

Fundamentals of Programming Languages

Assignment 4

Universal Types

Mestrado em (Engenharia) Informática
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1 Binary Sums in System F

1. $\text{type Sum} = \forall T U V . (T \rightarrow V) \rightarrow (U \rightarrow V) \rightarrow V$
2. $\text{inl} : \forall T U . T \rightarrow (\forall V . (T \rightarrow V) \rightarrow (U \rightarrow V) \rightarrow V)$
 $\text{inl} = \lambda T . \lambda U . \lambda l : T . \lambda V . \lambda t : (T \rightarrow V) . \lambda u : (U \rightarrow V) . t \ l$
3. $\text{inr} : \forall T U . U \rightarrow (\forall V . (T \rightarrow V) \rightarrow (U \rightarrow V) \rightarrow V)$
 $\text{inr} = \lambda T . \lambda U . \lambda r : U . \lambda V . \lambda t : (T \rightarrow V) . \lambda u : (U \rightarrow V) . u \ r$
4. $\text{cases} : \forall T U V . (\forall V . (T \rightarrow V) \rightarrow (U \rightarrow V) \rightarrow V) \rightarrow (T \rightarrow V) \rightarrow (U \rightarrow V) \rightarrow V$
 $\text{cases} = \lambda T . \lambda U . \lambda V . \lambda v : (\forall V . (T \rightarrow V) \rightarrow (U \rightarrow V) \rightarrow V) . \lambda lc : (T \rightarrow V) . \lambda rc : (U \rightarrow V) . v \ [V] \ lc \ rc$
5. $\text{cases} [T] [U] [T] (\text{inl} [T] [U] v) (\lambda x_1 . t_1) (\lambda x_2 . t_2) =$
 $(\lambda V . \lambda t : (T \rightarrow V) . \lambda u : (U \rightarrow V) . t \ v) [T] (\lambda x_1 . t_1) (\lambda x_2 . t_2) =$
 $(\lambda t : (T \rightarrow T) . \lambda u : (U \rightarrow T) . t \ v) (\lambda x_1 . t_1) (\lambda x_2 . t_2) =$
 $(\lambda u : (U \rightarrow T) . (\lambda x_1 . t_1) \ v) (\lambda x_2 . t_2) =$
 $(\lambda x_1 . t_1) \ v$

This result is equivalent to $[x_1 \rightarrow v]t_1$, proving that it reduces to the latter result.

We can notice that this looks similar to a more generic implementation of what we have in fromL (2.2).

6. $\text{cases} [T] [U] [U] (\text{inr} [T] [U] v) (\lambda x_1 . t_1) (\lambda x_2 . t_2) =$
 $(\lambda V . \lambda t : (T \rightarrow V) . \lambda u : (U \rightarrow V) . u \ v) [U] (\lambda x_1 . t_1) (\lambda x_2 . t_2) =$
 $(\lambda t : (T \rightarrow U) . \lambda u : (U \rightarrow U) . u \ v) (\lambda x_1 . t_1) (\lambda x_2 . t_2) =$
 $(\lambda u : (T \rightarrow U) . u \ v) (\lambda x_2 . t_2) =$
 $(\lambda x_2 . t_2) \ v$

This result is equivalent to $[x_2 \rightarrow v]t_2$, proving that it reduces to the latter result.

We can notice that this looks similar to a more generic implementation of what we have in fromR (2.3).