# **Cerebras System**

remarkable hardware accelarator architecure & its programming model

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**Installation and Setup** 

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## How to get an access to SDK

Fill the form on this link.

- Human operators reply, thus it takes more than half a day.
- The reply includes **Dropbox link** to the **SDK files**.

Most of the things about **installation and setup** are <u>here</u>.

• The most important thing is that this SDK is for amd64 architecture.

## Overview of installation and setups

Some of the (critical) things below are **NOT** explicitly written in the guide.

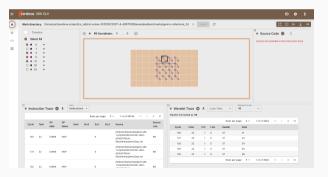
- Azure VM is highly recommended as an environment (Miyabi is not available here...)
  - Recommended instance type is Standard B4ms (4vcpu, 16GiB memory) with 64GiB disk.
- 2. When you connect via ssh,

```
ssh -i ~/.ssh/YOUR_PRIVATE_KEY -Y -L 8000:localhost:8000 YOUR_USERNAME@PUBLICIP
```

- transfer port 8000 to remote port 8000
- YOUR\_USERNAME is **NOT** a resource name.
- YOUR\_PUBLICIP can be seen on the resource via azure home
- You can also refer past spring-training by gotonao

## Overview of installation and setups

- 3. You can follow the guide at installation/setup, but some filenames are changed.
- 4. Finally, you can try remote GUI debug (step7 of the guide).
  - sdk\_debug\_shell visualize after running the test, open http://localhost:8000/sdk-gui at your local browser (eg. chrome)
  - Sometimes, version conflict occurs and show a kind of bugs.



A Conceptual View: Hardware

organization

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A Conceptual View: Hardware organization

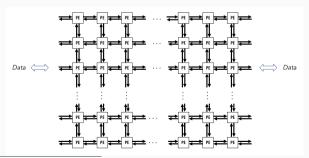
PE (Processing Elements)

A Conceptual View: Programming mode

# Wafer Scale Engine

Cerebras refers its hardware accelarator as a WSE (Wafer Scale Engine).

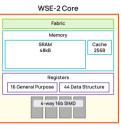
- WSE consists of hundreds (dies) of thousands (数千万個) of independent PE (processing element)s (~ cores).<sup>1</sup>
- The PEs are interconnected by communication links, and they form a two-dimensional rectangular mesh on one single silicon wafer.



<sup>&</sup>lt;sup>1</sup>uses TSMC 5nm processed node

## Characteristics of a PE: memory

- Each PE has its own physically-local SRAM (called local PE memory) <sup>a</sup> with single-cycle access.
  - is 48kB total, consists of 8 banks, and has full datapath bandwidth: 2 64(128)bit read + 1 64(128)bit write per cycle
  - each bank has 6kB, and 32bit wide, has single port
- All the code and data related to the execution on the PE are stored within this memory.
- This physically-local memory is logically local as well (i.e., No other PE are directly accessible to this memory)



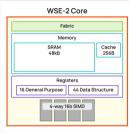
WSE-3 Core

Fabric		
Memory		
SRAM 48kB	Cache 512B	
Registers		
16 General Purpose 48 Data St	ructure	
8-way 16b SIMD		

<sup>&</sup>lt;sup>a</sup>200x normalized memory BW vs. GPU

## Characteristics of a PE: processor

- Each PE has a processor called CE (Compute Engine).
  - 16 general purpose registers, 48 data structure registers
  - Compact 6-stage pipeline
  - Flexible general ops (e.g., arithmetic, logical, load/store, compare, branch) for control processing
- CE has **SIMD** computing unit
  - In the WSE2, 4-way 16 bit SIMD<sup>a</sup>, each way has ALU for FADD, FMUL, FMAC<sup>b</sup> etc.,
  - In the WSE3, 8-way 16 bit SIMD and 16-way 8 bit SIMD



WSE-3 Core

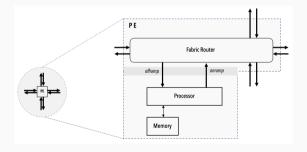
Fabric	
Memory	
SRAM 48kB	Cache 512B
Registers  16 General Purpose 48 Data Structure	
8-way 16b SIMD	

<sup>&</sup>lt;sup>a</sup>which means execute single instruction with 4 different data simultaneously

<sup>&</sup>lt;sup>b</sup>Fused Multiply-Add

## Characteristics of a PE: processor

- Each PE has its own independent PC (program counter).
  - Thus, each PE can execute codes asynchronously by default.



#### Characteristics of a PE: router

- Each PE has the hardware unit for communication (send, receive) called **Router**.
- A Router is directly connected to its own CE via bidirectional link called RAMP (Router ALU Messaging Path).<sup>2</sup>
- A Router is directly connected to the routers of the four nearest neighboring PEs (north, south, east, west)
  - Thus, a router has 4 ports
- A Router has 8 input queues and output queues inside the PE.
  - An **input queue** is a hardware buffer where data is temporarily stored before the entering the CE. (I will explain what this means later)

 $<sup>^2</sup>$ You may have heard this word as  $4 \vee 9 - 7 \times 10^{\circ}$  of highway

A Conceptual View: Programming

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## **Programming Language: CSL**

- To develop code for the WSE, write *device code* in the **CSL** (**Cerebras Software Language**), and *host code* in *Python*.
  - The host code is responsible for copying data to and from the device
  - CSL gives programmers full control of the WSE.
- Then, compile the device code with cslc, and run your program on either
   Cerebras fabric simulator (or the actual network-attached device).
  - The usage is here.

## **Programs and Tasks**

A **CSL program** consists of one or more *subprograms*.

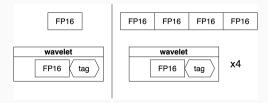
Subprogram has two types (declaration).

- function: callable
- task: a procedure that cannot be called from other code
  - something like atomic (i.e., unsplittable) code block managed by hardware
  - tasks are managed at the specialized hardware unit of CE similar to rich NPC generator
  - A task can be activated (i.e., ready for running) by some hardware trigger (imagine
    a flip of a flag bit)
  - tasks are started by PE hardware (specifically NPC generator of CE)
  - Only **one** task can be executed at a time on the CE
  - Once a task is started by hardware, it runs until it complete. Then, the NPC generator chooses a new task to run.

#### The unit of communication between PEs: wavelet

32-bit messages ( $\sim$  packets), called **wavelets**, can be sent to or received by neighboring PEs in a single clock cycle.

- Arrivals of wavelets trigger something inside the PE
  - Task Activation
  - Stored in a data struct managed with a fabric DSD
- transfering data of massive size (like array, tensor) is splitted into multiple
  wavelets, and data of wavelets that arrive at the destination PE first is buffered in
  the input queue until all wavelets arrive.



#### Virtual Communication Channel: Color

The virtual communication path (channel) through which **wavelets** (packets) travel is called **color**.

- There exists 24 virtual channels used by hardware.
- All colors transfer data on a single physical channel.
  - If multiple colors have wavelets to send via physical path from PE X to PE Y, those colors (not wavelets) are scheduled by hardware arbiter on router.
  - IMPORTANT: The congestion of one color does NOT block traffic of another color; Fairness between colors (not wavelets)
  - For example, RR algorighm could be used as a policy.
- Each wavelet has a 5-bit tag that encodes its color

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## Task IDs and Types of Tasks

Each task can be associated with task ID from 0 to 63.

- Data Task: the arrival of wavelet triggers its activation, its ID is associated with a input queue on the router<sup>3</sup>
- Local Task: the @activate(task\_id) in some other codes within the same PE triggers its activation,
- Control Task: controls other tasks on the same PE as follows, its ID can take any values from 0 to 63.
  - unblock other data task
  - conditional launch of local task

<sup>&</sup>lt;sup>3</sup>In the WSE2, task ID is directly associated with the **color** with implicit linkage between **input queue** and task ID

## The conditions to be ready for execution

There are two conditions for tasks be scheduled by task picker (hardware selector)

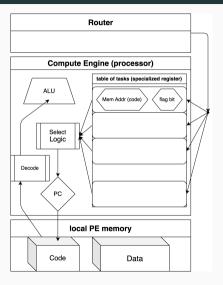
#### Activated

- every task is inactive by default
- programmers can activate the task within the same PE, with @activate(task\_id).
- programmers can activate the task in another PE, with @send\_to\_color(output\_queue\_id).

#### Unblocked

- every task is unblocked by default
- but, programmers can block the ID of a task at compile time, with @block(task\_id).

# Psuedo Image of hardware that manages Data Task



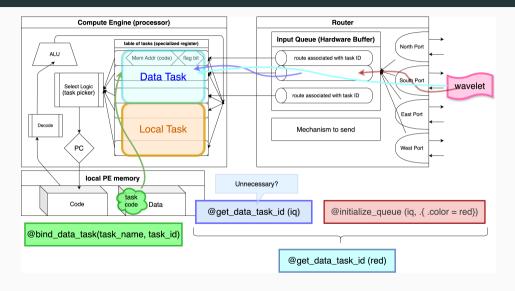
#### Communication via router: Task activation

- 1. Secure **color** for communication with const color\_name: color = @get\_color( . )
- 2. Secure **input queue** with const iq\_name: input\_queue = @get\_input\_queue( . )
- 3. Create data task ID from an input queue

```
const task_id_name: data_task_id = @get_data_task_id(iq_name)
```

- Each data task is bound to an input queue
- 4. Define what to do in the task with task task\_name(wavelet\_data) { ... }
- 5. Define comptime{ ... } block:
  - Bind defined task task\_name to created task ID task\_id\_name with @bind\_data\_task(task\_name, task\_id\_name)
  - Bind input queue(s) to which data task is bound to color @initialize\_queue(iq\_name, .{.color = color\_name});

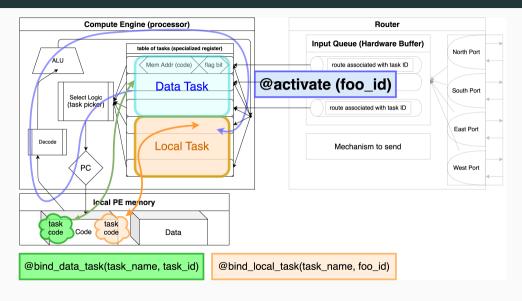
## Psuedo Image of task activation from Router



## Code Template: link computation and communication using Task

```
1 // 7 is the ID of a color with some defined data routing
 2 const red: color = @get color(7):
 4 // On WSE-3. 2 is the ID of an input queue which will be bound to our data task.
 5 const iq: input_queue = @get_input_queue(2);
 6
 7 // For WSE-2, the ID for this task is created from a color.
8 // For WSE-3, the ID for this task is created from an input queue.
9 const red task id: data task id =
       if (@is_arch("wse3")) @get_data_task_id(iq)
       else
                             @get_data_task_id(red);
13 var result: f32 = 0.0:
14
15 task main task(wavelet data: f32) {
16
       result = wavelet data:
17 }
18
19
   comptime {
20
       @bind data task(main task. red task id):
22
       // For WSE-3, input gueue to which our data task is bound must be bound to color red.
23
       if (@is arch("wse3")) @initialize queue(ig. .{ .color = red }):
24 }
```

## Psuedo Image of task activation from another task



## Code Template: link computations using Task

## A data task (main\_task) activates a local task (foo\_task)

```
const red: color = @get color(7):
2 const ig: input gueue = @get input gueue(2): // Used only on WSE-3
  const red task id: data task id =
      if (@is arch("wse3")) @get data task id(ig)
                                                else @get data task id(red):
6
  const foo_task_id: local_task_id = @get_local_task_id(8);
  var result: f32 = 0.0: var sum: f32 = 0.0:
  task main_task(wavelet_data: f32) {
      result = wavelet data:
13
      14 }
  task foo_task() {    sum += result; }
  comptime {
19
      Qbind data task(main task, red task id):
20
      // bind foo_task to the hardware to be managed
21
      @bind local task(foo task, foo task id):
22
      if (@is_arch("wse3")) @initialize_queue(iq, .{ .color = red });
23 }
```

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#### Preliminaries: the hardware instruction of Tensor

#### CSL supports several tensor ops

- Every tensor op has its options
  - asynchronous execution means, release software serialization
  - i.e., start execution even before the previous op is completed, considering whether or not hazard (data, structural) exist

• The format of these ops are (here):

```
1 // ARITHMETIC
2 // dst = src1 * src2 + acc
3 @fmacs(dst. acc. src1, src2 [, options])
4 // dst = src1 {+, -, *, /} src2
5 @fadds(dst, src1, src2 [, options])
6 @fsubs(dst, src1, src2 [, options])
 7 @fmuls(dst. src1. src2 [. options])
8 @fdivs(dst, src1, src2 [, options])
10 // COPY
11 // dst = src
12 Ofmovs(dst, src [, options])
13 // dst = src(size)
14 @copv(dst, src, size_in_bvtes)
  // ELEMENT-WISE OPS
17 @fexps(dst, src [, options])
                                   // exp
18 @flogs(dst. src [. options])
                                   // 1n
19 Ofrelus(dst, src [, options])
                                   // ReLII
20
  // REDUCTION OPS
  @fsum(dst, src [, options]) // sum of items
23 Ofmax(dst, src [, options]) // dst is scaler
```

## DSD: Abstraction of data consists of multiple values like tensor

### Data Structure Descriptors (DSDs) [More details are available here.]

- are a compact representation of
  - a chunk of memory (or)
  - a sequence of incoming or outgoing wavelets
- enable various repeated operations to be expressed using just one hardware instruction
- is an software object which consists of
  - dst\_type: One Of mem1d\_dsd, mem4d\_dsd, fabin\_dsd, Or fabout\_dsd
  - properties: different dst\_type has different properties

#### **ADVANCED:** hardware microthread

hardware **microthread** is an mechanism to distribute and manage hardware resources for enabling asynchronous execution [Details are Microthread IDs, Async DSD Ops]

- Arbitration of hardware resource like ALU, Memory Access Unit, Router Interface (~ scheduling)
- An asynchronous DSD operation can be assigned a microthread ID through .ut\_id = @get\_ut\_id(n)
  - microthread ID could be different from input/output queue ID<sup>4</sup>
  - If multiple DSR/DSD operands have the .ut\_id setting specified, the hardware will
    pick one of them according to the order: dst > src1 > src2
- programmers can attach priority to each thread (including main thread).
- **IMPORTANT:** The programmer is responsible for ensuring that no two concurrent DSD operations share a microthread.

 $<sup>^4</sup>$ In WSE2, the same as output queue ID when using fabs\_dsd, otherwise the same as input queue ID

# The hardware mechanisms to manage DSD: DSR

There have to be the hardware mechanism that utilize DSDs: **Data Structure Registers (DSRs)** 

- DSRs are physical registers that are used to store DSD values
- All DSD operations will actually operate on DSRs behind the scenes, thus all DSD operands to DSD operations must be loaded to DSRs before executing
- Each DSR belongs to one of three DSR files, namely dest, src0, src1 DSR files (i.e., physically distinguished)
  - dsr\_dest: DSR number (レジスタ番地) that can only be used to store a destination operand DSD of a DSD operation
  - dsr\_src0: DSR number that can be used to store a source DSD as well as a destination operand DSD
  - dsr\_src1: DSR number that can only be used to store a source operand DSD
- Basically, compiler allocate DSR (and extra DSR) automatically, but programmers can use them directly with dsr = @get\_dsr, @load\_to\_dsr(dsr, dsd [, option])

### Communication via router: DSD

#### sender

- set colors used to send/recv data between PEs
- 2. assign task ID used by a local task to unblock cmd stream

# Code Template: link computation and communication using DSD

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# Writing the top-level CSL file

There are a few things that programmers need for our device code to form a complete program

- Initialization the infrastructure of the memcpy library with @import\_module
  - In order to allow the host to launch kernels and copy data to and from the device
  - has to specify width and height parameters which correspond to the dimensions of the program rectangle
- A top-level "layout" file
  - define the program rectangle on which our kernel will run, with @set\_rectangle(columns\_dim, rows\_dim)
  - assign a code file to the single PE in our rectangle, with
     @set\_tile\_code(column\_idx, row\_idx, "pe\_program.csl" [, parameters])
  - pass memcpy parameters as a parameter, which are parameterized by the PE's column number(idx), with .{ .memcpy\_params = memcpy.get\_params(column\_idx)}

## Code template of the top-level CSL file

```
1 // Import memcpy layout module for 1 x 1 grid of PEs
 2 const memcpy = @import_module("<memcpy/get_params>",
                                     \{ . \{ . \} \}  width = 1. . height = 1 \}):
 5 layout {
     // Use just one 1 PE (columns=1, rows=1)
     @set rectangle(1, 1):
9
10
     // The lone PE in this program should execute the code in "pe_program.csl"
11
     Qset tile code(0. 0.
12
                        "pe_program.csl".
13
                        .{ .memcpv_params = memcpv.get_params(0) });
14
15
     // Export device symbol for array "v"
16
     // Last argument is mutability: host can read v. but not write to it
17
     @export name("v". [*]f32. false):
18
19
     // Export host-callable device function
20
     @export_name("init_and_compute", fn()void);
21 }
```

# Writing the host code

#### 1. **Import libraries** which is required

- SdkRuntime is the library containing the functionality necessary for loading and running the device code, as well as copying data on and off the wafer.
- MemcpyDataType and MemcpyOrder are enums containing types for use with memcpy calls

## 2. Instantiate (=construct) runner objects like

```
runner = SdkRuntime(args.name, cmaddr=args.cmaddr)
```

- name: specify the directory containing the compilation output
- cmaddr: attach IP address of targetted real accelarator obtained from command-line like --cmaddr \$CS\_IP\_ADDR:9000<sup>5</sup>

## 3. Load and Run device kernel (named init\_and\_compute here)

- Before loading the program, get symbol for copying y result off device
- runner.load() → runner.run() → runner.launch('init\_and\_compute' [, option])

<sup>&</sup>lt;sup>5</sup>CS use port 9000 to connect to the system and launch the program

## Writing the host code

- 4. Copy back result (with many arguments attached as follows)
  - Before copying, allocate space on the host to hold the result
  - 4.1 the array on the host to hold the result y\_result, the symbol on device that points to
    the y array y\_symbol
  - 4.2 To specify the location of rectangle of PEs from which to copy (called ROI), give the northwest corner of the ROI o, o and the width and height of the ROI 1, 1
  - 4.3 how many elements to copy back from each PE in the ROI M (if 2darray, M\*N)
  - 4.4 ROW\_MAJOR specifies that the data is ordered by (height, width, element)
  - 4.5 data\_type keyword specifies the width of the data copied back
  - 4.6 nonblock=False specifies that this call will not return control to the host until the copy into y\_result has finished

```
1 y_result = np.zeros([1*1*M], dtype=np.float32)
2 runner.memcpy_d2h(y_result, y_symbol, 0, 0, 1, 1, M, streaming=False,
3 order=MemcpyOrder.ROW_MAJOR, data_type=MemcpyDataType.MEMCPY_32BIT,
4 nonblock=False)
```

# [Appendix] ROW MAJOR and COLUMN MAJOR

"Learn by example" is the best way to understand the concept.

Here is the example of copying from multiple PEs.

- Configuration
  - Size of ROI: height = 2, width = 3
  - Calculated Data in each PE

```
PEO = [[1, 2, 3],

[4, 5, 6]]

PE1 = [[7, 8, 9],

[10, 11, 12]]
```

- ROW\_MAJOR case
  - Each PE data is copied continuously
  - Copied result (1d):[1, 2, 3, 4, 5,
    6, 7, 8, 9, 10, 11, 12]
- COLUMN\_MAJOR case
  - Elements with the same index in each PE are copied continuously
  - Copied result (1d):[1, 7, 2, 8, 3, 9, 4, 10, 5, 11, 6, 12]

# Finally, Compile and run it!

1. compile device code layout.csl

2. run the host code run.py

#### References

- 1. cerebras SDK Documentation (1.4.0)
- 2. Cerebras Al Day Deck: A closer look at the world's fastest Al Chip
- 3. Cerebras Architecture Deep Dive: First Look Inside the HW/SW Co-Design for Deep Learning