

# Applications of Graph Integration to Function Comparison and Malware Classification

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# Agenda

1. Overview of our Vectorization Method
2. The .NET Framework and Common Language Runtime (CLR)
3. Decompilation
4. Graph Integration
5. Results

# Overview of our Vectorization Method

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# Summary

## Overall Goal

Construct a vectorization method to be leveraged by a classifier on .NET files

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4. Compute component-wise mean/std of antiderivatives across all  $G$  resulting from decompilation of the given file

# **The .NET Framework and Common Language Runtime (CLR)**

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# .NET Framework - two main components

## Framework Class Library (FCL)

- user interface
- data access
- database connectivity
- cryptography
- web application development

## Common Language Runtime (CLR)

- is an application virtual machine which provides

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- compilation of high-level .NET code results in an Intermediate Language Binary
- the CLR JITs the code from IL to machine code run on the cpu

# Decompilation

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# Decompilation

## Definition

*Decompilation* is a program transformation by which compiled code is transformed into a high-level human-readable form.

## Definition

An *Abstract Syntax Tree* is a tree representation of the abstract syntactic structure of the source code, where each node denotes a construct occurring in the source code.

Program control flow is understood by studying the structure of two types of control flow graphs resulting from decompilation.

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- the function call graph describes the calling structure of the functions (subroutines) constituting the overall program
- *Shortsighted Data Flow Graphs* (SDFG) - each obtained by merging all paths through the AST corresponding to a constituent function

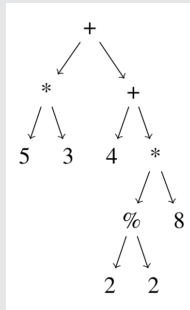
# Abstract Syntax Trees

## Example: Arithmetic Expressions

Consider the following BinaryOp expression:

$$5 * 3 + (4 + 2 \% 2 * 8)$$

The semantic structure of this expression can be distilled by considering the following binary tree:



Distilled semantic structure = order of operations



## Control Flow

- if - reference the conditional and execute accordingly
- break - immediately exit the enclosing loop
- [CLRWhile](#) - infinite loop

## Expressions

Code that when evaluated **does** yield a value. Valid in places such as tests, for loops, conditionals, or as the right-hand side of assignments.

- BinaryOp - expression computed from two operands and some operator
- [Call](#) - function call, including list of args pass to the function

## Statements

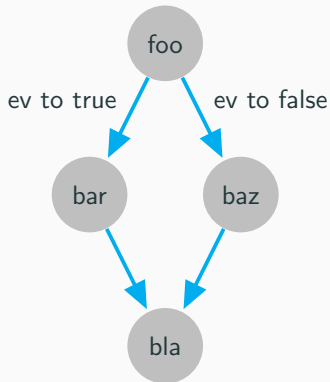
Code that when evaluated **does not** yield a value. E.g., a statement cannot be on the right-hand side of an assignment.

- [Assignment](#) - storage of rh variable to the location yielded by lh variable
- [CLRVariableWithInitializer](#) - declaration and subsequent initialization

# Construction of Shortsighted Data Flow Graph

Small code block resulting in a nonlinear SDFG.

```
if foo() {  
    bar();  
}  
else {  
    baz();  
}  
bla();
```



# Functions on SDFG Graphs - Motivation

We often study an object  $X$  by studying a set of functions defined on  $X$

$$X^* := \{f : X \longrightarrow \mathbb{R}\}$$

## Example 1

Consider the case of a distribution  $\mathcal{D}$  on a sample space  $\Omega$  defined by the measure  $\mu$ . We might choose to study  $\mathcal{D}$  by studying

$$f_n : \mathcal{D} \mapsto \int_{\Omega} x^n d\mu(x)$$

## Example 2

Consider the set of invertible  $n \times n$  matrices  $GL_n(\mathbb{F})$  on some field  $\mathbb{F}$ . We might choose to study  $GL_n(\mathbb{F})$  by studying

$$\text{tr}, \det : GL_n(\mathbb{F}) \longrightarrow \mathbb{R}$$

# Functions on SDFG Graphs

Let  $G$  be a SDFG graph resulting from traversing a given AST corresponding to some source code function.

## Example

Define

$$\text{NumPass2Call} : \text{Vert}(G) \longrightarrow \mathbb{R}$$

by

$$v \mapsto \# \text{args}_v$$

where  $\# \text{args}_v$  is the number of arguments passed to the function called at  $v$ .

Other Examples:

1.  $\text{BinaryOp} : v \mapsto \eta(\text{whichOpCode}_v)$

for some string-to-float hash function  $\eta$ .

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1.  $\text{BinaryOp} : v \mapsto \eta(\text{whichOpCode}_v)$
2.  $\text{CLRClassRef} : v \mapsto \eta(\text{ReferencedClass}_v)$

for some string-to-float hash function  $\eta$ .

# Graph Antiderivative - Ingredients

In order to define an integral of a function

$$f : \text{Vert}(G) \longrightarrow \mathbb{R}$$

for  $G$  a directed graph, we must define a measure  $\mu$  on  $\text{Vert}(G)$  in such a way that  $f$  is measurable.

We do this by imposing a Markov chain structure on  $G$  and taking  $\mu$  to be the PageRank measure

$$\mathbb{P} : \text{Vert}(G) \longrightarrow [0, 1]$$

$$v \mapsto \text{PageRank}(G)_v$$

where the PageRank vector is taken to be the steady-state probability distribution over the nodes resulting from the long-run behavior of the random-walk Markov Chain.

# Graph Integration

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# Markov Chains and the PageRank Vector

## Definition

A discrete-time *Markov chain* is a sequence of random variables  $X_1, X_2, \dots$  such that

$$P(X_{n+1} = x | X_1 = x_1, \dots, X_n = x_n) = P(X_{n+1} = x | X_n = x_n)$$

Given  $G$ , order the vertices  $\{v_i\}$  of the graph  $G$  and define the  $n \times n$  probability transition matrix  $T$  by

$$t_{ij} = \begin{cases} 1/|v_i^{\text{out}}| & \text{if } (v_i, v_j) \in \text{Edges}(G) \\ 0 & \text{otherwise} \end{cases}$$

where  $v_i^{\text{out}}$  is the set of edges emanating from vertex  $v_i$  and  $n = |\text{Vert}(G)|$ .

To ensure the irreducibility of our transition matrix, we smooth  $T$  to

$$M = (1 - p)T + pB \quad (\text{Perron-Frobenius})$$

where

$$B = \frac{1}{n} \begin{bmatrix} 1 & 1 & \dots \\ \vdots & \ddots & \\ 1 & & 1 \end{bmatrix}$$



# PageRank Measure

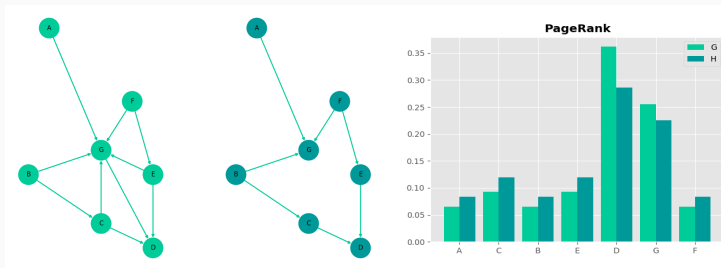
The PageRank vector  $\mathbb{P}$  is given by the left eigenvector of  $M$  and corresponds to

$$\lim_{n \rightarrow \infty} M^n \frac{1}{|\text{Vert}(G)|} \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix}$$

and can usually be adequately approximated with  $n = 10$ .

The corresponding Markov chain is defined by

$$P(X_t = v_i | X_{t-1} = v_j) = (1 - p)t_{ij} + p \frac{1}{n}$$



# Construction of the Graph Integral

Consider a function  $f : \text{Vert}(G) \rightarrow \mathbb{R}$  for  $G$  a finite directed graph.

Let  $\mathbb{P} = \{p_v\}$  be the PageRank measure on  $\text{Vert}(G)$ . We can then define a measure  $\nu_f$  on  $\text{Vert}(G)$  by

$$\begin{aligned}\nu_f(S) &= \int_S f \, d\mathbb{P} \\ &= \sum_{\alpha_j \in \text{image}(f)} \alpha_j \mathbb{P}(f^{-1}(\alpha_j) \cap S) \\ &= \sum_{v \in S} f(v) p_v\end{aligned}$$

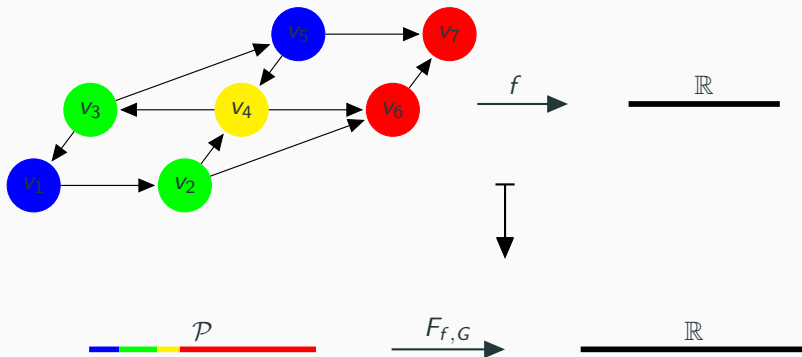
Let  $\mathcal{P}$  be a partition of  $[0, 1]$  and let  $G_q = \{v | p_v \leq q\}$ .

$$G_{q_1} \subseteq G_{q_2} \subseteq \cdots \subseteq G_{q_{|\mathcal{P}|}} = \text{Vert}(G)$$

allows us to define our **graph antiderivative**  $F_{f,G}$  of  $f$  by

$$\begin{aligned}F_{f,G} &:= (\nu_f(G_{q_1}), \nu_f(G_{q_2}), \dots, \nu_f(G_{q_{|\mathcal{P}|}})) \\ &= (\mathbb{E}[f|_{G_{q_1}}], \mathbb{E}[f|_{G_{q_2}}], \dots, \mathbb{E}[f|_{G_{q_{|\mathcal{P}|}}}] )\end{aligned}$$

# The Graph Antiderivative Visualized



## Antiderivative

$$\Gamma \times \text{Fun}\left(\bigsqcup_{\Gamma} \text{Vert}(G), \mathbb{R}\right) \longrightarrow \text{Fun}(\mathcal{P}, \mathbb{R})$$

$$(G, f) \mapsto (F_{f,G} : q \mapsto \mathbb{E}[f|_{G_q}]),$$

# Graph Integration: Example

Consider a SDFG  $G$  given by:

$$\text{Edge}(G) = \{(v_1, v_2), (v_1, v_3), (v_2, v_4), (v_3, v_4)\}$$

$$\text{PageRank}(G) = \langle p_{v_1} = 0.10, p_{v_2} = 0.15, p_{v_3} = 0.25, p_{v_4} = 0.50 \rangle$$

Assume the nodes  $v_1, v_4 \in \text{Vert}(G)$  both correspond to function calls  $\phi_{v_i}(\text{args}_{v_i})$ , where  $\text{args}_{v_i}$  represent the set of arguments passed to  $\phi_{v_i}$ . Define

$$\text{NumPass2Call} : \text{Vert}(G) \longrightarrow \mathbb{R}$$

by

$$v_i \mapsto \begin{cases} \#\text{args}_{v_i} & \text{if } i \in \{1, 4\} \\ 0 & \text{otherwise} \end{cases}$$

Let  $\mathcal{P} = (0.05, 0.12, 0.95)$ .

Then  $F_{\text{NumPass2Call}, G} : \mathcal{P} \longrightarrow \mathbb{R}$  takes the form

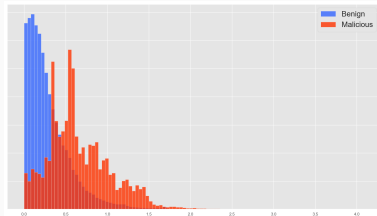
$$\begin{pmatrix} 0.05 \\ 0.12 \\ 0.95 \end{pmatrix} \mapsto \begin{pmatrix} 0 \\ 0.1 * \#\text{args}_{v_1} \\ 0.1 * \#\text{args}_{v_1} + 0.5 * \#\text{args}_{v_4} \end{pmatrix}$$

## Results

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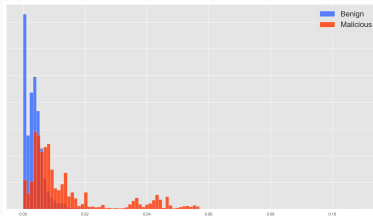
# Vectorization Efficacy

$$\int_0^{0.6} \text{ClassRefname} : v \mapsto \eta(\text{name}(v))d\mathbb{P}$$



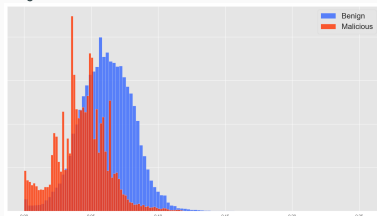
Name of referenced class at  $v$

$$\int_0^{0.4} \text{CLRLiteral} : v \mapsto \eta(\text{type}(v))d\mathbb{P}$$



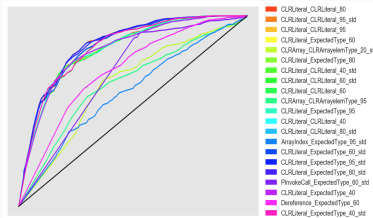
Type of literal occurring at  $v$

$$\int_0^{0.95} \text{ArgRefType} : v \mapsto \eta(\text{type}(v))d\mathbb{P}$$



Type of argument referenced at  $v$

Top Features by AUC



Value/type of literal expression at  $v$

# Model Results - Random Forest

**Table 1:** Graph Antiderivative-based vectorization

Class	Precision	Recall	F1-score	Support
Benign	97.88%	99.37%	98.62%	696827
Malware	98.94%	96.47%	97.69%	424420
avg/total	98.28%	98.27%	98.27%	1121247
False Positive Rate	1.10%			
False Negative Rate	1.72%			

**Table 2:** Text-only vectorization

Class	Precision	Recall	F1-score	Support
Benign	90.61%	87.04%	88.79%	696827
Malware	87.80%	91.18%	89.46%	424420
avg/total	89.19%	89.13%	89.13%	1121247
False Positive Rate	8.79%			
False Negative Rate	12.96%			

**Questions?**