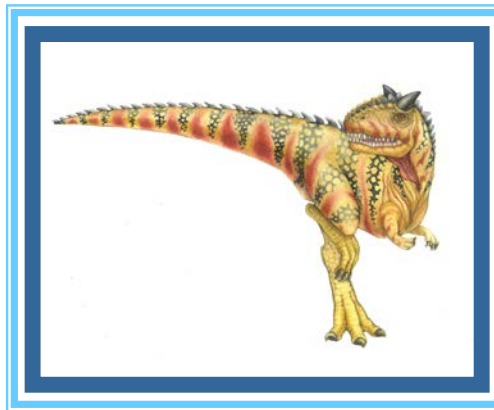


# Chapter 1: Introduction





# What is an Operating System?

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- 👁️ A program that acts as an intermediary between a user of a computer and the computer hardware
- 👁️ Operating system goals:
  - 🍃 Execute user programs and make solving user problems easier
  - 🍃 Make the computer system convenient to use
  - 🍃 Use the computer hardware in an efficient manner





# Computer System Structure

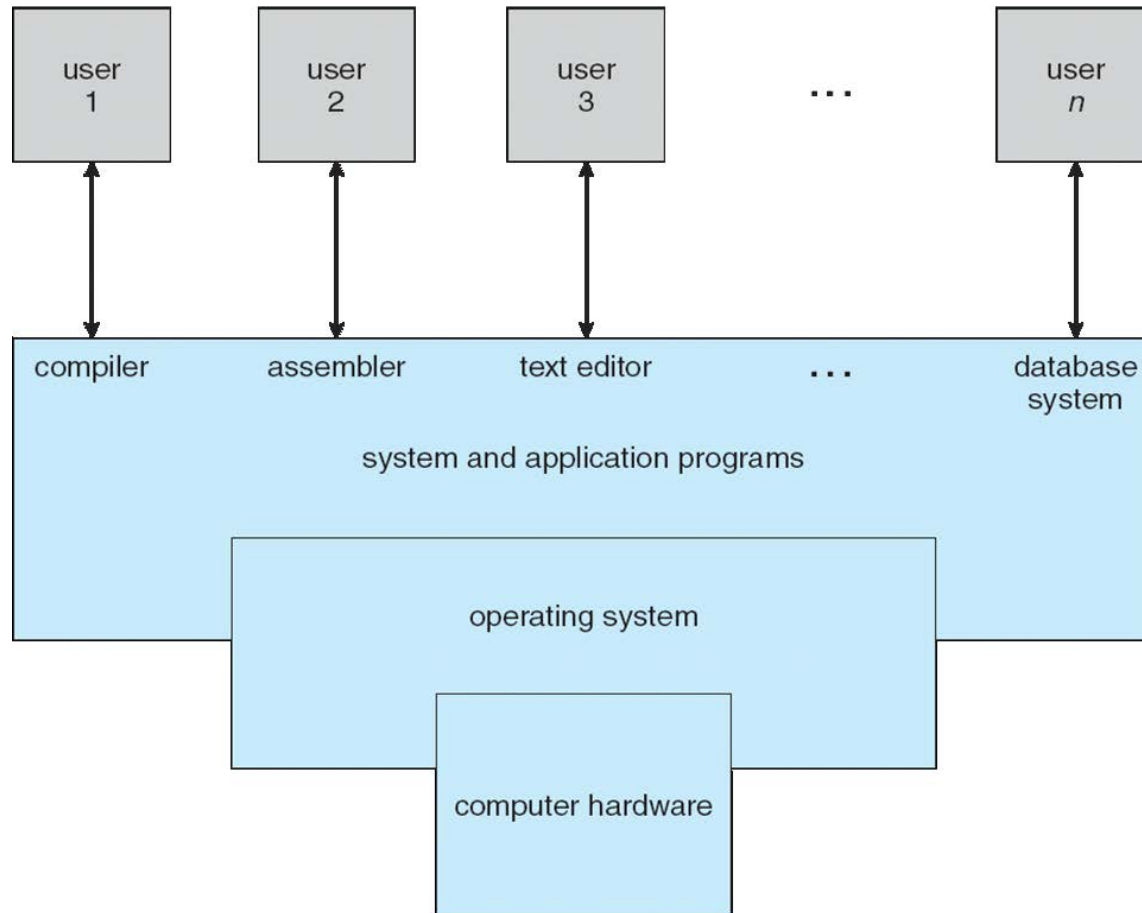
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- 👁 Computer system can be divided into four components:
  - 🍃 Hardware – provides basic computing resources
    - ▶ CPU, memory, I/O devices
  - 🍃 Operating system
    - ▶ Controls and coordinates use of hardware among various applications and users
  - 🍃 Application programs – define the ways in which the system resources are used to solve the computing problems of the users
    - ▶ Word processors, compilers, web browsers, database systems, video games
  - 🍃 Users
    - ▶ People, machines, other computers





# Four Components of a Computer System





# Operating System Definition

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👁️ OS is a **resource allocator**

- 🍃 Manages all resources

- 🍃 Decides between conflicting requests for efficient and fair resource use

👁️ OS is a **control program**

- 🍃 Controls execution of programs to prevent errors and improper use of the computer





# Computer Startup

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- 👁 **bootstrap program** is loaded at power-up or reboot
  - 🍃 Typically stored in ROM or EPROM, generally known as **firmware**
  - 🍃 Initializes all aspects of system
  - 🍃 Loads operating system kernel and starts execution

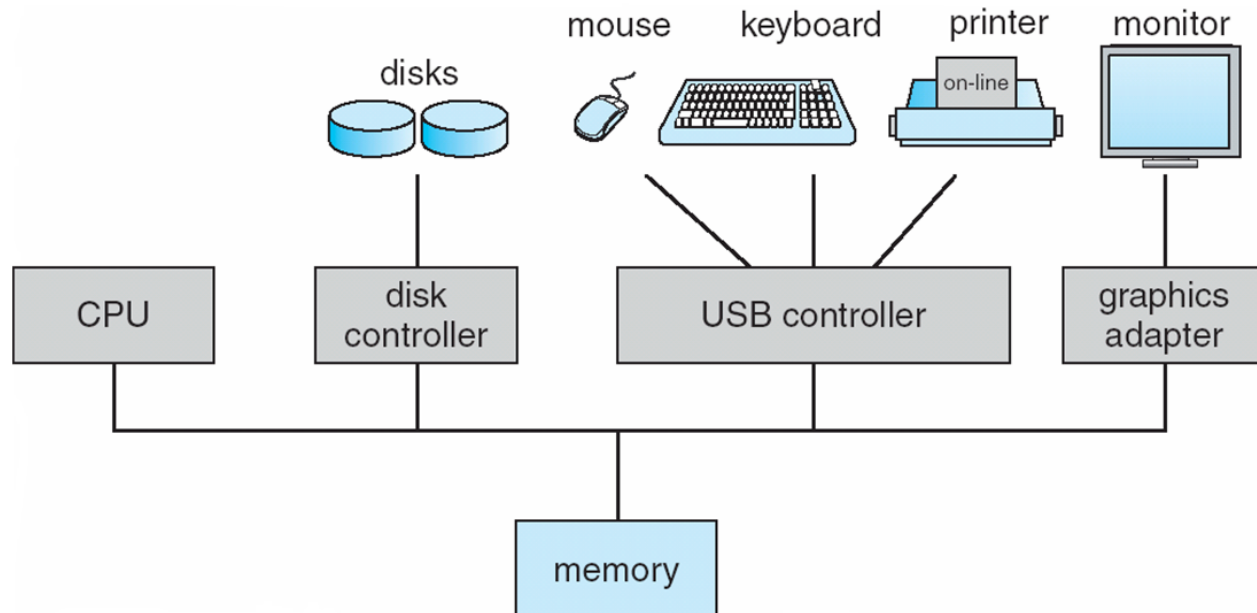




# Computer System Organization

## 👁 Computer-system operation

- 👉 One or more CPUs, device controllers connect through common bus providing access to shared memory
- 👉 Concurrent execution of CPUs and devices competing for memory cycles





# Computer-System Operation

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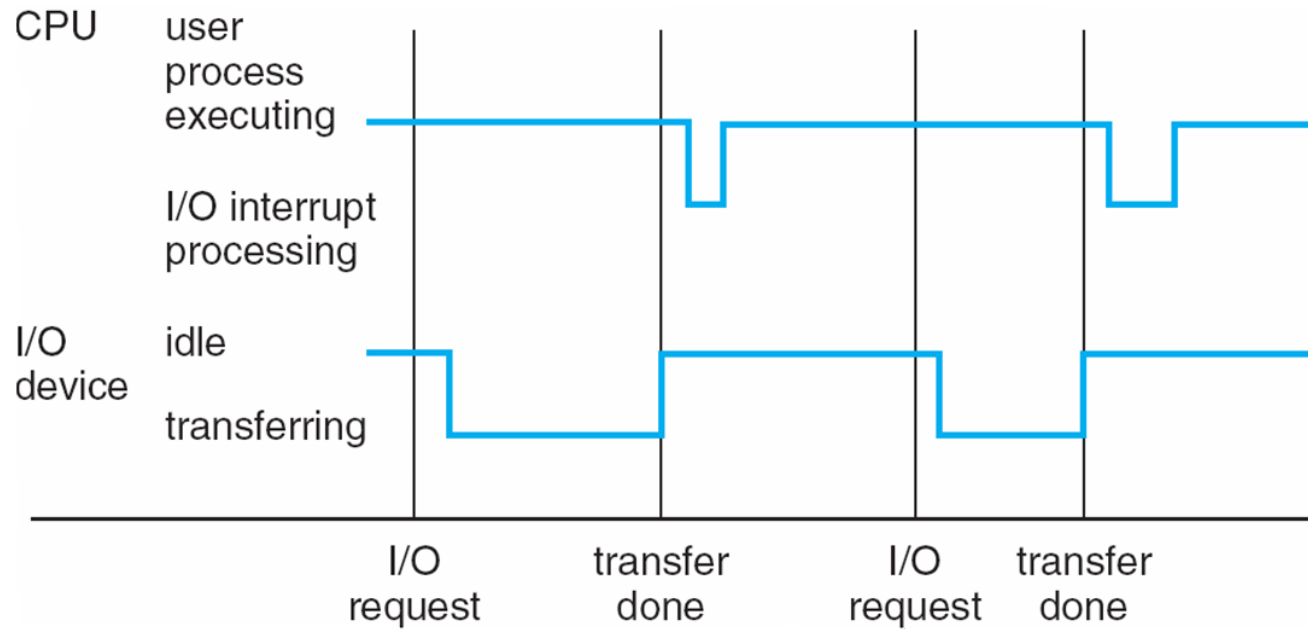
- 👁 I/O devices and the CPU can execute concurrently
- 👁 Each device controller is in charge of a particular device type
- 👁 Each device controller has a local buffer
- 👁 CPU moves data from/to main memory to/from local buffers
- 👁 I/O is from the device to local buffer of controller
- 👁 Device controller informs CPU that it has finished its operation by causing an **interrupt**
- 👁 An operating system is **interrupt driven**





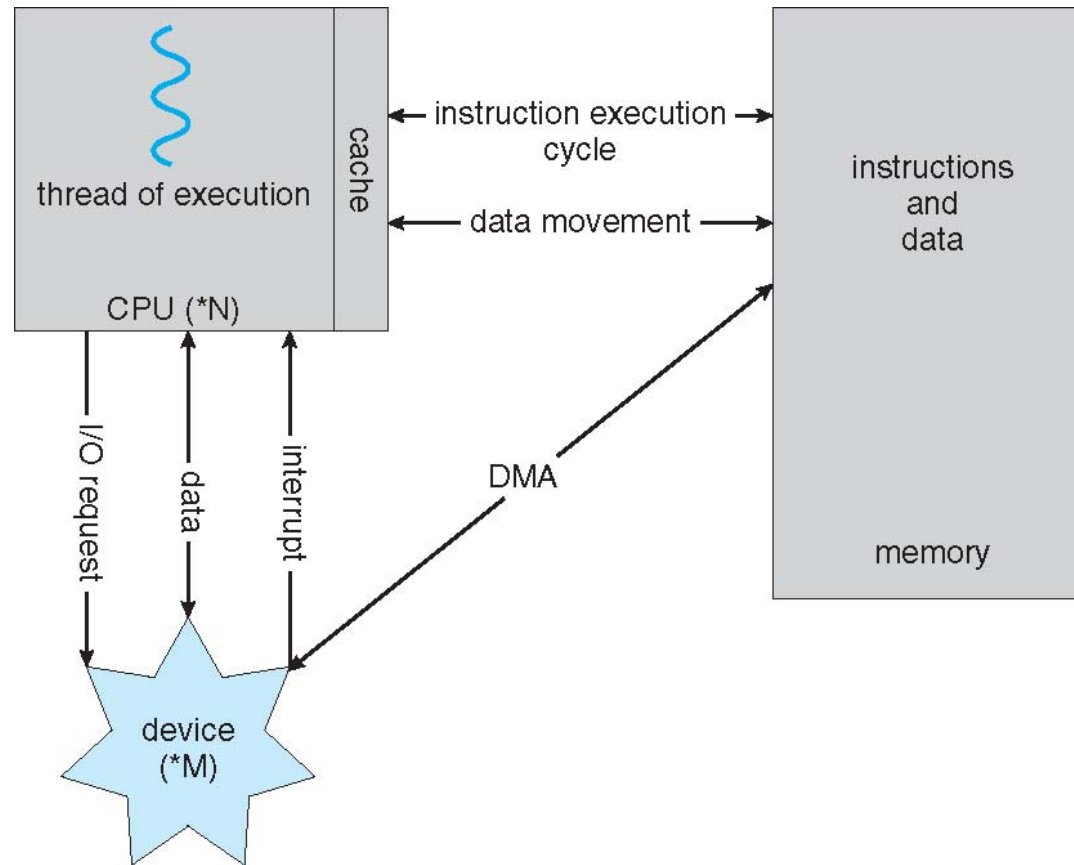


# Interrupt Timeline





# How a Modern Computer Works



*A von Neumann architecture*





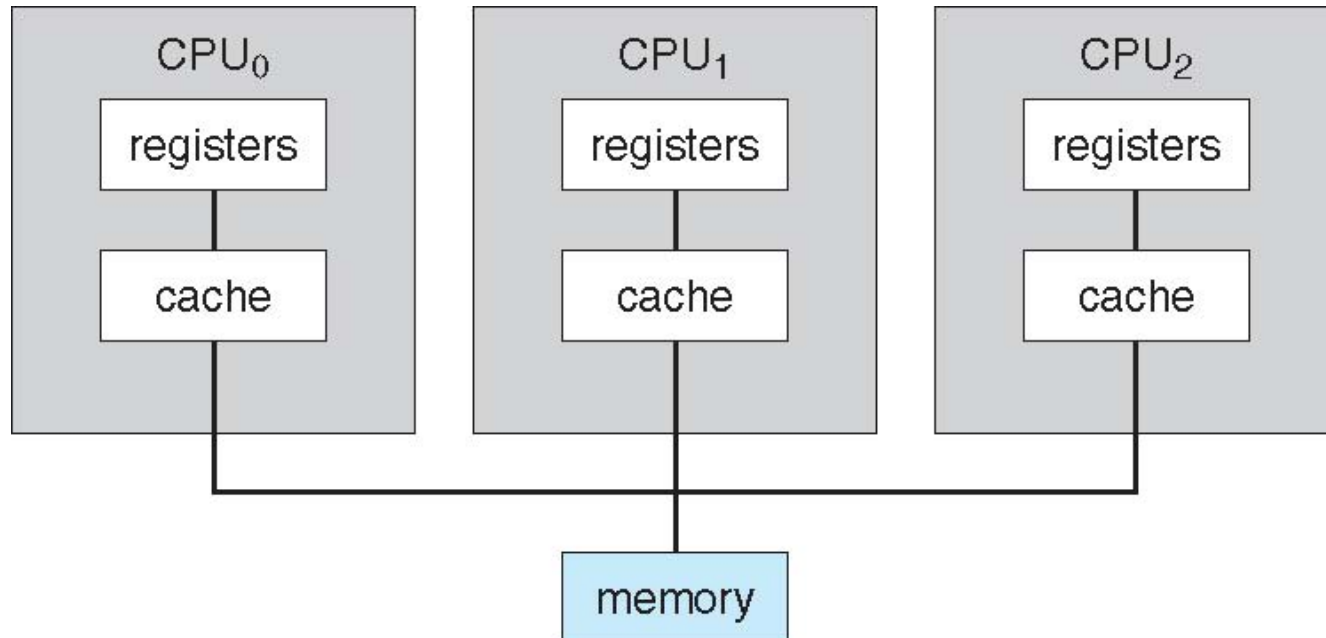
# Computer-System Architecture

- 👁 Most systems use a single general-purpose processor
  - 🍃 Most systems have special-purpose processors as well
- 👁 **Multiprocessors** systems growing in use and importance
  - 🍃 Also known as **parallel systems**, **tightly-coupled systems**
  - 🍃 Advantages include:
    1. **Increased throughput**
    2. **Economy of scale**
    3. **Increased reliability** – graceful degradation or fault tolerance
  - 🍃 Two types:
    1. **Asymmetric Multiprocessing** – each processor is assigned a specific task.
    2. **Symmetric Multiprocessing** – each processor performs all tasks





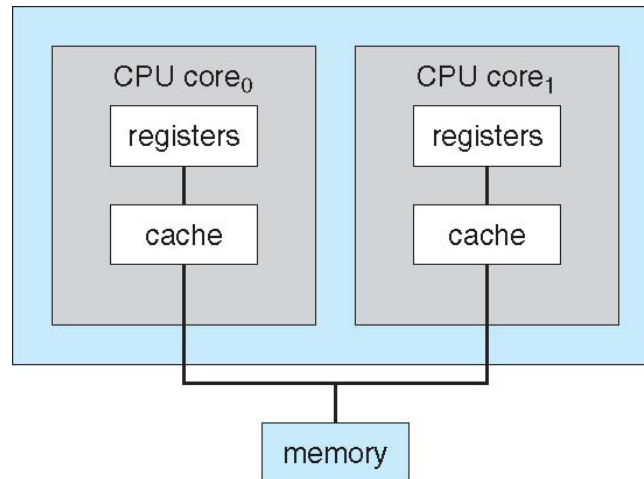
# Symmetric Multiprocessing Architecture





# A Dual-Core Design

- 👁 Multi-chip and **multicore**
- 👁 Systems containing all chips
  - 🍃 Chassis containing multiple separate systems





# Operating System Structure

## 👁 Multiprogramming

- 👉 Single user cannot keep CPU and I/O devices busy at all times
- 👉 Multiprogramming organizes jobs (code and data) so CPU always has one to execute
- 👉 A subset of total jobs in system is kept in memory
- 👉 One job selected and run via **job scheduling**
- 👉 When it has to wait (for I/O for example), OS switches to another job

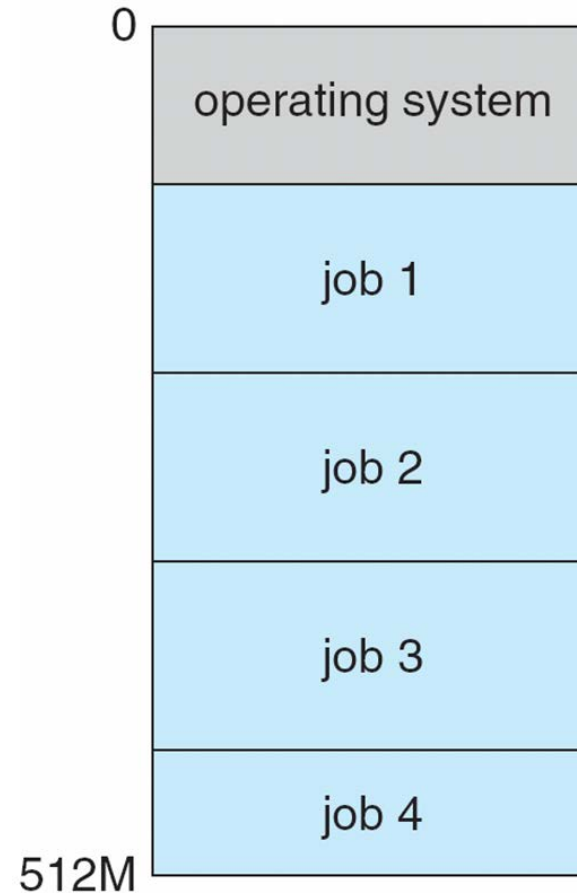
## 👁 Timesharing (**multitasking**) is logical extension in which CPU switches jobs so frequently that users can interact with each job while it is running, creating **interactive** computing

- 👉 **Response time** should be  $< 1$  second
- 👉 Each user has at least one program executing in memory ⇄ **process**
- 👉 If several jobs ready to run at the same time ⇄ **CPU scheduling**
- 👉 If processes don't fit in memory, **swapping** moves them in and out to run
- 👉 **Virtual memory** allows execution of processes not completely in memory





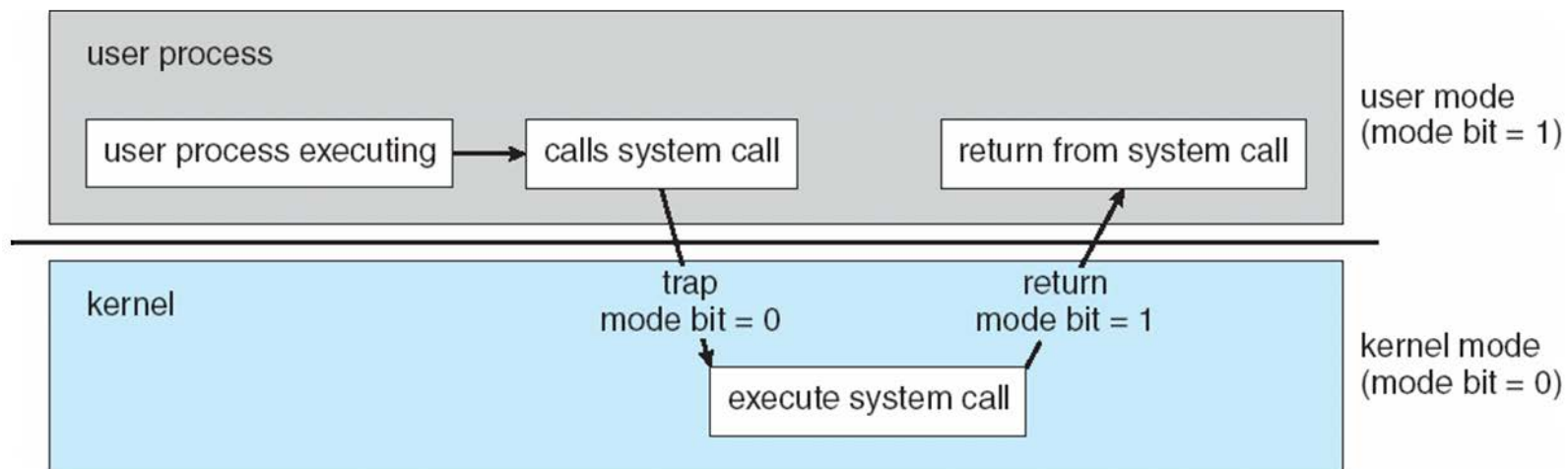
# Memory Layout for Multiprogrammed System





# Transition from User to Kernel Mode

- 👁️ Timer to prevent infinite loop / process hogging resources
  - 👉 Timer is set to interrupt the computer after some time period
  - 👉 Keep a counter that is decremented by the physical clock.
  - 👉 Operating system set the counter (privileged instruction)
  - 👉 When counter zero generate an interrupt
  - 👉 Set up before scheduling process to regain control or terminate program that exceeds allotted time







# Process Management

- 👁️ A process is a program in execution. It is a unit of work within the system. Program is a ***passive entity***, process is an ***active entity***.
- 👁️ Process needs resources to accomplish its task
  - 👉 CPU, memory, I/O, files
  - 👉 Initialization data
- 👁️ Process termination requires reclaim of any reusable resources
- 👁️ Single-threaded process has one **program counter** specifying location of next instruction to execute
  - 👉 Process executes instructions sequentially, one at a time, until completion
- 👁️ Multi-threaded process has one program counter per thread
- 👁️ Typically system has many processes, some user, some operating system running concurrently on one or more CPUs
  - 👉 Concurrency by multiplexing the CPUs among the processes / threads





# Process Management Activities

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The operating system is responsible for the following activities in connection with process management:

- 👁 Creating and deleting both user and system processes
- 👁 Suspending and resuming processes
- 👁 Providing mechanisms for process synchronization
- 👁 Providing mechanisms for process communication
- 👁 Providing mechanisms for deadlock handling



# End of Chapter 1

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