$$AB = \sum_{t=1}^{n} a_{kt} B_{tk} A = \begin{bmatrix} 1 \\ 1 \end{bmatrix} B = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

$$= \sum_{t=1}^{n} a_{kt} b_{tk}$$

Use independent & identically distributed (iid)
tirialsis to irondomly sample rank-one mots (Axi Bix)
Gives idea for algorithm.

Alg 1

Function:

before addressing the effectiveness, consider the Skyld = { NC Pit | k=it | j=t following: Define SERnxc by Then, C=AS & R=STB Simpling & Rescaling Approach can use any aborithm for) producing matrix multiplication 2) Approach unaffected by sparsity/density 3) Algorithm can be implemented in Huh? some pass over the input metrices (Given Epz)

Analysis of Algorithm

$$Var\left((CR)_{ij}\right) \leq \frac{1}{c} \sum_{k=1}^{n} \frac{A_{ik}^{2} B_{kj}^{2}}{P_{k}}$$

Then $X_t = \frac{Aii_t Bi_t j}{C Pi_t}$ Hence,

$$E[X_t] = E\left[\frac{A_{i} i_{k} B_{i_{k}}}{C P_{i_{k}}}\right] = \sum_{k=1}^{N} P_{k} \frac{A_{i_{k}} B_{k_{j}}}{C P_{k}} \sum_{p_{i_{k}} \in \mathbb{N}} P_{i_{k}} \left[X = x\right]$$

$$= \frac{1}{C} \sum_{k=1}^{n} A_{ik} B_{kj} = \frac{1}{C} (AB)_{ij}$$

$$(CR)_{i,i} = \sum_{t=1}^{c} \chi_{t} \Rightarrow E[CR)_{i,j} = \sum_{t=1}^{c} E[\chi_{t}] = (AB)_{i,j}$$

$$P_{y} = \mathbb{E}[X^{2}] - \mathbb{E}[X]^{2} \Rightarrow \mathbb{V}_{x}[X] \leq \mathbb{E}[X^{2}].$$

$$\Rightarrow Var\left[X_{t}\right] \leq E\left[X_{t}^{2}\right] = \sum_{k=1}^{n} \frac{Aik^{2}Bk_{F}^{2}}{C^{2}P^{2}}$$

$$\left((CR)_{ij} \right) \leq \left(\sum_{k=1}^{n} \frac{A_{ik}^{2}B_{kj}^{2}}{c^{2} p_{k}} = \frac{1}{c} \sum_{k=1}^{n} \frac{A_{ik}^{2}B_{kj}^{2}}{p_{k}}$$

Analysis Cont

Analysis Cont

Thm 22
$$E[NAB-CRN_F^2] \leq \sum_{i=1}^n \frac{||A + i_i||_2^2 ||B_{k+1}||_2^2}{CP_k}$$
 $E[NAB-CRN_F^2] = \sum_{i=1}^m \sum_{j=1}^n E[(AB-CR)_{ij}^2]$
 $Var[X] = E[X^2] - E[X]^2 & E[(AB-CR)_{ij}] = 0$
 $E[(AB-CR)_{ij}^2] = Var[(CR-AB)_{ij}] = Var[(CR)_{ij}]$
 $E[(AB-CR)_{ij}^2] = Var[(CR-AB)_{ij}] = Var[(CR)_{ij}]$
 $E[(AB-CR)_{ij}^2] = Var[(CR)_{ij}^2]$
 $E[(AB-CR)_{ij}^2] = Var[(CR)_{ij}^2]$

Hence,
$$\sum_{i=1}^{m} \sum_{j=1}^{p} E[(AB-CR)^{2}] \stackrel{1}{=} \frac{1}{c} \sum_{i=1}^{m} \sum_{j=1}^{p} \sum_{k=1}^{n} \frac{A_{i} k}{P_{ik}} B_{kj}^{2}$$

$$= \frac{1}{c} \sum_{k=1}^{n} \frac{1}{P_{ik}} \left(\sum_{i=1}^{m} A_{i} k \right) \left(\sum_{j=1}^{p} B_{kj}^{2} \right)$$

$$= \frac{1}{c} \sum_{k=1}^{n} \frac{1}{P_{ik}} \|A_{kk}\|_{2}^{2} \|B_{kk}\|_{2}^{2}$$

$$= E\left[\|AB-CR\|_F^2\right] \leq \sum_{k=1}^n \frac{\|A_{kk}\|_2^2}{C} \frac{\|B_{kk}\|_2^2}{C}$$

Analysis Cont

(5

We show this choice of Epr3 minimizes $E[||AB-CR||_F^2]$:

Minimire & subject to $\Sigma p_n=1$. Use Lag multiplier!

$$g(p_1,...,p_n) = f(p_1,...,p_n) + \lambda \left(\sum_{k=1}^{n} p_k - 1\right)$$

Diff wirt px & X, then set = 0.

$$\frac{\partial q}{\partial P_k} = -\frac{1}{P_k^2} \|A_{kk}\|_2^2 \|B_{k*}\|_2^2 + \lambda = 0$$

$$\frac{\partial q}{\partial x} = \sum_{R} -1 = 0 \implies \int \overline{\lambda} = \sum_{n=1}^{N} \|A_{n+n}\|_{2} \|B_{n+n}\|_{2}$$

(Minimizers $\frac{\partial^2 g}{\partial \rho_k} > 0 \quad \forall k$)

Other choices for Epz3 TH.2

1. Suppose I px = 1 & Px = \frac{B || A xx ||_2 || B x * ||_2}{\frac{D}{2} || A a x ||_2 || B x * ||_2}

for OKBEI. Them,

E(11 AB-CR 11 = BC (= 11 Axx 112 11 B kx 112)

E NAB-CRIP = = = = = = | A+K | 2 | BK + | 2

< \frac{1}{\beta c} \left(\sum_{\beta c}^{\beta} || A_{\kappa k} ||_2 || B_{\kappa k} ||_2 \right)^{-1}

2. $\sum p_{k} = 1 & p_{k} > \frac{\beta_{k} \|A_{kk}\|_{2}}{\|A\|_{2}^{2}}, OL\beta \leq 1$

Then, E[NAB-CRIF] = 1 BC NAND NBIP

E[NAB-Cellip] = I | Alp | Alp | Brah = I | Alp | Bla

3.
$$\sum P_{k}=1$$
 & $P_{k} \geq \frac{\beta \|B_{k} + \|^{2}}{\|B\|_{F}^{2}}$, $0 \leq \beta \leq 1$ \exists

Then $E[\|AB-CR\|_{F}^{2}] \leq \frac{1}{\beta c} \|A\|_{F}^{2} \|B\|_{F}^{2}$

P\$ identical to #2.

$$(\hat{\Sigma} \|A_{kk}\|_{2} \|B_{k} + \|_{2})^{2} \leq \hat{\Sigma} \|A_{kk}\|_{2}^{2} \hat{\Sigma} \|B_{k} + \|_{2}^{2}$$

So #I bound generally better than #283.

Suppose
$$col(U) \subseteq col(A)$$
. Space important part of the cal(N).

Interested in calculating UTU.

Let $U \in \mathbb{R}^{n \times d}$ ($n >> d$)

Use idea of sampling & rescaled 2005 of U.

By Thin $22!$
 $E[\|U^TU - R^TR\|_F^2] = E[\|I_d - R^TR\|_F^2] \le \frac{d^2}{RC}$

Puch chosen to said: $\sum R = 1$ & $\sum R = 1$ &

9 Since || All2 = ||AllF, can also conclude III-RTRN2≤ 2 ω/ 90% considence (& using the same c). Thm 36 Us Rand (n>>d), UTU=IA Construct R as described at bey. of §4.3. Let $c \ge \frac{96d}{B z^2} ln \left(\frac{96d}{B z^2 \sqrt{8}} \right)$ Then w/ prob = 1-8: | UTU-RTR | 2 = | Id-RTR | 2 = 2 = Nead Lemma for pol Lumma 38 let 21,..., x be i.i.d copies of a d-dinensional rand vector x, w/ $\|x\|_2 \leq M$ & $\|E[xx^T]\|_2 \leq 1$ Then 4270: Nt ≤ ziziT - E[zz] Na ≤ a holds w/ prob at least $1 - (2c^2) \exp\left(-\frac{c\alpha^2}{16M^2 + 8M^2\alpha}\right)$

Detine random row vec.
$$y \in \mathbb{R}^d$$
 as
$$P[y = \frac{1}{\sqrt{\rho_{k}}} U_{k,k}]^{2} P_{k} \geq \frac{\beta \|U_{k,k}\|_{2}^{2}}{d}$$

for
$$k = 1: N$$
, "y is rescaled with row of U w/ prob px

[[] [] [] [] i.i.d copies of

For
$$k = 1: N$$
, $y = 1: N$,

$$= \sum_{k=1}^{n} p_{k} \left(\frac{1}{p_{k}} U_{k*}^{T} \right) \left(\frac{1}{p_{k}} U_{k*} \right) = U^{T} U = I J.$$

$$(y i) called "isotropic") so $N = [y^{T}y] |_{2} = 1$

$$(an apply Lemma!)$$$$

By prev Lemma 125yty - UU 1/2 | RTR-Iall2 < E With prob > 1- (2c2) exp (- c 22). Let S be failure prob. Then we seek c s.t.

(20rexp(16M2+8M2) = S. Also require SE1= SO,

$$\frac{1}{2} \frac{(2^{2})}{(2^{2})^{16M^{2}+8M^{2}}} \leq \frac{(2c)}{16} = \frac{1}{2} \frac{1}{2} \frac{(2c)^{2}}{16M^{2}+8M^{2}} \leq 2n \left(\frac{2c}{18}\right)$$

$$= \frac{c}{2n(2^{c}/\sqrt{8})} \geq \frac{2}{2^{2}} \left(16M^{2}+8M^{2}\epsilon^{2}\right)$$

Pf court

$$\|y\|_2 = \frac{1}{\sqrt{p_k}} \|y\|_2$$
 $p_k \ge \frac{\beta \|y\|_2}{d}$

$$\frac{2}{\xi^{2}}\left(16M^{2}+8M^{2}\xi\right) \leq \frac{2}{\xi_{1}^{2}}\left(16\frac{d}{\beta}+8\frac{d}{\beta}\right) = \frac{48d}{\beta\xi^{2}}$$

$$\Rightarrow \frac{ac/\sqrt{6}}{en(ac/\sqrt{6})} \geq \frac{96d}{\beta \epsilon^2 \sqrt{6}}$$

$$\left(\frac{z}{\sin(z)}\right) = \frac{2 \operatorname{nenh}}{\operatorname{en}(z)} = \frac{2 \operatorname{nenh}}{2 \operatorname{enh}} = \frac{2 \operatorname{nenh}}{2 \operatorname{enh}} = \frac{2 \operatorname{nenh}}{2 \operatorname{enh}}$$

Set
$$2e = \frac{2c}{\sqrt{8}} d \eta = \frac{96 d}{\beta \epsilon^2 \sqrt{8}}$$

Then it suffices for c to sertisfy

$$\frac{2c}{\sqrt{8}} \ge 2 \frac{96d}{\beta 2\sqrt{8}} \ln \left(\frac{96d}{\beta 2\sqrt{8}} \right)$$
to guarantee