The assignment is done via bottleneck matching in an appropriately constructed bipartite graph. The matching algorithm is run everytime a certain "event" is detected.

Once the final assignment of drones to packages has been performed, we run the convex optimization solver desribed in subsection 1.3.1 to get the exact trajectory for each of the packages for the given assignment.

The rest of this section is devoted to making the above description more precise. First we give an outline of algo_matchmove that will be fleshed out in subsequent subsections.

```
\langle Algorithms 14a \rangle \equiv
      import networkx as nx
      def algo_matchmove(drone_info, sources, targets, plot_tour_p = False):
            ⟨ Sanity checks on input for algo_matchmove 14b⟩
            ⟨ Basic setup 16⟩
           while not all(package_delivered_p):
                  (Construct bipartite graph G on drone wavelets and package wavelets 17a)
                  ⟨ Get a bottleneck matching on G 17b⟩
                  (Expand drone wavelets till an event of either Type I or Type II is detected 18a)
                 ⟨ Update drone pool and package states 18b⟩
           ⟨ Run convex optimization solver to get improved tours for drones and packages 18c⟩
            ⟨ Plot movement of packages and drones if plot_tour_p == True 18d⟩
           #return pass pass pass pass
Fragment defined by 6, 7, 10a, 12a, 14a.
Fragment referenced in 3.
Defines: algo_matchmove 26.
```

2.1.1.1 Sanity Checks, and Basic Setup

To start things off, we enforce the constraint that the number of drones should be greater than the number of packages. Also the number of packages is equal to the number of sources which in turn is equal to the number of targets. These conditions are encoded as sanity checks in the code chunk below that is incorporated at the beginning of the function

```
⟨ Sanity checks on input for algo_matchmove 14b⟩ ≡

assert len(drone_info) >= len(sources),\
    "Num drones should be >= the num source-target pairs"

assert len(sources) == len(targets),\
    "Num sources should be == Num targets"
    ♦
Fragment referenced in 14a.
```

The lists sources, targets, drone_initposns and drone_speeds are all *constant* throughout the execution of the algorithm. They denote the list of sources, targets, initial position and speeds of the drones respectively. For package i, its source and destination are respectively source[i] and target[i].

Then, we create some lists that keep track of various states in the main while loop of the algorithm.

- A. package_delivered_p is a Boolean list where the ith element keeps track whether the package with id i has been delivered to its destination. The main while loop stops when all flags are set to True.
- **B. drone_pool** is a list of ids of drones that can be considered for the bottleneck matching process for the next iteration of the while loop.
- C. drone_wavelet_info is a list of list of dictionaries. Each outer list corresponds to a list of wavelets corresponding to each drone. Every wavelet is represented by a dictionary, with three attributes
 - 1. wavelet_center The cartesian coordinates of the center H of the wavelet about which the wavelet expands.
 - 2. clock_time The time on the global clock at which the wavelet started expanding around the wavelet center H.
 - **3.** matched_package_ids A list of the ids of the packages that were assigned to the wavelet during its expansion around center H. Each such list is initialized as the empty list.

See Figure 2.1.1.1 to see a sequence of wavelets corresponding to a single drone.

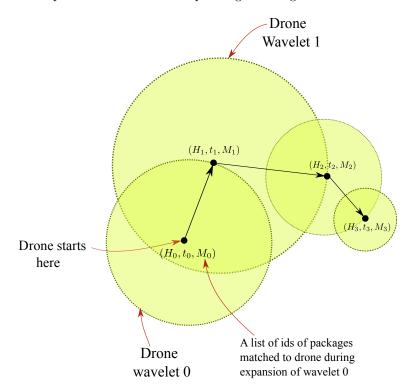


Figure 2.1: A sequence of wavelets corresponding to a drone which is stored as a list of dictionaries (each dictionary corresponding to a wavelet) [{'wavelet_center': H_i , 'clock_time': t_i , 'matched_package_ids': M_i }]. Each wavelet expands around H_i at speed u_i starting at time t_i on the global clock. Every drone has a sequence of wavelets. The H_i also represent handoff points where the drone either hands a package off to another drone or is handed a package by another drone. Note that the usage of the letter H to represent the center of a drone wavelet is suggestive of it also being a point where a package is handed to the drone or a package is handed off to another drone.

D. package_trail_info is a list of list of dictionaries, similar to **drone_wavelet_info**. The **i**th element of the outer list corresponds to the trail of a package. In each trail (itself a list), we record the positions of a package as well as the time on the clock t at which that position was recorded. We also note the id of the last drone known to have handled the package at time $t - \varepsilon$ for ε being an arbitrarily small positive number.

Two helper functions <code>get_current_position_of_package</code> and <code>get_current_speed_of_package</code> to extract information conveniently from <code>package_trail_info</code>. Again like <code>get_last_wavelet_of_drone</code> the result returned depends on the state of <code>package_trail_info</code> at the time the function is called.

```
\langle Basic\ setup\ 16 \rangle \equiv
    sources
                      = [np.asarray(source) for source in sources]
    targets
                      = [np.asarray(target) for target in targets]
    drone_initposns
                      = [ np.asarray(initposn) for (initposn, _) in drone_info ]
                                                      speed) in drone_info ]
    drone_speeds
                      = [ speed
                                            for (_,
    numpackages
                      = len(sources)
    numdrones
                      = len(drone_info)
    package_delivered_p = [ False for _ in range(numpackages) ]
    drone_pool
                    = range(numdrones)
    global_clock_time = 0.0
    #.....
    drone_wavelets_info = [ [{'wavelet_center' : posn,
                            'clock_time'
                                                : 0.0,
                           'matched_package_ids' : []}]
                          for posn in drone_initposns ]
    def get_last_wavelet_of_drone(i):
            return drone_wavelets_info[i][-1]
    #.....
    package_trail_info = [ [{'current_position' : source,
                            'clock_time'
                                             : 0.0,
                            'current_handler_id' : None }]
                          for source in sources ]
    def get_current_position_of_package(i):
            return package_trail_info[i][-1]['current_position']
    def get_current_speed_of_package(i):
            current_handler_id = package_trail_info[i][-1]['current_handler_id']
            if current_handler_id is None:
                return 0.0
            else:
                 return drone_speeds[current_handler_id]
```

Fragment referenced in 14a.

Defines: get_current_position_of_package, Never used, get_current_speed_of_package 17a, get_last_wavelet_of_drone, Never used.

The <code>get_wavelet_of_ith_drone</code> function defined at the end of the above chunk is a convenience function for returning the center of the last wavelet and the associated clock time recorded at that center for an arbitrary drone. **Note:** The result of calling this function depends on the time at which it is called because <code>drone_wavelet_info</code> list is mutated during the main while loop

2.1.1.2 Main Loop

```
\langle Construct \ bipartite \ graph \ G \ on \ drone \ wavelets \ and \ package \ wavelets \ 17a \rangle \equiv
      # Create nodes of the graph
     G = nx.Graph()
     G.add_nodes_from(['drone_' +str(didx) for didx in range(numdrones)])
     G.add_nodes_from(['package_'+str(pidx) for pidx in range(numpackages)])
      # Add edges between nodes
      for didx in range(numdrones):
          for pidx in range(numpackages):
                target = targets[pidx]
                       = get_current_position_of_package(pidx)
                       = get_current_speed_of_package(pidx)
                upkg
                       = get_last_wavelet_of_drone(didx)['wavelet_center']
                wav
                dro
                       = wav['wavelet_center']
                udro
                       = drone_speeds[didx]
                 if upkg < 1e-7: # zero testing for upkg
                    # Add edge to G
                    pass
                elif udro > upkg: # It is critical that the zero testing for upkg
                                     # has been done for time_target_to_solo to be
                                     # computed safely
                    time_to_target_solo = np.linalg.norm((target-pkg))/ upkg
                                          = get_interception_time(pkg, upkg, dro, udro,
                                                                     target, wav['clock_time'])
                    if tI < time_to_target_solo:</pre>
                           # Add edge to G
                           pass
                 else:
                     continue
      sys.exit()
     \Diamond
Fragment referenced in 14a.
Uses: get_current_speed_of_package 16, get_interception_time 10a.
\langle Get \ a \ bottleneck \ matching \ on \ G \ 17b \rangle \equiv
     pass
Fragment referenced in 14a.
```

There are exactly two types of (mutually exclusive) events in this algorithm.

Type I: A package gets delivered to its assigned destination.

Type II: A wavelet corresponding to a drone (not handling a package) meets up with a drone currently handling a package.

```
⟨ Expand drone wavelets till an event of either Type I or Type II is detected 18a⟩ ≡

pass

⟨ Update drone pool and package states 18b⟩ ≡

pass

⟨ Run convex optimization solver to get improved tours for drones and packages 18c⟩ ≡

pass

⟨ Plot movement of packages and drones if plot_tour_p == True 18d⟩ ≡

pass

⟨ Plot movement of packages and drones if plot_tour_p == True 18d⟩ ≡

pass

⟨ Pragment referenced in 14a.
```

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Appendices

Appendix A

History and Previous Work

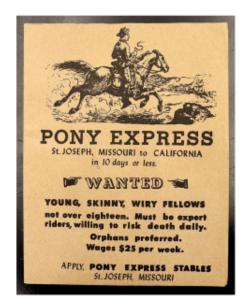




Figure A.1: A job application poster and a relay route used for the Pony Express. Images taken from [16] and [14] respectively.

A system of using relays for delivering packages is not a particularly new idea. A famous (and shortlived!) example of such a relay system was the Pony Express company which was used a system of a relay of horse riders to transport mail from St. Joseph, Missouri to Sacramento, California.

To quote the Wikipedia article

"Operated by Central Overland California and Pike's Peak Express Company, the Pony Express was a great financial investment to the U.S. During its 18 months of operation, it reduced the time for messages to travel between the Atlantic and Pacific coasts to about 10 days. It became the West's most direct means of east-west communication before the transcontinental telegraph was established (October 24, 1861), and was vital for tying the new U.S. state of California with the rest of the United States. The Pony Express demonstrated that a unified transcontinental system of communications could be established and operated year-round."

While the invention of the telegraph might have run the Pony Express out of business, the idea of using relay agents such as drones — instead of horses! — to transfer packages can have applications today for sending physical goods (which of course can't be telegraphed! ②) such as life-saving medicinces in under-developed countries or in disaster relief areas. ZipLine[15] [5] and Matternet [10] are just two of the companies which are involved in building networks of drones for precisely such missions.



Figure A.2: Dabbawallas exchanging lunchboxes (dabbas) at a relay point. Image from [9]

Another relay system for package deliveries (135 years old and still functioning!) is that of the dabbawallas 1 used for transporting lunch boxes from homes and restaurants to

people at work in Mumbai, India. To quote from [12]

Four thousand five hundred semi-literate dabbawalas collect and deliver 175,000 packages within hours. What should we learn from this unique, simple and highly efficient 120-year-old logistics system? [...] After the customer leaves for work, her lunch is packed into a tiffin provided by the dabbawala. A color-coded notation on the handle identifies its owner and destination. Once the dabbawala has picked up the tiffin, he moves fast using a combination of bicycles, trains and his two feet.

A BBC crew filming dabbawalas in action was amazed at their speed. "Following our dabbawala wasn't easy, our film crew quickly lost him in the congestion of the train station. At Victoria Terminus we found other fast moving dabbawalas, but not our subject... and at Mr Bhapat's ayurvedic pharmacy, the lunch had arrived long before the film crew," the documentary noted wryly. So, how do they work so efficiently?

The entire system depends on teamwork and meticulous timing. Tiffins are collected from homes between 7.00 am and 9.00 am, and taken to the nearest railway station. At various intermediary stations, they are hauled onto platforms and sorted out for area-wise distribution, so that a single tiffin could change hands three to four times in the course of its daily journey.

At Mumbai's downtown stations, the last link in the chain, a final relay of dabbawalas fan out to the tiffins' destined bellies. Lunch hour over, the whole process moves into reverse and the tiffins return to suburban homes by 6.00 pm.

See https://youtu.be/dX-0el2wuEU for a short video on the dabbawallas.

¹literally: lunchbox carriers

Appendix B

README.md

This README file can be read more clearly alongside its appropriate formatting at https://github.com/gtelang/PackageHandoff_Python/tree/master/packagehandoff_lit

"README.md" $22 \equiv$

To run this code, you will need a distribution of Python 2.7.12 along with the following libraries

- networkx
- matplotlib
- numpy
- scipy
- cgal-bindings
- cvxpy

All source code is contained in the .web file. If you modify the file, the resulting code and corresponding description file can be weaved and tangled with the script 'weave-tangle.sh'. You will need the

- [pdflatex](https://linux.die.net/man/1/pdflatex)
- [nuweb] (http://nuweb.sourceforge.net/)
- [asymptote](http://asymptote.sourceforge.net/)

executables to be somewhere on your system's path.

All source code is in the 'src' directory. The asy2d, asy3d and docs folders can be neglected since they contain images and documents referenced in the 'packagehandoff.pdf' file. From the point of view of _running_ the code they can be ignored.

0

Appendix C

utils_graphics.py

This file contains useful functions for visualization and plotting functions described in the previous chapters.

"src/utils_graphics.py" $23 \equiv$

```
from matplotlib import rc
from colorama import Fore
from colorama import Style
from scipy.optimize import minimize
from sklearn.cluster import KMeans
import argparse
import itertools
import math
import matplotlib as mpl
import matplotlib.pyplot as plt
import numpy as np
import os
import pprint as pp
import randomcolor
import sys
import time
xlim, ylim = [0,1], [0,1]
# Borrowed from https://stackoverflow.com/a/9701141
import numpy as np
import colorsys
def get_colors(num_colors, lightness=0.2):
    colors=[]
    for i in np.arange(60., 360., 300. / num_colors):
                  = i/360.0
        saturation = 0.95
        colors.append(colorsys.hls_to_rgb(hue, lightness, saturation))
    return colors
\Diamond
```

Appendix D

utils_algo.py

This file contains useful functions for writing algorithms described in the previous chapters.

```
"src/utils_algo.py" 24 \equiv
```

```
import numpy as np
import random
from colorama import Fore
from colorama import Style
def vector_chain_from_point_list(pts):
    vec_chain = []
    for pair in zip(pts, pts[1:]):
        tail= np.array (pair[0])
        head= np.array (pair[1])
        vec_chain.append(head-tail)
    return vec_chain
def length_polygonal_chain(pts):
    vec_chain = vector_chain_from_point_list(pts)
    acc = 0
    for vec in vec_chain:
       acc = acc + np.linalg.norm(vec)
    return acc
def pointify_vector (x):
    if len(x) \% 2 == 0:
        pts = []
        for i in range(len(x))[::2]:
            pts.append( [x[i],x[i+1]] )
        return pts
    else :
        sys.exit('List of items does not have an even length to be able to be pointifyed')
def flatten_list_of_lists(l):
      return [item for sublist in 1 for item in sublist]
def print_list(xs):
    for x in xs:
       print x
def partial_sums( xs ):
    psum = 0
    acc = []
    for x in xs:
```

```
psum = psum + x
        acc.append( psum )
    return acc
def are_site_orderings_equal(sites1, sites2):
    for (x1,y1), (x2,y2) in zip(sites1, sites2):
        if (x1-x2)**2 + (y1-y2)**2 > 1e-8:
            return False
    return True
def bunch_of_non_uniform_random_points(numpts):
    cluster_size = int(np.sqrt(numpts))
   numcenters = cluster_size
   import scipy
    import random
    centers = scipy.rand(numcenters,2).tolist()
    scale, points = 4.0, []
    for c in centers:
        cx, cy = c[0], c[1]
        # For current center $c$ of this loop, generate \verb|cluster_size| points uniformly in a square centered a
        sq_size
                     = min(cx, 1-cx, cy, 1-cy)
                     = np.random.uniform(low=cx-sq_size/scale, high=cx+sq_size/scale, size=(cluster_size,))
        loc_pts_x
                     = np.random.uniform(low=cy-sq_size/scale, high=cy+sq_size/scale, size=(cluster_size,))
        loc_pts_y
        points.extend(zip(loc_pts_x, loc_pts_y))
    # Whatever number of points are left to be generated, generate them uniformly inside the unit-square
    num_remaining_pts = numpts - cluster_size * numcenters
    remaining_pts = scipy.rand(num_remaining_pts, 2).tolist()
    points.extend(remaining_pts)
   return points
def write_to_yaml_file(data, dir_name, file_name):
   import yaml
   with open(dir_name + '/' + file_name, 'w') as outfile:
    yaml.dump( data, outfile, default_flow_style = False)
0
```

Appendix E

 $\langle Run \ Handlers \ 26 \rangle \equiv$

def clearAxPolygonPatches(ax):

Implementation of (Run Handlers)

```
This chunk contains code required for the interactive input of sites and agents onto the canvas.
```

```
# Set up logging information relevant to this module
logger=logging.getLogger(__name__)
logging.basicConfig(level=logging.DEBUG)
def debug(msg):
   frame,filename,line_number,function_name,lines,index=inspect.getouterframes(
       inspect.currentframe())[1]
   line=lines[0]
   indentation_level=line.find(line.lstrip())
   logger.debug('{i} [{m}]'.format(
        i='.'*indentation_level, m=msg))
def info(msg):
   frame,filename,line_number,function_name,lines,index=inspect.getouterframes(
        inspect.currentframe())[1]
   line=lines[0]
   indentation_level=line.find(line.lstrip())
   logger.info('{i} [{m}]'.format(
        i='.'*indentation_level, m=msg))
xlim, ylim = [0,1], [0,1]
def applyAxCorrection(ax):
      ax.set_xlim([xlim[0], xlim[1]])
      ax.set_ylim([ylim[0], ylim[1]])
      ax.set_aspect(1.0)
def clearPatches(ax):
   # Get indices cooresponding to the polygon patches
   for index , patch in zip(range(len(ax.patches)), ax.patches):
        if isinstance(patch, mpl.patches.Polygon) == True:
            patch.remove()
   ax.lines[:]=[]
   applyAxCorrection(ax)
```

```
# Get indices cooresponding to the polygon patches
   for index , patch in zip(range(len(ax.patches)), ax.patches):
        if isinstance(patch, mpl.patches.Polygon) == True:
           patch.remove()
   ax.lines[:]=[]
   applyAxCorrection(ax)
class Single_PHO_Input:
   def __init__(self, drone_info = [] , source = None, target=None):
          self.drone_info = drone_info
          self.source
                         = source
          self.target
                          = target
   def get_drone_pis (self):
          return [self.drone_info[idx][0] for idx in range(len(self.drone_info)) ]
   def get_drone_uis (self):
          return [self.drone_info[idx][1] for idx in range(len(self.drone_info)) ]
   def get_tour(self, algo, plot_tour_p=False):
          return algo( self.drone_info,
                       self.source,
                        self.target,
                       plot_tour_p
   # Methods for \verb|ReverseHorseflyInput|
   def clearAllStates (self):
         self.drone_info = []
         self.source = None
         self.target = None
def single_pho_run_handler():
   import random
   def wrapperEnterRunPoints(fig, ax, run):
     def _enterPoints(event):
        if event.name
                          == 'button_press_event'
                                                            and \
           (event.button == 1 or event.button == 3)
                                                            and \
           event.dblclick == True and event.xdata != None and event.ydata != None:
            if event.button == 1:
                 # Insert blue circle representing the initial position of a drone
                 print Fore.GREEN
                 newPoint = (event.xdata, event.ydata)
                        = np.random.uniform() # float(raw_input('What speed do you want for the drone at '+str(ne
                 run.drone_info.append( (newPoint, speed) )
                patchSize = (xlim[1]-xlim[0])/40.0
                 print Style.RESET_ALL
                 ax.add_patch( mpl.patches.Circle( newPoint, radius = patchSize,
                                                   facecolor='#b7e8cc', edgecolor='black' ))
                 ax.text( newPoint[0], newPoint[1], "{:.2f}".format(speed), fontsize=15,
                         horizontalalignment='center', verticalalignment='center')
                 ax.set_title('Number of drones inserted: ' +\
                              str(len(run.drone_info)), fontdict={'fontsize':25})
             elif event.button == 3:
```

```
# Insert big red circles representing the source and target points
             patchSize = (x\lim[1]-x\lim[0])/50.0
             if run.source is None:
                  run.source = (event.xdata, event.ydata)
                  ax.add_patch( mpl.patches.Circle( run.source, radius = patchSize,
                                                     facecolor= '#ffd9d6', edgecolor='black', lw=1.0 ))
                  ax.text( run.source[0], run.source[1], 'S', fontsize=15,
                           horizontalalignment='center', verticalalignment='center')
             elif run.target is None:
                  run.target = (event.xdata, event.ydata)
                  ax.add_patch( mpl.patches.Circle( run.target, radius = patchSize,
                                                    facecolor= '#ffd9d6', edgecolor='black', lw=1.0 ))
                  ax.text( \ run.target[0], \ run.target[1], \ 'T', \ fontsize=15,\\
                           horizontalalignment='center', verticalalignment='center')
             else:
                   print Fore.RED, "Source and Target already set", Style.RESET_ALL
         # Clear polygon patches and set up last minute \verb|ax| tweaks
         clearAxPolygonPatches(ax)
         applyAxCorrection(ax)
         fig.canvas.draw()
  return _enterPoints
# The key-stack argument is mutable! I am using this hack to my advantage.
def wrapperkeyPressHandler(fig, ax, run):
       def _keyPressHandler(event):
           if event.key in ['i', 'I']:
                # Select algorithm to execute
                algo_str = raw_input(Fore.YELLOW
                        "Enter algorithm to be used to compute the tour:\n Options are:\n"
                                    One Dimensional Wavefront \n"
                        Style.RESET_ALL)
                algo_str = algo_str.lstrip()
                # Incase there are patches present from the previous clustering, just clear them
                clearAxPolygonPatches(ax)
                    algo_str == 'odw':
                      tour = run.get_tour( algo_odw, plot_tour_p=True )
                else:
                      print "Unknown option. No horsefly for you! ;-D "
                      sys.exit()
                applyAxCorrection(ax)
                fig.canvas.draw()
           elif event.key in ['c', 'C']:
                # Clear canvas and states of all objects
                run.clearAllStates()
                ax.cla()
                applyAxCorrection(ax)
                ax.set_xticks([])
                ax.set_yticks([])
                fig.texts = []
                fig.canvas.draw()
       return _keyPressHandler
```

```
# Set up interactive canvas
   fig, ax = plt.subplots()
   run = Single_PHO_Input()
   from matplotlib import rc
   # specify the custom font to use
   plt.rcParams['font.family'] = 'sans-serif'
   plt.rcParams['font.sans-serif'] = 'Times New Roman'
   xlim = utils_graphics.xlim
   ylim = utils_graphics.ylim
   ax.set_xlim([xlim[0], xlim[1]])
   ax.set_ylim([ylim[0], ylim[1]])
   ax.set_aspect(1.0)
   ax.set_xticks([])
   ax.set_yticks([])
   ax.set_title("Enter drone positions, source and target onto canvas. \n \
(Enter speeds into the terminal, after inserting a drone at a particular position)")
   mouseClick = wrapperEnterRunPoints (fig,ax, run)
   fig.canvas.mpl_connect('button_press_event' , mouseClick)
                = wrapperkeyPressHandler(fig,ax, run)
   keyPress
   fig.canvas.mpl_connect('key_press_event', keyPress
   plt.show()
class Multiple_PHO_Input:
   def __init__(self, drone_info = [] , sources = [], targets=[]):
          self.drone_info = drone_info
          self.sources
                           = sources
          self.targets
                           = targets
   def get_drone_pis (self):
          return [self.drone_info[idx][0] for idx in range(len(self.drone_info)) ]
   def get_drone_uis (self):
          return [self.drone_info[idx][1] for idx in range(len(self.drone_info)) ]
   def get_tour(self, algo, plot_tour_p=False):
          return algo( self.drone_info,
                       self.sources,
                        self.targets.
                       plot_tour_p
   # Methods for \verb|ReverseHorseflyInput|
   def clearAllStates (self):
         self.drone_info = []
         self.sources = []
         self.targets = []
```

```
def multiple_pho_run_handler():
   import random
   def wrapperEnterRunPoints(fig, ax, run):
     def _enterPoints(event):
        if event.name
                         == 'button_press_event'
                                                            and \
           (event.button == 1 or event.button == 3)
                                                            and \
           event.dblclick == True and event.xdata != None and event.ydata != None:
            if event.button == 1:
                 # Insert circle representing the initial position of a drone
                 print Fore.GREEN
                 newPoint = (event.xdata, event.ydata)
                         = np.random.uniform() # float(raw_input('What speed do you want for the drone at '+str(ne
                 run.drone_info.append( (newPoint, speed) )
                 patchSize = (x\lim[1]-x\lim[0])/40.0
                 print Style.RESET_ALL
                 ax.add_patch( mpl.patches.Circle( newPoint, radius = patchSize,
                                                   facecolor='#EBEBEB', edgecolor='black' ))
                 ax.text( newPoint[0], newPoint[1], "{:.2f}".format(speed), fontsize=10,
                          horizontalalignment='center', verticalalignment='center' )
                 ax.set_title('Number of drones inserted: ' +\
                              str(len(run.drone_info)), fontdict={'fontsize':25})
             elif event.button == 3:
                 # distinct colors, obtained from https://sashat.me/2017/01/11/list-of-20-simple-distinct-colors/
                 cols = ['#e6194b', '#3cb44b', '#ffe119', '#4363d8', '#f58231',
                 '#911eb4', '#46f0f0', '#f032e6', '#bcf60c', '#fabebe',
                 '#008080', '#e6beff', '#9a6324', '#fffac8', '#800000',
                 '#aaffc3', '#808000', '#ffd8b1', '#000075', '#808080']
                 # Insert big colored circles representing the source and target points
                 patchSize = (xlim[1]-xlim[0])/50.0
                 if (len(run.sources) + len(run.targets)) % 2 == 0:
                        run.sources.append((event.xdata, event.ydata))
                        ax.add_patch( mpl.patches.Circle( run.sources[-1], radius = patchSize,
                                                   facecolor= cols[len(run.sources) % len(cols)], edgecolor='black'
                        ax.text( run.sources[-1][0], run.sources[-1][1], 'S'+str(len(run.sources)), fontsize=15,
                                 horizontalalignment='center', verticalalignment='center')
                 else:
                     run.targets.append((event.xdata, event.ydata))
                     ax.add_patch( mpl.patches.Circle( run.targets[-1], radius = patchSize,
                                                       facecolor= cols[len(run.sources)%len(cols)], edgecolor='blac
                      ax.text( run.targets[-1][0], run.targets[-1][1], 'T'+str(len(run.targets)), fontsize=15,
                               horizontalalignment='center', verticalalignment='center')
             # Clear polygon patches and set up last minute \verb|ax| tweaks
             clearAxPolygonPatches(ax)
             applyAxCorrection(ax)
             fig.canvas.draw()
     return _enterPoints
   # The key-stack argument is mutable! I am using this hack to my advantage.
   def wrapperkeyPressHandler(fig, ax, run):
```

```
def _keyPressHandler(event):
           if event.key in ['i', 'I']:
                # Select algorithm to execute
                algo_str = raw_input(Fore.YELLOW
                        "Enter algorithm to be used to compute the tour:\n Options are:\n"
                        " (matchmove)
                                          Repeated bottleneck matching \n"
                                                                                                       +\
                        Style.RESET_ALL)
                algo_str = algo_str.lstrip()
                # Incase there are patches present from the previous clustering, just clear them
                clearAxPolygonPatches(ax)
                if algo_str == 'matchmove':
                      tour = run.get_tour( algo_matchmove, plot_tour_p=True )
                else:
                      print "Unknown option. No horsefly for you! ;-D "
                      sys.exit()
                applyAxCorrection(ax)
                fig.canvas.draw()
           elif event.key in ['c', 'C']:
                # Clear canvas and states of all objects
                run.clearAllStates()
                ax.cla()
                applyAxCorrection(ax)
                ax.set_xticks([])
                ax.set_yticks([])
                fig.texts = []
                fig.canvas.draw()
       return _keyPressHandler
# Set up interactive canvas
fig, ax = plt.subplots()
run = Multiple_PHO_Input()
from matplotlib import rc
# specify the custom font to use
plt.rcParams['font.family'] = 'sans-serif'
plt.rcParams['font.sans-serif'] = 'Times New Roman'
xlim = utils_graphics.xlim
ylim = utils_graphics.ylim
ax.set_xlim([xlim[0], xlim[1]])
ax.set_ylim([ylim[0], ylim[1]])
ax.set_aspect(1.0)
ax.set_xticks([])
ax.set_yticks([])
ax.set_title("Enter drone positions, sources and targets onto canvas.")
mouseClick = wrapperEnterRunPoints (fig,ax, run)
fig.canvas.mpl_connect('button_press_event' , mouseClick)
keyPress
             = wrapperkeyPressHandler(fig,ax, run)
```

```
fig.canvas.mpl_connect('key_press_event', keyPress )
plt.show()
```

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 $\begin{array}{l} {\rm Fragment\ referenced\ in\ 3.} \\ {\rm Uses:\ algo_matchmove\ 14a,\ algo_odw\ 7.} \end{array}$

Appendix F

 $\langle Plotting 33 \rangle \equiv$

Implementation of (**Plotting**)

We typically plot the tours onto a separate window if the boolean switch plot_tour_p is set to True while calling the algorithm. The path of the package is shown in bold red. The paths of the drones from their initial positions to the point where they pick up the package from another drone are shown in blue.

An example output from the plot_tour function is shown below.

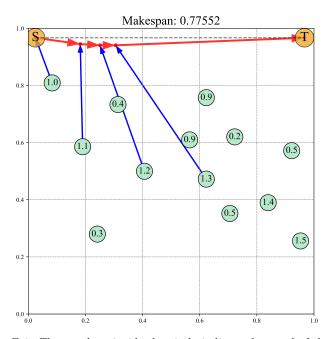


Figure F.1: The numbers inside the circle indicate the speed of the drone

```
import matplotlib.ticker as ticker
ax.set_aspect(aspect_ratio)
ax.set_xlim(xlims)
ax.set_ylim(ylims)
plt.rc('font', family='serif')
# Draw the package trail
xs, ys = extract_coordinates(package_trail)
ax.plot(xs,ys, 'ro', markersize=5 )
for idx in range(len(xs)-1):
     plt.arrow( xs[idx], ys[idx], xs[idx+1]-xs[idx], ys[idx+1]-ys[idx],
                **{'length_includes_head': True,
                   'width': 0.007,
                   'head_width':0.025,
                   'fc': 'r',
                   'ec': 'none',
                   'alpha': 0.8})
# Draw the source, target, and initial positions of the robots as bold dots
xs,ys = extract_coordinates([source, target])
ax.plot(xs,ys, 'o', markersize=sourcetargetmarkersize, alpha=1.0, ms=10, mec='k', mfc='#F1AB30')
#ax.plot(xs,ys, 'k--', alpha=0.6 ) # light line connecting source and target
ax.text(source[0], source[1], 'S', fontsize=sourcetargetmarkerfontsize,\
        horizontalalignment='center', verticalalignment='center')
ax.text(target[0], target[1], 'T', fontsize=sourcetargetmarkerfontsize.\
        horizontalalignment='center',verticalalignment='center')
xs, ys = extract_coordinates( [ drone_info[idx][0] for idx in range(len(drone_info)) ] )
ax.plot(xs,ys, 'o', markersize=speedmarkersize, alpha = 1.0, mec='None', mfc='#b7e8cc' )
# Draw speed labels
for idx in range(len(drone_info)):
     ax.text( drone_info[idx][0][0], drone_info[idx][0][1], format(drone_info[idx][1],'.2f'),
              fontsize=speedfontsize, horizontalalignment='center', verticalalignment='center' )
# Draw drone path from initial position to interception point
for pt, idx in zip(package_trail, used_drones):
     initdroneposn = drone_info[idx][0]
     handoffpoint = pt
     xs, ys = extract_coordinates([initdroneposn, handoffpoint])
     plt.arrow( xs[0], ys[0], xs[1]-xs[0], ys[1]-ys[0],
                **{'length_includes_head': True,
                   'width': 0.005,
                   'head_width':0.02,
                   'fc': 'b',
                   'ec': 'none'})
fig.suptitle(figtitle, fontsize=15)
ax.set_title('\nMakespan: ' + format(makespan(drone_info, used_drones, package_trail),'.5f'), fontsize=16)
startx, endx = ax.get_xlim()
starty, endy = ax.get_ylim()
```

Fragment referenced in 3. Uses: extract_coordinates 12a.

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Todo list

TODO!	 				 				 																4