# Embedded Systems with ARM Cortex-M Microcontrollers in Assembly Language and C

Chapter 5
Memory Access
Exercises ANS

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Fall 2025

Acknowledgement: Lecture slides based on Embedded Systems with ARM Cortex-M Microcontrollers in Assembly Language and C, University of Maine https://web.eece.maine.edu/~zhu/book/

#### Question: Endianness

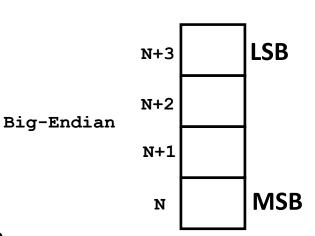
**32-bit** Bytes Addr. Words 0015 Addr 0014 Word 3 ?? 0013 0012 What are the memory address of 0011 these four words? Addr 0010 Word 2 33 0009 8000 0007 Addr 0006 Word 1 0005 ?? 0004 0003 Addr 0002 Word 0 0001 ?? 0000 2

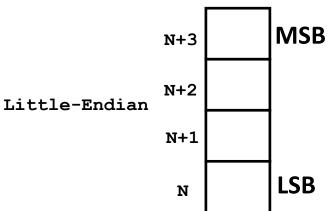
#### Answer: Endianness

		32-bit Words	Bytes	Addr.
				0015
	Word 3	Addr =		0014
	WOIG 3	0x0012		0013
				0012
What are the memory address of		Aslala		0011
these four words?	Word 2	0x0008 Addr		0010
Same as the address of the lowest-				0009
address Byte				0008
(this is true for either				0007
Little-Endian or Big-	Word 1			0006
Endian ordering)				0005
				0004
		Aslala		0003
	Word 0	0 <b>Addr</b> =		0002
		0x0000		0001
3				0000

#### Question: Endianness

- Q:Assume Big-Endian ordering. If a 32-bit word resides at memory address N, what is the address of:
  - (a) The MSB (Most Significant Byte)
  - (b) The 16-bit half-word corresponding to the most significant half of the word
- Q: Redo the question assuming Little-Endian ordering.

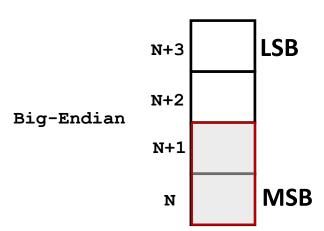




#### **Answer: Endianness**

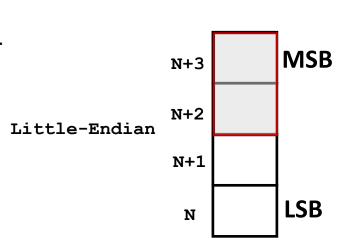
#### A:With Big-Endian ordering:

- (a) Address of MSB: N
- (b) Address of 16-bit half-word corresponding to the most significant half of the word: N (the halfword has address range of [N, N+1], so its address is N)



#### With Little-Endian ordering:

- (a) Address of MSB: N+3
- (b) Address of 16-bit half-word corresponding to the most significant half of the word: N+2 (the halfword has address range of [N+2, N+3], so its address is N+2)



#### Question: Endianness

The word stored at address 0x20008000 with Big-Endian ordering is



The word stored at address 0x20008000 with Little-Endian ordering is



Memory Address	Memory Data
0x20008003	0xA7
0x20008002	0x90
0x20008001	0x8C
0x20008000	0×EE

#### **Answer: Endianness**

The word stored at address 0x20008000 with Big-Endian ordering is

**0xEE8C90A7** 

The word stored at address 0x20008000 with Little-Endian ordering is

**0xA7908CEE** 

Memory Address	Memory Data
0x20008003	0xA7
0x20008002	0x90
0x20008001	0x8C
0x20008000	0×EE

Endianness only specifies byte order, not bit order in a byte!

#### Endianness

LDR r11, [r0]; r0 = 0x20008000

r11 before load

0x12345678

r11 after load w/
Big-Endian ordering

r11 after load w/

Little-Endian ordering

Memory Address	Memory Data
0x20008003	0xA7
0x20008002	0x90
0x20008001	0x8C
0x20008000	0xEE

#### **Endianness ANS**

LDR r11, [r0]; r0 = 0x20008000

rll before load

0x12345678

r11 after load w/
Big-Endian ordering

**0xEE8C90A7** 

r11 after load w/
Little-Endian ordering

**0**xA7908CEE

Memory Address	Memory Data
0x20008003	0xA7
0x20008002	0x90
0x20008001	0x8C
0x20008000	0×EE

# Review

#### Data Alignment

- Assume a byte-addressable memory with a data bus that is 32 bits (4 bytes) wide
- Consider 16 bytes of memory (addresses 0 to 15) arranged as four 32-bit words (4

	- `		
Address 15	Address 14	Address 13	Address 12
Address 11	Address 10	Address 9	Address 8
Address 7 (MSbyte)	Address 6	Address 5	Address 4 (LSbyte)
Address 3	Address 2	Address 1	Address 0

Well-aligned: each word begins on a mod-4 address, which can be read in a single memory cycle

The first read cycle would retrieve 4 bytes from addresses 4 through 7; of these, the bytes from addresses 4 and 5 are discarded, and those from addresses 6 and 7 are moved to the far right; The second read cycle retrieves 4 bytes from addresses 8 through 11; the bytes from addresses 10 and 11 are discarded, and those from addresses 8 and 9 are moved to the far left;

Finally, the two halves are combined to form the desired 32-bit operand:

Address 15	Address 14	Address 13	Address 12
Address 11	Address 10	Address 9 (MSbyte)	Address 8
Address 7	Address 6 (LSbyte)	Address 5	Address 4
Address 3	Address 2	Address 1	Address 0

Ill-aligned: a word begins on
address 6, not a mod-4 address,
which can be read in 2 memory
cycles

	Address 7	Address 6 (LSbyte)

Address 9 (MSbyte) Address 8

Address 9 (MSbyte)	Address 8	Address 7	Address 6 (LSbyte)
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### Question: Data Alignment

- Q: Assume a byte-addressable memory with a data bus that is 32 bits (4 bytes) wide. Consider 16 bytes of memory (addresses 0 to 15) arranged as four 32-bit words (4 bytes each). How many memory cycles are required to read each of the following from memory?
  - (a) A 2-Byte operand read from decimal address 5
  - (b) A 2-Byte operand read from decimal address 15
  - (c) A 4-Byte operand read from decimal address 10
  - (d) A 4-Byte operand read from decimal address 20

# Answer: Data Align

Address 15	Address 14	Address 13	Address 12
Address 11	Address 10	Address 9	Address 8
Address 7	Address 6	Address 5	Address 4
Address 3	Address 2	Address 1	Address 0

- Q: Assume a byte-addressable memory with a data bus that is 32 bits (4 bytes) wide. How many memory cycles are required to read each of the following from memory?
  - (a) A 2-Byte operand read from decimal address 5
  - (b) A 2-Byte operand read from decimal address 15
  - (c) A 4-Byte operand read from decimal address 10
  - (d) A 4-Byte operand read from decimal address 20
- A: (a) The operand contains memory content in address range [5,6]. It can be read in I memory cycle; the memory controller returns a word in address range [4,7]. The operand can be obtained via I-Byte offset addressing into the word.
- (b) The operand contains memory content in address range [15,16]. It can be read in 2 memory cycles; the memory controller returns 2 words in address ranges [12,15] and [16, 19], which can be combined to return a word in address range [14,17]. The operand can be obtained via 1-Byte offset addressing into the word.
- (c) The operand contains memory content in address range [10,13]. It can be read in 2 memory cycles; the memory controller returns 2 words in address ranges [8,11] and [12, 15], which can be combined to return a word with address range [10,13].
- (d) The operand contains memory content in address range [20,23]. Since 20%4=0, it is well-aligned, and can be read in 1 memory cycle.

# Question: Data Align

Address 111	Address 110	Address 109	Address 108
Address 107	Address 106	Address 105	Address 104
Address 103	Address 102	Address 101	Address 100
Address 99	Address 98	Address 97	Address 96

- Q: Assume a byte-addressable memory with a data bus that is 32 bits (4 bytes) wide. Consider 16 bytes of memory (addresses 0 to 15) arranged as four 32-bit words (4 bytes each).
  - (a) What is the address of MSB of the word at address 102, assuming Little-Endian ordering?
  - (b) What is the address of LSB of the word at address 102, assuming Little-Endian ordering?
  - (b) How many memory cycles are required to read the word at address 102?
  - (c) How many memory cycles are required to read the half word at address 102?

# Answer: Data Align

Address 111	Address 110	Address 109	Address 108
Address 107	Address 106	Address 105	Address 104
Address 103	Address 102	Address 101	Address 100
Address 99	Address 98	Address 97	Address 96

- Q: Assume a byte-addressable memory with a data bus that is 32 bits (4 bytes) wide. Consider 16 bytes of memory (addresses 0 to 15) arranged as four 32-bit words (4 bytes each).
  - (a) What is the address of MSB of the word at address 102, assuming Little-Endian ordering?
  - (b) What is the address of LSB of the word at address 102, assuming Little-Endian ordering?
  - (b) How many memory cycles are required to read the word at address 102?
  - (c) How many memory cycles are required to read the half word at address 102?

#### • A:

- (a) MSB of the word at address 102 is 105
- (b) LSB of the word at address 102 is 102
- (c) 2 cycles.
- (d) I cycle.

# Answer: Memory Cycles

Address 15	Address 14	Address 13	Address 12
Address 11	Address 10	Address 9	Address 8
Address 7	Address 6	Address 5	Address 4
Address 3	Address 2	Address 1	Address 0

- Q: Assume a byte-addressable memory with a data bus that is 32 bits (4 bytes) wide.
  - It takes \_\_\_\_\_ memory cycle(s) to read a Byte from memory
  - It takes \_\_\_\_\_ memory cycle(s) to read a half-word from memory
  - It takes \_\_\_\_\_ memory cycle(s) to read a word from memory
  - It takes \_\_\_\_\_ memory cycle(s) to read a double word from memory

# Answer: Memory Cycles

Address 15	Address 14	Address 13	Address 12
Address 11	Address 10	Address 9	Address 8
Address 7	Address 6	Address 5	Address 4
Address 3	Address 2	Address 1	Address 0

- Q:Assume a byte-addressable memory with a data bus that is 32 bits (4 bytes) wide.
  - It takes \_\_\_\_\_ memory cycle(s) to read a Byte from memory
  - It takes \_\_\_\_\_ memory cycle(s) to read a half-word from memory
  - It takes \_\_\_\_\_ memory cycle(s) to read a word from memory
  - It takes \_\_\_\_\_ memory cycle(s) to read a double word from memory
- A:
  - It takes \_\_\_I \_\_ memory cycle(s) to read a Byte from memory
  - It takes \_\_\_I or 2\_\_ memory cycle(s) to read a half-word from memory
  - It takes \_\_\_I or 2\_\_ memory cycle(s) to read a word from memory
  - It takes \_\_\_2 or 3\_\_\_ memory cycle(s) to read a double word from memory (a double word may span at most 3 consecutive words in memory)

# Question: Arrays

- Q: If the first element of a one-dimensional array x[] is stored at memory address 0x12345678, what is address of the second element if the array x[] contains
  - (a) chars
  - (b) shorts
  - (c) ints
  - (c) longs

# Answer: Arrays

- Q: If the first element of a one-dimensional array x[] is stored at memory address 0x12345678, what is address of the second element if the array x[] contains
  - (a) chars
  - (b) shorts
  - (c) ints
  - (c) longs
- A: x[1]'s address is x's address plus the data type size in Bytes
  - (a) chars:  $0 \times 12345678 + 1 = 0 \times 12345679$
  - (b) shorts:  $0 \times 12345678 + 2 = 0 \times 1234567A$
  - (c) ints:  $0 \times 12345678 + 4 = 0 \times 1234567C$
  - (c) longs:  $0 \times 12345678 + 8 = 0 \times 12345680$