Sistemi Operativi I

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SAPIENZA UNIVERSITÀ DI ROMA

Recap from Last Lecture

- Concurrent accesses to shared resources by multiple cooperating processes/threads can lead to unexpected behavior

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- Process/Thread cooperation must guarantee consistency of any shared data/resource, regardless of CPU scheduling
- Maintaining shared data consistency requires mechanisms to ensure synchronized execution of critical sections by processes/threads
- Critical sections are specific pieces of code which contain shared resources that need to be "protected"

Recap from Last Lecture

We need to have appropriate "tools" (i.e., primitive constructs) provided by programming languages used as atomic building blocks for synchronization

- Locks → At each time, only one process holds a lock, executes its critical section, and finally releases the lock
- Semaphores → A generalization of locks
- Monitors → To connect shared data to synchronization primitives

Require some HW support and waiting

Locks

- Provide mutual exclusion to shared data using 2 atomic primitives:
  - lock.acquire() → wait until the lock is free, then grab it
  - lock.release() → unlock and wake up any thread waiting in acquire()
- Rules for using a lock:
  - Always acquire the lock before accessing shared data
  - Always release the lock after finishing with shared data
  - Lock must be initially free
- Only one process/thread can acquire the lock, others will wait!

Too Much Milk: Solution Using Locks

Use lock primitives

```
# Thread Bob

lock.acquire()

if (!milk):
    buy_milk()

lock.release()
```

```
# Thread Carla

lock.acquire()

if (!milk):
    buy_milk()

lock.release()
```

This solution is clean and symmetric

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This solution is clean and symmetric

Q: How do we make acquire() and release() atomic?

HW Support for Synchronization

Implementing high-level synchronization primitives requires low-level hardware support

High-level atomic operations (SW)	lock, monitor, semaphore, send/receive
Low-level atomic operations (HW)	disabling interrupts, atomic instructions (test&set)

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  - external events → interrupts (e.g., time slice) cause the scheduler to take over the currently running thread

We want to prevent the CPU scheduler to take control while an acquire() operation is ongoing

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- On single-CPU systems, we can prevent the scheduler to take over by:
  - internal events → discouraging threads from requesting any I/O operation within a critical section
  - external event → disabling interrupts (i.e., telling the HW to delay the handling of any external event until the current thread is done with the critical section)

We cover all the possible cases where the current thread might loose control of the CPU, either voluntarily (due to internal events) or involuntarily (due to external events)

Implementing Locks: Disabling Interrupts

```
Class Lock {
  public void acquire(Thread t);
  public void release();

  Lock() {}
}
```

Implementing Locks: Disabling Interrupts

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Class Lock {
  public void acquire(Thread t);
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  Lock() {}
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```
public void acquire(Thread t) {
    disable_interrupts();
}
```

Implementing Locks: Disabling Interrupts

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Class Lock {
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  Lock() {}
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```
public void acquire(Thread t) {
    disable_interrupts();
}
```

```
public void release() {
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Implementing Locks: Disabling Interrupts

```
Class Lock {
  public void acquire(Thread t);
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  Lock() {}
}
```

We need both acquire and release being implemented as system calls

```
public void acquire(Thread t) {
    disable_interrupts();
}
```

```
public void release() {
    enable_interrupts();
}
```

Implementing Locks: Disabling Interrupts

```
Class Lock {
  public void acquire(Thread t);
  public void release();

  Lock() {}
}
```

We need both acquire and release being implemented as system calls

Why?

```
public void acquire(Thread t) {
    disable_interrupts();
}
```

```
public void release() {
    enable_interrupts();
}
```

Implementing Locks: Disabling Interrupts

- PROs:
  - Very simple!

Implementing Locks: Disabling Interrupts

- PROs:
  - Very simple!
- CONs:
  - Privileged instructions
  - Trust no one abuses this capability, e.g.:
    - A greedy program that gets the lock and hogs the CPU
    - A malicious program that gets the lock and goes into an infinite loop
  - Does not work on multiprocessors!
  - May loose relevant interrupts
  - Masking/Unmasking interrupts is inefficient

Implementing Locks: A First Attempt

```
Class Lock {
  public void acquire(Thread t);
  public void release();
  private int flag; // 0=free; 1=busy

  Lock() {
     this.flag = 0; // initially free
  }
}
```

Implementing Locks: A First Attempt

```
Class Lock {
  public void acquire(Thread t);
  public void release();
  private int flag; // 0=free; 1=busy

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```
public void release() {
  this.flag = 0; // set the flag to free
}
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  public void acquire(Thread t);
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  Lock() {
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  }
}
```

```
public void acquire(Thread t) {
    while(this.flag == 1) { // test flag
       // do nothing (spin-wait)
    }
    this.flag = 1; // set the flag to busy
}
```

```
public void release() {
  this.flag = 0; // set the flag to free
}
```

Does this solution work?

Implementing Locks: A First Attempt

An unlucky (yet plausible) trace

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time

Thread 1

Thread 2

Implementing Locks: A First Attempt

An unlucky (yet plausible) trace

time

Thread 1

flag = 0

Thread 2

```
call acquire()
while(this.flag == 1)
```

Implementing Locks: A First Attempt

An unlucky (yet plausible) trace

time

Thread 1

Thread 2

call **acquire()**
while(this.flag == 1)
interrupt: switch to Thread 2

Multiple machine-code instructions

Implementing Locks: A First Attempt

An unlucky (yet plausible) trace

time

Thread 1

flag = 0

Thread 2

Thread 1:

```
call acquire()
while(this.flag == 1)
interrupt: switch to Thread 2
```

Thread 2:

```
call acquire()
while(this.flag == 1)
this.flag = 1
```

Implementing Locks: A First Attempt

An unlucky (yet plausible) trace

time

Thread 1	Thread 2
call **acquire()** while(this.flag == 1) *interrupt:* switch to Thread 2	
	call **acquire()** while(this.flag == 1) this.flag = 1 *interrupt:* switch to Thread 1

Implementing Locks: A First Attempt

An unlucky (yet plausible) trace

time

Thread 1

flag = 1

Thread 2

Thread 1:

call **acquire()**
while(this.flag == 1)
interrupt: switch to Thread 2

Thread 2:

call **acquire()**
while(this.flag == 1)
this.flag = 1
interrupt: switch to Thread 1

Thread 1:

this.flag = 1 // set to busy too!

Violation of mutual exclusion!

Implementing Locks: A First Attempt

- Testing and setting a variable with MOV-like instructions does not work
- Spin-Waiting is bad!
  - The waiter thread will waste CPU cycles doing nothing
  - On uniprocessors such a waste could be even worse
  - The only thread holding the lock must take its turn on the CPU, otherwise the other(s) spin-waiting will never take it!
- What if the we have a single CPU with a non-preemptive scheduler?

HW Support for Synchronization

Implementing high-level synchronization primitives requires low-level hardware support

High-level atomic operations (SW)	lock, monitor, semaphore, send/receive
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Implementing Locks: Atomic Instructions

- An atomic read-modify-write instruction reads a value from memory into a register and writes a new value in one shot!

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- An atomic read-modify-write instruction reads a value from memory into a register and writes a new value in one shot!
  - On a uniprocessor → straightforward to implement adding a new instruction
  - On a multiprocessor → the processor issuing the instruction must also be able to invalidate any copies of the value other processes may have in their cache

Implementing Locks: Atomic Instructions

- Examples:
  - test&set → writes (sets) 1 to a memory location and returns its old value

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  - compare&swap (x86: compare&exchange) → compares the content of a memory location with a given value and, if they are the same, modifies the content of that location to a new value

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Implementing Locks: test&set

```
Class Lock {
  public void acquire();
  public void release();
  private int flag;

  Lock() {
    // lock is initially free
    this.flag = 0;
  }
}
```

Implementing Locks: test&set

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Class Lock {
  public void acquire();
  public void release();
  private int flag;

  Lock() {
    // lock is initially free
    this.flag = 0;
  }
}
```

```
public void acquire() {
  while(test&set(this.flag) == 1) {
    // while busy do nothing
  }
}
```

Implementing Locks: test&set

```
Class Lock {
  public void acquire();
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Case 1: if lock is free (flag = 0) test&set(flag) will read 0, set it to 1 and return 0

Implementing Locks: test&set

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Class Lock {
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}
```

Case 1: if lock is free (flag = 0) test&set(flag) will read 0, set it to 1 and return 0

The lock is now busy, the boolean expression in the while guard is false and acquire terminates

Implementing Locks: test&set

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Class Lock {
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  public void release();
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  Lock() {
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    this.flag = 0;
  }
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public void acquire() {
  while(test&set(this.flag) == 1) {
    // while busy do nothing
  }
}
```

```
public void release() {
  this.flag = 0;
}
```

Case 2: if lock is busy (flag = 1) test&set(flag) will read 1, set it to 1 and return 1

Implementing Locks: test&set

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Class Lock {
  public void acquire();
  public void release();
  private int flag;

  Lock() {
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  }
}
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```
public void acquire() {
  while(test&set(this.flag) == 1) {
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  }
}
```

```
public void release() {
  this.flag = 0;
}
```

Case 2: if lock is busy (flag = 1) test&set(flag) will read 1, set it to 1 and return 1

The lock is still busy, the boolean expression in the while guard is true and acquire continues to loop until release executes

Atomic Instructions: Any Issue?

```
public void acquire() {
  while(test&set(this.value) == 1) {
    // while busy do nothing
  }
}
```

- What’s (still) bad with the above implementation?

Atomic Instructions: Any Issue?

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Atomic Instructions: Any Issue?

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busy waiting
(*spin-lock*)

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Atomic Instructions: Any Issue?

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who is going to take the lock once released?

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  - What is the CPU doing?
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Masking Interrupts vs. Atomic Instructions

- 3 main problems with disabling interrupts:
  - overhead as it requires kernel privileges
  - trusted code (only OS kernel code can use this mechanism)
  - unfeasible with mulitprocessor architectures

Masking Interrupts vs. Atomic Instructions

- 3 main problems with disabling interrupts:
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  - trusted code (only OS kernel code can use this mechanism)
  - unfeasible with mulitprocessor architectures
- 3 main problems with atomic instructions:
  - busy waiting (*spin-lock*)
  - unfairness as there is no queue where threads wait for the lock to be released
  - performance as N-1 waiter threads may spin for one round each before the thread holding the lock gets its turn on the CPU

Reduce Busy-Waiting: Yield!

```
Class Lock {
  public void acquire();
  public void release();
  private int flag;

  Lock() {
    // lock is initially free
    this.flag = 0;
  }
}
```

```
public void acquire() {
  while(test&set(this.flag) == 1) {
    yield(); // give up the CPU
  }
}
```

```
public void release() {
  this.flag = 0;
}
```

An OS primitive (**yield**) that allows threads to give up the CPU

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An OS primitive (**yield**) that allows threads to give up the CPU

In this way, the waiter thread releases the CPU immediately

Still, some threads may starve in an endlessly yielding!

Reduce Busy-Waiting and Improve Fairness: Sleeping + Queues

```
Class Lock {
  public void acquire(Thread t);
  public void release();
  private int flag;
  private int guard;
  private Queue q;

  Lock() {
    // lock is initially free
    this.flag = 0;
    this.guard = 0;
    this.q = new Queue();
  }
}
```

```
public void acquire(Thread t) {
  while(test&set(this.guard) == 1) {
    // while busy do nothing
  }
  if(this.flag == 0) {
    this.flag = 1; // lock is taken
    this.guard = 0;
  }
  else {
    this.q.push(t);
    this.guard = 0;
    park(); // Solaris primitive
  }
}
```

```
public void release() {
  while(test&set(this.guard) == 1) {
    // while busy do nothing
  }
  if(q.is_empty()) {
    this.flag = 0; // lock is free
  }
  else {
    t = q.pop();
    unpark(t); // Solaris primitive
  }
  this.guard = 0;
}
```

park → Put the caller thread to sleep if it tries to acquire a busy lock

unpark → Wake up the waiter when the lock is free

Reduce Busy-Waiting and Improve Fairness: Sleeping + Queues

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    this.flag = 0; // lock is free
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  else {
    t = q.pop();
    unpark(t); // Solaris primitive
  }
  this.guard = 0;
}
```

NOTE: The flag is not set to 0 when another thread is woken up

The thread releasing the lock passes it directly to the next thread

Reduce Busy-Waiting and Improve Fairness: Sleeping + Queues

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  public void acquire(Thread t);
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  Lock() {
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public void acquire(Thread t) {
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}
```

```
public void release() {
  while(test&set(this.guard) == 1) {
    // while busy do nothing
  }
  if(q.is_empty()) {
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  }
  else {
    t = q.pop();
    unpark(t); // Solaris primitive
  }
  this.guard = 0;
}
```

We can't totally get rid of busy-waiting but we can make it **independent** on how long is the critical section delimited by acquire and release

Reduce Busy-Waiting and Improve Fairness: Sleeping + Queues

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  public void acquire(Thread t);
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```

What would happen if an unlucky thread switch happens right before calling park() and the incoming thread releases the lock?

Reduce Busy-Waiting and Improve Fairness: Sleeping + Queues

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  }
  if(this.flag == 0) {
    this.flag = 1; // lock is taken
    this.guard = 0;
  }
  else {
    this.q.push(t);
    setpark(); // Solaris primitive
    this.guard = 0;
    park(); // Solaris primitive
  }
}
```

```
public void release() {
  while(test&set(this.guard) == 1) {
    // while busy do nothing
  }
  if(q.is_empty()) {
    this.flag = 0; // lock is free
  }
  else {
    t = q.pop();
    unpark(t); // Solaris primitive
  }
  this.guard = 0;
}
```

Solution:

Use another system call to tell the OS the thread is *about* to park

Locks: Wrap Up

- Synchronization primitives ensure that only one process/thread at a time executes in a critical section (mutual exclusion)
- Locks allow protection of critical sections by atomically testing and taking/releasing the access to a critical section
- Locks can be implemented leveraging some HW support:
  - disabling interrupts (can miss or delay important events)
  - atomic instructions (busy waiting/spin-lock inefficient)

Higher-Level Synchronization Primitives

- More general synchronization mechanisms
  - Not only for safely accessing critical sections

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Higher-Level Synchronization Primitives

- More general synchronization mechanisms
  - Not only for safely accessing critical sections
- 2 common high-level synchronization primitives:
  - Semaphores: binary (mutex) and counting
  - Monitors: mutex and condition variables

Semaphores: Overview

- Another data structure that provides mutual exclusion to critical sections

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- Another data structure that provides mutual exclusion to critical sections
- Can also play the role of an atomic counter
- Generalization of locks invented by Dijkstra in 1965
- Special type of (integer) variable that supports 2 atomic operations
  - wait() (also P()): decrement, block until semaphore is open
  - signal() (also V()): increment, allow another thread to enter

Blocking in Semaphores

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  - If a thread is waiting on the queue the thread is unblocked, whilst if no threads are waiting on the queue, the signal is remembered for the next thread
- In other words, signal() is stateful and has "history"

Semaphores: Types

- Binary Semaphore a.k.a. Mutex (same as a Lock)
  - Guarantees mutually exclusive access to a resource (i.e., only one process/thread executes in a critical section)
  - Its associated integer variable can only take 2 values: 0/1
  - Initialized to open (e.g., value = 1)

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  - Initialized to open (e.g., value = 1)
- Counting Semaphore
  - To manage multiple shared resources
  - The semaphore is initially set to the number of resources
  - A process can access to a resource as long as at least one is available

Semaphores: Key Ideas

```
// Semaphore S

S.wait(); // wait until S is available

<critical section>

S.signal(); notify other processes that S is open
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If a process executes S.wait() and semaphore S is open (non-zero), it continues executing, otherwise the OS puts the process on the wait queue

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If a process executes S.wait() and semaphore S is open (non-zero), it continues executing, otherwise the OS puts the process on the wait queue

A S.signal() unblocks one process on semaphore S's wait queue

Binary Semaphore: Example

"Too Much Milk" Using Lock

```
# Thread Bob

lock.acquire()

if (!milk):
    buy_milk()

lock.release()
```

```
# Thread Carla

lock.acquire()

if (!milk):
    buy_milk()

lock.release()
```

Binary Semaphore: Example

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S.signal()
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S.signal()
```

Binary Semaphore: Example

"Too Much Milk" Using Lock		"Too Much Milk" Using Semaphore	
# Thread Bob	# Thread Carla	# Thread Bob	# Thread Carla
lock.acquire()	lock.acquire()	S.wait()	S.wait()

if (!milk):
    buy_milk()
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	    buy_milk()
		if (!milk):
		    buy_milk()
			if (!milk):
			    buy_milk()

lock.release()	lock.release()	S.signal()	S.signal()

Semaphore: Implementation

```
Class Semaphore {
  public void wait(Thread t);
  public void signal();
  private int value;
  private int guard;
  private Queue q;

  Semaphore(int val) {
    // initialize semaphore
    // with val and empty queue
    this.value = val;
    this.q = new Queue();
  }
}
```

```
public void wait(Thread t) {
  while(test&set(this.guard) == 1) {
    // while busy do nothing
  }
  this.value -= 1;
  if(this.value < 0) {
    q.push(t);
    park();
  }
  else {
    this.guard = 0;
  }
}
```

```
public void signal() {
  while(test&set(this.guard) == 1) {
    // while busy do nothing
  }
  this.value += 1;
  if(!q.isEmpty()) { // this.value < 0
    t = q.pop();
    unpark(t);
  }
  this.guard = 0;
}
```

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Class Semaphore {
  public void wait(Thread t);
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  Semaphore(int val) {
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public void wait(Thread t) {
  while(test&set(this.guard) == 1) {
    // while busy do nothing
  }
  this.value -= 1;
  if(this.value < 0) {
    q.push(t);
    park();
  }
  else {
    this.guard = 0;
  }
}
```

```
public void signal() {
  while(test&set(this.guard) == 1) {
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  }
  this.value += 1;
  if(!q.isEmpty()) { // this.value < 0
    t = q.pop();
    unpark(t);
  }
  this.guard = 0;
}
```

wait() and signal() are of course atomic!

Semaphore: Example

S = 2

A

```
S.wait()
S.wait()
S.signal()
S.signal()
```

B

```
S.wait()
S.signal()
```

Semaphore: Example

S = 2

A

```
S.wait()
S.wait()
S.signal()
S.signal()
```

B

```
S.wait()
S.signal()
```

A possible execution flow

S (value)	Queue	A	B
2	∅	ready to exec	ready to exec

Semaphore: Example

A

```
S.wait()
S.wait()
S.signal()
S.signal()
```

B

```
S.wait()
S.signal()
```

A possible execution flow

	S (value)	Queue	A	B
	2	∅	ready to exec	ready to exec
A: S.wait()				

Semaphore: Example

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```
S.wait()
S.wait()
S.signal()
S.signal()
```

B

```
S.wait()
S.signal()
```

A possible execution flow

	S (value)	Queue	A	B
	2	Ø	ready to exec	ready to exec
A: S.wait()	1	Ø	ready to exec	ready to exec

Semaphore: Example

A
```
S.wait()
S.wait()
S.signal()
S.signal()
```

B
```
S.wait()
S.signal()
```

A possible execution flow

	S (value)	Queue	A	B
	2	∅	ready to exec	ready to exec
A: S.wait()	1	∅	ready to exec	ready to exec
B: S.wait()				

Semaphore: Example

A

```
S.wait()
S.wait()
S.signal()
S.signal()
```

B

```
S.wait()
S.signal()
```

A possible execution flow

	S (value)	Queue	A	B
	2	∅	ready to exec	ready to exec
A: S.wait()	1	∅	ready to exec	ready to exec
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Semaphore: Example

A

```
S.wait()
S.wait()
S.signal()
S.signal()
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```
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S.signal()
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A possible execution flow

	S (value)	Queue	A	B
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Semaphore: Example

A

```
S.wait()
S.wait()
S.signal()
S.signal()
```

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```
S.wait()
S.signal()
```

A possible execution flow

	S (value)	Queue	A	B
	2	∅	ready to exec	ready to exec
A: S.wait()	1	∅	ready to exec	ready to exec
B: S.wait()	0	∅	ready to exec	ready to exec
A: S.wait()	-1	A	blocked	ready to exec

Semaphore: Example

A

```
S.wait()
S.wait()
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Semaphore: Example

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S.signal()

B

S.wait()
S.signal()

A possible execution flow

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A: S.wait()	1	∅	ready to exec	ready to exec
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A: S.wait()	-1	A	blocked	ready to exec
B: S.signal()	0	∅	ready to exec	ready to exec
A: S.signal()	1	∅	ready to exec	ready to exec
A: S.signal()	2	∅	ready to exec	ready to exec

Semaphores: Purposes

- Mutual Exclusion: used to guard critical sections
  - The initial value of the semaphore is set to 1
  - Call wait() before the critical section, signal() after the critical section

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  - Example → join() or waitpid()

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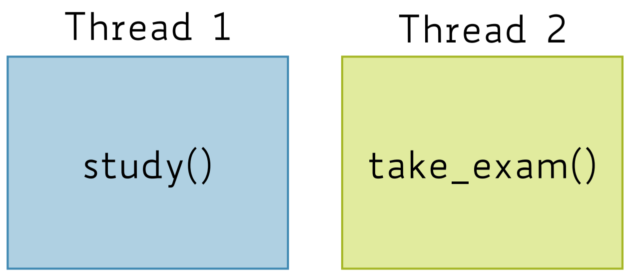

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  - Call wait() before the critical section, signal() after the critical section
- Scheduling Constraints: used to enforce threads to wait
  - The initial value of the semaphore is set to 0
  - Example → join() or waitpid()

S = 0

Thread 1	Thread 2
study() S.signal()	S.wait() take_exam()

Producer-Consumer

Producer Process:

```
while (true)
{
      /* produce an item in nextProduced */
      while (counter == BUFFER_SIZE)
         ; /* do nothing */
      buffer[in] = nextProduced;
      in = (in + 1) % BUFFER_SIZE;
      counter++;
}
```

Consumer Process:

```
while (true)
{
      while (counter == 0)
         ; /* do nothing */
      nextConsumed = buffer[out];
      out = (out + 1) % BUFFER_SIZE;
      counter--;
      /* consume the item in nextConsumed */
}
```

Both the producer and the consumer share a common buffer (of items)

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Both the producer and the consumer share a common buffer (of items)

counter keeps track of the number of items currently in the buffer

possible race condition as counter can be updated by the producer and consumer

Producer-Consumer: Race Condition

Producer:

$$register_1 = \texttt{counter}$$
$$register_1 = register_1 + 1$$
$$\texttt{counter} = register_1$$

Consumer:

$$register_2 = \texttt{counter}$$
$$register_2 = register_2 - 1$$
$$\texttt{counter} = register_2$$

Interleaving:

Assuming the initial value of counter is 5

T_0:	*producer*	execute	$register_1 = \texttt{counter}$	$\{register_1 = 5\}$
T_1:	*producer*	execute	$register_1 = register_1 + 1$	$\{register_1 = 6\}$
T_2:	*consumer*	execute	$register_2 = \texttt{counter}$	$\{register_2 = 5\}$
T_3:	*consumer*	execute	$register_2 = register_2 - 1$	$\{register_2 = 4\}$
T_4:	*producer*	execute	$\texttt{counter} = register_1$	$\{counter = 6\}$
T_5:	*consumer*	execute	$\texttt{counter} = register_2$	$\{counter = 4\}$

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T_4:	*producer*	execute	$\texttt{counter} = register_1$	$\{counter = 6\}$
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Q1: What would be the resulting value of counter if the order of statements T4 and T5 were reversed?

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$$register_1 = register_1 + 1$$
$$\texttt{counter} = register_1$$

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Q2: What should the value of counter be after one producer and one consumer, assuming the original value was 5?

Producer-Consumer: Desiderata

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  - Access to the shared buffer of items must be granted to a single thread at a time (either the producer or the consumer)

Producer-Consumer: Desiderata

- Mutual Exclusion
  - Access to the shared buffer of items must be granted to a single thread at a time (either the producer or the consumer)
- Scheduling Constraints
  - Producer can put a new item iff the buffer is not full
  - Consumer can take an item iff the buffer is not empty

Producer-Consumer in Java

Semaphores: Wrap Up

- Generalization of locks
- Can be used for 3 purposes:
  - To ensure mutually exclusive execution of a critical section as locks do (binary semaphore)
  - To control access to a shared pool of resources (counting semaphore)
  - To enforce scheduling constraints so as to execute threads according to some specific order

What's Wrong with Semaphores?

- Not easy to get the meaning of waiting/signaling on a semaphore
- They are essentially shared global variables
- There is no direct connection between the semaphore and the data which the semaphore controls access to
- They serve multiple purposes (e.g., mutex, scheduling constraints, etc.)
- Their correctness depends on the programmer's ability

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Solution: Use a higher level primitive called monitors

What is a Monitor?

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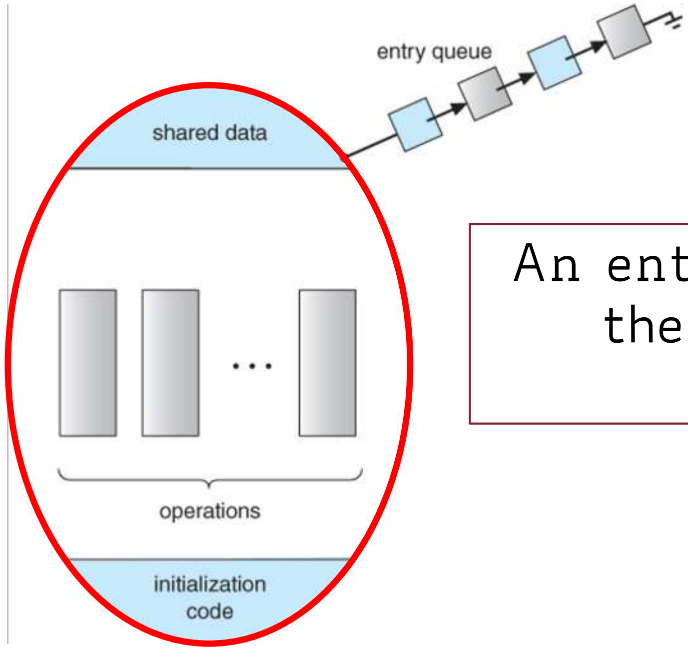
- A monitor is a programming language construct that controls access to shared data
- Similar to a (Java/C++) class that embodies all together: data, operations, and synchronization
- Synchronization code added by compiler, enforced at runtime

What is a Monitor?

- Unlike classes, monitors:
  - guarantee mutual exclusion, i.e., only one thread may execute a monitor's method at a time
  - require all data to be private

Monitor: A Schematic Overview



An entry queue of processes waiting their turn to execute monitor operations (methods)

Monitor: A Formal Definition

- Defines a lock and zero or more condition variables for managing concurrent access to shared data

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- Defines a lock and zero or more condition variables for managing concurrent access to shared data
- Uses the lock to ensure that only a single thread is active within the monitor at any time
- The lock provides of course mutual exclusion for shared data

Monitor: Java Implementation Example

- It is straightforward to turn a Java class into a monitor by just:
  - Making all the data private
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Monitor: Java Implementation Example

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  - Making all the data private
  - Making all methods (or non-private ones) synchronized
- The synchronized keyword indicates the method is subject to mutual exclusion

Monitor: Java Implementation Example

```
class Queue {
    …
    private ArrayList<Item> data;
    …

    public void synchronized add(Item i) {
        data.add(i);
    }

    public Item synchronized remove() {
        if (!data.isEmpty()) {
            Item i = data.remove(0);
            return i;
        }
    }
}
```

Condition Variables

- In the previous example, the remove() method should wait until something is available on the queue

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  - But if the thread sleeps while still holding a lock then no other threads can access the queue, add an item to it, and eventually wake up the sleeping thread

Condition Variables

- In the previous example, the remove() method should wait until something is available on the queue
  - Intuitively, the thread should sleep inside the critical section
  - But if the thread sleeps while still holding a lock then no other threads can access the queue, add an item to it, and eventually wake up the sleeping thread
  - Deadlock (more on this later…)

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Condition Variables

- Solution: condition variables
  - Conceptually a queue of threads, associated with a lock, on which a thread may wait for some condition to become true
  - Enable a thread to sleep within a critical section
  - Any lock held by the thread is atomically released before going to sleep

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- Each condition variable supports 3 operations:
  - wait → release lock and go to sleep atomically (queue of waiters)
  - signal → wake up a waiting thread if one exists, otherwise it does nothing
  - broadcast → wake up all waiting threads
- Rule: thread must hold the lock when doing condition variable operations
- Note: condition variables are not boolean objects!

Condition Variables in Java

- Use `wait()` to give up the lock *and* sleep **atomically**

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- Use `notify()` to signal that the condition a thread is waiting on is satisfied

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Condition Variables in Java

- Use `wait()` to give up the lock *and* sleep **atomically**
- Use `notify()` to signal that the condition a thread is waiting on is satisfied
- Use `notifyAll()` to wake up all waiting threads
- Concretely, one condition variable per object

Monitor: Java Implementation Example

```
class Queue {
    …
    private ArrayList<Item> data;
    …

    public void synchronized add(Item i) {
        data.add(i);
        notify();
    }

    public Item synchronized remove() {
        while (data.isEmpty()) {
            wait(); // give up the lock and sleep
        }
        Item i = data.remove(0);
        return i;
        }
    }
}
```

Condition Variables != Semaphores

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- Same operations yet entirely different semantics
- Access to the monitor is controlled by a lock
- wait() blocks the calling thread, and gives up the lock
  - to call wait(), the thread has to be in the monitor (hence, it has the lock!)
  - on a semaphore, wait() just blocks the thread on the queue
- signal() causes a waiting thread to wake up
  - If there is no waiting thread, the signal is lost though!
  - on a semaphore, signal increases the counter, allowing future entry even if no thread is currently waiting

signal(): Mesa- vs. Hoare-style

- Mesa-style (Nachos, Java, and most real OSs)
  - The signaling thread places a waiter on the ready queue, but signaler continues inside monitor
  - Condition is not necessarily true when waiter runs again
  - Returning from wait() is only a hint that something changed
  - Must re-check the conditional case

signal(): Mesa- vs. Hoare-style

- Hoare-style (most textbooks)
  - The signaling thread immediately switches to a waiting thread
  - The condition that the waiter was anticipating is guaranteed to hold when waiter executes

Mesa vs. Hoare Monitors

- Mesa-style

```
while (empty) {
    wait(condition);
}
```

Mesa vs. Hoare Monitors

- Mesa-style

```
while (empty) {
    wait(condition);
}
```

- Hoare-style

```
if (empty) {
    wait(condition);
}
```

Mesa vs. Hoare Monitors

- Mesa-style

```
while (empty) {
    wait(condition);
}
```

Easier to use and more efficient

- Hoare-style

```
if (empty) {
    wait(condition);
}
```

Easier to reason about the program's behaviour

Mesa vs. Hoare

Mesa

```
class Queue {
    ...
    private ArrayList<Item> data;
    ...

    public void synchronized add(Item i) {
        data.add(i);
        notify();
    }

    public Item synchronized remove() {
        while (data.isEmpty()) {
            wait(); // give up the lock and sleep
        }
        Item i = data.remove(0);
        return i;
        }
    }
}
```

The waiting thread may need to wait again after it is awakened, because some other thread could grab the lock and remove the item before it gets to run

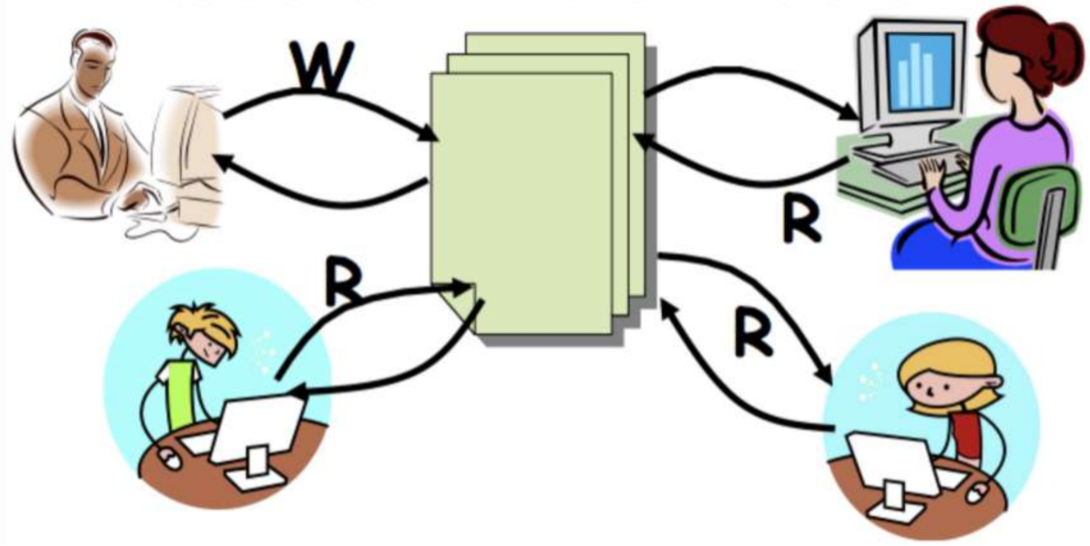
Hoare

```
class Queue {
    ...
    private ArrayList<Item> data;
    ...

    public void synchronized add(Item i) {
        data.add(i);
        notify();
    }

    public Item synchronized remove() {
        if (data.isEmpty()) {
            wait(); // give up the lock and sleep
        }
        Item i = data.remove(0);
        return i;
        }
    }
}
```

The waiting thread runs immediately after an item is added to the queue

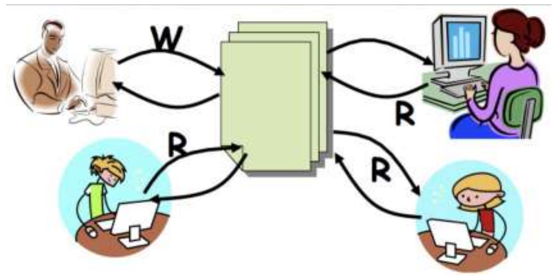
Readers-Writers Problem

Motivation: Consider a shared database system
(more generally, any shared resource)



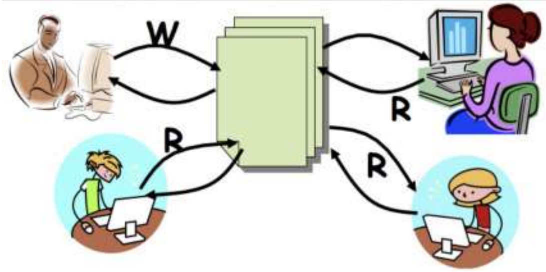
Readers-Writers Problem

Two classes of users:
- Readers → never modify the DB



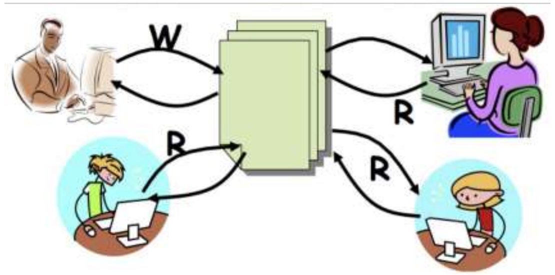
Readers-Writers Problem

Two classes of users:
- Readers → never modify the DB
- Writers → read and modify the DB



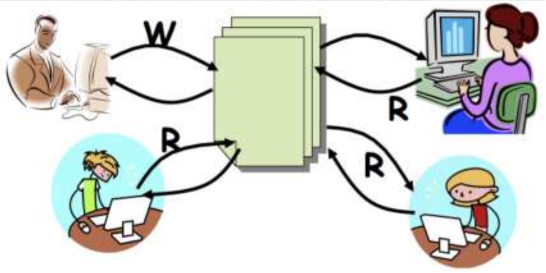
Readers-Writers Problem

Simplest solution:
- Use a single lock on the data object for each operation



Readers-Writers Problem

Simplest solution:
- Use a single lock on the data object for each operation
- May be too restrictive!



Readers-Writers Problem

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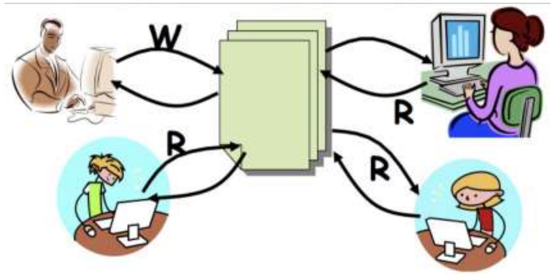
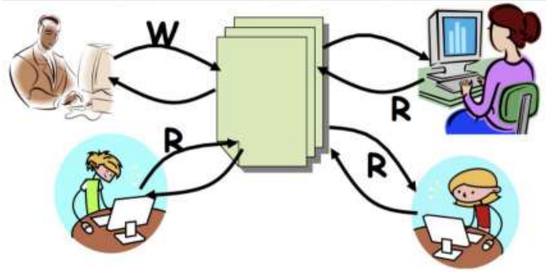


Only one writer at a time but, possibly, multiple readers

Readers-Writers Problem

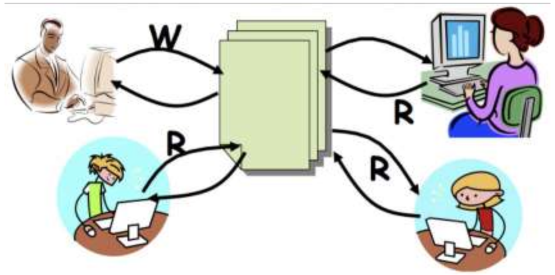
<u>Constraints:</u>

- Readers can access DB when no writers

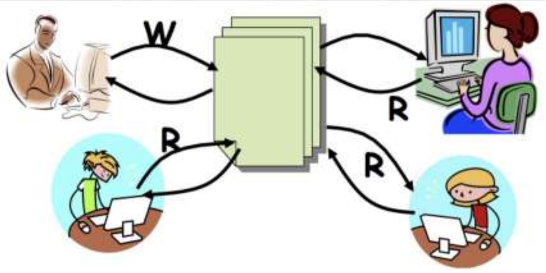


Readers-Writers Problem

<u>Constraints:</u>
- Readers can access DB when no writers
- Writers can access DB when no readers or writers



Readers-Writers Problem



Constraints:
- Readers can access DB when no writers
- Writers can access DB when no readers or writers
- Only one thread manipulates state variables at a time

Readers-Writers Problem

- 2 variations of the problem depending on whether priority is on readers or writers:

Readers-Writers Problem

- 2 variations of the problem depending on whether priority is on readers or writers:
  - first readers-writers problem (priority to the readers)

Readers-Writers Problem

- 2 variations of the problem depending on whether priority is on readers or writers:
  - first readers-writers problem (priority to the readers)
  - second readers-writers problem (priority to the writers)

First Readers-Writers Problem

- Priority to the readers
- If a reader wants access to the data, and there is not already a writer accessing it, then access is granted to the reader
- Possible starvation of the writers, as there could always be more readers coming along to access the data

Second Readers-Writers Problem

- Piority to the writers
- When a writer wants access to the data it jumps to the head of the queue
- Possible starvation of the readers, as they are all blocked as long as there are writers

Readers-Writers in Java Using Lock

Readers-Writers in Java Using Monitors

Summary

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  - Readers-Writers