

开拓 创新 诚信 求实

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Data Structures and Algorithms I

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The List Abstract Data Type - Singly-Linked List

Learning Outcomes

- After this lecture and the related practical, students should...
 - understand the operations of the list and position abstract data types
 - be able to implement a singly-linked list
 - know the complexity of all of the operations of the singly-linked list implementation

Table of Contents

- The List Abstract Data Type
- Singly-Linked List Implementation
 - Algorithmic Complexity

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- List Abstract Data Type
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List Abstract Data Type

■ Concept

- The list abstract data type models a sequence of positions
- Each position store one piece of data
- There is a before/after relationship between the positions
 - ◆ This allows for efficient insertion into and deletion from anywhere in the list

Position Abstract Data Type

- Before we talk about the list abstract data type, let's first look at the position abstract data type
- Concept
 - Position abstract data type models the idea of one place in a data structure into which a single data item is stored
 - Position provides a simple view of many different ways of storing a single data item
 - ◆ An element in an array
 - ◆ A node in a linked list

Position Abstract Data Type

■ Specification

● Operation

- ◆ `element()`: this returns the piece of data that is stored in a position

■ Interface

```
public interface Position {  
    public int element();  
}
```


List Abstract Data Type: Specification

■ Operations

- first(): return the first position in the list
- last(): return the last position in the list
- before(p): return the position in the list before p
- after(p): return the position in the list after p
- insertBefore(p, d): insert the value d into the position in the list before p
- insertAfter(p, d): insert the value d into the position in the list after p
- insertFirst(d): insert the value d into the first position in the list
- insertLast(d): insert the value d into the last position in the list
- remove(p): remove the position p from the list
- size(): return the number of elements stored in the list
- isEmpty(): return true if the list is empty and false otherwise

List Abstract Data Type

■ Interface

```
public interface List {  
    public Position first();  
    public Position last();  
    public Position before(Position p);  
    public Position after(Position p);  
    public Position insertBefore(Position p, int d);  
    public Position insertAfter(Position p, int d);  
    public Position insertFirst(int d);  
    public Position insertLast(int d);  
    public int remove(Position p);  
    public int size();  
    public boolean isEmpty();  
}
```

List Abstract Data Type

■ Implementation strategies

- Array based implementation
- Link based implementation

◆ Singly-linked list

- Each Position object keeps a reference to the next Position in the sequence

◆ Doubly-linked list

- Each Position object keeps a reference to the next and the previous Positions in the sequence

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Singly-Linked List Implementation

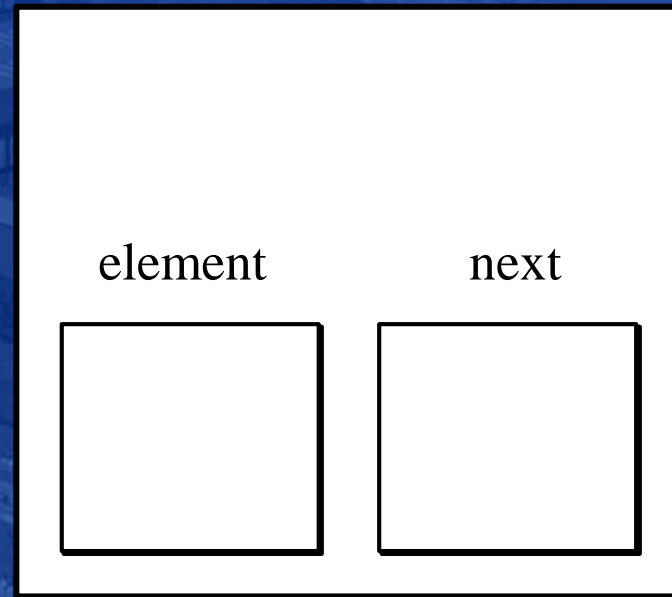
- We create a Node class that implements the Position interface
- We add functionality to the class to store a Node in the sequence

```
public class Node implements Position {  
    private int element;  
    Node next;  
  
    public Node(int e) {  
        this.element = e;  
    }  
  
    public int element() {  
        return element;  
    }  
}
```

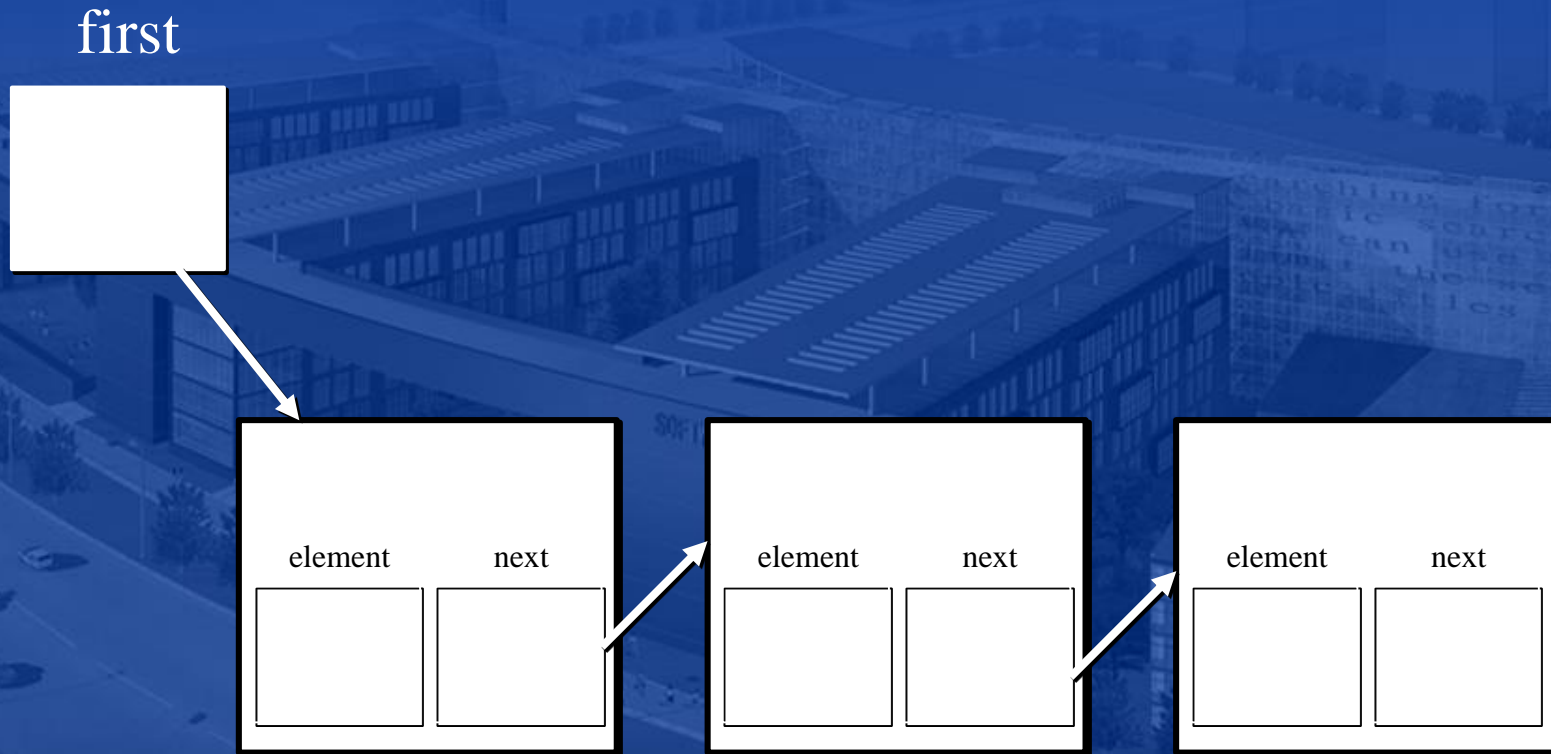
Singly-Linked List Implementation

- We keep a reference to the first position in the list
- We update the references when necessary
- We keep count of the number of positions in the list
- Variables
 - A reference to the first position in the list
private Node first;
 - A number to keep track of the number of nodes in the list
private int size;

Representation of a Node Object



Representation of a List



Singly-Linked List Operations

■ first()

- Return the reference that is stored in the variable first

■ size()

- Return the value of the variable size

■ isEmpty()

- Return the result of the expression $size == 0$

Singly-Linked List Operations

- after(p)
 - Return the next reference of p
- last()
 - Return the last position in the list
 - We only have a reference to the first position in the list
 - Each position only knows how to get to the position after it
 - We must follow the list from position to position until we reach the last position of the list
 - The last position is reached when there is no reference stored in the variable next of a position

Singly-Linked List Operations

■ Pseudocode of last()

Algorithm last()

Output: the last position in the list

if isEmpty() then
 return null;

i ← first;

while (i.next != null) do

 i ← i.next;

return i;

Singly-Linked List Operations

■ Pseudocode of before(p)

Algorithm before(p)

Input: the reference position

Output: the position before p

if isEmpty() OR $p = \text{first}$ then
 return null;

$i \leftarrow \text{first};$

while ($i.\text{next} \neq \text{null} \ \&\& \ p \neq i.\text{next}$) do

$i \leftarrow i.\text{next};$

return i;

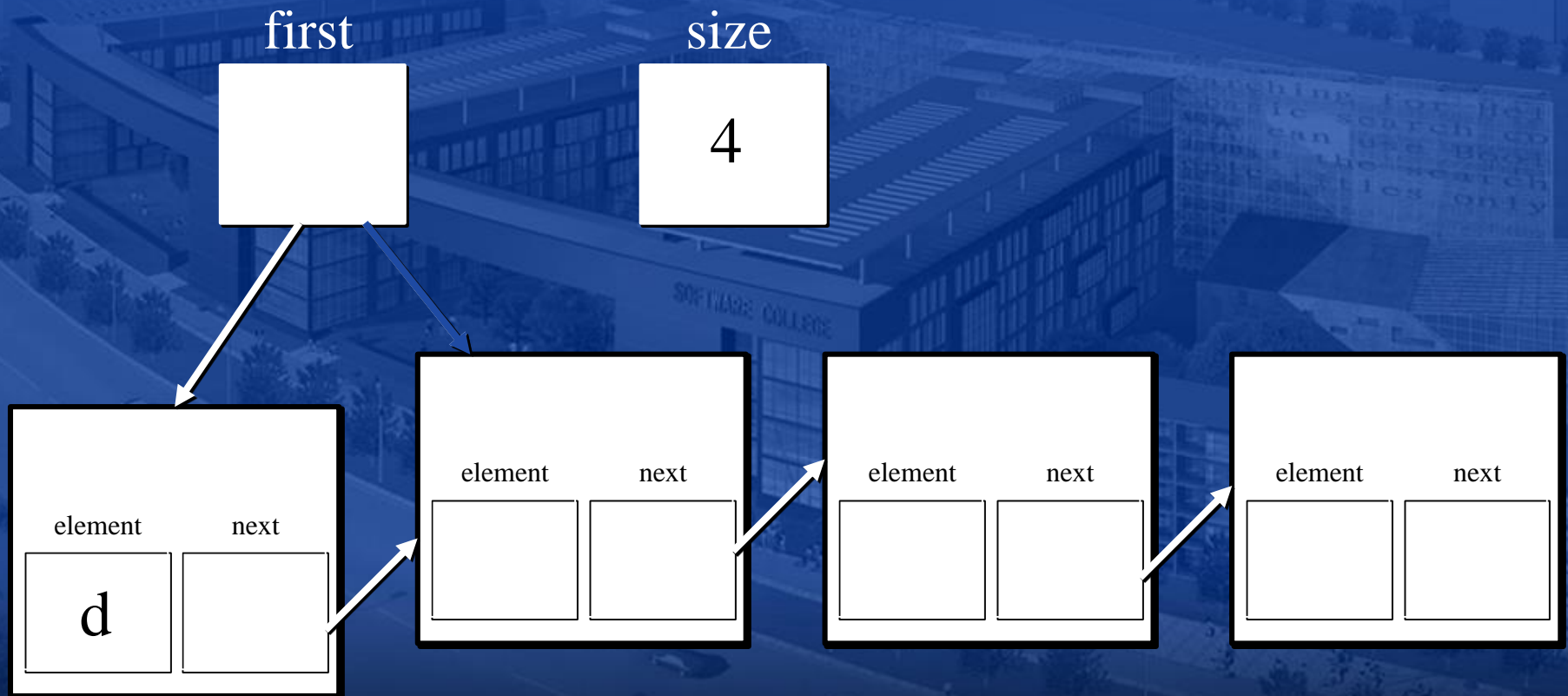
Singly-Linked List Operations

■ insertFirst(d)

- Add a Node with value d as the first position in the list
 - ◆ Construct a new Node object, called n, that contains value d
 - ◆ Change the next reference in n to make it point to the first node in the list
 - ◆ Change the first reference to make it point to n
 - ◆ Increment the size

Singly-Linked List Operations

■ Example: insertFirst(d)



Singly-Linked List Operations

■ Pseudocode of insertFirst(d)

Algorithm insertFirst(d)

Input: the value d to be inserted into the list

Output: the position that contains d

Create node n containing d;

n.next \leftarrow first;

first \leftarrow n;

size \leftarrow size+1;

return n;

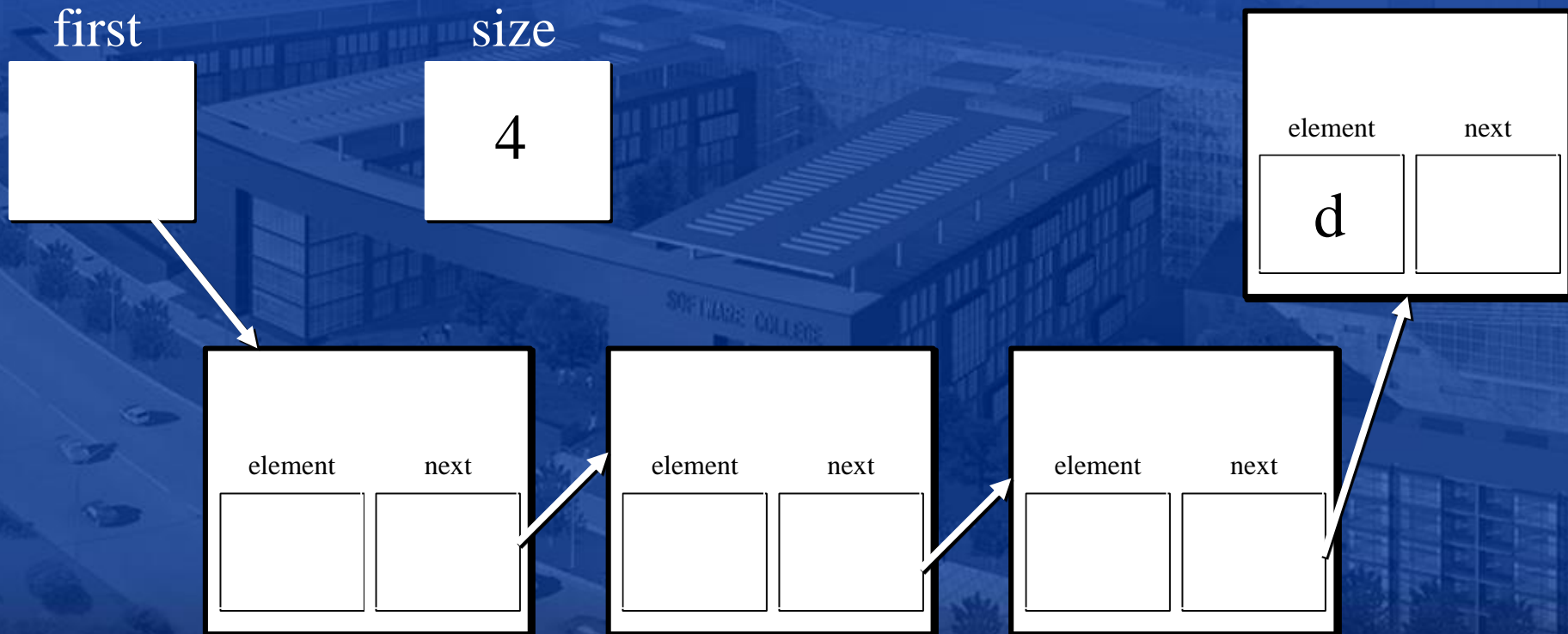
Singly-Linked List Operations

■ insertLast(d)

- Add a Node with value d as the last position in the list
 - ◆ Construct a new Node object, called n, that contains value d
 - ◆ Find the last position, called i, in the list using the last operation
 - ◆ Change the next reference of i to make it point to n
 - ◆ Increment the size
- What would happen if the list is empty?

Singly-Linked List Operations

■ Example: insertLast(d)



Singly-Linked List Operations

■ Pseudocode of insertLast(d)

Algorithm insertLast(d)

Input : the value to be inserted

Output: the position that contains d

if isEmpty() then

 return insertFirst(d);

Create node n containing d;

i ← last();

i.next ← n;

size ← size + 1;

return n;

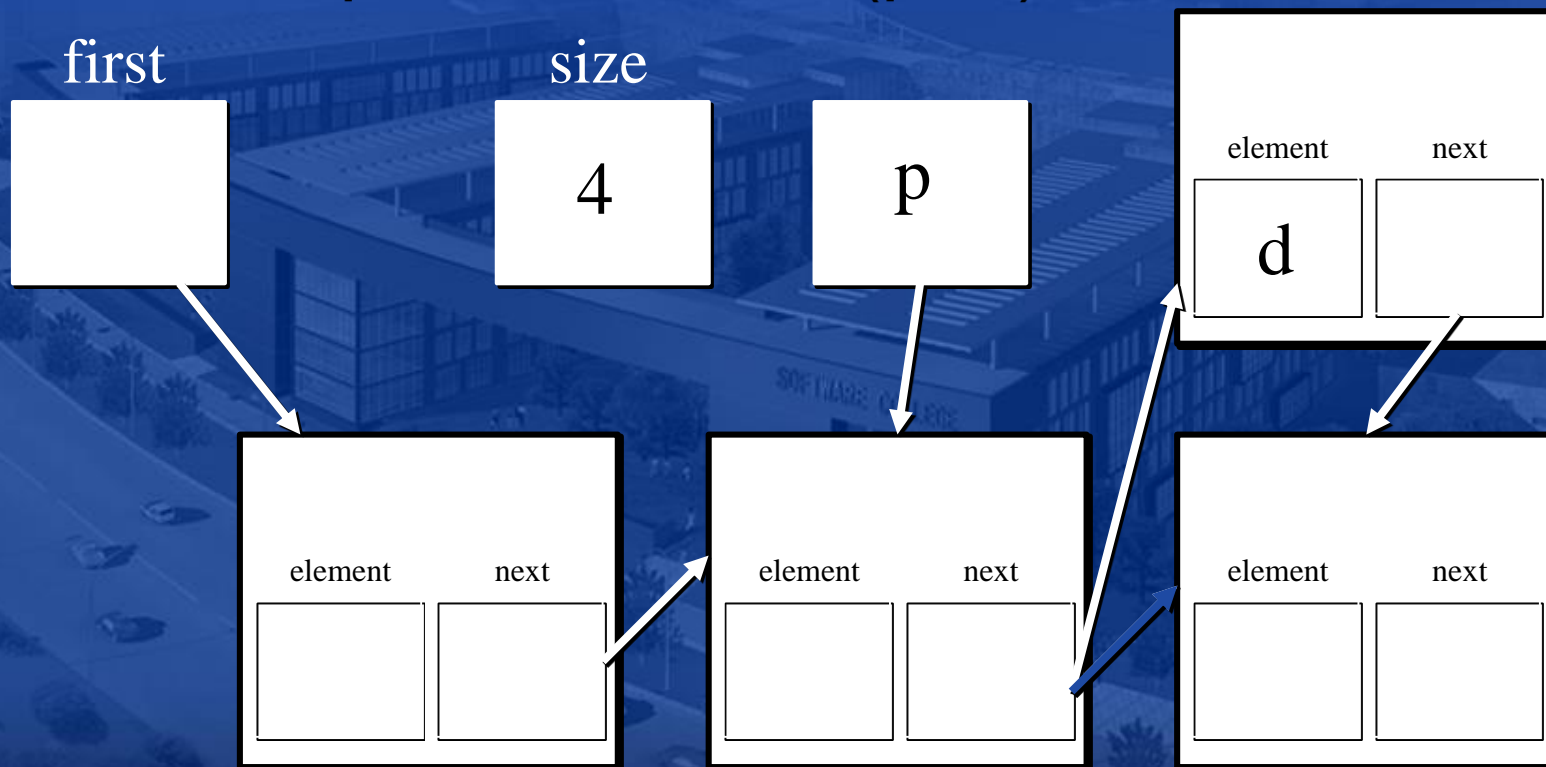
Singly-Linked List Operations

■ insertAfter(p, d)

- Insert a node containing d into the list at the position after p
 - ◆ Construct a new Node object, called n, that contains the value d
 - ◆ Change next of n to make it point to the next of p
 - ◆ Change the next of p to make it point to n
 - ◆ Increment the size

Singly-Linked List Operations

■ Example: insertAfter(p, d)



Singly-Linked List Operations

■ Pseudocode of insertAfter(p, d)

Algorithm insertAfter(p, d)

Input : the value to be inserted and the position it should
be inserted after

Output: the position that contains d

Create node n containing d;

$n.next \leftarrow p.next$;

$p.next \leftarrow n$;

$size \leftarrow size + 1$;

return n;

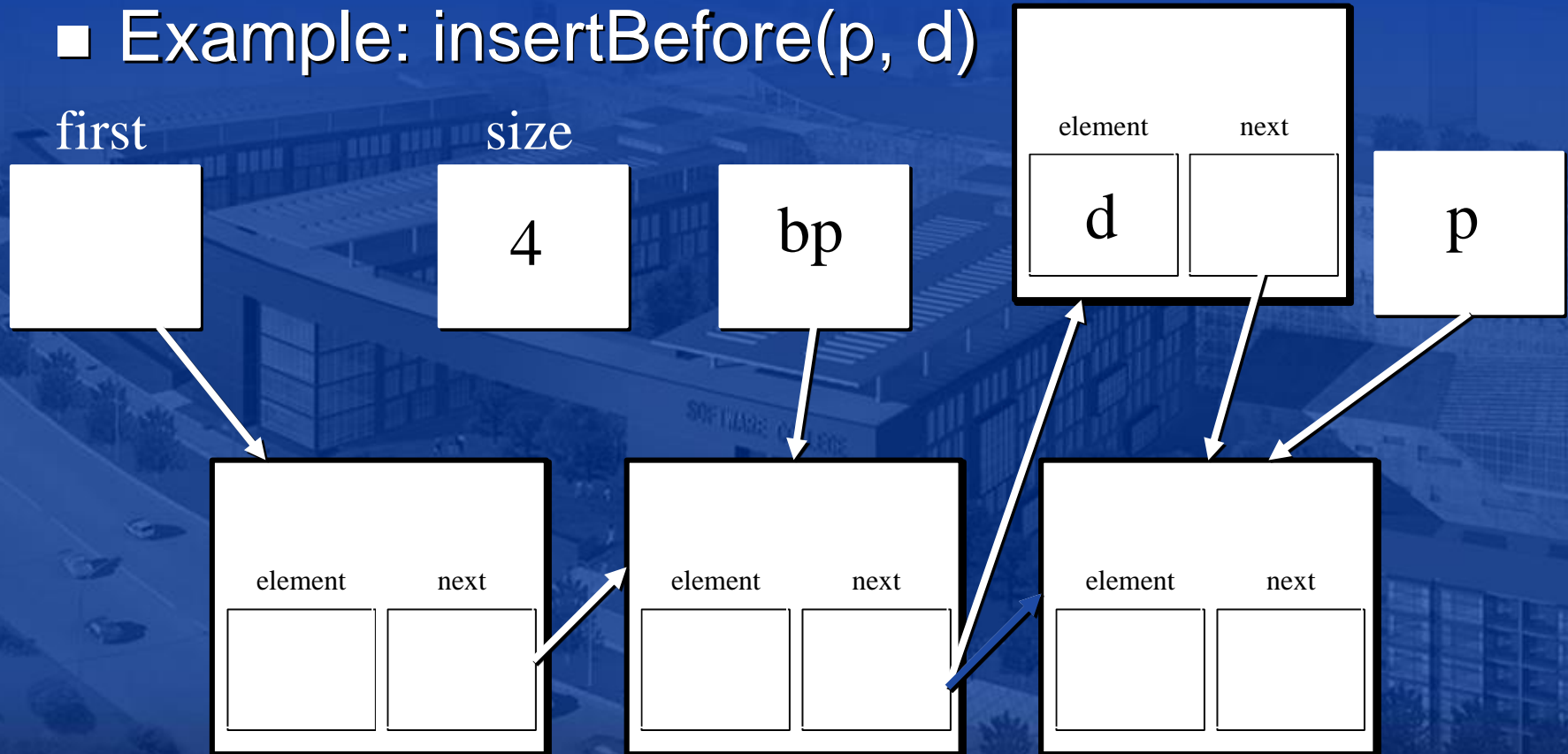
Singly-Linked List Operations

■ insertBefore(p, d)

- Insert a node containing d into the list at the position before p
 - ◆ Construct a new Node object, called n, that contains the value d
 - ◆ Use the before operation to get to the position before p, call it bp
 - ◆ Change next of n to make it point to p
 - ◆ Change the next of bp to make it point to n
 - ◆ Increment the size
- What happens if p is the first position?

Singly-Linked List Operations

■ Example: insertBefore(p, d)



Singly-Linked List Operations

■ Pseudocode of insertBefore(p, d)

Algorithm insertBefore(p, d)

Input : the value to be inserted and the position it should be inserted before

Output: the position that contains d

if p=first then

 return insertFirst(d);

Create node n containing d;

bp \leftarrow before(p);

n.next \leftarrow p;

bp.next \leftarrow n;

size \leftarrow size + 1;

return n;

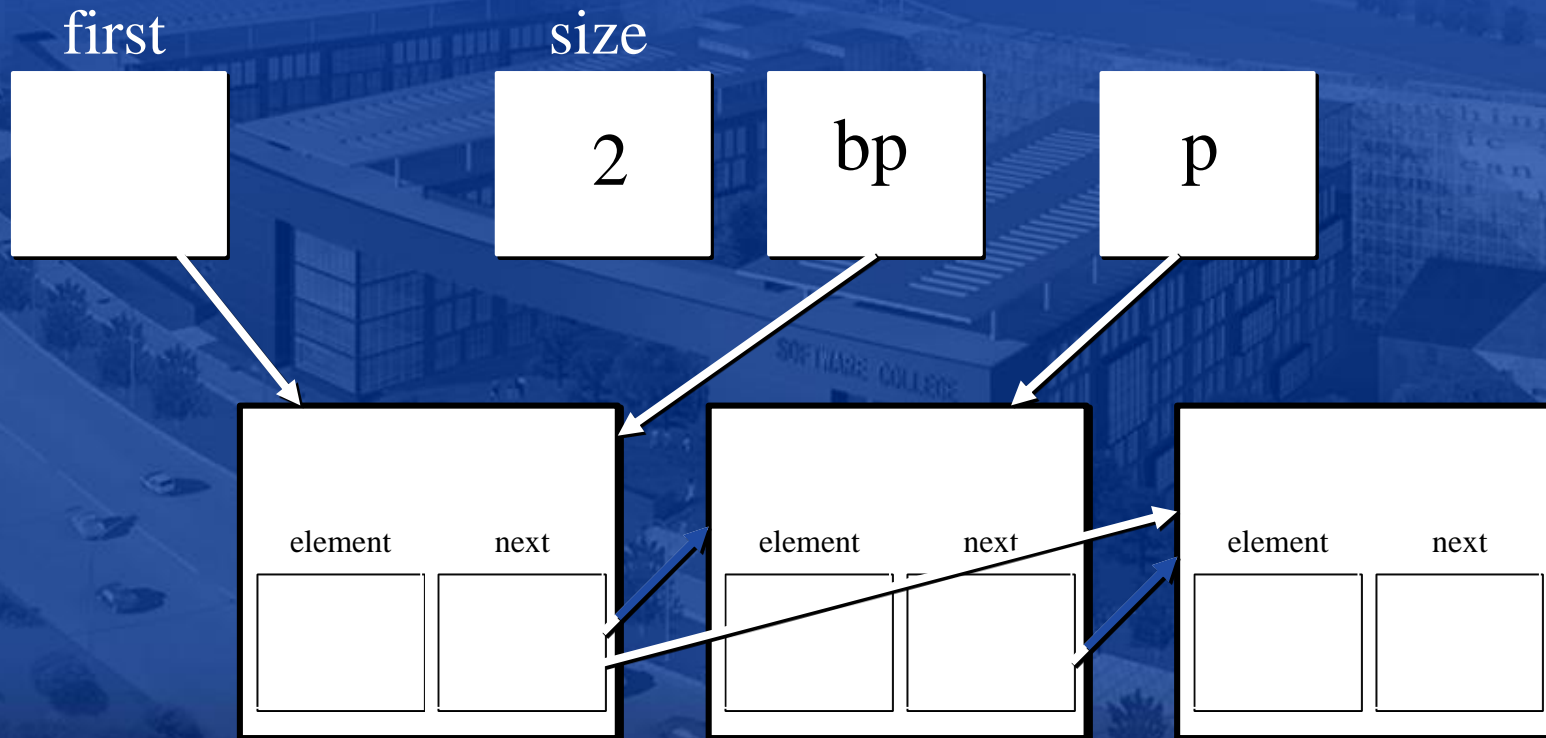
Singly-Linked List Operations

■ remove(p)

- remove node p from the list
 - ◆ Copy the data inside position p to variable d
 - ◆ Use the before operation to get the position before p, call it bp
 - ◆ Change next of bp to make it point to p.next
 - ◆ Decrement the size
 - ◆ return d
- What happens if p is the first position?

Singly-Linked List Operations

■ Example: remove(p)



Singly-Linked List Operations

■ Pseudocode of remove(p)

Algorithm move(p)

Input : the position to be removed from the list

Output: the value that is removed

Copy data of p to variable d;

if p = first then

 first = p.next;

else

 bp ← before(p);

 bp.next ← p.next;

size ← size - 1;

return d;

Singly-Linked List Operations

■ remove(d)

- Remove the node that has value d from the list
 - ◆ Note: d is the value in the node, but not the node itself
- We need to find the object(s) with value d, then remove it/them

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Algorithmic Complexity

Algorithm	Complexity
first()	$O(1)$
size()	$O(1)$
isEmpty()	$O(1)$
after(p)	$O(1)$
last()	$O(n)$
before(p)	$O(n)$
insertFirst(d)	$O(1)$
insertLast(d)	$O(n)$
insertAfter(p, d)	$O(1)$
insertBefore(p, d)	$O(n)$
remove(p)	$O(n)$
remove(d)	$O(n)$

Further Information and Review

- If you wish to review the materials covered in this lecture or to get further information, read the following sections in the text book *Data Structures and Algorithms*
 - 7 - List and Iterator ADTs

Q & A

