Racket: Macros

Advanced Functional Programming

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Today

- Macros
 - o pattern-based macros
 - Hygiene
 - Syntax objects and general macros
 - Examples

Macros

(According to the Racket Guide...) A macro is a new syntactic form that can be expanded into existing forms through a transformer.

Let's start with an example. Say we do not have the or construct and we want to implement it in terms of if.

Implementing or: First try

```
(define (or x y)
  (if x x y))
```

What are the problems?

```
(or 3 (/ 2 0))
```

• y is always evaluated.

A function does not work. We need something that would transform the or into a if before running time.

Using define-syntax-rule

define-syntax-rule defines a macro.

```
(define-syntax-rule (or x y)
     (if x x y))
```

(or 3 (/ 2 0)) is effectively transformed into (if 3 3 (/ 2 0)).

The problem is solved. But we introduced another one:

```
(or (begin (display 3) 3) (/ 2 0))
```

x is evaluated twice.

Introducing let

```
(define-syntax-rule (or x y)
  (let ([z x])
      (if z z y)))

(or 3 (/ 2 0))
(or (begin (display 3) 3) (/ 2 0))
```

Seems like we are good.

Hygiene

Consider another example

```
(define-syntax-rule (swap! x y)
  (let ([tmp x])
        (set! x y)
        (set! y tmp)))
```

What if we pass a variable named tmp (or even set! or let) to swap!?

The naive rewriting would produce an erroneous code.

Hopefully, the racket macro system is "hygienic", that is it makes sure to avoid clashes between identifiers and it uses lexical scoping.

More complicated macros

- syntax-rules defines several rules by pattern matching.
- ... in the pattern matches several (zero, one or more) occurences of the last pattern (y in this case).
- ... in the template provides all the bound occurences.

Using literal keywords

do is part of the macro definition.

A wrong version of while

```
(define-syntax while
  (syntax-rules (do)
   [(while cond do body ...)
        (when cond
        body ...
        (while cond do body ...))]))
        Why?
```

The macro are expanded before runtime.

In this case, the expansion never completes.

Examples

- Memoization
- Rewriting

Example: Implementing Memoization

Memoization is to record the result of previous calls to a procedure to avoid recomputing them again.

For instance, a (naive) fibonacci procedure takes exponential time to complete.

```
(define (fib x)
  (case x
     [(0 1) 1]
  [else (+ (fib (- x 1)) (fib (- x 2)))]))
```

Memoization would allow us to reuse already computed results.

We will define a macro define-memo to define a procedure with memoization.

define-memo

Previous results are stored in a (mutable) hash-table.

The keys are the arguments of the function.

The actual body of the function is executed only if the arguments are not already in the hash-table.

define-memo, application

```
(define-syntax-rule (define-memo (f args ...) body ...)
  (define f
    (let ([memo (make-hash)])
      (lambda (args ...)
        (cond [(hash-has-key? memo (list args ...))
               (hash-ref memo (list args ...))]
              [else
               (let ([res (begin body ...)])
                 (hash-set! memo (list args ...) res)
                 res)])))))
(define-memo (fib/m x)
  (display x); to see when we actually execute the body
  (case x
    [(0 \ 1) \ 1]
    [else (+ (fib/m (- x 1)) (fib/m (- x 2)))])
(fib/m 10)
(fib/m 10)
```

Example: Expressions Simplification

It is often useful to simplify expressions.

We will do it here by rewriting rules.

Consider this target piece of code

simplifier returns a function of one argument: the expression to simplify.

Simplifying one expression

Each rule is rewriten as a clause in a match expression.

The whole code is enclosed in a procedure.

-> is part of the macro definition, not a variable.

Making it recursive

The previous code only simplifies the root of the expression.

We need to call the simplification on each sub-expression.

We need to repeat the simplification until fix-point.

simplify is the function produced by simplifier.

Putting it together

Going Further

Up to now, we have seen macros based on patterns, rewriten to a template.

Macros can do more than that, by manipulating the syntax (almost) however you want.

For this, we need to introduce syntax objects.

Syntax Objects

(syntax (+ 1 2)) returns a syntax object representing the data (+ 1 2) with lexical information and source code location.

Syntax objects can be manipulated by the macro system.

```
(define-syntax Hello
  (lambda (stx)
      (display stx)
      (datum->syntax stx (cdr (syntax->datum stx)))))
(Hello + 1 3)
```

The body of define-syntax must be a procedure taking a syntax object and returning a syntax object.

syntax-case

syntax-rules is a shortcut for "simple" macros: one writes directly the resulting code.

In the general form, one can use the whole power of Racket to modify the syntax.

syntax-case allows one to match the syntax against patterns.

#' is a shortcut for syntax

There is also #` and #, for quasisyntax and unsyntax.

syntax-case: Example

The swap example with syntax.

syntax-case: Example

The shortcut version to avoid the lambda.

This is not better than syntax-rules but wait...

syntax-case: Example (2)

Checking (before runtime) that the arguments make sense.

```
(define-syntax (swap stx)
  (syntax-case stx ()
    [(swap a b)
     (cond [(and (identifier? #'a) (identifier? #'b))
            #'(let ([tmp a])
                (set! a b)
                (set! b tmp))]
           [else (raise-syntax-error
                  #f
                  "not an identifier"
                  stx
                  (if (identifier? #'a)
                      #'a
                      #'b))])))
```

More syntax manipulation

syntax-e returns the data in the syntax object.

syntax->datum recursively calls syntax-e.

datum->syntax creates a syntax object from some data and the lexical information of another syntax object.

syntax-list is like syntax-e but it unrolls a whole list.

with-syntax is similar to a let statement but for syntax objects.

syntax? tells whether the given value is a syntax object.

identifier? tells if the syntax object represents an identifier (syntax-e returns a symbol).

Example: Extended Rewriting

Now we want to be able to put guards on our rewriting rules.

Extended Rewriting with syntax

```
(define-syntax (simplifier stx)
  (define (clause expr)
    (syntax-case expr (? ->)
      [(x -> y) #'[`x `y]]
      [(x ? z -> y) #'[`x (=> fail)]
                           (when (not z) (fail)) `y]]))
  (syntax-case stx ()
    [(_ expr ...)
     (with-syntax ([(res ...)
                    (map clause (syntax->list
                                  #'(expr ...)))])
       #'(simplify (lambda (form)
                     (match form
                        res
                        [else form])))))))
```

Breaking Hygiene

Sometimes we want to create identifiers that are available outside of the macro.

This happens when creating a struct, for instance.

(struct box (val)) creates e.g. the procedures box-val and box?.

There is a simple way to break hygiene:

Use (datum->syntax lex-cont name) where lex-cont is an identifier that is given to the macro.

In that case, the new name inherits the same lexical scope.

Breaking Hygiene: Example

Define a for loop with it bound to the current value.

Breaking Hygiene: Building names

This piece of code shows how to build a name from another one:

One need to go from a syntax object to a symbol to a string, manipulate the string, and then all the other way around.

Phases

We can use normal procedures when defining macros.

But macros are expanded before runtime.

How is that possible?

There are several phases.

The identifiers in each phase are only available in that phase, unless imported.

One can have even more than two phases.

Conclusion: word of caution

Macros can be very powerful.

As any powertool, they must be used with care.

In particular, do not use a macro if a procedure can make the work.