

Chapter 1

Introduction

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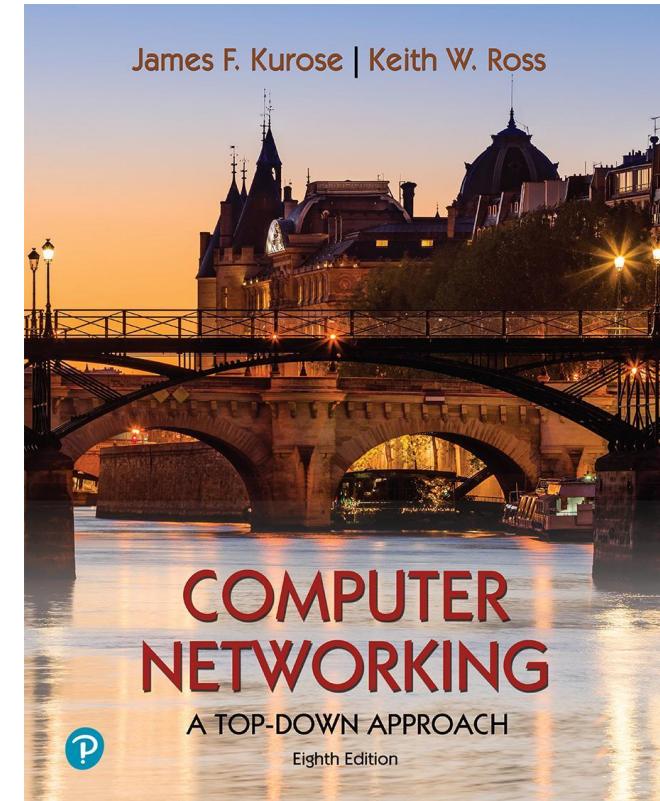
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*Computer Networking:
A Top-Down Approach*
8th edition
Jim Kurose, Keith Ross
Pearson, 2020

Chapter 1: introduction

Chapter goal:

- Get “feel,” “big picture,” introduction to terminology
 - more depth, detail *later* in course



Overview/roadmap:

- What *is* the Internet? What *is* a protocol?
- **Network edge:** hosts, access network, physical media
- **Network core:** packet/circuit switching, internet structure
- **Performance:** loss, delay, throughput
- Protocol layers, service models
- Security
- History

The Internet: a “nuts and bolts” view



Billions of connected computing **devices**:

- **hosts** = end systems
- running **network apps** at Internet's “edge”

Packet switches: forward packets (chunks of data)

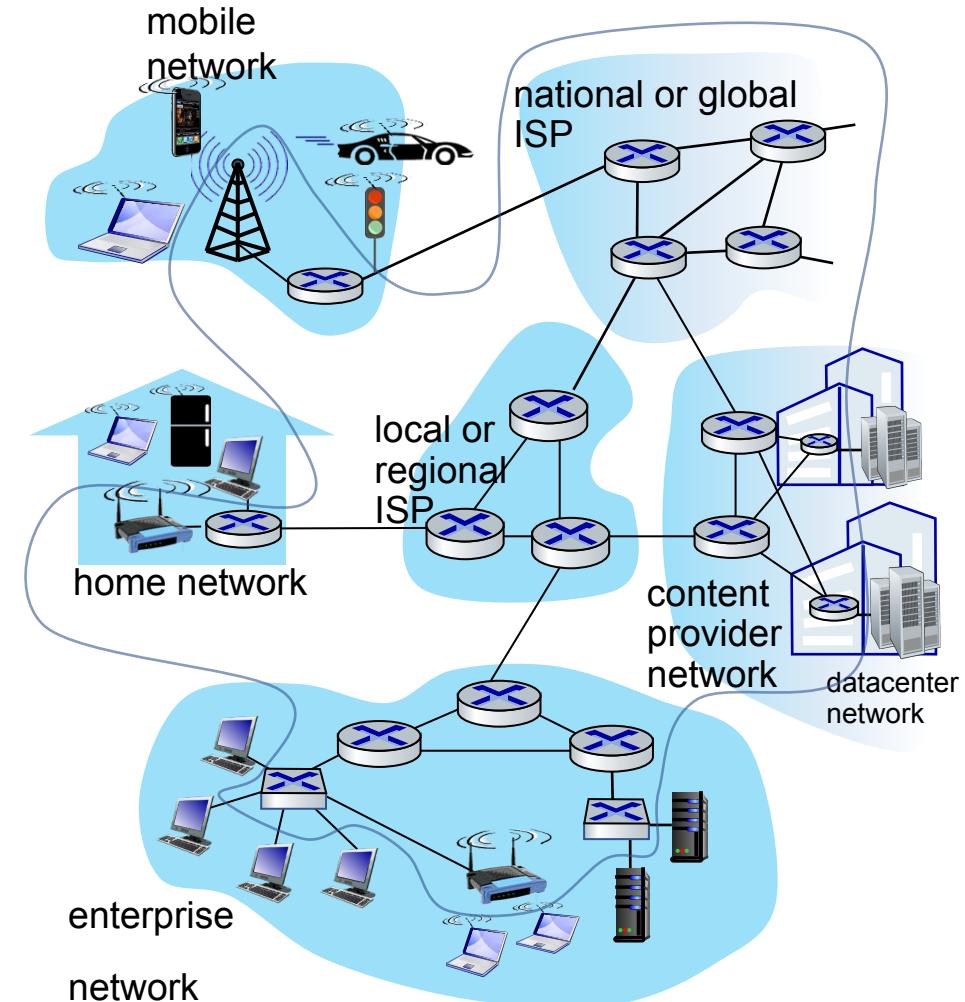
- routers, switches

Communication links

- fiber, copper, radio, satellite
- transmission rate: *bandwidth*

Networks

- collection of devices, routers, links: managed by an organization



“Fun” Internet-connected devices



Amazon
Echo



Internet
refrigerator



Security
Camera



Internet
phones



IP picture
frame



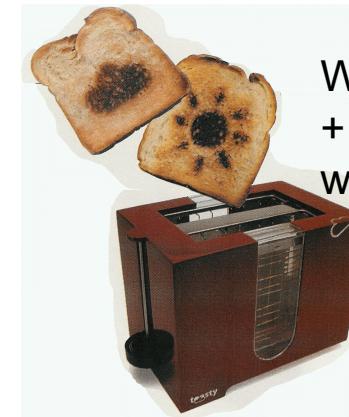
Slingbox:
remote
control cable
TV



Gaming
devices



Pacemaker &
Monitor



Web-enabled toaster
+
weather forecaster



sensorize
d,
bed
mattress



Fitbit



Tweet-a-watt:
monitor energy
use



bike
s



car
s



scooter
s

Others
?

The Internet: a “nuts and bolts” view

Rede de redes. Conjunto de provedores de acesso a internet e a conexão entre eles.

- *Internet: “network of networks”*

- Interconnected ISPs

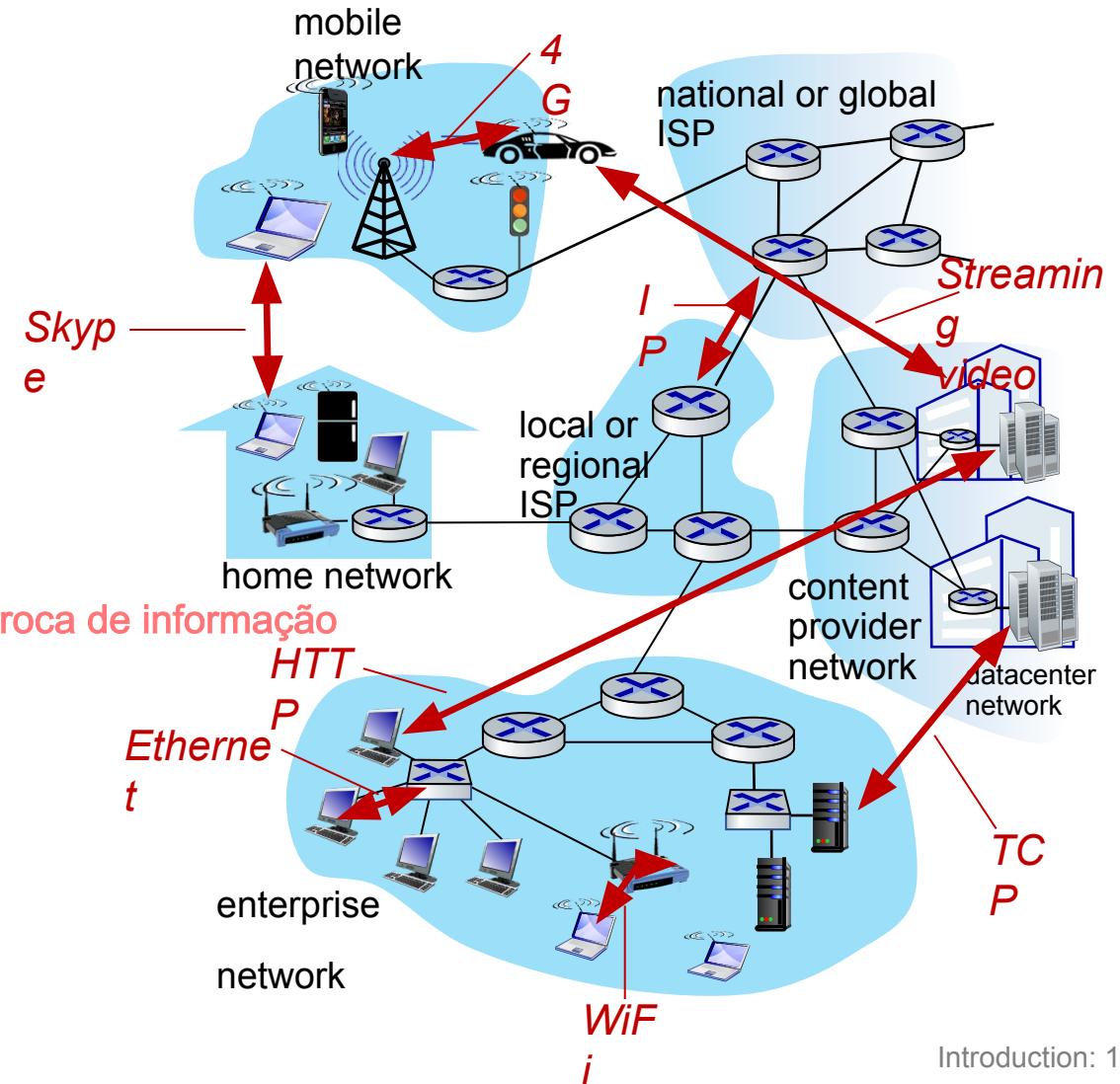
- *protocols are everywhere*

- control sending, receiving of messages
 - e.g., HTTP (Web), streaming video, Skype, TCP, IP, WiFi, 4G, Ethernet

Conjunto de regras para troca de informação
Privado - Skype
Publico - HTTP, 4G

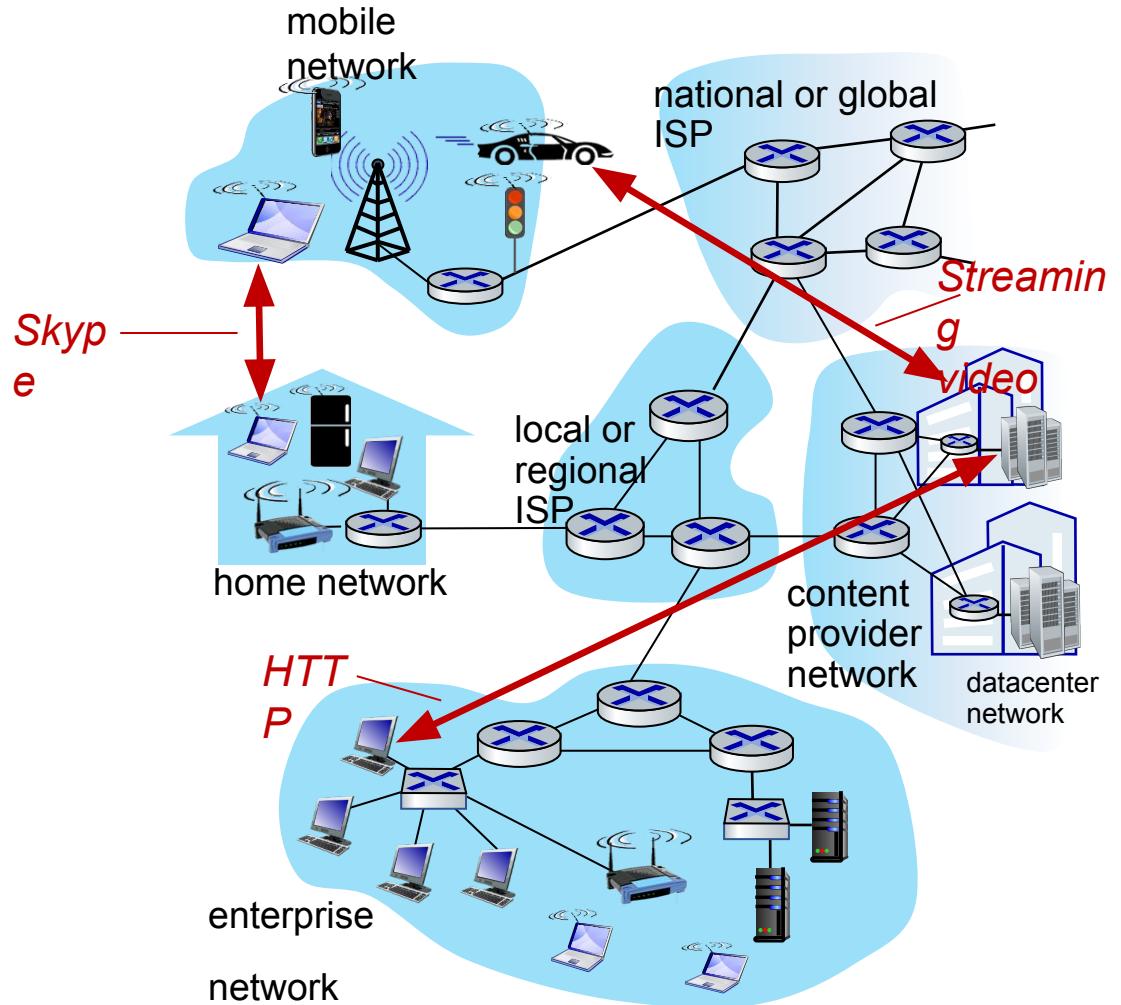
- *Internet standards*

- RFC: Request for Comments
 - IETF: Internet Engineering Task Force



The Internet: a “services” view

- **Infrastructure** that provides services to applications:
 - Web, streaming video, multimedia teleconferencing, email, games, e-commerce, social media, inter-connected appliances
- provides **programming interface** to distributed applications:
 - “hooks” allowing sending/receiving apps to “connect” to, use Internet transport service
 - provides service options, analogous to postal service



What's a protocol?

Human protocols:

- “what’s the time?”
- “I have a question”
- introductions

Rules for:

- ... specific messages sent
- ... specific actions taken when message received, or other events

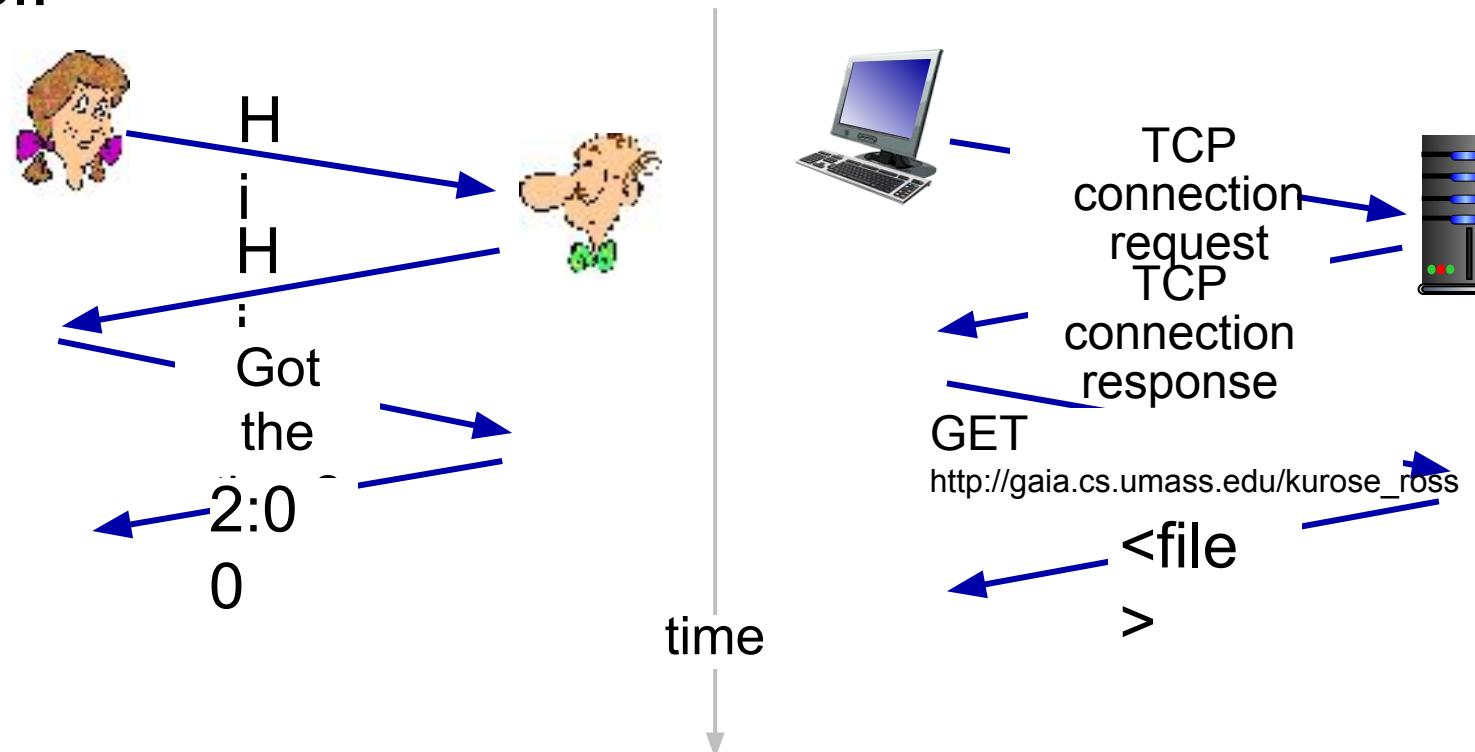
Network protocols:

- computers (devices) rather than humans
- all communication activity in Internet governed by protocols

Protocols define the format, order of messages sent and received among network entities, and actions taken on message transmission, receipt

What's a protocol?

A human protocol and a computer network protocol:



Q: other human protocols?

Chapter 1: roadmap

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- What *is* a protocol?
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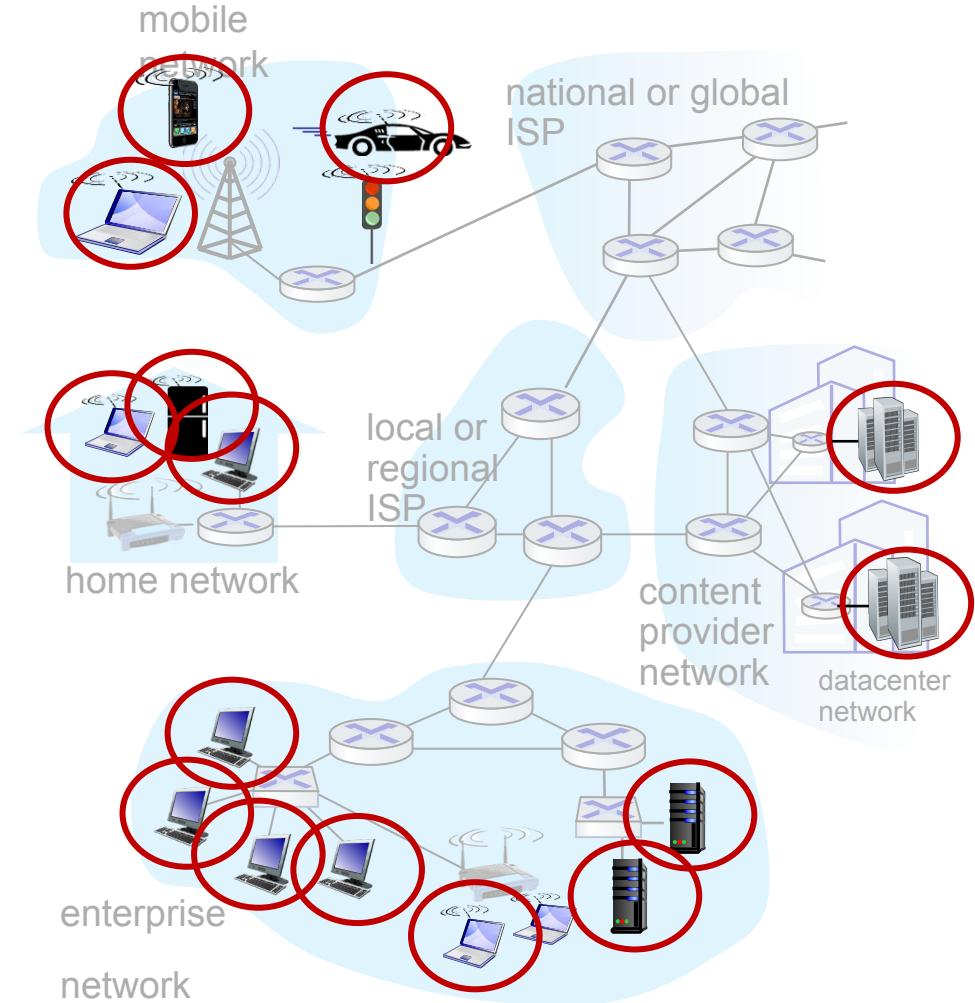


A closer look at Internet structure

Network edge:

- hosts: clients and servers
- servers often in data centers

- Sistemas finais na borda (Notebooks, desktops, maquinas servidoras)



A closer look at Internet structure

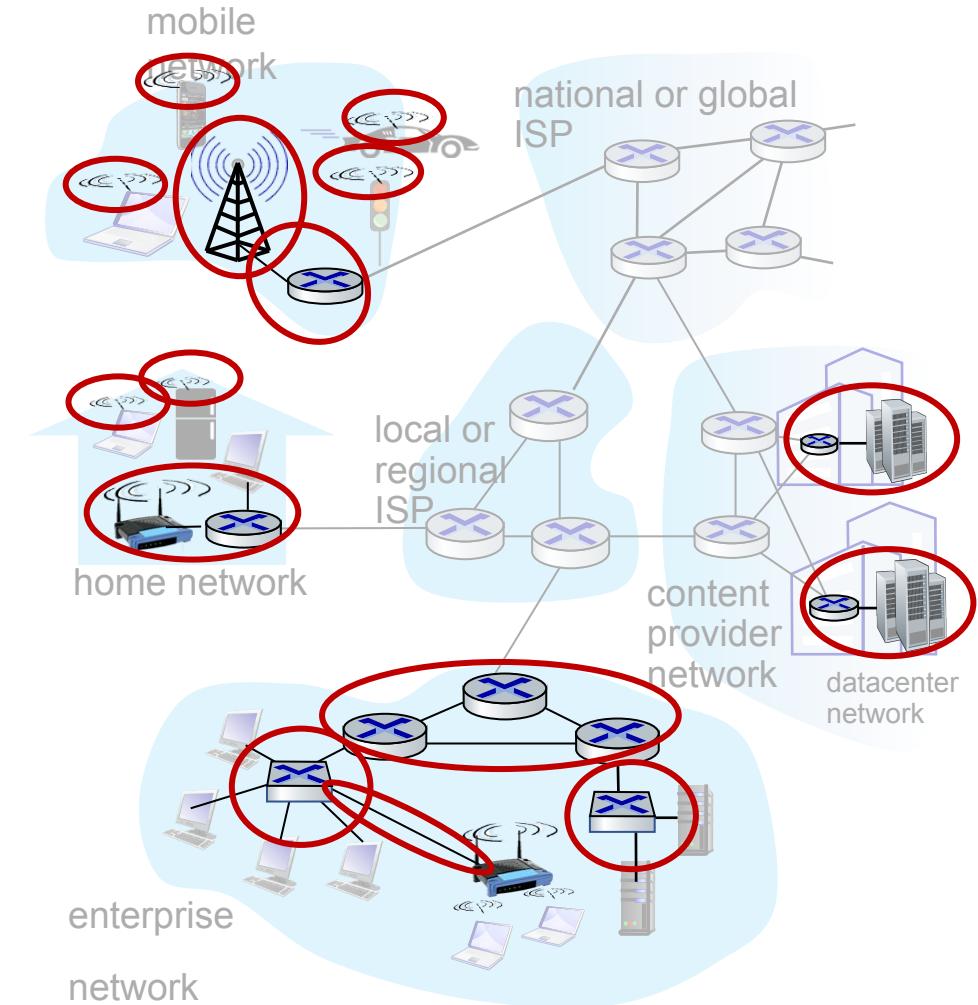
Network edge:

- hosts: clients and servers
- servers often in data centers

Access networks, physical media:

- wired, wireless communication links

- roteadores dao acesso ao núcleo da internet



A closer look at Internet structure

Network edge:

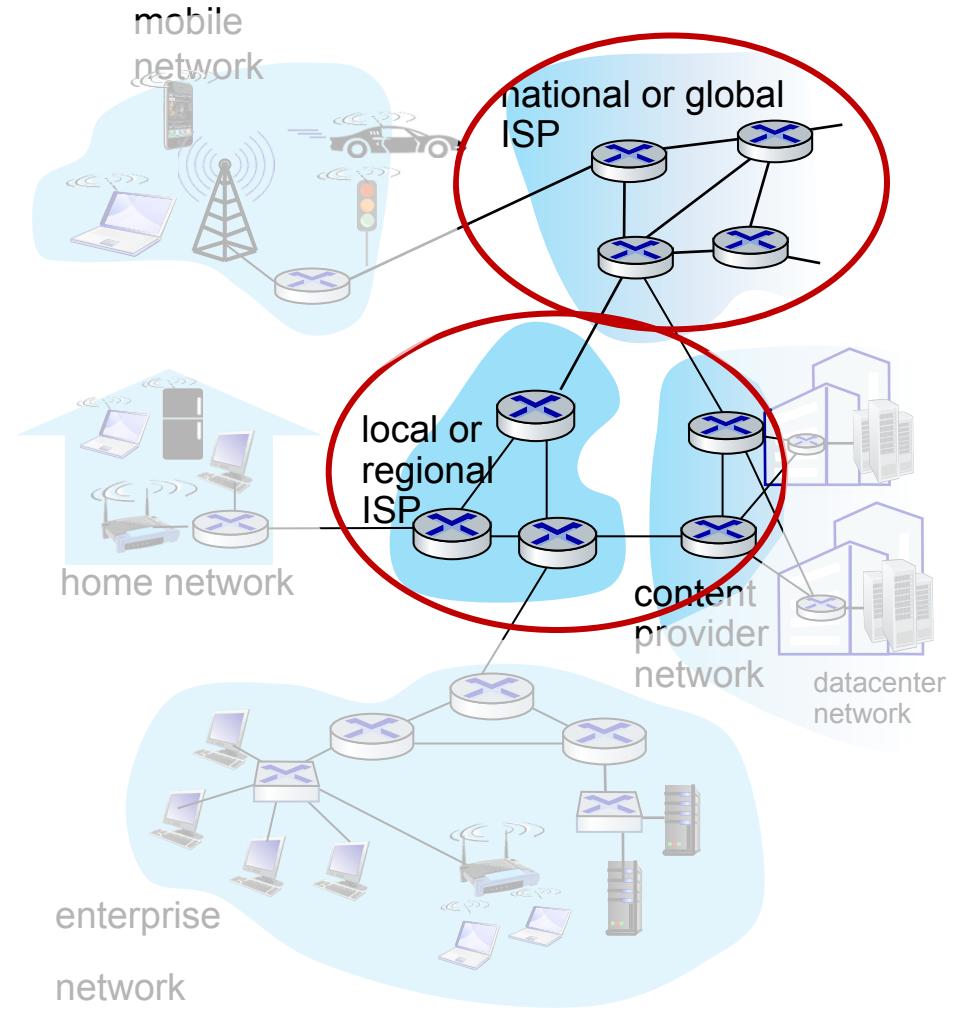
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Access networks, physical media:

- wired, wireless communication links

Network core:

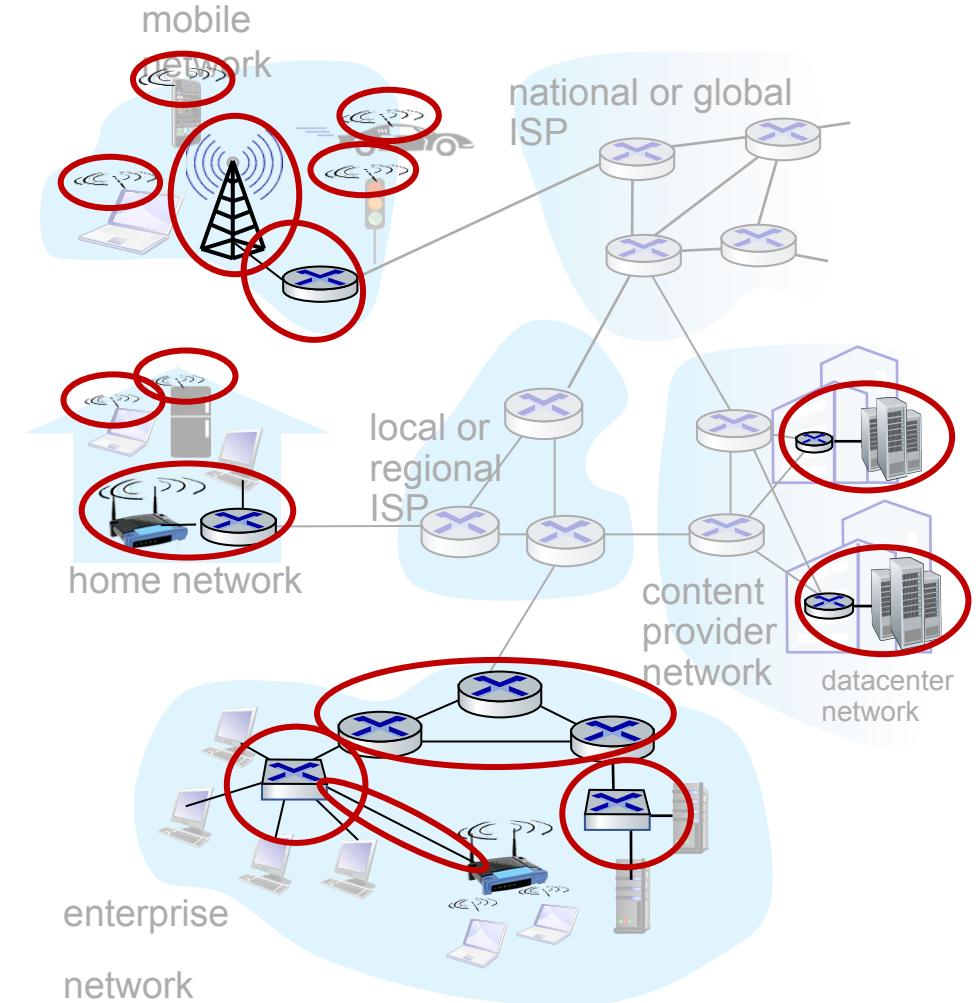
- interconnected routers
- network of networks



Access networks and physical media

Q: How to connect end systems to edge router?

- residential access nets
- institutional access networks (school, company)
- mobile access networks (WiFi, 4G/5G)



Quais os elementos listados abaixo que compõem a borda da rede:

A sua resposta :

Dispositivos hospedeiros

Retorno:

Correto.

Redes de acesso

Retorno:

Correto.

Roteadores de borda

Retorno:

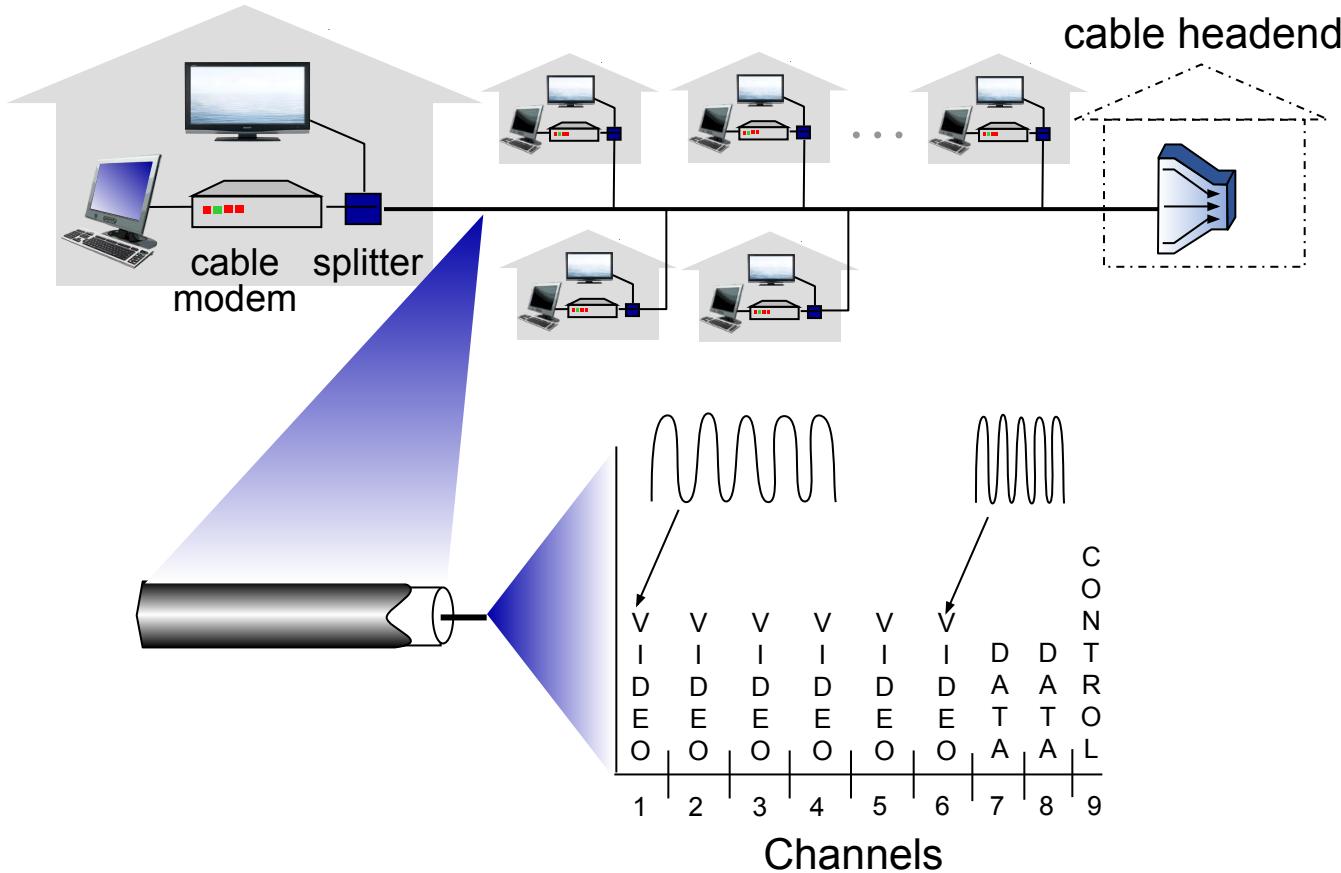
Correto.

Enlaces de comunicação

Retorno:

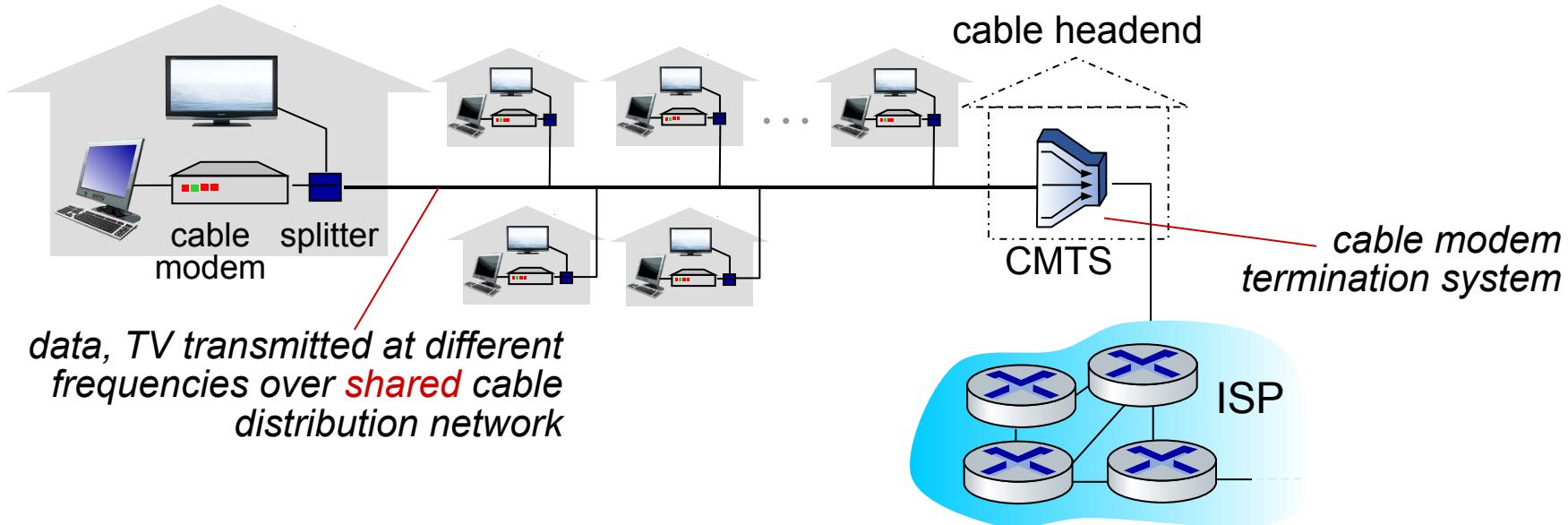
Correto.

Access networks: cable-based access



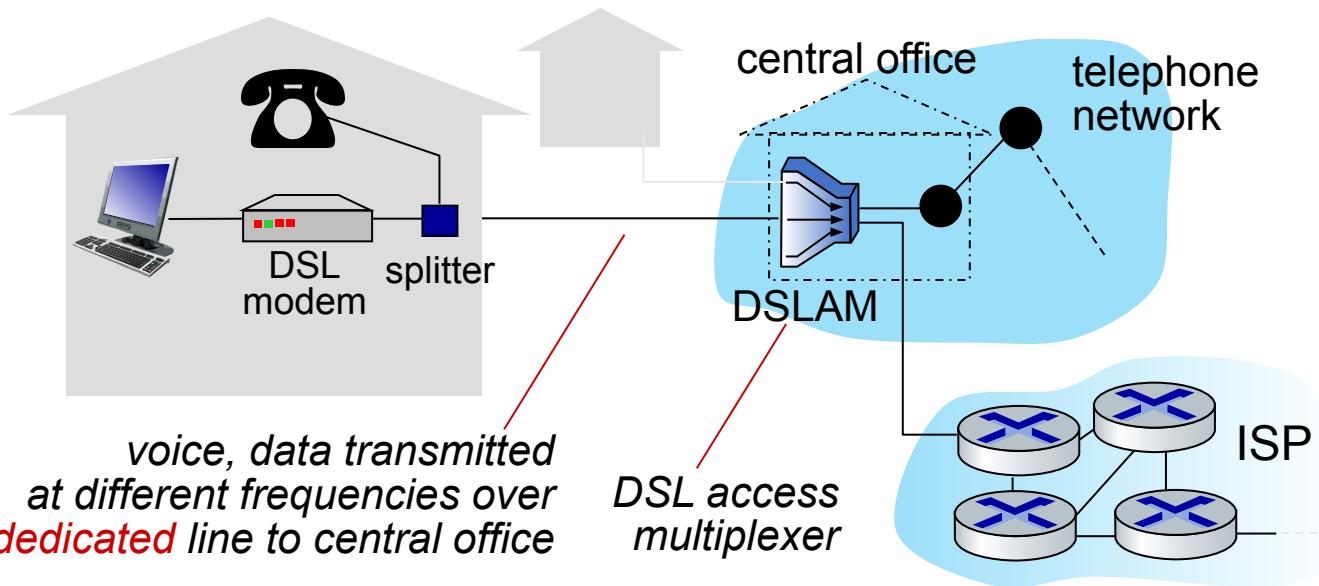
frequency division multiplexing (FDM): different channels transmitted in different frequency bands

Access networks: cable-based access



- HFC: hybrid fiber coax
 - asymmetric: up to 40 Mbps – 1.2 Gbps downstream transmission rate, 30-100 Mbps upstream transmission rate
- network of cable, fiber attaches homes to ISP router
 - homes **share access network** to cable headend

Access networks: digital subscriber line (DSL)



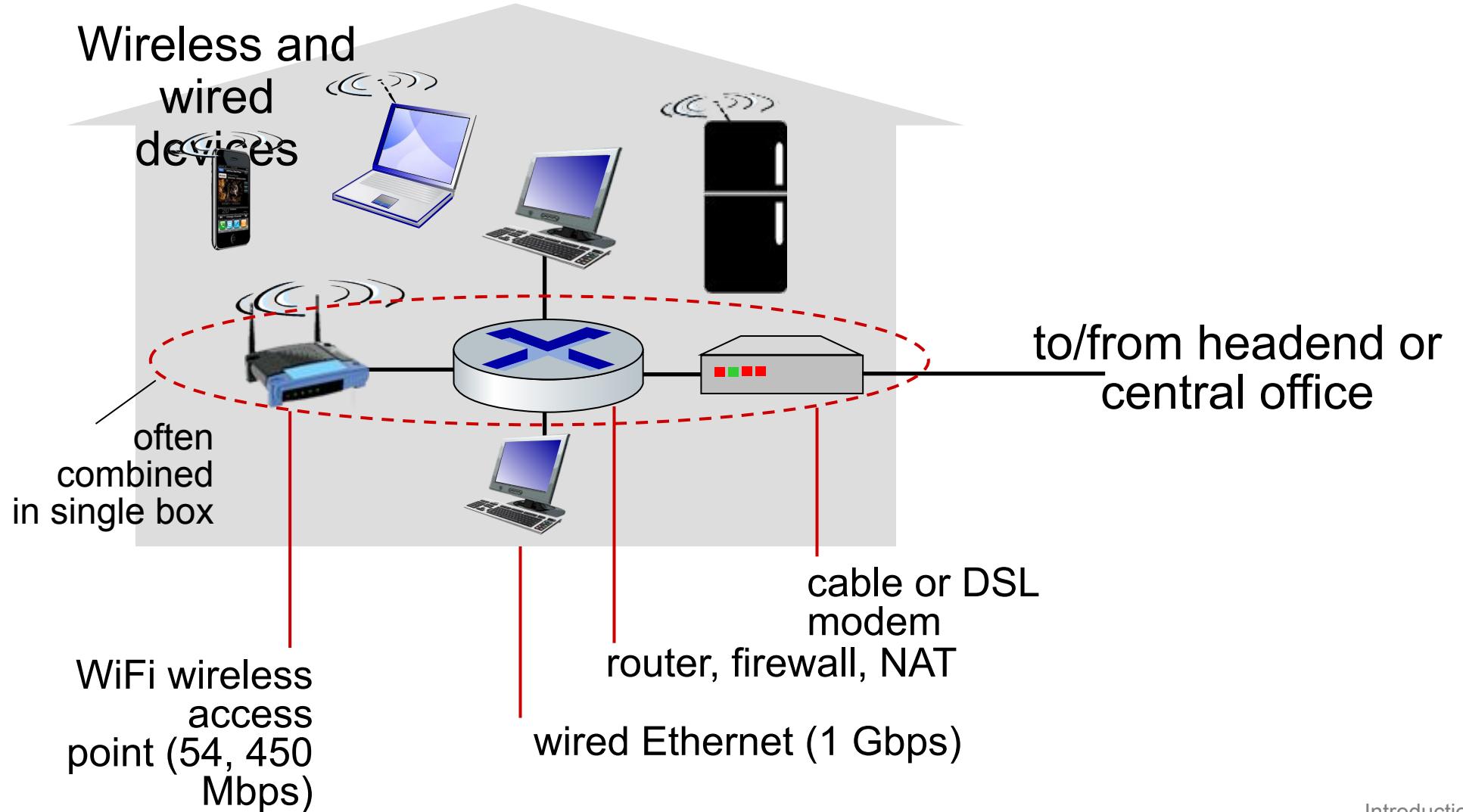
DSL -> Cabos telefônicos são aproveitados para transmissão do sinal de internet. Usam a diferença de frequência.

DSL ->

downstream - download
upstream - upload

- use *existing* telephone line to central office DSLAM
 - data over DSL phone line goes to Internet
 - voice over DSL phone line goes to telephone net
- 24-52 Mbps dedicated downstream transmission rate
- 3.5-16 Mbps dedicated upstream transmission rate

Access networks: home networks



Wireless access networks

Shared *wireless* access network connects end system to router

- via base station aka “access point”

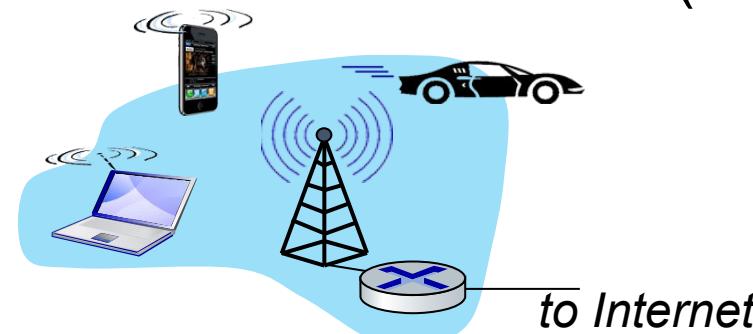
Wireless local area networks (WLANs)

- typically within or around building (~100 ft)
- 802.11b/g/n (WiFi): 11, 54, 450 Mbps transmission rate



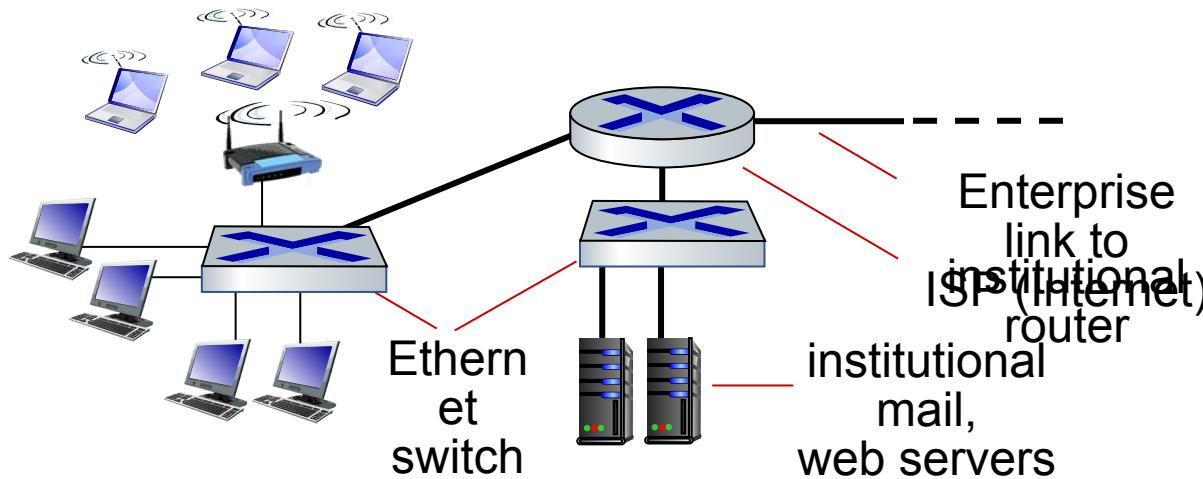
Wide-area cellular access networks

- provided by mobile, cellular network operator (10's km)
- 10's Mbps
- 4G cellular networks (5G coming)



Access networks: enterprise networks

Ethernet e Wifi s o protocolos



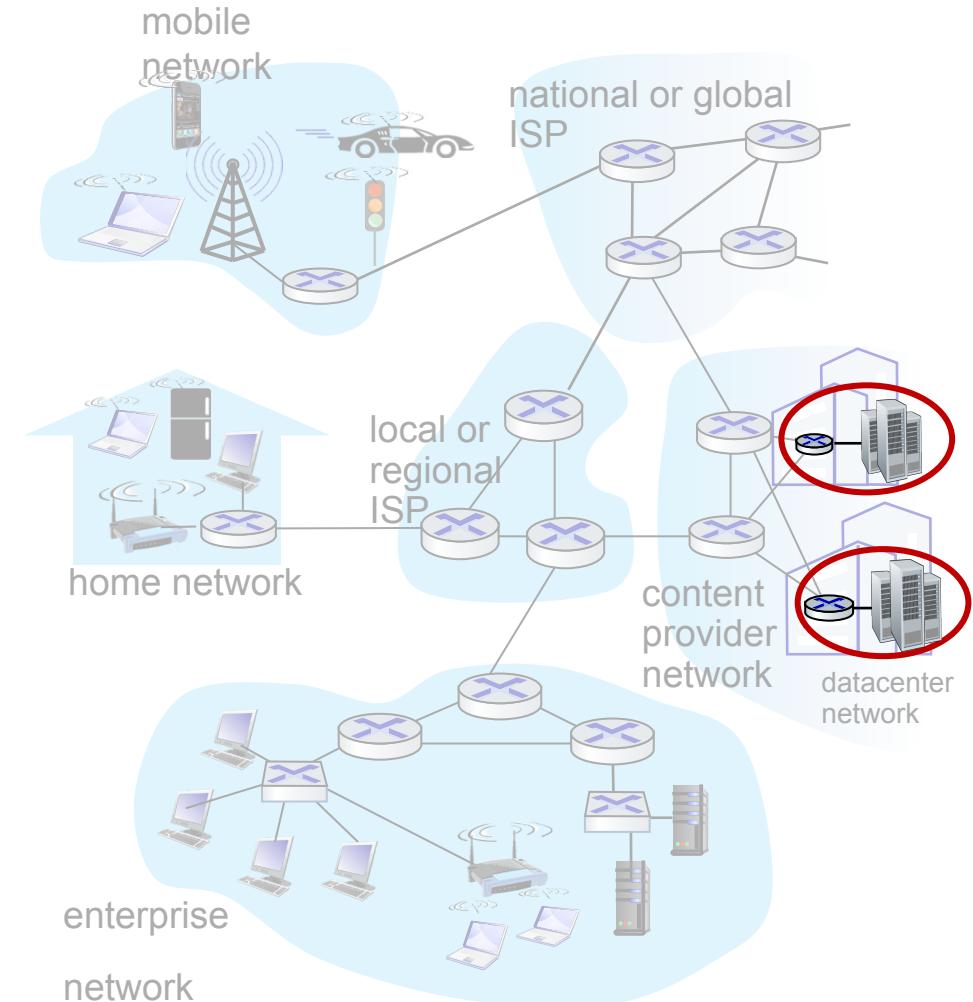
- companies, universities, etc.
 - mix of wired, wireless link technologies, connecting a mix of switches and routers (we'll cover differences shortly)
 - Ethernet: wired access at 100Mbps, 1Gbps, 10Gbps
 - WiFi: wireless access points at 11, 54, 450 Mbps

Access networks: data center networks

- high-bandwidth links (10s to 100s Gbps) connect hundreds to thousands of servers together, and to Internet



Courtesy: Massachusetts Green High Performance Computing Center (mghpcc.org)



Host: sends packets of data

host sending function:

$$\begin{aligned}L &= 10 \text{ bits} \\R &= 5 \text{ bits/sec} \\ \text{Delay} &= 2 \text{ s}\end{aligned}$$

- takes application message
- breaks into smaller chunks, known as *packets*, of length L bits
- transmits packet into access network at *transmission rate R*

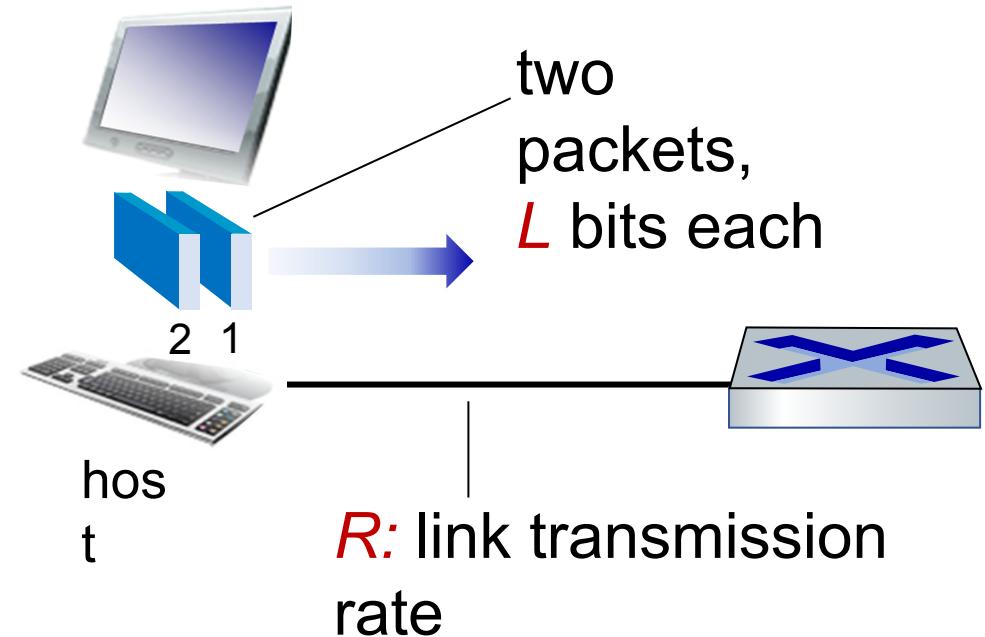
- link transmission rate, aka link *capacity, aka link bandwidth*

packet transmission

delay

time needed
to
transmit L -bit
packet into
link

$$= \frac{L}{R} \text{ (bits/sec)}$$

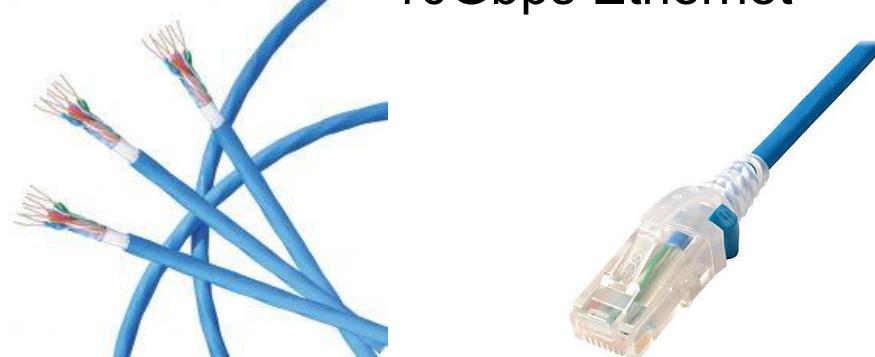


Links: physical media

- **bit**: propagates between transmitter/receiver pairs
- **physical link**: what lies between transmitter & receiver
- **guided media**:
 - signals propagate in solid media: copper, fiber, coax
- **unguided media**:
 - signals propagate freely, e.g., radio

Twisted pair (TP) Par trançado

- two insulated copper wires
 - Category 5: 100 Mbps, 1 Gbps Ethernet
 - Category 6: 10Gbps Ethernet



Links: physical media

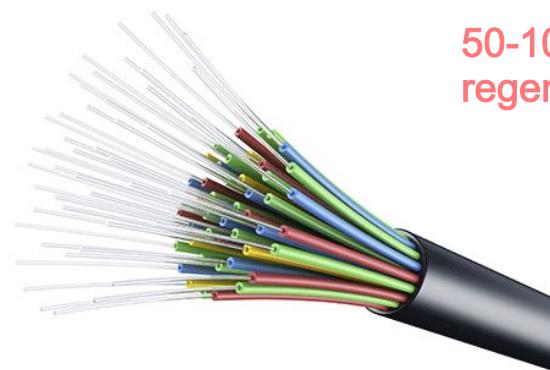
Coaxial cable:

- two concentric copper conductors
- bidirectional
- broadband:
 - multiple frequency channels on cable
 - 100's Mbps per channel



Fiber optic cable:

- glass fiber carrying light pulses, each pulse a bit
- high-speed operation:
 - high-speed point-to-point transmission (10's-100's Gbps)
- low error rate:
 - repeaters spaced far apart
 - immune to electromagnetic noise



Links: physical media

Wireless radio

- signal carried in various “bands” in electromagnetic spectrum
- no physical “wire”
torre de acesso hora recebe hora envia
- broadcast, “half-duplex” (sender to receiver)
- propagation environment effects:
 - reflection
 - obstruction by objects
 - Interference/noise

Radio link types:

- Wireless LAN (WiFi)
 - 10-100's Mbps; 10's of meters
- wide-area (e.g., 4G cellular)
 - 10's Mbps over ~10 Km
- Bluetooth: cable replacement
 - short distances, limited rates
- terrestrial microwave
 - point-to-point; 45 Mbps channels
Esse delay de 270 não era para estar contido no 45 Mbps? Não era pra taxa de transmissão ser menor?
- satellite
 - up to 45 Mbps per channel
 - 270 msec end-end delay

Dados o tamanho de um pacote (L) e a taxa de transmissão (R) de um enlace. A equação L/R define:

A sua resposta :

O retardo de transmissão do enlace

Quais os elementos listados abaixo que compõem a borda da rede:

A sua resposta :

Dispositivos hospedeiros

Retorno:

Correto.

Redes de acesso

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Retorno:

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Enlaces de comunicação

Retorno:

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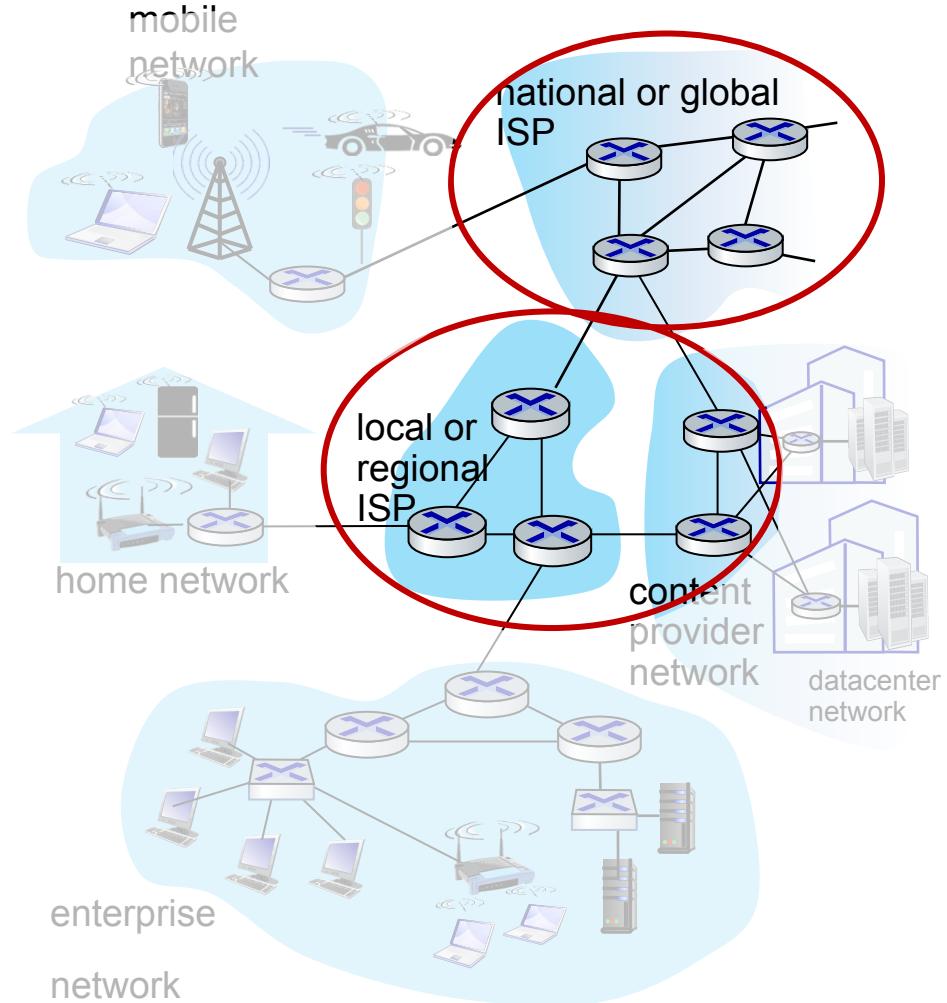
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The network core

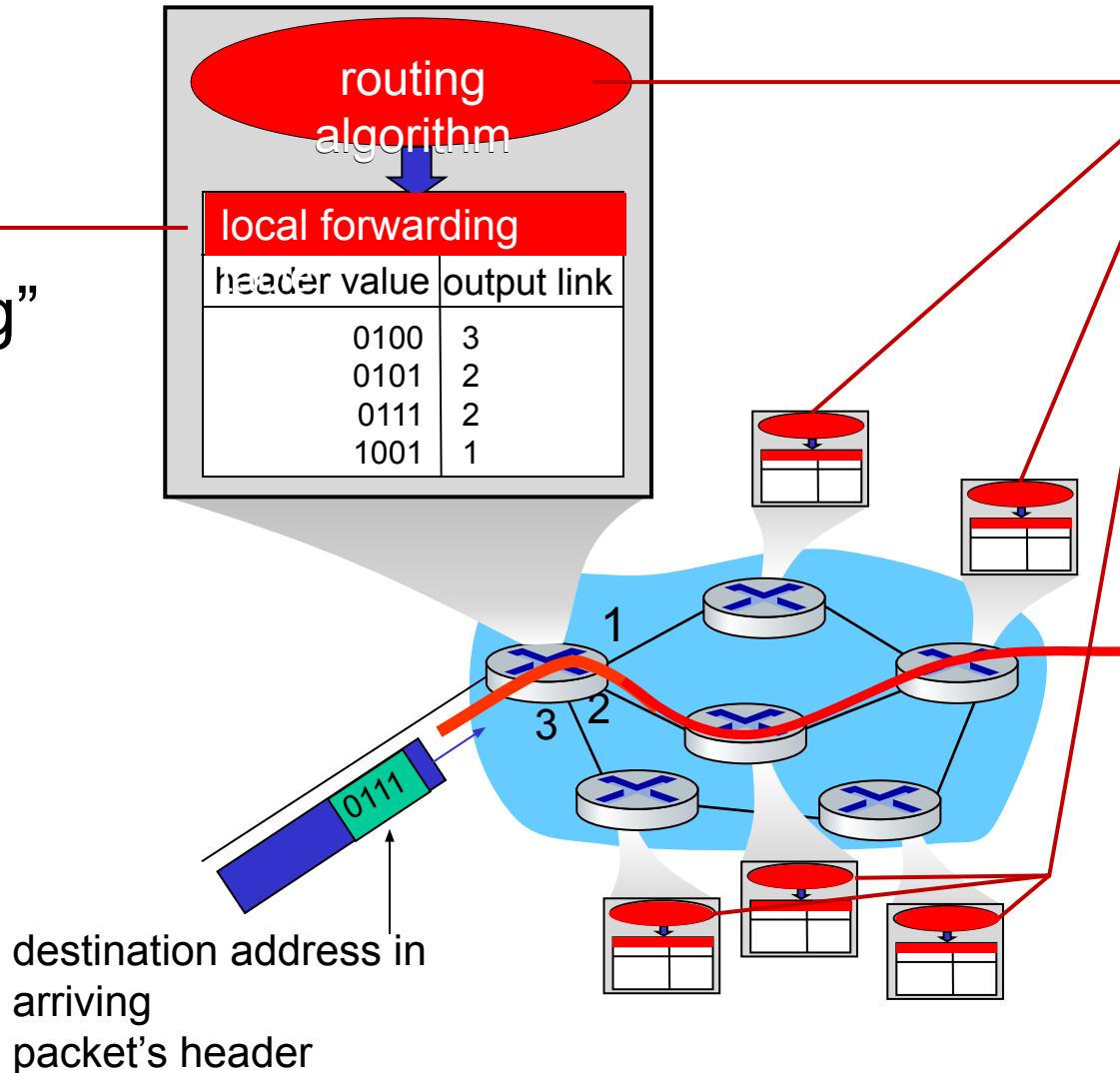
- mesh of interconnected routers
- **packet-switching**: hosts break application-layer messages into *packets*
 - network **forwards** packets from one router to the next, across links on path from **source to destination**



Two key network-core functions

Forwarding:

- aka “switching”
- *local* action: move arriving packets from router’s input link to appropriate router output link



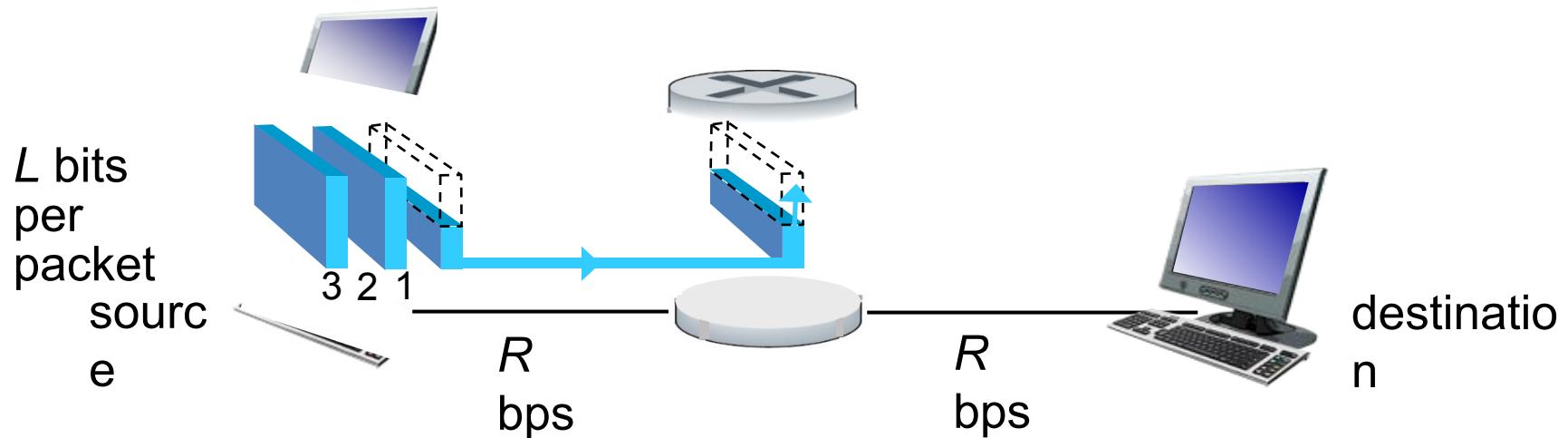
Routing:

- *global* action: determine source-destination paths taken by packets
- routing algorithms





Packet-switching: store-and-forward

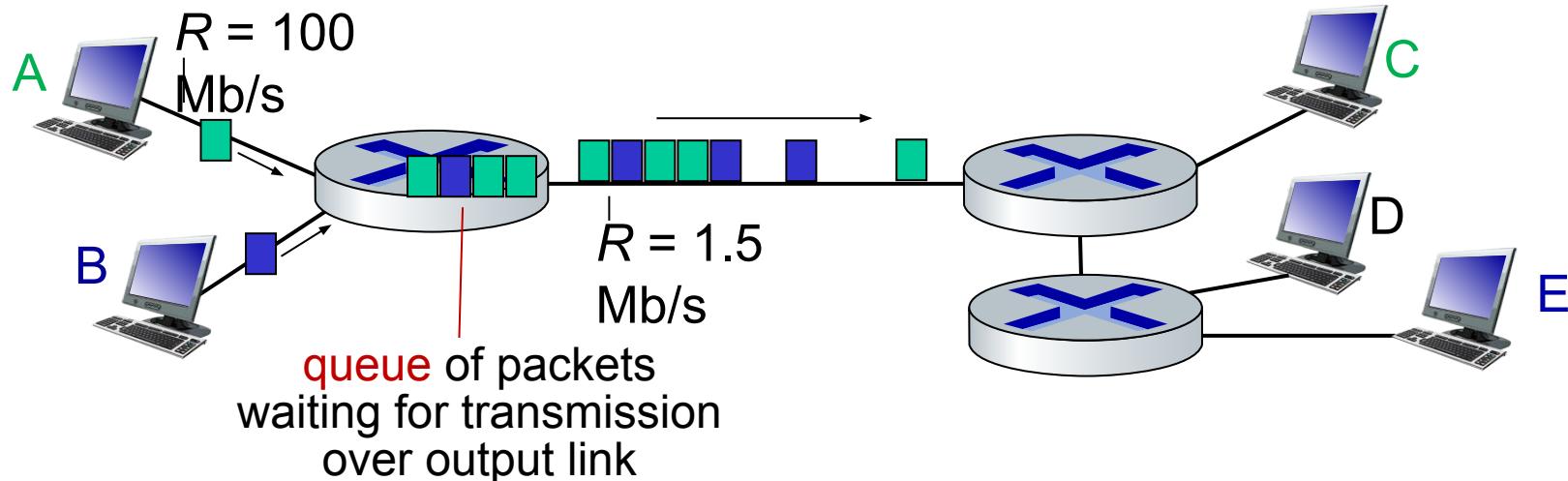


- **packet transmission delay:** takes L/R seconds to transmit (push out) L -bit packet into link at R bps
- **store and forward:** entire packet must arrive at router before it can be transmitted on next link

One-hop numerical example:

- $L = 10$ Kbits
- $R = 100$ Mbps
- one-hop transmission delay = 0.1 msec

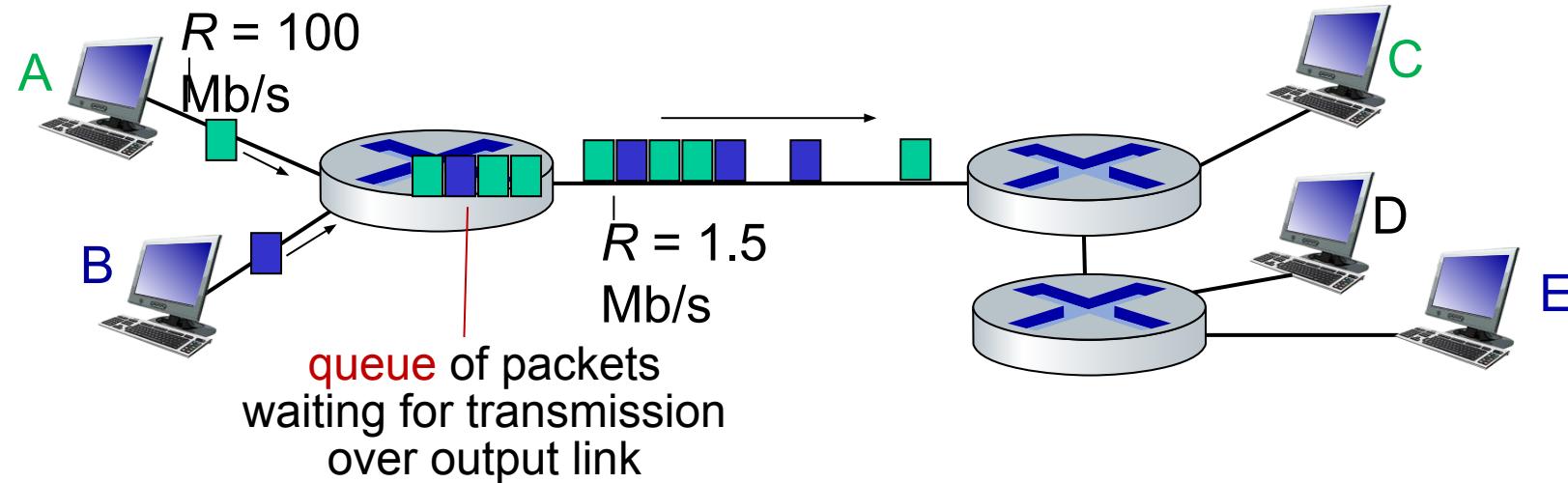
Packet-switching: queueing



Queueing occurs when work arrives faster than it can be serviced:



Packet-switching: queueing



Packet queuing and loss: if arrival rate (in bps) to link exceeds transmission rate (bps) of link for some period of time:

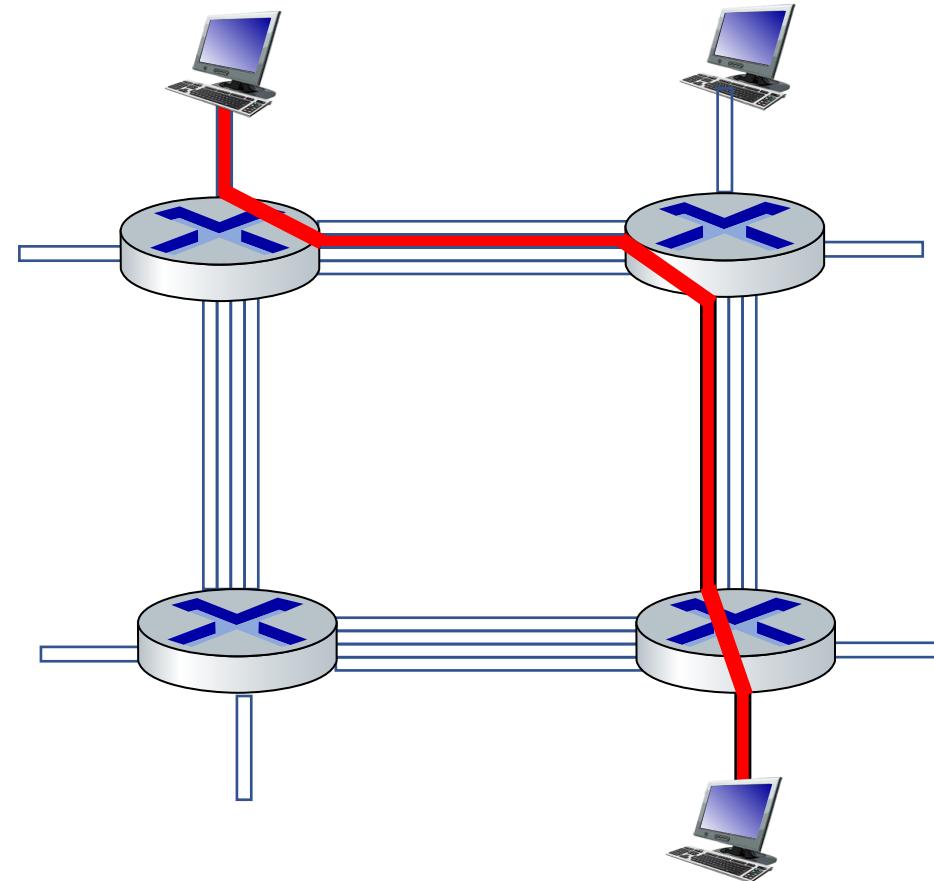
- packets will queue, waiting to be transmitted on output link
- packets can be dropped (lost) if memory (buffer) in router fills up

Alternative to packet switching: circuit switching

end-end resources allocated to, reserved for “call” between source and destination

- in diagram, each link has four circuits.
 - call gets 2nd circuit in top link and 1st circuit in right link.
- dedicated resources: no sharing
 - circuit-like (guaranteed) performance
- circuit segment idle if not used by call
- **(no sharing)**
- commonly used in traditional telephone networks

- pacotes tem a informação sob a qual circuito virtual ele pertence em vez de qual o prox roteador ele vai.



* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive

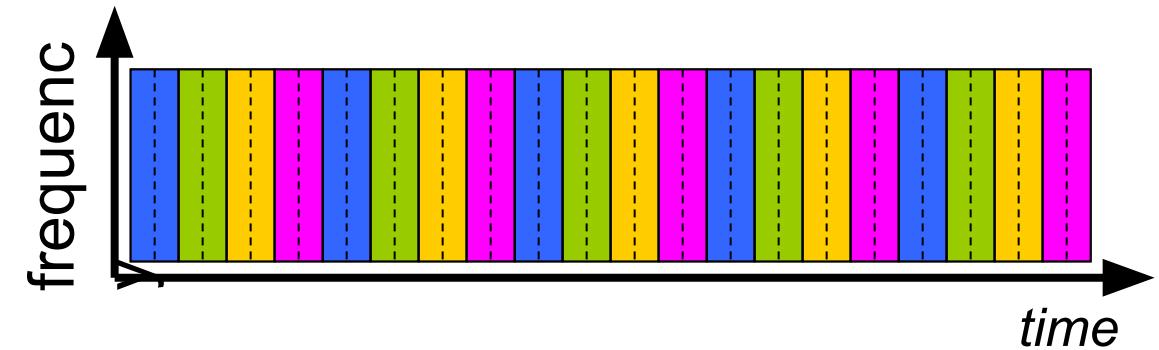
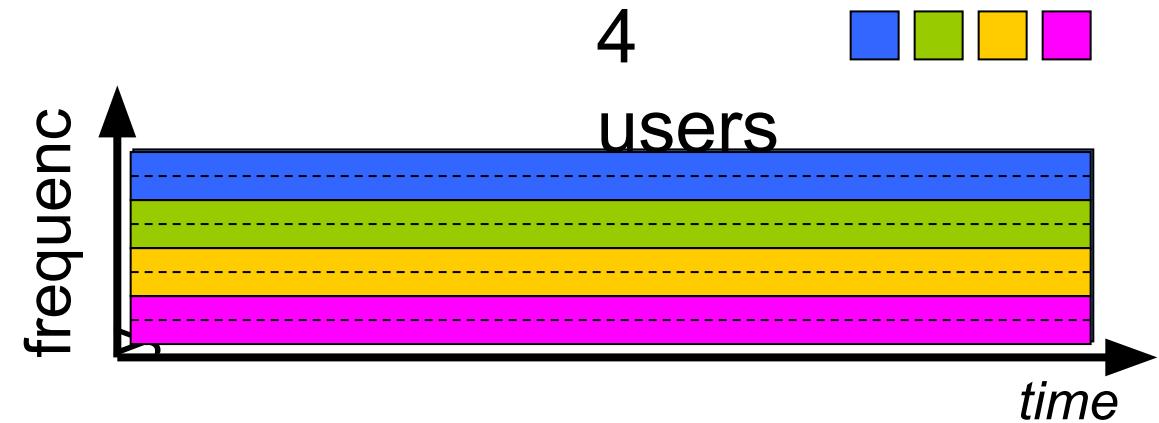
Circuit switching: FDM and TDM

Frequency Division Multiplexing (FDM)

- optical, electromagnetic frequencies divided into (narrow) frequency bands
- each call allocated its own band, can transmit at max rate of that narrow band

Time Division Multiplexing (TDM)

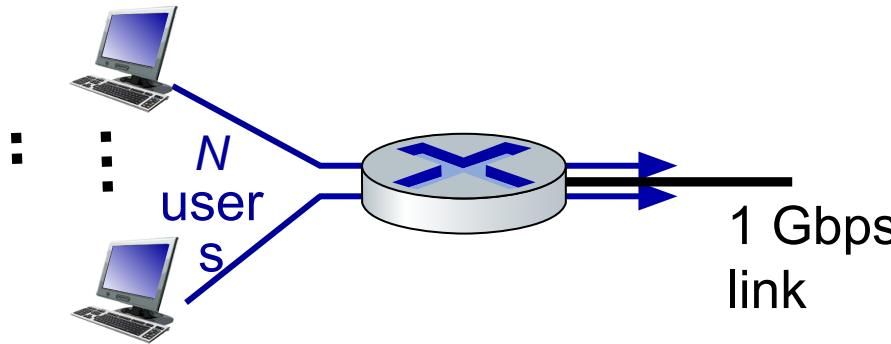
- time divided into slots
- each call allocated periodic slot(s), can transmit at maximum rate of (wider) frequency band (only) during its time slot(s)



Packet switching versus circuit switching

example:

- 1 Gb/s link
- each user:
 - 100 Mb/s when “active”
 - active 10% of time



Q: how many users can use this network under circuit-switching and packet switching?

- **circuit-switching:** 10 users
- **packet switching:** with 35 users, probability > 10 active at same time is less than .0004 *

Q: how did we get value
A:.0004 problem (for those with course in probability only)

* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive

Packet switching versus circuit switching

Is packet switching a “slam dunk

• winner?

- great for “bursty” data – sometimes has data to send, but at other times not

por pacote é bom quando é pouco tempo o uso do recurso de cada computador.

- resource sharing
- simpler, no call setup

- **excessive congestion possible:** packet delay and loss due to buffer overflow

- protocols needed for reliable data transfer, congestion control

- **Q: How to provide circuit-like behavior with packet-switching?**

- “It’s complicated.” We’ll study various techniques that try to make

Q: Human analogies of reserved resources (circuit switching) versus on-demand allocation (packet switching)?

O núcleo da Internet é composto por uma malha de roteadores, comutadores e sistemas finais.

A sua resposta :

Falso

Retorno:

Resposta correta. O núcleo da rede não possui comutadores ou sistemas finais em sua infraestrutura.

Pacotes transmitidos no núcleo da Internet são recebidos, analisados e encaminhados, em cada roteador, ao longo do seu caminho, entre a origem e o destino da comunicação.

A sua resposta :

Verdadeiro

Retorno:

Resposta correta.

As principais funcionalidades do núcleo da Internet são:

A sua resposta :

Encaminhamento de pacotes

Retorno:

Resposta correta.

Roteamento de pacotes

Retorno:

Resposta correta.

Analise afirmação abaixo e marque as opções que melhor se aderem aos conceitos vistos sobre com
Identifique vantagens da comutação por pacotes.

A sua resposta :

Não requer o estabelecimento de uma conexão na rede.

Retorno:

Resposta correta.

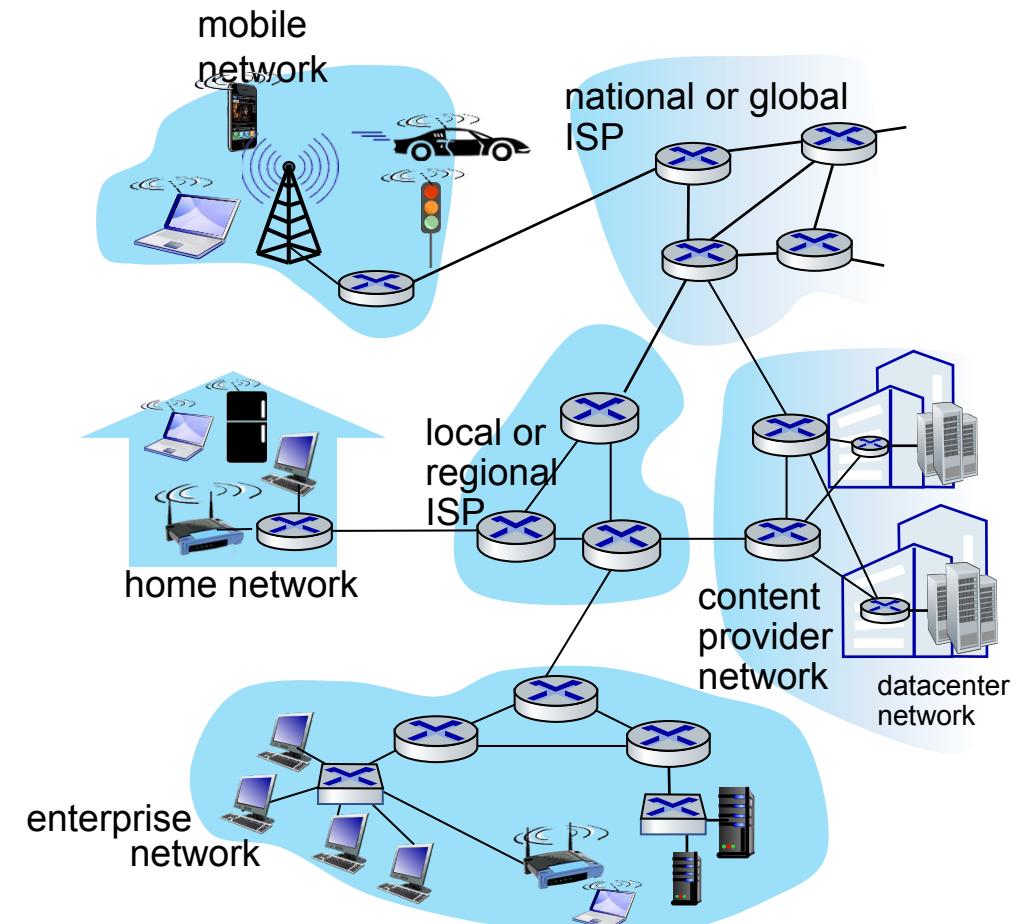
Evita o congestionamento na rede.

Retorno:

Resposta incorreta. Este é um conceito associado a alocação de recurso prévio, visto no modelo cir

Internet structure: a “network of networks”

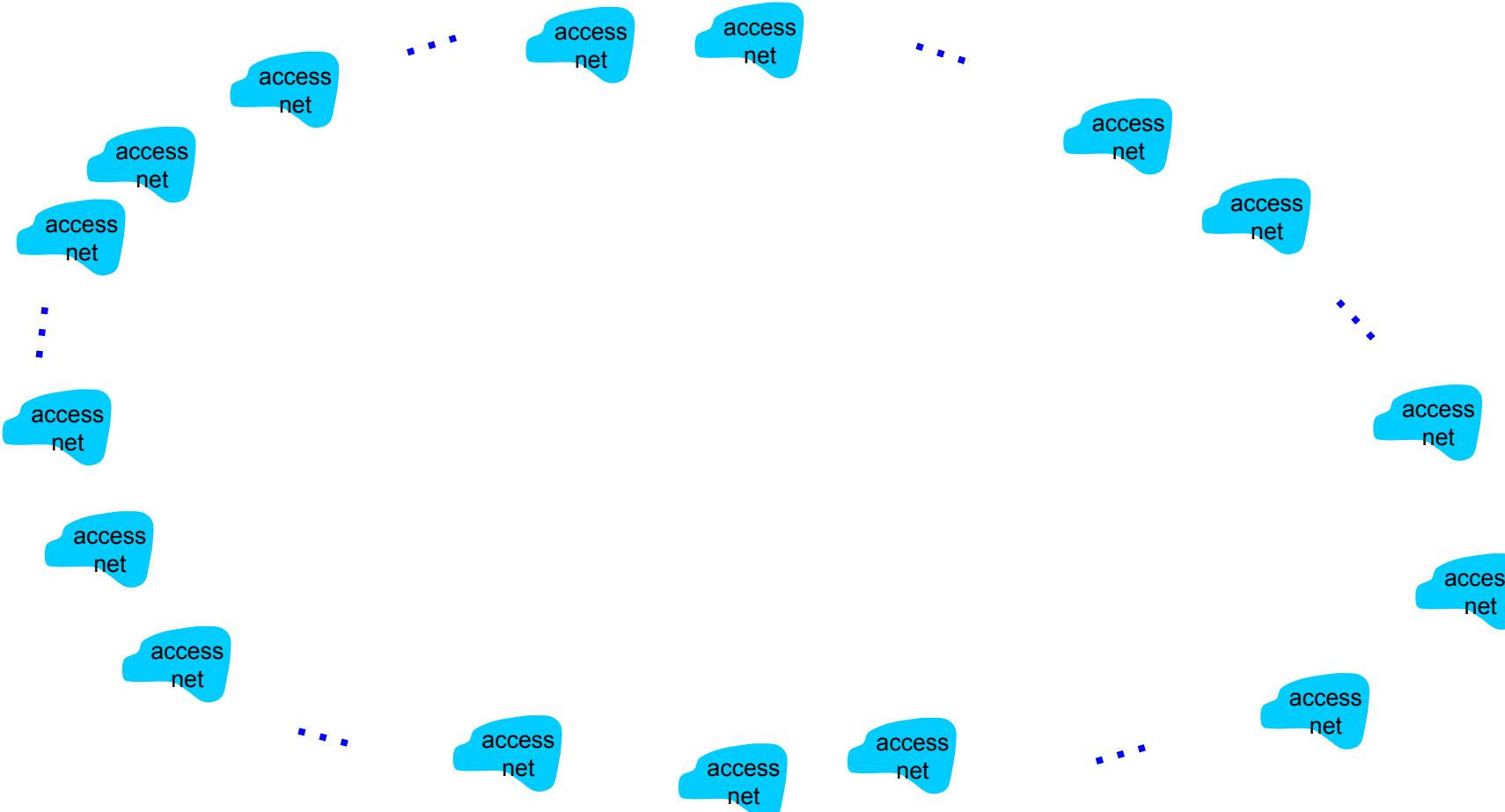
- hosts connect to Internet via **access** Internet Service Providers (ISPs)
- access ISPs in turn must be interconnected
 - so that *any two hosts (anywhere!)* can send packets to each other
- resulting network of networks is very complex
 - evolution driven by **economics, national policies**



Let's take a stepwise approach to describe current Internet structure

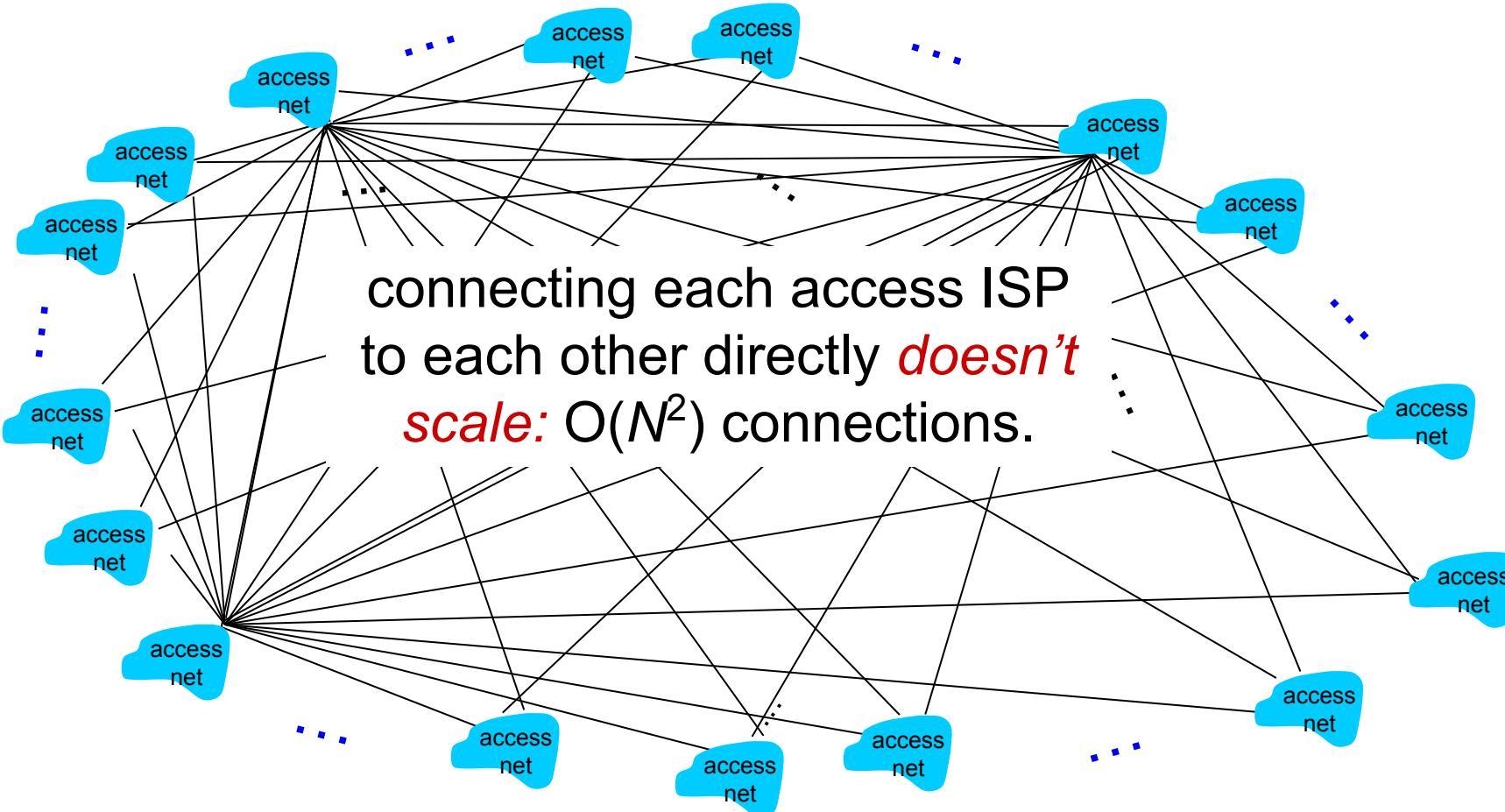
Internet structure: a “network of networks”

Question: given *millions* of access ISPs, how to connect them together?



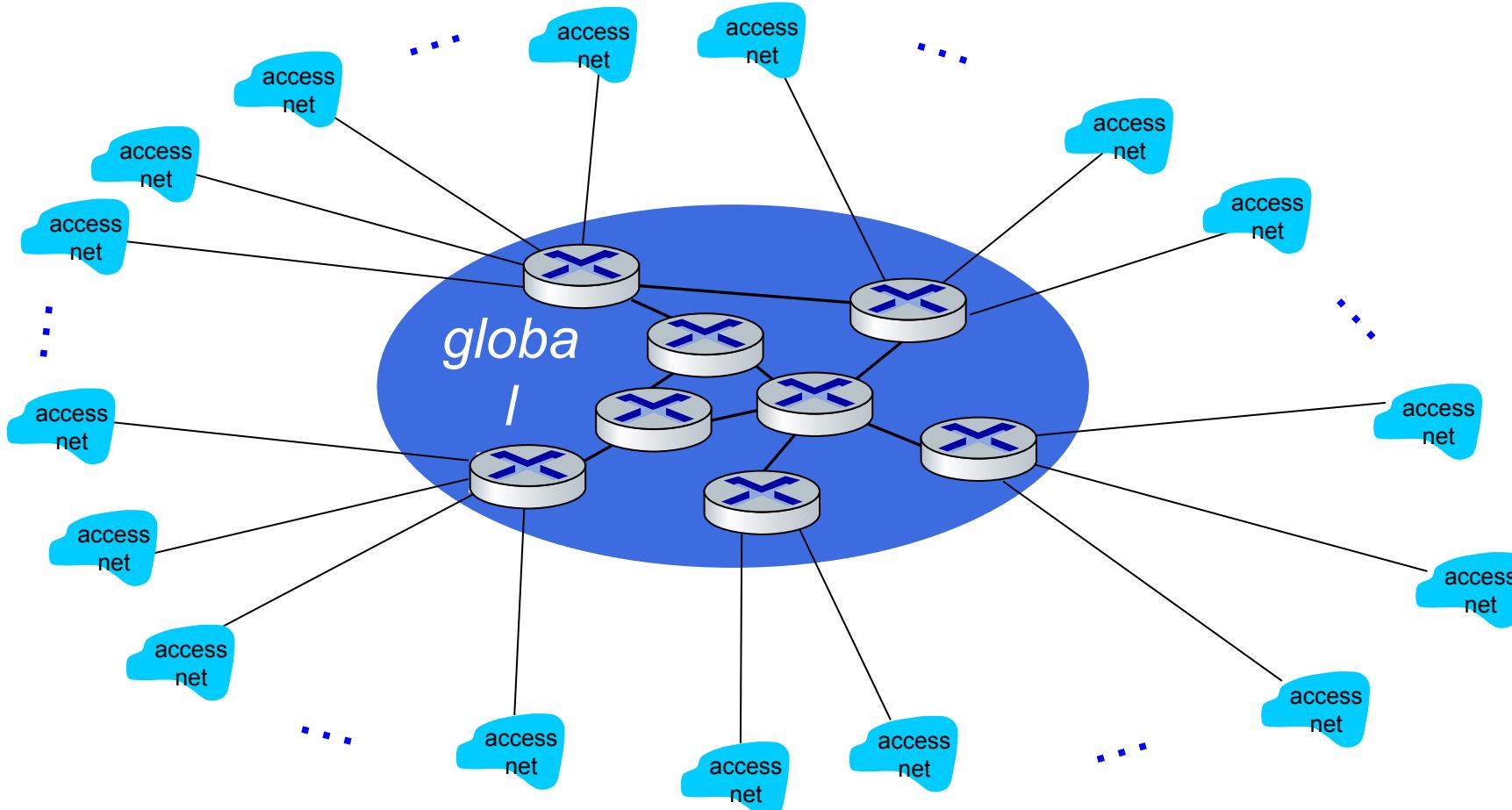
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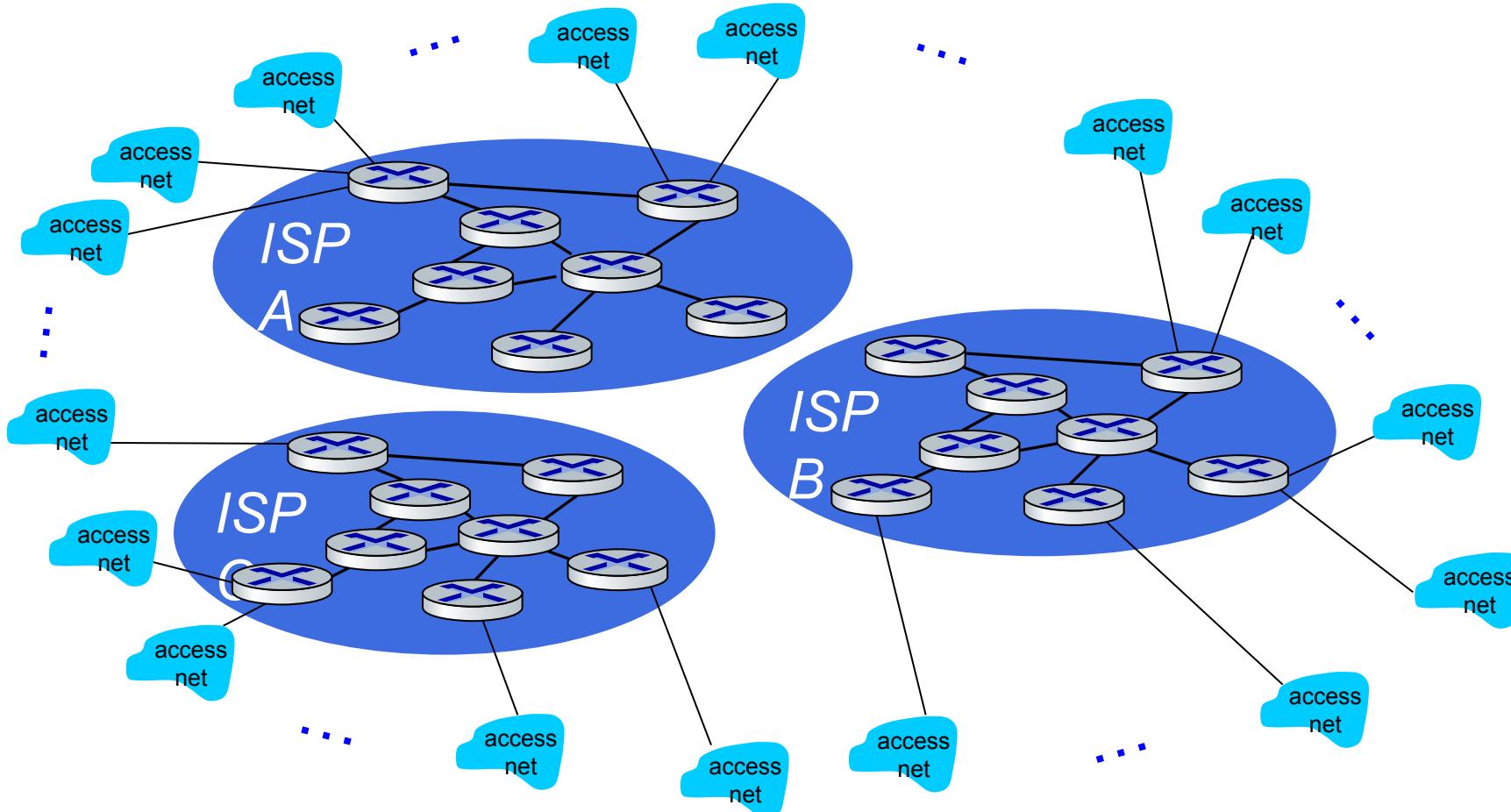
Internet structure: a “network of networks”

Option: connect each access ISP to one global transit ISP?
Customer and provider ISPs have economic agreement.



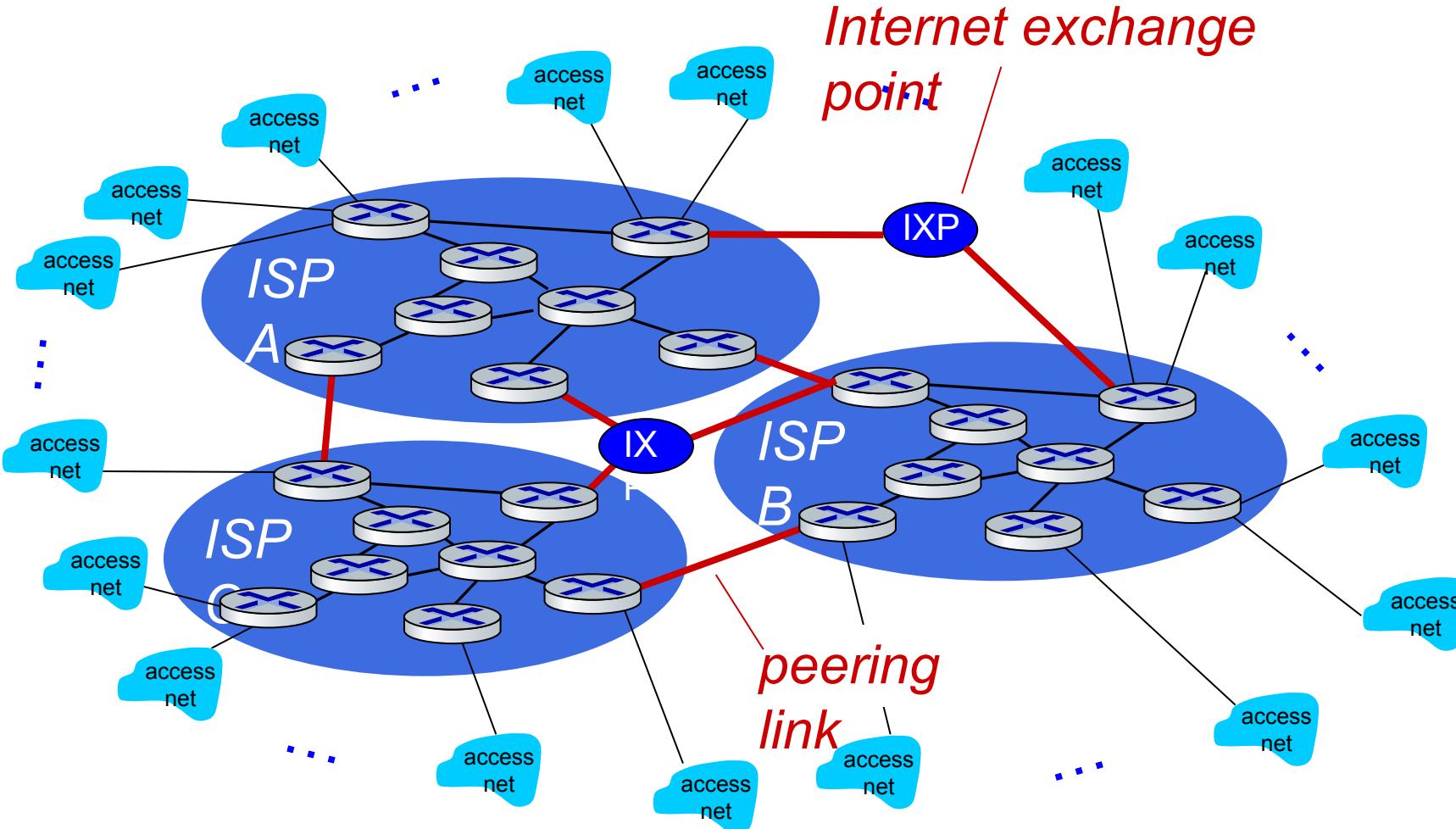
Internet structure: a “network of networks”

But if one global ISP is viable business, there will be competitors



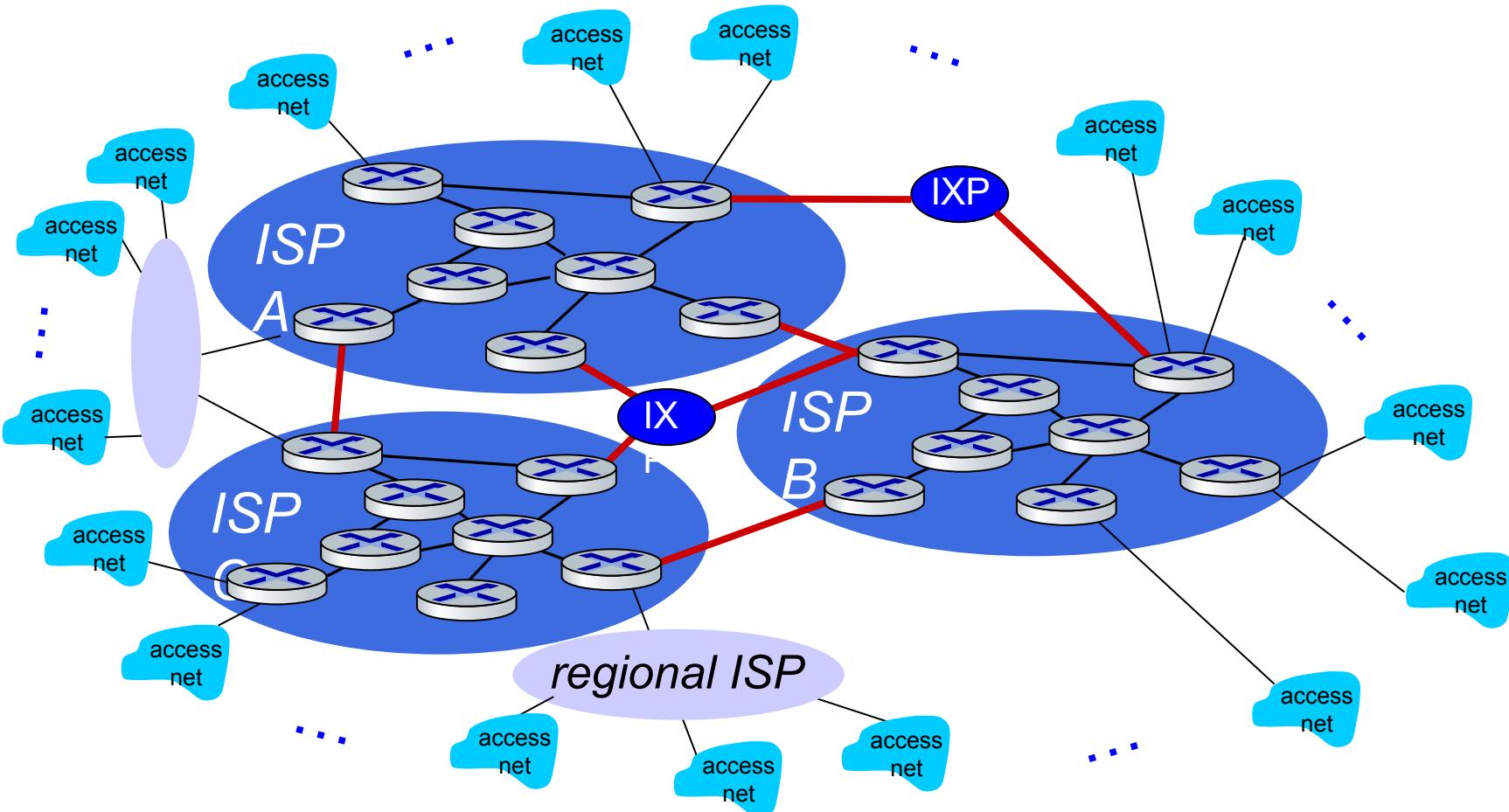
Internet structure: a “network of networks”

But if one global ISP is viable business, there will be competitors
who will want to be connected



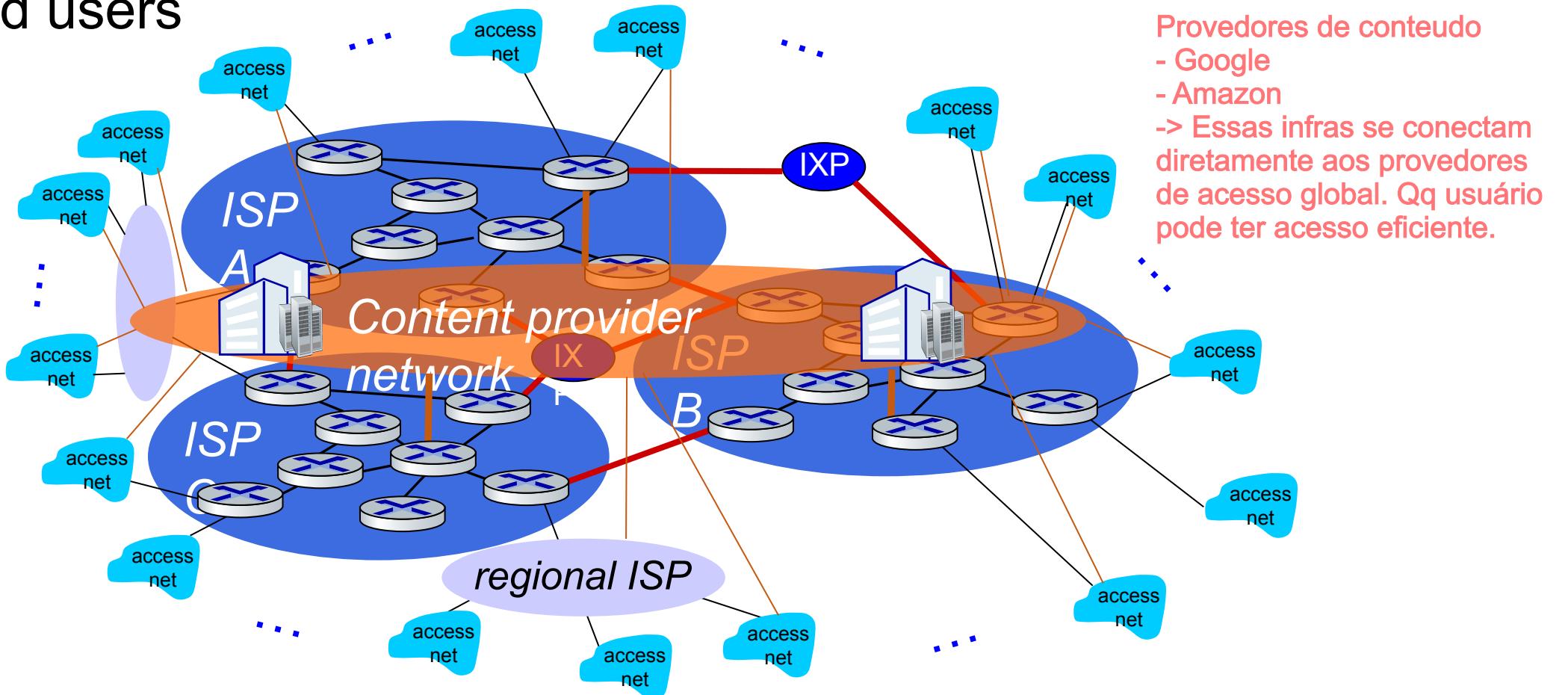
Internet structure: a “network of networks”

... and regional networks may arise to connect access nets to ISPs

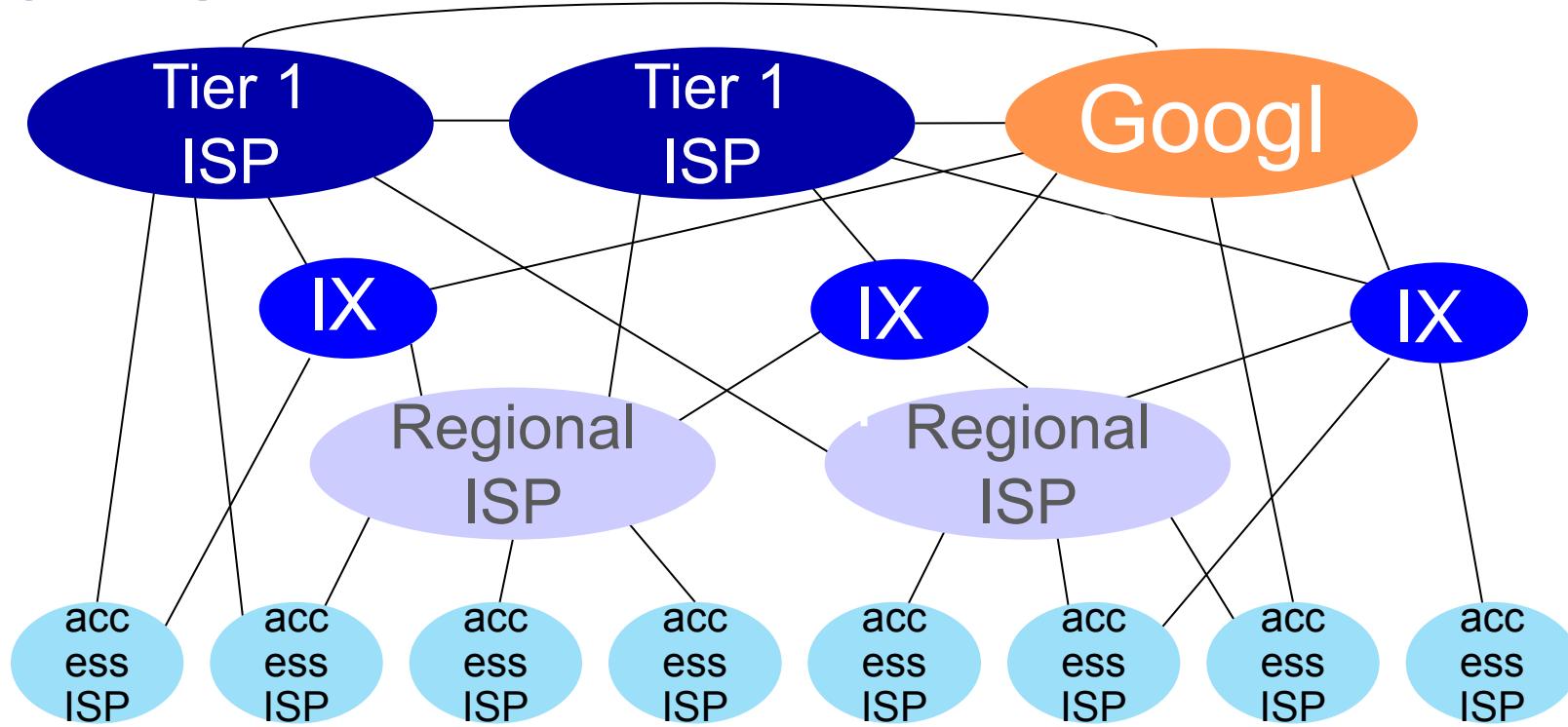


Internet structure: a “network of networks”

... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users



Internet structure: a “network of networks”



At “center”: small # of well-connected large networks

- **“tier-1” commercial ISPs** (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
- **content provider networks** (e.g., Google, Facebook): private network that connects its data centers to Internet, often bypassing tier-1, regional ISPs

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How do packet delay and loss occur?

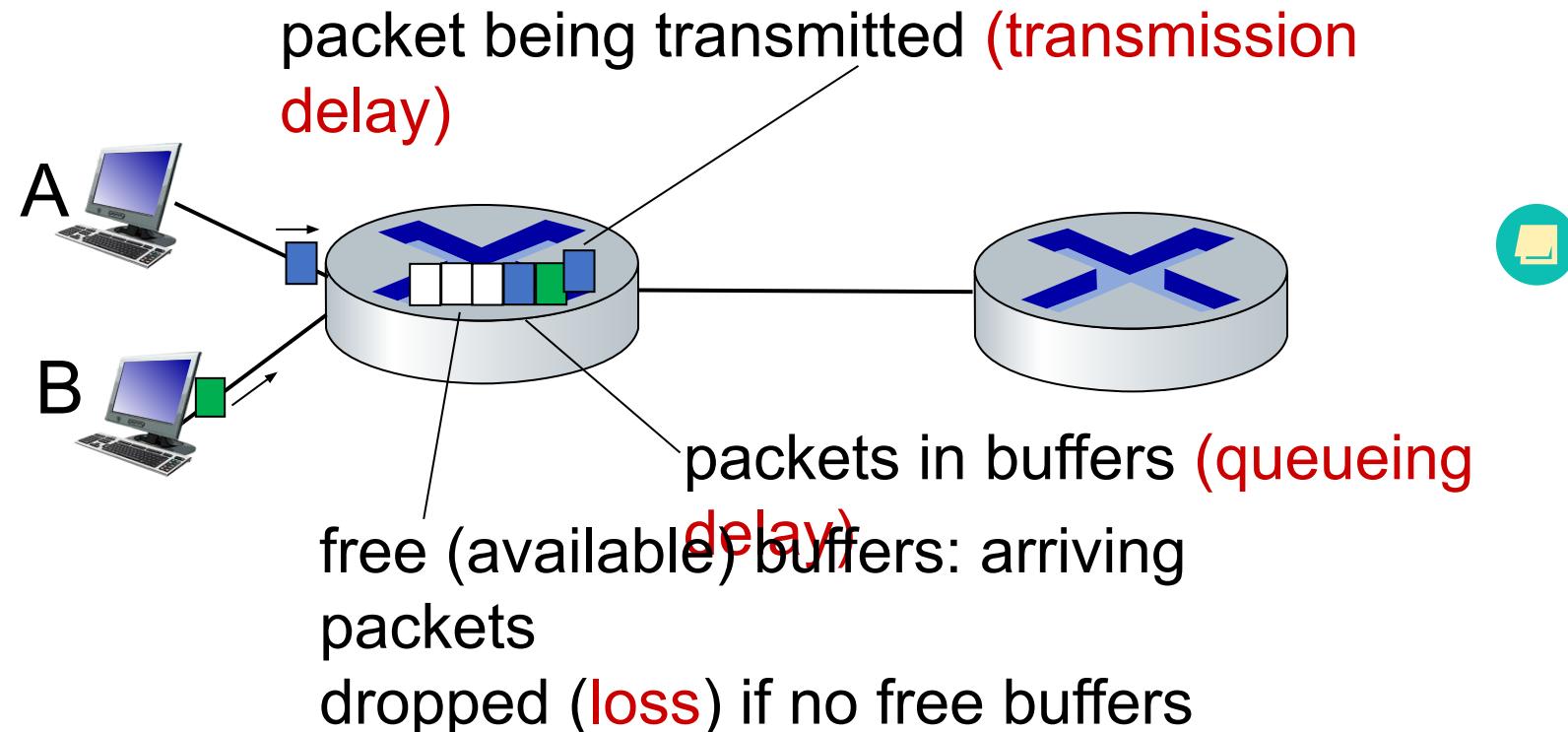
- packets *queue* in router buffers, waiting for turn for transmission
 - queue length grows when arrival rate to link (temporarily) exceeds output link capacity
- packet *loss* occurs when memory to hold queued packets fills up

processamento: checar a integridade, endereço de saída (tempo pequeno)

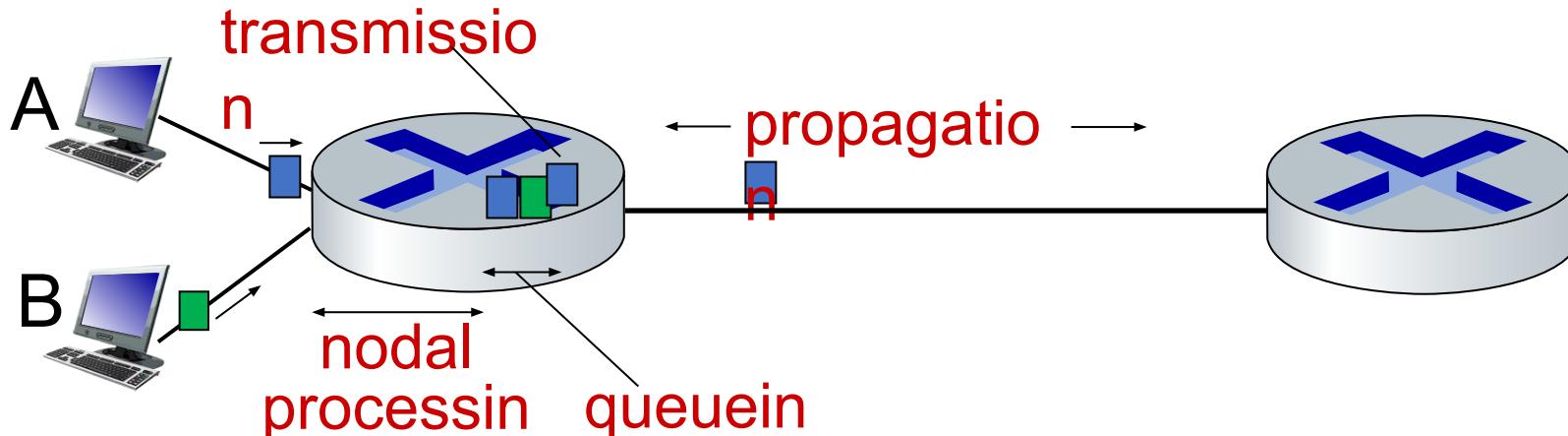
enfileiramento: fila esperando na fila

transmissão: tempo necessário para transmitir o pacote no enlace $t = L/R$ (taxa de transmissão do enlace)

propagação: tempo necessário para que os bits se propaguem até o destino. comprimento do enlace e características. quanto maior o enlace, maior o tempo



Packet delay: four sources



$$d_{\text{nodal}} = d_{\text{proc}}^{\text{g}} + d_{\text{queue}}^{\text{g}} + d_{\text{trans}} + d_{\text{prop}}$$

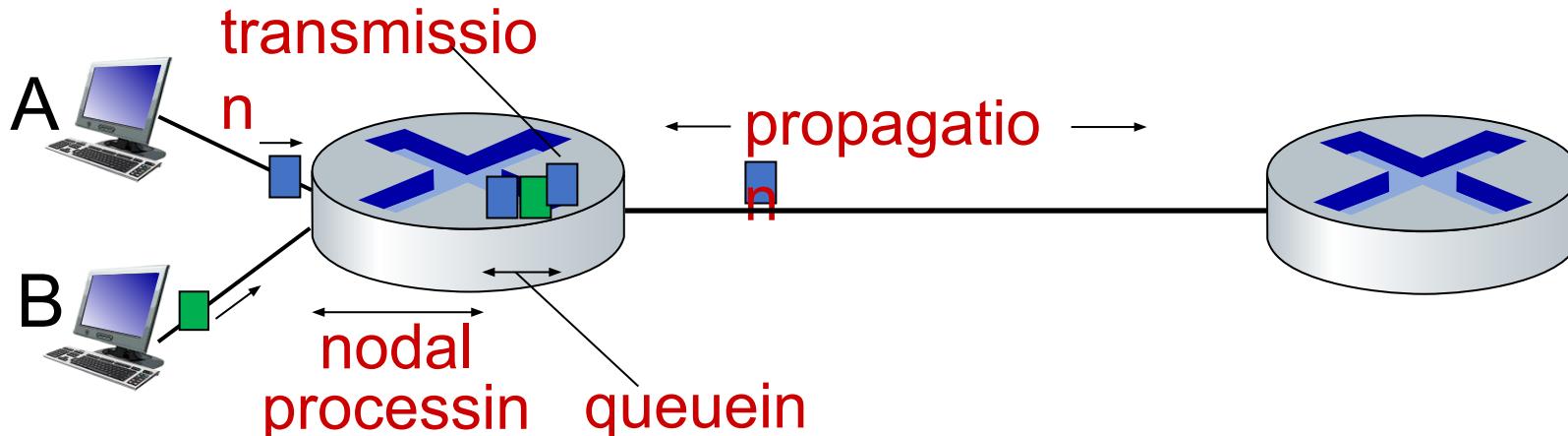
d_{proc} : nodal processing

- check bit errors
- determine output link
- typically < microsecs

d_{queue} : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

Packet delay: four sources



$$d_{\text{nodal}} = d_{\text{proc}}^{\text{g}} + d_{\text{queue}}^{\text{g}} + d_{\text{trans}} +$$

d_{trans} : **transmission delay:**

- L : packet length (bits)
- R : link *transmission rate*

(bps)

$$\bullet d_{\text{trans}} = L/R$$

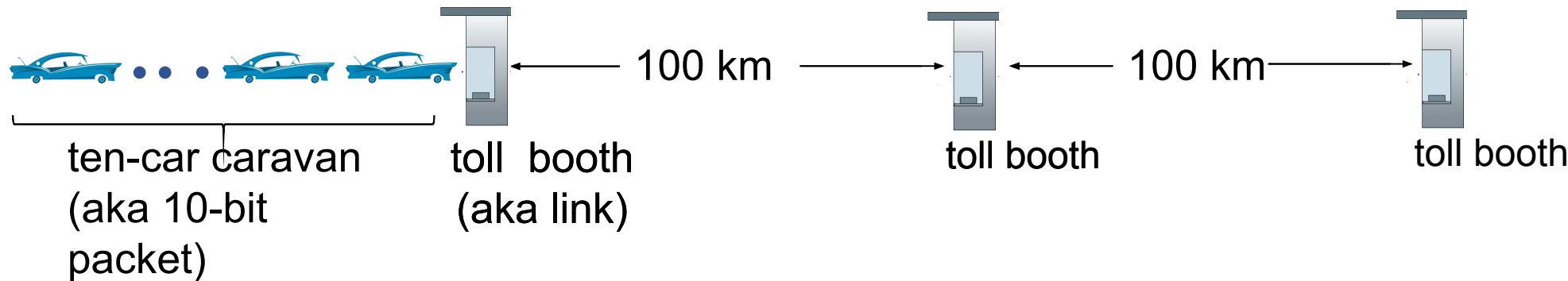
d_{trans} and d_{prop}
very different

d_{prop} : **propagation delay:**

- d : length of physical link
- s : propagation speed ($\sim 2 \times 10^8$ m/sec)

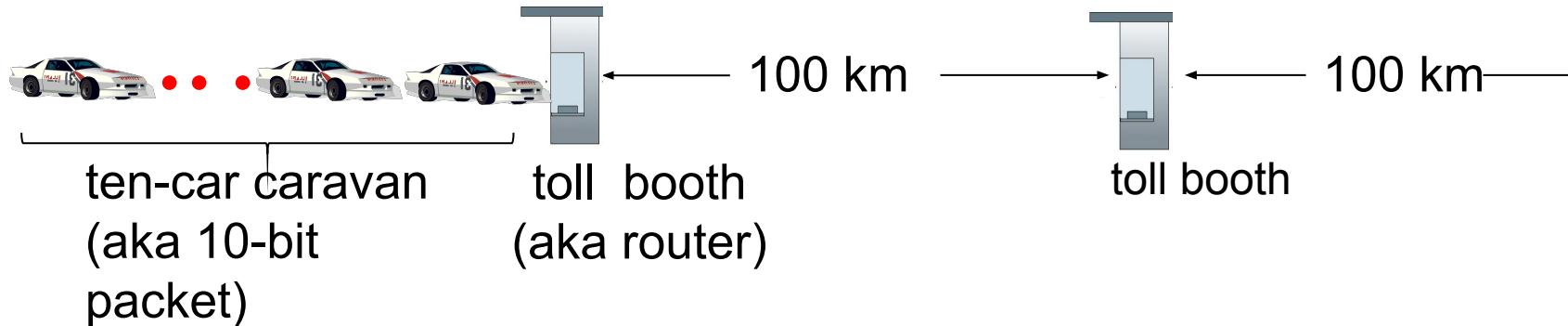
$$\bullet d_{\text{prop}} = d/s$$

Caravan analogy



- car ~ bit; caravan ~ packet; toll service ~ link transmission
- toll booth takes 12 sec to service car (bit transmission time)
- “propagate” at 100 km/hr
- **Q: How long until caravan is lined up before 2nd toll booth?**
- time to “push” entire caravan through toll booth onto highway = $12 \times 10 = 120$ sec
- time for last car to propagate from 1st to 2nd toll both: $100\text{km}/(100\text{km/hr}) = 1$ hr
- **A: 62 minutes**

Caravan analogy



- suppose cars now “propagate” at 1000 km/hr
- and suppose toll booth now takes one min to service a car
- **Q: Will cars arrive to 2nd booth before all cars serviced at first booth?**

pacote de 10 bits

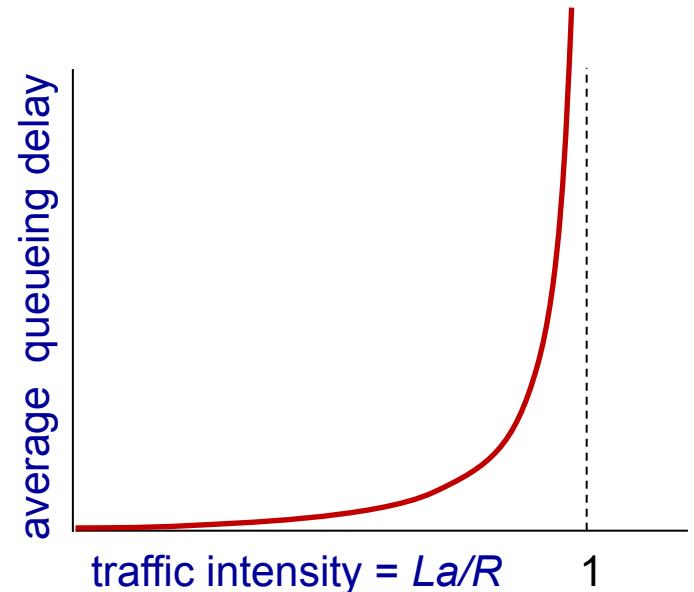
retardo de enfileiramento: como se comporta quando comparamos a taxa de entrada de tráfego com a capacidade de transmissão de saída.

Packet queueing delay (revisited)

- a : average packet arrival rate
- L : packet length (bits)
- R : link bandwidth (bit transmission rate)

$$\frac{L \cdot a}{R} : \frac{\text{arrival rate of bits}}{\text{Service rate of bits}}$$

“traffic intensity”



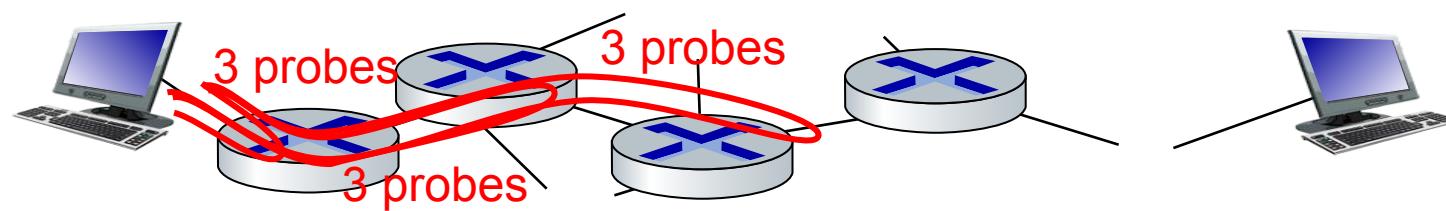
- $La/R \sim 0$: avg. queueing delay small
- $La/R \rightarrow 1$: avg. queueing delay large
- $La/R > 1$: more “work” arriving is more than can be serviced -



$La/R \sim 0$
 $La/R \rightarrow 1$

“Real” Internet delays and routes

- what do “real” Internet delay & loss look like?
- **traceroute** program: provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination (with time-to-live field value of i)
 - router i will return packets to sender
 - sender measures time interval between transmission and reply



Real Internet delays and routes

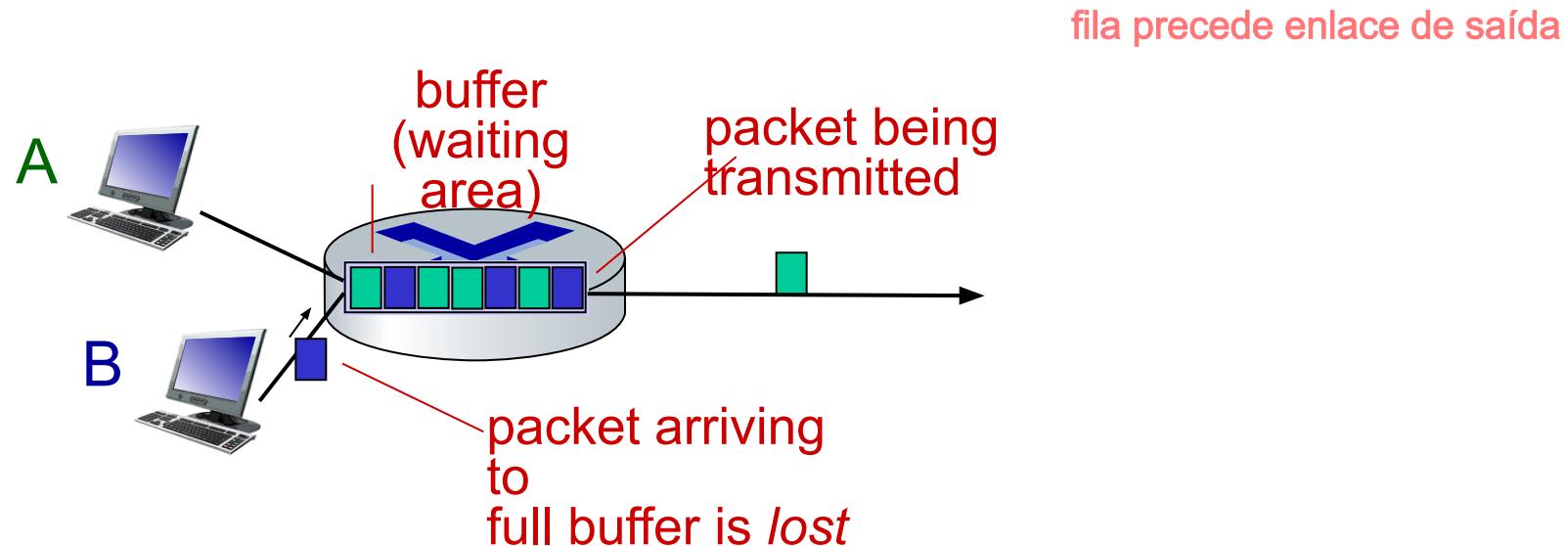
traceroute: gaia.cs.umass.edu to www.eurecom.fr

		3 delay measurements from gaia.cs.umass.edu to cs-gw.cs.umass.edu			
1	cs-gw (128.119.240.254)	1 ms	1 ms	2 ms	
2	border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145)	1 ms	1 ms	2 ms	3 delay measurements to border1-rt-fa5-1-0.gw.umass.edu
3	cht-vbns.gw.umass.edu (128.119.3.130)	6 ms	5 ms	5 ms	
4	jn1-at1-0-0-19.wor.vbns.net (204.147.132.129)	16 ms	11 ms	13 ms	
5	jn1-so7-0-0-0.wae.vbns.net (204.147.136.136)	21 ms	18 ms	18 ms	
6	abilene-vbns.abilene.ucaid.edu (198.32.11.9)	22 ms	18 ms	22 ms	
7	nycm-wash.abilene.ucaid.edu (198.32.8.46)	22 ms	22 ms	22 ms	trans-oceanic link
8	62.40.103.253 (62.40.103.253)	104 ms	109 ms	106 ms	
9	de2-1.de1.de.geant.net (62.40.96.129)	109 ms	102 ms	104 ms	
10	de.fr1.fr.geant.net (62.40.96.50)	113 ms	121 ms	114 ms	
11	renater-gw.fr1.fr.geant.net (62.40.103.54)	112 ms	114 ms	112 ms	looks like delays decrease! Why?
12	nio-n2.cssi.renater.fr (193.51.206.13)	111 ms	114 ms	116 ms	
13	nice.cssi.renater.fr (195.220.98.102)	123 ms	125 ms	124 ms	
14	r3t2-nice.cssi.renater.fr (195.220.98.110)	126 ms	126 ms	124 ms	
15	eurecom-valbonne.r3t2.ft.net (193.48.50.54)	135 ms	128 ms	133 ms	retardo de propagação mto grande "paulatinamente"
16	194.214.211.25 (194.214.211.25)	126 ms	128 ms	126 ms	
17	***				
18	***	* means no response (probe lost, router not replying)			
19	fantasia.eurecom.fr (193.55.113.142)	132 ms	128 ms	136 ms	

* Do some traceroutes from exotic countries at
www.traceroute.org

Packet loss

- queue (aka buffer) preceding link in buffer has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not at all



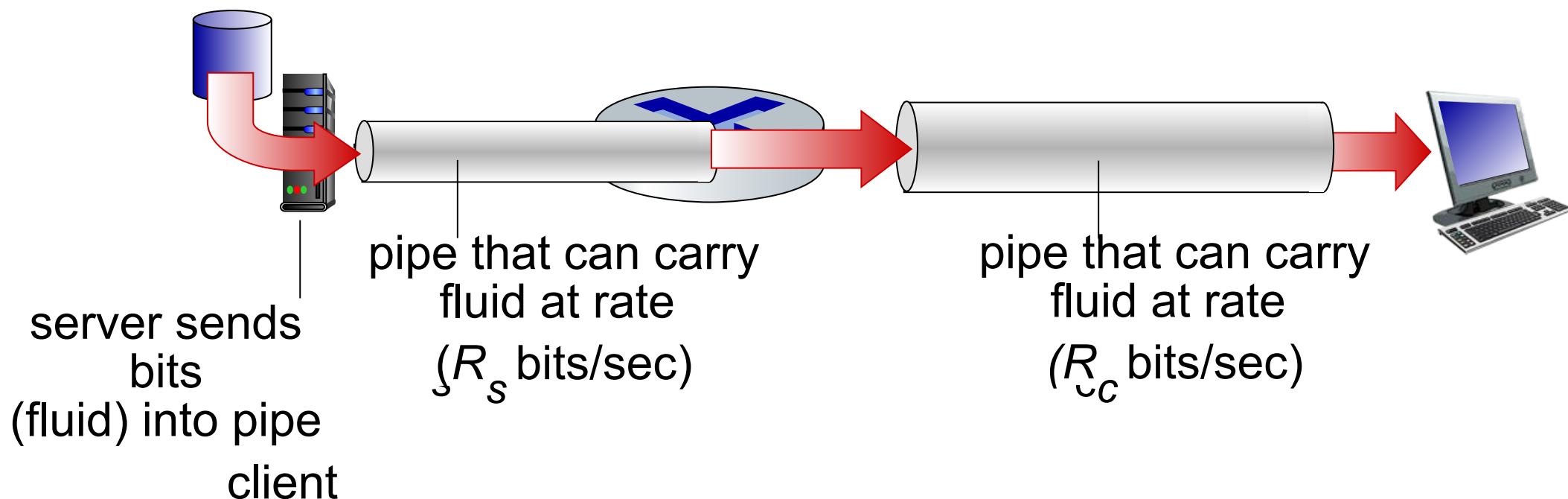
* Check out the Java applet for an interactive animation (on publisher's website) of queuing and loss

Throughput

- *throughput*: rate (bits/time unit) at which bits are being sent from sender to receiver
 - *instantaneous*: rate at given point in time
 - *average*: rate over longer period of time

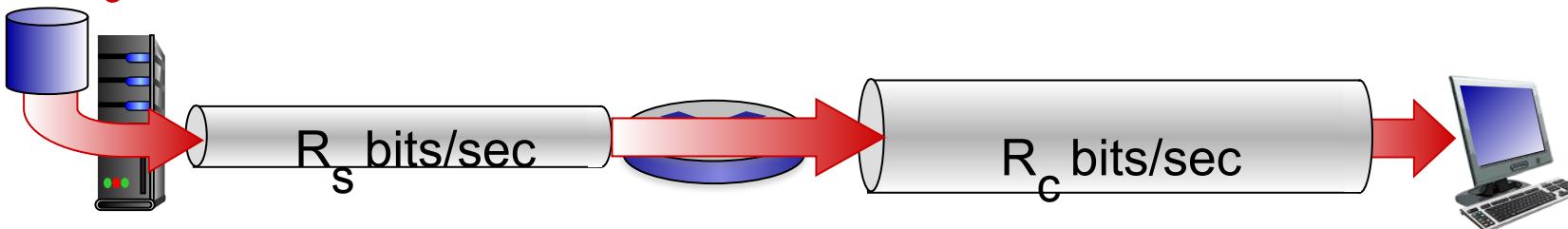
vazão é a taxa de transmissão.

qual a taxa de transmissão do enlace
x
taxa de transmissão efetiva,
experimentada por uma conexão.

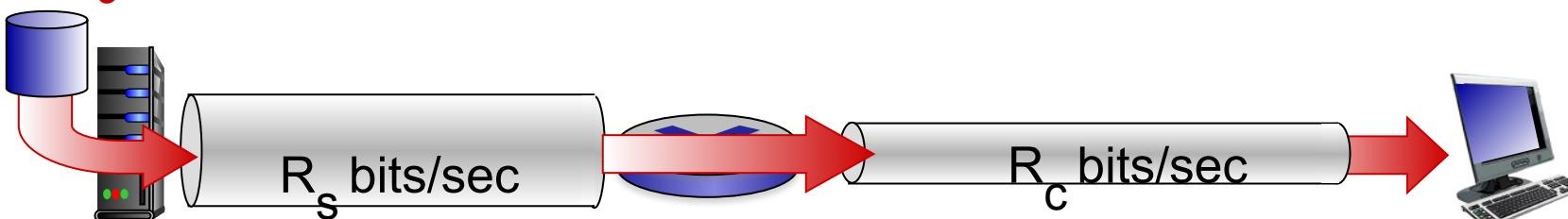


Throughput

$R_s < R_c$ What is average end-end throughput?



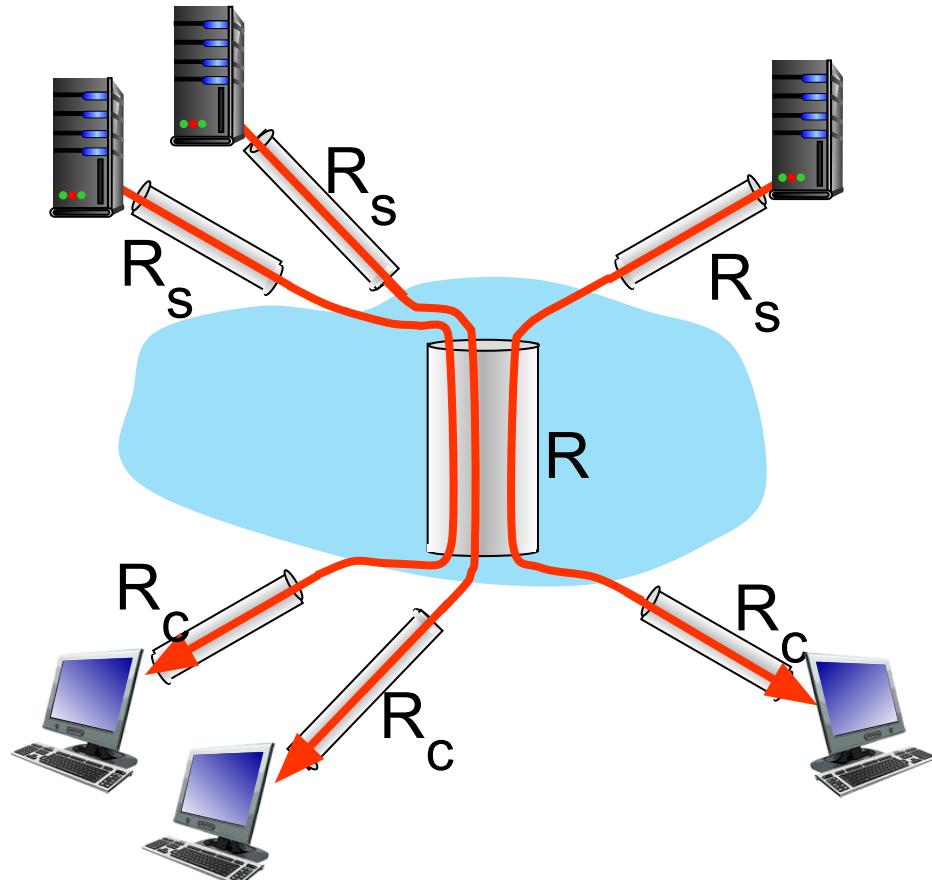
$R_s > R_c$ What is average end-end throughput?



— *bottleneck link* —

link on end-end path that constrains end-end throughput

Throughput: network scenario



10 connections (fairly) share
backbone bottleneck link R
bits/sec

- per-connection end-end throughput: $\min(R_c, R_s, R/10)$
- in practice: R_c or R_s is often bottleneck

* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/

Chapter 1: roadmap

- What *is* the Internet?
- What *is* a protocol?
- Network edge: hosts, access network, physical media
- Network core: packet/circuit switching, internet structure
- Performance: loss, delay, throughput
- **Security**
- Protocol layers, service models
- History



Network security

- Internet not originally designed with (much) security in mind
 - *original vision*: “a group of mutually trusting users attached to a transparent network” ☺
 - Internet protocol designers playing “catch-up”
 - security considerations in all layers!
- We now need to think about:
 - how bad guys can attack computer networks
 - how we can defend networks against attacks
 - how to design architectures that are immune to attacks

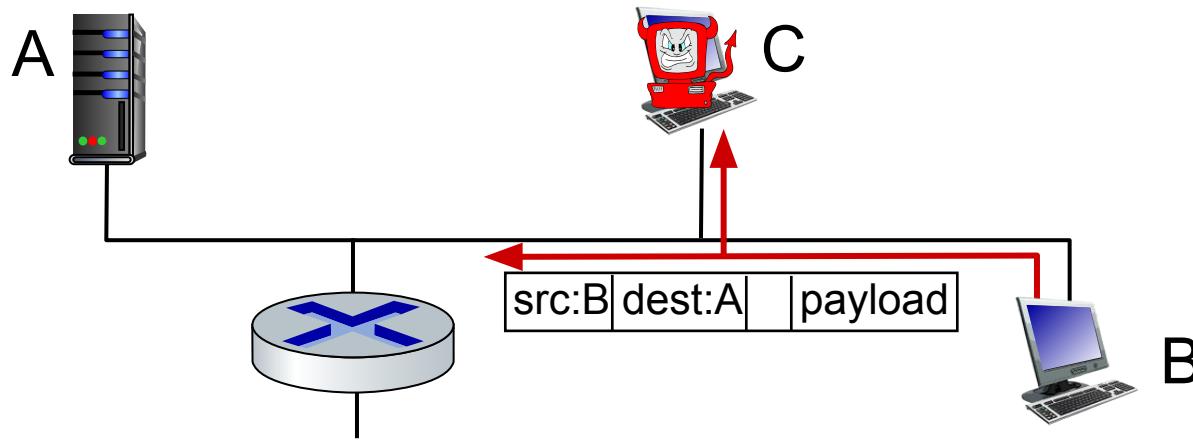
Network security

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Bad guys: packet interception

packet “sniffing”:

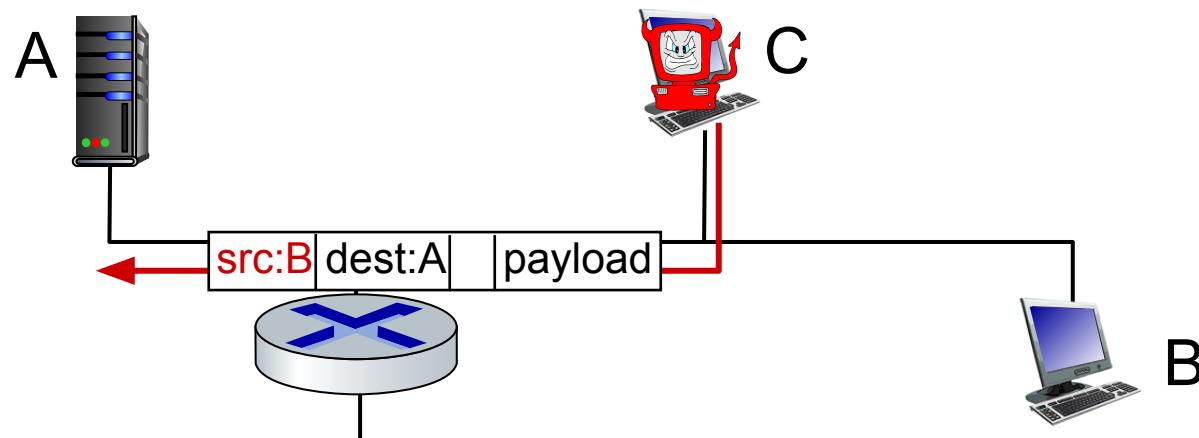
- broadcast media (shared Ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords!) passing by



Wireshark software used for our end-of-chapter labs is a (free) packet-sniffer

Bad guys: fake identity

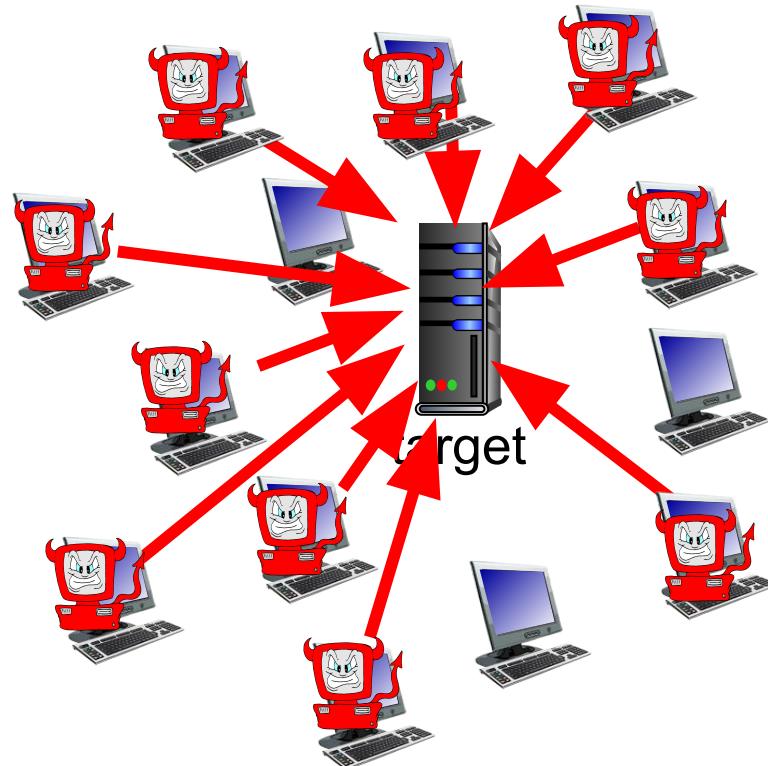
IP spoofing: injection of packet with false source address



Bad guys: denial of service

Denial of Service (DoS): attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

1. select target
2. break into hosts around the network (see botnet)
3. send packets to target from compromised hosts



Lines of defense:

- **authentication**: proving you are who you say you are
 - cellular networks provides hardware identity via SIM card; no such hardware assist in traditional Internet
 - **confidentiality**: via encryption
 - **integrity checks**: digital signatures prevent/detect tampering
 - **access restrictions**: password-protected VPNs
 - **firewalls**: specialized “middleboxes” in access and core networks:
 - off-by-default: filter incoming packets to restrict senders, receivers, applications
 - detecting/reacting to DOS attacks
- ... lots more on security (throughout, Chapter 8)*

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Protocol “layers” and reference models

Networks are complex, with many “pieces”:

- hosts
- routers
- links of various media
- applications
- protocols
- hardware, software

Question: is there any hope of *organizing* structure of network?

- and/or our *discussion* of networks?

Example: organization of air travel



end-to-end transfer of person plus
ticket (purchase)

baggage (check)

gates (load)

runway takeoff

airplane routing

ticket

(complain)

baggage

(claim)

gates (unload)

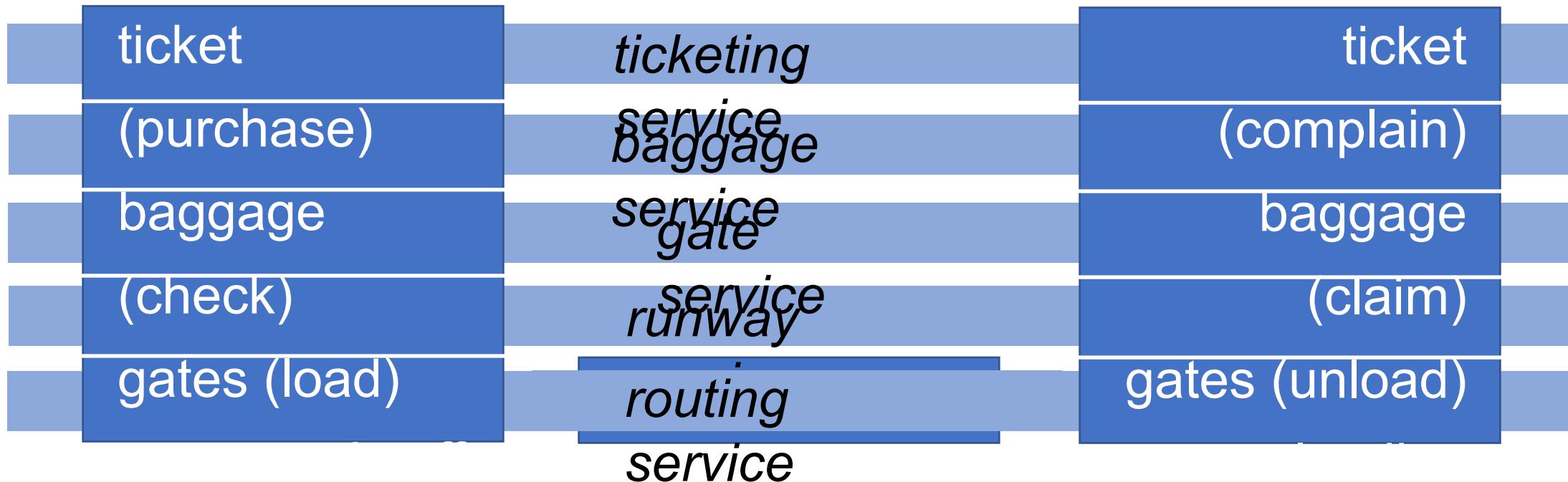
runway landing

airplane routing

airplane routing

How would you *define/discuss* the system of airline travel? a series of steps, involving many services

Example: organization of air travel



layers: each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

Why layering?

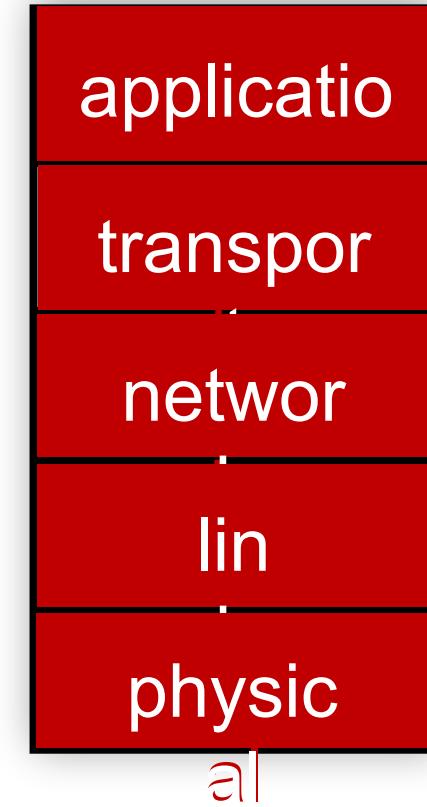
Approach to designing/discussing complex

systems:

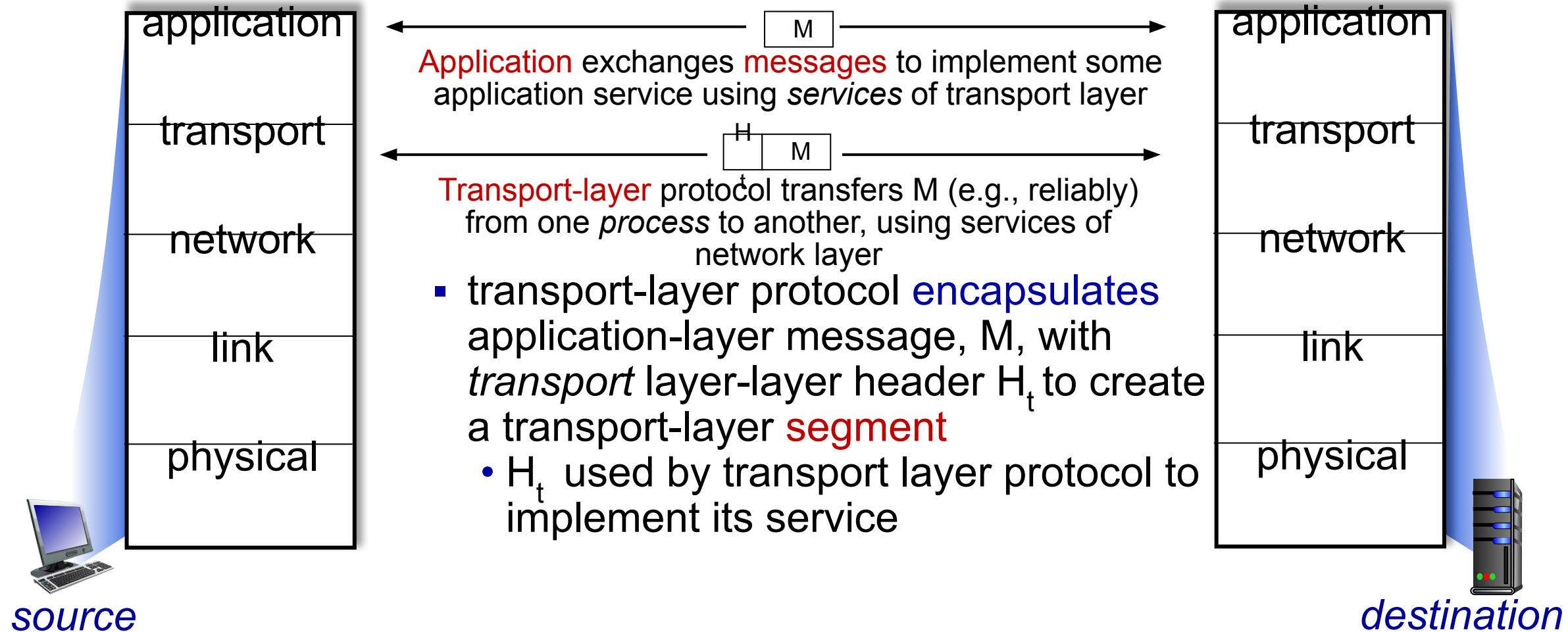
- explicit structure allows identification, relationship of system's pieces
 - layered *reference model* for discussion
- modularization eases maintenance, updating of system
 - change in layer's service *implementation*: transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system

Layered Internet protocol stack

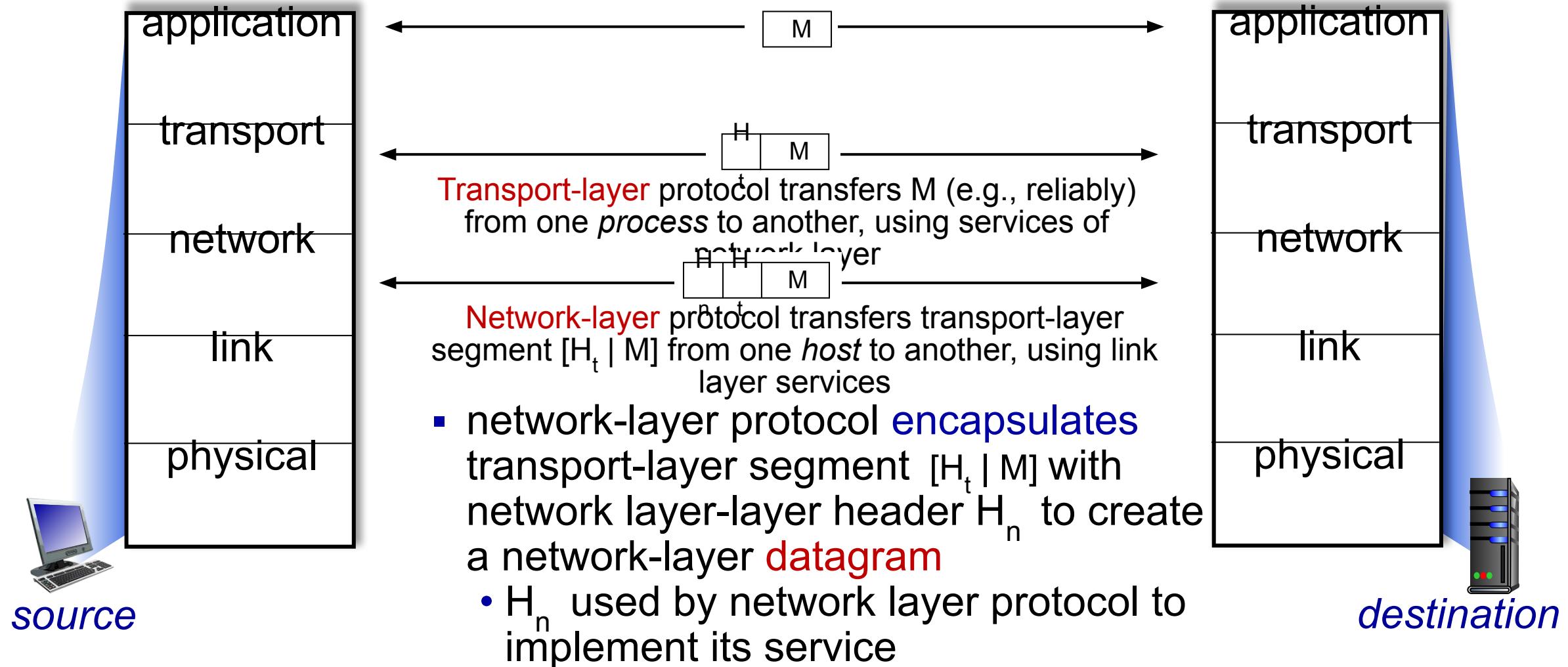
- ***application:*** supporting network applications
 - HTTP, IMAP, SMTP, DNS
- ***transport:*** process-process data transfer
 - TCP, UDP
- ***network:*** routing of datagrams from source to destination
 - IP, routing protocols
- ***link:*** data transfer between neighboring network elements
 - Ethernet, 802.11 (WiFi), PPP
- ***physical:*** bits “on the wire”



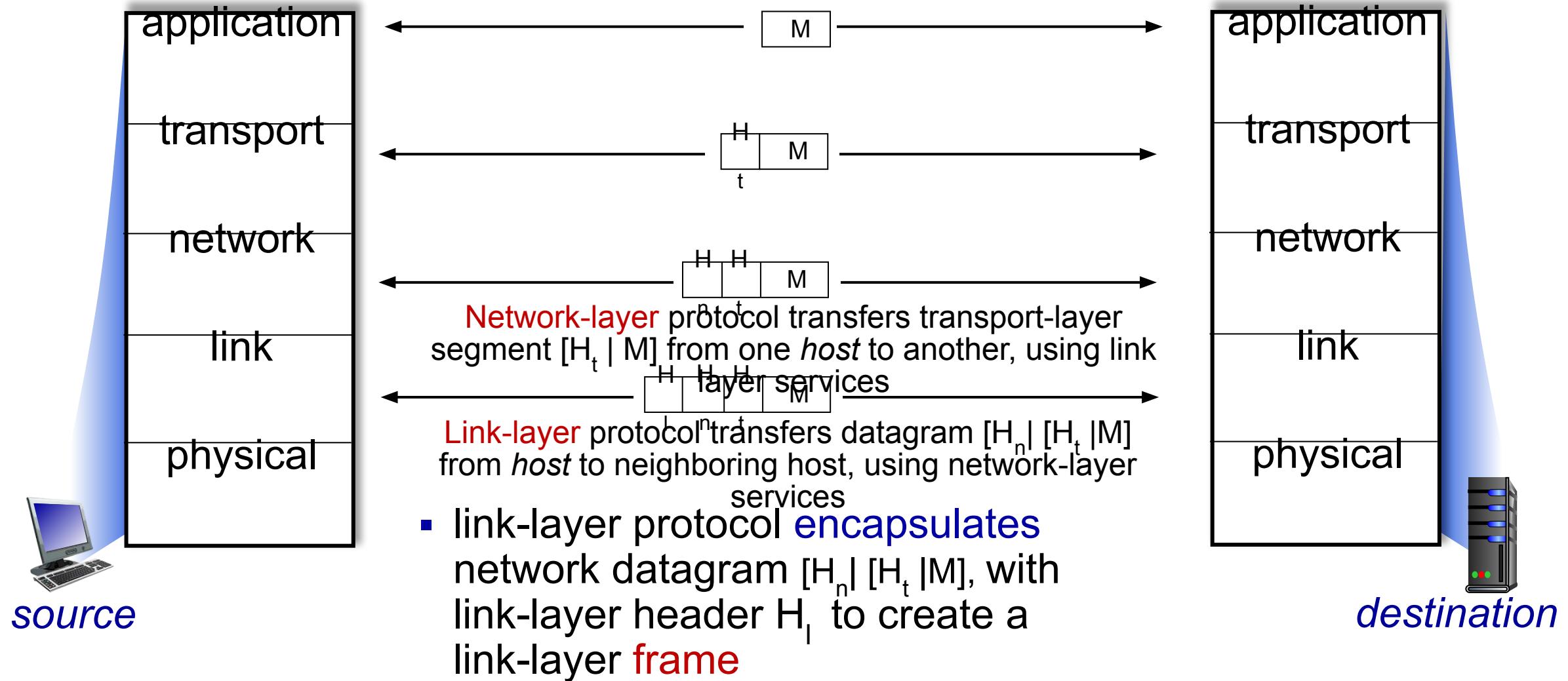
Services, Layering and Encapsulation



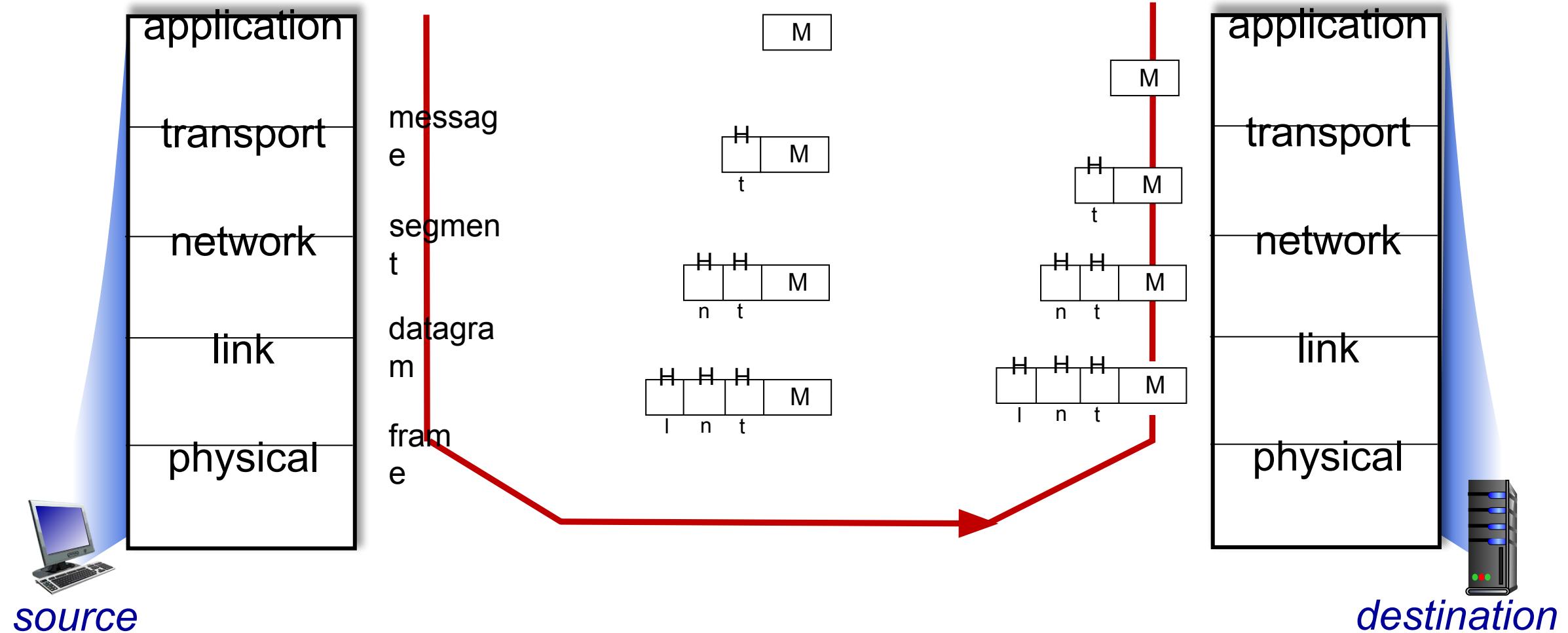
Services, Layering and Encapsulation



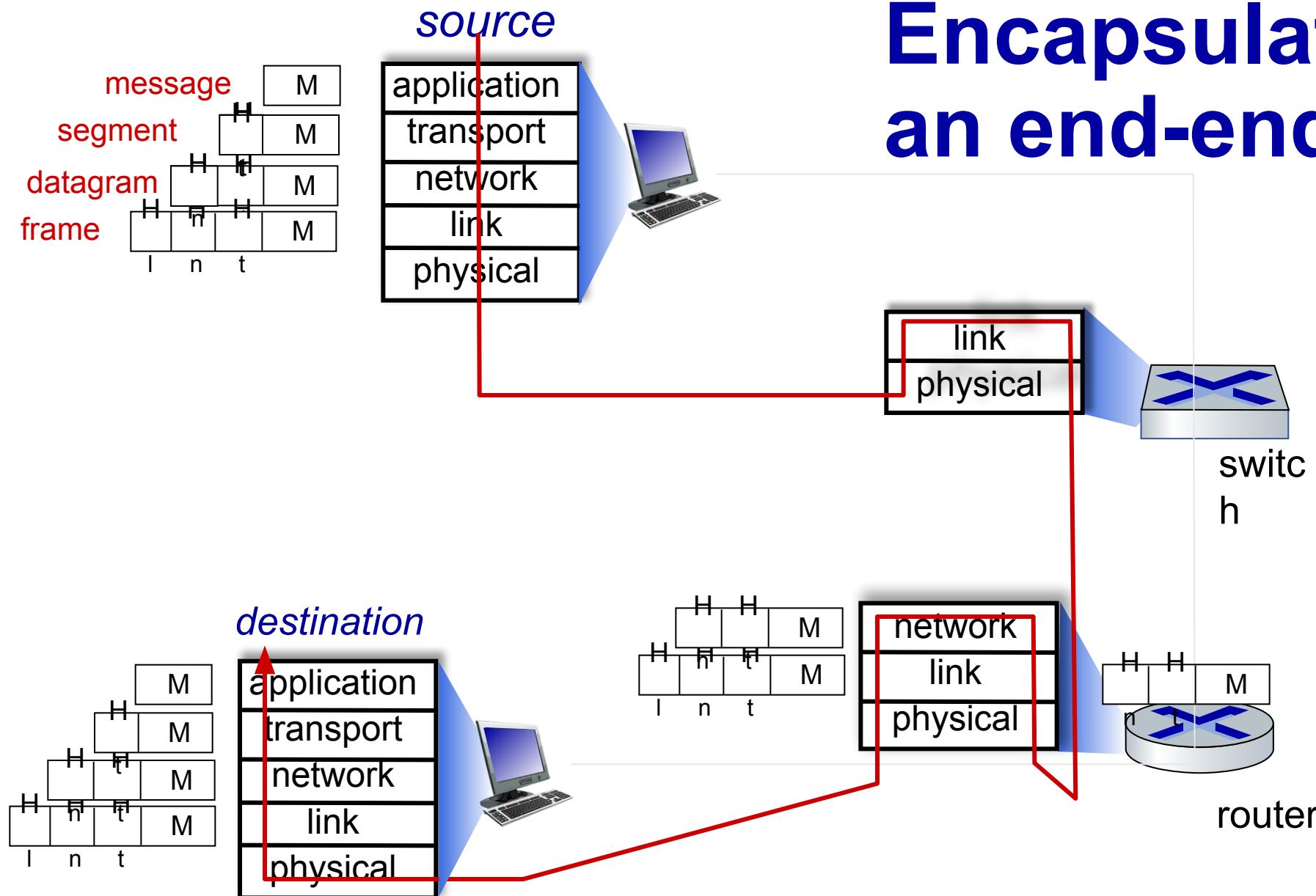
Services, Layering and Encapsulation



Services, Layering and Encapsulation



Encapsulation: an end-end view



Chapter 1: roadmap

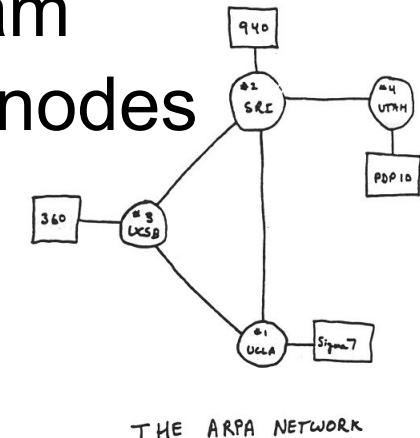
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- Protocol layers, service models
- History



Internet history

1961-1972: Early packet-switching principles

- 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- 1964: Baran - packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational
- 1972:
 - ARPAnet public demo
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program
 - ARPAnet has 15 nodes



THE ARPANET NETWORK

Internet history

1972-1980: Internetworking, new and proprietary networks

- 1970: ALOHAnet satellite network in Hawaii
- 1974: Cerf and Kahn - architecture for interconnecting networks
- 1976: Ethernet at Xerox PARC
- late70's: proprietary architectures: DECnet, SNA, XNA
- 1979: ARPAnet has 200

Cerf and Kahn's internetworking principles:

- minimalism, autonomy - no internal changes required to interconnect networks
- best-effort service model
- stateless routing
- decentralized control

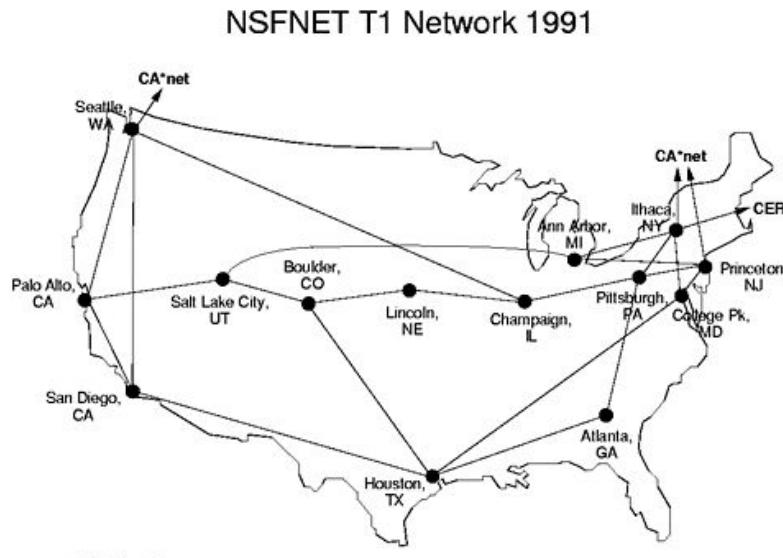
define today's Internet architecture

Internet history

1980-1990: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: smtp e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: ftp protocol defined
- 1988: TCP congestion control

- new national networks: CSnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks



Internet history

1990, 2000s: commercialization, the Web, new applications

- early 1990s: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990s: commercialization of the Web

late 1990s – 2000s:

- more killer apps: instant messaging, P2P file sharing
- network security to forefront
- est. 50 million host, 100 million+ users
- backbone links running at Gbps

Internet history

2005-present: scale, SDN, mobility,

cloud

- aggressive deployment of broadband home access (10-100's Mbps)
- 2008: software-defined networking (SDN)
- increasing ubiquity of high-speed wireless access: 4G/5G, WiFi
- service providers (Google, FB, Microsoft) create their own networks
 - bypass commercial Internet to connect “close” to end user, providing “instantaneous” access to social media, search, video content, ...
- enterprises run their services in “cloud” (e.g., Amazon Web Services, Microsoft Azure)
- rise of smartphones: more mobile than fixed devices on Internet (2017)

Chapter 1: summary

We've covered a "ton" of material!

- Internet overview
- what's a protocol?
- network edge, access network, core
 - packet-switching versus circuit-switching
 - Internet structure
- performance: loss, delay, throughput
- layering, service models
- security

You now have:

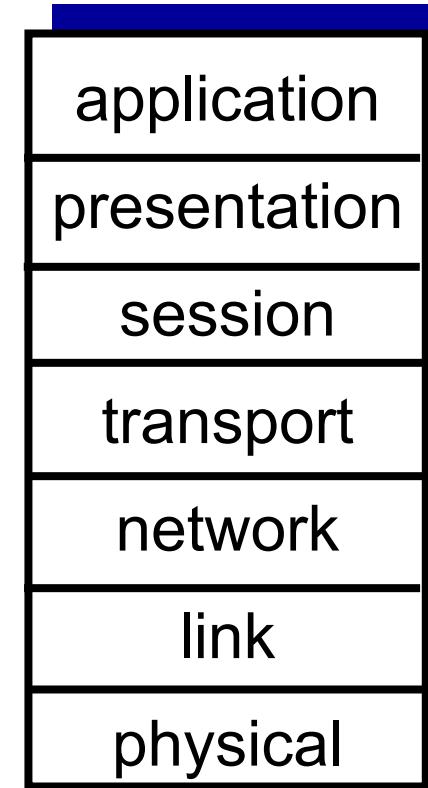
- context, overview, vocabulary, "feel" of networking
- more depth, detail, *and fun* to follow!

Additional Chapter 1 slides

ISO/OSI reference model

Two layers not found in Internet protocol stack!

- ***presentation:*** allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- ***session:*** synchronization, checkpointing, recovery of data exchange
- Internet stack “missing” these layers!
 - these services, *if needed*, must be implemented in application
 - needed?



The seven layer
OSI/ISO
reference model

Wireshark

