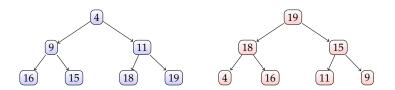
### **CSCI 2270: Data Structures**

#### Lecture 25: Priority Queues and Heaps

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#### **Priority Queues**

Heaps

### **Priority Queues**

- 1. A priority queue is an abstract data type (ADT) similar to queues, where the elements have a priority field attached to them.
- 2. In a priority queue, the elements with high priority are served before the elements with low priority.
- 3. In addition, in some implementations, the order amongst the elements with equal priority follows the order they were enqueued.
- 4. Stacks and queues may be modeled as priority queues, where in a stack the priority of each inserted element is monotonically increasing, while in a queue the priority of each inserted element is monotonically decreasing.
- Applications: Job-scheduling in operating systems, load-balancing problems, emergency-room patient priority, Dijkstra's algorithm to find shortest path, best-first search algorithms in graphs, Prim's minimum spanning tree algorithm.

## **Priority Queue Implementations**

- Key operations:
  - Adding an element (push)
  - Deleting an Element (pop).
- Implementation as an unsorted array:
  - push: O(1)
  - pop: O(n)
- Implementation as a sorted array:
  - push: O(n)
  - pop: O(1)
- Similar complexities for sorted and unsorted linked-list implementations.
- Can we do better?

## **Priority Queue Implementations**

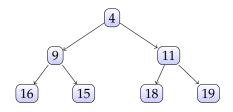
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Binary Heaps!

**Priority Queues** 

Heaps

### Binary Heaps: Min Heap

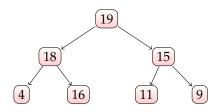


#### Min Heap Property:

- A complete binary-tree, i.e. difference in height between two branches is at most 1.
- If x is a node and y is its (either left or right) child then

 $x.priority \leq y.priority.$ 

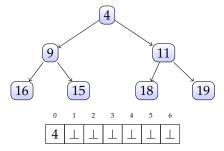
### Binary Heaps: Max Heap

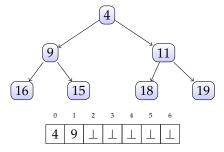


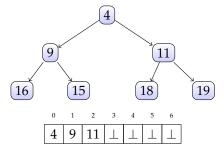
#### Max Heap Property:

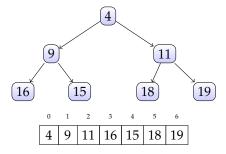
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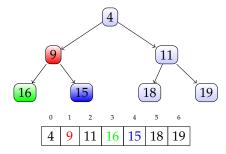
 $x.priority \ge y.priority$ .

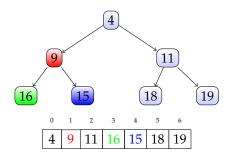












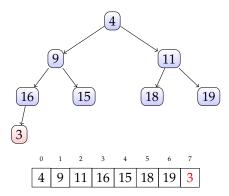
$$\begin{array}{rcl} \textit{leftChild}(i) & = & 2*i+1 \\ \textit{rightChild}(i) & = & 2*i+2 \\ \textit{parent}(i) & = & \textit{floor}((i-1)/2) \end{array}$$

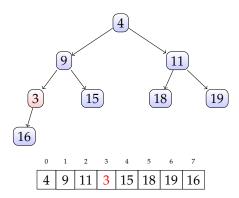
## Min Heap: Abstract DataType

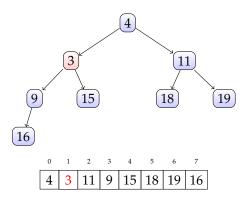
```
class MinHeap {
private:
  int* heap;
  int capacity;
  int currentSize;
  MinHeap();
 MinHeap(int s);
  ~MinHeap();
  void push (int value);
  int pop();
  int peek();
  void printHeap();
  void minHeapify(int index);
  int parent(int index) {return (index-1)/2;}
  int leftChild(int index) {return 2*index+1;}
  int rightChild(int index) {return 2*index+2;}
  void swap(int &x, int &y) {int z = x; x = y; y = z;}
};
```

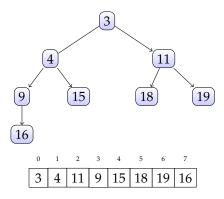
## Max Heap: Abstract DataType

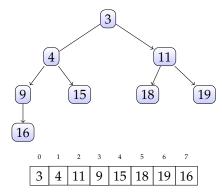
```
class MaxHeap {
private:
  int* heap;
  int capacity;
  int currentSize;
 MaxHeap();
  MaxHeap(int s);
  -MaxHeap();
  void push (int value);
  int pop();
  int peek();
  void printHeap();
  void maxHeapify(int index);
  int parent(int index) {return (index-1)/2;}
  int leftChild(int index) {return 2*index+1;}
  int rightChild(int index) {return 2*index+2;}
  void swap(int &x, int &y) {int z = x; x = y; y = z;}
};
```

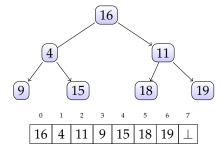


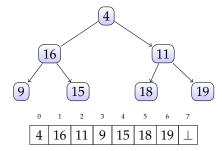


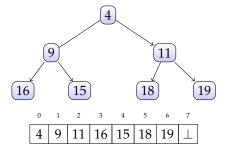




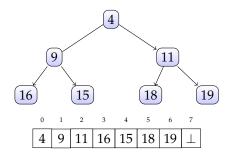








## Min Heap: Complexity



- Push: O(log(n))

- Pop: O(log(n))