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1 Security: Introduction into my interests

I believe security is a good entrance into my current interests. It awards both attention to detail and high-level systematic view; its subfields connects directly with both theoretical and practical areas that fascinates me, and it offers effective tools to solve the imminent problems I face in real life.

The term *security* in computer science generally carries the sense of achieving *security goals* through *mechanisms or properties* of a system, despite the presence of adversaries in a *threat model*. The threat model has to be carefully chosen and reviewed; there would be no security against an omnipotent adversary.

Unlike in math, one cannot arbitrarily wiggle the threat model to one's desire. On the contrary, one's rights are increasingly threatened by the powerful digitally [21, 48, 44]. A prominent example is the continuous push for exceptional access to people's data and communication in various forms by governments¹, ignoring expert opinions against such measures [2, 1, 3]. I unfortunately has suffered from more aggressive attacks, primarily systematic, comprehensive, and far-reaching censorship [4] [53, Sect. 5], because of my background.

How my security goal of being able to exercise my human rights in the presence of these powerful adversaries drives my interests will be explained by the following subsections, some of which also describes how my more pure academic excitements are intertwined with this realistic drive. Finally, the last subsection, 1.3, will lead the transition into the other areas of my interest.

1.1 Secure Systems: Understanding the mechanisms

1.1.1 A high-level view of the Internet

A high-level understanding of how the digital infrastructures work is essential to understand what enables these threats. The current networking infrastructure, in my opinion, has conflicting properties. On one hand, the fundamental problem of the impossibility to link every two computers requires sharing of links, a

¹For instance, EU revived the Chat Control proposal in 2025. See <https://eutechloop.com/time-is-running-chat-control/>.

solution that welcomes centralization. On the other hand, the poor scalability of simple sharing schemes of a link (usually via a switch) necessitates a better scheme to extend a small network globally. To achieve this, the Internet relies heavily on delegation and distributed structures, such as topological divisions in its address and delegation of most routing responsibilities to each individual networks (e.g. ISPs). Its BGP, operating in between them, is mainly concerned exchanging reachability information, whereas its IGPs handle routing paths and allow different networks to implement different routing policies.

Therefore, the Internet has become a mixture of centralization and decentralization, where each end node is managed by an ISP network, yet no single ISP runs the whole Internet. Perhaps surprisingly, I found this hybrid structure more optimal for localized sabotage at the state level, creating “sub-Internet”s, each of which is crippled at a different level.

Unfortunately, the lower layers of the Internet have proven to have a great inertia for change. Handley gave a nice discussion on it in 2006 [26], and the trend he described has mostly been the same since then: *“the core Internet protocols have not changed significantly in more than a decade, in spite of exponential growth in the number of Internet users and the speed of the fastest links.”* [26]. On the other hand, the protocols in the higher layers evolved to a much greater degree. As the transport layer welcomes QUIC[27], the application layer has accumulated an enormous amount of innovation and progress, in particular, onion routing and Tor[24, 12], Bitcoin[37], VPN protocols[39, 14] and decentralized instant chat[33, 47], that fight for digital rights and/or promote decentralization. Yet one must not overlook the crypto-constructs that lay the foundation for all of them, which I discuss in detail in Section 1.3.

Unfortunately again, because all of these are built on the Internet, a state-level censor could easily abuse its local authority to target them. A few regimes are notorious to have blocked a vast amount of them to various extents, employing complicated passive analysis and active probing techniques [16, 29, 50, 51, 18, 17]. Although a state would need to consider the collateral damage, a totalitarian regime would not hesitate to block an entire protocol before it can figure out how to block it selectively [55, 51]. Apart from state-level actors, because commercial local ISPs additionally have an incentive in income and profit, not only do they perform surveillance and censorship like state-actors [6, 5], they also implement unjust policies easily with their local control of the infrastructure, like unfairly limiting the use of certain P2P protocols [13, 40, 35, 6]. Although ISPs argue that these P2P protocols can consume too much bandwidth, the other side of the story is that ISPs often oversubscribe and fraudulently advertise the bandwidth of Internet service they provide [42, 8]. This essentially is a probabilistic exploit on its customers — when almost all of the users happen to use the maximal bandwidth the ISP sells to them, congestion and throttling occur, and the users, not the ISP, ultimately pay the price — this is also the same kind of injustice imposed by airlines who oversell tickets.

The question of how to build protocols and systems that preserves one’s rights on top of the Internet that powerful adversaries control, fighting against surveillance, censorship, and overall other unfair practices which the infrastruc-

tures of the current networks happen to enable, is central to my interest in secure distributed systems.

1.1.2 Low-level system details

Beside the threat model, when we talk about security, we also implicitly assume another model, environment, or host, that encapsulates the problem. For instance, in the context of isolation,

- The host that encapsulates process isolation is the operating system kernel.
- The host that encapsulates virtual machine isolation is the hypervisor, often including hardware support.
- The host that encapsulates air-gapped machine isolation is the physical world, or more abstractly, the physical laws.

These hosts form a hierarchy. For instance, processes inside an OS may run inside a virtual machine, on hardware, within the physical world.

As a result, a security researcher has to understand these hosts, each to a different degree, depending on how secure she wants her systems to be. Even if a system is formally verified to be secure within a host, it can still be attacked from interactions with an outer host. That is well exemplified by the Meltdown attack, which breaks both kernel and hypervisor’s isolation of memory, because an outer host, the hardware, has a vulnerability [32]. In fact, Meltdown touches more than logical flaws in hardware; it relies on the timing difference of cache accesses to extract information from it [54], one that directly connects with the physical world. At this point, I feel that I should also mention the Spectre attacks family [30], which, unlike Meltdown, does not exploit a clear logic bug or fault in the hardware. Instead, they work by observing the side effects of various implementations of speculative execution, which is arguably not a bug but simply a feature implemented without caution to how information may be leaked via the outer hosts.

Indeed, there is little hope to understand and model the universe thoroughly and accurately in one’s life-time, and one can expect endless discoveries of how an otherwise secure system can leak information in unexpected ways, such as power and radiation via physical effects, which opens to a large range of attacks [31, 7, 19, 54], collectively defined as side-channel attacks. Then, one must ask, is it futile trying to cope with every such possible attack? My opinion is that it depends on the current understanding of the whole cybersecurity community. One novel attack just discovered may be increasingly understood and gain popularity by the whole community and thus gradually become a necessary part every secure system designer should take into account. I view this as an everlasting dynamics where people continuously push the boundary of our understanding, and consequently expanding the requirements for all in the community.

Therefore, a good understanding of the most common and well-understood hosts, the operating systems and the hardware (including the architecture), will

be essential for me as a security researcher. Additionally, figuring out how to perform clever hacks (as in the hacking culture of MIT, not the mainstream meaning of cracking) based on knowledge of details of a system gives a great sense of achievement to me. Moreover, this continuous expansion of knowledge frontier carries a special implication in security. Whereas a mathematician is perfectly fine to specialize in one direction without knowing every details in other fields, a computer system can never be secure if only designed with one perspective. For now, a good system is usually designed by collaboration; and it will be a very interesting challenge to figure out how one single researcher can be empowered to do this task better; could the LLMs be a potential assistant who fills the knowledge gaps? Finally, this is one major place where cybersecurity connects directly with a broad range of other fields, including software engineering, systems architecture, networking, and physics. These together contribute to my interest in the low-level details of secure systems.

1.2 Data-driven analysis: Lights through artificial clouds

Despite my intention to understand systems, many of them are not open for analysis, for various reasons. In particular, the aforementioned systems for digital surveillance and censorship are not only proprietary but often state secrets. On the other hand, similar to side channel attacks, these systems inevitably leak information in different ways.

By carefully observing their behavior and probing their responses under controlled conditions, we can gain insight. Combined with data science techniques, these observations reveal hidden structures and enable us to draw meaningful conclusions.

The process usually involves 1) raising questions and claims about some properties of a target system 2) conducting experiments (often involving purchased servers within such as Aliyun) 3) analyzing data 4) answering questions and verifying claims. As a simple example, Sheffey et al. very recently studied the IP addresses injected by the GFW censorship system. By first finding injected IPs and then probing them within the GFW, they found three categories of such IPs [46].

Thanks to continuous work from both academic scholars and dedicated organizations, (e.g. [52, 38, 17]), I and a small set of others who suffer from such systems and who are fortunate enough to know these works, could know what we are facing everyday better. Conversely, I strongly hope to contribute to the free-side of the arms race, making the free Internet reachable to everyone.

1.3 Cryptography: A mathematical savior

Security is often described as an arms race — that who controls more resources tends to discover more vulnerabilities, devise more attacks, and harden their systems more. Where is hope, then, when most self-censor and give up fighting the regime [9, 36, 49] and even fewer of the fighters have sufficient skills to join in the crypto war [41, 28]?

I believe one solution to what might sound like a power struggle lies within (modern) cryptography. For thousands of years since the use of the earliest symmetric encryption methods like the Caesar cipher, the ability to securely communicate through an insecure channel had still been mostly limited to the privileged who could afford persistent access to physical secure channels to exchange the keys and reliable safeguarding of the keys. A stunning turning point was found in what was widely regarded as the beginning of modern cryptography, Diffie and Hellman’s *New Directions in Cryptography* [11], where they gave a practical mathematical procedure to Merkle’s original idea for establishing a key known only to both parties over an insecure channel.

At first glance, the idea sounded so impossible that Merkle himself faced rejections when presenting the idea to his then professor Hoffman and to the CACM [34]:

“I am sorry to have to inform you that the paper is not in the main stream of present cryptography thinking and I would not recommend that it be published in the Communications of the ACM.”

“Experience shows that it is extremely dangerous to transmit key information in the clear.” [34]

The rejections represented a consensus of the old cryptography community that even Shannon concurred with: “*The key must be transmitted by non-interceptible means from transmitting to receiving points*” [45, p. 670], which demonstrates how “paradoxical” the idea was.

Such “paradoxical” ideas of modern cryptography are what attract me most to it, because they are the enabler that underlies those systems mentioned in section 1.1: TOR, Bitcoin, OpenVPN, Matrix, etc., and the catalyst that leads to the massive adoption of society advancements such as e-commerce. Remarkably, in less than 50 years since 1976, we already have a number of such discoveries in addition, and here’s a list of some of them.

Pseudorandom functions The kind of deterministic algorithms whose outputs look like that of a random oracle, by block ciphers like AES[10] or other ways [22], has a strong link with various other essential constructs like one-way functions.

Zero-knowledge proofs By interactively leveraging challenges that only a true prover could easily solve, zero-knowledge proof [25] enables one to verify that the prover knows a witness of a problem in NP [23] without revealing the witness.

Homomorphic encryptions Doing computation on encrypted data has been a very desirable property for many years. Although many earliest public-key encryption schemes, such as RSA [43] and ElGamal [15], already natively supported limited homomorphism like modular multiplication, by their design, the first full scheme that allowed arbitrary computation on encrypted data was only proposed in 2009 [20].

These discoveries have many applications and implications, two of which I pay most attention to. One is that they update our understanding of certain theoretical lower bounds of how much security one can achieve in the face of strong adversaries, no matter the power difference between them. The other is how they help minimize the trust required to perform global-level of collaboration that involves people from vastly different backgrounds and beliefs; such a collaboration that would typically require a strong central authority to organize, now only requires the participants to trust a few fundamental assumptions of cryptography.

Personally, I would go a step further. I believe modern cryptography offers not only technical tools, but a possible solution for addressing one of the big challenges our democracies face today: fragmentation, polarization, and the erosion of shared trust. By design, cryptography redistributes power — it makes mass surveillance prohibitively costly, even for state-level actors, by forcing them to target individuals rather than populations. For example, suppose breaking one person’s encryption, on average, required a single day of dedicated effort (through side-channels, vulnerabilities, or social engineering rather than mathematics), that constraint alone would prevent massive surveillance, since no adversary has that many days or resources to do this on everyone. In this way, encryption does something profound: it automatically and passively unites individuals, protecting each not by active coordination that is increasingly difficult, but by the collective shield of widespread adoption. Fortunately, as many modern infrastructures already deploy encryption and other crypto schemes by default, this “passive solidarity” is not a dream but a reachable reality, one that shows how mathematics can serve as a safeguard for democracy in an increasingly divided world.

To end this section, cryptography not only flows towards my passion for purer philosophical understanding, because of its deep connection with theoretical computer science and mathematics, to which I turn in Section 2, but also carries my hope to break free from the strong grasp of the tyrannies today, one that represents my free will which refuses to lose my agency and independence, in a world that too often seeks to reduce us into interchangeable screws for the system.

2 Mathematics: My purer philosophical pursuits

As much as I am driven by real-life problems to understand systems better, the force behind my sometimes unusually deep dive into even more primitive principles is my strong desire to understand the foundational reasonings. In this section, I will describe how that has led me to explore the purer mathematical areas and my current interests in them.

2.1 Type theory and metamathematics

In the manuals of various proof assistant languages, technical terms from type theory such as universe, sort, and introduction & elimination rules, are frequently mentioned. For me who is driven to understand what's going on behind, contact with type theory is my destiny. Furthermore, I ventured deeper into its history and the philosophical doctrines it was born from.

I have thought of a few ways to present this section, and in the end I decided that I will present it as a narrative of the basic history. If you share my curiosity for foundational reasonings, then after seeing the history I presented in my style, I believe you will also see my passion. A good starting point is probably Cantor's work on his theory of abstract sets, particularly his ingenious abstraction of sets to cardinal numbers for their comparison. In 1638 Galileo noted that there is a bijection between the natural numbers and their squares: $f : n \mapsto n^2$, raising questions about the principle that the whole is greater than any of its parts. Cantor took this further to systematically compare sets, using bijections as the primary tool.

An essential mental construct that Cantor used was his diagonal argument, which is often presented when proving the real numbers is not countable. Yet the central philosophical reasoning behind this argument applies to general sets which we don't know are countable or not, and this leads to this important theorem, whose proof I present because it's short and raises questions about many things.

Cantor's Theorem. *There is no bijection from any set S to the set of all of its subsets $\mathcal{P}(S)$. In Cantor's words, the cardinal number of $\mathcal{P}(S)$ is always larger than that of S .*

Proof. When S is empty, it is trivial that there's no surjection between \emptyset and $\{\emptyset\}$. Then we consider when S is not empty.

Suppose for now there exists a bijection $f : S \rightarrow \mathcal{P}(S)$. f 's being a bijection enables us to apply Cantor's diagonal argument, because now we can make sure the number of rows and columns are the same: arrange such a matrix that each column corresponds to an element $s \in S$, and each row is arranged to correspond to the sequence $f(s)$ for all s , in the same order as the columns. Along the diagonal, flip each element to define this subset $T \subset S$ such that $s \in T$ iff $s \notin f(s)$. Although we used the bijection fact to get to this definition, the definition itself does not require a f 's being a bijection; it can be any function $f : S \rightarrow \mathcal{P}(S)$. But we will need it to be a surjection for the argument that follows. So we now demote f to be any surjection.

Because f is a surjection, there is such s_0 that $f(s_0) = T$. Now we ask if $s_0 \in T$. Note that here we implicitly assume either $s_0 \in T$ or $s_0 \notin T$. In either case, by definition of T , we will arrive at a contradiction. Thus, by *reductio ad absurdum*, it is not the case that there exists such a surjection f . \square

Cantor and Russell arrived at two similar paradoxes, both involving the set of all sets.

Cantor’s paradox Consider the set of all sets Ω and its power set $\mathcal{P}(\Omega)$. By Cantor’s theorem, the latter should be larger than the former. On the other hand, $\mathcal{P}(\Omega)$, being a set of sets, should be a subset of Ω , a contradiction.

Russell’s paradox When Russell was later analyzing Cantor’s theorem, he wondered what if S itself was Ω and the supposedly existed f was the simple $f : s \mapsto s$. In that case, Cantor’s diagonal subset T becomes the famous form that we know of: $T : \{s \mid s \notin s\}$.

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