

Chapter 14§ : Fusc.

In which we apply the techniques we learned in Chapter 13.

Suppose we are given the fusc function defined as follows.

$\text{fusc} : \text{natural} \rightarrow \text{natural}$

* (0) $\text{fusc}.0 = 0$

* (1) $\text{fusc}.1 = 1$

* (2) $\text{fusc}.(2*n) = \text{fusc}.n, 0 < n$

* (3) $\text{fusc}.(2*n+1) = \text{fusc}.n + \text{fusc}.(n+1), 0 < n$

We are given a natural number N and asked to write a program to compute $\text{fusc}.N$

Choosing invariants.

We propose as invariants

$P0 : \alpha * \text{fusc}.n + \beta * \text{fusc}.(n+1) = \text{fusc}.N$

$P1 : 0 \leq n \leq N$

Establishing the Invariants.

The following assignment establishes the invariants

$n, \alpha, \beta := N, 1, 0$

Termination.

We note the following.

$$\begin{aligned}
 & P0 \wedge P1 \wedge n=0 \\
 = & \quad \{\text{definitions of } P0, P1\} \\
 & \alpha * \text{fusc}.n + \beta * \text{fusc}.(n+1) = \text{fusc}.N \wedge 0 \leq n \leq N \wedge n=0 \\
 \Rightarrow & \quad \{\text{leibniz}\} \\
 & \alpha * \text{fusc}.0 + \beta * \text{fusc}.1 = \text{fusc}.N \wedge \text{true} \\
 = & \quad \{\text{predicate calculus}\} \\
 & \alpha * \text{fusc}.0 + \beta * \text{fusc}.1 = \text{fusc}.N
 \end{aligned}$$

$$\begin{aligned}
&= \quad \{\text{defn. fusc}\} \\
&\quad \alpha * 0 + \beta * 1 = \text{fusc.N} \\
&= \quad \{\text{arithmetic}\} \\
&\quad \beta = \text{fusc.N}
\end{aligned}$$

Guard.

$$n \neq 0$$

Variant.

$$n$$

Calculate the loop body.

$$\begin{aligned}
&P0 \\
&= \quad \{\text{definition P0}\} \\
&\quad \alpha * \text{fusc.n} + \beta * \text{fusc.(n+1)} = \text{fusc.N} \\
&= \quad \{\text{case analysis, even.n i.e. } n = 2*p\} \\
&\quad \alpha * \text{fusc.(2*p)} + \beta * \text{fusc.(2*p+1)} = \text{fusc.N} \\
&= \quad \{\text{definition fusc (2) (3)}\} \\
&\quad \alpha * \text{fusc.p} + \beta * (\text{fusc.p} + \text{fusc.(p+1)}) = \text{fusc.N} \\
&= \quad \{*/+\} \\
&\quad \alpha * \text{fusc.p} + \beta * \text{fusc.p} + \beta * \text{fusc.(p+1)} = \text{fusc.N} \\
&= \quad \{\text{gather like terms together}\} \\
&\quad (\alpha + \beta) * \text{fusc.p} + \beta * \text{fusc.(p+1)} = \text{fusc.N} \\
&= \quad \{\text{WP}\} \\
&\quad (n, \alpha, \beta := n \text{ div } 2, \alpha + \beta, \beta).P0
\end{aligned}$$

Giving us the program fragment

$$\text{if even.n} \rightarrow n, \alpha, \beta := n \text{ div } 2, \alpha + \beta, \beta$$

Let us analyse the other case

$$\begin{aligned}
&P0 \\
&= \quad \{\text{definition P0}\} \\
&\quad \alpha * \text{fusc.n} + \beta * \text{fusc.(n+1)} = \text{fusc.N} \\
&= \quad \{\text{case analysis, odd.n i.e. } n = 2*p+1\} \\
&\quad \alpha * \text{fusc.(2*p+1)} + \beta * \text{fusc.(2*p+2)} = \text{fusc.N} \\
&= \quad \{\text{definition fusc (2) (3)}\} \\
&\quad \alpha * (\text{fusc.p} + \text{fusc.(p+1)}) + \beta * \text{fusc.(p+1)} = \text{fusc.N} \\
&= \quad \{*/+\} \\
&\quad \alpha * \text{fusc.p} + \alpha * \text{fusc.(p+1)} + \beta * \text{fusc.(p+1)} = \text{fusc.N} \\
&= \quad \{\text{gather like terms together}\}
\end{aligned}$$

$$\begin{aligned}
& \alpha * \text{fusc.p} + (\alpha + \beta) * \text{fusc.(p+1)} = \text{fusc.N} \\
= & \quad \{WP\} \\
& (n, \alpha, \beta := (n-1) \text{ div } 2, \alpha, \alpha + \beta).P0
\end{aligned}$$

Which gives us the program fragment

if odd.n \rightarrow n, α , $\beta := (n-1) \text{ div } 2, \alpha, \alpha + \beta$

Finished program.

And so the entire program is as follows.

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n,  $\alpha$ ,  $\beta := N, 1, 0$  {P0  $\wedge$  P1}
;do n  $\neq 0 \rightarrow$  {P0  $\wedge$  P1  $\wedge$  n  $\neq 0$ }

    if even.n  $\rightarrow$  n,  $\alpha$ ,  $\beta := n \text{ div } 2, \alpha + \beta, \beta$ 
    [] odd.n  $\rightarrow$  n,  $\alpha$ ,  $\beta := (n-1) \text{ div } 2, \alpha, \alpha + \beta$ 
    fi

    {P0  $\wedge$  P1}
od
{ $\beta = \text{fusc.N}$ }

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We note that this solution has temporal complexity $O(\log(N))$