

I Can See Your Future! Analysing the Risk of Observable Callbacks in Android System

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ABSTRACT

The feature of event-driven acts as a key role that makes Android application differentiate from traditional PC software. For many of those events are hardly predicted and could not be observed by other applications, attackers are similarly impossible to engage corresponding attacks by finding the vulnerabilities of such an event-driven mechanism. However, of various kinds of events offered by either user or system, there are still events that can be received by more than one applications and further, could offer important basic resources to predict specific behaviours of targeted application.

PEC-threats(Public Events Callback)

In this paper, we perform a systematic study of such "public events" and their callback functions, and analyse potential security threats inside them. By exploiting the prediction capability of such callbacks, we demonstrate typical kinds of proof-of-concept attack examples, including spoofing, phishing and privilege escalation. To evaluate the real influence of such threats, we implement a static analysis tool to detect a large scale of real world applications achieved from Google Play market. The detection result reveals such callback threats are pervasively existing in Android market applications. Since the threats involve in the fundamental design of the Android framework, it is very hard for ordinary developers and users to defense them. Given this challenge, we also proposed a migration strategy.

CCS Concepts

•Computer systems organization → Embedded systems; *Redundancy*; Robotics; •Networks → Network reliability;

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ability;

Keywords

Android security; Information leaks; Permissions

1. INTRODUCTION

Android platform has been already proven to be prevalent as a typical operating system installed in nowadays smart-phone and tablet. Different from traditional PC program and software, applications within Android are event-driven and without a launching main function. That is to say, with different events, the program control flow would change correspondingly. Thus, call back mechanism that is used to response an event has played an important role in Android framework.

Callback mechanism is originated from PC programs, which provides an asynchronous way for a specific class to invoke another's method when a certain event happens. In Android, callback functions reside everywhere of program structure in a typical app, especially in UI handling. For instance, after user click a button, a new activity would be popped up, presenting corresponding functionalities expected by user. In this case, a typical callback mechanism is used to response this specific user-click event. That is to say, the new activity would not be activated until user event happens. Besides, the events from system, other apps or life-cycle have similar callback responding mechanism.

On the other side, Android sandbox is known as an effect mechanism to protect apps from other one's attack. Sandbox restricts the resources to be offered for certain independent entities, leaving only very few ways for them to communicate with others. Actually, through the usage of conventional so-called side-channels, certain indirect inferring gradually evolves into a typical sort of attacks. In general, this sort of attacks threat a broad range of apps causing side-channels(such as process status) exist in system level. In other words, most of developed apps would be unable to avoid the influence brought by features from system level. Nevertheless, the impact appears limited for the deviation posed by sophisticated algorithms from current side-channel

approaches. For instance, [9](UI state) treats several selected process statuses as side-channels, and utilizes them to infer a proper emergence time for the fake phishing Activity devised by attackers. However, the critical inference bases on a series of data collection and computations, leading to an inevitable derivation in term of device and targeted app itself. A typical example is that the typical Activity transition signal presence in [9] is quite weak and the shape is also changed in our reproducing experiment. Thus, using the same procedures as proposed in [9] would hardly perform an ideal attack. The similar situation exists in other existing side-channel attacks aiming to Android apps as well [41][18].

Compared to conventional side-channel attacks, we introduce a new type of potential threats called public event callback(PEC) threats, which leverage the weakness of public event callback mechanism within Android apps. Generally, callback functions expect a typical event to trigger their execution, and correspondingly the event is specified to sent to these callbacks. For example, the explicit intent specifies the single address(package name) of the target recipient and sends the intent out through a broadcast event. The broadcast only can be received by the specified single recipient. The entire broadcast process is properly protected by such single receiver mechanism. It is hard to obtain any data of such intent by other apps from the application level. However, not all the cases emerge so ideally. There still exists a large spectrum of public events that can trigger the execution of callbacks from different components and even different apps. We call this kind of callbacks PECs(public event callbacks). Examples includes system event, implicit intent, service running status, etc., which are insensitive on surface but offer attack materials for other apps with evil purpose. In particular, as the source of the expected event is possible to be exposed to public, the invoking time of a typical callback function could be inferred by other apps. As a result, the inferring result evolves to a variety of corresponding attacks directed against the display-sink, such as spoofing, phishing, privilege escalation, etc.

However, the inferring process might not be as such simple as imaging. One challenge is that attackers have to fall into code level to gain necessary materials of victim apps overcoming the code confusion. Moreover, the attackers have to detect out at least one feasible inferring path across numerous control flow paths within targeted app. Such path has to contain a callback function that can be invoked by an exposed public event and a display-sink(e.g., startActivity or Toast.makeText) point with exposure features. The other challenge is that it should to be proven that there exists a determinable path from the start of such callback function to the display-sink, which involves more fine-grain analysis of program control flow. To achieve this goal, branch conditions and method invocation have to be properly handled to avoid the restriction of permission request and user-interaction. Aiming at above challenges, we implement a statical analysis prototype tool CSDroid with 4k lines python codes to automatically conduct this inferring process to 2375** apps collected from Googly Play. Of all the targeted apps, we find 783** such sensitive paths within 354** targeted apps, including some popular apps like **. Moreover, it's surprising that 69% ** of apps with service component suffer from such threats, which alerts us a high risk of such potential security threats within background

components widely existing in generic apps.

Given such an app with certain PEC weaknesses, an adversary attacker could devise a rich variety of attacks. Since such attacks are mainly based on the display-sink, the attack types are somehow similar with traditional UI attacks. However, as the display-sink should be reasonably inferred from the PEC, adversary app needs to concern more details about time-delay and program flow conditions. Besides that, traditional UI attacks mainly rely on activity display of victim app; display-sink contains more types of displays such as dialog and notification, which allow attackers to flexibly craft more attacks contents. To validate the feasibility of such PEC weaknesses exploitation, we explore and implement a set of proof-of-concept attacks based on such threats, including not only traditional phishing and spoofing to sensitive pages, but also newly dialog and notification exploitation for privilege escalation and other privacy steal.

To migrate the impact of such PEC threats, we also discuss the defence strategy from two parts: Android system and developers.

To summarize, the main contribution of this paper are:

1. We systematically study vulnerabilities of Android app lead by abuse of public event callbacks as well as discover various **four** types of attack vectors leading to such PEC threats.
2. By a set of corresponding proof-of-concept attacks, we validate the feasibility of such PEC exploitation and illustrate the severity of such threats.
3. We introduce a tool implemented to automatically detect the PEC threats within targeted apps. By our study, a wide range of apps are affected by this kind of PEC attacks.
4. We provide a mitigation strategy to the PEC threats, under which a more security callback invoking mechanism can be established.

2. BACKGROUND

In order to better understand the mechanism of this kind of attacks, firstly some relative background knowledge needs to be introduced.

Android platform is a kind of open source operating system established by Linux core. Developers could design and develop their own app product on Android in a flexible way, causing Android freely provides a variety of open source libraries and frameworks for developers. Of which, four types of components plays the key role in Android app's construction, namely *Activity*, *Service*, *Broadcast Receiver* and *Content Provider*. *Activity* is designed as a graphical interface for interacting with users. Nevertheless, *Service* is mainly used for handling complex computing job in background without a graphical user interface. *Broadcast Receiver* provides a global listener that could trigger specific code logic when receive relative message. *Content Provider* providers a universal interface for data storage and management, and for sharing data between different components.

2.0.1 Callback Categories

As a side-effect of user-driven mechanism, Android provides numerous of callback interfaces to be implemented by generic developers. Different from conventional PC program launched by the *main* function, callback function could only be executed when a specific event happens. The trigger-events originated by different sources are equipped with different trigger features. For instance, a *low-battery* event

would be triggered only when the rest quantity of battery becomes as low as a pre-defined level, which is irrelevant to user's interaction. We categorize the callback functions as three groups in term of their trigger events, namely *life-cycle* callbacks, *user-driven* callbacks, and *system-driven* callbacks.

Life-cycle callbacks Android apps would spontaneously arouse corresponding callback function according to the life-cycle step it resides. For instance, an Activity executes its *onCreate* function only when it is initially loaded, and executes the *onPause* callback after it loses the screen focus. Life-cycle callbacks are somehow executed along with user's interactions and impacted by users. However, it performs fairly transparent to users, as a part of internal mechanism of Android system. Due to constraint of sandbox mechanism, the life-cycle phrases (Create, Pause, Destroy, etc.) of an given activity are entirely isolated from components belong to other apps. That is to say, any app component is unable to capture the life-cycle statuses from other apps. This sets up a huge obstacle for adversary app that wants to conduct attacks over target benign app.

User-driven callbacks User-driven callbacks interact with users in a direct way. User regularly are guided to perform diverse actions on screen, such as tapping, dropping up and down, click, double-click, etc., to operate a new app. Right after capturing an action from users, Android system would react following the corresponding callback function that reflects the functionality and logic designed by developer. Generally,

System-driven callbacks System-driven callbacks are normally triggered by a given system event which could be captured when a given system status changes. Not being the dependency of any apps, system status could be obtained by any app in theory so long as one is granted proper permissions. Generally, the permission granting could hardly attract user's attention for system event such as battery status is insensitive to user. To the end, system event provides an indirect entry for adversary app getting parts of the running process of other apps that implement system-driven callbacks interfaces.

2.0.2 Callback and Components

Callback in Activity Most of callbacks in Activity are user-driven and life-cycle, causing activity is mainly used for providing a visible graphic interfaces so as to interact with users. As aforementioned, upon the protection of sandbox, these callback statuses are hardly obtained by other apps only if the event triggered by user is irrelevant to screen focus. For instance, a shake action from users would lead to *sensor* related events, which have no especial association with the screen. Besides, still some system-driven callbacks are implemented in activity for feeding the requirements of functionalities.

Callback in Service Handling program logic and computing tasks in background, Service has no direct interaction with users, which is typically different from Activity. By this means, callbacks from service only come from life-cycle and system statuses. Yet, user action without screen focus could be an exception, which we will discuss later.

Service life-cycle consists of five callback functions(*onCreate*, *onStartCommand*, *onBind*, *onUnbind* and *onDestroy*) but has a simpler life-cycle procedure compared with Activity. Being the same as Activity, these life-cycle callbacks are isolated from components of other apps so that has little utility

value for attacker.

Regularly, most of *Services* are designed to wait for a typical signal and then invoke corresponding callback functions to push their computing result to an output interface, such as a *dialog* or an *Activity*. In general, the signal could be any type of event, invocation from other component or a program condition. If the trigger event belongs to system events, it can be easily obtained by component from other apps. This makes possible for attackers to perceive the execution process of a target app.

It is also possible to trigger a callback through an user-driven event in service, which is irrelevant to visible screen (e.g., shaking the phone to trigger *SensorChanged* event). Without specified receiver, this kind of events could be easily obtained by other apps either.

Callback in Broadcast Receiver The *Broadcast* has a wide spectrum of sources, from system or other components and can be delivered across all the installed apps through a common medium *Intent*. Generally, callback functions in Broadcast Receiver only act as a channel that responses the coming Broadcast and then launches corresponding *service* or *activity* to handle it. However, once lacking of strong constraint, Broadcast Receiver could suffer from the threats of repeatedly receiving trash message from an adversary Broadcast sender. Similarly, it is also possible that a typical Broadcast is intercepted by a deliberately designed Receiver. Therefore, developers normally add some protection measures, such as setting the "export" property as false or using self-defined permission, to prevent from such attacks.

Table1 lists typical system broadcast event examples.

3. EXPLOITING CALLBACKS

Android framework provides a variety of callback functions for developers conducting their apps a particular responding after a specific event happens. Yet, abusing these callbacks would incur undesirable consequence as well.

3.1 Motivating Example

We take the well-known app Wechat as a motivating example. Wechat is known as a popular IM(Instant Messenger) tool for users sharing information and connecting with each other. Of diverse functionalities it provides, "Shake" is rather welcomed, which enables users to randomly obtain an online chatting partner so long as slightly shaking the device several times.

Although harmony in surface, there exists potential crisis in the user shaking action which can be easily perceived by any other app through the **SENSOR_SERVICE**, a system service for managing inner device sensor. That is to say, only with the **SystemSensorManager** object instantiation and the arrangement of shaking action size, an adversary app could clearly receive the same action event signal as Wechat. Further more, associating with the running status of Wechat through **AMS(Activity Manager Service)**, adversary app could naturally make a judgement about the accurate time when the Wechat responding page would be popped up after user shaking. Although there exists an exception that user could shake the device in other page of "Wechat", the possibility of the exception's emergence is scarcely small in practice because user shakes device in such wide margin normally with intensive intention. **Figure 1** shows how the "Shake" event is "stolen" by others.

It seems that the PEC threats is easily ignored by de-

Table 1: System broadcast event callback

Event	Callback	Permission	Usage
Intent.ACTION_BATTERY_LOW	onReceive	N/A	The battery level has fallen below a threshold
Intent.ACTION_BOOT_COMPLETED	onReceive	N/A	Android is up and running
Intent.ACTION_DEVICE_STORAGE_LOW	onReceive	N/A	Storage space on the device is getting limited

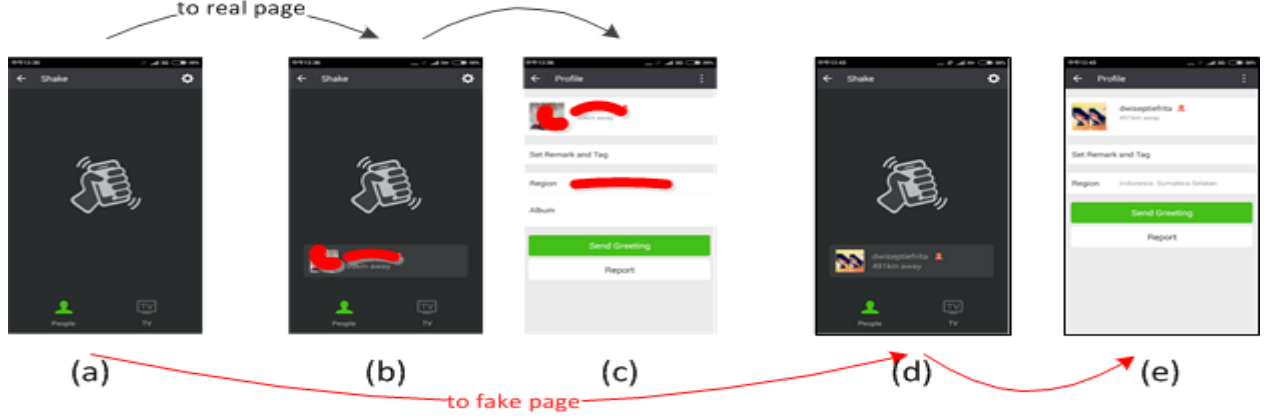


Figure 1: Motivating example about the PEC threats. In (a), a user prepare to try the "Shake"; (b) shows a normal case that a real result popped up after shake; (c) presents the further information of such real page after user clicks the result; (d) shows the case that the shake action is perceived by an adversary app, which pop up a fake "result" page after it receives the shake action; (e) presents the further fake information crafted by attacker, which is used to induce user to convince it.

velopers and app markets. We see the callback time leakage severely because it offers a big change for attackers. A wily attacker could devise diverse attacks utilizing the PEC threats. For instants, attackers could deliberately construct a similar result page that links to a malicious third-party chatting tool, or even mimic a fake chatting page to pretend to chat with user for malicious purpose. We presents greater details about the attacks in section **.

3.2 PEC State Model

In order to clearly illustrate the callback threats, we view the threats process as a kind of callback state model. Each callback within a installed app can be described as (S,F,E,C), where "S" denotes current state of app, which can be observed by other apps; "F" denotes the final sink contained by the callback function; "E" denotes the event that invokes executing of the callback function; "C" denotes the proper control flow conditional values towards the "F".

State.

State(s belongs to S). The State refers to current program state of an app, where the running location can be inferred. Similar to the Event, the State here also needs to be exposed to other apps.

Before **Android 5.0** emerging, the execution state of component like Activity and Service can be easily observed by the **getRunningTasks** method in **ActivityManager** object. A more detailed observation needs to request the "GET_TASK" permission, which could somehow attract the attention of most of normal users. Unfortunately, Android starts to constraint the usage of this permission from recent **Android 5.0** for the security consideration. The "REAL_GET_TASK"

permission is only available to system apps; while normal third-party app is forbidden to request. A substitute solution is to request new "PACKAGE_USAGE_STATS" permission, and to disable warning that permission is system level and will not be granted to third-party apps. Again, this approach needs extra permission request.

Although it is difficult to obtain the accurate running task above the version of 5.0, the running application package name is still easily obtained by **getAppRunningProcess** in **ActivityManager**. The good news is that the usage of "getAppRunningProcess" and "getRunningServices" are still available in the newest version and need not request any extra permission. Thus, we can judge if a typical app is running through the "getAppRunningProcess"[101]. Through a continuing such observation to a target app, we can determine if its launching activity is running on foreground. More fine-grained, the "getRunningServices" can be used to distinguish which service is running in background.

There is a kind of **AccessibilityService** that can capture other apps' states like "click, window change, input, etc.". However, it needs to let user grant the "service" in setting. So we could design a permission-escalation spoofing to enable the **AccessibilityService**.

Display-sink.

Display-sink(f belongs to F). Display Sink is based on the conception of leakage detection, which refers to a set of label functions that can expose data or information to outside. However, different from traditional conception, Display Sink here specifically refers to state output that could be perceived by users rather than information leakage, including UI page, dialog or notification, etc.

As to the Display-sink, only those output vectors which can be leveraged by attackers are considered. Most of which are related to UI activity, causing the UI is designed as the main channel to interact with users. Some of the UI vectors are mentioned in recent research[102](e.g., startActivity API, Toast message), yet we also collect some other UI related vectors, e.g., AlertDialog.Builder API, NotificationManager API. Generally, both of Dialog frame and Notification need response from users(e.g., click) and then engage corresponding reaction functionalities. This process offers chance for attackers to interfere the users behaviours, which is insecure.

Besides, some non-UI vectors are also contained, e.g., audio, vibrate related API, sendBroadcast. Similar to UI related vectors, audio and vibrate related API also engage information output although as different information forms("voice and vibration"). The "SendBroadcast" is able to deliver information between different apps, breaking the constraint of sandbox. Therefore, the app that receives the broadcast could be operated by others.

1 for Dialog

Generally speaking, it goes against Android design and UI guidelines that directly launching a Dialog from Service. Therefore, if AlertDialog is roughly created from a Service there would pop up an error report causing Dialog is implemented relying on a particular Activity. However, still it would pretty work that setting a system alert Dialog with the "SYSTEM_ALERT_WINDOW" permission grant. This kind of Dialog has global feature that can be invoked by a Service or a system event rather than an Activity.

2 for Toast

Similar to Dialog, the Toast, in general, also should be constructed based on an existing Activity. The reason of that is one of the parameters needed by Toast function refers to the main UI context the one belongs to. Therefore, directly popping up from Service would not work causing the UI context is erroneously referred. In principle, the communication between main thread and sub thread mainly relies on the message queue and message cycle in Android, which follows the communication mechanism in traditional Windows. In practice, a common-used tool class "Handler" can be used to insert the self-defined thread message into the main thread queue, so that the main thread UI context can be obtained by the Toast.

3 for StartActivity

Some Activity or Service is set self-defined permission to improve the security protection. For those low risk "protectionLevel", namely "normal" and "dangerous", we can still apply corresponding self-defined permission to access the target components. Nevertheless, for those high risk level, namely "signature" and "signatureOrSystem", currently there is no any effect way to access the target components by a third party app. /////// This kind of cases also need to be taken enough attentions in real attacks.

Event.

Event(e belongs to E). The event here refers to those can be observed by public. For example, normally user-click event does not belong to the set, causing it generally only can be preserved by local app components; other app is forbidden to obtain this event due to the isolation of sandbox.

Following introduction aforementioned, Event here contains three categories.

1. Those Events originated from system are loaded by either Broadcast or serviceManagers. For serviceManagers ones, they normally need corresponding low-risk permission request(e.g., "android.permission.BATTERY_STATS", "android.permission.VIBRATE ") to manipulate the event information. However, system broadcast need not any permission, leading to more threats.

2. Those Events originated from other apps are loaded by implicit Intent or pendingIntent(contains implicit Intent). For explicit Intent, the destination is clearly defined as a typical app so that it is impossible for other app to obtain this kind of Intents. However, implicit Intent and pending-Intent could be sent to a set of apps so as long they match the Intent features.

3. Those Events originated from users are loaded by serviceManager. These events generally are equipped with global characteristics that impacts not only single app. For example, user click event could only affect the activity that occupies focus; while sensor change event would impact all the apps installed. Some of serviceManagers like sensorManager, audioManager need not permission request.

Program Control Flow.

Reverse engineering

Static Analysis

Control Flow(C). Control Flow is a collection of conditional values over the path from a typical State to a Final Sink. In theory, as for a callback threat vulnerable app, the "C" needs to guarantee an observable and feasible path that can reach the Display-sink. However, the "C" that strictly conforms above definition rarely emerges in real apps, so that we adopt a more flexible mechanism in practice which will be introduced in next section.

Attackers need to find a feasible path from the State to the Final Sink within the given operable source code of a target app. For judging the existing of such path, we first collect the condition set of the paths to the Final Sink. Then a signal used to refer to if it is visible for each value, event and function emerging in the condition set is created. To describe more clearly, we denote the following concepts:

1. c: the condition value of each branch.

2. CS: the condition set collected from a program control flow path.

3. cv(a) : the condition value of variable a and a = variable | function | event

4. s(a): the signal value of variable a and a = variable | function | event, whose type is boolean.

s(a) represents if the variable a can be observed by components from other apps. s(a) = true represents yes, s(a) = false represents no.

5 $CS = \cap (cv(a) \& s(a)) = (\cap(c)) \& (\cap s(a))$

Above equation means the value of condition set of a particular path consists of two parts, 1. the value of each variable; 2. the signal value of it. A feasible path needs not only has a solvable condition set, but also observable for each variable. However, practically we normally use the second equation to compute the feasibility of a target path. First, computing out if it is feasible for each branch ; then , considering if it is able to be observed for each variable with the branch condition.

A traditional technique for above computation is symbolic execution. However, it needs unacceptable overhead in practice. **To the end, we using an improved SMT approach to**

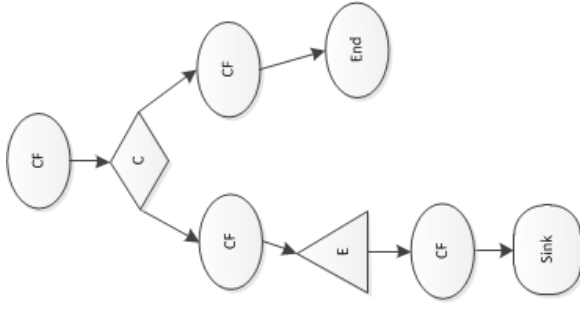


Figure 2: control-flow

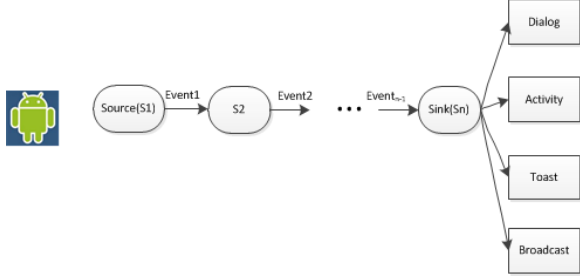


Figure 3: principle0

solve the constraint conditions. Normally there should be feasible from source to sink in a target app, otherwise there exists a designation shortage with in the target app. Under this prerequisite, we only need to consider $s(a)$, namely if all the variables are observable.

Table 2 lists the main vulnerability vectors with the callback threats.

4. ATTACK

4.1 Exploiting Condition

App Description.

Timer.

4.2 Threat Model

In order to clearly illustrate the attack process, first the threat model need to be depicted. Suppose that we have a victim app installed in a given device, and it contains corresponding callback threats vulnerabilities. Besides, suppose that we have an adversary app installed in the same device and start its service running on the background stealthily. The goal of our adversary app is to engage diverse attacks

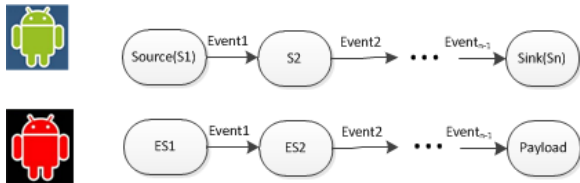


Figure 4: principle1

through the victim app. In theory, the adversary app need not apply any permission when it is installed. However, in practice, it needs to apply some common-used permissions like "Internet" and insensitive permissions like "batter related permission" according to the attack type. The adversary needs to visually conceal its behaviours from users and also minimize its impact on the device performance.

The adversary is supposed having the capability of using public available resources and analysing them on the fly. Again, it should stay aware of the running states of apps and some particular services. As introduced in "Callback State Model", this capability can be obtained by particular methods within ActivityManager class.

4.3 Attack Overview

The PEC could be used to construct various types of attacks, including mainstream attack methods like spoofing, phishing, privilege lifting, etc. An adversary app could perform a customized response in terms of a targeted event. Moreover, a more complex attack would rely on the specific installed victim apps. By predicting the behaviours of such victim apps using PEC, the adversary app could engage appropriate attacks avoiding user's awareness. Base on whether utilizing victim apps, we divide attacks as two main categories.

4.4 Utilizing Public Events

4.4.1 Spoofing

1. Adversary apps could be designed as a kind of spoofing pages, which would be popped up right after a particular event happens. The spoofing page content should be closely related to such event so that users have no doubt about it.

4.4.2 Phishing

4.5 Utilizing Victim App

Following represents typical attacks forms that utilizes the PEC of victim app. According to the display window a PEC ends to, we representatively list four main types of attack examples, namely dialog & toast privilege lifting, notification forgery, shaking-display spoofing and activity phishing.

4.5.1 Dialog & Toast Privilege Lifting

Dialog and Toast used to perform as the notice windows to notify users some intermediate results or alert information after apps successfully receive events. Once the received events are open to public, an adversary app would be easy to predict the time when a dialog or toast really emerges as long as the targeted app lacks of complex condition constrains in its control flow of the key callbacks. A common case comes from the notice related to those system resources like battery, memory, voice, etc. Normally, these kinds of notices are sent by pre-installed or system manager apps. Also, users are with fully trust for these apps in security part because they are professional and have numerous of users. However, this trust acts as a big opportunity that can be exploited to achieve the attack goal in attacker's eyes.

For example, when the memory rests as low as a given level, some system manager apps(360 mobilesafes, Battery Doctor, etc.) would pop up corresponding dialog with alarm information("Dear user, your device has little memory available, clean the data residue to achieve more free memory?"),

Table 2: Typical vectors of PEC

State	Display-sink	Event	Control Flow
1.Activity running state(Launching page)	1.startActivity	1.Broadcast from system: {Intent.ACTION_BATTERY_LOW}	1.branch or cycle condition analysis
2. Service running state	2. Toast message	2.Implicit Intent& pending-Intent	2.method invocation analysis
	3.AlertDialog.Builder	3.ServiceManager(need not permission grant): sensor-Manager, memoryManager	
	4.sendBroadcast (intent)	4.ServiceManager(need permission grant): battery-Manager	
	5.AudioManager API		
	6.VibrateManager API		
	7.NotificationManager API		

```

private void displayBriefMemory() {
    final ActivityManager activityManager =
        (ActivityManager) getSystemService(ACTIVITY_SERVICE);
    ActivityManager.MemoryInfo info = new ActivityManager.MemoryInfo();
    activityManager.getMemoryInfo(info);
    if(info.availMem < threshold){
        AlertDialog.create()
    }
}

```

Figure 5: control-flow

and if user clicks agree button, a new process launched by the manager app starts running. Then the user could be at ease at the memory trouble.

A typical implementation logic is represented as following:

The key point is the memory status(info.availMem) are public resources and can be obtained by any app installed in the device without any permission request. Moreover, the implementation control flow is designed so simple that the condition logic and variable "threshold" could be easily reconstructed in an adversary app. A experienced attacker could also obtain the these key logic and variables through reverse-engineering and binary analysis. As a result, the condition "info.availMem < threshold" as a public events makes the function "AlertDialog.create()" a PEC, and attackers are able to judge the creating time of the AlertDialog from its own reconstructed adversary app installed in the same device.

This time capture offers a rich variety of opportunities to engage attacks escaping user's notice, of which a typical one is privilege lifting. Figure * presents an attack procedure for an attacker getting the "Accessibility Service" grant by user. The Accessibility Service is a special functionality offered to help disabled people to get the component status and handle the device operations more conveniently. Since Android 4.0, it supports richer interfaces for developers implementing their Accessibility functionalities like listening user actions. Undoubtedly, this functionality is so powerful that it is impossible to be ignored by attackers. However, all the implementations have to be engaged under the prerequisite that user needs to open the Accessibility Service in setting options for a specific app. It seems infeasible for attackers to get the grant directly. After all, it performs very suspicious if such a setting window suddenly opens on the

foreground against user's intention.

By exploiting the PEC of such a trusted app, all the attack procedure would become more reasonable. First, the adversary app opens a service on the background to monitor if the PEC service is running through the getRunningServices() offered by the ActivityManager. Fortunately, the invoking of getRunningServices() needs not any permission request and has not been constrained like the getRunningTasks()(mentioned in Section *) so far. Then, the adversary app construct the same logic as trusted app to wait for the time when the available memory size falls below the given threshold. Once it happens, the adversary app would predict the notice dialog will pop up in a few seconds. After that, it shows a toast tip to tell user the corresponding Accessibility Service have to be grant to complete the data clean task, and then opens the Accessibility Service setting page using "Setting.ACTION_ACCESSIBILITY_SETTINGS" intent like Figure *(b) did. In the setting page, an attack service with a confused name "360 Data Clean" makes user fool about this attack. At last, the adversary app gets the grant and could engage further attacks.

Adversary apps also could utilize the callback logic within other apps. This is a more stealthy way. This kind of spoofing has various forms corresponding to the displayed windows of victim apps after a PEC happens. 1. Dialog. A victim system app is designed to give a data notice dialog when the month data rest below 150M.

The display logic is located in the activity life-cycle function OnDestroy() which could be observed by adversary app outside.

4.5.2 Notification Forgery

From perspective of practise, the Notification works for showing the received message, notifying the updated information and representing the progress of running tasks, which is somehow similar as Dialog or Toast. Different from the Dialog, it is not necessary to handle a Notification immediately. Again, a Notification would not disappear in seconds as the Toast does until user handles it. The user could handle it at any time they like, which gives attackers enough time to forge a fake one escaping user's attention. Here we introduce two common used situations.

1. Push mechanism. Most of the market apps currently are based on CS(Client-Server) design, and the server needs

to frequently push the updated data or content to the client in user's device. The client keeps running a service in the background receiving the update from server. Normally, this push-receive mechanism is implemented by specialized module with high-level security protection that makes it hard to get any status from the push process by other apps. However, causing the status of client service for receiving is publicly opened, attacker also could exploit it to fool the user in an indirect way. For example, the popular social app Google+ frequently uses the Notification to notify users to check its updated backup data, with two services `PendingIntentCallbackService` and `DispatchingService` running on the background. The status of these services could be easily observed by a third-party adversary app. Further more, the attacker has already analysed these services control flow offline and found they connect to the Notification display. A fake Notification would be constructed using the same resources and framework as original Google+ did. The difference is the fake Notification are connected to a carefully constructed "Google login" page that needs user input their Google account. When the user open the Notification panel, there are several very similar Notifications listing there. If the user chooses the fake one and tries to log in the Google account, unfortunately, the username and password will be delivered to attacker.

2. Notice. Another important usage of the Notification is to notify users update status from a specific app or the device system. For instance, the QPython (An app for python debugging on Android) would perform a log Notification when the `OnDestroy()` lifecycle function in the `MainActivity` happens. However, since the `OnDestroy()` is invoked when the `MainActivity` is destroyed, a third-party app could use the `getRunningAppProcesses()` (offered by `ActivityManager`) function to judge it. Thus, the `OnDestroy()` belongs to the PEC as well. Attackers could capture the time of such Notification emergence and even the time user triggers the Notification (when the targeted Notification is dismissed and corresponding app starts running.) by analysis from `getActiveNotifications()` offered by `NotificationManager` (after Android 4.3)

Notification forgery. Figure * represents a notification forgery example.

4.5.3 Shaking-display Spoofing

Shaking-display refers to the display (Dialog, Toast, Notification and Activity) led by a shaking action. A typical case comes from the motivation example we introduced in section 3.1**. The shaking action could be captured by the sensor in-set in the device. App developers normally use the `sensorManager` listener and its corresponding callbacks `onSensorChanged()` and `onAccuracyChanged()` to listen and handle the shaking action. Normally, the value of `SensorEvent` (parameter type of sensor callbacks) should be large enough to avoid the misrecognition of unexpected action without user's participating.

The shaking-display spoofing attack is based on the assumption that the shaking action always aims at the foreground activity. Since the shaking action always fits user's intention and in general cases, users only concern the foreground activity, the assumption is reasonable.

Consider the motivating example the YaoYiYao page pops up under two basic conditions. One is the WeChat is running on the foreground, and the other is shaking action event hap-

pens. As the assumption mentioned before, here we don't consider extreme situation that a device shaking happens accidentally in other page like the chatting page or setting page. Fortunately, both of the two conditions could be observed by public even without any of the permission request.

Then let's consider another shaking-display spoofing example with the targeted app named "Digital Vulture". This app implements a novel usage that users could conveniently open the camera only needing slightly shaking the device several times. As analysis before, the handling function for responding the shaking event acts as a PEC. Therefore, an adversary app could capture when the camera is popped up by "Digital Vulture". At that time, a fake camera is launched by the adversary background service and covers on the face of original benign camera. The user is hard to distinguish the two camera surfaces and he/she is very likely to choose the fake one because it is on the foreground. Then adversary could stealthily deliver the picture to the web or store it in a secret location for other usage. The entire attack procedure is hard perceived by user. The only challenge that the adversary app has to overcome is to request the "android.permission.CAMERA" permission. After all, the camera permission is not as sensitive as location or phone status for a user, so it is not hard for an attacker to get this non-sensitive permission.

4.5.4 Activity Phishing

Activity Phishing represents a more direct way to conduct the attack by the PEC. A traditional common used public event is the launching event of targeted app whose `MainActivity` contains private information input, e.g. login page and sign up page. An background adversary service keeps monitoring if the targeted app process begins running. Once the targeted page is launched, a forgery page with the same outlook would cover on the top surface and phish user's private data.

Here, we introduce a more complex PEC example related to implicit Intent. The implicit Intent is known as a bridge to connect another component without specifying a single receiver. Instead, it limits a range of receivers through making a operation statement including the intent data and the intended action. If an app sends an implicit Intent request and there are more than one apps fitting the Intent statement, the Android system would pop up a system dialog listing all the fitting apps. Then the user could choose one of them to handle the Intent data.

There exists a series of iterative PECs threats within the implicit Intent procedure. We still take the Wechat as an example. The Wechat provides a share feature that allows users to conveniently handle the received picture from other chatters. When the user long-presses the picture, a system dialog listing all the apps that can handle this picture pops up. Supposing that the user chooses the QQ, which means the user wants to send this picture to one of his/her QQ friends. As a result, the QQ login page pops up and the user could complete his/her intent after logging in the QQ.

As the analysis before, the running status of foreground app acts as a public event, so that the intent sender ("com.tencent.mm" in this example) could be observed by an adversary app. Actually, the event that pops up the intent receiver QQ is also a public event and could be observed. Moreover, when the system dialog emerges for listing the receivers, the foreground process infers to the system dialog and the observable pro-

cess name changes to "system:ui", which indicates the system dialog emerging is a public event as well. Therefore, there exists a PEC chain that contains sender app process, system dialog process and receiver process in sequence. By capturing the PEC chain, the adversary app could predict the emerging time of QQ login page and cover a phishing login page on the authentic one. Consequently, the adversary app could successfully bypass the unrelated pages(e.g., introduction page) and phish the login user name and password in an accurate way.

4.5.5 Other Attacks

The attacks discussed before are just typical attack types. There are still other attacks could be conducted using the PEC features. For instance, the Notification PEC also could be used to conduct the privilege lifting, only needing a reasonable spoofing page emerging when the user clicks the Notification icon. Again, shaking action could also be used to do phishing if the display is a sensitive login page. In addition, attackers even could try more flagrant attack forms, like ransoming user's critical data and ask for something they are interested for. To achieve this goal, attackers need to design a mal-block and pop it up once the targeted PEC happens. Actually, the attack is designed relying on the PEC type contained by victim app. Table ** list the possible associations of attack and the PEC.

From this perspective, the PEC indeed exists inevitable risk for normal apps.

5. DETECTION AND EVALUATION

This section we first analyse the detection details for the PEC and then provider an evaluation on a large scale of market apps.

5.1 The PEC Detection

To measure the PEC threats in real world apps, we designed a prototype tool CSDroid for detecting and analysing the PEC contained by market apps, which is implemented with 4k lines python code and based on the static analysing framework Androguard[11].

Our analysis includes four steps. First, we match the risk display with the source code of detected apps. We collect four types of risk display: Dialog, Toast, Notification, and Activity, whose function features are listed in Table **.

Then, the CSDroid would recognize the PEC vectors within the source code, including some of lifecycle callbacks of Activity & Service, the callbacks of statically registered BroadcastReceiver and the callbacks defined as observable, e.g., onAccuracyChanged and onSensorChanged.

In the third step, the CSDroid backward traverses feasible method paths from risk displays to the PEC vectors. Generally, the callee methods invoked by the PECs are treated as observable, since these methods would be executed right after the execution of corresponding PECs according to control flow.

However, there are two challenges needed to overcome. First, it is inevitable that some apps contain iterative such method traces. For example, a background service is waiting for a public event and then would launch a new activity. Similarly, in the new activity, the execution of onCreate() lead to a notice dialog popping up. This case exists an iterative method path that the PEC of service first transmits to the activity, and then the activity PEC transmits to the di-

alog. We use redirecting PEC to handle this iterative cases. Considering the joints of an iterative path are always code-block for activity launching(code-block of method startActivity), such as the case introduced above, we model four types of intents used to launch a new activity and flag the target activity name within the intent. After that, in the rest of the method traces, if there is a PEC activity name being the same as the flagged one, the CSDroid associates the two methods, and as a result, the two method traces(one is ended by the startActivity referring to the flagged activity; the other is started by the flagged activity) are integrated as an iterative one.

The second challenge is that not all the extracted PEC method traces are observable for other apps. On the one hand, the adversary app is supposed to be installed without sensitive permission grant. In this way, the PEC path with a system API that needs permission grant is hard to be observed to an adversary app. On the other hand, features involving private resources of the targeted app are also forbidden to observe. For instance, the statical-registered BroadcastReceiver is regarded as a PEC. Yet, an explicit intent that specifies the single receiver's name, received by the objective BroadcastReceiver, could not be observed by other app at all.

The CSDroid analysing such violation situations in a fine-grain way. To solve the permission API problem, considering that the key impact of such a permission grant API is able to confuse the direction of control flow, CSDroid focuses on handling the branch conditions containing such API, with three steps. First, to a method node N whose predecessor method node is denoted as M, CSDroid tells accurate position P in the body of M where N is invoked. Then, from position P, CSDroid conduct a backward traversing to detect whether P potentially depends on a branch condition(we denote such branch condition as B if it exists). If such branch condition exists, CSDroid would recognize whether it depends on a permission grant API through dependency flow analysis. Our permission grant API database comes from the project of PScout[4] which offers a complete mapping list for Android system permission and API. If a permission dependency branch condition is recognized, the CSDroid would remove it from existing method traces.

Compared to the permission API problem, CSDroid just easily build a whitelist to exclude such method traces with any of private resources feature.

handling of inner class

5.2 Market-scale Study

Our analysis is targeted on 2,723 real-world apps crawled by Google Play in 2016, covering all the 27 categories. The final detection results are listed in Table **.

5.3 Vulnerability Discuss

6. MIGRATION STRATEGY

The conditions that enable the attacks successful come from three parts. First, Android system allows the existence of such PECs (mentioned in section **). Second, app developers ignore the PEC problem during their developing process. Third, users have little awareness about such PEC attacks. The violation of any condition would make the attacks fail. Therefore, to migration work should start with above three parts.

System Level. The observation to activity status is known to be prevented from other task stacks in Android 5 and newer versions. It brings substantial benefits for protection of app status from being stolen during runtime. However, the status of components like service, launching activity and broadcastreceiver still can be observed or predicted. Besides that, it is free to access some insensitive-in-surface system resources such as the sensor data. Again, system events can be received by more than one entities at the same time. All above system design makes the PEC attack possible. Therefore, Android system could conduct effective improvement aiming to above features. Compared with changing the event receiving mechanism which involves entire design of Android framework, adding reasonable access to obtain particular resources or observe specific components appears more feasible. However, it is a trade-off between security threats and existing functionalities after all. We expect Android has a concrete improvement for this concern in later versions.

Application Level. From our market survey of the PEC threats, numbers of real-world apps leave the possibility of being attacked by malware through such threats. One clear way to address such problem is to provide a reasonable guideline for developers to avoid this threats during their developing process. Some safety developing habits would be intensively suggested, such as don't set the log-in page in the launching activity, don't link a service to a sensitive activity or dialog, don't pop up a sensitive window right after a system event happens, etc. The developers should also clearly consider related factors based on such tradeoff between security and functionality before developing. However, it seems impossible for a vendor to constrain it's app products as such suggests, which appears overstrict for developers since only a small part of such "unqualified products" suffer from real attacks in practice. To address this challenge, users are expected to be presented a notification tip asking them to whether approve or not when a new display is built after a certain public event. Such notification guides users to be aware of the PEC threats, which is what we strive to in the next stage.

7. RELATED WORK

Android Event-driven Analysis Based on event-driven mechanism, Android application adopts a much flexible execution way, which makes it hard to be directly analysed like the one on conventional personal computer. Researchers has done much efforts on tackling this challenge. [38] fundamentally proposes a program representation built by context-sensitive static analysis of callback methods to model the app's GUI. [20] provides an event-driven analysis manner to measure the power dissipation. [8] centres around so-called Permission Event Graph built with static analysis and uses model checking to illegal interaction between an application and the Android event system. Again, [19] handle event-driven analysis to reach challenging source code for application testing. [26] presents Dynodroid to monitor the reaction of an app upon each event either from system or user. [16] presents a race detection tool named CAFA for event-driven mobile systems. Although prior works present different insights of event-driven analysis in Android system, there are few efforts concerning the security impact raised by public event abuse. Our work fills this gap via systematic study of Android system design and statical analysis for a large scale of

real-world apps.

Android Defects Exploitation The exploration and exploitation of weaknesses within either Android system or apps are intensively focused by current related researchers. The study of memory side-channel analysis and exploitation[41, 9, 18]uncovers an indirect attack way aiming to the weak architecture of Android public data exposure. Several recent researches focus on the vulnerabilities raised by Android advertising libraries[35][34]. Besides that, [39] proposes data residue problem during app uninstallation. As a direct interaction channel with users, the Android UI also suffers from severity security threats, which is carefully discussed in [32][5][1][25][17][33][23]. Again, in the work of [28][37][27], the sensor data could also be regarded as a sort of attack resources predicting some isolated private information including user input, fingerprints. In addition, [2] introduces a denial-of-service attack approach utilizing the AMS design defect. [36] and [6] present detailed studies about privilege escalation for malware through defects of Android os. As to Android permission mechanism, [14][4][40][13] reveal security threats arised by the abuse of permission requests and propose corresponding migration strategy. In contrast to existing work, this paper represents the impact of public event defect of Android that affects a much wider range of system components and services. The attacks relying on such public event callbacks are firstly studied systematically.

Android Vulnerability Detection As a critical security concerning on Android system, privacy data leakage incurs widely related research efforts. Of all the detection works, dynamic taint tracking [12][21][31][30] receives much eyesights for its sound accuracy. Another important complement for dynamic approach relies on static analysis[24][3][15], which can effectively reduce the false negative rate. Further works concern the impacts of ICC(inter-component communication) mechanism[7][29][22]. In addition, [10] measures implicit leak of user password using symbolic execution. Our work is devoted to uncover the public event threats within real market app, which can be easily exploited by attackers to build a wide spectrum of new malwares.

8. CONCLUSIONS

9. ACKNOWLEDGMENTS

10. ADDITIONAL AUTHORS

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APPENDIX

A. HEADINGS IN APPENDICES

A.1 References

B. MORE HELP FOR THE HARDY