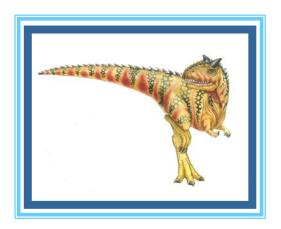
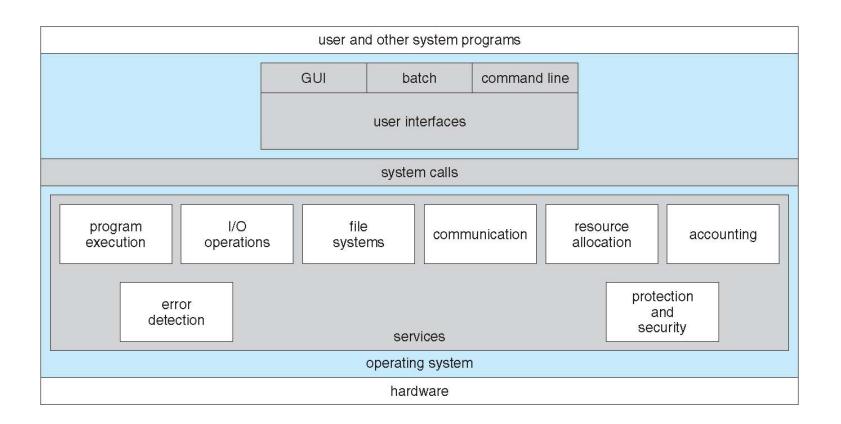
# Chapter 2: Operating-System Structures

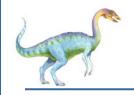




# **A View of Operating System Services**





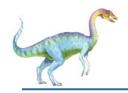


### **System Calls**

- Programming interface to the services provided by the OS
- Typically written in a high-level language (C or C++)
- Mostly accessed by programs via a high-level Application Programming Interface (API) rather than direct system call use
- Three most common APIs are Win32 API for Windows, POSIX API for POSIX-based systems (including virtually all versions of UNIX, Linux, and Mac OS X), and Java API for the Java virtual machine (JVM)

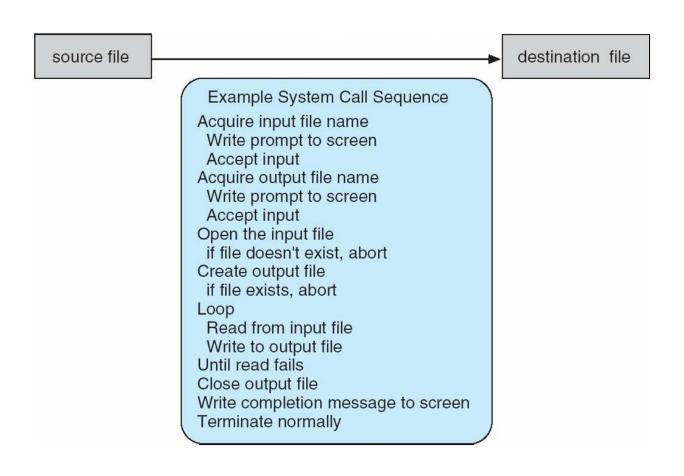
Note that the system-call names used throughout this text are generic





#### **Example of System Calls**

System call sequence to copy the contents of one file to another file







#### **Example of Standard API**

#### EXAMPLE OF STANDARD API

As an example of a standard API, consider the read() function that is available in UNIX and Linux systems. The API for this function is obtained from the man page by invoking the command

#### man read

on the command line. A description of this API appears below:

```
#include <unistd.h>

ssize_t read(int fd, void *buf, size_t count)

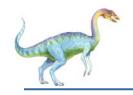
return function parameters
value name
```

A program that uses the read() function must include the unistd.h header file, as this file defines the ssize\_t and size\_t data types (among other things). The parameters passed to read() are as follows:

- int fd—the file descriptor to be read
- void \*buf —a buffer where the data will be read into
- size\_t count—the maximum number of bytes to be read into the buffer

On a successful read, the number of bytes read is returned. A return value of 0 indicates end of file. If an error occurs, read() returns -1.





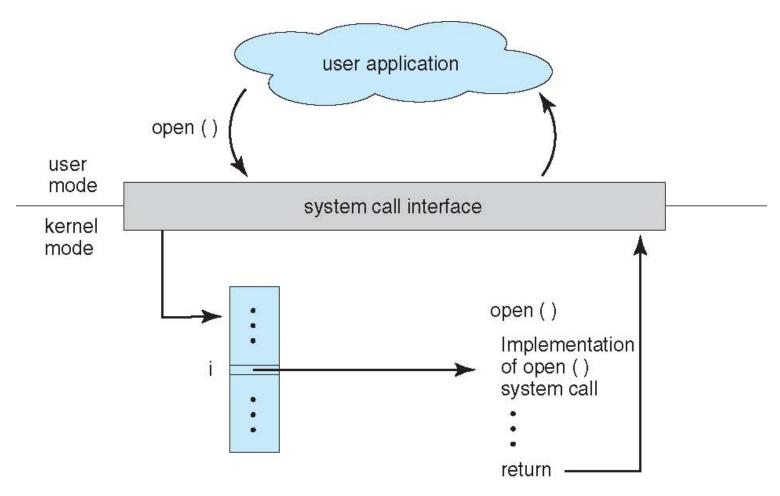
#### **System Call Implementation**

- Typically, a number associated with each system call
  - System-call interface maintains a table indexed according to these numbers
- The system call interface invokes the intended system call in OS kernel and returns status of the system call and any return values
- The caller need know nothing about how the system call is implemented
  - Just needs to obey API and understand what OS will do as a result call
  - Most details of OS interface hidden from programmer by API
    - Managed by run-time support library (set of functions built into libraries included with compiler)





## API – System Call – OS Relationship



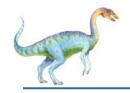




#### **Types of System Calls**

- Process control
  - create process, terminate process
  - end, abort
  - load, execute
  - get process attributes, set process attributes
  - wait for time
  - wait event, signal event
  - allocate and free memory
  - Dump memory if error
  - Debugger for determining bugs, single step execution
  - Locks for managing access to shared data between processes

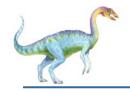




#### **Types of System Calls**

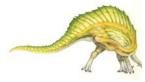
- File management
  - create file, delete file
  - open, close file
  - read, write, reposition
  - get and set file attributes
- Device management
  - request device, release device
  - read, write, reposition
  - get device attributes, set device attributes
  - logically attach or detach devices

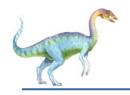




### **Types of System Calls (Cont.)**

- Information maintenance
  - get time or date, set time or date
  - get system data, set system data
  - get and set process, file, or device attributes
- Communications
  - create, delete communication connection
  - send, receive messages if message passing model to host name or process name
    - From client to server
  - Shared-memory model create and gain access to memory regions
  - transfer status information
  - attach and detach remote devices

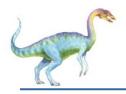




## **Types of System Calls (Cont.)**

- Protection
  - Control access to resources
  - Get and set permissions
  - Allow and deny user access





#### **Examples of Windows and Unix System Calls**

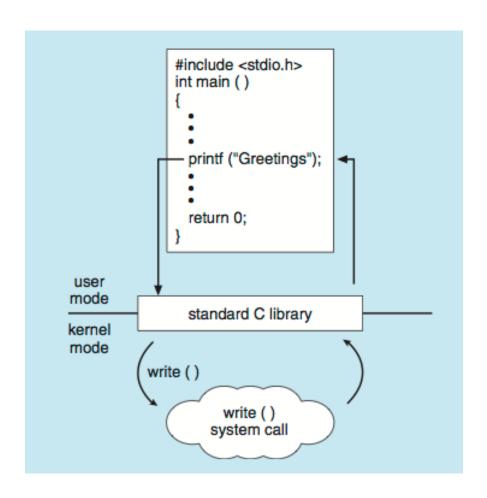
	Windows	Unix
Process Control	<pre>CreateProcess() ExitProcess() WaitForSingleObject()</pre>	<pre>fork() exit() wait()</pre>
File Manipulation	<pre>CreateFile() ReadFile() WriteFile() CloseHandle()</pre>	<pre>open() read() write() close()</pre>
Device Manipulation	SetConsoleMode() ReadConsole() WriteConsole()	ioctl() read() write()
Information Maintenance	<pre>GetCurrentProcessID() SetTimer() Sleep()</pre>	<pre>getpid() alarm() sleep()</pre>
Communication	<pre>CreatePipe() CreateFileMapping() MapViewOfFile()</pre>	<pre>pipe() shmget() mmap()</pre>
Protection	<pre>SetFileSecurity() InitlializeSecurityDescriptor() SetSecurityDescriptorGroup()</pre>	<pre>chmod() umask() chown()</pre>



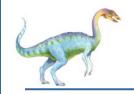


#### **Standard C Library Example**

C program invoking printf() library call, which calls write() system call







#### **Example: MS-DOS**

- Single-tasking
- Shell invoked when system booted
- Simple method to run program
  - No process created
- Single memory space
- Loads program into memory, overwriting all but the kernel
- Program exit -> shell reloaded

free memory command interpreter kernel (a)

At system startup

process

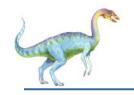
command interpreter

kernel

(b)

running a program





#### **Example: FreeBSD**

- Unix variant
- Multitasking
- User login -> invoke user's choice of shell
- Shell executes fork() system call to create process
  - Executes exec() to load program into process
  - Shell waits for process to terminate or continues with user commands
- Process exits with:
  - $\nearrow$  code = 0 no error

process D

free memory

process C

interpreter

process B

kernel



# **End of Chapter 2**

