# Geometry Morphing

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## The objective

The objective of this internship is to morph some tetrahedral meshes preserving some properties (for now, the volume). We want to achieve this in two ways:

- Using (semi-discrete) Optimal Transport
- Using Variational Autoencoders

#### Some notation

Let A a tetrahedral mesh, composed by N tetrahedrals, then  $A_i$  will denote the volume matrix of the i-th tetrahedra:

$$a_1 - d_1$$
  $b_1 - d_1$   $c_1 - d_1$   
 $a_2 - d_2$   $b_2 - d_2$   $c_2 - d_2$   
 $a_3 - d_3$   $b_3 - d_3$   $c_3 - d_3$ 

where a,b,c,d are the points of the i-th tetrahedral. It is called volume matrix because the volume of the tetrahedra is  $\frac{|\det(A_i)|}{6}$ , by swapping a,b,c,d we can assume that  $\det(A_i)>0$ . The volume of the mesh so can be calculated with  $\frac{\sum_{i=1}^N \det(A_i)}{6}$ 

### Optimal Transport map I

Given  $\Omega$  a Borel set and two measures  $\mu$  and  $\nu$  on  $\Omega$  such that  $\mu(\Omega) = \nu(\Omega)$ , c a convex function.  $T:\Omega \to \Omega$  such that  $\begin{cases} \nu(X) = \mu(T^{-1}(X)) \text{for any Borel (i.e. measurable) subset } X \text{ of } \Omega \\ \int_{\Omega} c(x,T(x)) d\mu \text{ is minimal is called the optimal transport map from } \mu \text{ to } \nu. \text{ We are interested in } \mu \text{ continuos and } \nu \text{ discrete.} \end{cases}$ 

## Optimal Transport map II

The Voronoi diagram Vor (P) is the partition of  $\mathbb{R}^d$  into the subsets Vor  $(p_i)$  defined by :

Vor  $(p_i) := \{x | ||x - p_i||^2 < ||x - p_j||^2 \quad \forall j \neq i \}$  the power diagram  $\text{Pow}_W(P)$  is the partition of  $\mathbb{R}^d$  into the subsets  $\text{Pow}_W(p_i)$  defined by:

 $\operatorname{Pow}_W(p_i) := \left\{ x \mid \|x - p_i\|^2 - w_i < \|x - p_j\|^2 - w_j \quad \forall j \neq i \right\}$  the map  $T_W$  defined by  $\forall i, \forall p \in \operatorname{Pow}_W(p_i), T_W(p) = p_i$  is called the assignment defined by the power diagram  $\operatorname{Pow}_W(P)$  and is an optimal transport map.

Knowing this we can approximate the optimal transport map in the following algorithm:

Data: Two tetrahedral meshes M and M', and k the desired number of vertices in the result

Result: A tetrahedral mesh G with k vertices and a pair of points  $p_i^0$  and  $p_i^1$  attached to each vertex. Transport is parameterized by time  $t \in [0,1]$  with  $p_i(t) = (1-t)p_i^0 + tp_i^1$ .

- (1) Sample M' with a set Y of k points
- (2) Compute the weight vector W that realizes the optimal transport between M and Y
- (3) Compute  $E = \text{Del}(Y) \mid M' \text{ and } F = \text{Pow}_W(Y) \mid M$  $\text{Tets}(G) \leftarrow E \cap F$  where Del is the Delaunay triangulation.
- (4) Foreach  $i \in [1 ... k], (p_i)^0 \leftarrow \operatorname{centroid}(\operatorname{Pow}_W(y_i) \cap M)$ ,  $(p_i)^1 \leftarrow y_i$

### **Problems**

There are two main problems with this approach:

- Step 3 may modify substantially the two meshes, so they may end up having different volumes (this however can be resolved by using more tetrahedras and rescaling at the end)
- Let A and B the two meshes obtained with  $p_i^0$  and  $p_i^1$ , then the map f(t) = t \* A + (1 t) \* B will preserve volume with only very restrictive assumptions.

#### $\mathsf{Theorem}$

The volume is preserved for every  $t \in [0,1]$  iff  $\sum_{i=1}^{M} det(B_i)(Tr(B_i^{-1}A_i) - 3) = 0$  and  $\sum_{i=1}^{M} det(A_i)(Tr(A_i^{-1}B_i) - 3) = 0$ 

A weaker version also holds

#### Theorem

$$|Vol(t*A+(1-t)*B)-Vol(B)| \leq \frac{4}{27}|\sum_{i=1}^{M} det(B_i)(Tr(B_i^{-1}A_i)-det(A_i)Tr(A_i^{-1}B_i))| + \frac{1}{2}|\sum_{i=1}^{M} det(A_i)(Tr(A_i^{-1}B_i)-3)|$$