

# Lecture 8

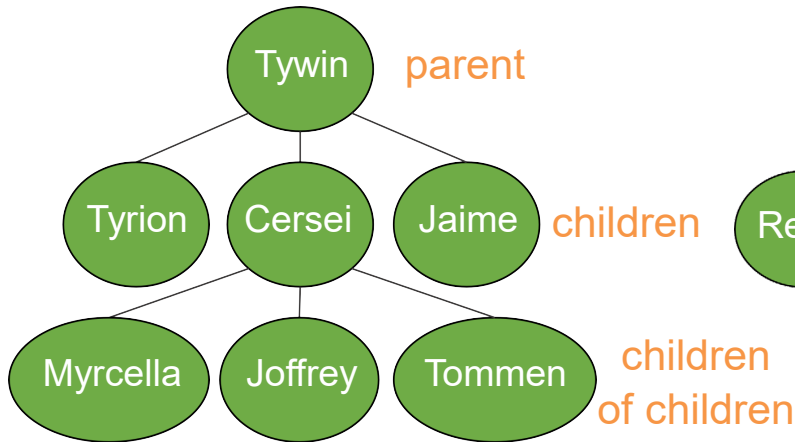
## Binary Search Tree and Trie

Department of Computer Science  
Hofstra University

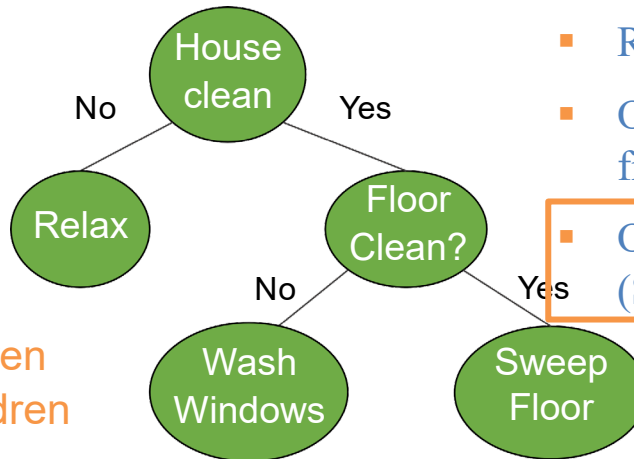
# Lecture Goals

- Describe the **value** of trees and their data structure
- Explain the need to visit data in different **orderings**
- Perform pre-order, in-order, post-order and level-order **traversals**
- Define a **Binary Search Tree**
- Perform **search, insert, delete** in a Binary Search Tree
- Explain the running time **performance** to find an item in a BST
- Compare the **performance** of linked lists and BSTs
- Explain what a **trie** data structure is
- Describe the algorithm for **finding** keys in and **adding** keys to a trie
- **Compare** the time to find a key in a BST to a trie
- **Implement** a trie data structure in Java

# Different Trees in Computer Science



Family Trees

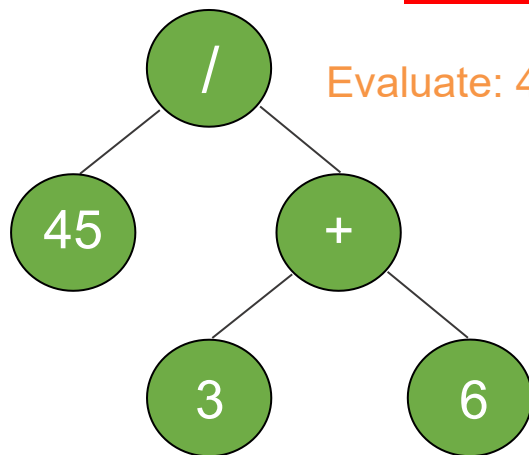


Decision Trees

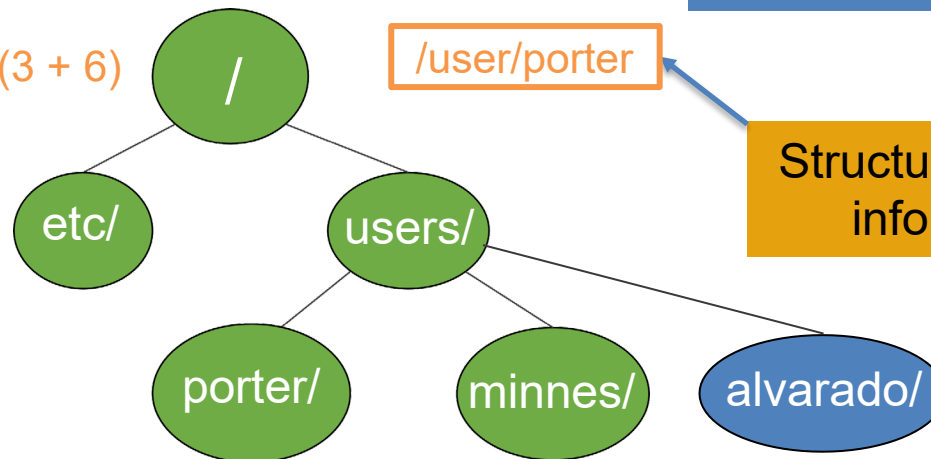
- Root is most important (Heap)
- Organized by character frequency (Huffman Tree)
- Organized by node ordering (Search Trees)
- Etc...

Why trees?

Different Organizations  
→ Different Trees



Expression Trees

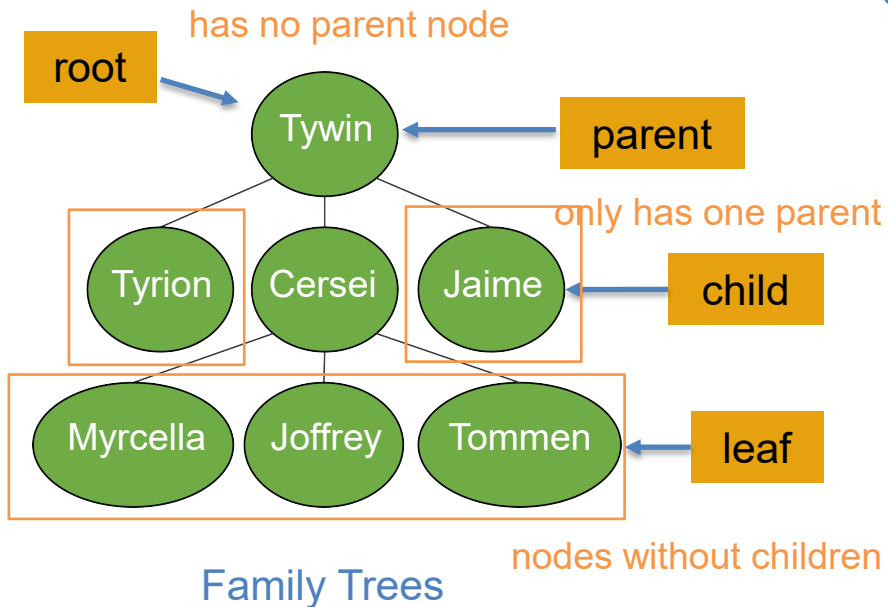


File System

Structure conveys  
information

Dynamic Data Structure

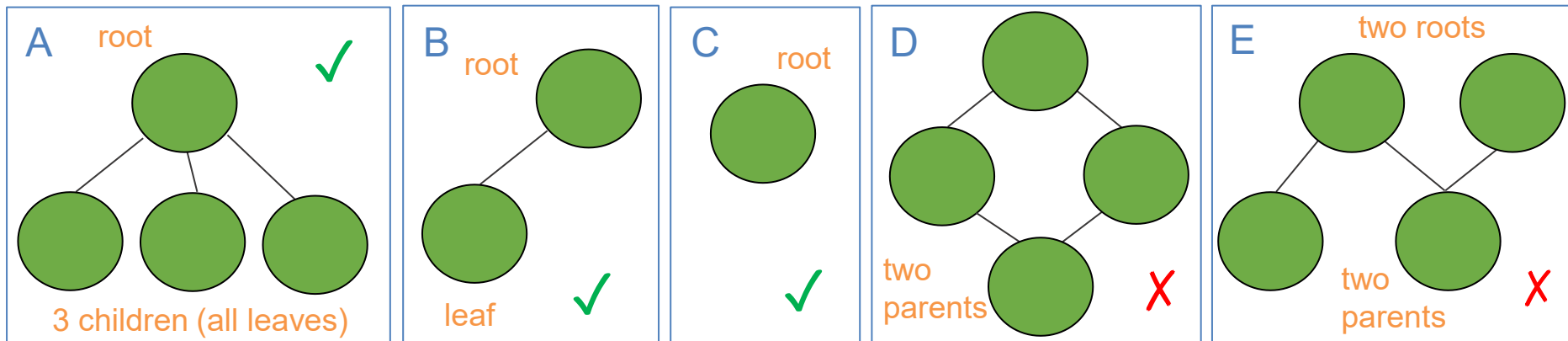
# Defining Trees



What defines a tree?

- Single root
- Each node can have only one parent (except for root)
- No cycles in a tree

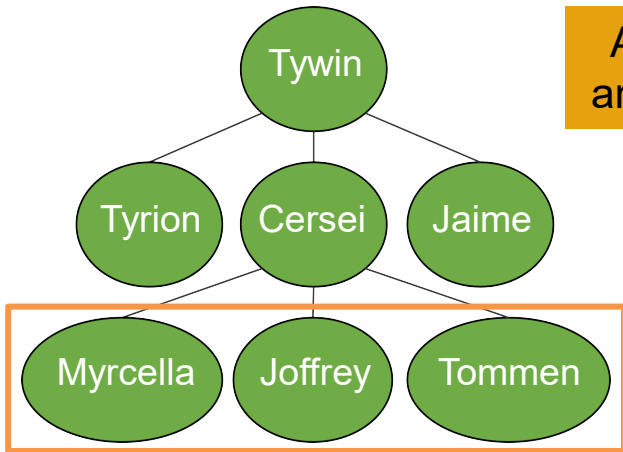
Which are trees?



Cycle: two different paths  
between a pair of nodes

# Binary Trees

## Generic Tree



Any Parent can have any number of children

How would a general tree node differ?

A general tree would just have a list for children

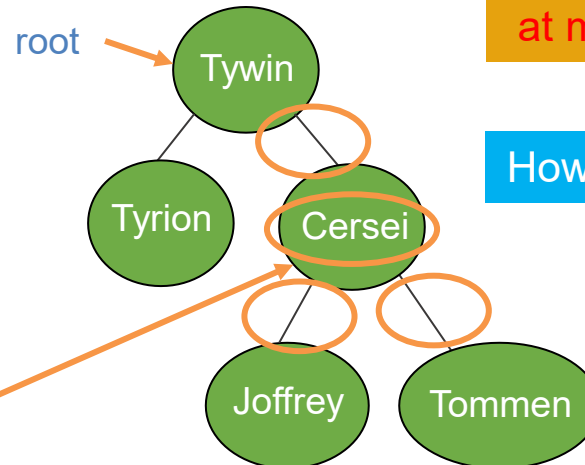
A tree just needs a root node

like the head and tail for linked list

Each node needs:

1. A value
2. A parent
3. A left child
4. A right child

## Binary Tree



Any Parent can have **at most** two children

How do we construct a tree?

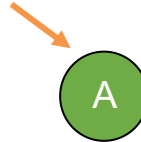
Like Linked Lists, Trees have a "Linked Structure"

nodes are connected by references

# Write Code for Binary Tree

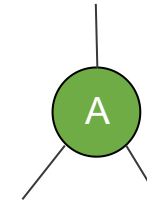
```
public class BinaryTree<E> {  
    TreeNode<E> root;  
    // more methods  
}
```

root



```
public class TreeNode<E> {  
    private E value;  
    private TreeNode<E> parent;  
    private TreeNode<E> left;  
    private TreeNode<E> right;  
    public TreeNode(E val, TreeNode<E> par) {  
        this.value = val;  
        this.parent = par;  
        this.left = null;  
        this.right = null;  
    }  
    public TreeNode<E> addLeftChild(E val) {  
        this.left = new TreeNode<E>(val, this);  
        return this.left;  
    }  
}
```

For root: `TreeNode(val, null)`



Let's write a constructor together

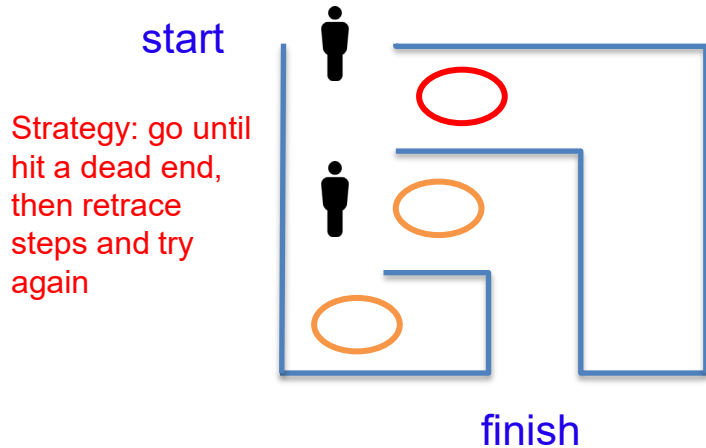
Next Step is to be able to set/get children

Fill in the blank:

- A. `this.parent`
- B. `this.left`
- C. `this.right`
- D. `this`

# Tree Traversal - Motivation

Warning: These first examples are really graphs. We'll visit graphs in detail in the next course. Here they are used as motivating examples



## Maze Traversal

Suppose you have a list of your friends and each of your friends have lists

How closely are you connected with D?

What's my next step?

Strategy: look at all of your friends first, and then branch out.

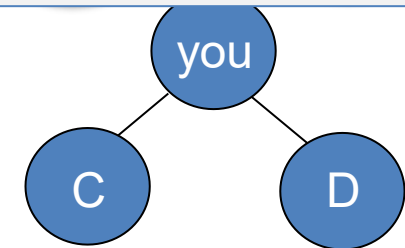
This problem benefits from "Breadth First Traversals"

Imagine this is a hedge maze

What's my next step?

Mazes benefit from "Depth First Traversals"

Bottom line: Order we visit matters and we'll make choices based on our needs

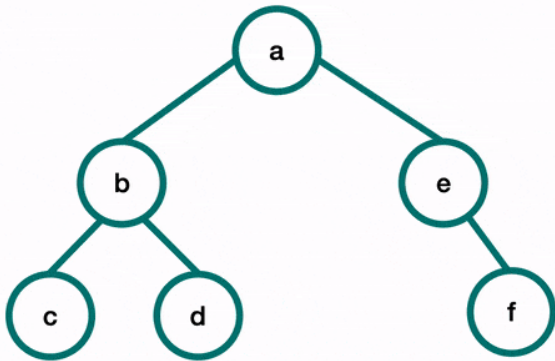


## Social Network

# Graph traversal with DFS: in-order, pre-order, post-order

```
function inOrderTraversal(node) {  
  if (node !== null) {  
    inOrderTraversal(node.left);  
    visitNode(node);  
    inOrderTraversal(node.right);  
  }  
}
```

**In-Order Traversal**

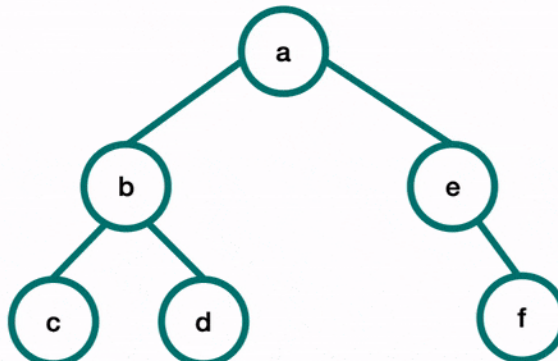


Print ""

cbdaef

```
function preOrderTraversal(node) {  
  if (node !== null) {  
    visitNode(node);  
    preOrderTraversal(node.left);  
    preOrderTraversal(node.right);  
  }  
}
```

**Pre-Order Traversal**

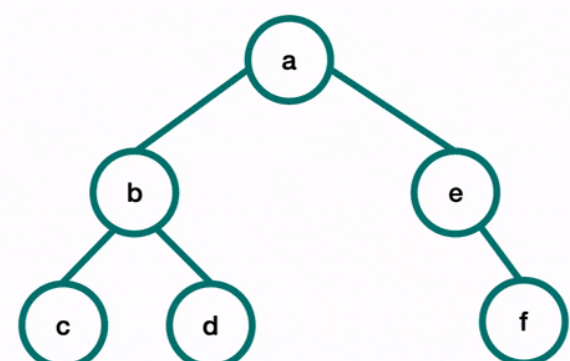


Print ""

abcdef

```
function postOrderTraversal(node) {  
  if (node !== null) {  
    postOrderTraversal(node.left);  
    postOrderTraversal(node.right);  
    visitNode(node);  
  }  
}
```

**Post-Order Traversal**



Print ""

cdbfea

<https://skilled.dev/course/tree-traversal-in-order-pre-order-post-order>



# Motivation for Binary Search Tree

Agra	Beijing	Chicago	Essen	Lagos	Montreal	Quito
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Binary Search -  $O(\log n)$  search:  
get rid of half each time

toFind

Chicago

Sorted arrays are good for search,  
but bad for insertion/removal

root

Essen

Beijing

Montreal

Agra

Chicago

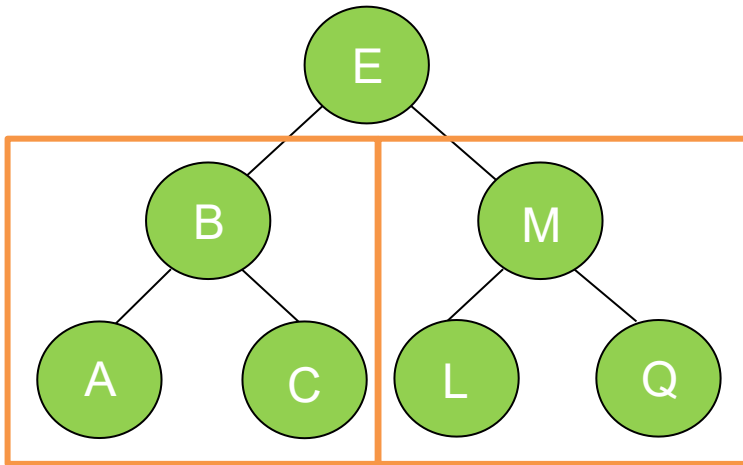
Lagos

Quito

So now we can do the same kind of fast searching we did within an array, but we can also get the benefit of a fast insert and a fast removal that a tree provides.

Binary Search Trees (BST) Explained in Animated Demo  
<https://www.youtube.com/watch?v=mtvbVLK5xDQ>

# Defining a Binary Search Tree



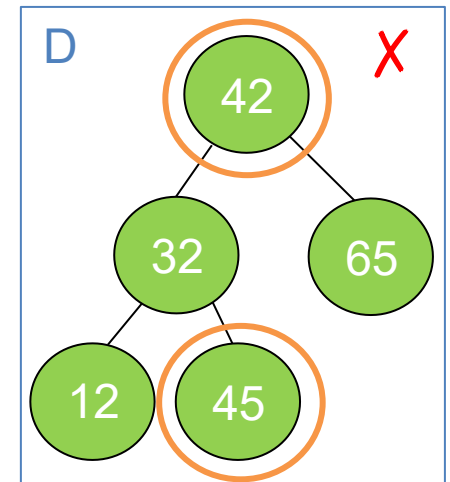
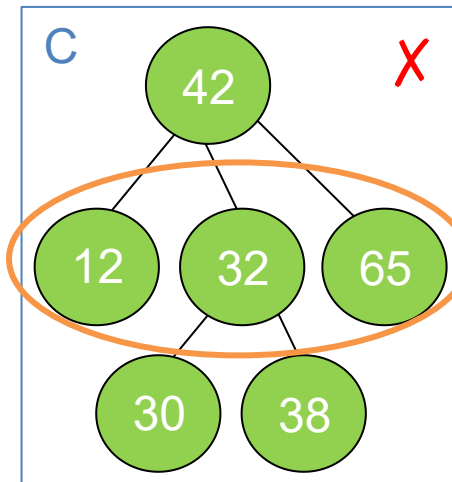
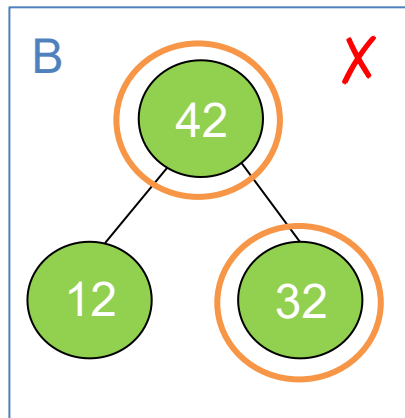
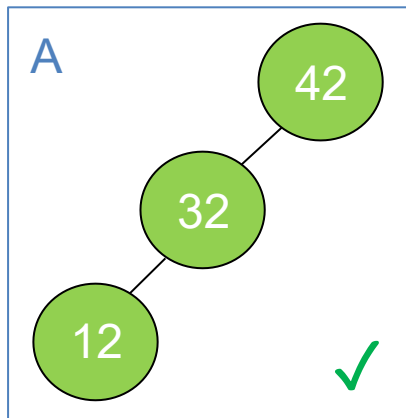
Left subtree's values  
must be lesser

Right subtree's values  
must be greater

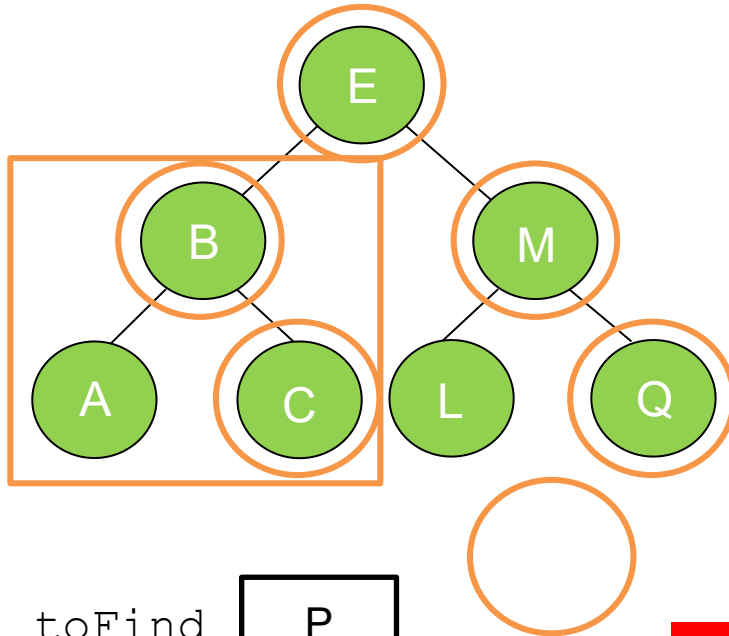
Binary Search Tree:

1. Binary Tree
2. Left subtrees are less than parent
3. Right subtrees are greater than parent

Which of these are binary search trees?



# Searching a BST



toFind

P

Compare: E and P

Compare: M and P

Compare: Q and P

Node is null

Not Found!

Same fundamental idea as binary search of an array

toFind

C

Compare: E and C

Compare: B and C

Compare: C and C

Found it!

How to implement this?

You could solve this with **recursion**.

You could also solve it with **iteration** by keeping track of your current node.

# Searching a BST Iteratively

```
public class BinaryTree<E> {  
    <E extends Comparable<? super E>> {  
        TreeNode<E> root;  
        public boolean search(E toSearch) {  
            TreeNode<E> curr = root;  
            while (curr != null) {  
                int comp = toSearch.compareTo(curr.getValue());  
                if (comp < 0)  
                    curr = curr.getLeftChild();  
                else if (comp > 0)  
                    curr = curr.getRightChild();  
                else // comp = 0  
                    return true;  
            }  
            return false;  
        }  
    }  
}
```

Do NOT change root pointer!

It means that either the class E itself or one of its super classes implements Comparable

Doesn't work with objects

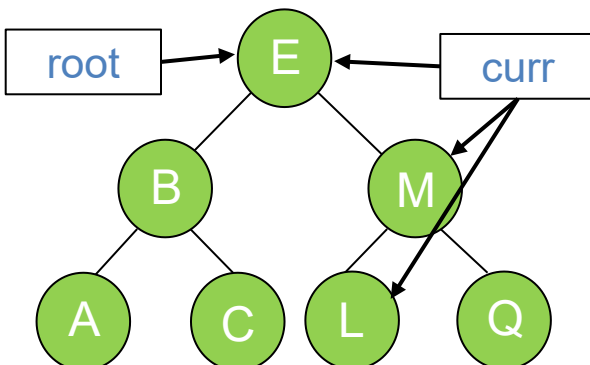
We need to do this over and over if not found

if calling object is less than parameter, compareTo returns a value < 0

if calling object is greater than parameter, compareTo returns a value > 0

if calling object is equal to parameter, compareTo returns 0

Are we done?



t.search('L')

Traverse down tree until:

- a) end is reached
- b) element is found

# Searching a BST Recursively

```
public class BinaryTree<E extends Comparable<? super E>> {  
    TreeNode<E> root;
```

Root of the tree we look at

```
    private boolean search(TreeNode<E> p, E toSearch) {
```

```
        if (p == null)
```

```
            return false;
```

Tree is empty

```
        int comp = toSearch.compareTo(p.getValue());
```

```
        if (comp == 0)
```

```
            return true;
```

Found it!

```
        else if (comp < 0)
```

```
            return search(p.left, toSearch);
```

look left

```
        else // comp > 0
```

```
            return search(p.right, toSearch);
```

look right

```
    }
```

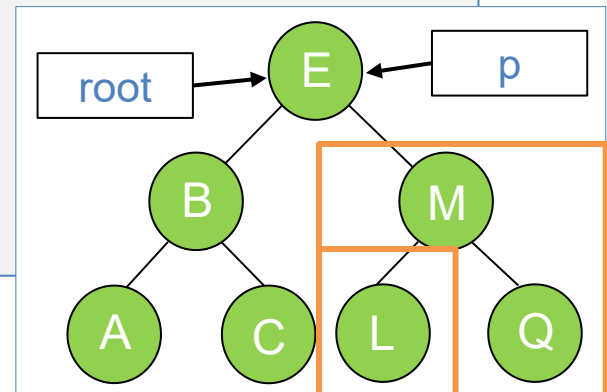
```
    public boolean search(E toSearch) {
```

```
        return search(root, toSearch);
```

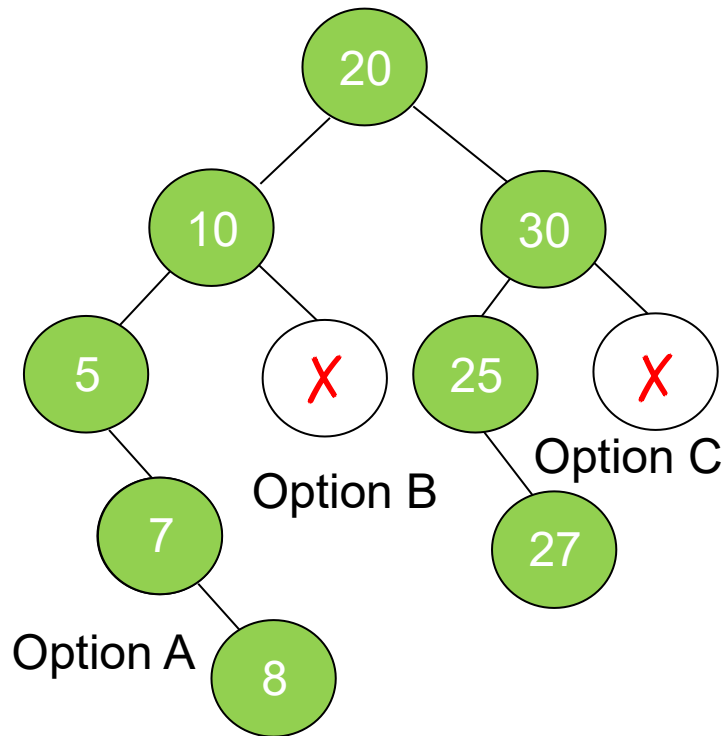
```
    }
```

```
}
```

```
t.search('L')
```



# Inserting into a BST



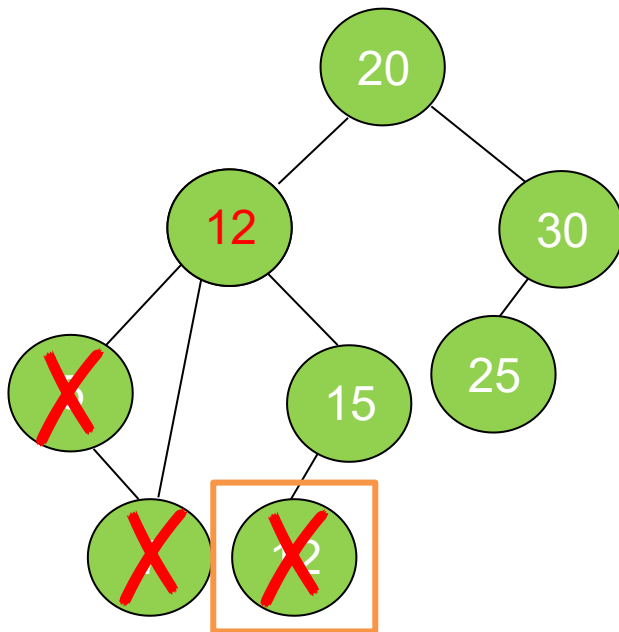
Where should we insert 7?

Insert 27?

Insert 8?

**X** Option D: Either Option A or Option B are fine.

# Deleting from a BST



Which of the following is true about the smallest element in a node's right subtree?

- A. Its left child is null
- B. Its right child is null
- C. Both of its children are null

Delete 7

If leaf node: Delete parent's link 7

Delete 5

If only one child, hoist child

Delete 10

When a deleted node has two children, this gets tricky.

Find smallest value in right subtree

Replace deleted element with  
smallest right subtree value

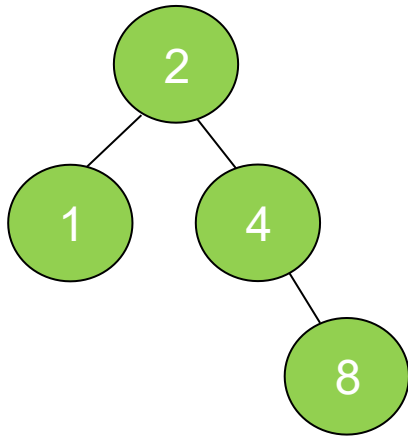
Then delete right subtree duplicate (12)

# Binary Search Tree Shape

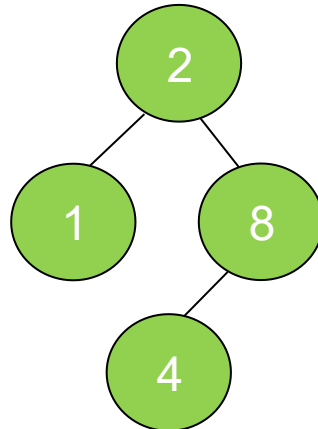
Which of the following Binary Search Trees could be the result of adding elements: 1, 2, 4, and 8 in some order.

These are all valid binary search trees!

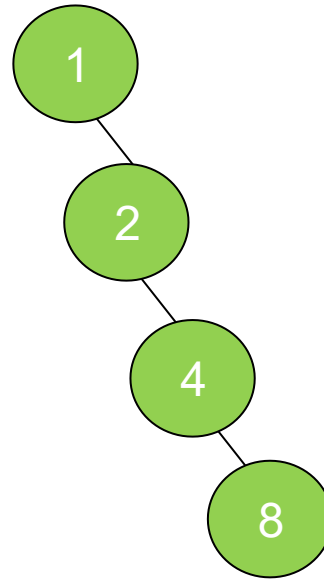
A



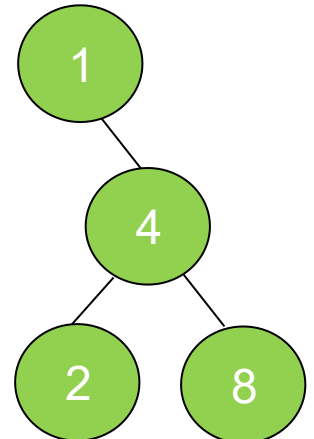
B



C



D





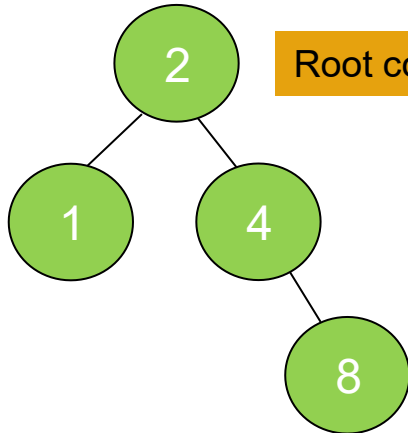
# Binary Search Tree Shape (Contd.)

Inserting a node means making it a child of an existing node

A



Insert nodes as leaves



Root comes first

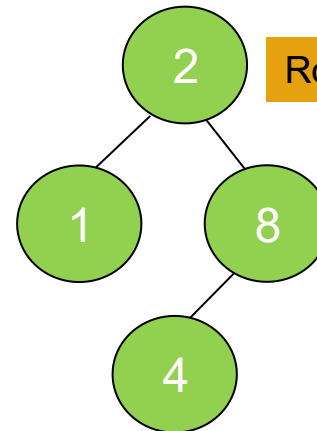
8 needs to be inserted AFTER 4



B

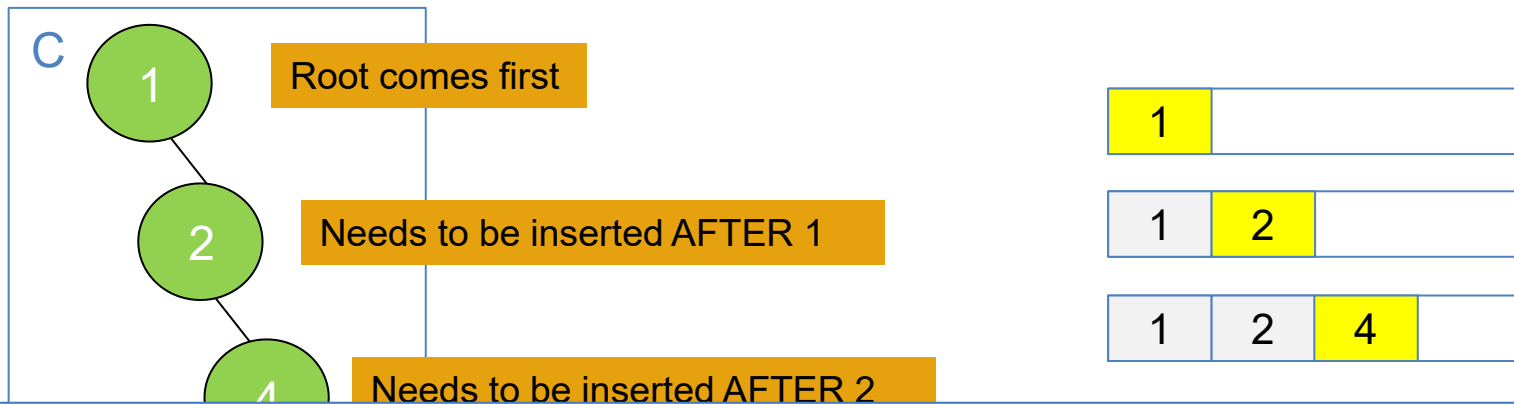


Root comes first

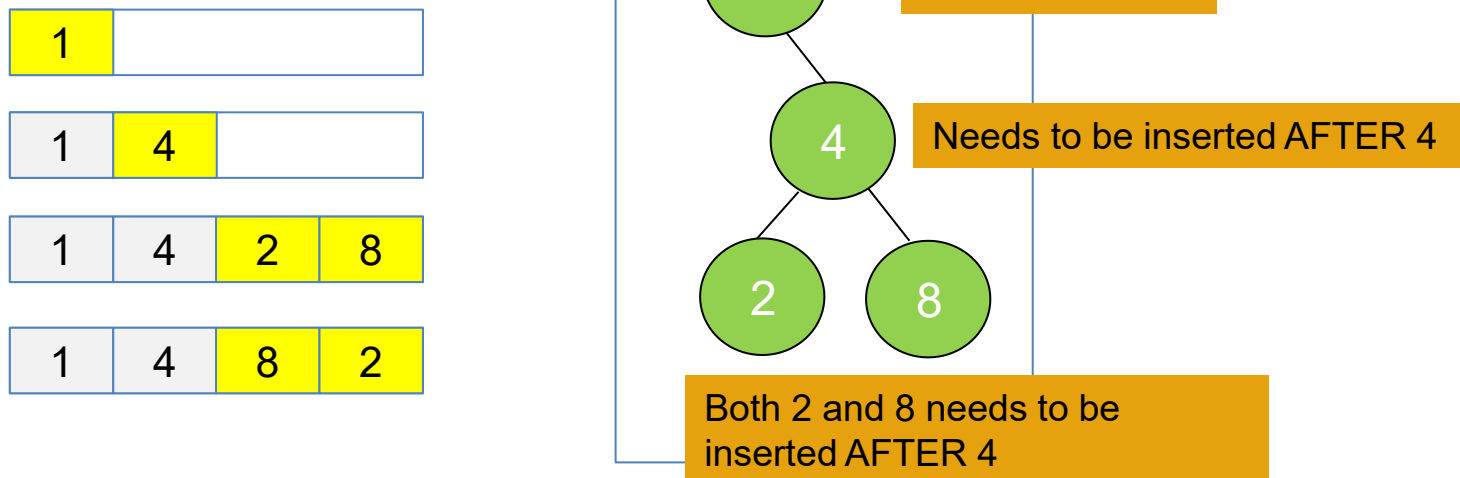


4 needs to be inserted AFTER 8

# Binary Search Tree Shape (Contd.)



The order in which we put elements into a BST impacts the shape, and what you'll see is that the shape of BST will have a huge impact on the performance of operations.



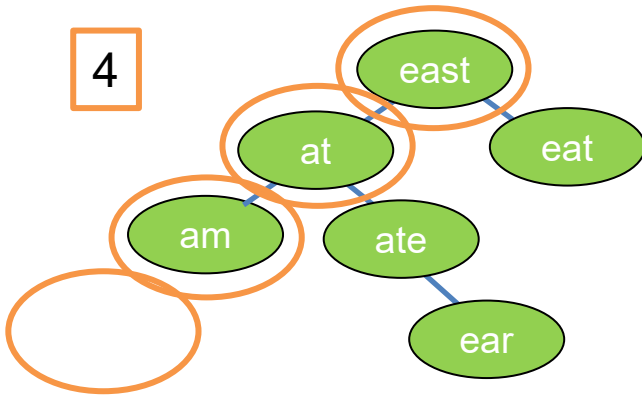
# Performance Analysis of BST

Storing a dictionary as a BST

{ am, at, ate, ear, eat, east }

Structure of a BST depends on the order of insertion

4



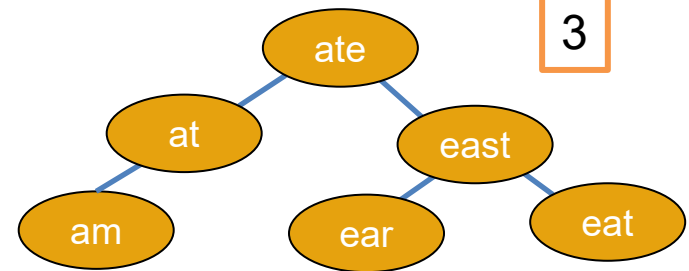
`isWord(east)`

Best case:  $O(1)$

`isWord(a)`

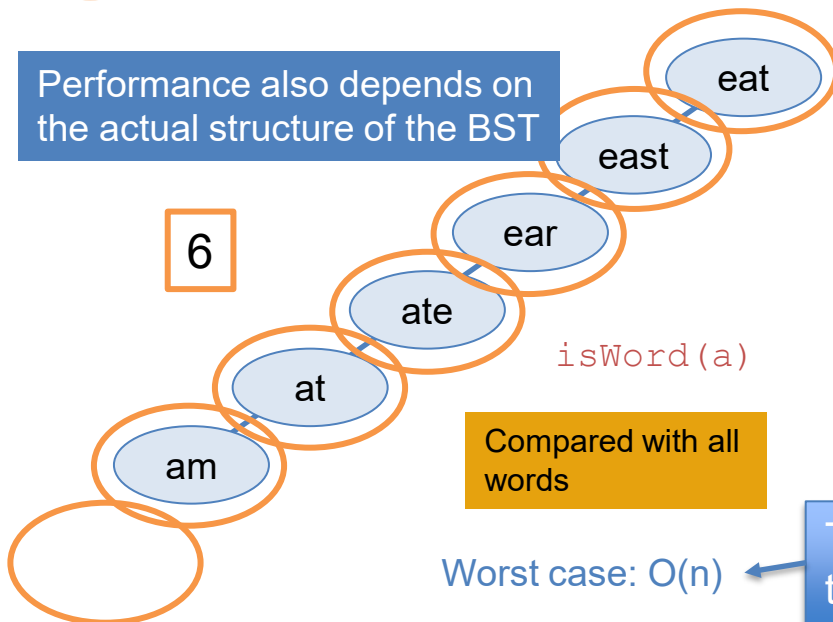
Compared with 3 out of 7 words

3



Performance also depends on the actual structure of the BST

6



`isWord(a)`

Compared with all words

Worst case:  $O(n)$

How does the performance of `isWord` relate to input size  $n$ ?

`isWord(String wordToFind)`

1. Start at root
2. Compare word to current node
  1. If current node is null, return false
  2. If `wordToFind` is less than word at current node, continue searching in left subtree
  3. If `wordToFind` is greater than word at current node, continue searching in right subtree
  4. If `wordToFind` is equal to word at current node, return true

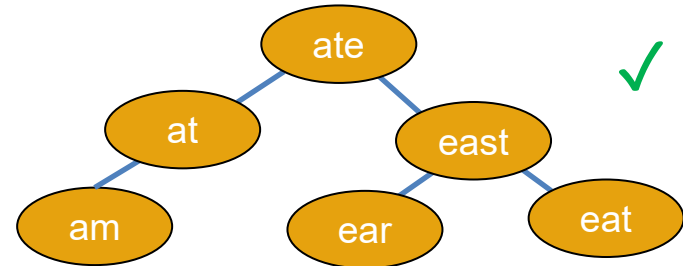
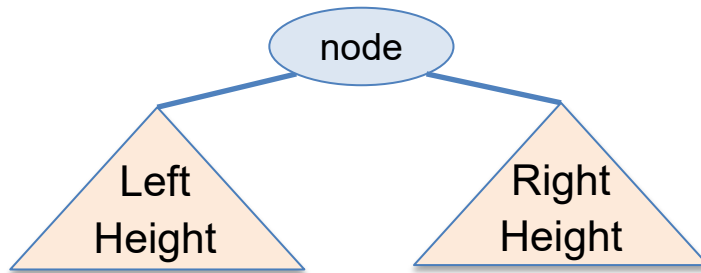
To optimize the worst case, we can modify the tree to control the max distance until leaf

height

# Balanced BST

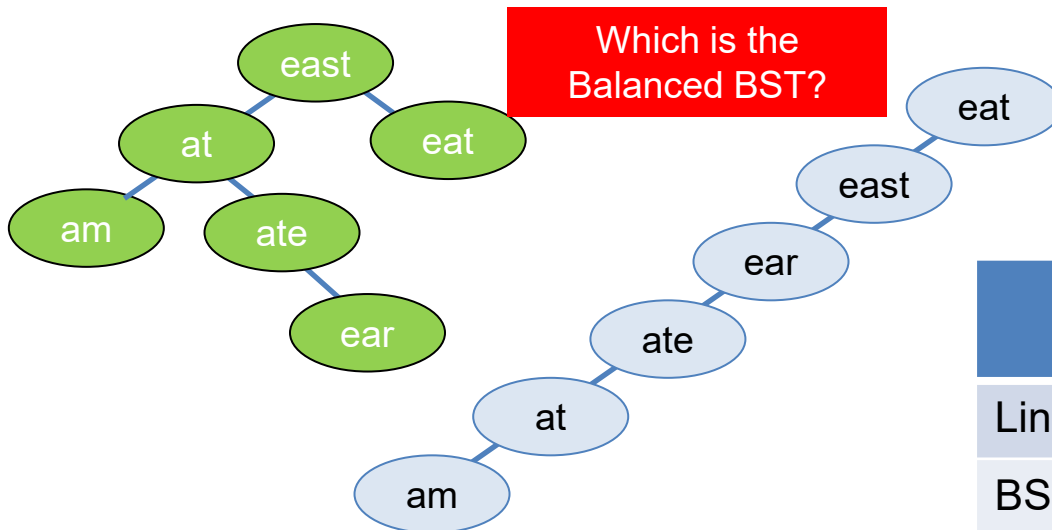
We want to keep the height down as much as we can while still maintaining the same number of nodes.

$$| \text{LeftHeight} - \text{RightHeight} | \leq 1$$



$$\text{height} \approx \log(n)$$

Especially if insert to BST in order!



	Best case	Average case	Worst case
Linked List	O(1)	O(n)	O(n)
BST	O(1)	O(log n)	O(n)
Balanced BST	O(1)	O(log n)	O(log n)

How to keep balanced? TreeSet and TreeMap in Java API

`isWord(String wordToFind)`

# BST vs. Hash Table

- Time Complexity
  - Average case:
    - Hash Tables generally offer  $O(1)$  average time complexity for insertion, deletion, and search operations.
    - BSTs provide  $O(\log n)$  time complexity for these operations, assuming the tree is balanced.
  - Worst case
    - Hash Tables can degrade to  $O(n)$  performance in cases of poor hash function design or many collisions.
    - BSTs maintain  $O(\log n)$  performance even in the worst-case for self-balancing BST.
- Ordered Operations
  - BSTs excel at operations requiring ordered data
    - In-order traversal yields sorted elements.
    - Efficient range searches and finding closest elements.
  - Hash Tables do not inherently maintain order, making these operations more difficult.

# Tree vs. Trie

- Structure and Purpose

- Trees:

- General-purpose data structure for representing hierarchical relationships
    - Each node can contain any type of data
    - Nodes typically have a value and references to child nodes

- Tries:

- Specialized tree structure for storing and retrieving strings efficiently
    - Also known as a prefix tree
    - Optimized for operations on strings or sequences

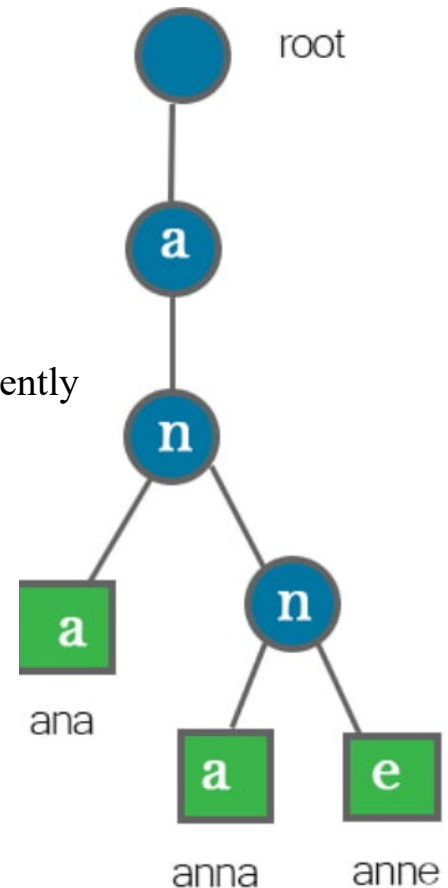
- Node Content

- Trees:

- Each node stores a value directly

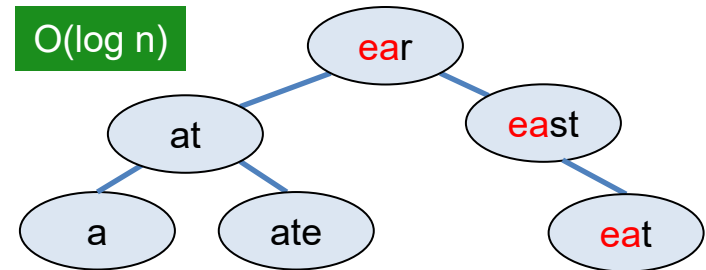
- Tries:

- Nodes typically do not store complete strings
    - The path from the root to a node represents a string or prefix
    - Characters are stored along the edges between nodes



# re(TRY)ve

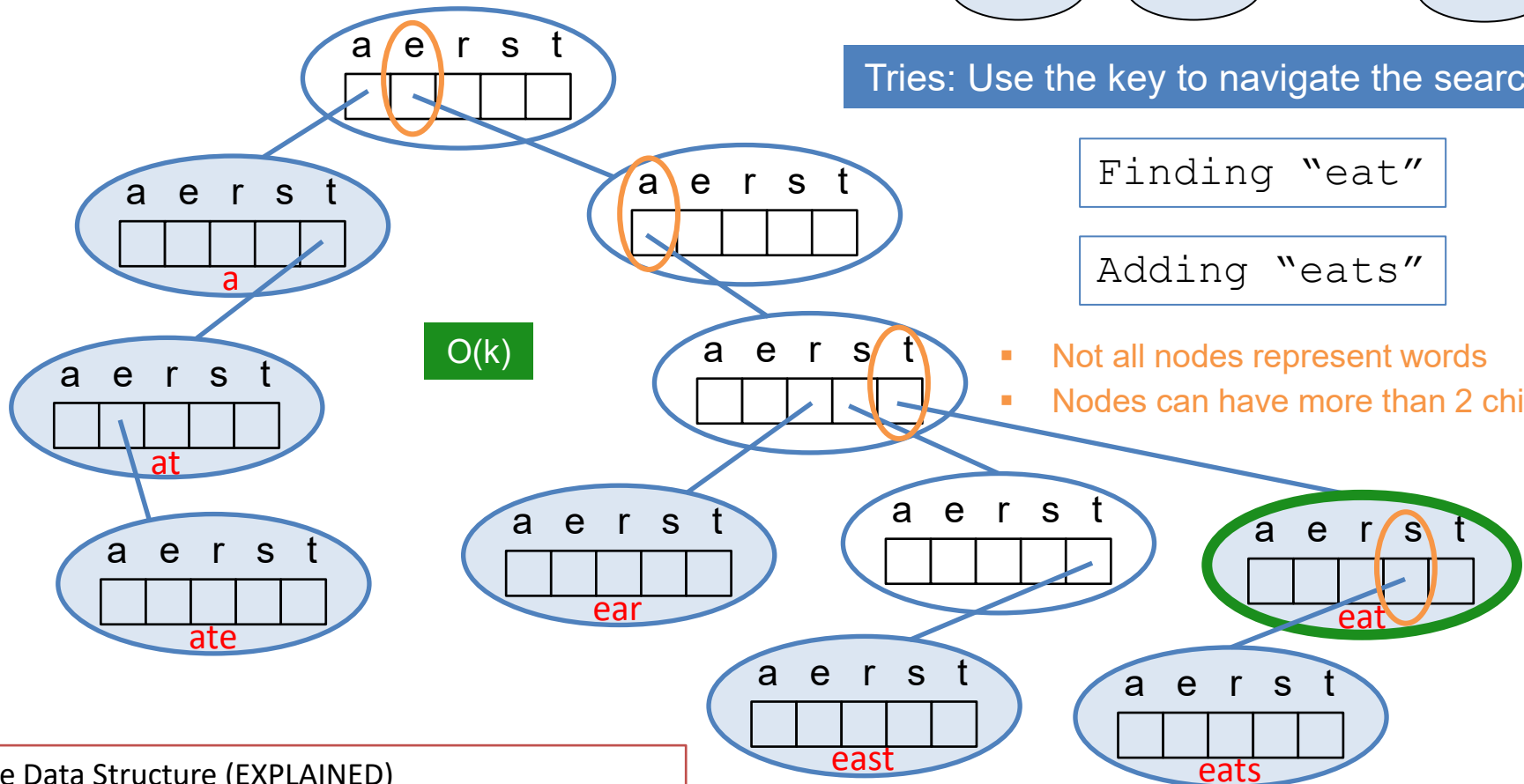
## BSTs don't take advantage of shared structure



Tries: Use the key to navigate the search

## Adding "eats"

- Not all nodes represent words
- Nodes can have more than 2 children



# Trie Data Structure (EXPLAINED)

<https://www.youtube.com/watch?v=-urNrIAQnNo>

$$\log_2(250000) \approx 18$$

# Additional Resources

- Trees and Binary Search Trees

- <http://www.openbookproject.net/thinkcs/archive/java/english/chap17.htm> -- explains trees, how to build and traverse it
- <http://algs4.cs.princeton.edu/32bst/> -- about binary search trees
- Data structures: Binary Search Tree
  - [https://www.youtube.com/watch?v=pYT9F8\\_LFTM](https://www.youtube.com/watch?v=pYT9F8_LFTM)

- Tries

- <https://www.toptal.com/java/the-trie-a-neglected-data-structure> -- explains with solid example
- <https://www.topcoder.com/community/data-science/data-science-tutorials/using-tries/> -- explains as well as providing code