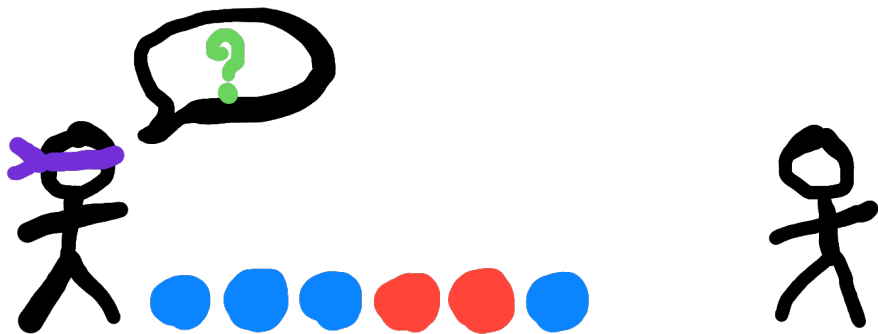


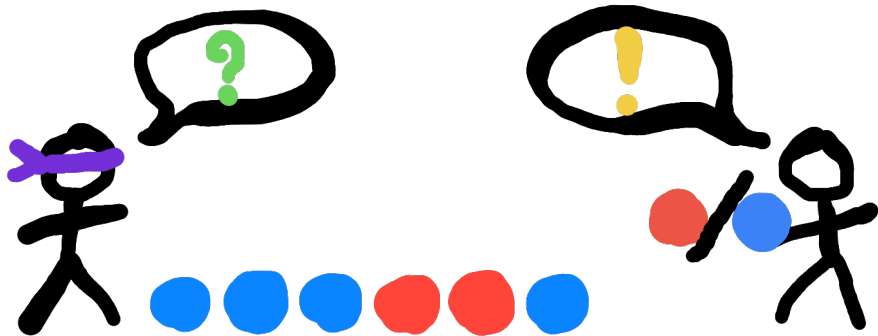
# Characteristic Boolean algebras in classical realizability: a completeness result

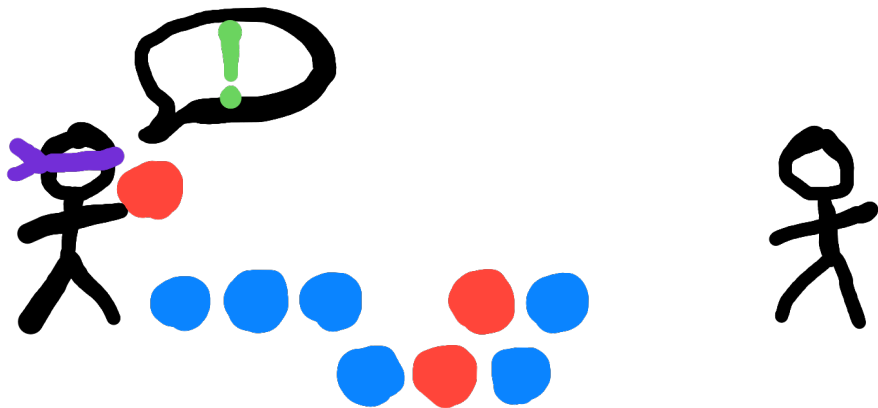
Guillaume Geoffroy

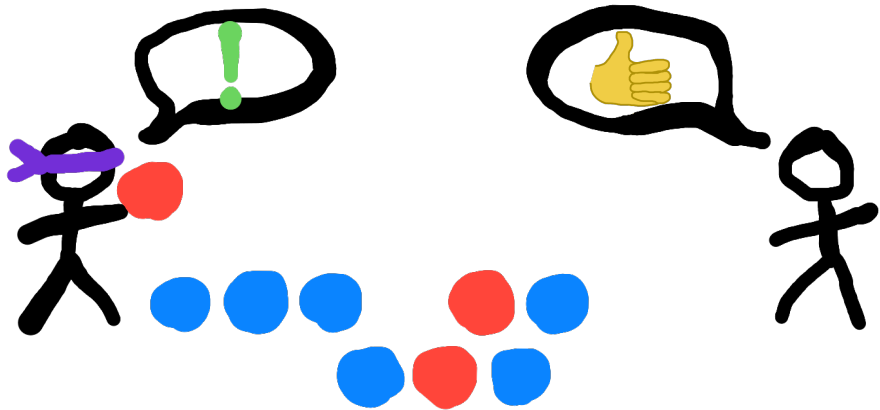
IRIF, Université Paris Cité

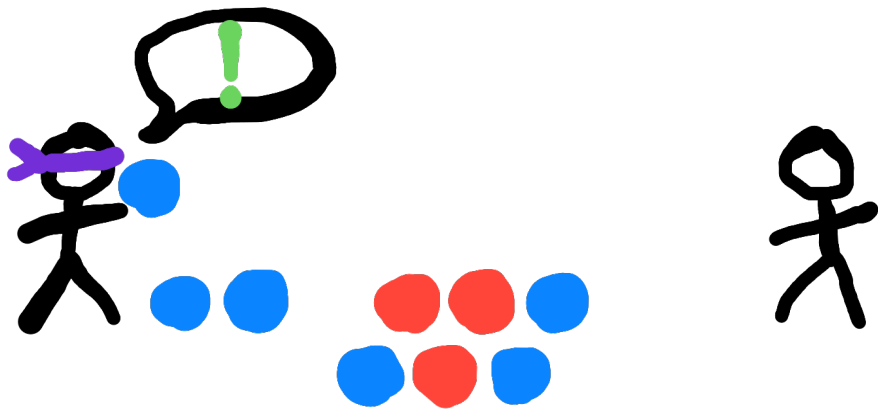


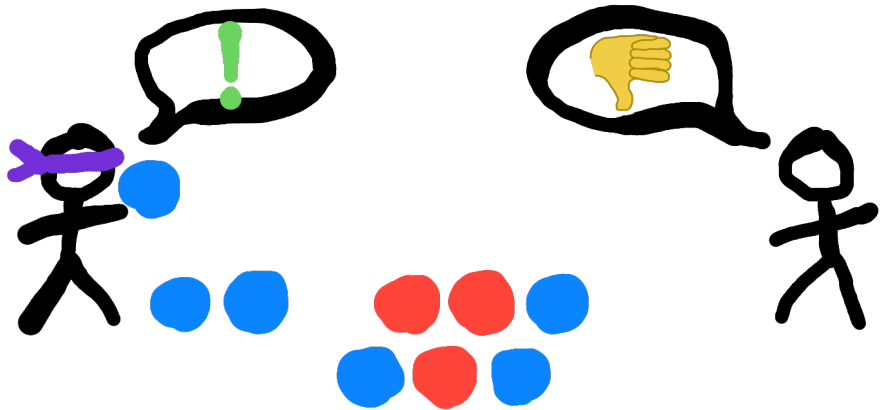










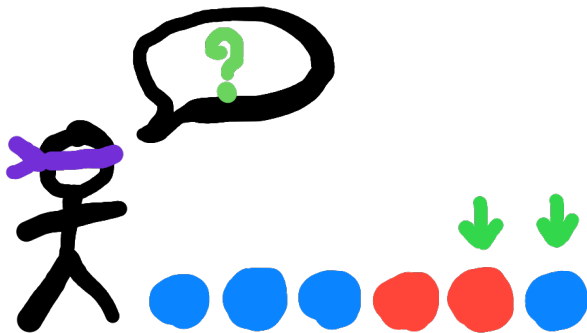




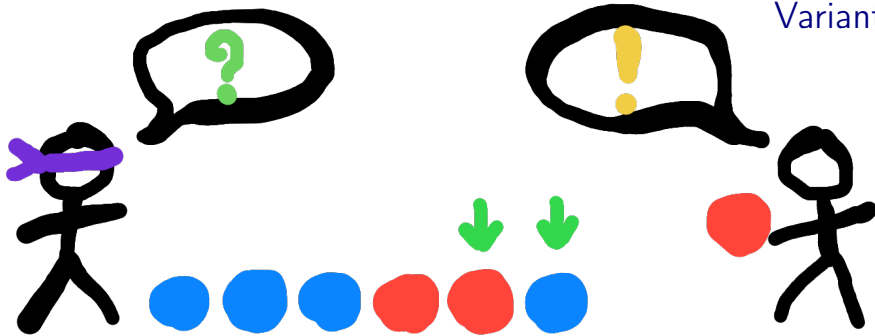
Variant  $\frac{1}{2}$



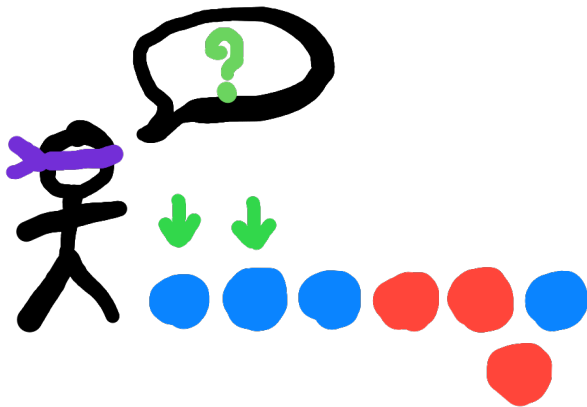
Variant  $\frac{1}{2}$



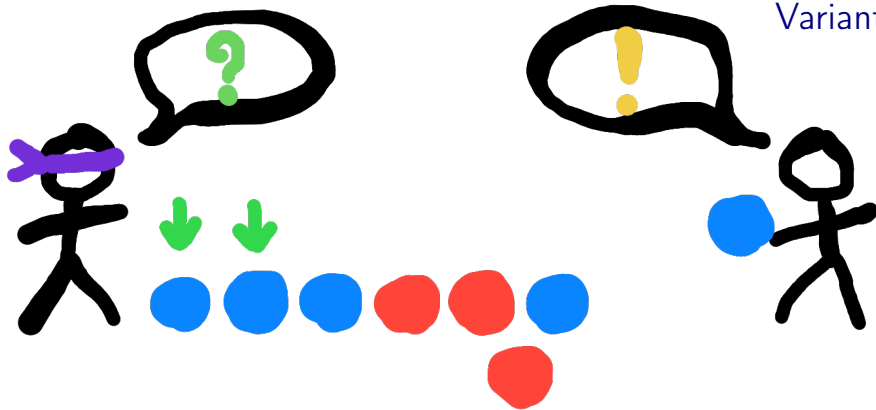
Variant  $1/2$



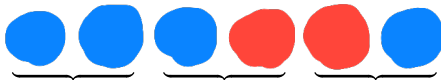
Variant <sup>1</sup>/<sub>2</sub>



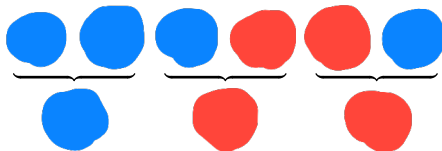
Variant <sup>1</sup>/<sub>2</sub>



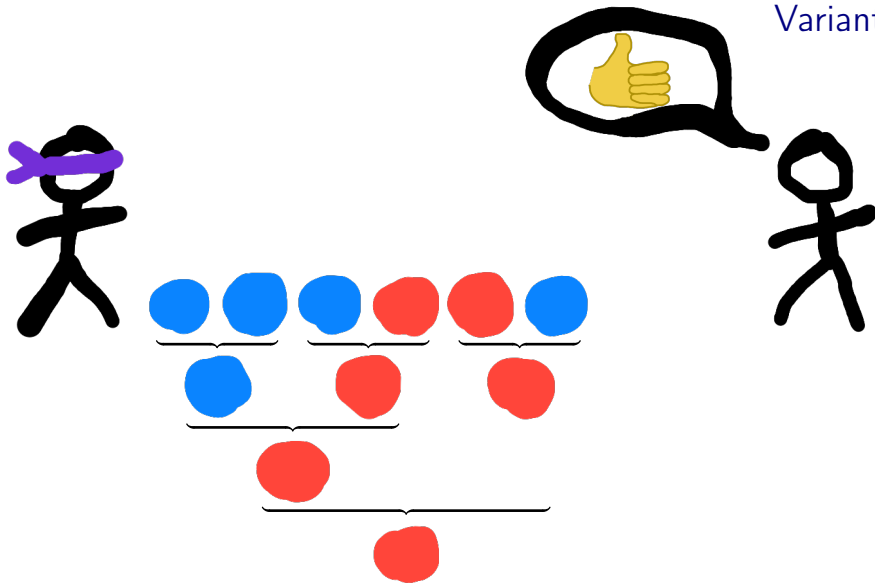
Variant <sup>1</sup>/<sub>2</sub>



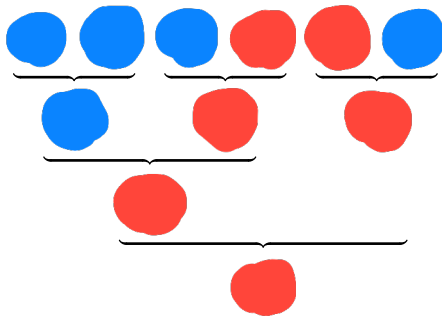
Variant <sup>1</sup>/<sub>2</sub>



Variant  $\frac{1}{2}$



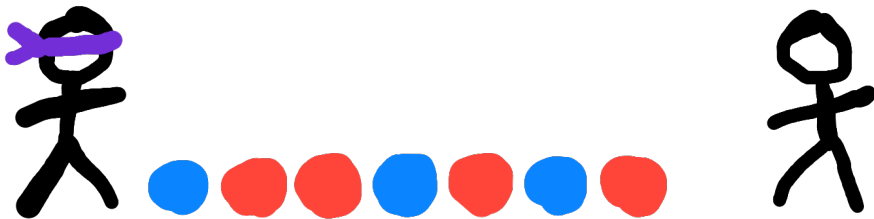




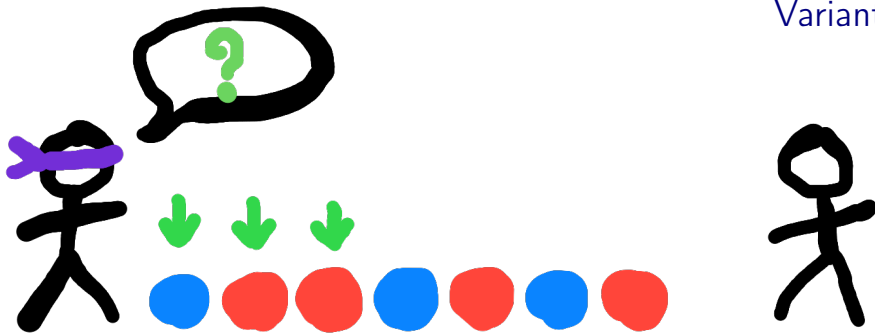
Variant  $\frac{1}{2}$   
Win if  $r > 0$



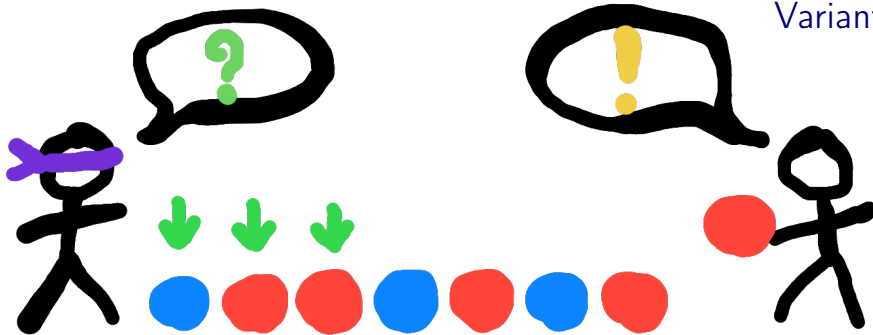
Variant <sup>2</sup>/<sub>3</sub>



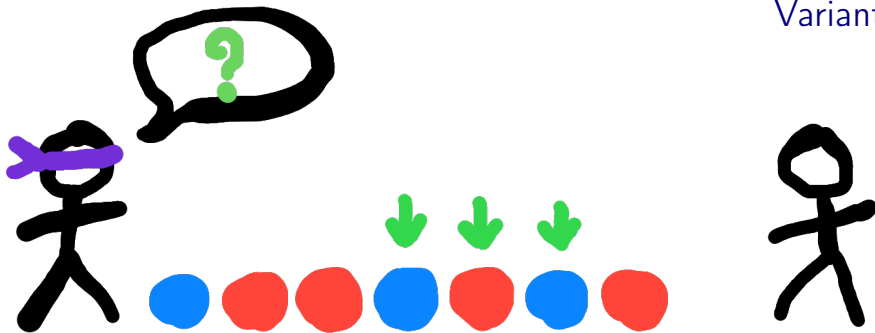
Variant <sup>2</sup>/<sub>3</sub>



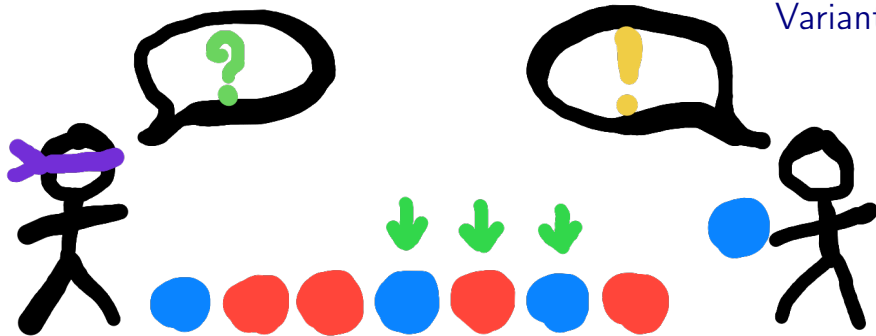
Variant <sup>2</sup>/<sub>3</sub>

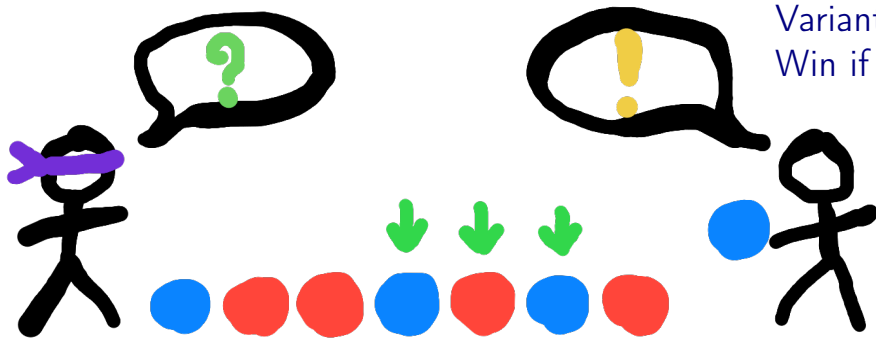


Variant <sup>2</sup>/<sub>3</sub>



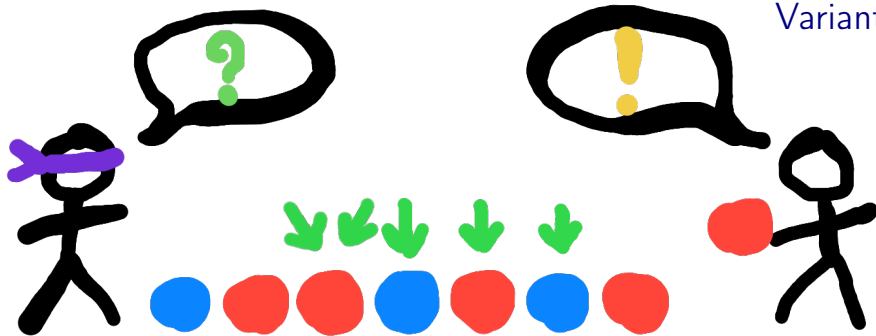
Variant <sup>2</sup>/<sub>3</sub>



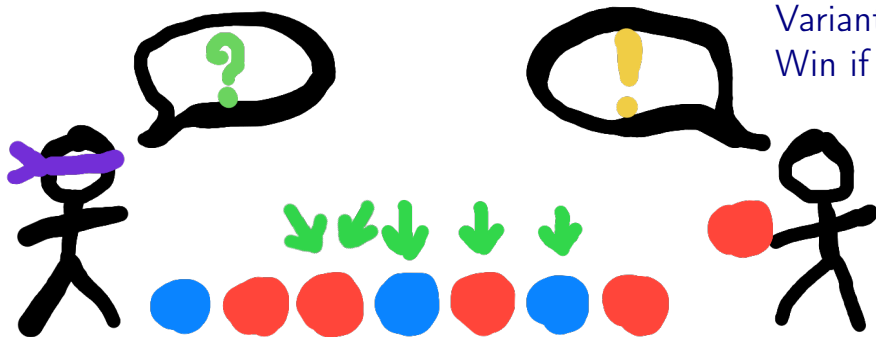


Variant  $\frac{2}{3}$   
Win if  $r > \frac{1}{2}$

Variant 3/5







Variant  $k/n$

Win if  $r > ???$



Puzzle for you!



# Classical realizability

Formulas:  $(A, B, \dots)$

- ▶ first-order
- ▶ language of Boolean algebras

Programs:  $t, u ::= x \mid \lambda x. t \mid tu$


Realizability relation: “ $t$  realizes  $A$ ”

# Classical realizability

Formulas:  $(A, B, \dots)$

- ▶ first-order
- ▶ language of Boolean algebras

Programs:  $t, u ::= x \mid \lambda x. t \mid tu$

Realizability relation: “ $t$  realizes  $A$ ”  
  
e.g.  $\lambda x. \lambda y. xyy$   
realizes  $(A \rightarrow A \rightarrow B) \rightarrow A \rightarrow B$   
for all  $A, B$

# A completeness result

Adequacy (standard): if  $A$  is true in all Boolean algebras,  
then there exists a  $\lambda$ -term  $t$  that realizes  $A$

# A completeness result

Adequacy (standard): if  $A$  is true in all Boolean algebras,  
then there exists a  $\lambda$ -term  $t$  that realizes  $A$

Theorem:  $A$  is true in all Boolean algebras  
iif there exists a  $\lambda$ -term  $t$  that realizes  $A$

# Back to the game

Allowed questions  
(variant  $k/n$ )



Formula  $A_{k,n}$

Initial situation  
( $N_{\text{red}}, N_{\text{balls}}$ )



Formula  $A_{N_{\text{red}}, N_{\text{balls}}}$

# Back to the game

Allowed questions  
(variant  $k/n$ )



Formula  $A_{k,n}$

Initial situation  
( $N_{\text{red}}, N_{\text{balls}}$ )



Formula  $A_{N_{\text{red}}, N_{\text{balls}}}$

Winning strategy



$\lambda$ -term that realizes  
 $A_{k,n} \rightarrow A_{N_{\text{red}}, N_{\text{balls}}}$



# Back to the game

Allowed questions  
(variant  $k/n$ )



Formula  $A_{k,n}$

Initial situation  
( $N_{\text{red}}, N_{\text{balls}}$ )



Formula  $A_{N_{\text{red}}, N_{\text{balls}}}$

$\exists$  Winning strategy



In all Boolean algebras,  
 $A_{k,n} \rightarrow A_{N_{\text{red}}, N_{\text{balls}}}$  holds

# Back to the game

Allowed questions  
(variant  $k/n$ )

Formula  $A_{k,n}$   
“ $\exists$  at most  $f(k/n)$  elements”

Initial situation  
( $N_{\text{red}}, N_{\text{balls}}$ )

Formula  $A_{N_{\text{red}}, N_{\text{balls}}}$   
“ $\exists$  at most  $f(N_{\text{red}}/N_{\text{balls}})$  elements”

$\exists$  Winning strategy

In all Boolean algebras,  
 $A_{k,n} \rightarrow A_{N_{\text{red}}, N_{\text{balls}}}$  holds

# Back to the game

Allowed questions  
(variant  $k/n$ )

Formula  $A_{k,n}$   
“ $\exists$  at most  $f(k/n)$  elements”

Initial situation  
( $N_{\text{red}}, N_{\text{balls}}$ )

Formula  $A_{N_{\text{red}}, N_{\text{balls}}}$   
“ $\exists$  at most  $f(N_{\text{red}}/N_{\text{balls}})$  elements”

$\exists$  Winning strategy

In all Boolean algebras,  
 $A_{k,n} \rightarrow A_{N_{\text{red}}, N_{\text{balls}}}$  holds  
 $f(k/n) \leq f(N_{\text{red}}/N_{\text{balls}})$

## Translating variant <sup>2</sup>/<sub>3</sub>

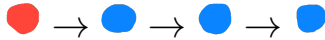
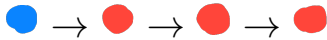
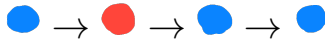
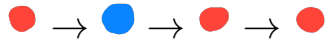
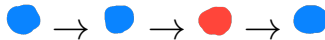
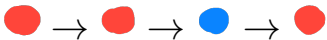
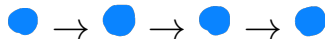
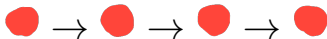
► The key:    ● ●  $\Leftrightarrow \perp$       ●  $\Leftrightarrow \top$

## Translating variant <sup>2</sup>/<sub>3</sub>

- The key:
-   $\Leftrightarrow \perp$
  -   $\Leftrightarrow \top$
  - universal quantification  $\Leftrightarrow$  intersection type

# Translating variant $2/3$

- ▶ The key:  $\bullet \text{ (red)} \Leftrightarrow \perp$        $\bullet \text{ (blue)} \Leftrightarrow \top$ 
  - universal quantification  $\Leftrightarrow$  intersection type
- ▶ Variant  $2/3$  is the intersection of:



# Translating variant $2/3$

- ▶ The key:  $\bullet \text{ (red)} \Leftrightarrow \perp$        $\bullet \text{ (blue)} \Leftrightarrow \top$ 
  - universal quantification  $\Leftrightarrow$  intersection type
- ▶ Variant  $2/3$  is the intersection of:

$\perp \rightarrow \perp \rightarrow \perp \rightarrow \perp$

$\perp \rightarrow \perp \rightarrow \top \rightarrow \perp$

$\perp \rightarrow \top \rightarrow \perp \rightarrow \perp$

$\top \rightarrow \perp \rightarrow \perp \rightarrow \perp$

$\top \rightarrow \top \rightarrow \top \rightarrow \top$

$\top \rightarrow \top \rightarrow \perp \rightarrow \top$

$\top \rightarrow \perp \rightarrow \top \rightarrow \top$

$\perp \rightarrow \top \rightarrow \top \rightarrow \top$

# Translating variant $2/3$

- ▶ The key:  $\bullet \text{ (red)} \Leftrightarrow \perp$        $\bullet \text{ (blue)} \Leftrightarrow \top$ 
  - universal quantification  $\Leftrightarrow$  intersection type
- ▶ Variant  $2/3$  is the intersection of:

$$\perp \rightarrow \perp \rightarrow \perp \rightarrow \perp$$

$$\perp \rightarrow \perp \rightarrow \top \rightarrow \perp$$

$$\perp \rightarrow \top \rightarrow \perp \rightarrow \perp$$

$$\top \rightarrow \perp \rightarrow \perp \rightarrow \perp$$

$$\top \rightarrow \top \rightarrow \top \rightarrow \top$$

$$\top \rightarrow \top \rightarrow \perp \rightarrow \top$$

$$\top \rightarrow \perp \rightarrow \top \rightarrow \top$$

$$\perp \rightarrow \top \rightarrow \top \rightarrow \top$$

- ▶  $A_{2,3}$  is the formula:

$$\forall x \forall y \forall z (x \neq 0 \rightarrow y \neq 0 \rightarrow z \neq 0 \rightarrow (x \wedge y \vee y \wedge z \vee z \wedge x) \neq 0)$$