

Big Data Computing

Master's Degree in Computer Science

2025-2026



SAPIENZA
UNIVERSITÀ DI ROMA

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Analysis of (Large) Graphs

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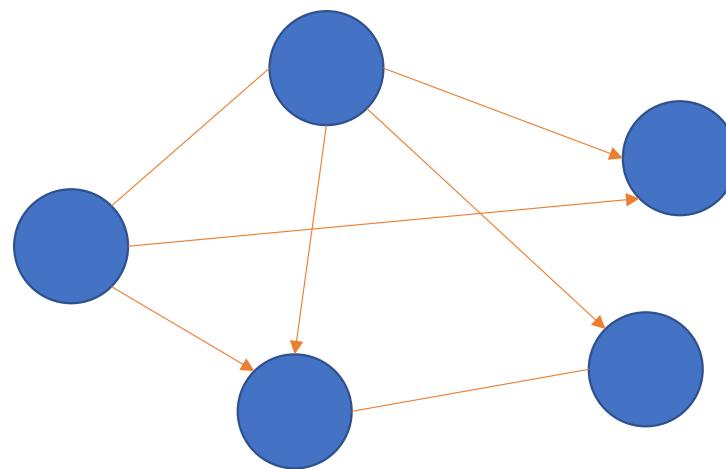
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 - **Social Networks** (i.e., the set of social connections between people)
 - ...

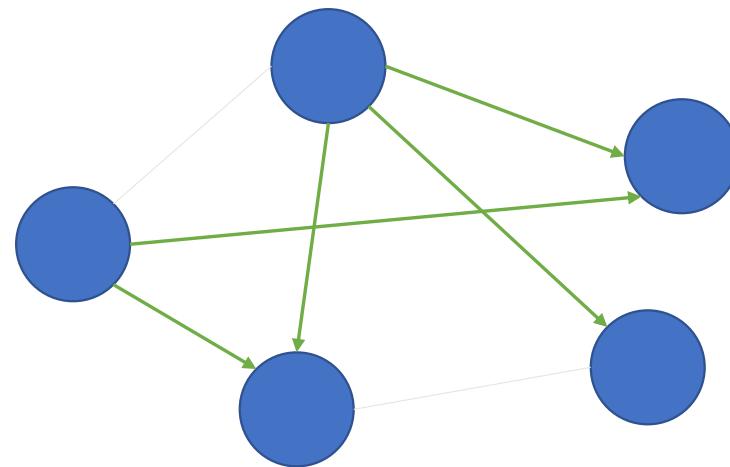
What is a Graph?

Informally, a set of **vertices** (nodes)
connected by a set of **edges** (links)



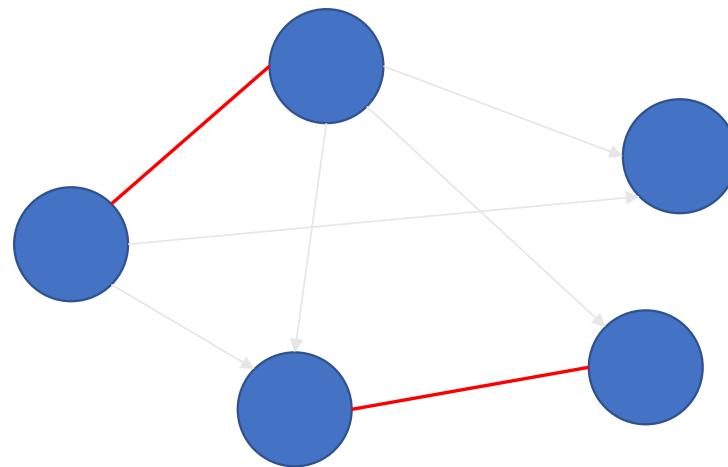
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edges may be directed



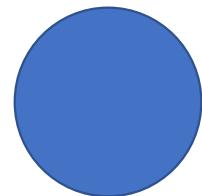
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edges may be **undirected**



Directed vs. Undirected

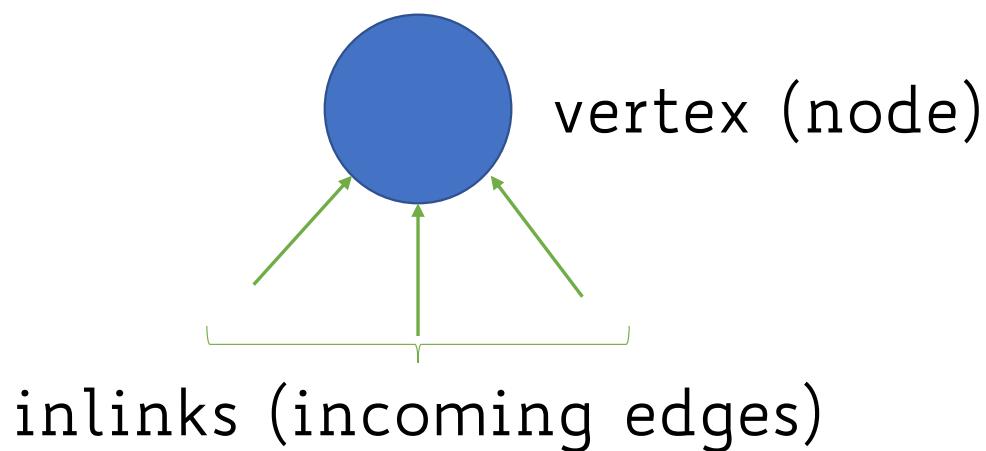
Directed



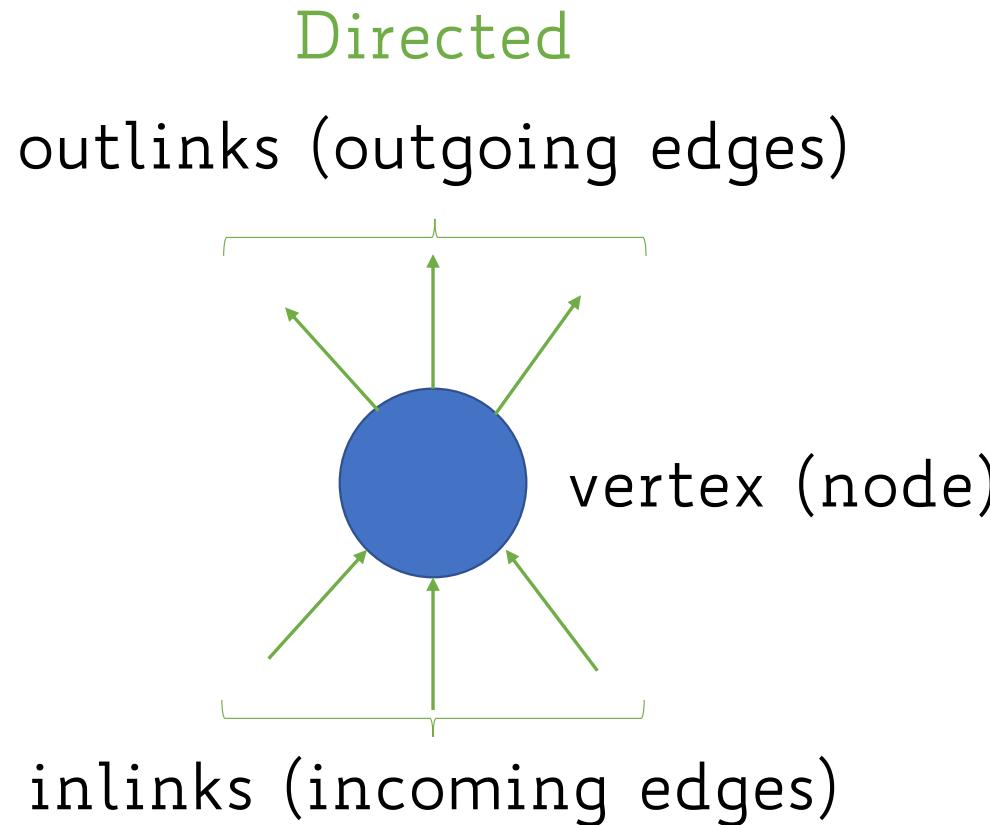
vertex (node)

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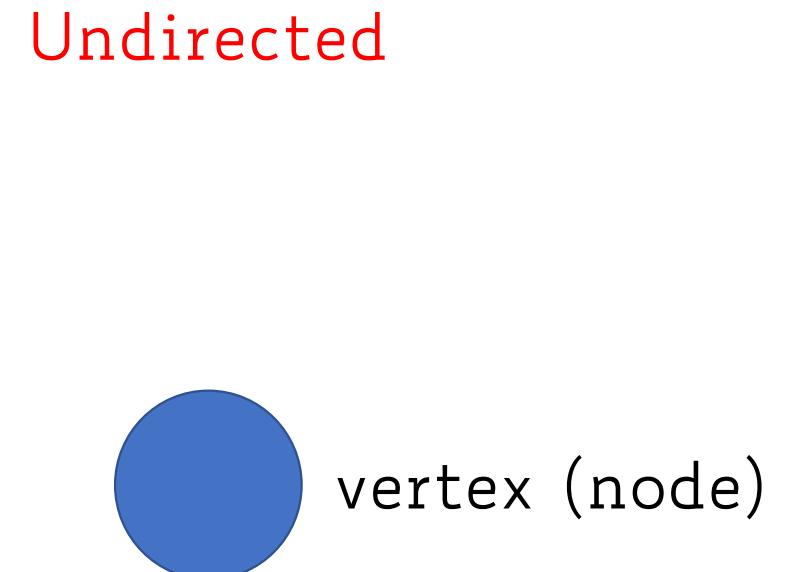
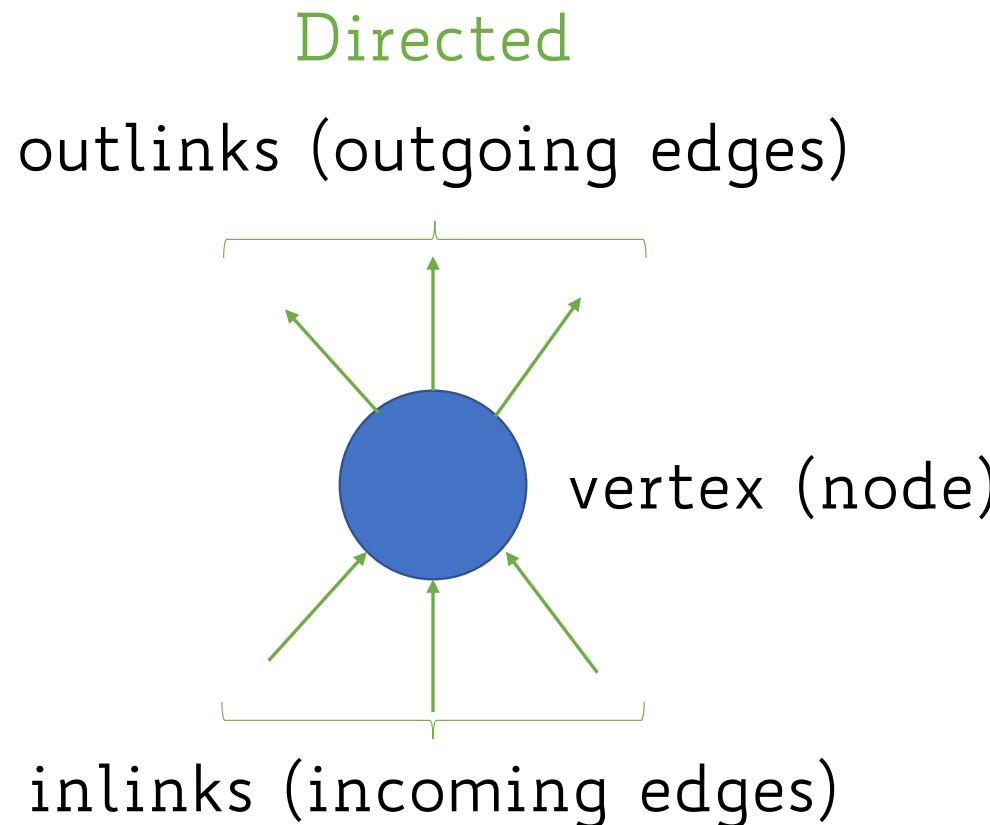
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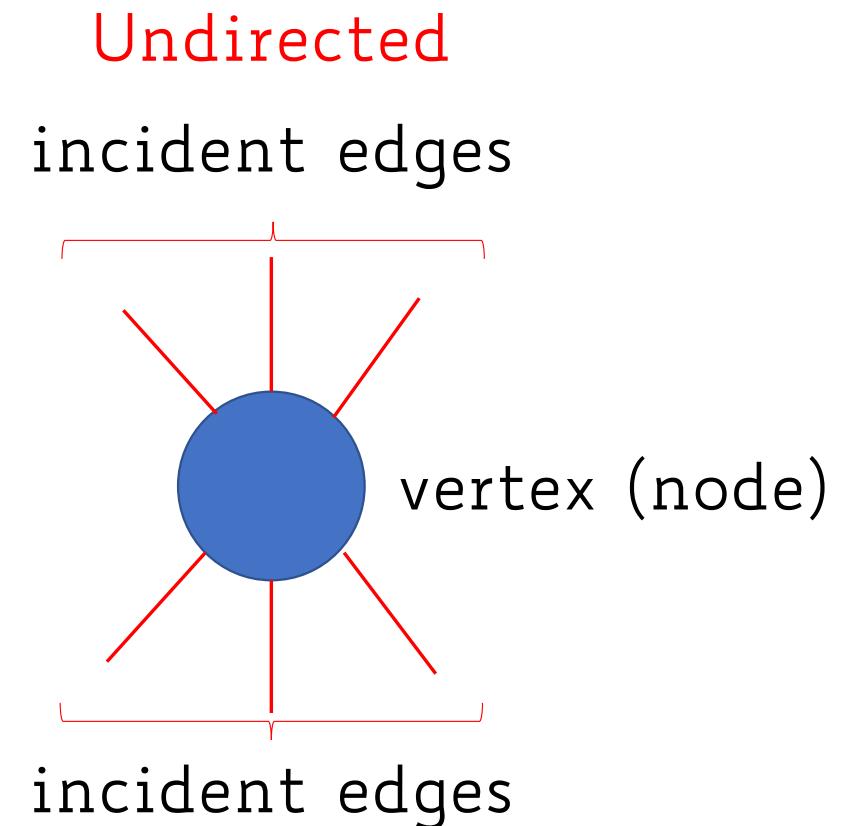
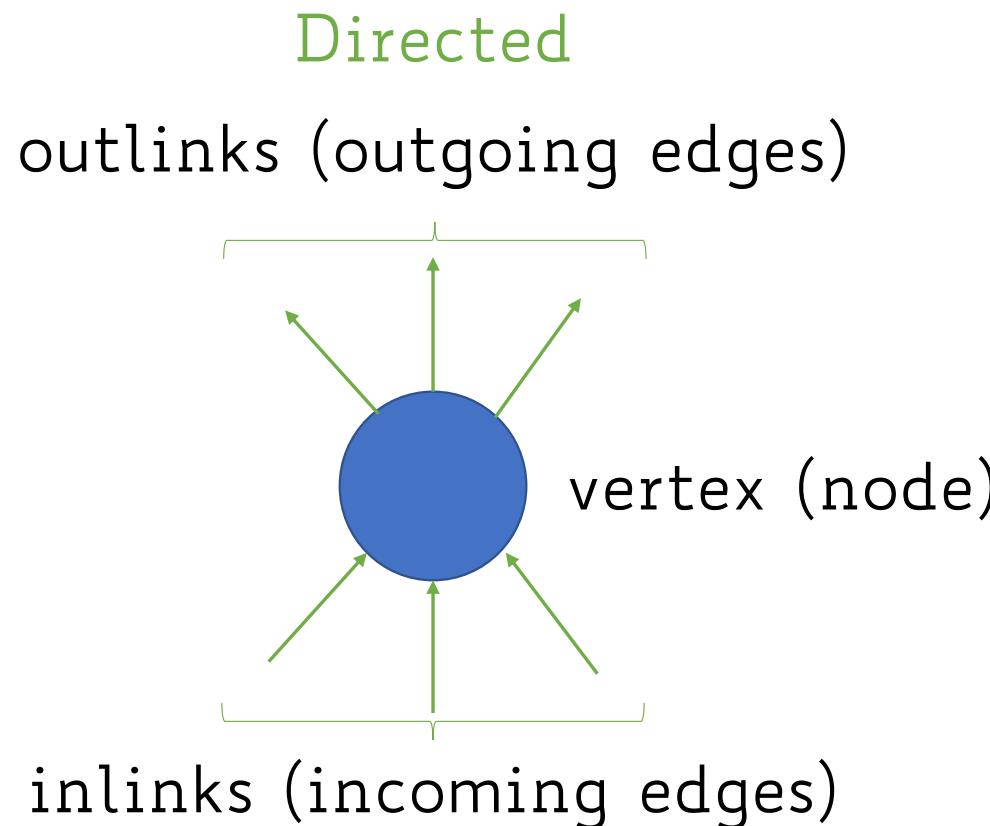
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Graph Formalism

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Note that an **undirected** graph is just a special case of a **directed** graph where the set of edges contain symmetric pairs of vertices

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Intuitively, the number of inbound/incident links to a node

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To be more explicit, in the case of a directed graph sometimes we distinguish between **in-degree** and **out-degree**

$$\text{in-deg}(v) = |\{u \in V | (u, v) \in E\}|$$

$$\text{out-deg}(v) = |\{u \in V | (v, u) \in E\}|$$

How Do We Represent Graphs?

3 main ways of representing graphs

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Adjacency
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Edge Lists

Adjacency Matrix

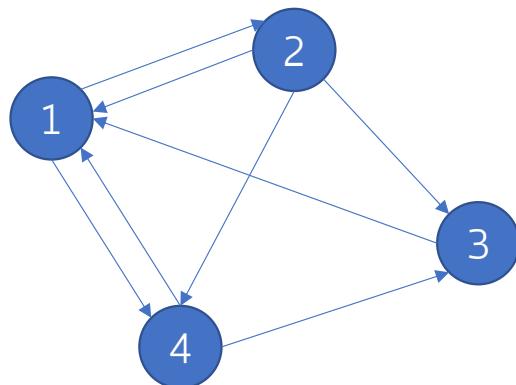
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 - $M[i, j] = 1$ iff there exists an edge from vertex v_i to vertex v_j

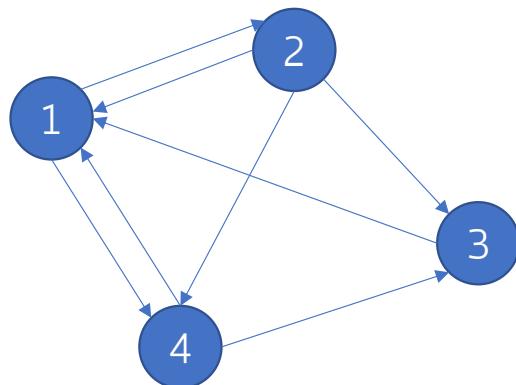
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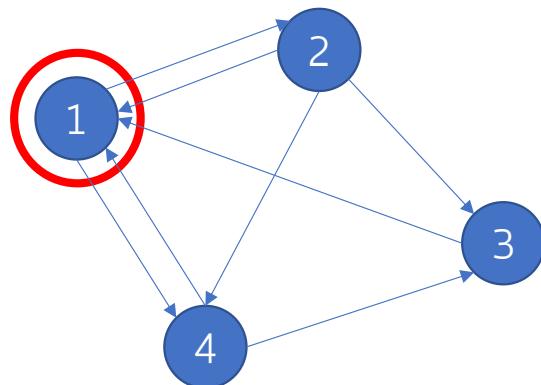
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- CONs:
 - Space inefficient (especially for loosely connected graphs, i.e., sparse matrices)
 - Easy to write yet hard to compute

Adjacency Lists

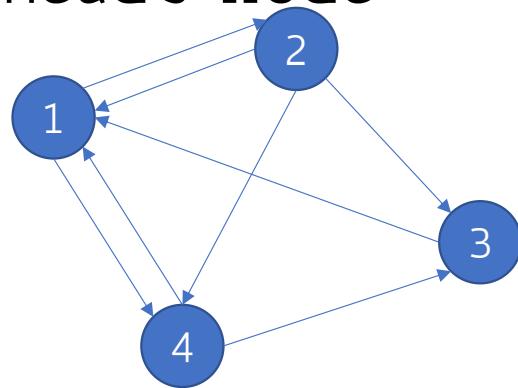
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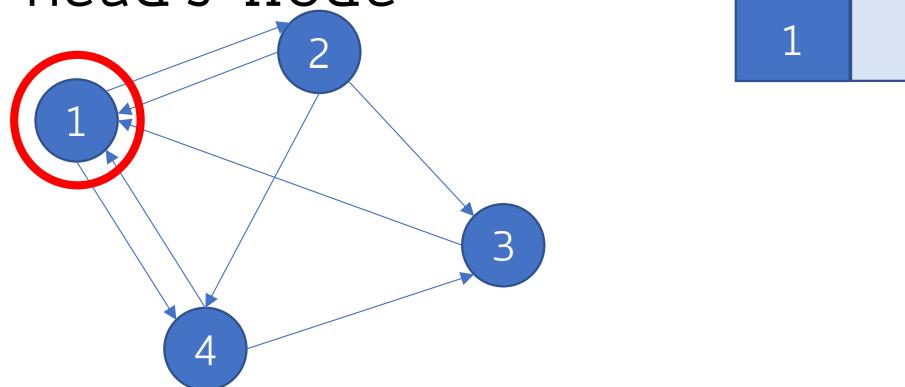
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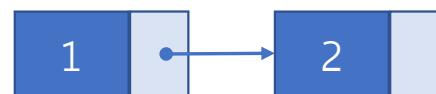
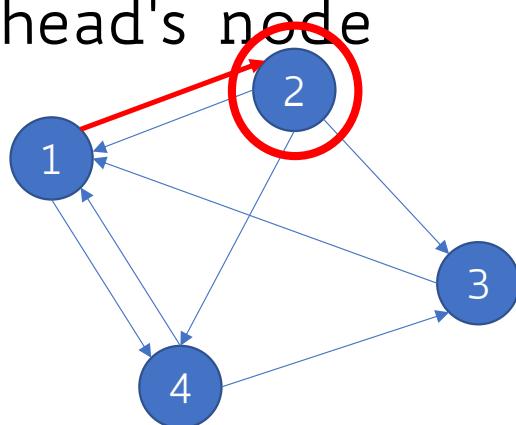
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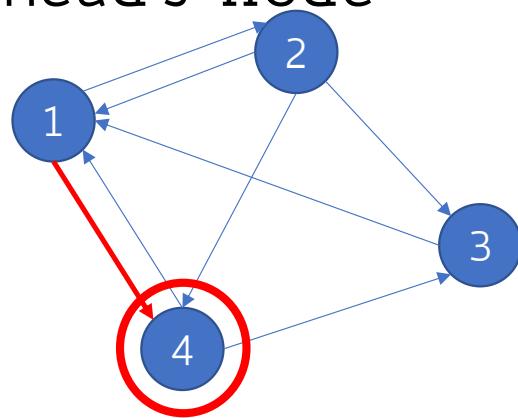
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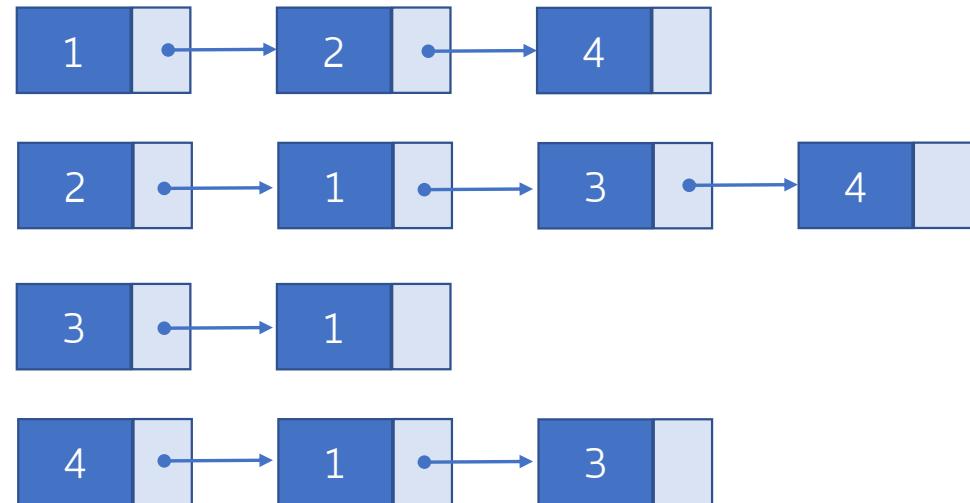
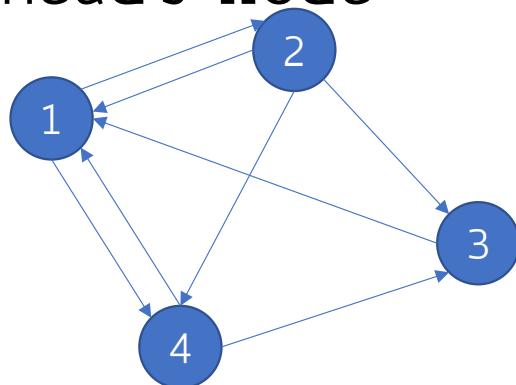
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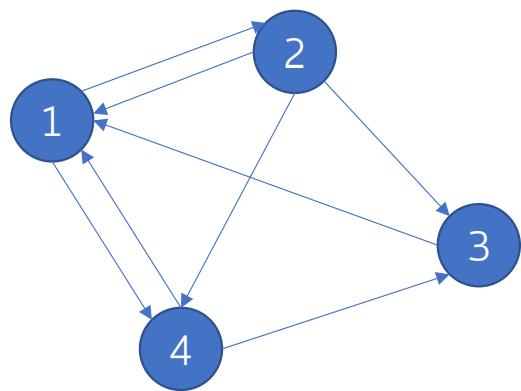
Note that with adjacency matrix, any computation over incoming (outgoing) links reduces to a column (row) scan of the matrix

Edge Lists

- Explicitly enumerates all the edges

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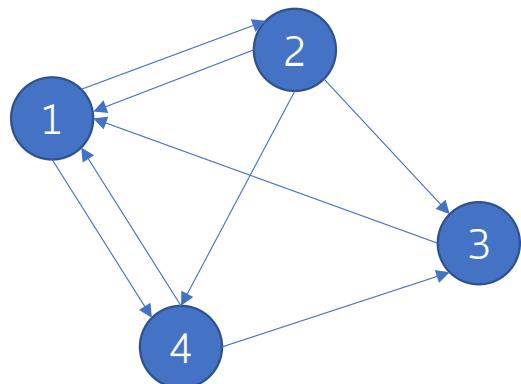
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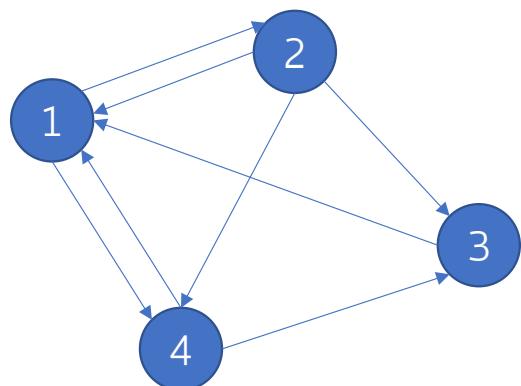
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Waste of space

Some Famous Graph Problems

Problems

Applications

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Finding Shortest Paths	Routing IP packets, GPS navigation systems
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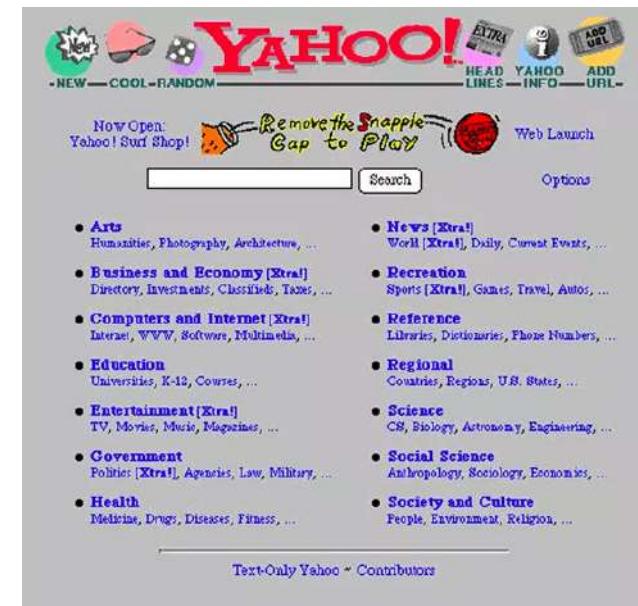
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- The **Web graph** is a great test bed for link analysis

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- Other attempts: DMOZ, LookSmart



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- The list of top- k documents most similar to a query are returned (e.g., measuring **cosine similarity** between each query-document pair)

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The Web is huge and full of untrusted documents!

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Trustworthy web pages should point to each other

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Academic Home News HERCULE Lab Publications Teaching Contact CV



Biography

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The screenshot shows the official website of Sapienza Università di Roma. At the top, there is a search bar and a navigation menu with links for STUDENTI, LAUREATI, TERRITORIO, and CONTATTI. Below the menu, there is a large image of a person using a laptop, with the text "Lezioni, esami e lauree a distanza" overlaid. To the right of the image are several colored boxes representing different services: COORS E ORIZZONTE (orange), RICERCA SCIENTIFICA (green), INTERNAZIONALE (blue), ATENEO (dark green), DOCENTI (red), and PERSONALE (maroon). At the bottom, there is a search bar with the placeholder "Cerca il tuo corso" and a magnifying glass icon.

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[Email](#) [Twitter](#) [LinkedIn](#) [Google Scholar](#) [ORCID](#) [cv](#)

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STUDENTI LAUREATI TERRITORIO CONTATTI



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CORSI E INIZIATIVE

- CORSI E INIZIATIVE
- RICERCA SCIENTIFICA
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Cerca il tuo corso

Welcome to the United Nations

United Nations | Peace, dignity and equality on a healthy planet

LIVE NOW | SEARCH | 

About the UN | What we do | What we work | News and Media | Documents | Resources | Contact Us | English | Français | Español | 中文 | 日本語 | 简体中文 | 繁體中文

Climate change and COVID-19: Call to 'recover better'

At the worldwide call to 'post-pandemic recovery' the United Nations calls Governments to seize the opportunity to "Build back better" by creating more sustainable, resilient and inclusive economies. The UN is defining a blueprint for a sustainable and inclusive recovery that one billion and counting are being taken to ensure a more resilient future. Secretary-General António Guterres proposed six climate-related actions to shape the recovery. While UNEP works closely to build scientific knowledge on links between ecosystem stability and human health.



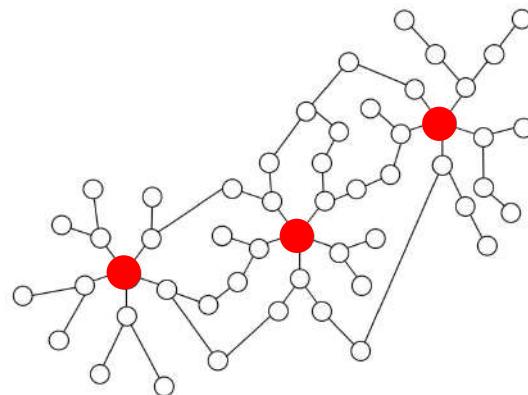
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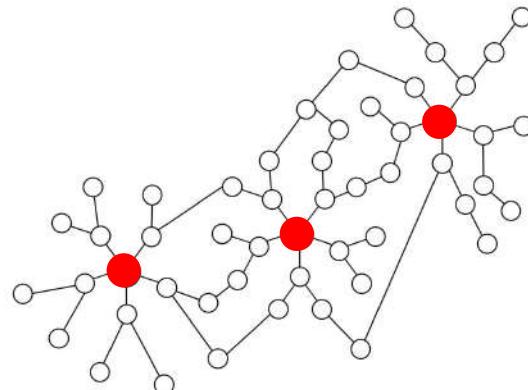
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Rank nodes (i.e., assign them an importance score) on the basis of their connectivity

The Web as a "Scale-Free" Network

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They observed the degree distribution follows a **power law**

They refer to graphs (i.e., networks) exhibiting such a behavior as **scale-free networks**

Barabási, A.L. and Albert,R. Emergence of scaling in random networks, Science, 286:509-512, October 15, 1999

Random Graph: The Erdős–Rényi Model

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The degree of a node follows a **Binomial distribution**

$$\deg(u) = \text{Binomial}(n-1, p)$$

The Erdős-Rényi Model: Approximations

Poisson (sparse graphs):

If n is large and p is small, s.t. $\lambda = p(n-1) \rightarrow \deg(u) \approx \text{Poisson}(\lambda)$

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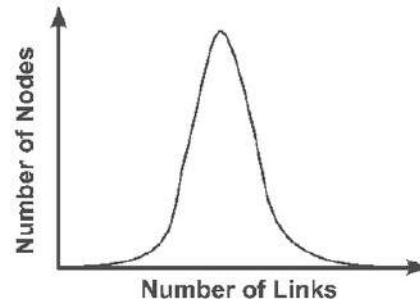
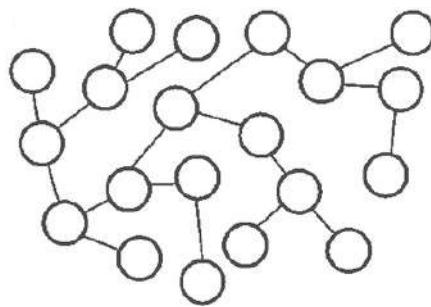
Poisson (sparse graphs):

If n is large and p is small, s.t. $\lambda = p(n-1) \rightarrow \deg(u) \approx \text{Poisson}(\lambda)$

Normal (dense graphs):

If n is large and p is not too small $\rightarrow \deg(u) \approx \text{Normal}(\mu, \sigma^2)$,
where $\mu = p(n-1)$; $\sigma^2 = (n-1)p(1-p)$

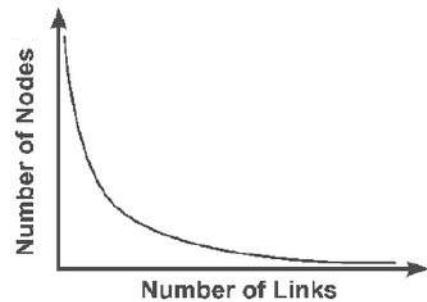
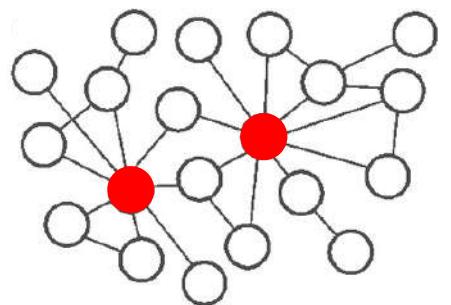
The Web as a Scale-Free Network



Random Graph

Most nodes have approximately the same number of links producing a bell-shaped curve of the degree distribution

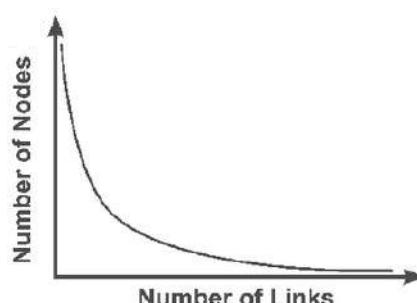
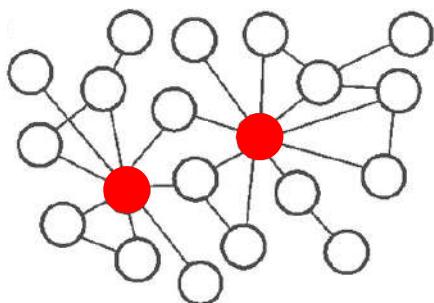
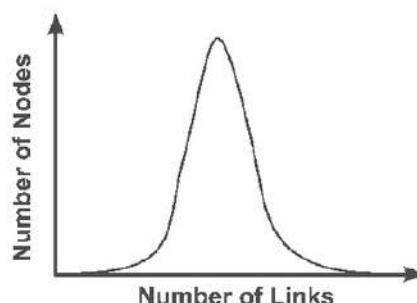
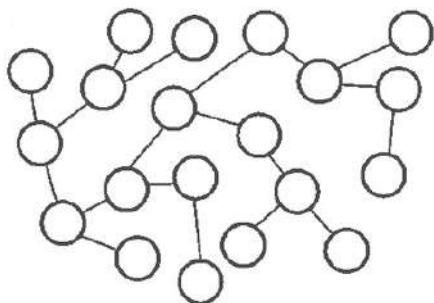
The Web as a Scale-Free Network



Scale-Free Graph

Most nodes have few links, and few nodes (i.e., red ones) have a large number of links, resulting into a power law degree distribution

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The ratio of very connected nodes to the number of nodes in the rest of the network remains constant as the network changes in size

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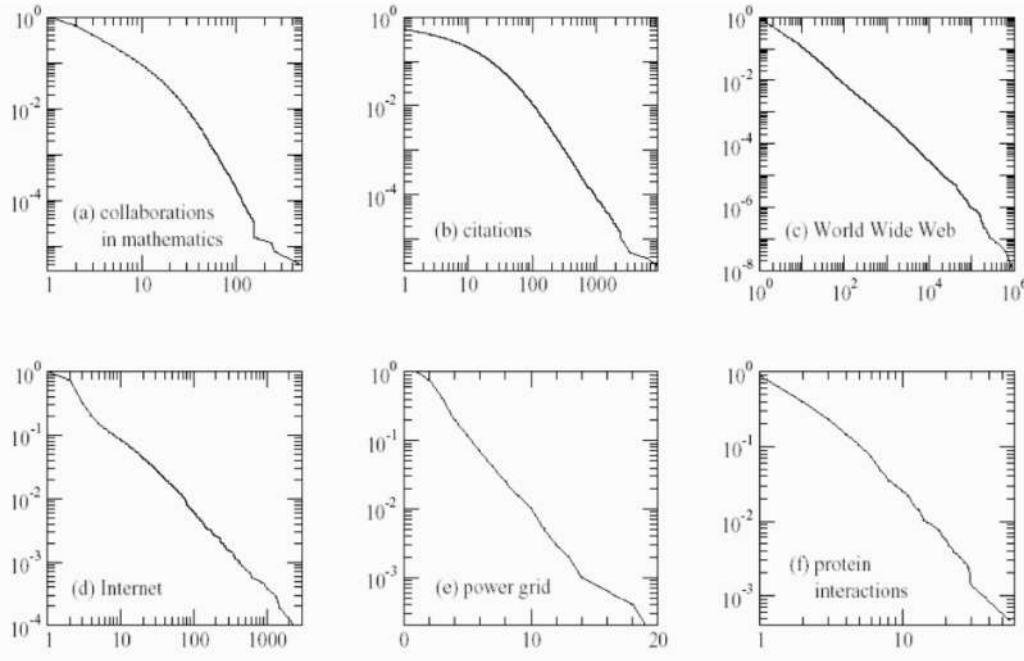
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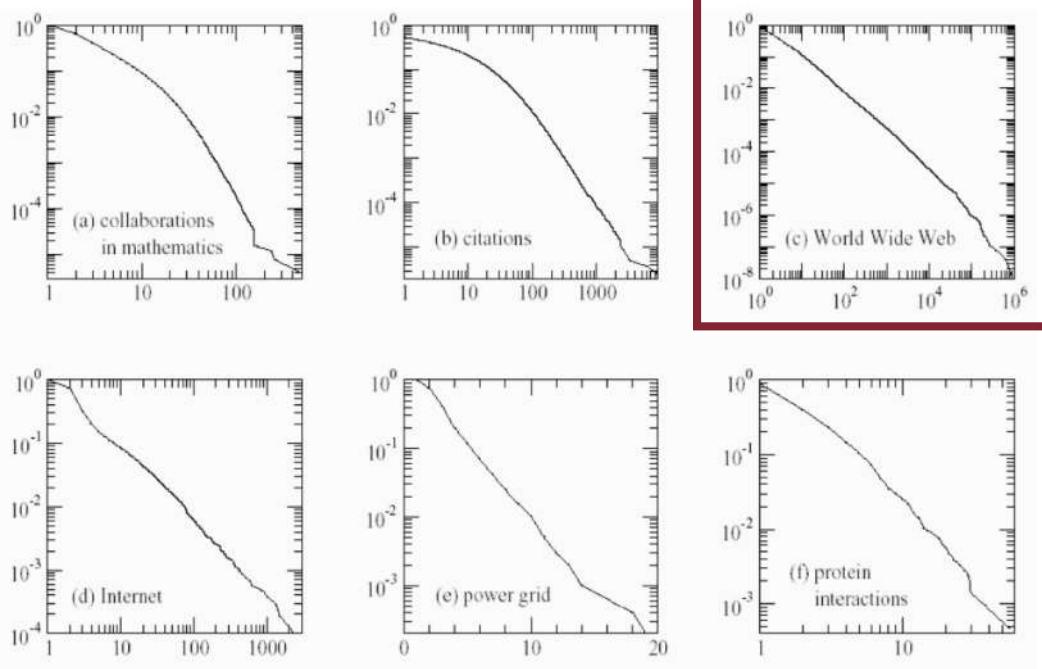
$$p(\text{linking to node } i) \propto \frac{k_i}{\sum_j k_j}$$

Scale-Free Networks: Examples



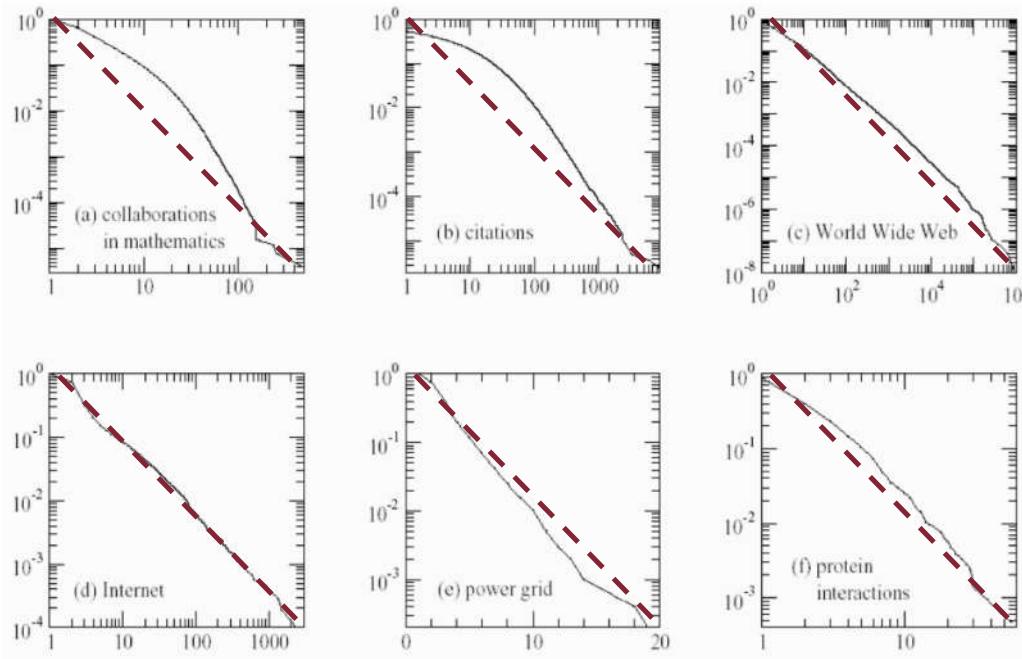
Many real-world networks
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Scale-Free Networks: Examples



The Web is one of
those!

Scale-Free Networks: Examples



On log-log scale power law distributions look like straight lines

$$\log(p(k)) = \log(\alpha k^{-\gamma}) = \underbrace{\log(\alpha)}_{\text{constant } q} + \log(k^{-\gamma}) = q - \gamma \log(k)$$

Computing Node Importance

Several link analysis approaches to compute web page importance

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PageRank

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Hubs and Authorities
(HITS)

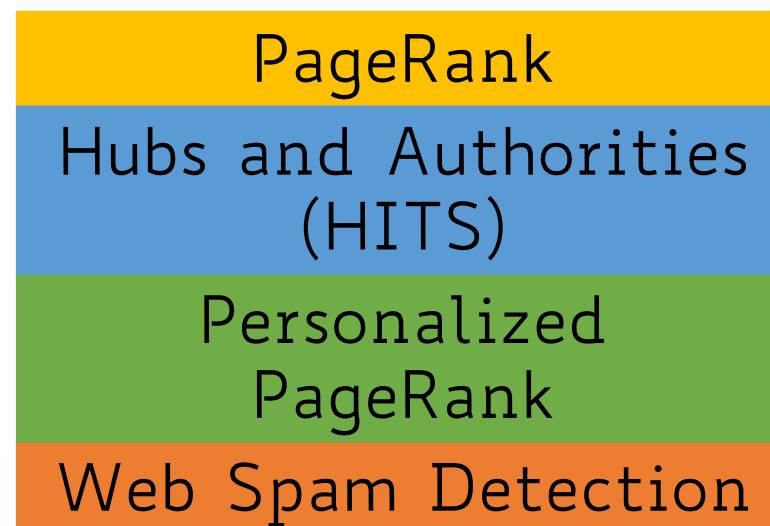
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Take-Home Message of Today

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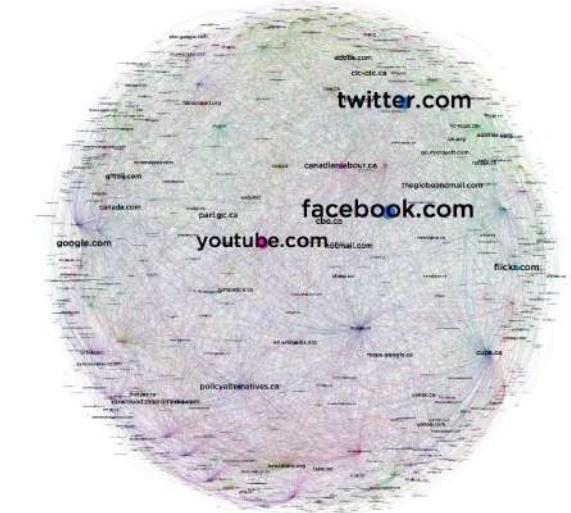
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- Examples:
 - Finding the shortest path on a map
 - Computing the importance of a page in the Web graph
 - Suggesting friends in a social network graph
- Many techniques exist to approach the problems above

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- Idea: Use node's connectivity to determine the **importance of a node**