# Systems and Networking I

Applied Computer Science and Artificial Intelligence 2023–2024



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File System API

OS Implementation

Physical Implementation

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File creation, manipulation, protection, etc.

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OS internal data structures and algorithms

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Second storage structure, disk scheduling algorithms

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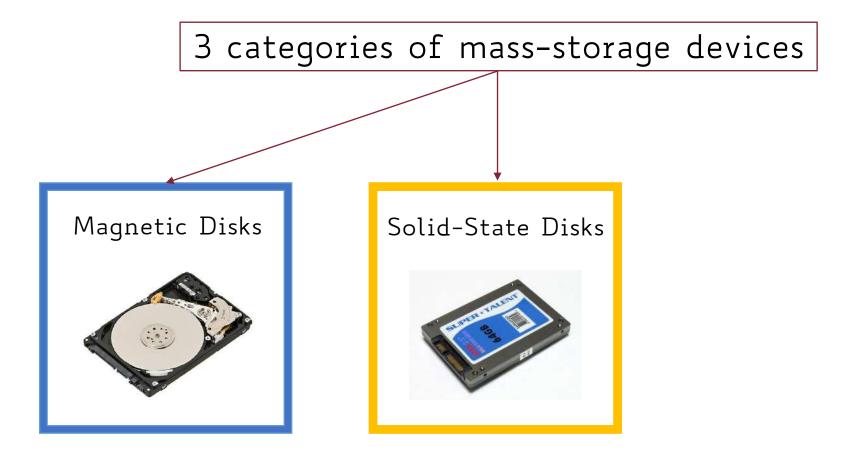
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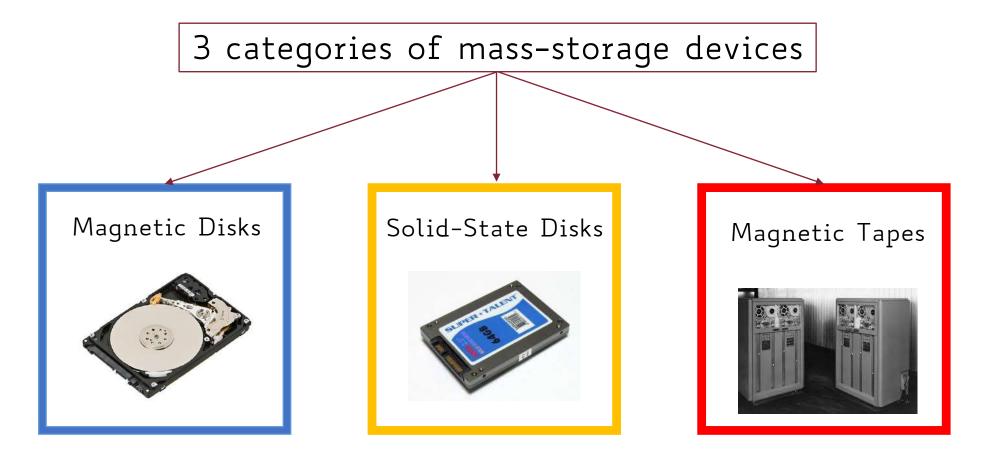
# Part V: Storage Management

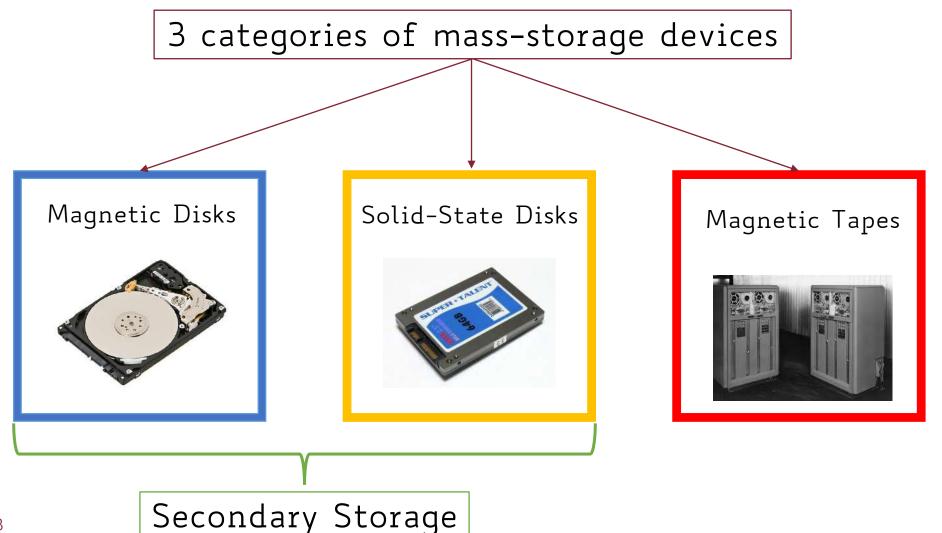
3 categories of mass-storage devices

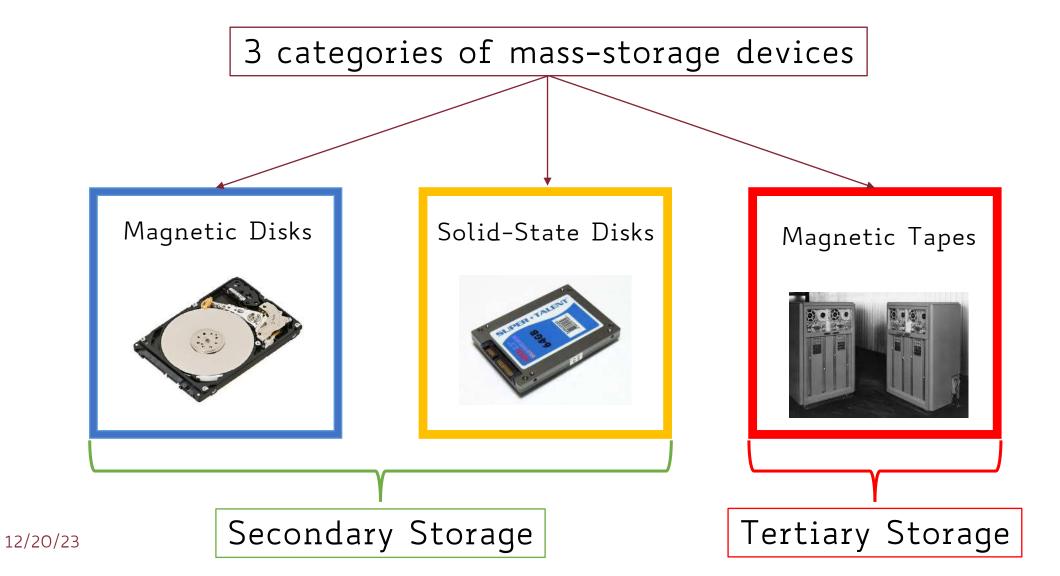
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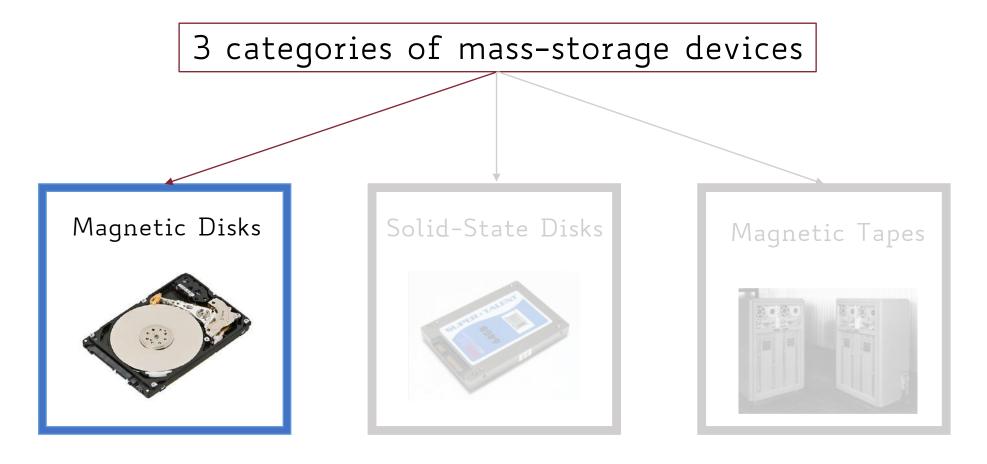


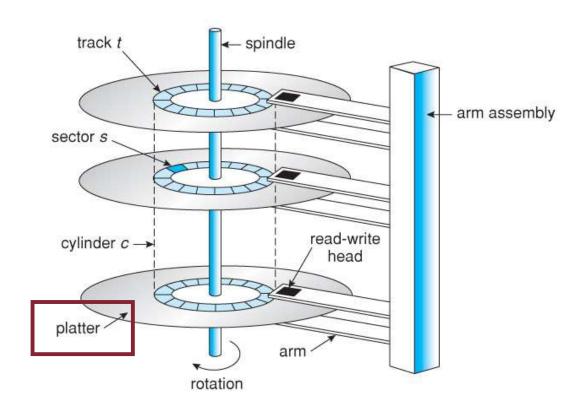




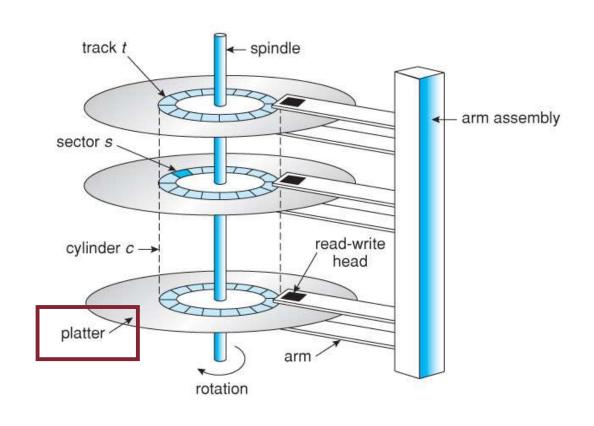


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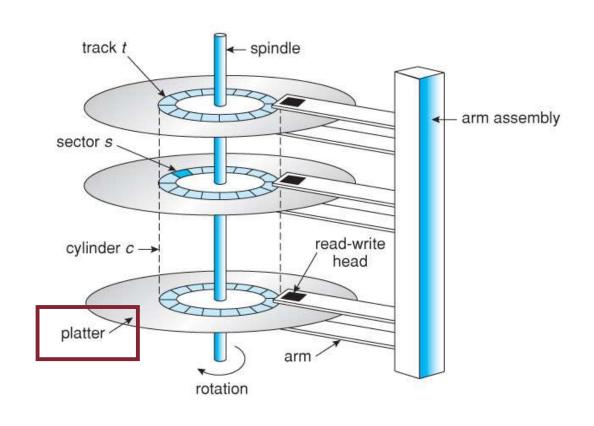


One or more platters covered with magnetic media



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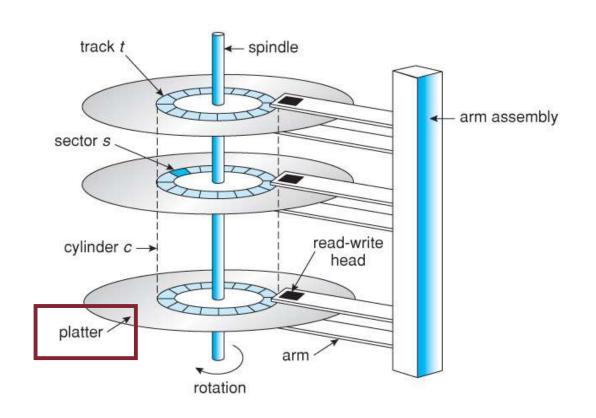
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Floppy disk flexible plastic



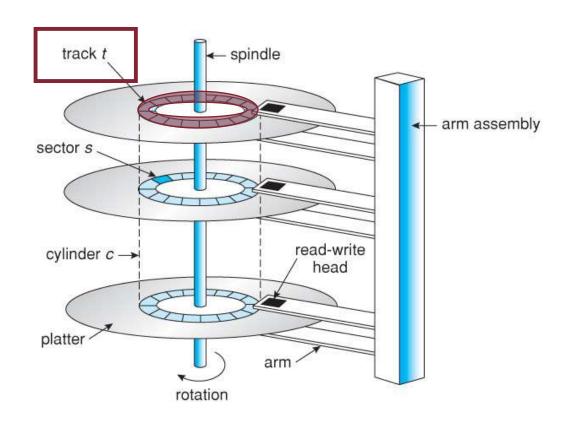
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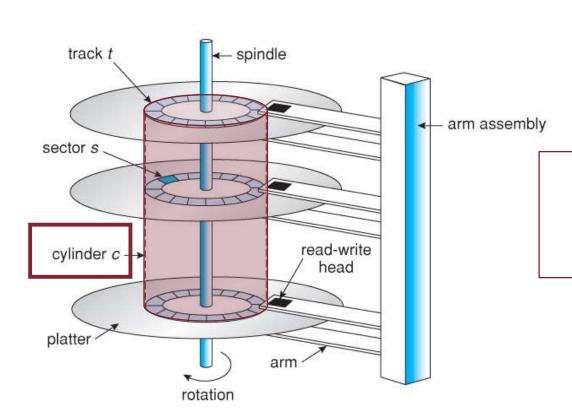
Each platter has 2 working surfaces

# Magnetic Disks: Tracks and Cylinders



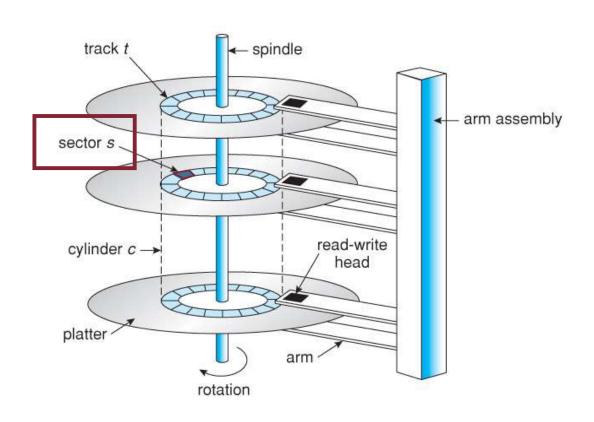
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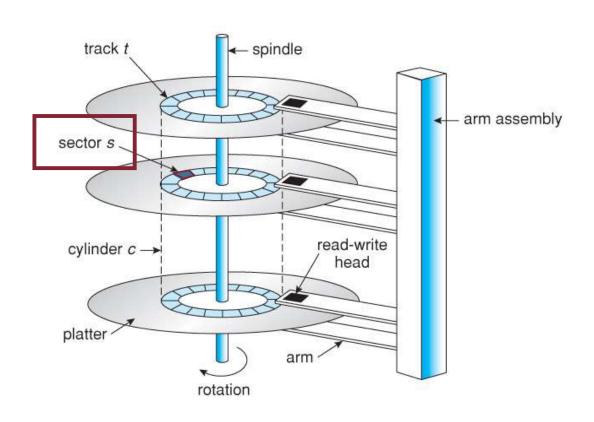


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The set of all tracks that are the same distance from the edge of the platter is called a cylinder

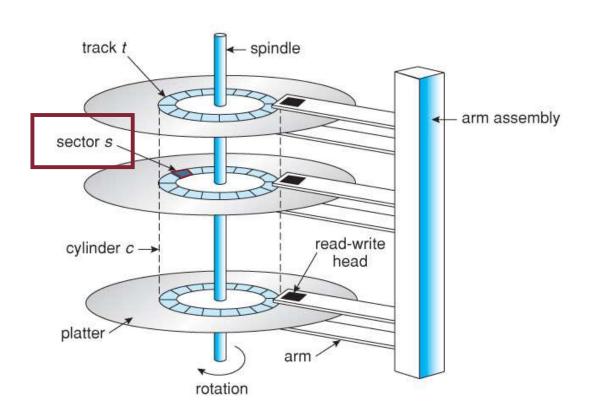


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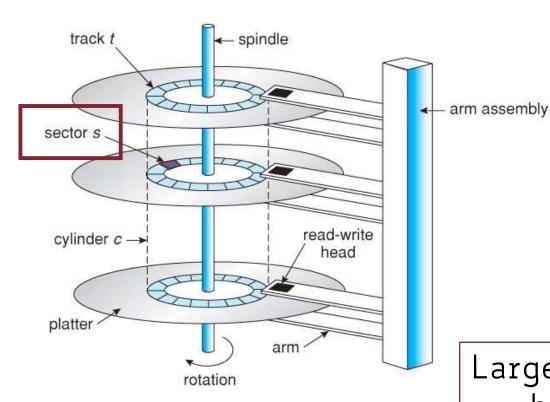
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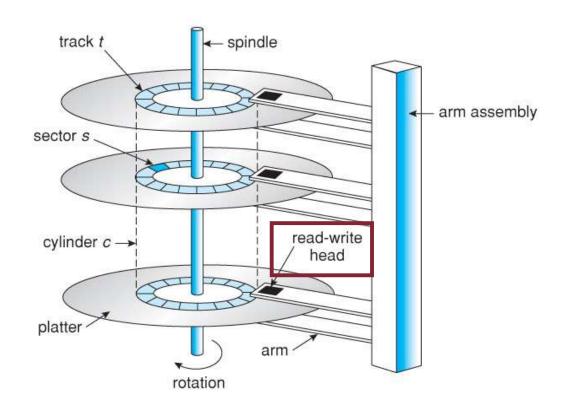


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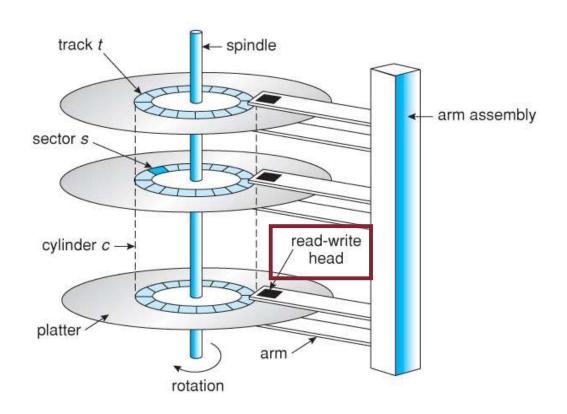
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Larger sector sizes reduce the space wasted by headers and trailers, but increase internal fragmentation

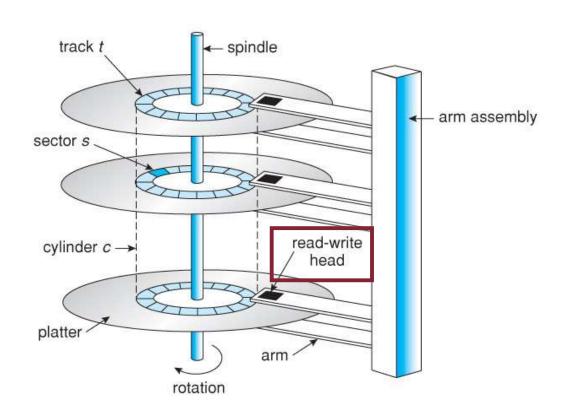


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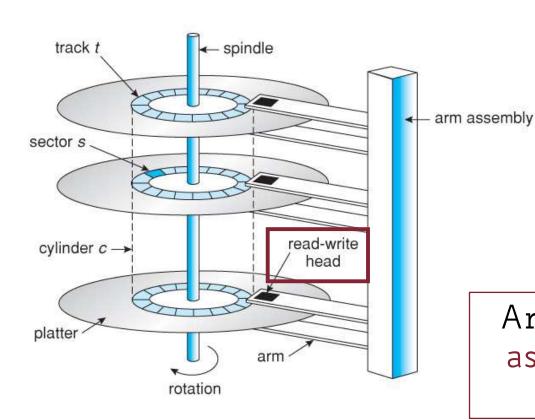
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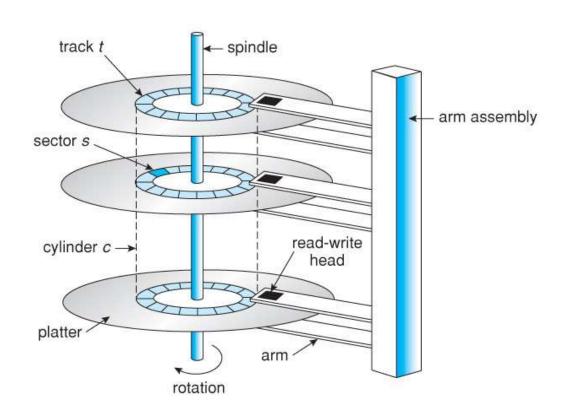


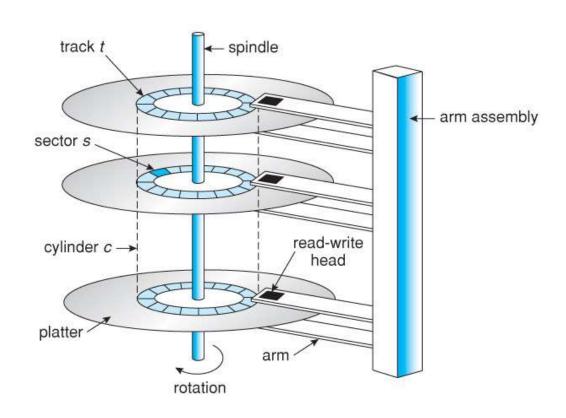
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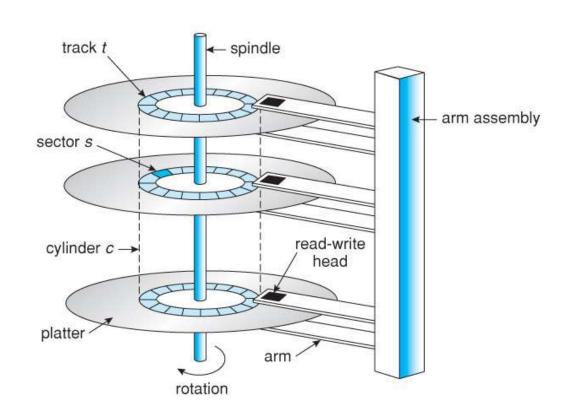
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Arms are controlled by a common arm assembly moving simultaneously from one cylinder to another



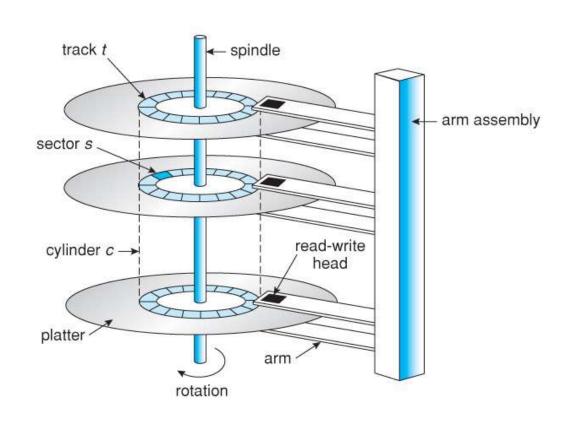


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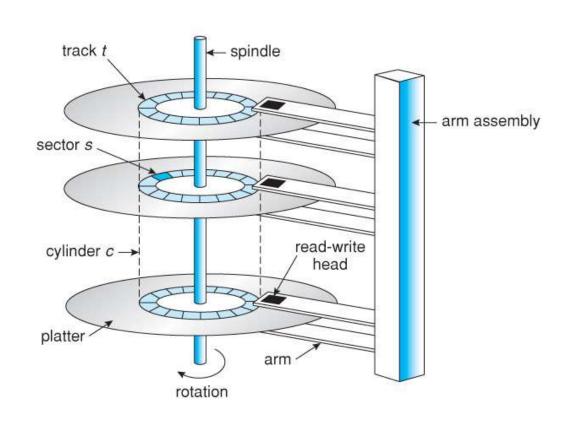
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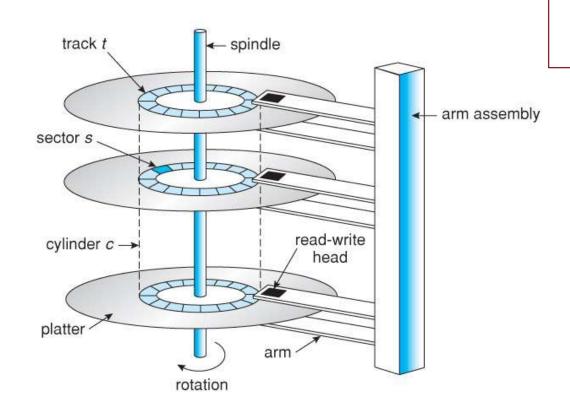
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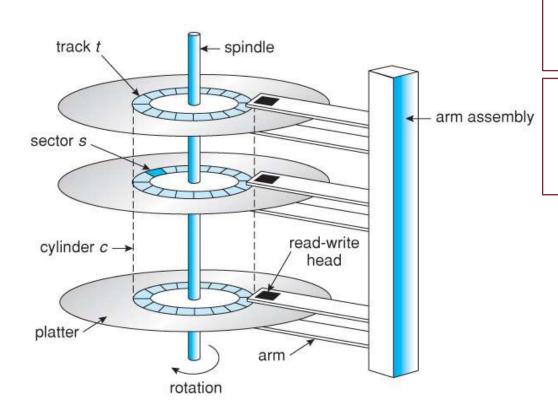
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C = H \* T \* S \* B

OVERALL CAPACITY

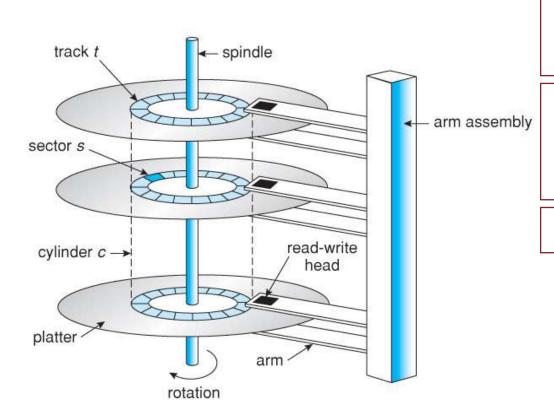


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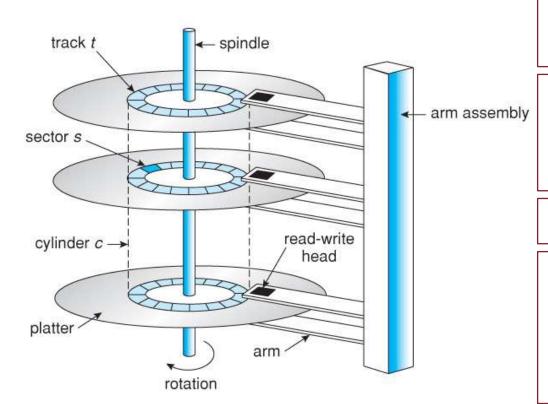
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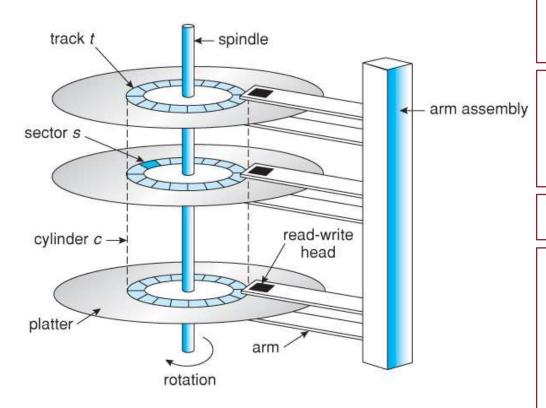
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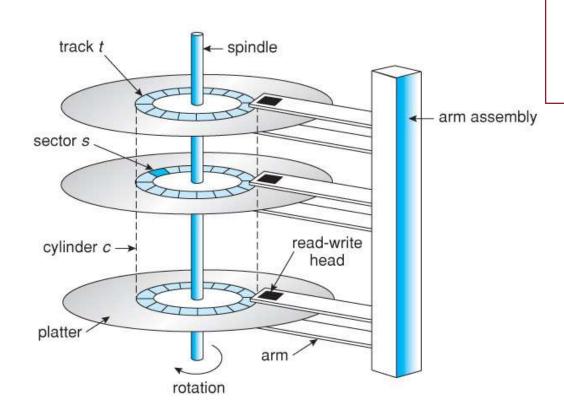
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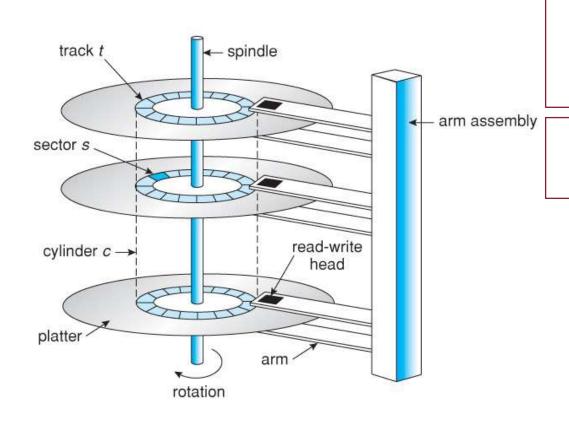
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#### Drawbacks:

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- Different frequencies and timing from innermost to outermost tracks

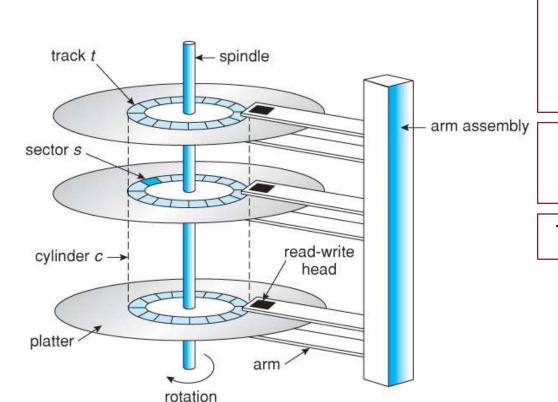


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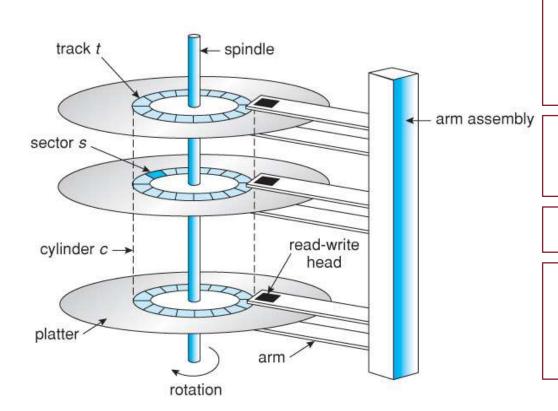
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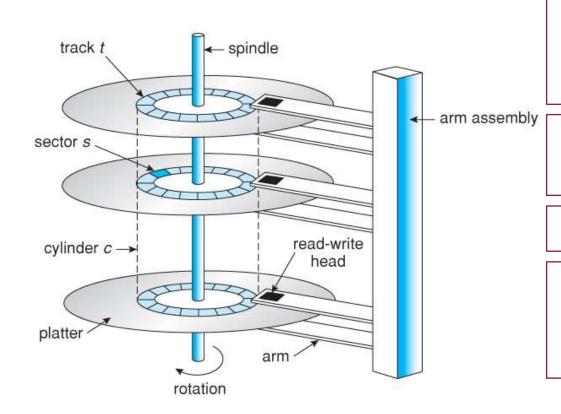


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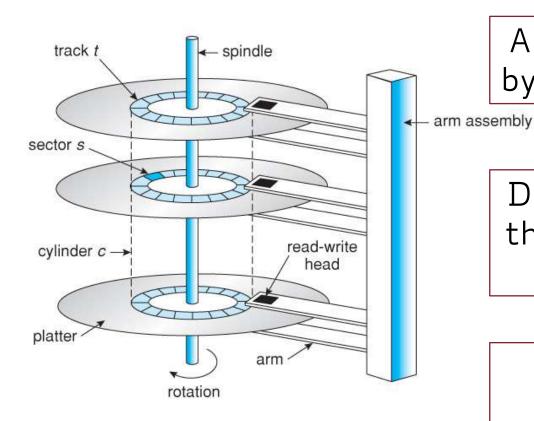
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Zone Bit Recording (ZBR)

# Magnetic Disks: (Logical) Referencing



A physical block of data is specified by the (head, cylinder, sector) number

Disk blocks are numbered starting at the outermost cylinder, identified by O

Note that cylinder coincides with track

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   7200 rpm = 120 rps)
- Outer tracks spin faster than inner tracks (more sectors traversed in the same amount of time due to larger radius → more sectors per zone in ZBR)

- Data transfer from the disk to memory is made of 3 steps:
  - positioning time (seek time or random access time)
  - rotational delay
  - transfer time

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- Typically, the slowest step in the entire process

Bottleneck of overall disk data transfer

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- On average, O.5 revolutions (r)
  - E.g., for a 7200 rpm (120 rps) disk this equals to 0.5 r/120 rps ~4 msec

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Data Transfer Time = Seek Time + Rotational Delay + Transfer Time

Sometimes the term transfer rate is used to refer to the overall data transfer time

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- Then through other cylinders (from the outermost to innermost)

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- Head crash may permanently damage the disk or even destroy it
- To avoid such a risk, disk heads are "parked" when the computer is turned off

# Magnetic Disks: Interfaces

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- Some of the common interface formats include:
  - Enhanced Integrated Drive Electronics (EIDE);
  - Advanced Technology Attachment (ATA) and Serial ATA (SATA);
  - Universal Serial Bus (USB);
  - Fiber Channel (FC);
  - Small Computer Systems Interface (SCSI)

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- Finally, data is transferred from that cache to the host controller and the motherboard memory at electronic speeds

- Mechanical components of magnetic disks cause bottleneck
  - Seek Time
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  - Seek Time
  - Rotational Delay
- To minimize data transfer time from disk we need to minimize those

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- Fast-spinning disks → lower rotational delay

Hardware Optimization

- How can the OS help minimize data transfer time?
- Schedule disk operations so as to minimize head movement
- Lay out data on disk so that related data are located on close tracks
- Place commonly-used data on a specific portion of the disk
- Pick carefully the block size contained on each sector:
  - Too small → more seeks are needed to transfer the same amount of data
  - Too large → more internal fragmentation and space wasted

# Summary

- Disks are slow devices compared to CPUs (and main memory)
- Manage those device efficiently is crucial
- Minimize seek and rotational delay on magnetic disks
- HW optimizations are limited → OS needs to take the lead here!