

# Databases and Info Systems

## Structured Query Language (SQL)

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March 10, 2020

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## majors

major_id	major_name
1	Internet of Things Engineering
2	Software Engineering
3	Finance
4	Electronic and Information Engineering

## modules

module_code	module_name	major_id
COMP2013J	Databases and Information Systems	2
COMP1001J	Introduction to Programming 1	1
EENG2002J	Circuits and Systems	4
EENG2003J	Digital Circuits	4

## students

student_id	student_name	major_id
06373313	Sean Russell	2
04123123	David Lillis	1
12453234	Abey Campbel	4

## grades

student_id	module_code	grade
06373313	COMP2013J	A+
06373313	COMP1001J	A+
04123123	EENG2002J	B
04123123	COMP1001J	A
12453234	COMP1001J	A
12453234	EENG2002J	C

# Inter-Relational Constraints

- Inter-Relational Constraints means that we are talking about constraints that apply between relations
- Previous constraints we have seen (UNIQUE, NOT NULL, PRIMARY KEY) have all been **intra-relational** constraints that operate just within one table
- What we are talking about here are constraints on **Foreign Keys**

# Foreign Keys

- A foreign key is an attribute (or group of attributes) in one table that is linked to the primary key of another table
  - They can also be linked to the same table
- Foreign keys act as a way to combine tables by representing the relationship between them
  - Foreign keys are usually used in join conditions

# Inter-Relational Constraints

- This type of constraint is used to enforce **referential integrity**
  - **Integrity**: The data stored in the database is accurate and consistent
  - **Referential**: The attribute references another attribute
- **Referential Integrity**: When an attribute refers to another attribute, it can only store values that also exist in that other attribute

# Referential Integrity Example

- **Referential Integrity:** When an attribute refers to another attribute, it can only store values that also exist in that other attribute
  - An attribute `module_code` in the **grades** table might refer to a `module_code` attribute in a **modules** table
  - Only codes that actually exist in the modules table can be stored in the results table



# Defining Inter-Relational Constraints

- SQL allows the definition of referential integrity constraints in two ways:
  - Using REFERENCES
  - Using FOREIGN KEY
- When we are defining these constraints we can associate "reaction policies" to violations
  - What should happen when the rule is broken

# REFERENCES

- The REFERENCES keyword allows us to associate a single attribute as a foreign key of another table
- This is added after the domain of the attribute in the create table statement, the syntax is like this:

```
module_code CHAR(10) REFERENCES modules(module_code),
```

- NOTE: In MySQL this will not be enforced so we will not use this one

# FOREIGN KEY

- The FOREIGN KEY syntax allows us to associate a single or multiple attributes as a foreign key of another table
- This is added in the "other constraints" section of the create table statement, the syntax is like this:

```
FOREIGN KEY ( Attribute {, Attribute } ) REFERENCES table(  
    Attribute {, Attribute } )
```

```
FOREIGN KEY (module_code) REFERENCES modules(module_code),
```

# Example

- Setting an inter-relation constraint based on a value in one table being used as a key in another table
- Here, each student has a `major_id`, which references the `major_id` attribute in the `majors` table

```
CREATE TABLE students (  
  student_id CHAR(8) PRIMARY KEY,  
  student_name VARCHAR(30),  
  major_id INT,  
  FOREIGN KEY(major_id) REFERENCES majors(major_id)  
);
```

- Only values for `major_id` that are stored in the `majors` table can be used

# Adding Without Major

## Available Values

```
select * from majors;
```

major_id	major_name
1	Internet of Things Engineering
2	Software Engineering
3	Finance
4	Electronic and Information Engineering

## Inserting A Student

```
INSERT INTO students VALUES ("06373313", "Sean Russell", 5);
```

```
ERROR 1452 (23000): Cannot add or update a child row: a foreign key constraint fails  
('week6'. 'students', CONSTRAINT 'students_ibfk_1' FOREIGN KEY ('major_id')  
REFERENCES 'majors' ('major_id'))
```

# Changing A Major

## Current Students

```
mysql> SELECT * FROM students;
```

student_id	student_name	major_id
06373313	Sean Russell	2

## Modifying Major

```
mysql> UPDATE majors SET major_id = 5 WHERE major_id = 2;
```

```
ERROR 1451 (23000): Cannot delete or update a parent row: a foreign key constraint  
fails ('week6'. 'students', CONSTRAINT 'students_ibfk_1' FOREIGN KEY ('major_id')  
REFERENCES 'majors' ('major_id'))
```

# Deleting a Major

## Current Students

```
mysql> SELECT * FROM students;
```

student_id	student_name	major_id
06373313	Sean Russell	2

## Deleting Major

```
mysql> DELETE FROM majors WHERE major_id = 2;
```

```
ERROR 1451 (23000): Cannot delete or update a parent row: a foreign key constraint  
fails ('week6'. 'students', CONSTRAINT 'students_ibfk_1' FOREIGN KEY ('major_id')  
REFERENCES 'majors' ('major_id'))
```

# Viewing Constraints

- To see what inter-relational constraints are on a table, we cannot use describe

```
DESCRIBE students;
```

Field	Type	Null	Key	Default	Extra
student_id	char(8)	NO	PRI	NULL	
student_name	varchar(30)	YES		NULL	
major_id	int	YES	MUL	NULL	

3 rows in set (0.00 sec)

- Instead, we can use SHOW CREATE TABLE
- This has a more complex output, but is more detailed
- IT gives you the SQL command that can create this table



# SHOW CREATE TABLE

```
SHOW CREATE TABLE students;
```

Table	Create Table
	<pre>CREATE TABLE 'students' (   'student_id' char(8) NOT NULL,   'student_name' varchar(30) DEFAULT NULL,   'major_id' int DEFAULT NULL,   PRIMARY KEY ('student_id'),   KEY 'major_id' ('major_id'),   CONSTRAINT 'students_ibfk_1' FOREIGN KEY ('major_id') REFERENCES 'majors' ('major_id') ) ENGINE=InnoDB DEFAULT CHARSET=utf8</pre>
1 row in set	(0.00 sec)

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# Reaction Policies

- In our example, making a change to `major_id` in the majors table may break our rules
  - A student with a major that does not exist
  - The FOREIGN KEY constraint prevents this from happening
  - This can be because of insertions, updates or deletions
- Depending on the database design, we may want to allow this type of change, but for the database to automatically react to make sure that integrity is maintained

# Reaction Policies

- There are different reactions that can happen when constraint is broken
  - CASCADE** The change that was made in the external table should also be made here.
  - SET NULL** The current value of the field is set to NULL so we are no longer connected to the other table
  - SET DEFAULT** Whatever the default value is for the attribute is assigned in place of the current value
  - NO ACTION** or **RESTRICT** This prevents the change from taking place on the external table
- **RESTRICT** is the default, when no reaction policy is set

# What to React To?

- The reaction we want may depend on what happened to break the rules
  - It might be an UPDATE operation where the key is changed
  - It might be a DELETE operation where the row is removed
- The FOREIGN KEY constraint allows us to specify what we want to happen in each of these circumstances

# Updated

- Choosing What should happen when the key is updated
  - ON UPDATE CASCADE
  - ON UPDATE SET NULL
  - ON UPDATE SET DEFAULT
  - ON UPDATE NO ACTION
  - ON UPDATE RESTRICT

# Deleted

- Choosing What should happen when the row is deleted
  - ON DELETE CASCADE
  - ON DELETE SET NULL
  - ON DELETE SET DEFAULT
  - ON DELETE NO ACTION
  - ON DELETE RESTRICT

# Cascade Example

```
CREATE TABLE modules (  
  module_code CHAR(9) PRIMARY KEY,  
  module_name VARCHAR(40),  
  major_id INT,  
  FOREIGN KEY(major_id) REFERENCES majors(major_id)  
  ON DELETE CASCADE  
  ON UPDATE CASCADE  
);
```

- When the major is deleted, then the rows matching this `major_id` are removed from the modules table
- When the value of `major_id` is changed in the majors table, then the value of `major_id` is changed for every module with the matching value



# Cascade Example

# Set NULL Example

```
CREATE TABLE students (  
  student_id CHAR(8) PRIMARY KEY,  
  student_name VARCHAR(30),  
  major_id INT,  
  FOREIGN KEY(major_id) REFERENCES majors(major_id)  
  ON DELETE SET NULL  
  ON UPDATE SET NULL  
);
```

- When the major is deleted, then the `major_id` of each student with that value is set to `NULL`
- When the value of `major_id` is changed in the majors table, then the value of `major_id` is set to `NULL` for every student with the matching value

# SET NULL EXAMPLE

# More Complex Example

```
CREATE TABLE grades(  
  student_id CHAR(8),  
  module_code CHAR(9),  
  grade VARCHAR(2),  
  PRIMARY KEY(student_id,module_code),  
  FOREIGN KEY(student_id) REFERENCES students(student_id)  
    ON DELETE RESTRICT ON UPDATE CASCADE,  
  FOREIGN KEY(module_code) REFERENCES modules(module_code)  
    ON UPDATE CASCADE ON DELETE RESTRICT  
);
```

- What are the effects here?

# Effects

- 1 If we try to delete a student from the students table, but they have grades in the grades table it will be blocked (ON DELETE RESTRICT)
- 2 If we change a students id in the students table, the matching student id will be changed in the grades table (ON UPDATE CASCADE)
- 3 If we try to delete a module from the modules table, but there are grades for that module in the grades table it will be blocked (ON DELETE RESTRICT)
- 4 If we change a modules' code in the modules table, then the same module code in the grades table will be change to match (ON UPDATE CASCADE)

# Testing 1

- A student with grades

```
DELETE FROM students WHERE student_id = "06373313";
```

```
ERROR 1451 (23000): Cannot delete or update a parent row: a foreign key  
constraint fails ('week6'. 'grades', CONSTRAINT 'grades_ibfk_1' FOREIGN KEY  
( 'student_id' ) REFERENCES 'students' ( 'student_id' ) ON DELETE RESTRICT ON  
UPDATE CASCADE)
```

- A student with no grades

```
DELETE FROM students WHERE student_id = "14232232";
```

```
Query OK, 1 row affected (0.01 sec)
```

# Testing 2

```
SELECT * from grades where module_code = "COMP1001J";
```

student_id	module_code	grade
04123123	COMP1001J	A
06373313	COMP1001J	A+
12453234	COMP1001J	A

3 rows in set (0.00 sec)

```
UPDATE students SET student_id = "16373313" WHERE student_id = "06373313";
```

```
SELECT * from grades where module_code = "COMP1001J";
```

student_id	module_code	grade
04123123	COMP1001J	A
12453234	COMP1001J	A
16373313	COMP1001J	A+

3 rows in set (0.00 sec)

# Testing 3

- A module with grades

```
DELETE FROM modules where module_code = "COMP2013J";
```

```
ERROR 1451 (23000): Cannot delete or update a parent row: a foreign key  
constraint fails ('week6'. 'grades', CONSTRAINT 'grades_ibfk_2' FOREIGN KEY  
( 'module_code' ) REFERENCES 'modules' ( 'module_code' ) ON DELETE RESTRICT ON  
UPDATE CASCADE)
```

- A module with no grades

```
DELETE FROM modules WHERE module_code = "EENG2003J";
```

```
Query OK, 1 row affected (0.01 sec)
```



# Testing 4

```
SELECT * from grades WHERE grade = "A";
```

student_id	module_code	grade
04123123	COMP1001J	A
12453234	COMP1001J	A
2 rows in set (0.00 sec)		

```
UPDATE modules SET module_code = "COMP1004J" WHERE module_code =  
"COMP1001J";
```

```
SELECT * from grades WHERE grade = "A";
```

student_id	module_code	grade
04123123	COMP1004J	A
12453234	COMP1004J	A
2 rows in set (0.00 sec)		

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## employees

emp_id	name	title	salary	dept_id	join_date
1234	Sean Russell	Trainer	50000	10	2018-03-01
4567	Jamie Heaslip	Manager	47000	10	2004-10-21
6542	Leo Cullen	Trainer	45000	10	2012-12-01
1238	Brendan Macken	Technician	25000	20	2001-09-10
1555	Sean O'Brien	Designer	50000	20	1999-06-24
1899	Brian O'Driscoll	Manager	45000	20	1998-02-27
2525	Peter Stringer	Designer	25000	30	2017-01-16
1585	Denis Hickey	Architect	20000	30	2009-08-07
1345	Ronan O'Gara	Manager	29000	30	2019-12-25

## departments

dept_id	dept_name	office	division	manager_id
10	Training	Lansdowne	D1	4567
20	Design	Belfield	D2	1899
30	Implementation	Donnybrook	D1	1345
40	Strategy	Terenure	D2	NULL

# Nested Queries

- A **nested query** is an SQL query that is contained within another query
  - This is sometimes called a **subquery**
- Usually used in the WHERE clause
- Subqueries can return:
  - A single value (called a **scalar**)
  - A single column
  - A single row
  - A table (multiple columns/rows)

- When a subquery return a single value (scalar), it can be used in the same way as an ordinary single value
- Find the name of the employee(s) who earn the most money

```
SELECT name, salary FROM employees WHERE salary = (SELECT MAX(salary)
FROM employees);
```

name	salary
Sean Russell	50000
Sean O'Brien	50000

2 rows in set (0.01 sec)

- A nested query returning a single column can be thought of as a set of values.
- We can compare attributes with values from this set to see if its equal, greater than, less than, etc.
  - Using = > < >= <= <> !=
- Two other keywords are important:
  - ANY** returns true if the comparison is true for any value in the set
  - ALL** returns true if the comparison is true for all the values in the set

# Example 1

- Find the names of employees who work in departments in Division 'D1'

```
SELECT name FROM employees WHERE dept_id = ANY(SELECT dept_id FROM  
departments WHERE division='D1');
```

- The subquery finds a list of all the department numbers (`dept_id`) for departments that are in division D1
- When selecting from the employees table, it will match any employee whose `dept_id` is in the set returned by the subquery
- We could also have solved this with a join

## Example 2

- Find the employees of department number 10 who have the same first name as a member of department 20
- We can use the SUBSTRING\_INDEX function from <https://dev.mysql.com/doc/refman/8.0/en/string-functions.html>
- SUBSTRING\_INDEX(name, ' ', 1) returns a substring of name, starting at the beginning and ending just before the first space character

```
SELECT name FROM employees WHERE dept_id=10 AND  
SUBSTRING_INDEX(name, ' ', 1)=ANY(SELECT SUBSTRING_INDEX(name,  
' ', 1) FROM employees WHERE dept_id=20 );
```



# Example 3

- Find the name of the Department in which there is no employee named Sean

```
SELECT dept_name FROM departments WHERE dept_id != ALL(SELECT  
DISTINCT dept_id FROM employees WHERE name LIKE 'Sean %');
```

- The subquery finds the dept\_id for departments who do have an employee named Sean.
- We want to find departments whose numbers are different to all of these
- Note that we change from **ANY** to **ALL** because this is a negative match

# Example 4

- Find the name of all employees who earn more money than everybody in the Implementation department

```
SELECT name, salary FROM employees WHERE salary > ALL(SELECT salary
FROM employees JOIN departments USING(dept_id) WHERE
dept_name='Implementation');
```

- Remember the Closure Property: the output of an operation is a relation
- The output of a SELECT query is a relation, so another query can be run on it
- In SQL, this is OK if you use an alias for the (temporary) table
- A silly example:

```
SELECT name FROM (SELECT * FROM employees WHERE job='Manager') AS managers;
```

- Here, the managers table doesn't actually exist, but can be used in a more complex SELECT query just like any real table

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- Creating Views
- Changing Views
- Style

- SQL provides the ability to use **views** in your schema
- A view is a **virtual table** based on a query
- It looks just like a normal table, but it is not stored directly in the database
- We can allow users to see a subset of one or more tables, without giving them access to the table itself

- Every view has a name and a SELECT query that defines it
- Example:

```
CREATE VIEW managers AS SELECT * FROM employees WHERE emp_id =  
    ANY(SELECT manager_id FROM departments);
```

- After this, managers looks just like an ordinary table.
- Changes in the employees table will automatically be seen in the manager table
- To delete a view: `DROP VIEW managers;`

# Changing Data

- For some **simple** views, you can update and delete just like a normal table, and these changes will be made to the underlying tables
- Full details:  
(<https://dev.mysql.com/doc/refman/8.0/en/view-updatability.html>)
- Some examples of When you can't
  - If it contains aggregate functions
  - If it uses DISTINCT
  - If it uses GROUP BY
  - Subquery in SELECT list (sometimes)
  - Certain joins
  - If it refers to a non-updatable view in FROM

# Style Guide

- As a database designer and programmer, SQL gives you a lot of freedom about how you choose names for tables/databases/variables and how you choose to capitalise.
- It's a good idea to follow a consistent style, which makes your queries more understandable and is more professional overall.
- A good style guide is by Simon Holywell at <https://www.sqlstyle.guide> (linked on Moodle).