

MonetDB:

Reaching the stars step by step

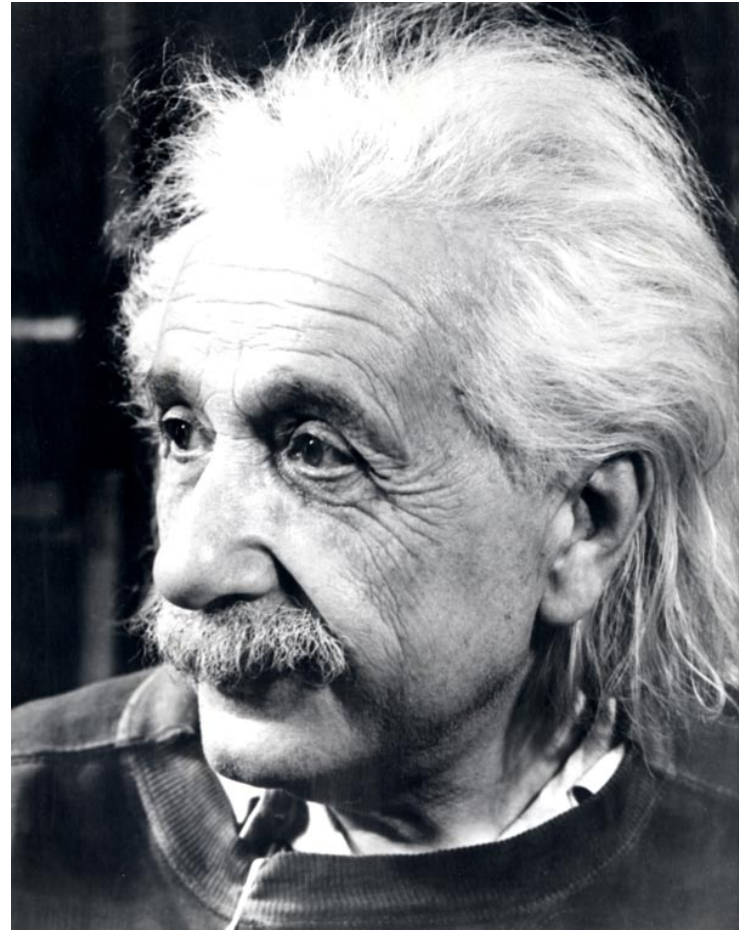
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<http://www.monetdb.org/>

“We can't solve problems by using the same kind of thinking we used when we created them.”



The world of column stores

Functionality and performance of MonetDB

Roadmap for a Science Database System

The landscape

Motivation

- Relational DBMSs dominate since the late 1970's / 1980's
 - Transactional workloads (OLTP, row-wise access)
 - I/O based processing
 - Ingres, Postgresql, MySQL, Oracle, SQLserver, DB2, ...
- Column stores dominate product development since 2008
 - Datawarehouses and business intelligence applications
 - Startups: Infobright, Aster Data, Greenplum, LucidDB, ...
 - Commercial: Microsoft, IBM, SAP, ...

MonetDB, the pioneer

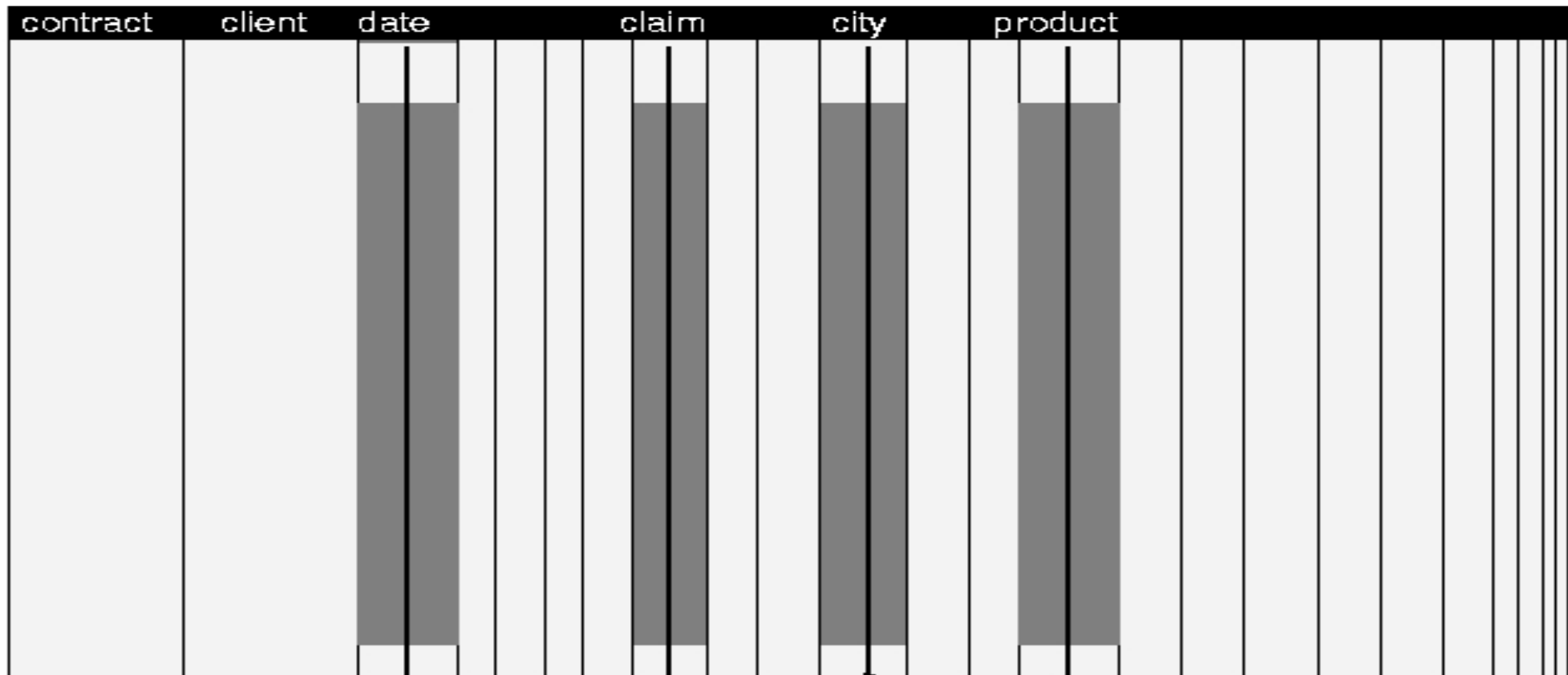
Workload changes: Transactions (OLTP) vs ... ↩

[illegible]

```
find client
10032112
```

OLTP queries:
access all
columns of
just one row.

Workload changes: ... vs OLAP, BI, Data Mining, ...



select those tuples
sold after march 21

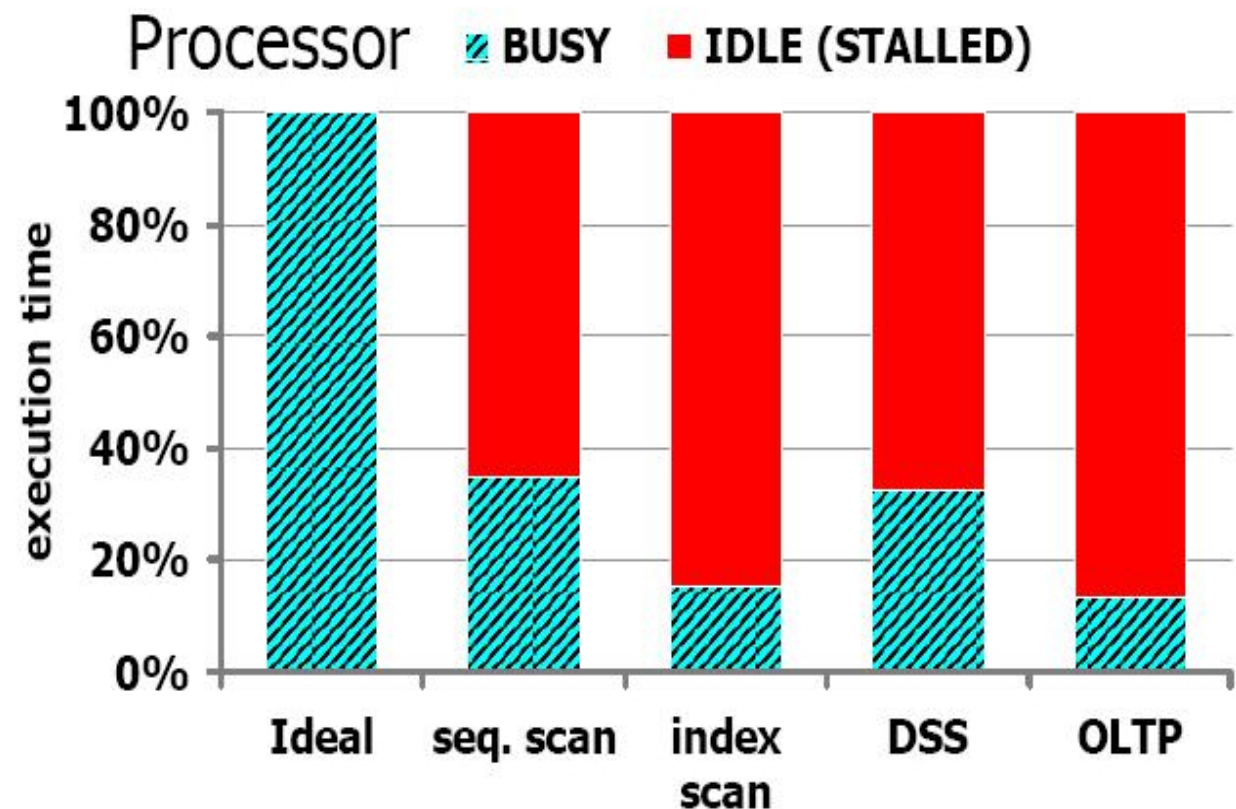
sum
claims

while grouping by
city and product

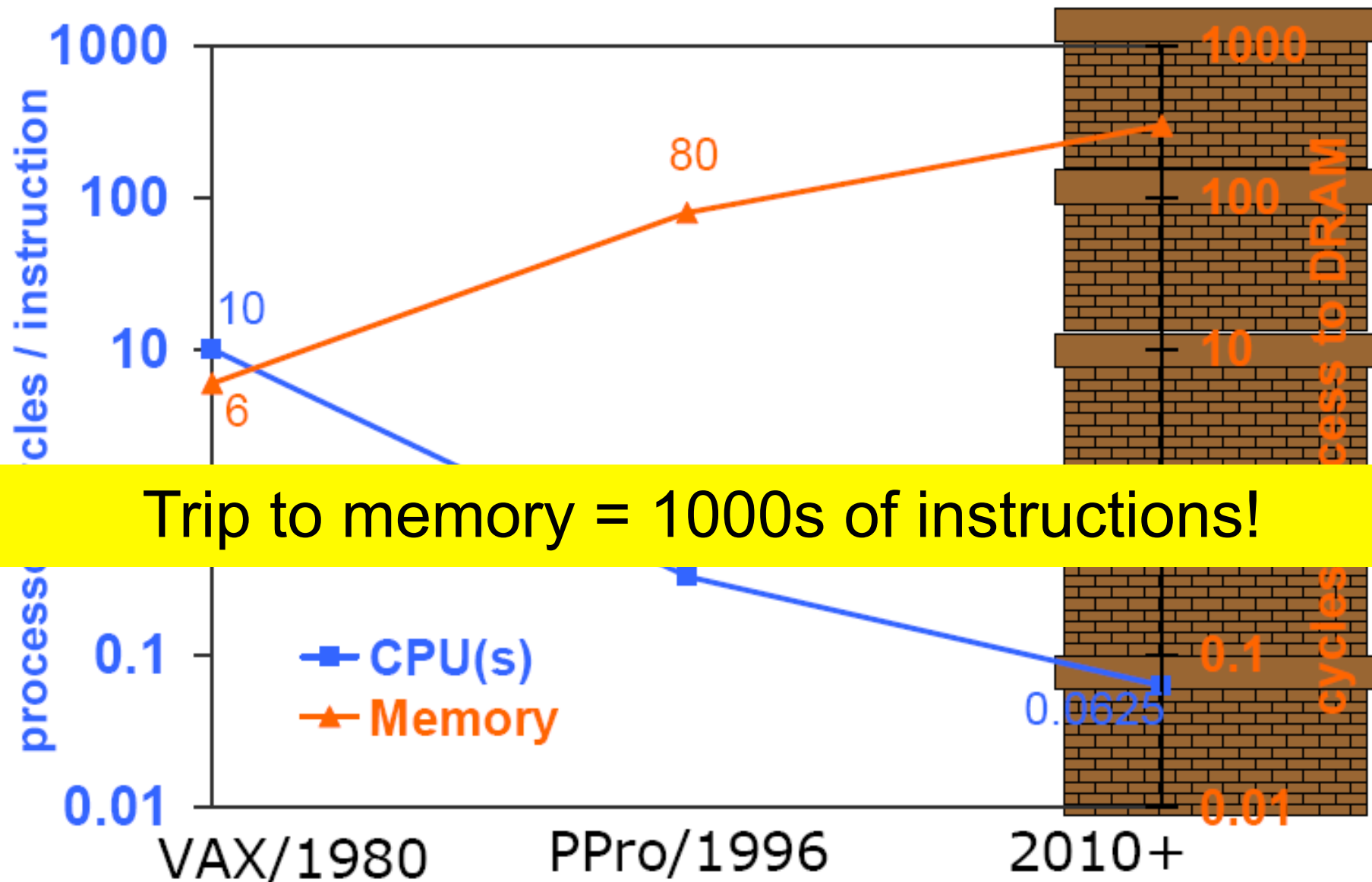
OLAP query:
accesses only a
few columns of
almost all rows.

Databases hit The Memory Wall

- Detailed and exhaustive analysis for different workloads using 4 RDBMSs by Ailamaki, DeWitt, Hill,, Wood in VLDB 1999: *"DBMSs On A Modern Processor: Where Does Time Go?"*
- CPU is 60%-90% idle, waiting for memory:
 - L1 data stalls
 - L1 instruction stalls
 - L2 data stalls
 - TLB stalls
 - Branch mispredictions
 - Resource stalls



Hardware Changes: The Memory Wall



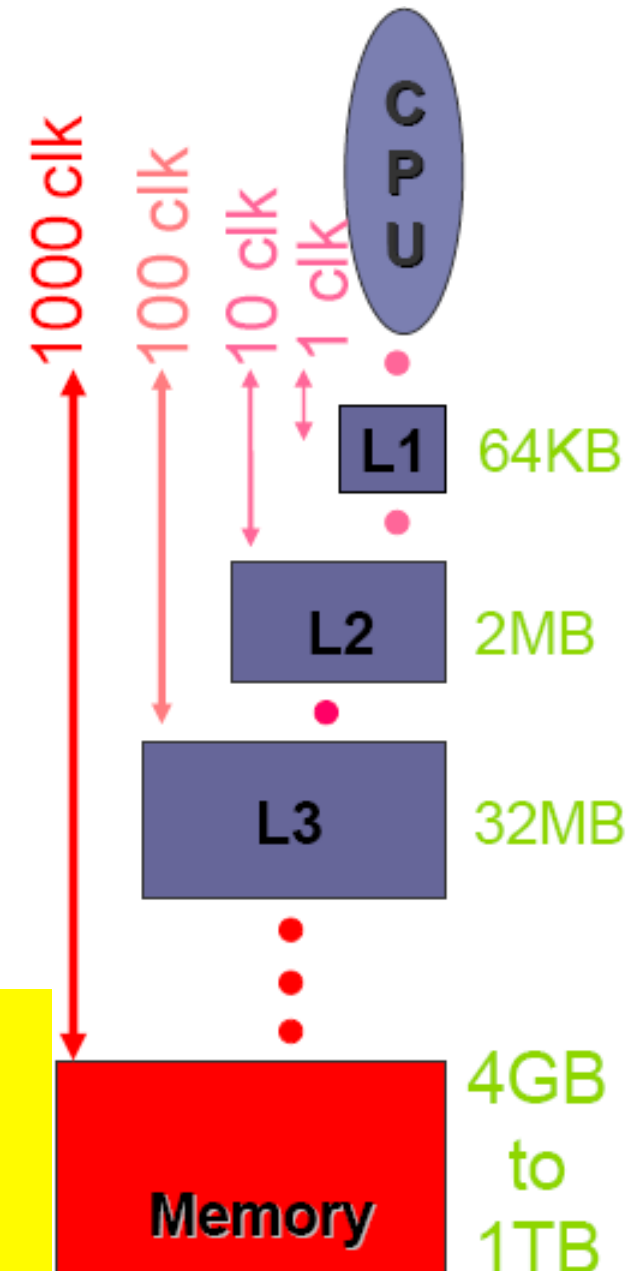
Hardware Changes: Memory Hierarchies

- Caches trade off capacity for speed
- Exploit instruction/data locality
- Demand fetch/wait for data

[ADH99]:

- Running top 4 database systems
- **At most 50% CPU utilization**

+Transition Lookaside Buffer (TLB)
Cache for VM address translation →
only 64 entries!



Evolution

It is not the strongest of the species that survives, nor the most intelligent, but the one most responsive to change.

Charles Darwin (1809 - 1882)

- Database kernel developed at CWI since 1993
 - *Research prototype turned into open-source product*
- Pioneering columnar database architecture
 - *Complete Relational/SQL & XML/XQuery DBMS*
- Design focus on large memory
 - *Data is kept persistent on disk and can exceed memory limits*
- Aiming at OLAP, BI & Data Mining workloads (“read-dominated”)
 - *Supporting ACID transactions (WAL, optimistic CC)*
- Platform for database architecture research
 - *Used in academia (research & teaching) & commercial environments*
- Back-end for various DB research projects:
 - Multi-Media DB & IR (“Tijah”), XML/XQuery (“Pathfinder”),
 - Data Mining (“Proximity”), Digital Forensics (“XIRAF”),
 - GIS (“OSM”), ...

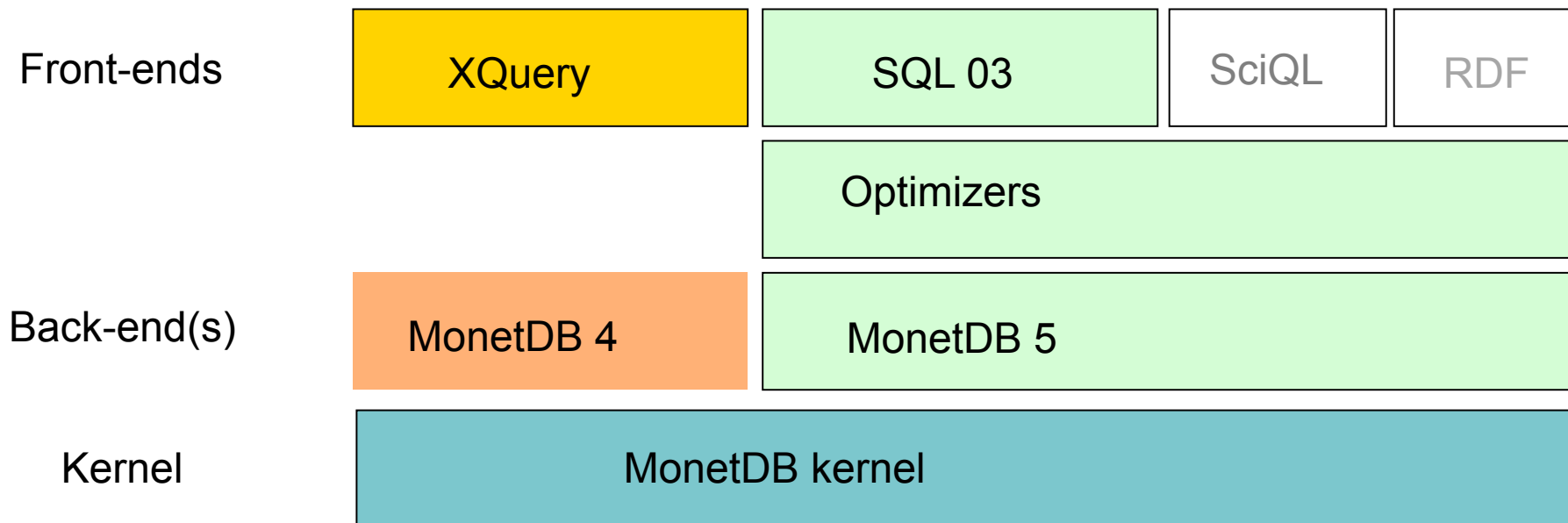
- full vertical fragmentation: always!
 - everything in binary (2-column) tables (**B**inary **A**ssociation **T**able)
 - saves you from table scan hell in OLAP and Data Mining
- RISC approach to databases
 - simple back-end data model
 - simple back-end query language (binary/columnar relational algebra)
 - don't need (to pay for) a buffer manager => manage virtual memory
 - explicit transaction management => DIY approach to ACID
 -
- Multiple user data models & query languages
 - SQL, XML/XQuery, SciQL, RDF/SPARQL
 - front-ends map data models and query languages

- optimized for large memory hierarchies
 - cache-conscious algorithms
 - exploit the persistence storage (disk, network, SSD)
- operator-at-a-time bulk processing
 - avoids tuple-at-a-time management overhead
- CPU and memory cache optimized
 - programming team experienced in main memory DBMS techniques
 - use of scientific programming optimizations (loop unrolling)

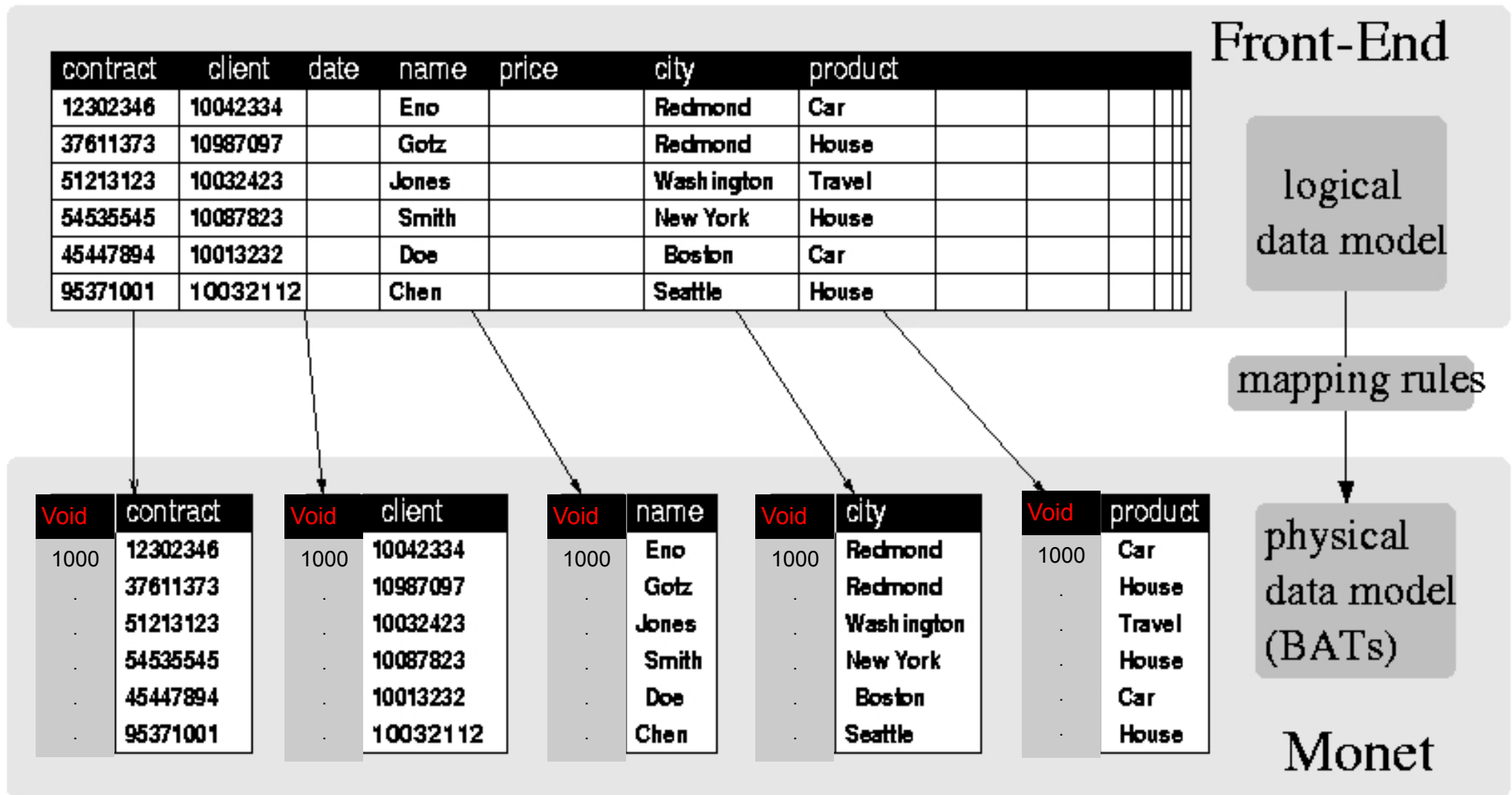
- **Architecture-Conscious Query Processing**
 - vs **Magnetic disk I/O conscious processing**
 - *Data layout, algorithms, cost models*
- **RISC Relational Algebra (operator-at-a-time)**
 - vs **Tuple-at-a-time Iterator Model**
 - *Faster through simplicity: no tuple expression interpreter*
- **Multi-Model: ODMG, SQL, XML/XQuery, ..., RDF/SPARQL**
 - vs **Relational with Bolt-on Subsystems**
 - *Columns as the building block for complex data structures*
- **Decoupling of Transactions from Execution/Buffering**
 - vs **ARIES integrated into Execution/Buffering/Indexing**
 - *ACID, but not ARIES.. Pay as you need transaction overhead.*
- **Run-Time Indexing and Query Optimization**
 - vs **Static DBA/Workload-driven Optimization & Indexing**
 - *Extensible Optimizer Framework;*
 - *cracking, recycling, sampling-based runtime optimization*



The MONETDB Software Stack

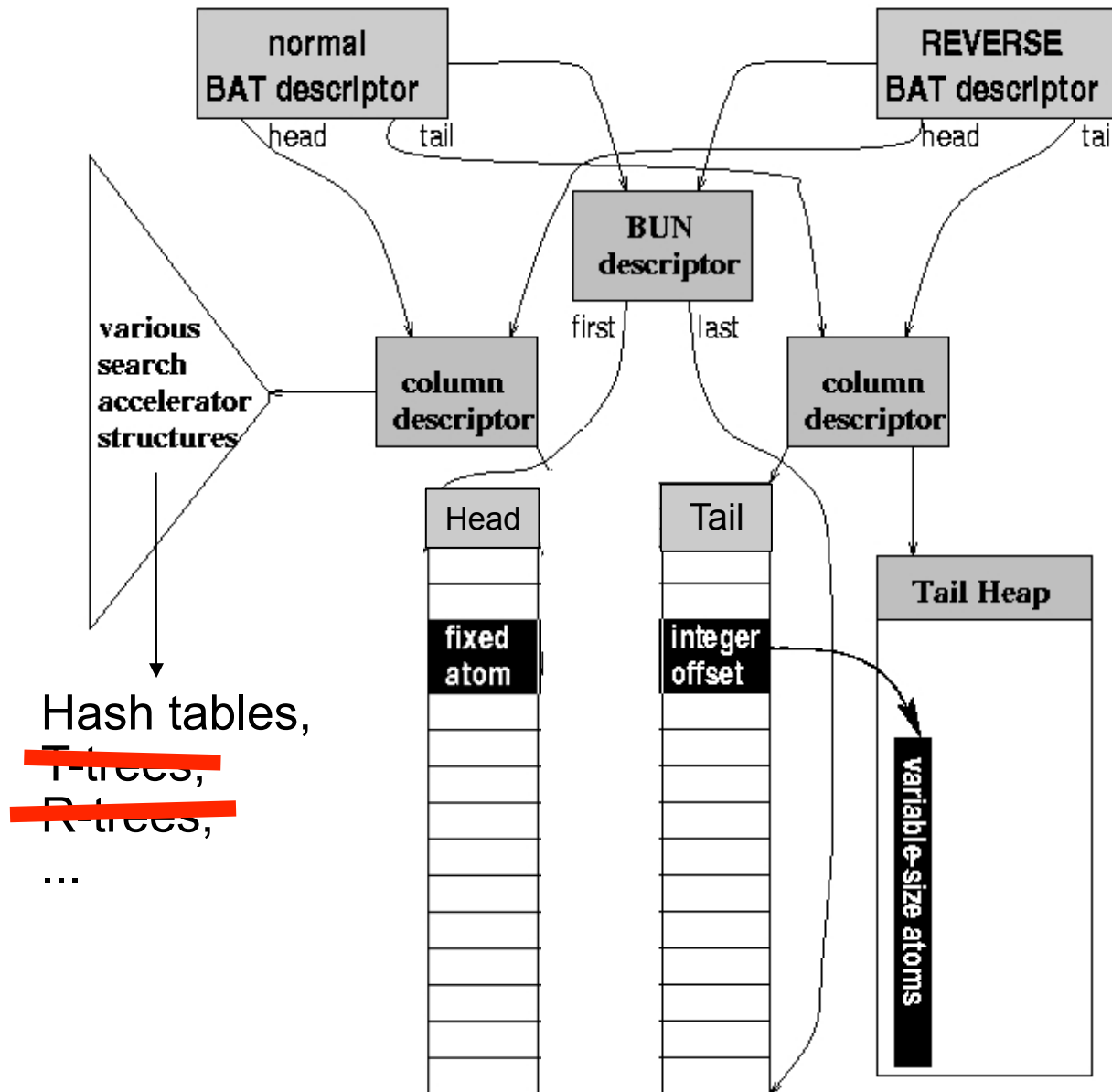


Storing Relations in MonetDB



Virtual **OID**: seqbase=1000 (increment=1)

BAT Data Structure



BAT:
binary association table

BUN:
binary unit

Head & Tail:

- consecutive memory blocks (arrays)
- memory-mapped files

Tail Heap:

- best-effort duplicate elimination for strings (~ dictionary encoding)

```
SELECT  id, name, (age-30)*50 as bonus
FROM    people
WHERE   age > 30
```

```
batcalc_minus_int(int* res,
                  int* col,
                  int val,
                  int n)
{
    for(i=0; i<n; i++)
        res[i] = col[i] - val;
}
```

CPU 😊? Give it “nice” code !

- few dependencies (control,data)
- CPU gets out-of-order execution
- compiler can e.g. generate SIMD

One loop for an entire column

- no per-tuple interpretation
- arrays: no record navigation
- better instruction cache locality

**Simple, hard-coded
semantics in operators**

**MATERIALIZED
intermediate results**

Processing Model (MonetDB Kernel)

- **Bulk processing:**
 - full materialization of all intermediate results
- **Binary (i.e., 2-column) algebra core:**
 - select, join, semijoin, outerjoin
 - union, intersection, diff (BAT-wise & column-wise)
 - group, count, max, min, sum, avg
 - reverse, mirror, mark
- **Runtime *operational optimization*:**
 - Choosing optimal algorithm & implementation according to input properties and system status

- Heavy use of code expansion to reduce cost

1 algebra operator

`select()`

3 overloaded operators

`select(“=”,value)`

`select(“between”,L,H)`
`(“fcn”,parm)`

`select`

10 operator algorithms

scan

hash-lookup

bin-search
lookup

bin-tree

pos-

~1500(!) routines

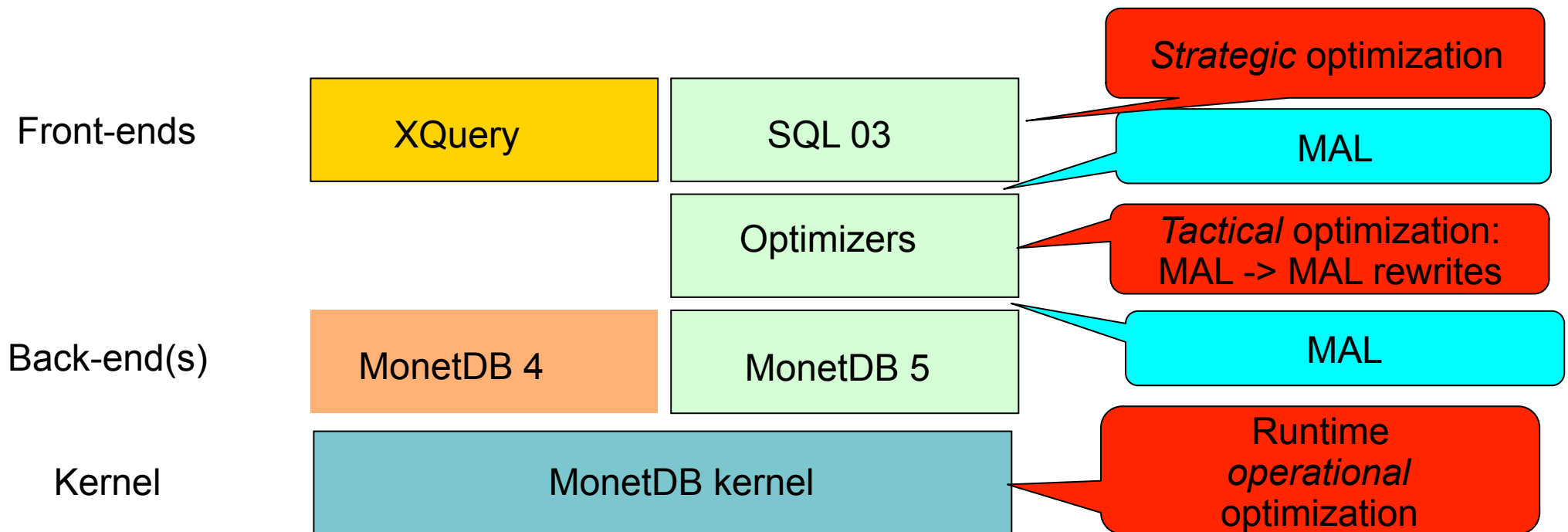
`scan_range_select_oid_int()`,
`hash_equi_select_void_str()`, ...

(macro expansion)

- ~1500 selection routines
- 149 unary operations
- 335 join/group operations
- ...



The MONETDB Software Stack



MonetDB/5 Back-end: MAL

- **MAL: Monet Assembly Language**
 - textual interface
 - Interpreted language
- **Designed as system interface language**
 - Reduced, concise syntax
 - Strict typing
 - Meant for automatic generation and parsing/rewriting/processing
 - Not meant to be typed by humans
- **Efficient parser**
 - Low overhead
 - Inherent support for *tactical optimization*: MAL -> MAL
 - Support for optimizer plug-ins
 - Support for runtime schedulers
- **Binary-algebra core**
- **Flow control** (MAL is computational complete)

```
EXPLAIN SELECT a, z FROM t, s WHERE t.c = s.x;
```

```
function user.s2_1():void;
barrier _73 := language.dataflow();
  _2:bat[:oid,:int] := sql.bind("sys","t","c",0);
  _7:bat[:oid,:int] := sql.bind("sys","s","x",0);
  _10 := bat.reverse(_7);
  _11 := algebra.join(_2,_10);
  _13 := algebra.markT(_11,0@0);
  _14 := bat.reverse(_13);
  _15:bat[:oid,:int] := sql.bind("sys","t","a",0);
  _17 := algebra.leftjoin(_14,_15);
  _18 := bat.reverse(_11);
  _19 := algebra.markT(_18,0@0);
  _20 := bat.reverse(_19);
  _21:bat[:oid,:int] := sql.bind("sys","s","z",0);
  _23 := algebra.leftjoin(_20,_21);
exit _73;
  _24 := sql.resultSet(2,1,_17);
  sql.rsColumn(_24,"sys.t","a","int",32,0,_17);
  sql.rsColumn(_24,"sys.s","z","int",32,0,_23);
  _33 := io.stdout();
  sql.exportResult(_33,_24);
end s2_1;
```

MonetDB: MAL Optimizers

- **General front-end independent MAL -> MAL rewriting**
 - Implemented once, shared by all (future) front-ends
- **Examples:**
 - Constant propagation
 - Scalar expression evaluation
 - Dead-code elimination
 - Common sub-expression elimination
 - Reordering to optimize intermediate result usage
 - Reordering of linear (projection-) join chains
 - **Parallelization:**
 - Dataflow analysis
 - Horizontal partitioning
 - Remote execution
 - *Cracking*
 - *Recycling*
 - ...

MonetDB Front-end: SQL

- SQL 2003
- Parse SQL into logical n-ary relational algebra tree
- Translate n-ary relational algebra into logical 2-ary relational algebra
- Turn logical 2-ary plan into physical 2-ary plan (MAL program)
 - Generate internal tree representation, not textual MAL program
- Front-end specific **strategic optimization**:
 - Heuristic optimization during all three previous steps
- Primary key and distinct constraints:
 - Create and maintain hash indices
- Foreign key constraints
 - Create and maintain foreign key join indices
- Exploit both above indices during query evaluation



Sloan Digital Sky Survey / SkyServer

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Welcome to the **DR6 site!!**

The Sixth Data Release is dedicated to **Jim Gray** for his fundamental contribution to the SDSS project and the extraordinary energy and passion he shared with everybody!

This website presents data from the Sloan Digital Sky Survey, a project to make a map of a large part of the universe. We would like to show you the beauty of the universe, and share with you our excitement as we build the largest map in the history of the world.

The site hosts data from **Data Release 6 (DR6)**. What's new in DR6, what's new on this site, and known problems. [More...](#)

For Astronomers

A separate branch of this website for professional astronomers (English)

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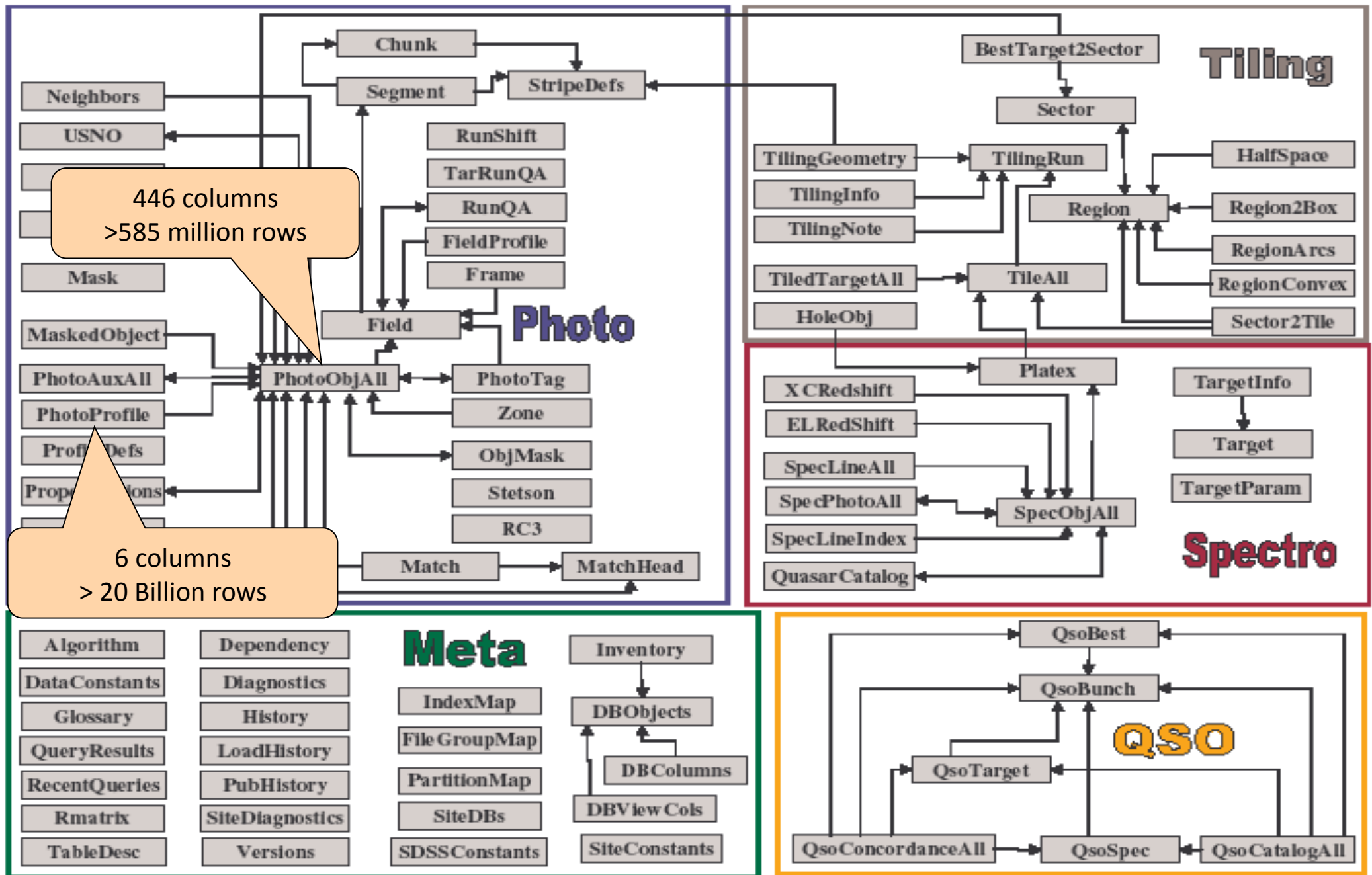
[Getting Started](#)[FAQ](#)[How To](#)[Glossary](#)[Schema Browser](#)[Sample SQL Queries](#)[Details of SDSS Data](#)

The contours for boundaries of the six different regions

Coordinates are represented as the equation of a 2D plane intersecting the unit sphere. These intersections are given in right ascension. The intersection is in terms of a specific RA value, which is the RA of the center of the SDSS field. In the left column, the number of the region is given. In the right column, the RA value is given. In the third column, the RA value is given. In the fourth column, the RA value is given. In the fifth column, the RA value is given. In the sixth column, the RA value is given.

Region	RA	Dec	RA	Dec	RA	Dec
1	0	0	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0
6	0	0	0	0	0	0

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Motivation:

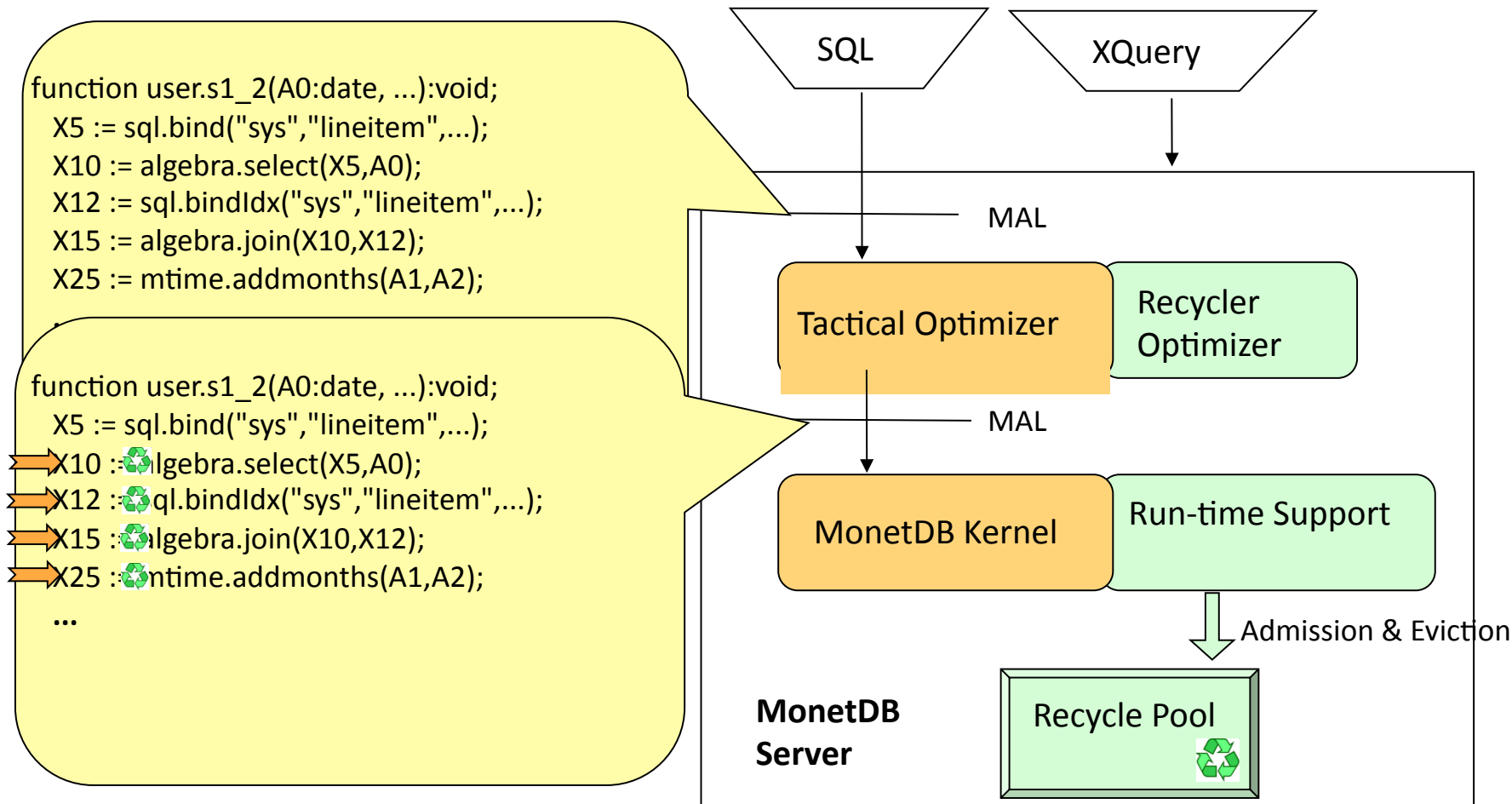
- scientific databases, data analytics
- Terabytes of data (observational , transactional)
- Prevailing read-only workload
- Ad-hoc queries with **commonalities**

Background:

- Operator-at-a-time execution paradigm
 - **Automatic materialization of intermediates**
- Canonical column-store organization
 - Intermediates have **reduced dimensionality** and **finer granularity**
 - **Simplified overlap analysis**

Recycling idea:

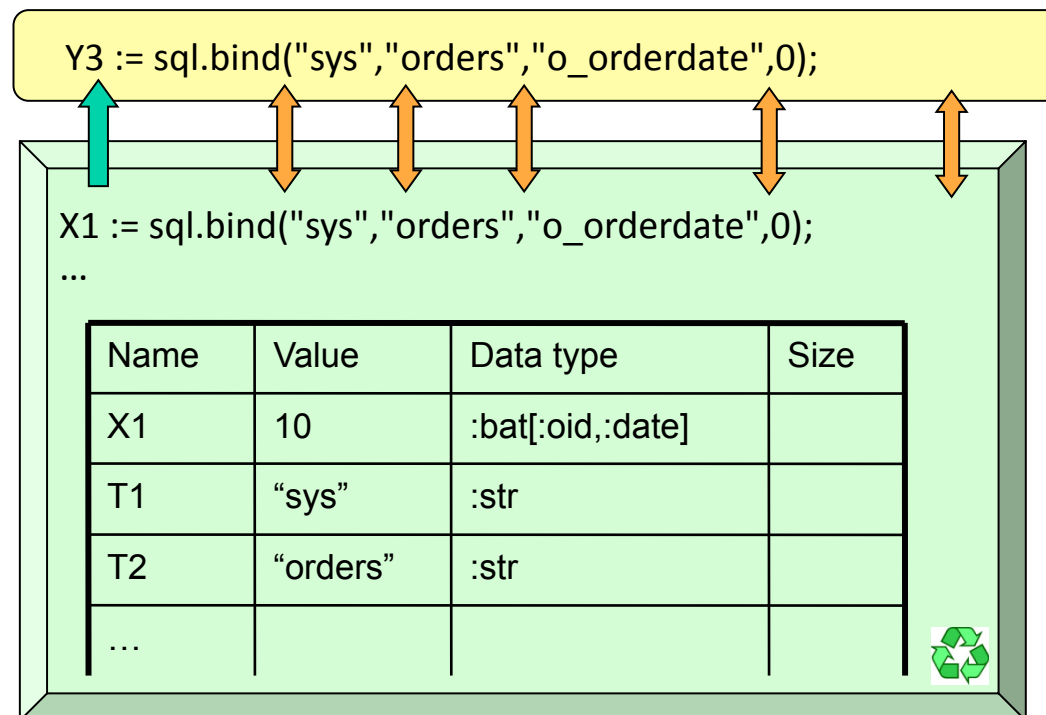
- instead of garbage collecting,
keep the intermediates and reuse them
 - speed up query streams with commonalities
 - low cost and self-organization



Run time comparison of

- instruction types
- argument values

Exact
matching



Y3 := algebra.select(X1,20,45);

X3 := algebra.select(X1,10,80);

...

X5 := algebra.select(X1,20,60);

Name	Value	Data type	Size
X1	10	:bat[:oid,:int]	2000
X3	130	:bat[:oid,:int]	700
X5	150	:bat[:oid,:int]	350
...			



Decide about storing the results

- KEEPALL
 - all instructions advised by the optimizer
- CREDIT
 - instructions supplied with credits
 - storage 'paid' with 1 credit
 - reuse returns credits
 - lack of reuse limits admission and resource claims

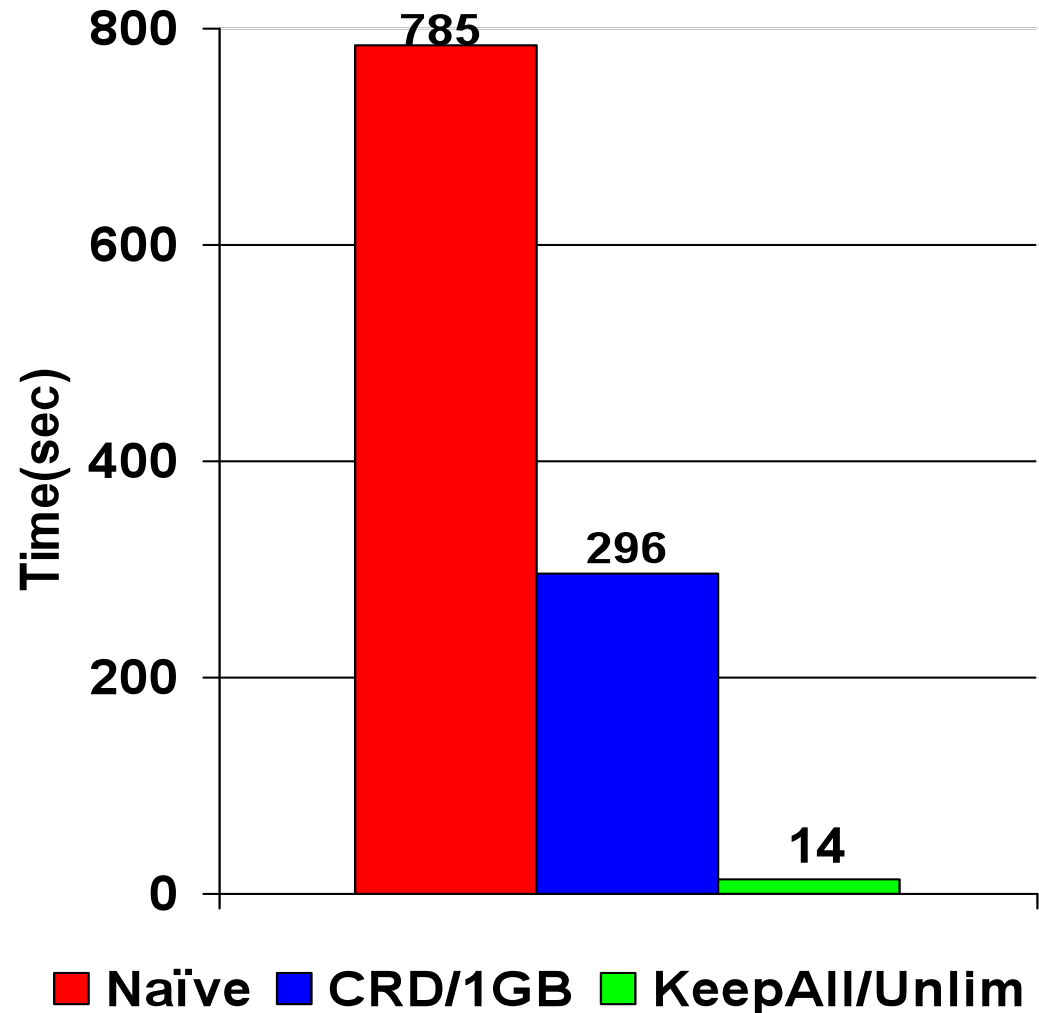
Decide about eviction of intermediates

- Pick instructions with smallest utility
 - LRU : time of computation or last reuse
 - BENEFIT : estimated contribution to performance:
CPU and I/O costs, recycling
- Triggered by resource limitations (memory or entries)

Sloan Digital Sky Survey / SkyServer

<http://cas.sdss.org>

- 100 GB subset of DR4
- 100-query batch from January 2008 log
- 1.5GB intermediates, 99% reuse
- Join intermediates major consumer of memory and major contributor to savings



Project portfolio

Commercial: 0.5 PB telco warehouse

Commercial: DNA warehouses

LOFAR : Transient detection

Emili: Streaming in sensor-based systems

TELEIOS: Remote sensing virtual observatory

Planetdata,Lod2: Semantic web, linked open data

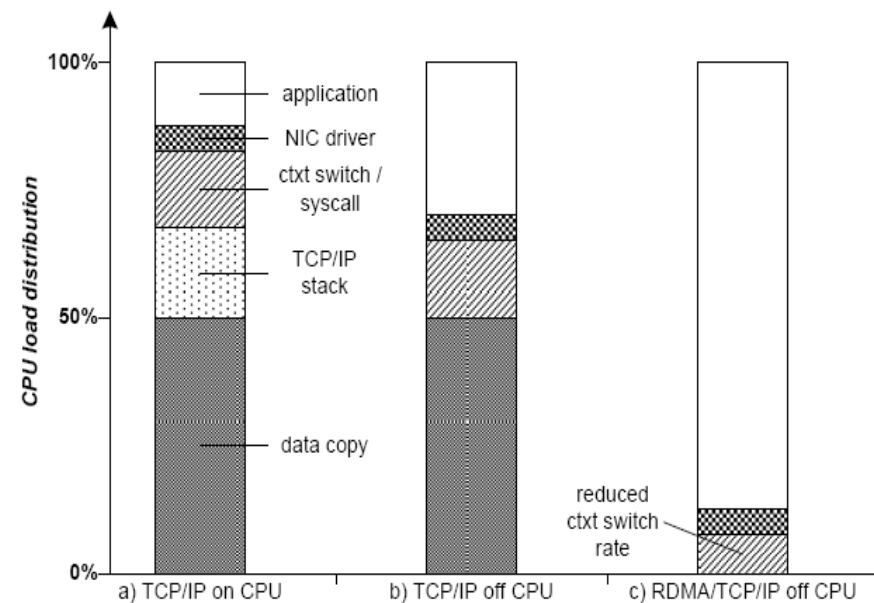
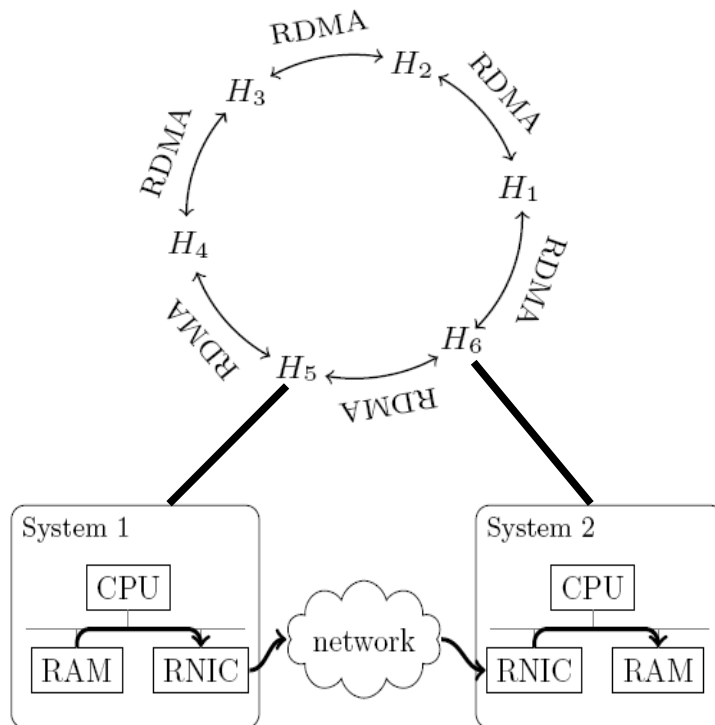
NWO: Biased sampling for science database

SciLens: Dissemination and coordination

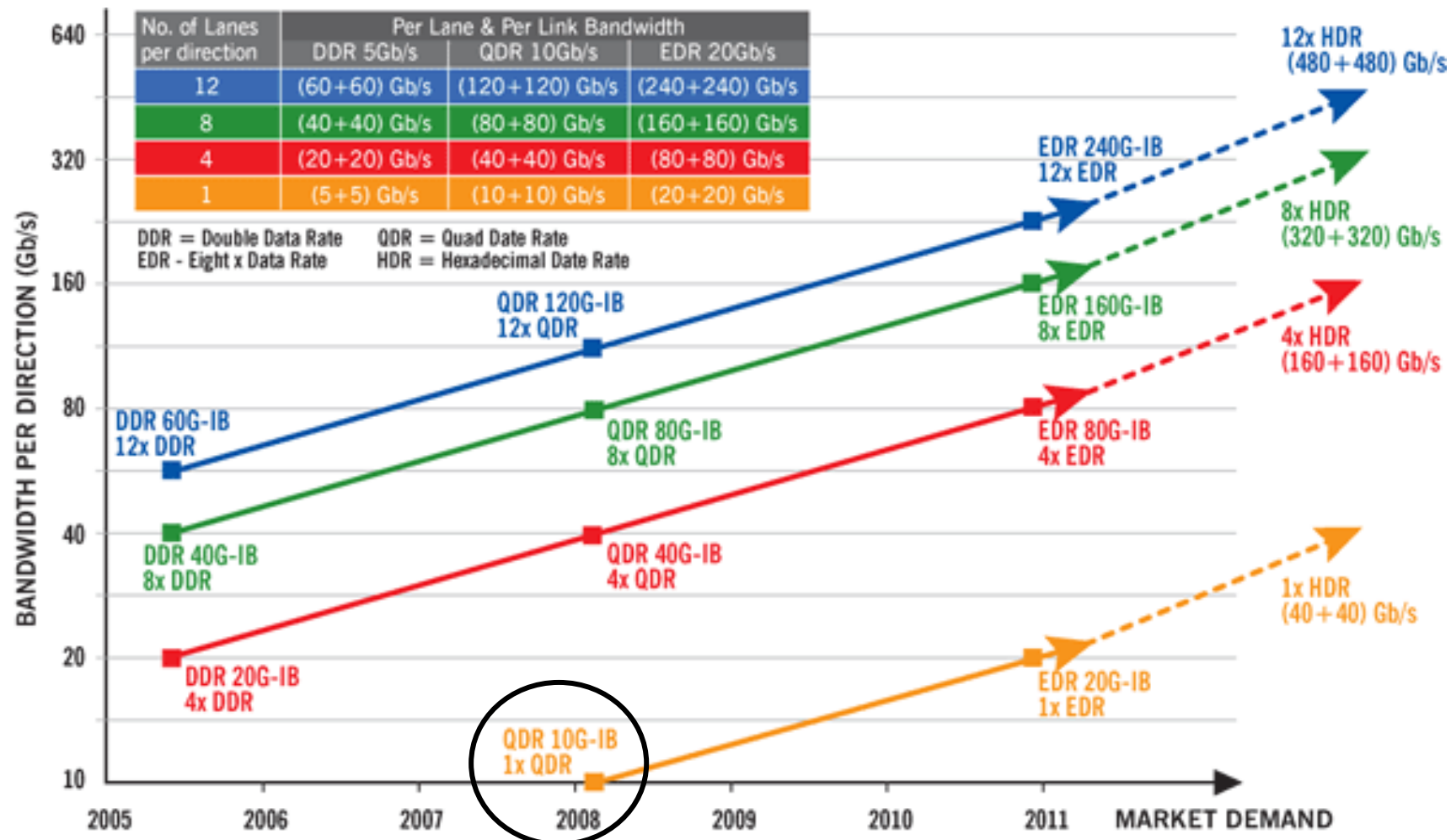
Datacyclotron: Novel distributed architectures

Remote Direct Memory Access (RDMA)

- Remote Memory at Your Finger Tips.
- RDMA Benefits.
 - Cpu Load
 - Reduced Memory Bus Traffic

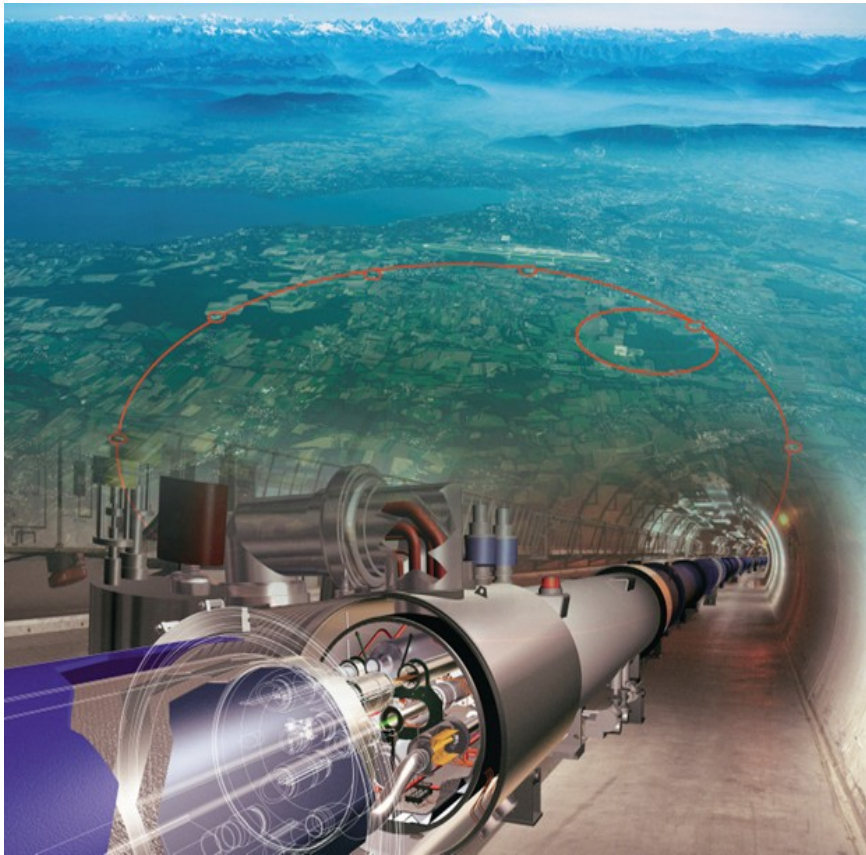


Road-map for RDMA



The topology.

- Swiss one (LHC)

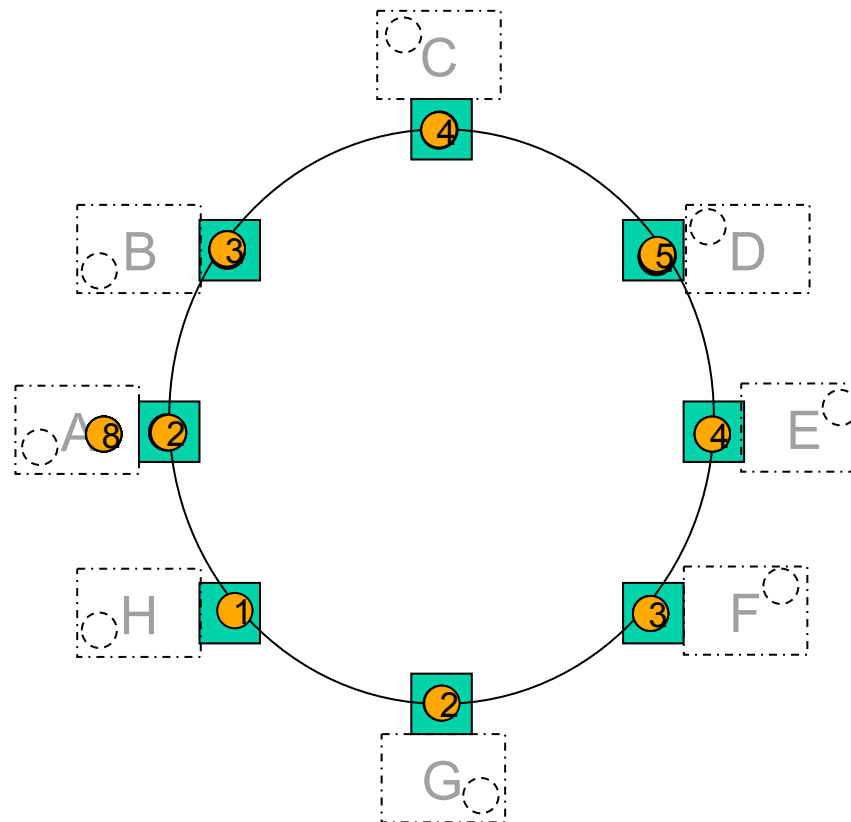


- Dutch one (DaCy)



The data acceleration.

- A chunk is loaded by a node into the ring.
- It flows clockwise...
- It continuously hops from node to node... until it is removed...



MonetDB SciQL

SciQL (*pronounced 'cycle'*)

- A backward compatible extension of SQL'03
- Symbiosis of relational and array paradigm
- Flexible structure-based grouping
- Capitalizes the MonetDB array storage
 - Recycling, an adaptive 'materialized view'
 - Zero-cost attachment contract for cooperative clients

MonetDB Vaults

A contract between MonetDB and file repository of voluminous scientific data

- provide seamless SQL access to foreign file formats using SciQL views
- zero cost, adaptive loading and replication
- Capitalize libraries as UDFs (linpack, R,..)
- Short term targets:
 - MSEED, FITS, NETCDF, csv

MonetDB Octopus

- Distributed SQL processing without a DBA
- Merovingian, managing a cluster of servers
- Cloud-based infrastructure with fail-over
- Partial/full replication adaptive to query needs.
- Recycling, a basis for distribution and load scheduling

eScience- landscape

Science Domain Workbench

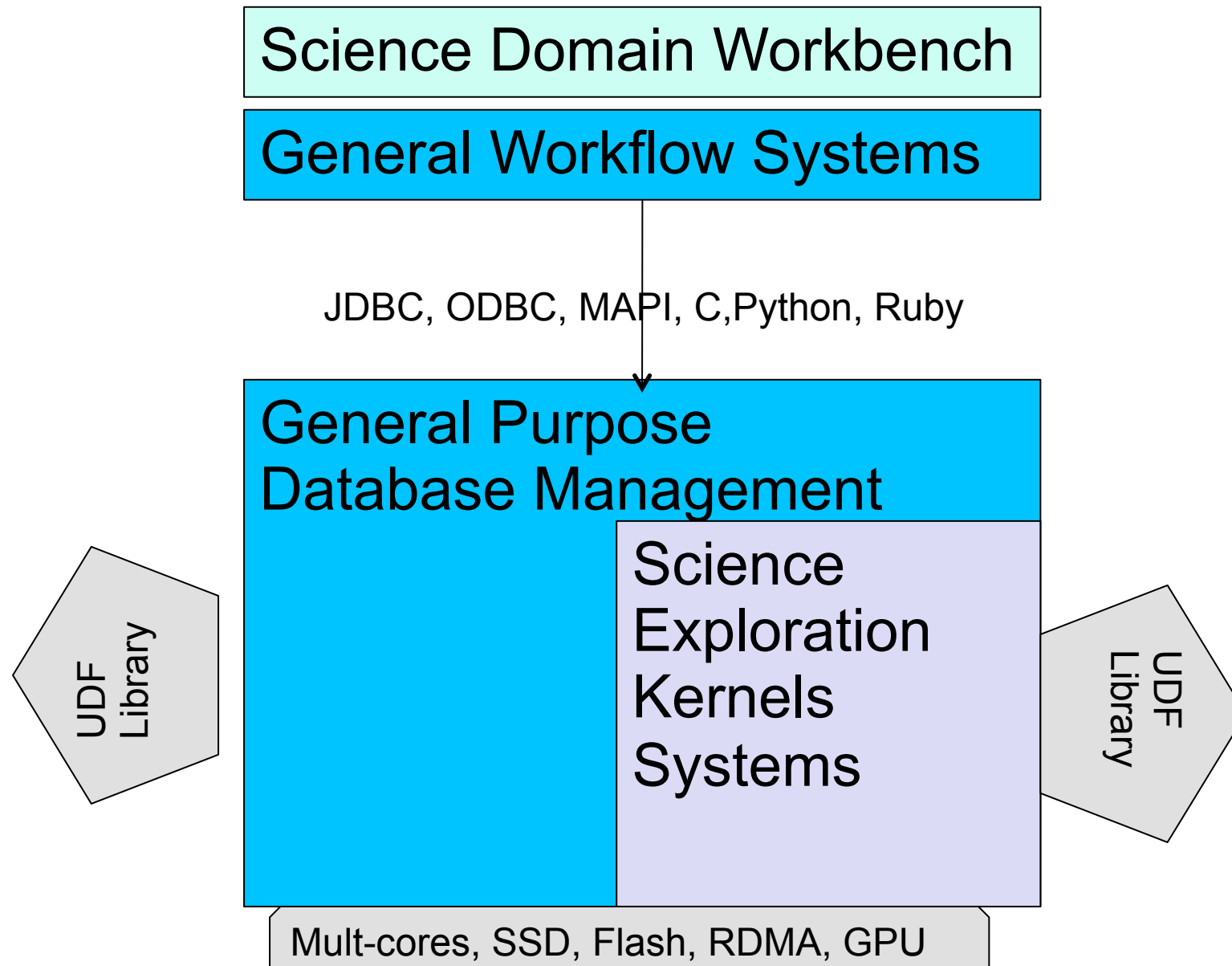
General Workflow Systems

General Purpose
Database Management

Data acquisition systems
Data scrubbing, cleaning
Data refinement, enrichment
Data catalogues, meta-data
Data exploration, mining
Data visualisation

Science
Exploration
Database
Systems

MonetDB for Science



	MonetDB 5.23	SciDB 0.5
Open source	Mozilla License	GPL 3.0 + Commercial
Downloads	>12.000 /month	Tens up to now
SQL compliance	SQL 2003	??
Interoperability	JDBC, ODBC, MAPI, C, Python, Ruby, C++	C++ UDF
Array model	SciQL	AQL
Science support	Linked libraries	Linked libraries
Foreign files	Vaults to FITS, NETCDF, MSEED	??
Distribution	50 node cluster Octopus 200 node cluster Cyclotron	4 node cluster
Distribution tech	Dynamic partial replication	Static fragmentation
	Distributed query, map-reduce, streaming, multi-core	Map-reduce
Largest local demo	Skyserver SDSS 6 3TB	---

Open-Source Development

- Feature releases: 3-4 per year
 - Research results
 - User requests
- Bug-fix releases: monthly
- QA
 - Automated nightly testing on >20 platforms
 - Ensure correctness & stability
 - Ensure portability
 - Bug reports become test cases
 - Semi-automatic performance monitoring
 - Passed static code verification by Coverity with only minor problems

Below: Complete list showing all tests. (Alternative: Shortened list showing only tests with unexpected outcome on at least one platform.

 2 °C
  6 °C
 Sun Jan 17, 16:28:32

MonetDB,

full-functional open-source product

a mature modern column-store

proven track record in (science) applications

strong and committed development team

close interaction with application developers

Reaching the stars step by step