

Definition of Safety

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1 Language

1.1 Term and Pools

We inductively define terms as

- all numerals, symbolic constants, and variables are terms
- $f(\mathbf{t})$ is a term, if f is a symbolic constant and \mathbf{t} is a pool
- $(t_1 \star t_2)$ is a term, if t_1 and t_2 are terms and \star is one of the symbols $(+ - \times / ..)$
- $\langle \mathbf{t} \rangle$ is a term, if \mathbf{t} is a pool, which can have a possibly empty set of terms.

A tuple of terms has the following form t_1, \dots, t_n where t_i is a term.

A pool is an expression of the form $\mathbf{t}_1; \dots; \mathbf{t}_n$ where $n \geq 1$ and each \mathbf{t}_i is a tuple of terms. In particular, every tuple of terms is a pool.

1.2 Constants

We inductively define a term to be *constant* if:

- it is a numeral
- it has form $t \star u$ where t and u are *constant* and \star is one of the symbols $(+ - \times /)$

1.3 Atoms and Literals

An atom has form $p(\mathbf{t})$ where p is a predicate symbol and \mathbf{t} is a pool.

A literal is either an atom or atom preceded by not.

2 Safety

We define the function *provide()* to check if a statement is safe. The *provide* function returns a set of tuples $\{t_1, \dots, t_n\}$. Each tuple t_i has the form (p, d) , where p is a set of variables the statement provides if d variables are provided. If 2 tuples share the same d , the p can be merged. If 2 tuples share the same p , the d can be merged.

For example:

$$\{(\{X\}, \{\}), (\{Y\}, \{\})\} = \{(\{X, Y\}, \{\})\} \quad (1)$$

$$\{(\{\}, \{X\}), (\{\}, \{Y\})\} = \{(\{\}, \{X, Y\})\} \quad (2)$$

2.1 Helper Functions

The *depend* function takes a set of tuples $\{(p_1, d_1), \dots, (p_n, d_n)\}$ and sets $d_i = p_i \cup d_i$ and $p_i = \emptyset$ for each tuple (p_i, d_i)

For Example:

$$\begin{aligned} \text{depend}(\{(\{X\}, \{\}), (\{Y\}, \{Z\})\}) &= \{(\{\}, \{X\}), (\{\}, \{Y, Z\})\} \\ &= \{(\{\}, \{X, Y, Z\})\} \end{aligned} \quad (3)$$

The *merge* function takes a set of tuples $\{(p_1, d_1), \dots, (p_n, d_n)\}$ and a set of variables v

For Example:

$$\text{merge}(\{(\{X\}, \{\})\}, \{Y\}) = \{(\{X\}, \{Y\})\} \quad (4)$$

The output of the *depend* function will also always be in the form $\{(\{\}, d)\}$ where d will contain all the variables. The output of the *depend* function can thus easily be converted to the input parameter v of the *merge* function.

For Example:

$$\begin{aligned} \text{merge}(\{(\{X\}, \{\})\}, \text{depend}(\{(\{Y\}, \{\})\})) \\ = \text{merge}(\{(\{X\}, \{\})\}, \{Y\}) \\ = \{(\{X\}, \{Y\})\} \end{aligned} \quad (5)$$

2.2 Terms

2.2.1 Constants

For any numeral n and symbolic constant f :

$$\text{provide}(n) = \text{provide}(f) = \emptyset \quad (6)$$

2.2.2 Variables

For any variable X :

$$\text{provide}(X) = \{(\{X\}, \{\})\} \quad (7)$$

2.2.3 Tuples

For any tuple of terms t_1, \dots, t_n :

$$provide(t_1, \dots, t_n) = provide(t_1) \cup \dots \cup provide(t_n) \quad (8)$$

For Example:

$$\begin{aligned} provide(X, Y) &= provide(X) \cup provide(Y) \\ &= \{(\{X\}, \{\})\} \cup \{(\{Y\}, \{\})\} \\ &= \{(\{X\}, \{\}), (\{Y\}, \{\})\} \\ &= \{(\{X, Y\}, \{\})\} \end{aligned} \quad (9)$$

2.2.4 Pools

For any pool of terms $t_1; \dots; t_n$:

$$provide(t_1; \dots; t_n) = provide(t_1) \cap \dots \cap provide(t_n) \quad (10)$$

For Example:

$$\begin{aligned} provide(X, Y; Y) &= provide(X, Y) \cap provide(Y) \\ &= provide(X, Y) \cap \{(\{Y\}, \{\})\} \\ &= (provide(X) \cup provide(Y)) \cap \{(\{Y\}, \{\})\} \\ &= (\{(\{X\}, \{\})\} \cup \{(\{Y\}, \{\})\}) \cap \{(\{Y\}, \{\})\} \\ &= \{(\{X, Y\}, \{\})\} \cap \{(\{Y\}, \{\})\} \\ &= \{(\{Y\}, \{\})\} \end{aligned} \quad (11)$$

2.2.5 Terms

For a term of form $f(\mathbf{t})$, where f a symbolic constant and \mathbf{t} a pool:

$$provide(f(\mathbf{t})) = provide(\mathbf{t}) \quad (12)$$

For Example:

$$\begin{aligned} provide(f(X)) &= provide(X) \\ &= \{(\{X\}, \{\})\} \end{aligned} \quad (13)$$

For a term of form $t_1 \star t_2$, where t_1 and t_2 are terms which are not *constants* and \star is one of the symbols $(+ - \times / \dots)$:

$$provide(t_1 \star t_2) = depend(provide(t_1) \cup provide(t_2)) \quad (14)$$

For Example:

$$\begin{aligned} provide(X * Y) &= depend(provide(X) \cup provide(Y)) \\ &= depend(\{(\{X\}, \{\})\} \cup \{(\{Y\}, \{\})\}) \\ &= depend(\{(\{X, Y\}, \{\})\}) \\ &= \{(\{\}, \{X, Y\})\} \end{aligned} \quad (15)$$

$$\begin{aligned}
provide((X * Y) + Z) &= depend(provide(X * Y) \cup provide(Z)) \\
&= depend(\{\{\}, \{X, Y\}\} \cup \{\{\{Z\}, \{\}\}\}) \\
&= depend(\{\{\{Z\}, \{X, Y\}\}\}) \\
&= \{\{\}, \{X, Y, Z\}\}
\end{aligned} \tag{16}$$

For a term of form $t \star c$, where c is a *constant* and t is a term the following holds for \star being one of the symbols $(+ - \times)$:

$$provide(t \star c) = provide(c \star t) = provide(t) \tag{17}$$

For Example:

$$\begin{aligned}
provide(X + 1) &= provide(X) \\
&= \{\{\{X\}, \{\}\}\}
\end{aligned} \tag{18}$$

For a term of form $-t$, where t is a term:

$$provide(-t) = provide(t) \tag{19}$$

For Example:

$$\begin{aligned}
provide(-X) &= provide(X) \\
&= \{\{\{X\}, \{\}\}\}
\end{aligned} \tag{20}$$

2.3 Atoms and Literals

2.3.1 Atoms

For an atom of form $p(\mathbf{t})$, where \mathbf{t} is a pool:

$$provide(p(\mathbf{t})) = provide(\mathbf{t}) \tag{21}$$

For Example:

$$\begin{aligned}
provide(p(X; Y)) &= provide(X; Y) \\
&= provide(X) \cap provide(Y) \\
&= \{\{\{X\}, \{\}\}\} \cap \{\{\{Y\}, \{\}\}\} \\
&= \emptyset
\end{aligned} \tag{22}$$

2.3.2 Literals

For an literal of form $not \ a$, where a is an atom:

$$provide(not \ a) = depend(provide(a)) \tag{23}$$

For Example:

$$\begin{aligned}
provide(not \ p(X)) &= depend(provide(p(X))) \\
&= depend(provide(X)) \\
&= depend(\{\{\{X\}, \{\}\}\}) \\
&= \{\{\{\}, \{X\}\}\}
\end{aligned} \tag{24}$$

For an literal of form a , where a is an atom:

$$provide(a) = provide(a) \quad (25)$$

For an literal of form $not \ l$, where l is another literal:

$$provide(not \ l) = depend(provide(l)) \quad (26)$$

2.4 Comparisons

2.4.1 Comparisons

For an comparison of form $t_1 \star t_2$, where t_1 and t_2 are terms and \star is one of the symbols ($\leq, \geq, <, >, \neq$):

$$provide(t_1 \star t_2) = depend(provide(t_1) \cup provide(t_2)) \quad (27)$$

For Example:

$$\begin{aligned} provide(X \geq Y) &= depend(provide(X) \cup provide(Y)) \\ &= depend(\{(\{X\}, \{\})\} \cup \{(\{Y\}, \{\})\}) \\ &= depend(\{(\{X, Y\}, \{\})\}) \\ &= \{(\{\}, \{X, Y\})\} \end{aligned} \quad (28)$$

2.4.2 Assignments

For an comparison of form $t_1 = t_2$, where t_1 and t_2 are terms:

$$\begin{aligned} provide(t_1 = t_2) &= merge(provide(t_1), depend(provide(t_2))) \\ &\quad \cup merge(provide(t_2), depend(provide(t_1))) \end{aligned} \quad (29)$$

For Example:

$$\begin{aligned} provide(X = Y) &= merge(provide(X), depend(provide(Y))) \\ &\quad \cup merge(provide(Y), depend(provide(X))) \\ &= merge(\{(\{X\}, \{\})\}, depend(\{(\{Y\}, \{\})\})) \\ &\quad \cup merge(\{(\{Y\}, \{\})\}, depend(\{(\{X\}, \{\})\})) \\ &= merge(\{(\{X\}, \{\})\}, Y) \cup merge(\{(\{Y\}, \{\})\}, X) \\ &= \{(\{X\}, \{Y\})\} \cup \{(\{Y\}, \{X\})\} \\ &= \{(\{X\}, \{Y\}), (\{Y\}, \{X\})\} \end{aligned} \quad (30)$$

Hypothetical Example (I don't know how to create this situation in clingo):

$$\begin{aligned}
& \text{provide}(\{(\{X\}, \{\}), (\{\}, \{Y\})\}) = \{(\{Z\}, \{\}), (\{T\}, \{W\})\}) \\
& = \text{merge}(\{(\{X\}, \{\}), (\{\}, \{Y\})\}, \text{depend}(\{(\{Z\}, \{\}), (\{T\}, \{W\})\})) \\
& \quad \cup \text{merge}(\{(\{Z\}, \{\}), (\{T\}, \{W\})\}, \text{depend}(\{(\{X\}, \{\}), (\{\}, \{Y\})\})) \\
& = \text{merge}(\{(\{X\}, \{\}), (\{\}, \{Y\})\}, \{T, W, Z\}) \\
& \quad \cup \text{merge}(\{(\{Z\}, \{\}), (\{T\}, \{W\})\}, \{X, Y\}) \\
& = \{(\{X\}, \{T, W, Z\}), (\{\}, \{T, W, Y, Z\})\} \cup \{(\{Z\}, \{X, Y\}), (\{T\}, \{W, X, Y\})\} \\
& = \{(\{X\}, \{T, W, Z\}), (\{\}, \{T, W, Y, Z\}), (\{Z\}, \{X, Y\}), (\{T\}, \{W, X, Y\})\}
\end{aligned} \tag{31}$$

2.5 Rule

For a rule r in the form of

$$H_1 \vee \dots \vee H_k \leftarrow B_1 \wedge \dots \wedge B_m \tag{32}$$

The following holds:

$$\begin{aligned}
\text{provide}(r) &= \text{depend}(\text{provide}(H_1)) \cup \dots \cup \text{depend}(\text{provide}(H_k)) \\
&\quad \cup \text{provide}(B_1) \cup \dots \cup \text{provide}(B_m)
\end{aligned} \tag{33}$$

For Example:

$$\begin{aligned}
\text{provide}(a(X) \leftarrow b(X)) &= \text{depend}(\text{provide}(a(X))) \cup \text{provide}(b(X)) \\
&= \text{depend}(\text{provide}(X)) \cup \text{provide}(X) \\
&= \text{depend}(\{(\{X\}, \{\})\}) \cup \{(\{X\}, \{\})\} \\
&= \{(\{\}, \{X\})\} \cup \{(\{X\}, \{\})\} \\
&= \{(\{\}, \{X\}), (\{X\}, \{\})\}
\end{aligned} \tag{34}$$

3 Other Examples

$$\begin{aligned}
\text{provide}(p(X, Y + Y)) &= \text{provide}(X, Y + Y) \\
&= \text{provide}(X) \cup \text{provide}(Y + Y) \\
&= \{(\{X\}, \{\})\} \cup \text{depend}(\text{provide}(Y) \cup \text{provide}(Y)) \\
&= \{(\{X\}, \{\})\} \cup \text{depend}(\{(\{Y\}, \{\})\} \cup \{(\{Y\}, \{\})\}) \\
&= \{(\{X\}, \{\})\} \cup \text{depend}(\{(\{Y\}, \{\})\}) \\
&= \{(\{X\}, \{\})\} \cup \{(\{\}, \{Y\})\} \\
&= \{(\{X\}, \{\}), (\{\}, \{Y\})\}
\end{aligned} \tag{35}$$