
Advanced Networks

(EENGM4211)

Rasheed Hussain

rasheed.hussain@bristol.ac.uk

Best contact : Microsoft Team

1. Introduction
2. Internet Routing and Switching
3. IP Multicast
- 4. Networking for Realtime Applications**

Part 4:

Networking for Real-Time Applications

Multimedia Networking Applications

Classes of MM applications:

- 1) Streaming stored audio and video
- 2) Streaming live audio and video
- 3) Real-time interactive audio and video

Jitter is the variability of packet delays within the same packet stream

Fundamental characteristics:

- Typically **delay sensitive**
 - end-to-end delay
 - delay jitter
- But **loss tolerant**: infrequent losses cause minor glitches
- Antithesis of data, which are *loss intolerant but delay tolerant*.

Examples:

- Internet radio talk show
- Live sporting event

Streaming:

- playback buffer
- playback can lag tens of seconds after transmission
- still have timing constraint

Interactivity:

- fast forward impossible
 - rewind, pause possible!
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Interactive, Real-Time Multimedia

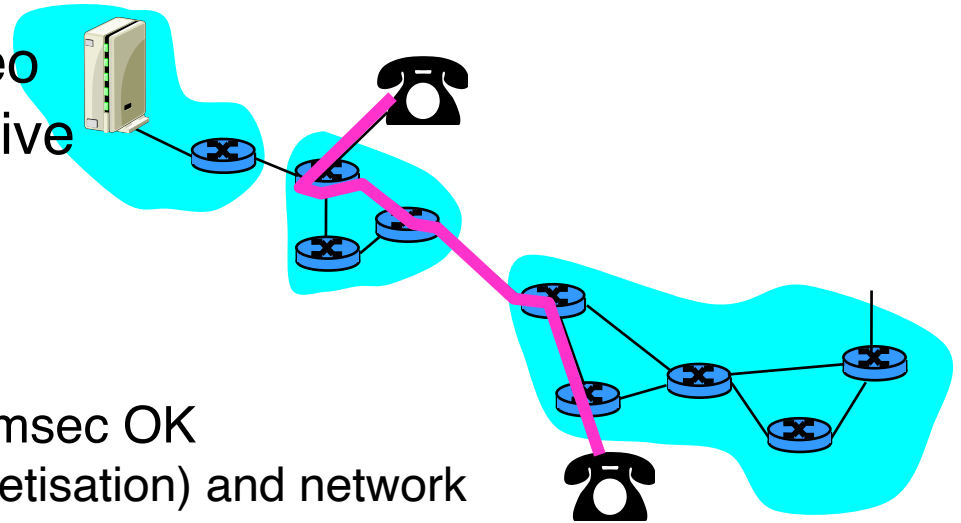
Applications: IP telephony, video conference, distributed interactive worlds

End-end delay requirements:

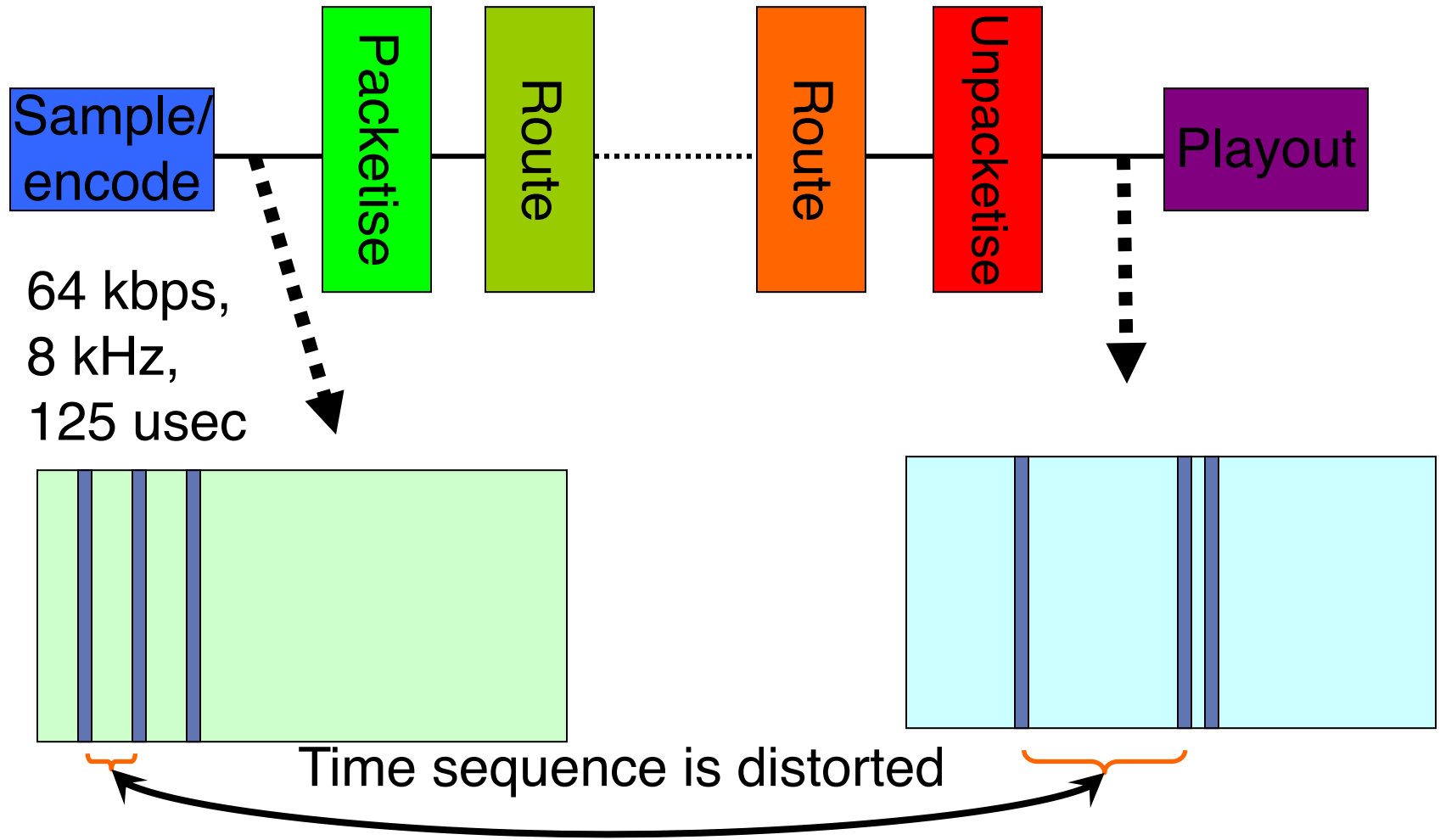
- audio: < 150 msec good, < 400 msec OK
 - includes application-level (packetisation) and network delays
 - higher delays noticeable, impair interactivity

Session initialisation

- how does callee advertise its IP address, port number, encoding algorithms?



Real Time Delivery over IP



Multimedia Over Today's Internet

TCP/UDP/IP: “best-effort service”

- *no* guarantees on delay, loss



? ? ? ? ?
But you said multimedia apps requires ?
QoS and level of performance to be
? effective! ? ?



Today's Internet multimedia applications
use application layer techniques to mitigate
(as best possible) effects of delay, loss

Real Time Delivery – Synchronicity

- Solutions?
 - Try to share or schedule
 - TDM solution like the telephone system
 - wastes bandwidth
 - too complex
 - Timestamps on everything
 - Receiver reconstructs relationship
 - Uses lots of memory

- Average rates must be equal
 - identical at sender and receiver (encode/playout)
 - need estimate of variance
- Delivery delay
 - variable but not unbounded
 - depends on queue sizes of routers
 - most variability $< 40\text{ms}$

User requirements

- Adjust to talkspurt
 - Audio = { talkspurts, silent period }
 - talkspurt is a period of time in which audio data must be acquired and played
- Conferences (high interactivity) require low delay and responsiveness vs. loss and distortion.
- Seminars (low interactivity) require low loss, low distortion but can tolerate delay.

Real Time – Packets and Sessions

- For each packet one needs to know:
 - Format and encoding of content
 - Sequence numbers
 - Timestamps – when should the packet be played?
 - Where (which source) did the packet come from?
 - Q: Why are both timestamps and sequence numbers required?
- For a session one needs to know:
 - Quality of delivery
 - So that one can adapt to network changes, and diagnose faults
 - Who is sending?
 - Who is participating in the session
 - Which sessions are associated
 - So that they can be synchronised.

Real-Time Transport Protocol (RTP) & Application Level Framing (ALF) Principle

(Clark & Tennenhouse 1990)

- Application knows how to process data, therefore
 - Data to be sent over the network should be organised into units that fit application requirements best:
 - What is the best network data unit for a stock-update application?
 - What is the most suitable data unit for a video application?
 - (media format, synchronisation requirements, etc.)
 - What is the most suitable data unit for an audio conference application?
 - Why do we need another protocol?
 - IP does not provide mechanisms for ensuring reliable and bounded delay delivery; content packets can be lost, re-ordered, delivered too late or too early.
 - Application is closer to content processing than the network.
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RTP: Real-Time Transport Protocol

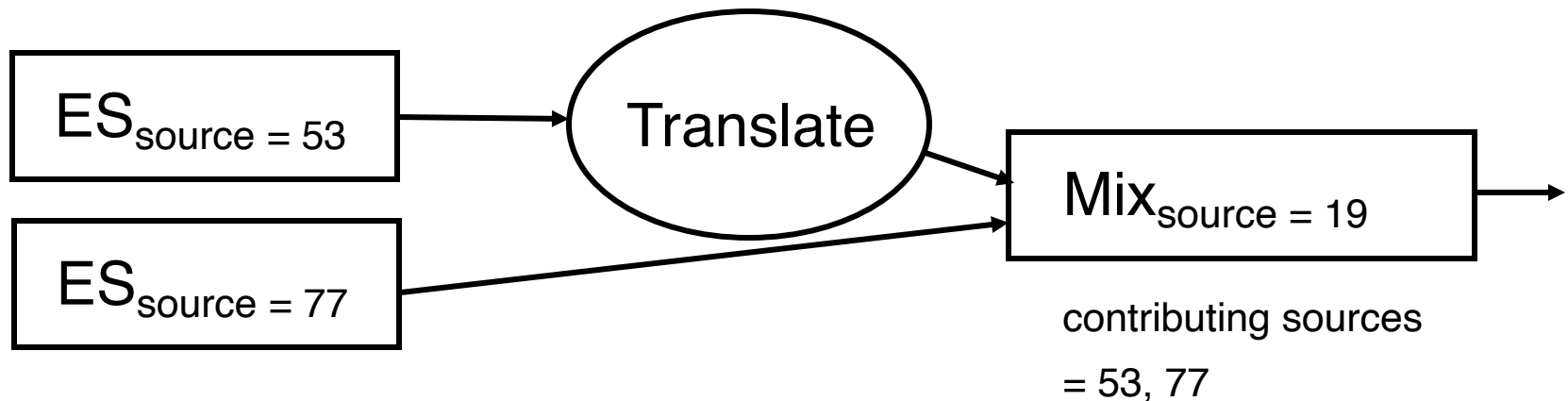
(Schulzrinne et al. RFC 1889 and after)

- Product of IETF Audio Video Transport Working Group (AVT WG)
 - Functional goals
 - lightweight, interoperable
 - easy integration with application
 - mechanism - not policy
 - scalable - from unicast to multipoint with 2 - 1000s participants
 - separation of control from data
 - adaptable to many types of client media
 - independent of underlying network
 - transported in UDP/IP but not necessarily IP
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- Multiple participants
 - May use multicast
 - Multiple streams per participant
 - Audio, video, whiteboard, ...
 - Dynamic membership
 - Session participants may come and go as they want.
 - No central control
 - Participants may not know about each-other.
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RTP Entities

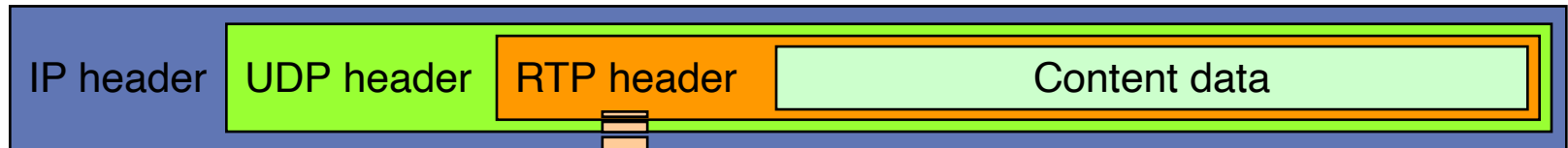
- *End System* (ES) - generates content
- *Mixer* - combines multiple RTP sources into a new RTP packet, re-times and re-formats
- *Translator* - recodes



RTP Functions (1)

- Timing information for playout
 - Reordering information
 - Loss detection for quality estimation, recovery
 - Synchronization
 - network jitter
 - clock drift
 - intermedia (lip sync)
 - QoS feedback
 - Source identification
 - Cannot ensure real-time delivery
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RTP Header Format (2)

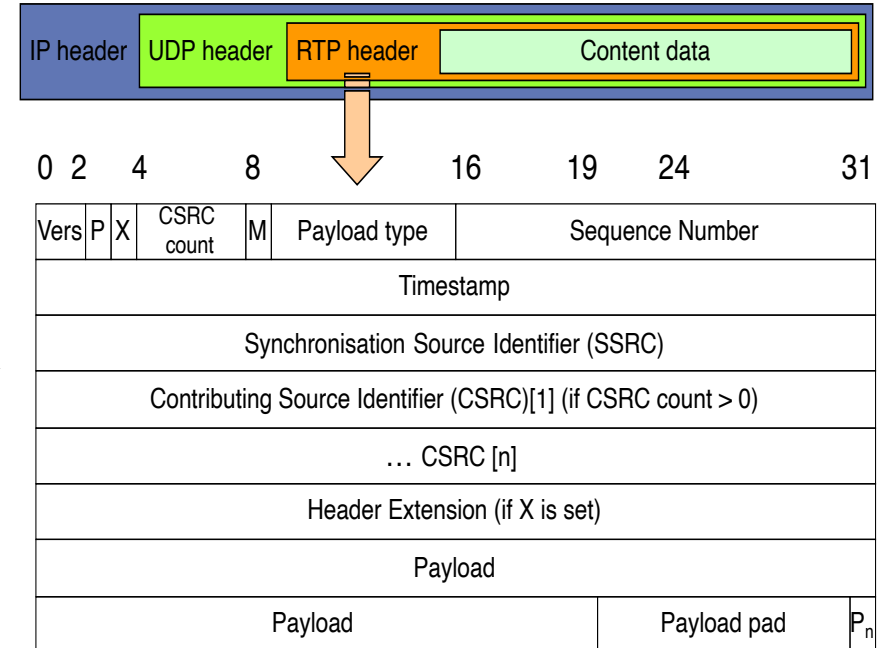


0 2 4 8 16 19 24 31

Vers	P	X	CSRC count	M	Payload type	Sequence Number		
Timestamp								
Synchronisation Source Identifier (SSRC)								
Contributing Source Identifier (CSRC)[1] (if CSRC count > 0)								
... CSRC [n]								
Header Extension (if X is set)								
Payload								
Payload						Payload pad		P _n

RTP Header Fields (3)

- Version is always 2 ()
- **P**: Last field contains padding
 - if lower layer has fixed frame sizes
 - last octet contains size of padding block
- **X**: Header extension is present
- **CSRC count**: Number of contributing sources for a mixed stream
- **M**: Marker used by application to identify significant events, e.g. frame boundaries in a stream



RTP Header Fields (4)

Payload type

- G.721 (2), GSM speech (3), JPEG (26)

Sequence Number

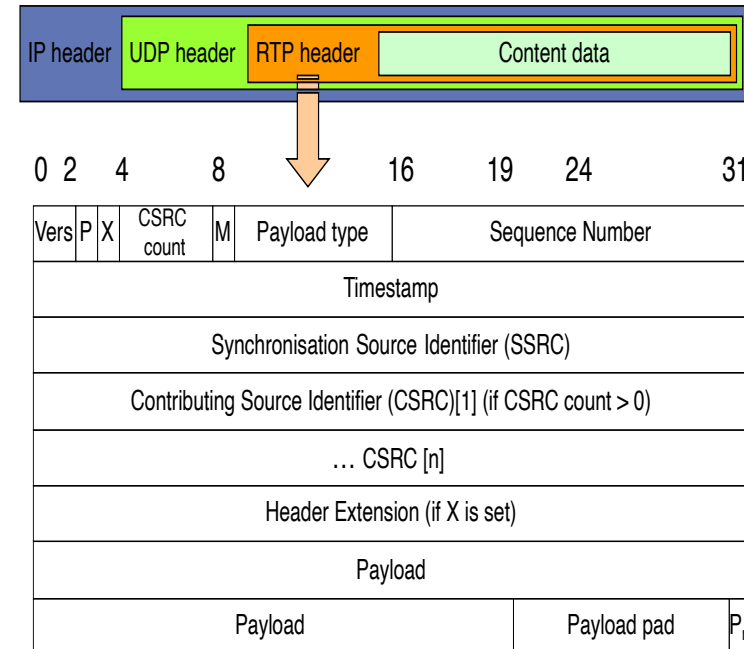
- Increments by 1 for each RTP packet, random initial value, detect loss
- Relate packets to each-other

Timestamp

- Timestamp relates packet to real time
- Timestamp values are sampled from media specific clock!
 - $\text{Timestamp} = \text{packetisation interval} * \text{sampling rate}$
 - E.g. = 160 for 20ms audio sampled at 8kHz
- fragmented frames have same timestamp
- controls playout rate

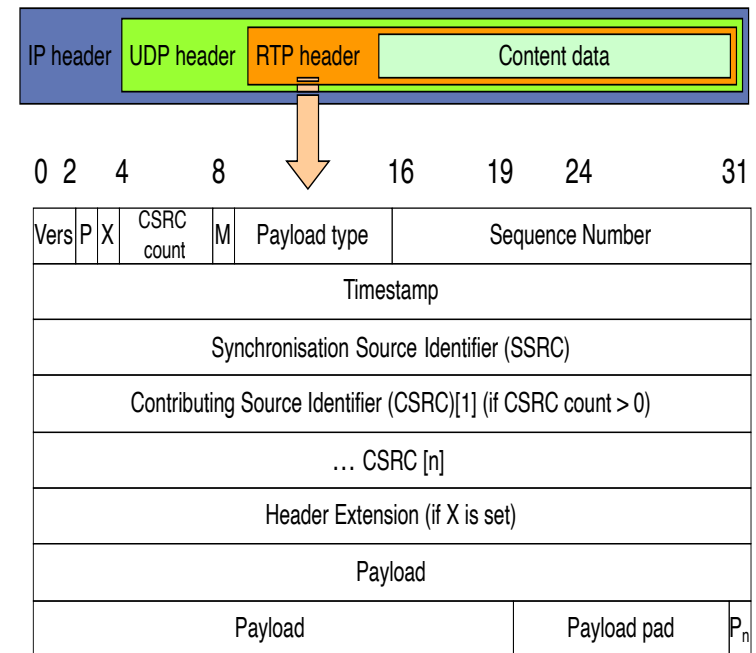
Timestamp vs. Sequence Number

- E.g. video: a frame is send over several packets. All packets will have same timestamp but different sequence numbers.



RTP Header Fields (5)

- Synchronisation source - **SSRC**
 - Identifies the entity that generated the timestamp and sequence number at the sender
 - It is chosen randomly
 - Identifiers are unique
 - Identifies a stream: all packets with the same SSRC go together.
 - What may be a problem?
 - Collision of source identifiers.
- Contributing source - **CSRC**
 - It is used by receiver to identify all sources in a stream



- Audio samples are packetised every 20msec, timestamped and sent with monotonically increasing sequence numbers.
- What about silence periods?
 - Do not send.
 - Problem: How can the receiving side tell if it is loss of packets or it is a silence?
 - The mechanism (sample, encode, timestamp, packetise, sequence number + 1, send) works in this case as well. How?

- Provides monitoring capabilities
 - quality of routes
 - number/state of participants (talker indication)
 - identification and type of sources
 - Feedback to application
 - QoS feedback – sender adjusts rate based on feedback
 - Scalable
 - Randomized control traffic : the rate of control messages decreases as number of participants increases
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RTCP Messages (1)

- Encoded as RTP messages with pre-defined Payload Type fields
 - Sender reports (SR)
 - Real-time of RTP timestamp + sequence number, number of packets and bytes sent
 - Receiver reports (RR)
 - proportion and total of packets lost
 - highest Sequence Number
 - jitter (variance in arrival time)
 - time of last Sequence Number and delay since last Sequence Number
 - Who uses reports?
 - Applications: sender may modify transmission based on feedback
 - Third-party monitors to locate problems.
 - BYE
 - Sent by sources that leave a session.
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- Source description (SDES)
 - maps SSRC/CSRC values to real objects
 - source types include
 - [user*]domain_name (CNAME)
 - name
 - e-mail address
 - phone number
 - location
 - application name

Synchronization of Streams

- Each RTCP sender-report packet contains (for the most recently generated packet in the associated RTP stream):
 - timestamp of the RTP packet
 - wall-clock time for when packet was created.
 - Receivers can use this association to synchronize the playout of audio and video.
 - RTCP can synchronize different media streams within a RTP session.
 - Consider videoconferencing app for which each sender generates one RTP stream for video and one for audio.
 - Timestamps in RTP packets tied to the video and audio sampling clocks
 - Not tied to the wall-clock time
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RTCP Bandwidth Scaling

- RTCP attempts to limit its traffic to 5% of the session bandwidth.

Example

- Suppose one sender, sending video at a rate of 2 Mbps. Then RTCP attempts to limit its traffic to 100 Kbps.
- RTCP gives 75% of this rate to the receivers; remaining 25% to the sender
- The 75 kbps is equally shared among receivers:
 - With R receivers, each receiver gets to send RTCP traffic at $75/R$ kbps.
- Sender gets to send RTCP traffic at 25 kbps.
- Participant determines RTCP packet transmission period by calculating avg RTCP packet size (across the entire session) and dividing by allocated rate.

References

- RTP: A Transport Protocol for Real-Time Applications.
Audio-Video Transport Working Group, H. Schulzrinne, S. Casner, R. Frederick, V. Jacobson. January 1996.
RFC 1889 (<http://www.ietf.org/rfc.html>)
- Some Frequently Asked Questions about RTP
<http://www.cs.columbia.edu/~hgs/rtp/faq.html>