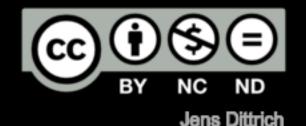
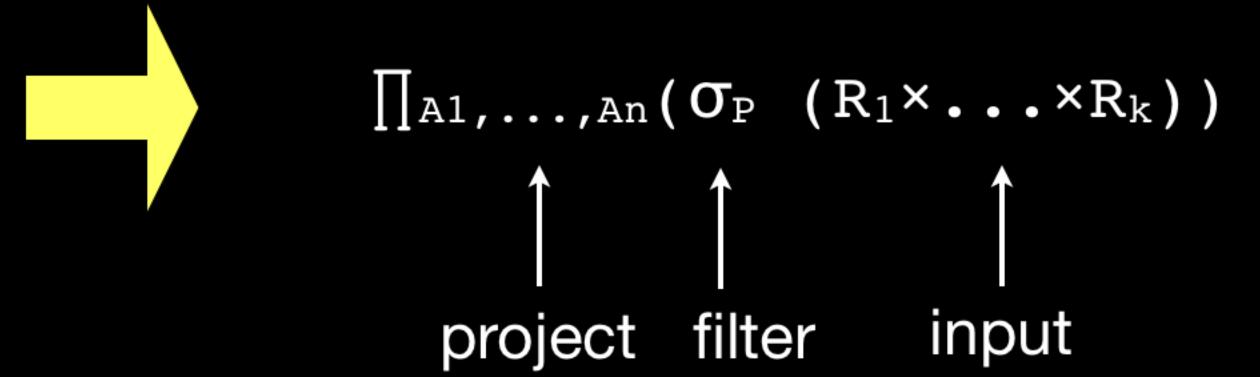
# Query Parser



### Input:

SELECT A<sub>1</sub>, ..., A<sub>n</sub>
FROM R<sub>1</sub>, ..., R<sub>k</sub>
WHERE P;

#### query parser



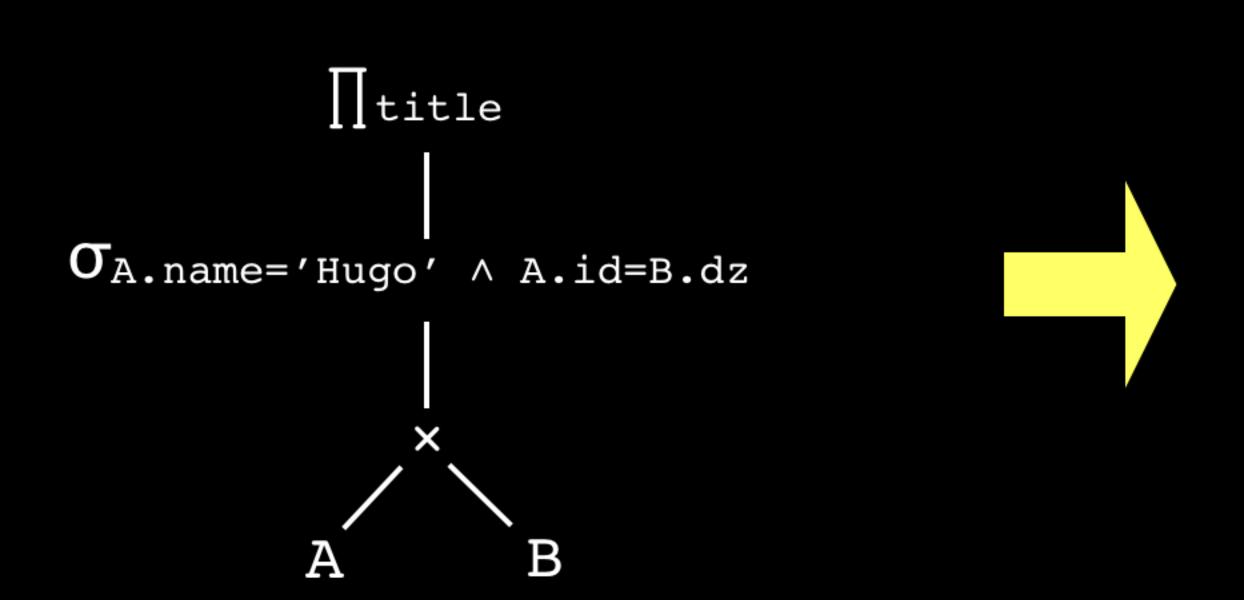
# Canonical Form of a Query

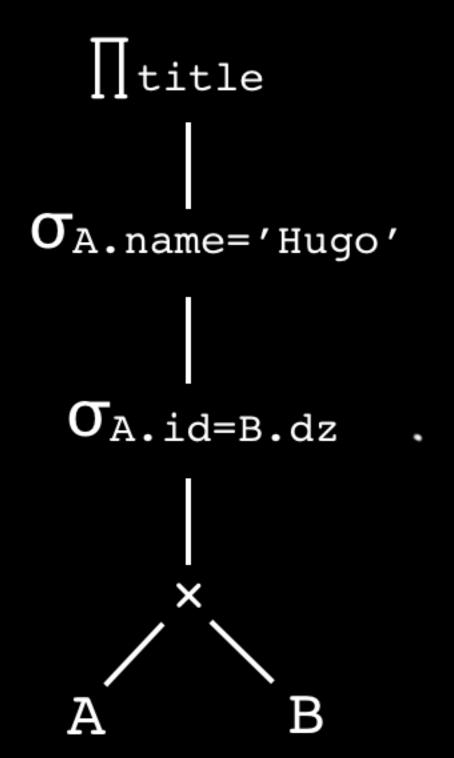
Input:  $\prod_{A1,\ldots,An}$  $\sigma_{\mathrm{P}}$  $\prod_{A1,\ldots,An} (\sigma_P (R_1 \times \ldots \times R_k))$  $R_{\mathbf{k}}$  $R_3$ DAG, e.g. selt-voits, Lested grevier

## Example for Both Transformation Steps

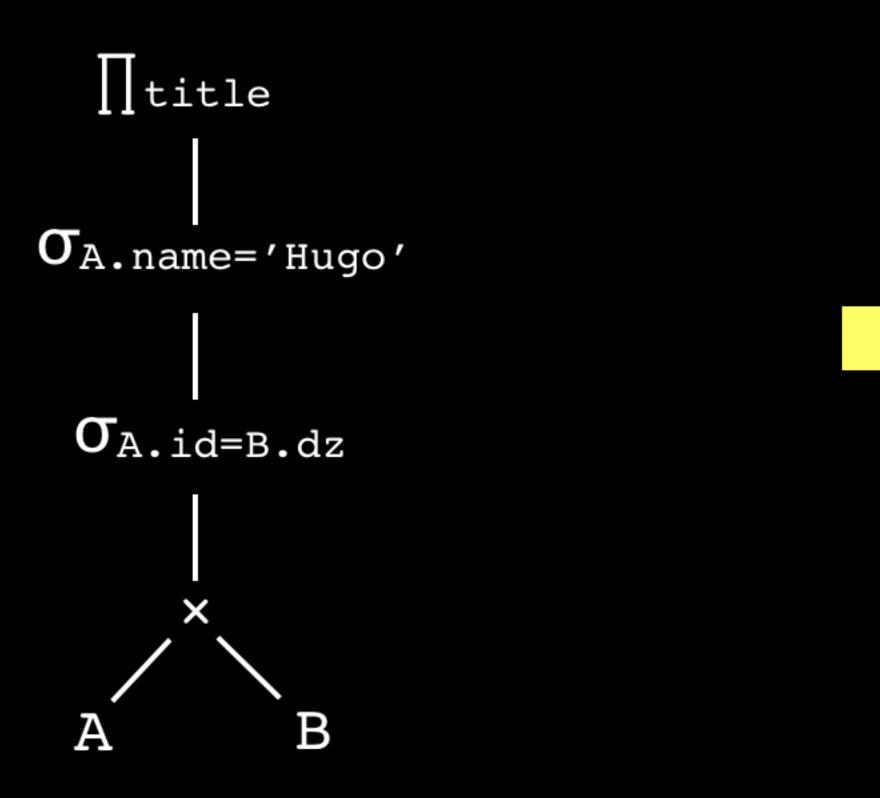
```
SELECT title
               FROM
                         A,B
                        A.name = 'Hugo' AND A.id = B.dz;
               WHERE
                      Ititle (\sigma_{A.name='Hugo'} \wedge A.id=B.dz (A×B))
                                       ∏title
Rule 42 (Nzy Lzy)
                              OA.name='Hugo' ∧ A.id=B.dz
rule-bæed optimization
(Leuristin)
                                               В
```

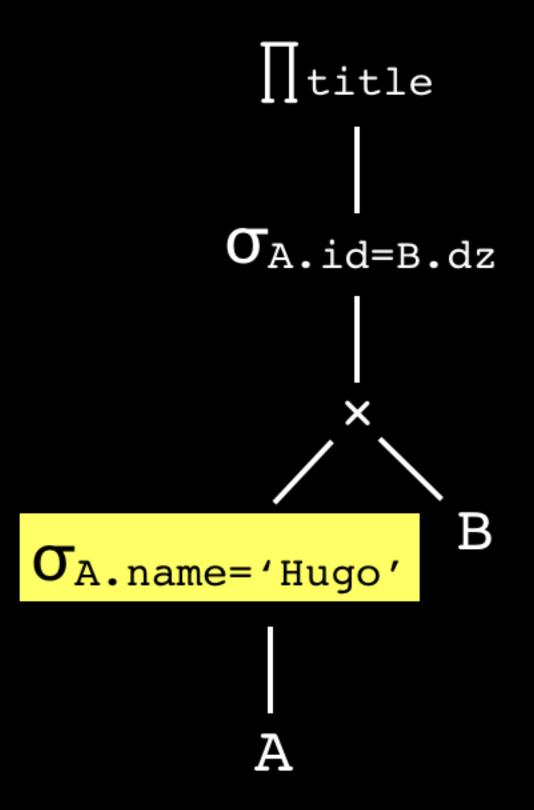
# Break Up Predicates



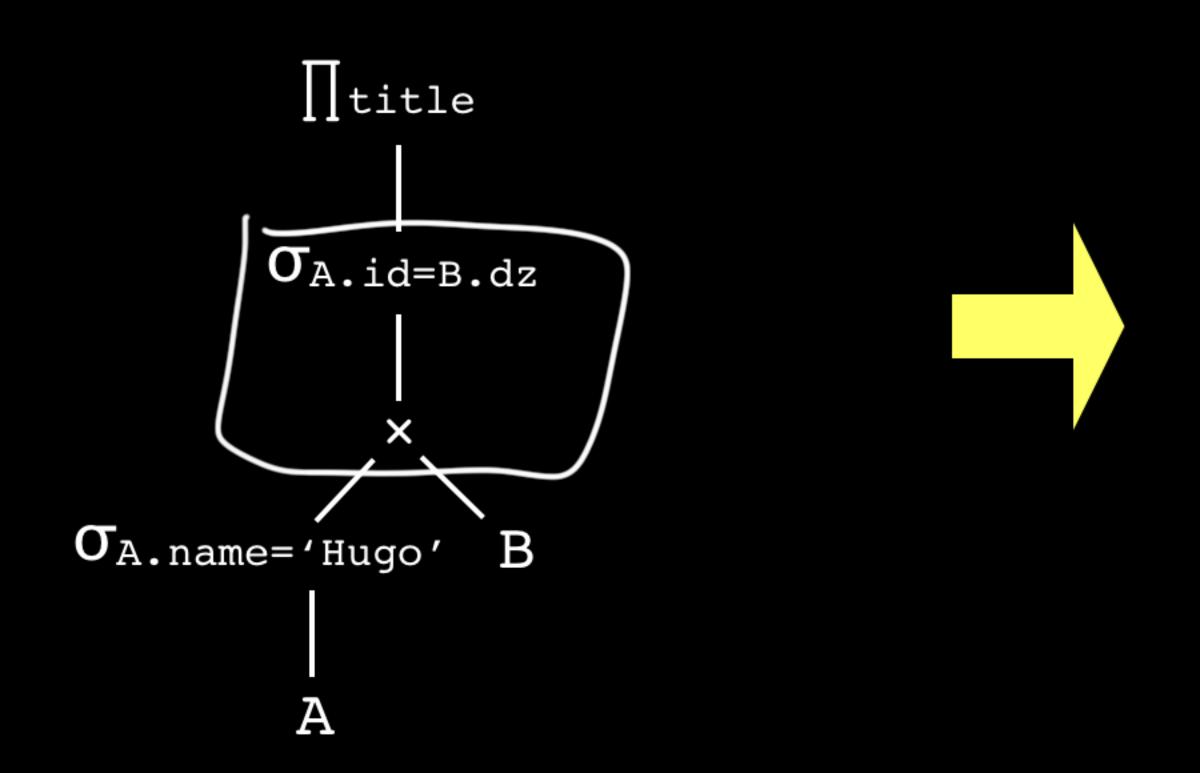


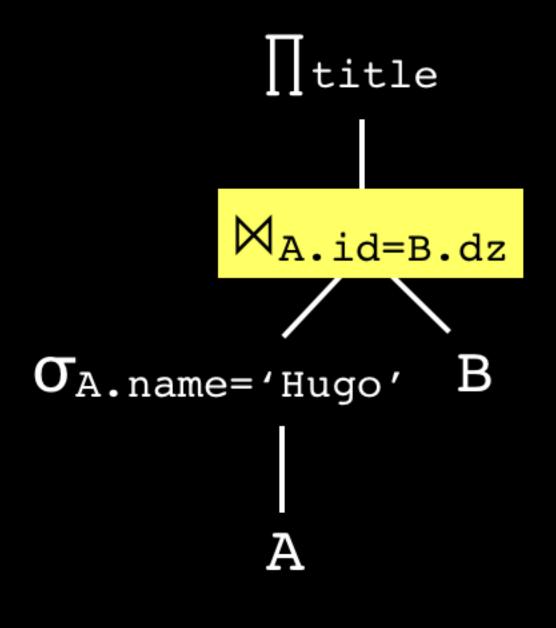
# Push Down Selections



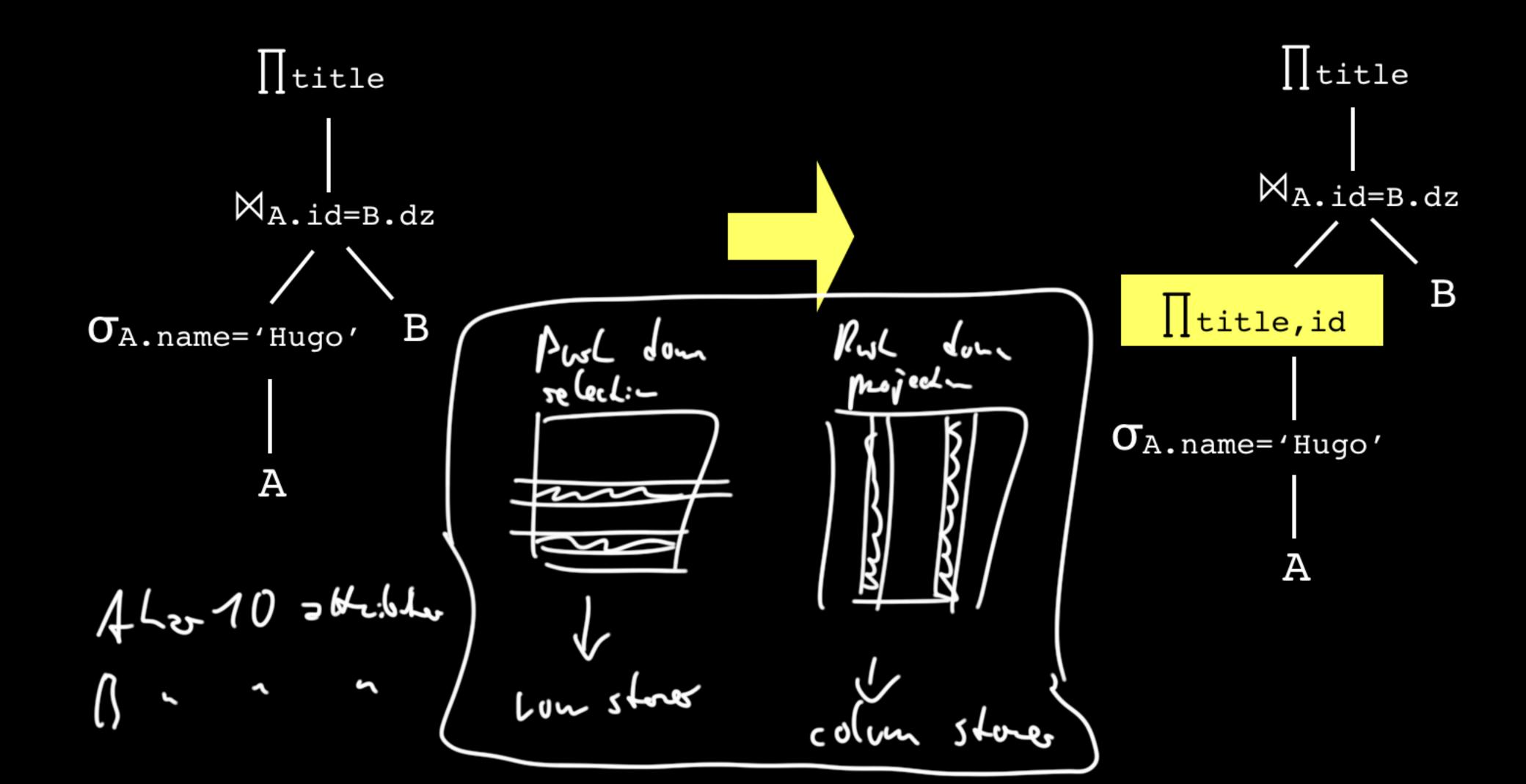


# **Build Joins**





# Push Down Projections



## Most Important Rules

- 1. push down selections and projections
- 2. combine selections and cross products into joins
- 3. insert additional projections