

Digital Filters & Spectral Analysis

Lecture 9

FFT

Problem sheet

Solutions

1. A time domain signal, $x(t)=\cos(2\pi ft)$, where $f=4\text{Hz}$, is sampled at 16Hz for 0.25s . Draw the appropriate decimation in time (DIT) FFT flowgraph, showing all complex coefficients and use this to compute the DFT for this signal. Comment on the following:

- Is there any smearing in the resulting spectrum? Explain your answer.
- What is the complexity of your FFT?

Period of cosine:

$$T = 1/f = 1/4 = 0.25\text{s}$$

Sampling time of 0.25s equals one period of the cosine function.

Sampling period:

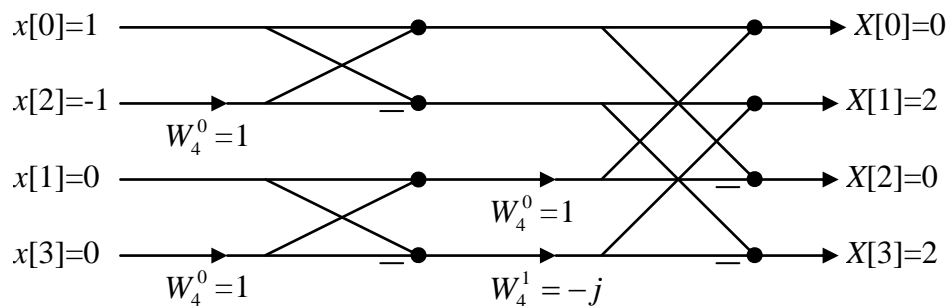
$$T_s = 1/f_s = 1/16 = T/4$$

Sampling at 16 Hz gives 4 samples at $nT_s : n = 0,1,2,3$ $T_s = 1/16 \Rightarrow$

$$x[n] = x(nT_s) = \cos(2\pi f n T_s) = \cos\left(\frac{\pi}{2}n\right) = \{1, 0, -1, 0\}$$

We need a $N=4$ point FFT, which has 2 stages.

Twiddle factor : $W_4 = e^{-j2\pi/4} = -j$



$$X[k] = \{0, 2, 0, 2\}$$

Normalised frequency is $\Omega = 2\pi f/f_s = 2\pi/4$. We sample the spectrum at $k 2\pi/N$ where $N=4$ here. The frequency of the cosine function coincides with an FFT sample and thus there is no smearing.

The FFT requires 4 trivial multipliers 3 of which are by $W^0=1$ and the remaining one is by $W^1=-j$. The FFT also uses 8 complex additions.

2. Show that we can calculate the DFT $G[k]$ of a sequence $g[n]$ of length $2N$ using an N -point DFT and some additional computation. Use the following properties of the DFT/FFT to derive the solution:
- The FFT algorithm is designed to perform complex multiplications and additions
 - The DFT is a linear operation
 - Relevant properties of the DFT
 - The FFT follows a decimation and merge approach
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A. Define two real valued sequences $x_1[n]$ and $x_2[n]$ of length N as:

$$x_1[n]=g[2n], \quad x_2[n]=g[2n+1], \quad n=0,1,\dots,N-1$$

B. Define a complex-valued sequence $x[n]$ as :

$$x[n] = x_1[n] + jx_2[n] \Rightarrow x_1[n] = \frac{x[n] + x^*[n]}{2}, \quad x_2[n] = \frac{x[n] - x^*[n]}{2j}$$

C. Linearity of DFT means that: $X_1[k] = \frac{DFT\{x[n]\} + DFT\{x^*[n]\}}{2}$, $X_2[k] = \frac{DFT\{x[n]\} - DFT\{x^*[n]\}}{2j}$

$$X_1[k] = \frac{X[k] + X^*[N-k]}{2}, \quad X_2[k] = \frac{X[k] - X^*[N-k]}{2j}$$

D. Express the $2N$ -point DFT $G[k]$ in terms of the two N -point DFTs $X_1[k]$ and $X_2[k]$ and proceed as in the DIT FFT algorithm:

$$G[k] = \sum_{n=0}^{N-1} g[2n] w_{2N}^{2nk} + \sum_{n=0}^{N-1} g[2n+1] w_{2N}^{(2n+1)k} = \sum_{n=0}^{N-1} x_1[n] w_N^{nk} + w_{2N}^k \sum_{n=0}^{N-1} x_2[n] w_N^{nk}$$

$$\begin{aligned} G[k] &= X_1[k] + w_{2N}^k X_2[k] \\ G[k+N] &= X_1[k] - w_{2N}^k X_2[k] \end{aligned} \quad k = 0, 1, \dots, N-1$$

2N-DFT for the price of a N-DFT
(plus some additional computation)