

Notation

a	Scalar.
\mathbf{a}	Vector.
\mathbf{A}	Matrix.
\mathbf{A}^{-1}	Inverse of the matrix \mathbf{A} .
\mathbf{A}^\dagger	Pseudo-inverse of the matrix \mathbf{A} .
\mathbf{A}^T	Transpose of a matrix.
\mathbf{A}^*	Complex conjugate of a matrix.
\mathbf{A}^H	Transpose and complex conjugate of a matrix (Hermitian).
$\mathbf{0}$	Vector/matrix of zeros of the appropriate dimensions.
\mathbf{I}	Identity matrix of the appropriate dimensions.
$\text{Tr}(\mathbf{A})$	Trace of the matrix \mathbf{A} .
$ \mathbf{A} $	Determinant of the matrix \mathbf{A} .
$\text{rank}(\mathbf{A})$	Rank of the matrix \mathbf{A} .
$\ker(\mathbf{A})$	Kernel/null-space of the matrix \mathbf{A} .
$\text{blkdiag}(\mathbf{A}_1, \dots, \mathbf{A}_N)$	Block diagonal matrix formed with the matrices $\{\mathbf{A}_1, \dots, \mathbf{A}_N\}$.
$\text{diag}(a_1, \dots, a_N)$	Diagonal matrix whose main diagonal is $\{a_1, \dots, a_N\}$.
$\ \mathbf{a}\ _2$	Norm-2 (Euclidean norm) of the vector \mathbf{a} , see (1.49).
$\nabla_{\mathbf{x}} f(\mathbf{x})$	Gradient of a function $f(\mathbf{x})$, (1.53)
$\mathbb{E}\{\cdot\}$	Statistical expectation.
$\log(\cdot)$	Natural (base e) logarithm.
$\log_2(\cdot)$	Base 2 logarithm.
$[a]^+$	$\max(0, a)$.

Acronyms

AWGN Additive White Gaussian Noise.

BD Block Diagonalization.

BS Base Station.

CDF Cumulative Distribution Function.

iid independent identically distributed.

KKT Karush-Kuhn-Tucker.

MMSE Minimum Mean Squared Error.

PAPC Per Antenna Power Constraint.

PBPC Per Base Station Power Constraint.

pdf probability density function.

SNR Signal to Noise Ratio.

SVD Singular Value Decomposition.

TPC Total Power Constraint.

Chapter 1

System Model

1.1 System Model

The system that will be considered throughout this work aims to represent the downlink of a canonical cellular network, comprising several identical cells, layed out over a regular hexagonal grid.

When studying a cellular network, the cells located at the edge of the network will not experience the same conditions as the cells in the center of the network. A typical way to deal with this situation is to consider a scenario that wraps around (Figure 1.2a) so that cells on one side of the scenario affect cells on the opposite side. Another option is to consider a scenario with more cells than necessary, and then analyze the behavior of the cells located within the center of the network (Figure 1.2b), so that the exterior cells account for the interference, equaling the conditions of all the cells in the network.

Each cell in the system under study will be served with a single Base Station (BS) that is equipped with t transmit antennas. Each of the users considered in the system has r receive antennas. Figure 1.1 shows a schematic representation of such a network, where R_{cell} is the cell radius, and d_{ij} is the distance from the j -th BS to the i -th user.

A system with M BS and N users can then be modelled as

$$\mathbf{y} = \mathbf{H}\mathbf{x} + \mathbf{n} \tag{1.1}$$

 \mathbf{y} represents the signal received at all the users and is defined as

$$\mathbf{y} = \begin{bmatrix} \mathbf{y}_1 \\ \vdots \\ \mathbf{y}_N \end{bmatrix} \in \mathbb{C}^{Nr \times 1} \tag{1.2}$$

where $\mathbf{y}_i \in \mathbb{C}^{r \times 1}$ is the signal received at the i -th user.

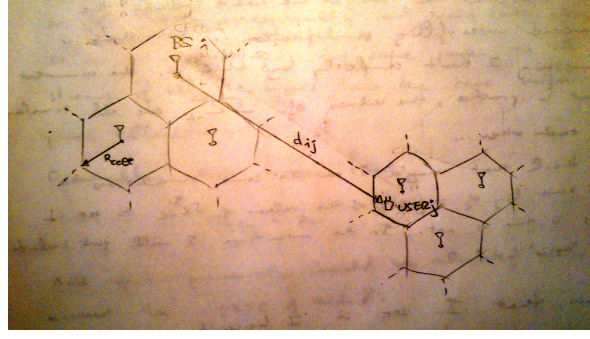
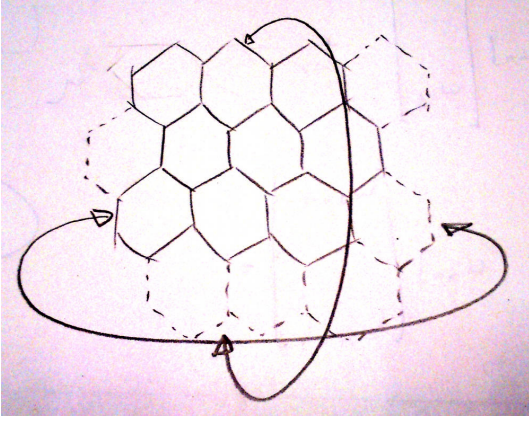
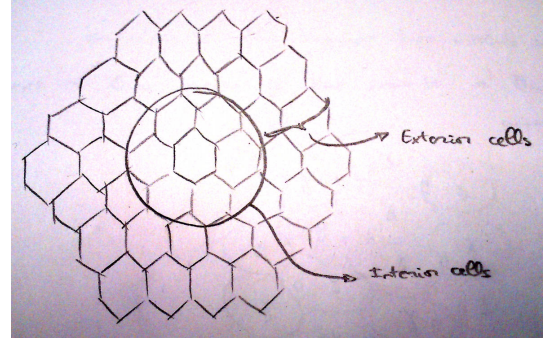


Figure 1.1: Schematic representation of the cellular network that will be used throughout this dissertation.



(a) Wrap around scenario. Cells on one side of the scenario influence the cells on the other side as if they were next to each other.



(b) Oversized scenario. The behavior of the network is analyzed in the central cells of the network, and the exterior cells compensate the network edge effects.

\mathbf{H} is the channel matrix from all the BS to all the users, with the following structure

$$\mathbf{H} = \begin{bmatrix} \mathbf{H}_{11} & \cdots & \mathbf{H}_{1M} \\ \vdots & \ddots & \vdots \\ \mathbf{H}_{N1} & \cdots & \mathbf{H}_{NM} \end{bmatrix} \in \mathbb{C}^{Nr \times Mt} \quad (1.3)$$

where $\mathbf{H}_{ij} \in \mathbb{C}^{r \times t}$ represents the channel matrix from the j -th BS to the i -th user. It will include the path loss due to propagation, small scale fading, shadowing, and any other characteristic of the radio channel that needs to be taken into consideration.

\mathbf{x} is the signal transmitted from all the BS, and it is composed of

$$\mathbf{x} = \begin{bmatrix} x_1 \\ \vdots \\ x_M \end{bmatrix} \in \mathbb{C}^{Mt \times 1} \quad (1.4)$$

where $\mathbf{x}_j \in \mathbb{C}^{t \times 1}$ is the signal transmitted by the j -th BS. Additionally, the power transmitted by the j -th BS can be calculated from the transmitted signal as

$$P_{j,\text{tx}} = \text{Tr}(\mathbf{x}_j \mathbf{x}_j^H) = \mathbf{x}_j^H \mathbf{x}_j \quad (1.5)$$

and each BS will have an independent power constraint

$$P_{j,\text{tx}} \leq P_{j,\text{max}} \quad (1.6)$$

Finally \mathbf{n} represents the Additive White Gaussian Noise (AWGN) at all the receivers

$$\mathbf{n} = \begin{bmatrix} \mathbf{n}_1 \\ \vdots \\ \mathbf{n}_N \end{bmatrix} \in \mathbb{C}^{Nr \times 1} \quad (1.7)$$

with $\mathbf{n}_i \in \mathbb{C}^{r \times 1}$ accounts for the Gaussian noise at the i -th receiver. Throughout this work \mathbf{n}_i is considered to be formed by *iid* entries, drawn from a zero mean, σ_i^2 variance Gaussian distribution, $\mathbf{n}_i \sim \mathcal{N}(\mathbf{0}, \sigma_i^2 \mathbf{I})$. The noise variance will be assumed the same for all the receivers.

In a general scenario, there may be cooperation among the BS in the system so that the information intended for a particular user will be transmitted by several or all the BS. Or, equivalently, each BS transmits a combination of the information of several users

$$\mathbf{x}_j = \mathbf{W}_{j1}^{(\text{tx})} \mathbf{s}_1 + \dots + \mathbf{W}_{jN}^{(\text{tx})} \mathbf{s}_N \quad (1.8)$$

where $\mathbf{s}_i \in \mathbb{C}^{\ell \times 1}$ is the vector of information symbols to be transmitted to user i , with ℓ being the number of simultaneous symbols or streams to be transmitted to that user. $\mathbf{W}_{ji}^{(\text{tx})} \in \mathbb{C}^{t \times \ell}$ is the precoding matrix used at the j -th transmitter for the data of the i -th user.

It will be assumed that the information symbols are independent and drawn from a Gaussian distribution such that $\mathbf{s}_i \sim \mathcal{N}(\mathbf{0}, \mathbf{R}_{\mathbf{s}_i})$, where $\mathbf{R}_{\mathbf{s}_i} = \text{diag}\{p_{i1}, \dots, p_{i\ell}\} \in \mathbb{R}^{\ell \times \ell}$ contains the power allocated to each of the symbols in \mathbf{s}_i . The transmitted power can be expressed as

$$P_{j,\text{tx}} = \sum_{i=1}^N \text{Tr}(\mathbf{W}_{ji}^{(\text{tx})} \mathbf{R}_{\mathbf{s}_i} \mathbf{W}_{ji}^{(\text{tx})H}) \quad (1.9)$$

The choice of the precoding matrices and of the receiving filter will determine the transmission strategy used. This Thesis will focus mainly on Block Diagonalization (BD) [1] which is described in Section 1.3.

On the receiver side, no cooperation among the users will be considered, so each user may perform, independently, additional processing of the received signal by applying a linear filter or equalizer

$$\hat{\mathbf{s}}_i = \mathbf{W}_i^{(\text{rx})} \mathbf{y}_i \quad (1.10)$$

where $\mathbf{W}_i^{(\text{rx})} \in \mathbb{C}^{r \times \ell}$ is the equalizer used at the i -th receiver.

Combining (1.1)–(1.8) it is possible to rewrite (1.1) as

$$\mathbf{y} = \mathbf{H} \mathbf{W}^{(\text{tx})} \mathbf{s} + \mathbf{n} \quad (1.11)$$

where

$$\mathbf{s} = \begin{bmatrix} \mathbf{s}_1 \\ \vdots \\ \mathbf{s}_N \end{bmatrix} \in \mathbb{C}^{N\ell \times 1} \quad (1.12)$$

and the global precoding matrix is

$$\mathbf{W}^{(\text{tx})} = \begin{bmatrix} \mathbf{W}_{11}^{(\text{tx})} & \cdots & \mathbf{W}_{1N}^{(\text{tx})} \\ \vdots & \ddots & \vdots \\ \mathbf{W}_{M1}^{(\text{tx})} & \cdots & \mathbf{W}_{MN}^{(\text{tx})} \end{bmatrix} \in \mathbb{C}^{Mt \times N\ell} \quad (1.13)$$

and it can be partitioned as

$$\mathbf{W}^{(\text{tx})} = [\mathbf{W}_1^{(\text{tx})}, \dots, \mathbf{W}_N^{(\text{tx})}] \quad (1.14)$$

with

$$\mathbf{W}_i^{(\text{tx})} = \begin{bmatrix} \mathbf{W}_{11}^{(\text{tx})} \\ \vdots \\ \mathbf{W}_{M1}^{(\text{tx})} \end{bmatrix} \in \mathbb{C}^{Mt \times \ell} \quad (1.15)$$

The channel matrix can then be partitioned as

$$\mathbf{H} = \begin{bmatrix} \mathbf{H}_1 \\ \vdots \\ \mathbf{H}_N \end{bmatrix} \quad (1.16)$$

where $\mathbf{H}_i \in \mathbb{C}^{r \times Mt}$ is the channel matrix from all the BS to the i -th user.

As it has already been mentioned, receiver cooperation is not going to be considered, so looking at a particular user, *e.g.* the i -th user, the signal that is received will be

$$\mathbf{y}_i = \mathbf{H}_i \mathbf{W}^{(\text{tx})} \mathbf{s} + \mathbf{n}_i \quad (1.17)$$

which can be rewritten as

$$\mathbf{y}_i = \mathbf{H}_i \mathbf{W}_i^{(\text{tx})} \mathbf{s}_i + \underbrace{\sum_{\substack{j=1 \\ j \neq i}}^N \mathbf{H}_i \mathbf{W}_j^{(\text{tx})} \mathbf{s}_j}_{\text{Interference}} + \mathbf{n}_i \quad (1.18)$$

so that it can be readily seen how other users' data appear as an interference term that degrades the received signal.

Defining the term of interference plus noise as

$$\mathbf{z}_i = \sum_{\substack{j=1 \\ j \neq i}}^N \mathbf{H}_i \mathbf{W}_j^{(\text{tx})} \mathbf{s}_j + \mathbf{n}_i \quad (1.19)$$

and using (1.18), the ergodic (mean) rate for the i -th user is given by [2], [3]

$$R_i = \mathbb{E} \left\{ \log_2 \left| \mathbf{I} + \mathbf{H}_i \mathbf{W}_i^{(\text{tx})} \mathbf{R}_{\mathbf{s}_i} \mathbf{W}_i^{(\text{tx})H} \mathbf{H}_i^H \mathbf{R}_{\mathbf{z}_i}^{-1} \right| \right\} \quad (1.20)$$

where $\mathbf{R}_{\mathbf{z}_i} \in \mathbb{C}^{r \times r}$ is the covariance matrix of the noise plus the interference term in (1.18)

$$\mathbf{R}_{\mathbf{z}_i} = \mathbf{z}_i \mathbf{z}_i^H = \left(\sum_{\substack{j=1 \\ j \neq i}}^N \mathbf{H}_i \mathbf{W}_j^{(\text{tx})} + \mathbf{n}_i \right) \left(\sum_{\substack{j=1 \\ j \neq i}}^N \mathbf{H}_i \mathbf{W}_j^{(\text{tx})} + \mathbf{n}_i \right)^H \quad (1.21)$$

In the rest of the work, there are a set of assumptions that will be made, mainly to guarantee the feasibility of some of the results obtained:

- The number of users will be the same as the number of BS, *i.e.* $N = M$.
- The total number of antennas transmitting will be greater or equal than the total number of antennas at the receiver side, this is $Mt \geq Nr$.
- The number of streams transmitted to each user must be $\ell \leq r$, and in general it will be assumed the equality.

- There is no correlation, neither at the transmitters nor at the receivers, so that \mathbf{H} is full rank or, equivalently, $\text{rank}(\mathbf{H}) = \min(Nr, Mt)$. As $Mt \geq Nr$, then $\text{rank}(\mathbf{H}) = Nr$.
- The power available at each of the BS will be assumed the same, *i.e.* $P_{j,\max} = P_{\max}$ for all j .

1.2 Channel Model

As it has been said, each component of \mathbf{H}_{ij} accounts for the propagation path loss, small scale fading, shadowing, and other characteristics of the channel.

In terms of propagation, the channel is typically decomposed as

$$\mathbf{H}_{ij} = \mathbf{R}_{\text{rx},ij} \bar{\mathbf{H}}_{ij} \mathbf{R}_{\text{tx},ij} \quad (1.22)$$

where $\mathbf{R}_{\text{tx},ij} \in \mathbb{C}^{t \times t}$ is the spatial correlation at the transmitter, $\bar{\mathbf{H}} \in \mathbb{C}^{r \times t}$ is the fading channel, and $\mathbf{R}_{\text{rx},ij} \in \mathbb{C}^{r \times r}$ is the spatial correlation at the receiver. In the current work, unless stated otherwise, the channel is considered spatially uncorrelated, both at the transmitter and at the receiver, that is $\mathbf{R}_{\text{rx},ij} = \mathbf{I}$ and $\mathbf{R}_{\text{tx},ij} = \mathbf{I}$, as in [4].



The small scale characteristics considered are Rayleigh, so that the entries of $\bar{\mathbf{H}}_{ij}$ are **independent identically distributed (iid)** complex Gaussian random variables with zero mean and a variance given by the power path loss between the j -th BS and the i -th user.



The attenuation that the signal **experiments** due to the propagation varies according to an exponential power decay with exponent γ , so that the path loss¹ is calculated as

$$\text{pl}_{ij} = \text{pl}_0 \left(\frac{d_{ij}}{d_0} \right)^{-\gamma} \quad (1.24)$$

where pl_0 represents the attenuation at a reference distance d_0 . For the analysis done in this work, but without any loss of generality, pl_0 and d_0 are assumed equal to 1.

1.2.1 Signal to Noise Ratio (SNR)

With the definition of the propagation model that is used in this work, it can also be explained the definition of SNR that is used for the theoretical analyses and the simulations.

¹In natural units, although it is also common to represent the path loss in decibels as

$$\text{PL}_{ij}(\text{dB}) \triangleq 10 \log_{10}(\text{pl}_{ij}) \quad (1.23)$$

Analogously to other works such as [5], the SNR, denoted as ρ , is defined with reference to the power received at the three-way corner of the cell, at a distance of R_{cell} , when the BS transmits at full power, so that the relationship between SNR and the noise power is given by

$$\rho = \frac{P_{\max} R_{\text{cell}}^{-\gamma}}{\sigma_n^2} \quad (1.25)$$

where it is assumed that all the BS have the same maximum transmission power. Using (1.25) it is possible to calculate the noise power for a given SNR and vice versa.

1.3 Block Diagonalization

One possibility to cancel the inter-user interference is to diagonalize the channel matrix. Perfect diagonalization is only possible if $Mt \geq Nr$ [6], and it is achieved using the following precoding matrix

$$\mathbf{W}^{(\text{tx})} = \mathbf{H}^\dagger \quad (1.26)$$

This solution is optimum only when every user has only one antenna. In the case under study, with multiantenna receivers, complete diagonalization of the channel matrix is suboptimal since each user is able to coordinate the processing of its received signal.

In [1] is stated that the optimum solution under the constraint that all inter-user interference be zero is obtained with $\mathbf{H}\mathbf{W}^{(\text{tx})}$ being block diagonal. In [1], BD is proposed as an algorithm to obtain a precoding matrix that is able to block diagonalize the channel matrix. This algorithm is described next.

In order to meet the condition of zero inter-user interference, it is necessary to cancel the interference term in (1.18), and this is equivalent to meet the following

$$\mathbf{H}_i \mathbf{W}_j^{(\text{tx})} = \mathbf{0} \quad \forall j \neq i \quad (1.27)$$

Let $\widetilde{\mathbf{H}}_i$ be the channel matrix \mathbf{H} without \mathbf{H}_i , i.e.

$$\widetilde{\mathbf{H}}_i = \begin{bmatrix} \mathbf{H}_1 \\ \vdots \\ \mathbf{H}_{i-1} \\ \mathbf{H}_{i+1} \\ \vdots \\ \mathbf{H}_N \end{bmatrix} \in \mathbb{C}^{(N-1)r \times Mt} \quad (1.28)$$

Then the condition (1.27) can be obtained making $\mathbf{W}_i^{(\text{tx})}$ lie in the null space or kernel of $\widetilde{\mathbf{H}}_i$. This is possible only if the dimension of the null space is greater than zero, *i.e.* $\text{rank}(\ker(\widetilde{\mathbf{H}}_i)) > 0$.

Now, with the dimensions of $\widetilde{\mathbf{H}}$, the rank of its null space is

$$\text{rank}(\ker(\widetilde{\mathbf{H}}_i)) = Mt - \text{rank}(\widetilde{\mathbf{H}}_i) \quad (1.29)$$

But it is assumed that \mathbf{H} is full rank, ergo $\text{rank}(\widetilde{\mathbf{H}}_i) = (N-1)r = \widetilde{L}_i$, and then

$$\text{rank}(\ker(\widetilde{\mathbf{H}}_i)) = Mt - \widetilde{L}_i > 0 \quad (1.30)$$

so it is guaranteed that a precoding matrix $\mathbf{W}_i^{(\text{tx})}$ that lies in the null space of $\widetilde{\mathbf{H}}_i$ exists.

The simplest way to obtain such $\mathbf{W}_i^{(\text{tx})}$ involves using the Singular Value Decomposition (SVD) of the matrix $\widetilde{\mathbf{H}}_i$.

Let $\widetilde{\mathbf{H}}_i$ be decomposed as

$$\widetilde{\mathbf{H}}_i = \widetilde{\mathbf{U}}_i \widetilde{\mathbf{A}}_i [\widetilde{\mathbf{V}}_i^{(1)}, \widetilde{\mathbf{V}}_i^{(0)}]^H \quad (1.31)$$

where $\widetilde{\mathbf{V}}_i^{(0)} \in \mathbb{C}^{Mt \times (Mt - \widetilde{L}_i)}$ contains the last $Mt - \widetilde{L}_i$ right singular vectors of $\widetilde{\mathbf{H}}_i$, corresponding to the singular values equal to zero. $\widetilde{\mathbf{V}}_i^{(0)}$ forms an orthonormal basis of the null space of $\widetilde{\mathbf{H}}_i$, and thus its columns can be used to cancel the inter-user interference

$$\widetilde{\mathbf{H}}_i \widetilde{\mathbf{V}}_i^{(0)} = \mathbf{0} \quad (1.32)$$

Using these matrices as precoding the result is

$$\widehat{\mathbf{H}} = \mathbf{H} [\widetilde{\mathbf{V}}_1^{(0)}, \dots, \widetilde{\mathbf{V}}_N^{(0)}] = \begin{bmatrix} \mathbf{H}_1 \widetilde{\mathbf{V}}_1^{(0)} & & \mathbf{0} \\ & \ddots & \\ \mathbf{0} & & \mathbf{H}_N \widetilde{\mathbf{V}}_N^{(0)} \end{bmatrix} \quad (1.33)$$

that, as it can be seen, has a block diagonal structure, which gives the name to the algorithm proposed in [1].

The next problem that BD solves is the maximization of the sum rate of the system, given the block diagonal structure in (1.33). The precoding matrix $\mathbf{W}_i^{(\text{tx})}$ will be considered to be

$$\mathbf{W}_i^{(\text{tx})} = \widetilde{\mathbf{V}}_i^{(0)} \mathbf{W}_i' \quad (1.34)$$

where $\mathbf{W}'_i \in \mathbb{C}^{(Mt-\widetilde{L}_i) \times \ell}$ will take care of the rate maximization.

Introducing (1.34) into (1.20) the ergodic capacity simplifies to

$$R_i^{\text{no interf}} = \mathbb{E} \left\{ \log_2 \left| \mathbf{I} + \frac{1}{\sigma_i^2} \widehat{\mathbf{H}}_i \mathbf{W}'_i \mathbf{R}_{s_i} \mathbf{W}'_i{}^H \widehat{\mathbf{H}}_i^H \right| \right\} \quad (1.35)$$

where $\widehat{\mathbf{H}}_i = \mathbf{H}_i \widetilde{\mathbf{V}}_i^{(0)} \in \mathbb{C}^{r \times (Mt-\widetilde{L}_i)}$.

In order to maximize the rate, consider the SVD

$$\widehat{\mathbf{H}}_i = \widehat{\mathbf{U}}_i \begin{bmatrix} \widehat{\mathbf{\Lambda}}_i & \mathbf{0} \\ \mathbf{0} & \mathbf{0} \end{bmatrix} [\widehat{\mathbf{V}}_i^{(1)}, \widehat{\mathbf{V}}_i^{(0)}]^H \quad (1.36)$$

where $\widehat{\mathbf{\Lambda}}_i = \text{diag} \{ \hat{\lambda}_{i1}^{1/2}, \dots, \hat{\lambda}_{ir}^{1/2} \} \in \mathbb{C}^{r \times r}$ contains the non-zero singular values of $\widehat{\mathbf{H}}_i$, which has $\text{rank}(\widehat{\mathbf{H}}_i) = r$. And $\widehat{\mathbf{V}}_i^{(1)} \in \mathbb{C}^{(Mt-\widetilde{L}_i) \times r}$ contains the first r right singular vectors of $\widehat{\mathbf{H}}_i$, and it will be used as \mathbf{W}'_i , yielding the following precoding matrix

$$\mathbf{W}_i^{(\text{tx})} = \widetilde{\mathbf{V}}_i^{(0)} \widehat{\mathbf{V}}_i^{(1)} \quad (1.37)$$

The BD also provides the receiver filter to be used at each user which will be

$$\mathbf{W}_i^{(\text{rx})} = \widehat{\mathbf{U}}_i^H \quad (1.38)$$

Using all of the above, the rate that the i -th user can obtain is given by the expression

$$R_i^{\text{BD}} = \mathbb{E} \left\{ \sum_{k=1}^{\ell} \log_2 \left(1 + \frac{\hat{\lambda}_{ik} p_{ik}}{\sigma_i^2} \right) \right\} \quad (1.39)$$

where the only parameters left to be computed are the power allocated to each of the ℓ streams of each user, and different options to do it will be discussed in Section 1.4.

1.4 Power Allocation

In Section 1.3, the BD algorithm has been described to get the precoding matrix to be used at the transmitter and the equalization filter to be used at the receiver side. After BD has been used, the power should be allocated to each of the data streams of each user, this is, the p_{ij} in (1.39) should be calculated in order to achieve a given performance, and subject to particular constraints.

The need for a power allocation algorithm comes from the restriction on the maximum power available for transmission, which may be due to physical limitations, or regulatory issues.

There are several different power constraints:

- **Per Antenna Power Constraint (PAPC):** The maximum power is constrained for each antenna at the transmitter. This option is specially well suited for distributed antenna systems [7], [8].
- **Per Base Station Power Constraint (PBPC):** In this case the maximum power is limited per base station instead of per antenna. This option is more appropriate for scenarios where all the transmitting antennas are collocated and may share a power budget, so that the transmission power can be arbitrarily allocated to each of the transmitter antennas.

A Total Power Constraint (TPC) can also be considered, which assumes that the maximum power is shared among all the transmitters in the system. Although this system is more easily analyzed, it is very unrealistic so it will not be considered in this work, except in Subsection 1.4.3 where a TPC is used to obtain an intermediate result.

For the sake of simplicity, PBPC will be used for the different analyses. In any case, PAPC can be seen as a particularization of PBPC as the derivations shown in this section can be applied to a PAPC system considering instead of each BS to have t transmit antennas, t single antenna BS.

The problem that needs to be solved is, in general, the maximization of some function of the rate of each of the users subject to a PBPC

$$\begin{aligned} & \underset{\{\mathbf{R}_{\mathbf{s}_i}\}}{\text{maximize}} && f\left(R_i^{\text{BD}}(\mathbf{R}_{\mathbf{s}_1}), \dots, R_N^{\text{BD}}(\mathbf{R}_{\mathbf{s}_N})\right) \\ & \text{subject to} && P_{j,\text{tx}} \leq P_{j,\text{max}}, \quad j = 1, \dots, M \end{aligned} \tag{1.40}$$

One common metric used for the maximization is the weighted sum-rate of the system, so that the function $f(\cdot)$ is equal to

$$f(R_i^{\text{BD}}, \dots, R_N^{\text{BD}}) = \sum_{i=1}^N \alpha_i R_i^{\text{BD}} \tag{1.41}$$

where $\alpha_i \in [0, 1]$ can be seen as different priorities for different users, and they are assumed to be

$$\sum_{i=1}^N \alpha_i = 1 \tag{1.42}$$

and in the particular case where all the $\alpha_i = 1/N$, then the function $f(\cdot)$ represents the sum-rate of the system.

Calling

$$\overline{\mathbf{W}}_j^{(\text{tx})} = [\mathbf{W}_{j1}^{(\text{tx})}, \dots, \mathbf{W}_{jN}^{(\text{tx})}] \in \mathbb{C}^{t \times N\ell} \quad (1.43)$$

the precoding matrix of the j -th BS, and

$$\mathbf{R}_s = \text{blkdiag}(\mathbf{R}_{s_1}, \dots, \mathbf{R}_{s_N}) \in \mathbb{C}^{N\ell \times N\ell} \quad (1.44)$$

the matrix containing the power assigned to all the streams of all the users, the power constraint in (1.40) can then be reformulated as

$$\text{Tr}(\overline{\mathbf{W}}_j^{(\text{tx})} \mathbf{R}_s \overline{\mathbf{W}}_j^{(\text{tx}),H}) \leq P_{j,\max} \quad (1.45)$$

Now the term inside the trace operator can be written explicitly as a function of p_{ik} in order to make it easier to analyze. First define

$$\overline{\mathbf{W}}_j^{(\text{tx})} = [\bar{\mathbf{w}}_{j,11}, \dots, \bar{\mathbf{w}}_{j,1\ell}, \dots, \bar{\mathbf{w}}_{j,N\ell}] \quad (1.46)$$

where $\bar{\mathbf{w}}_{j,ik} \in \mathbb{C}^{t \times 1}$ is the ik -th column of $\overline{\mathbf{W}}_j^{(\text{tx})}$, *i.e.* the precoding that is used at the j -th BS for the k -th stream of the i -th user. And then:

$$\begin{aligned} \overline{\mathbf{W}}_j^{(\text{tx})} \mathbf{R}_s \overline{\mathbf{W}}_j^{(\text{tx}),H} &= [\bar{\mathbf{w}}_{j,11}, \dots, \bar{\mathbf{w}}_{j,N\ell}] \begin{bmatrix} p_{11} & & 0 \\ & \ddots & \\ 0 & & p_{N\ell} \end{bmatrix} \begin{bmatrix} \bar{\mathbf{w}}_{j,11}^H \\ \vdots \\ \bar{\mathbf{w}}_{j,N\ell}^H \end{bmatrix} = \\ &= p_{11} \bar{\mathbf{w}}_{j,11} \bar{\mathbf{w}}_{j,11}^H + \dots + p_{N\ell} \bar{\mathbf{w}}_{j,N\ell} \bar{\mathbf{w}}_{j,N\ell}^H \\ &= \sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik} \bar{\mathbf{w}}_{j,ik} \bar{\mathbf{w}}_{j,ik}^H \end{aligned} \quad (1.47)$$

and the trace is

$$\text{Tr}(\overline{\mathbf{W}}_j^{(\text{tx})} \mathbf{R}_s \overline{\mathbf{W}}_j^{(\text{tx}),H}) = \sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik} \|\bar{\mathbf{w}}_{j,ik}\|_2^2 \quad (1.48)$$

where

$$\|\bar{\mathbf{w}}_{j,ik}\|_2^2 = \text{Tr}(\bar{\mathbf{w}}_{j,ik} \bar{\mathbf{w}}_{j,ik}^H) = \bar{\mathbf{w}}_{j,ik}^H \bar{\mathbf{w}}_{j,ik}. \quad (1.49)$$

The sum-rate maximization problem can then be formulated in *standard form* [9] as

$$\begin{aligned}
& \underset{p_{ik}}{\text{minimize}} && - \sum_{i=1}^N \sum_{k=1}^{\ell} \log_2 \left(1 + \frac{\hat{\lambda}_{ik} p_{ik}}{\sigma_i^2} \right) \\
& \text{subject to} && \sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik} \|\bar{\mathbf{w}}_{j,ik}\|_2^2 - P_{j,\max} \leq 0, \quad j = 1, \dots, M \\
& && -p_{ik} \leq 0, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array}
\end{aligned} \tag{1.50}$$

In the next sections, different alternatives for obtaining these powers are presented and described.

1.4.1 Optimal Power Allocation

In (1.50), the function $\log_2(\cdot)$ is convex on p_{ik} and the sum of convex functions is also convex, so the objective function in (1.50) is convex. The constraints are affine and therefore convex too. The optimization problem in (1.50) is a convex optimization problem that can be solved using a myriad of numerical technics [9]. Nonetheless, it would be interesting to analyze a bit further the problem in order to get some insight about it.

The problem in (1.50) satisfies *Slater's condition* [9] since the objective function is convex and all the inequality constraints are affine, hence *strong duality* holds. This means that the optimum value of the primal problem is equal to the optimum value of the *Lagrange dual problem*, so that this can be used to find out the solution to the primal, original, problem.

Under these conditions, and considering that the objective function is differentiable with respect to p_{ik} , Karush-Kuhn-Tucker (KKT) conditions [9] are necessary and sufficient for optimality of a solution, and they can be used to analyze the optimization problem in search for an optimal solution. The KKT conditions for (1.50) are

$$\begin{aligned}
\sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik}^* \|\bar{\mathbf{w}}_{j,ik}\|_2^2 - P_{j,\max} &\leq 0, & j = 1, \dots, M \\
p_{ik}^* &\leq 0, & \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \\
\nu_j^* &\geq 0, & j = 1, \dots, M \\
\mu_{ik}^* &\geq 0, & \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \\
\nu_j^* \left(\sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik} \|\bar{\mathbf{w}}_{j,ik}\|_2^2 - P_{j,\max} \right) &= 0, & j = 1, \dots, M \\
-\mu_{ik}^* p_{ik}^* &= 0, & \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \\
\nabla_{\mathbf{p}} \mathcal{L}(\mathbf{p}^*, \boldsymbol{\nu}^*, \boldsymbol{\mu}^*) &= \mathbf{0}
\end{aligned} \tag{1.51}$$

where the superscript $*$ represents a feasible solution of the optimization problem, $\nabla_{\mathbf{p}}$ is the gradient with respect to the powers $\mathbf{p} = [p_{11}, \dots, p_{N\ell}]^T$, $\boldsymbol{\nu} = [\nu_1, \dots, \nu_M]^T$ and $\boldsymbol{\mu} = [\mu_{11}, \dots, \mu_{N\ell}]^T$ are the Lagrange multipliers, and \mathcal{L} represents the *Lagrangian* associated with the problem (1.50), and it is defined as

$$\begin{aligned} \mathcal{L}(\mathbf{p}, \boldsymbol{\nu}, \boldsymbol{\mu}) = & - \sum_{i=1}^N \sum_{k=1}^{\ell} \log_2 \left(1 + \frac{\hat{\lambda}_{ik} p_{ik}}{\sigma_i^2} \right) + \\ & \sum_{j=1}^M \nu_j \left(\sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik} \|\bar{\mathbf{w}}_{j,ik}\|_2^2 - P_{j,\max} \right) - \\ & \sum_{i=1}^N \sum_{k=1}^{\ell} \mu_{ik} p_{ik} \end{aligned} \quad (1.52)$$

The gradient of a function $f(\mathbf{x})$ with respect to $\mathbf{x} \in \mathbb{C}^{n \times 1}$ is defined as

$$\nabla_{\mathbf{x}} f(\mathbf{x}) = \begin{bmatrix} \frac{\partial}{\partial x_1} f(\mathbf{x}) \\ \vdots \\ \frac{\partial}{\partial x_n} f(\mathbf{x}) \end{bmatrix} \quad (1.53)$$

First the gradient of the objective function is calculated, by computing the partial derivatives

$$\frac{\partial}{\partial p_{ik}} \left\{ - \sum_{i=1}^N \sum_{k=1}^{\ell} \log_2 \left(1 + \frac{\hat{\lambda}_{ik} p_{ik}}{\sigma_i^2} \right) \right\} = \frac{-\hat{\lambda}_{ik}}{\log(2) (\sigma_i^2 + \hat{\lambda}_{ik} p_{ik})} \quad (1.54)$$

And the same for the inequality constraints

$$\begin{aligned} \frac{\partial}{\partial p_{ik}} \left\{ \sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik} \|\bar{\mathbf{w}}_{j,ik}\|_2^2 - P_{j,\max} \right\} &= \|\bar{\mathbf{w}}_{j,ik}\|_2^2 \\ \frac{\partial}{\partial p_{ik}} \{p_{ik}\} &= 1 \end{aligned} \quad (1.55)$$

So that the condition of the gradient of the Lagrangian vanishing, in (1.50) can be written as

$$\frac{-\hat{\lambda}_{ik}}{\log(2) (\sigma_i^2 + \hat{\lambda}_{ik} p_{ik}^*)} + \sum_{j=1}^M \nu_j^* \|\bar{\mathbf{w}}_{j,ik}\|_2^2 - \mu_{ik}^* = 0, \quad \begin{matrix} i = 1, \dots, N \\ k = 1, \dots, \ell \end{matrix} \quad (1.56)$$

It can be seen that μ_{ik} is a slack variable that takes into account the non-negativeness of the powers p_{ik} , and it can be omitted to get the equation

$$\frac{-\hat{\lambda}_{ik}}{\log(2)(\sigma_i^2 + \hat{\lambda}_{ik}p_{ik}^*)} + \sum_{j=1}^M \nu_j^* \|\bar{\mathbf{w}}_{j,ik}\|_2^2 \geq 0, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \quad (1.57)$$

Calling

$$L_{ik} = \sum_{j=1}^M \nu_j^* \|\bar{\mathbf{w}}_{j,ik}\|_2^2 \quad (1.58)$$

(1.57) can be solved for p_{ik}^*

$$p_{ik}^* \leq \frac{1}{\log(2)L_{ik}} - \frac{\sigma_i^2}{\hat{\lambda}_{ik}}, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \quad (1.59)$$

The result in (1.59) resembles the classical *water-filling* solution, except that now the water level is not fixed, and it depends on the precoders. The coupling existing among the power constraints of the different BS makes it impossible to find a closed-form solution for the values of p_{ik} .

Nevertheless, this analysis motivates the development of suboptimal schemes that are described in the following sections.

1.4.2 Modified Water-Filling

[10] proposes a simplification to the original problem, in order to make it more tractable. The coupling of the power constraints in (1.50) makes it impossible to get a simple solution for the optimal power allocation problem. [10] approaches the problem by first considering an equivalent virtual BS so that the problem is cast with a single power constraint.

In order to do so, instead of having a power constraint for each of the BS consider a single power constraint given by the most restrictive BS in the original problem. Define

$$\Omega_{ik}^{\text{BS}} = \max_{j=1, \dots, M} \|\bar{\mathbf{w}}_{j,ik}\|_2^2 \quad (1.60)$$

as the weights of the single virtual BS corresponding to each of the users' streams. The optimization problem becomes then

$$\begin{aligned}
& \underset{\mathbf{p}_{ik}}{\text{minimize}} && - \sum_{i=1}^N \sum_{k=1}^{\ell} \log_2 \left(1 + \frac{\hat{\lambda}_{ik} p_{ik}}{\sigma_i^2} \right) \\
& \text{subject to} && \sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik} \Omega_{ik}^{\text{BS}} - P_{\text{BS}, \max} \leq 0 \\
& && -p_{ik} \leq 0, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array}
\end{aligned} \tag{1.61}$$

where $P_{\text{BS}, \max}$ represents the most restrictive power constraint among all of the BS.

This problem meets the same conditions as the original problem so that a similar analysis can be used. First formulate the Lagrangian of the new problem as

$$\begin{aligned}
\mathcal{L}(\mathbf{p}, \nu, \boldsymbol{\mu}) = & - \sum_{i=1}^N \sum_{k=1}^{\ell} \log_2 \left(1 + \frac{\hat{\lambda}_{ik} p_{ik}}{\sigma_i^2} \right) + \\
& \nu \left(\sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik} \Omega_{ik}^{\text{BS}} - P_{\text{BS}, \max} \right) - \\
& \sum_{i=1}^N \sum_{k=1}^{\ell} \mu_{ik} p_{ik}
\end{aligned} \tag{1.62}$$

And its gradient is given by

$$\nabla_{\mathbf{p}} \mathcal{L}(\mathbf{p}^*, \nu^*, \boldsymbol{\mu}^*) = \frac{-\hat{\lambda}_{ik}}{\log(2) (\sigma_i^2 + \hat{\lambda}_{ik} p_{ik}^*)} + \nu^* \Omega_{ik}^{\text{BS}} - \mu_{ik}^* = 0, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \tag{1.63}$$

Using the KKT condition that the gradient of the Lagrangian should vanish, and considering μ_{ik}^* a slack variable, and solving for p_{ik}^* , the following inequality is obtained

$$p_{ik}^* \leq \frac{1}{\log(2) \nu^* \Omega_{ik}^{\text{BS}}} - \frac{\sigma_i^2}{\hat{\lambda}_{ik}}, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \tag{1.64}$$

which, together with the constraint of the powers being non-negative, can be written as

$$p_{ik}^* = \left[\frac{1}{\log(2) \nu^* \Omega_{ik}^{\text{BS}}} - \frac{\sigma_i^2}{\hat{\lambda}_{ik}} \right]^+, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \tag{1.65}$$

The solution of this simplified problem is given by the water-filling solution with a variable water level and, in this case, an uncoupled solution for each of the data streams for each user. This allows for the use of standard and efficient methods to find the power allocation [11].

Clearly, the definition of the new problem makes it more restrictive than the original, and its solution will be also a feasible solution for the original problem, albeit not the optimal. The results in [10] show how under some conditions, the solution achieved like this can be rather close to the optimum one.

Standard Water-Filling

One further simplification that is done in [10] is to consider that, in practice, the values of all the Ω_{ik}^{BS} are very similar to each other, so that it is possible to consider them equal. This turns the problem into a standard water-filling problem, with constant water level, where the solution is given by

$$p_{ik}^* = \left[K_{\text{WF}} - \frac{\sigma_i^2}{\hat{\lambda}_{ik}} \right]^+, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \quad (1.66)$$

where $K_{\text{WF}} = \frac{1}{\log(2)\nu^*\Omega_{\text{WF}}}$, is constant, as $\Omega_{ik}^{\text{BS}} = \Omega_{\text{WF}} \forall i = 1, \dots, N; \forall k = 1, \dots, \ell$ is the same for all the data streams of all the users.

1.4.3 Scaled Water-Filling

In [5] the same power allocation problem as in (1.50) is dealt with by considering a TPC, so that the optimization problem becomes

$$\begin{aligned} \underset{p_{ik}}{\text{minimize}} \quad & - \sum_{i=1}^N \sum_{k=1}^{\ell} \log_2 \left(1 + \frac{\hat{\lambda}_{ik} p_{ik}}{\sigma_i^2} \right) \\ \text{subject to} \quad & \text{Tr}(\mathbf{W}^{(\text{tx})} \mathbf{R}_{\mathbf{s}} \mathbf{W}^{(\text{tx}, H)}) - M P_{\max} \leq 0 \\ & -p_{ik} \leq 0, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \end{aligned} \quad (1.67)$$

where it has been assumed that $P_{j,\max} = P_{\max} \forall j$.

Under this TPC, the solution is readily derived by water-filling [11], but the resulting $\mathbf{R}_{\mathbf{s}}^{\text{TPC}}$ may violate the individual power constraints of each BS.

In order to meet each PBPC, the matrix $\mathbf{R}_{\mathbf{s}}^{\text{TPC}}$ must be scaled so that the final power allocation is given by

$$\mathbf{R}_{\mathbf{s}}^{\text{SWF}} = \beta \mathbf{R}_{\mathbf{s}}^{\text{TPC}} \quad (1.68)$$

where the scaling factor $\beta \in (0, 1)$ is calculated as

$$\beta = \frac{P_{\max}}{\max_{j=1,\dots,M} \text{Tr} \left(\mathbf{W}_j^{(\text{tx})} \mathbf{R}_{\mathbf{s}}^{\text{TPC}} \mathbf{W}_j^{(\text{tx}),H} \right)} \quad (1.69)$$

The results in [5] show, as well, that this simplified approach can deliver near-optimum performance.

1.4.4 Uniform Power allocation

The simplest, both conceptually and computationally, alternative that can be considered to solve the power allocation in (1.50) consists on considering a uniform power allocation.

This approach assigns the same power to all the data streams of all the users. Formally this means

$$p_{ik} = p_s, \quad \begin{array}{l} i = 1, \dots, N \\ k = 1, \dots, \ell \end{array} \quad (1.70)$$

where the power p_s should be computed taking into account the PBPC for each of the BS.

Recall from (1.45) the power transmitted by the j -th BS, where now the matrix $\mathbf{R}_{\mathbf{s}}$ is given by

$$\mathbf{R}_{\mathbf{s}} = p_s \mathbf{I} \quad (1.71)$$

and (1.45) becomes

$$p_s \text{Tr} \left(\overline{\mathbf{W}}_j^{(\text{tx})} \overline{\mathbf{W}}_j^{(\text{tx}),H} \right) \leq P_{j,\max}, \quad j = 1, \dots, M \quad (1.72)$$

The new power allocation problem can be formulated as

$$\begin{array}{ll} \underset{p}{\text{maximize}} & p \\ \text{subject to} & p \text{Tr} \left(\overline{\mathbf{W}}_j^{(\text{tx})} \overline{\mathbf{W}}_j^{(\text{tx}),H} \right) \leq P_{j,\max}, \quad j = 1, \dots, M \end{array} \quad (1.73)$$

which is a *linear programming* optimization problem, and it can be solved efficiently using classical methods, *e.g.* bisection method [12].

Chapter 2

Mean Achievable Rates in Clustered Coordinated Base Station Transmission with Block Diagonalization¹

2.1 Introduction

In [\refc{ch:state_art}](#) it has been discussed how global coordination in a cellular network is not feasible for practical application. Research has shown that grouping cells in clusters may help to alleviate some of the problems of global coordination. The clustering solution is not unique, though, and multiple alternatives and numerous parameters have to be chosen in order to meet different objectives.

The objective of this chapter is to analyze the mean achievable rate in a cellular network where clusters have been formed, and where BD is used within each cluster to manage the interference.

With this analysis, further research can be made in order to be able to choose from one of the most important parameters in clustering, the cluster size.

Numerical simulations are also performed in order to validate the theoretical analysis developed.

2.2 Clustered Network

The network to be considered is organized as independent groups of M cells, each of which is a cluster where its cells coordinate in order to transmit to the N users that are located within

¹The work shown in this chapter has been published in [13].

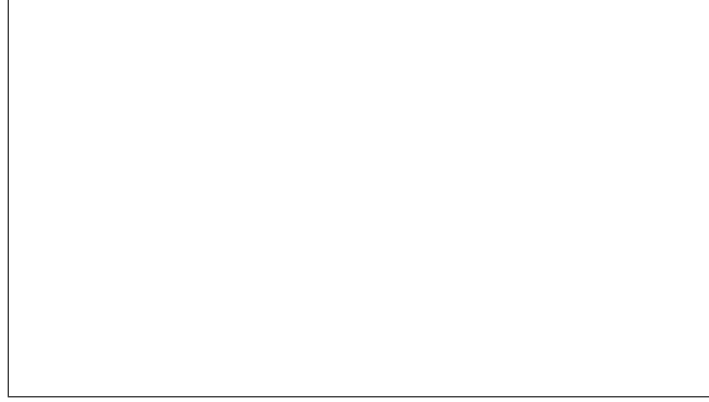


Figure 2.1: System layout with clusters of seven cells of radius R_{cell} (the radius of the circle circumscribing the cell) with an example of two users, with the distances $D_{\text{tier } 1}$ and $D_{\text{tier } 2}$, respectively, from their BS to the interfering BS.

the cluster.

The clusters that form the network are non-overlapping, *i.e.* each cell in the system belongs to one and only one cluster. Overlapping, user-centric clusters have been shown to provide, in some situations, better performance than disjoint clusters [14], but this approach gives rise to a dramatic increase of the management complexity.

The clusters, then, are defined by the network planner, and they are kept fixed, grouping the BS according to a distance criterion, so that the cells belonging to a cluster must form, borrowing the name from graph theory, a *connected component* of the graph containing all the cells in the system.

In the setup under study, all the BS in the cluster are considered cluster members, that is, no scheduling or adaptive selection of active BS is addressed in this work. In any case, they could be considered as a special case of the optimization problem to obtain the power allocation scheme.

Figure 2.1 shows such a clustered network, where the disjoint clusters can be observed, together with other parameters that will be used in considering the interference for the analysis.

2.3 Interference Model

In Figure 2.1 an example of a network is shown, with three complete clusters of $M = 7$ cells each, and with some cells belonging to clusters not completely shown.

The user i , at a distance $d_i \triangleq d_{ii}$ from its serving BS, is affected by the interference originated in the neighboring clusters. In this case, the closest interfering cells are located at a distance $D_{\text{tier } 1} = \sqrt{3}R_{\text{cell}}$ from the i -th BS.



Similarly, for the user j , at a distance $d_j \triangleq d_{jj}$ from its serving BS, the nearest interfering cells are located at a distance $D_{\text{tier } 2} = 3R_{\text{cell}}$ from the j -th BS.

Due to the cellular geometry, for a cluster size of up to 18, only these two possibilities exist: the closest interfering cell is at a distance of either $D_{\text{tier } 1}$ or $D_{\text{tier } 2}$ from the serving BS of each user.

The hexagonal cell can be approximated by a circular one with radius R_{cell} , see Figure 2.1, and then, assuming a uniform distribution of the users over each cell, the probability density function (pdf) of the distance of a user to its serving (closest) BS is given by

$$f_{d_i}(d_i) = \frac{2d_i}{R_{\text{cell}}^2} \quad (2.1)$$

The interference power received at the user i is equal to

$$I_i(d_i) = \sum_{m=1}^{M_{\text{interf}}} P_{\text{max}} \hat{d}_{im}^{-\gamma} \quad (2.2)$$

where M_{interf} is the number of interfering BS, *i.e.* the total number of cells in the system minus the M cells that form the cluster, and \hat{d}_{im} is the distance from the m -th BS outside the cluster to the i -th user in the cluster. All the interfering BS are assumed to be transmitting at full power.

In order to simplify the computation of the interference power, an equivalent model is introduced. In this model, the interference comes from $M_{\text{eq},i}$ cells, all of which are located at the same distance $(D_i - d_i)$ from the i -th user, the one being interfered.

The distance D_i takes the value of the distance from the serving BS to the closest interfering BS

$$D_i \in \{D_{\text{tier } 1}, D_{\text{tier } 2}\} \quad (2.3)$$

The equivalent number of interfering BS, $M_{\text{eq},i}$, is such that the total interference power is the same as in the original layout.

With all this, (2.2) can be written as

$$I_i(d_i) = P_{\text{max}} M_{\text{eq},i} (D_i - d_i)^{-\gamma} \quad (2.4)$$

This approach is similar to that followed in [15], where a fluid model network is used. This model assumes that there is a continuum of BS interfering, but this will not be considered in the current work for the sake of simplicity.

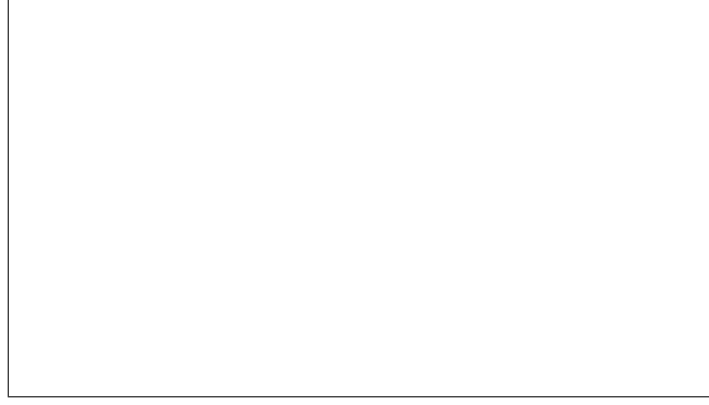


Figure 2.2: Cluster with $M = 4$ cells and neighbor interfering cells.

The real and equivalent model produce the same total interference, provided that $M_{\text{eq},i}$ is adequately selected.

In order to determine $M_{\text{eq},i}$ the only interference that is accounted for is the one coming from the first tier of neighboring cells. This implies that different cluster configurations may have different number of interfering BS for each of the cells in the cluster, and this number for the i -th cell is denoted as $M_{\text{int},i} \leq M_{\text{interf}}$.

This is made clear in Figure 2.2 where a cluster with $M = 4$ cells is surrounded by $M_{\text{interf}} = 12$ cells. In this cluster, cells 1 and 3 experience an interference coming from $M_{\text{int},1} = M_{\text{int},3} = 4$ neighboring cells, while cells 2 and 4 receive the interference from $M_{\text{int},2} = M_{\text{int},4} = 3$ cells.

In order to approximate the value of equivalent interfering BS for a general network setup, first consider a simple scenario Figure 2.3 with a cluster of $M = 1$ cells, where a single user, $i = 1$, in the cell is affected by an interference power I_1 coming from all the $M_{\text{int},1} = 6$ belonging to the first tier.

An assumption that can be made is that half of the $M_{\text{int},1}$, *i.e.* three, BS are located at a distance $(D_1 - d_1)$ and the other half are located at a distance $(D_1 + d_1)$. Accordingly, the interference power can be expressed as

$$I_1(d_1) = P_{\max} \left[\frac{3}{(D_1 - d_1)^\gamma} + \frac{3}{(D_1 + d_1)^\gamma} \right] \quad (2.5)$$

which is equal to (2.4) if the equivalent number of interfering BS is defined as

$$M_{\text{eq},1} = 3 \left[1 + \left(\frac{D_1 - d_1}{D_1 + d_1} \right)^\gamma \right] \quad (2.6)$$

In general, if a user i is considered to have $M_{\text{int},i}$ interfering BS in the first tier at a distance D_i , defined as in (2.3), then the equivalent number of interfering BS is given by

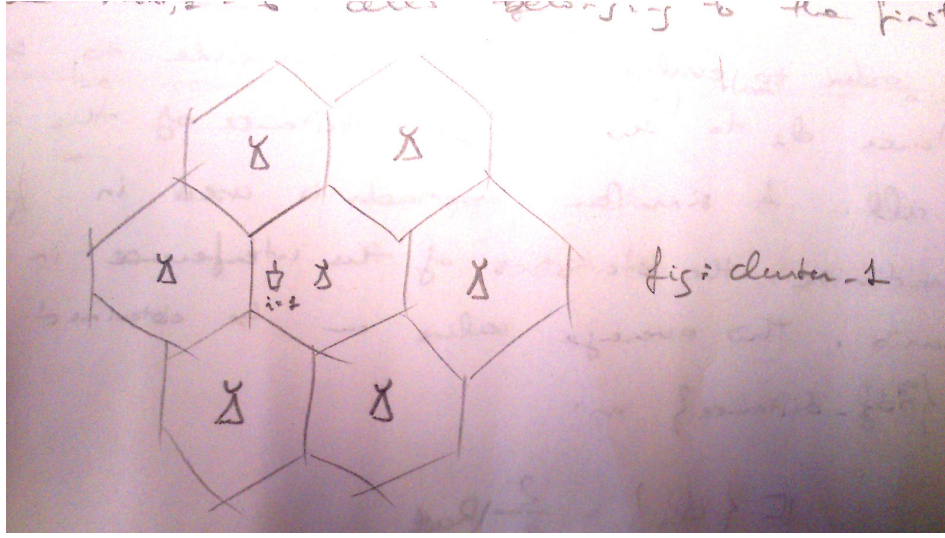


Figure 2.3: Simple scenario with a single cluster of $M = 1$ cell, and the six interfering cells surrounding it.

$$M_{\text{eq},i} \approx \frac{M_{\text{int},i}}{2} \left[1 + \left(\frac{D_i - d_i}{D_i + d_i} \right)^\gamma \right] \quad (2.7)$$

In order to evaluate (2.7), it is possible to set the distance d_i to the average distance of the user within the cell. A similar approach is used in [16] to characterize the statistics of the interference in a multi-cell scenario. This average distance can be obtained from the uniform spatial distribution in (2.1) as

$$\mathbb{E}\{d_i\} = \frac{2}{3}R_{\text{cell}} \quad (2.8)$$

and then

$$M_{\text{eq},i} \approx \frac{M_{\text{int},i}}{2} \left[1 + \left(\frac{D_i - \frac{2}{3}R_{\text{cell}}}{D_i + \frac{2}{3}R_{\text{cell}}} \right)^\gamma \right] \quad (2.9)$$

Section 2.5 deals with the comparison between simulations and the analytical results, showing that the approximation in (2.9) is rather good.



2.4 Analysis of the Rate

In the case of having global coordination or, equivalently, having only one cluster including all the BS, the interference among the users is completely eliminated through the use of BD, cf. Section 1.3. On the other hand, in a multicluster environment, it is necessary to consider the

effect of the interference coming from the cells outside the cluster. Hence, the mean achievable rate in (1.39) becomes (dropping the expectation notation)

$$R_i^{\text{BD}} = \sum_{k=1}^{\ell} \log_2 \left(1 + \frac{\hat{\lambda}_{ik} p_{ik}}{\sigma_i^2 + I_i} \right) \quad (2.10)$$

where the parameter I_i represents the average power of the total interference contributions received in each data stream of user i from the interfering BS.

It can be seen in (2.10) that the rate depends on the distance d_i of each user's equipment from the center of its cell. Using the pdf in (2.1) the average of the rate over all possible locations, making explicit the dependence on the distance d_i , can be expressed as

$$\begin{aligned} \bar{R}_i^{\text{BD}} &= \int_0^{R_{\text{cell}}} R_i^{\text{BD}}(u) \frac{2u}{R_{\text{cell}}^2} du \\ &= \int_0^{R_{\text{cell}}} \sum_{k=1}^{\ell} \log_2 \left(1 + \frac{\hat{\lambda}_{ik}(u) p_{ik}}{\sigma_i^2 + I_i(u)} \right) \frac{2u}{R_{\text{cell}}^2} du \end{aligned} \quad (2.11)$$

In (2.11) there are three parameters that determine the overall average rate, namely

- The interference $I_i(d_i)$ coming from outside the cluster.
- The effect of the channel fading and of the path loss, represented by the term $\hat{\lambda}_{ik}(d_i)$.
- The power p_{ik} assigned to the k -th data stream of user i .

In the following, the characterization of each of these parameters will be approached separately.

2.4.1 Interference

As described in Section 2.3 the contribution of interference, $I_i(d_i)$ on each data stream of user i , coming from the cells outside the cluster, can be considered as generated by an equivalent number of BS located all of them at a distance of $D_i - d_i$ from the user. Recall that, for clusters of size up to 18, D_i can take one of two values as in (2.3).

The interfering distance D_i is then normalized to the cell radius R_{cell} by setting

$$\bar{d}_i = \frac{D_i}{R_{\text{cell}}} \quad (2.12)$$

In expression (2.10) it is assumed that the interference can be treated as Gaussian noise, so that the power is calculated as the variance of that noise. Throughout the simulations that were performed, it has been found that, on average, at least 25 out-of-cluster cells contributed with significant interference.²

This number of 25 is dependent on the simulation parameters, but it gives an idea of the order of magnitude of interferers present, and it justifies the treatment of the interference as Gaussian noise, by virtue of the central limit theorem [17].

2.4.2 Fading Effect

The terms $\hat{\lambda}_{ij}$ are the squared diagonal values of the matrix $\hat{\mathbf{A}}_i$, cf. Section 1.3. This matrix is obtained in (1.36) by the combination of the channel matrix \mathbf{H}_i and a unitary matrix, $\tilde{\mathbf{V}}_i^{(0)}$.

The channel matrix \mathbf{H}_i is composed of the submatrices \mathbf{H}_{ij} , where the fading elements have a power path loss of $d_{ij}^{-\gamma}$, and the elements of \mathbf{H}_{ij} are independent from the elements of \mathbf{H}_{ik} for all j different from k .

It is possible to define an alternative set of coefficients κ_{ij} that do not include the path loss effect

$$\kappa_{ij} \triangleq \frac{\hat{\lambda}_{ij}}{d_i^{-\gamma}} \quad (2.13)$$

These coefficients are the elements of the main diagonal of the matrix

$$d_i^\gamma \hat{\mathbf{A}}_i \hat{\mathbf{A}}_i^H \quad (2.14)$$

And it can be seen that these diagonal elements are the singular values of the matrix

$$d_i^\gamma \mathbf{H}_i \tilde{\mathbf{V}}_i^{(0)} \tilde{\mathbf{V}}_i^{(0),H} \mathbf{H}_i^H \quad (2.15)$$

where \mathbf{H}_i has Gaussian entries, and $\tilde{\mathbf{V}}_i^{(0)}$ is a unitary matrix.

In the case of having $Mt = Nr$, the coefficients κ_{ij} are the eigenvalues of a Wishart matrix while, in the general case of $Mt \geq Nr$ the matrix in (2.15) can be approximated by a Wishart matrix. Through simulations, it has been verified that the mean of the eigenvalues of both matrices, the original in (2.15) and the approximate Wishart, is the same and the difference between the two Cumulative Distribution Function (CDF) is less than 10 %.

The joint pdf of the eigenvalues κ_{ij} of a Wishart matrix can be obtained when the columns of the corresponding Gaussian matrix have an identity covariance matrix

²Being significant defined as being greater than the power received at the cell edge.

$$\boldsymbol{\Sigma}_i = \mathbf{I} \quad (2.16)$$

and it is given by [18]

$$f_{\kappa_{i1}, \dots, \kappa_{i\ell}}(\kappa_{i1}, \dots, \kappa_{i\ell}) = e^{-\sum_{k=1}^{\ell} \kappa_{ik}} \prod_{k=1}^{\ell} \frac{1}{[(\ell - k)!]^2} \prod_{m=k+1}^{\ell} (\kappa_{im} - \kappa_{ik})^2 \quad (2.17)$$

However, in the evaluation of the rate the complete pdf is not needed, and only the sum

$$\sum_{i=1}^N \sum_{k=1}^{\ell} \log(\kappa_{ik}) \quad (2.18)$$

is required, which represents the expectation of the logarithm of the determinant, for which results are available, also for the general case when the covariance matrix is different from the identity, $\boldsymbol{\Sigma} \neq \mathbf{I}$, and it is given by [18]

$$\mathbb{E} \left\{ \sum_{i=1}^N \sum_{k=1}^{\ell} \log(\kappa_{ik}) \right\} = \sum_{m=1}^{N\ell} \psi(N\ell - m) + \log(|\boldsymbol{\Sigma}|) \quad (2.19)$$

where $\psi(\cdot)$ is the Euler's digamma function [19], and where the matrix $\boldsymbol{\Sigma}$ is

$$\boldsymbol{\Sigma} = \begin{bmatrix} \boldsymbol{\Sigma}_1 & \mathbf{0} & \dots & \mathbf{0} \\ \mathbf{0} & \boldsymbol{\Sigma}_2 & \dots & \mathbf{0} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0} & \mathbf{0} & \dots & \boldsymbol{\Sigma}_N \end{bmatrix} \quad (2.20)$$

with $\boldsymbol{\Sigma}_i$ the covariance matrix of the columns of the matrix $d_i^{\gamma/2} \mathbf{H}_i \widetilde{\mathbf{V}}_i^{(0)}$

$$\boldsymbol{\Sigma}_i = \mathbf{I} \left[1 + \sum_{\substack{j=1 \\ j \neq i}}^M \left(\frac{d_i}{d_{ij}} \right)^{\gamma} \right] \quad (2.21)$$

It is possible to define a parameter G_i

$$G_i \triangleq 1 + \sum_{\substack{j=1 \\ j \neq i}}^M \left(\frac{d_i}{d_{ij}} \right)^{\gamma} \quad (2.22)$$

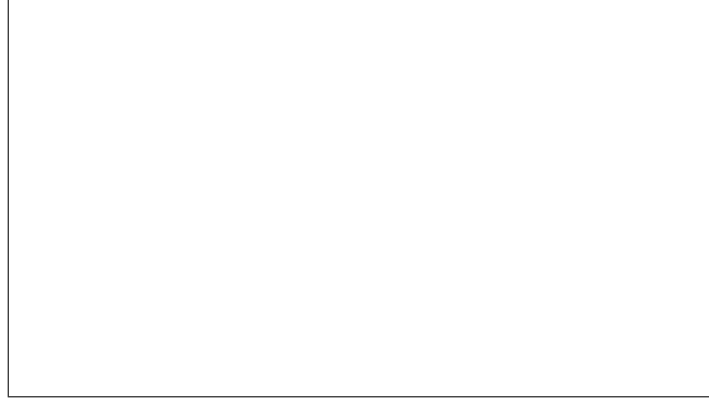


Figure 2.4: Sum-log of the terms κ_{ij} . Comparison between simulations and the values obtained using (2.19) for $t = r = 2$.

that can be considered as a cluster gain. Its value can be approximated considering the average distance of a user to its BS (2.8), and the distance from the rest of the BS in the cluster to be either $D_{\text{tier } 1}$ or $D_{\text{tier } 2}$. If additionally these two distances are normalized by the average distance in (2.8)

$$\begin{aligned}\overline{D}_{\text{tier } 1} &\triangleq \frac{3D_{\text{tier } 1}}{2R_{\text{cell}}} \\ \overline{D}_{\text{tier } 2} &\triangleq \frac{3D_{\text{tier } 2}}{2R_{\text{cell}}}\end{aligned}\tag{2.23}$$

then the gain factor G_i can be approximated for different cluster sizes as

$$G_i = \begin{cases} 1 + \frac{M-1}{2} \left[\left(\frac{1}{\overline{D}_{\text{tier } 1-1}} \right)^\gamma \left(\frac{1}{\overline{D}_{\text{tier } 1+1}} \right)^\gamma \right] & M \leq 7 \\ 3 \left[\left(\frac{1}{\overline{D}_{\text{tier } 1-1}} \right)^\gamma \left(\frac{1}{\overline{D}_{\text{tier } 1+1}} \right)^\gamma \right]^{\frac{M-7}{2}} \left[\left(\frac{1}{\overline{D}_{\text{tier } 2-1}} \right)^\gamma \left(\frac{1}{\overline{D}_{\text{tier } 2+1}} \right)^\gamma \right] & 7 < M \leq 18 \end{cases}\tag{2.24}$$

In Figure 2.4 a fixed distance for the users equal to $d_i = 2/3R_{\text{cell}}$, (2.8), is considered and the values of the sum of the natural logarithm of the values κ_{ij} are calculated through simulations and using the expression (2.19), for the case of $t = r = 2$.

It can be seen how the sum of the log values of κ_{ij} presents a diminishing increase as the cluster size M increases. This can be explained by a reduced contribution of the BS, that are farther than in smaller clusters, which becomes negligible due to the path loss.

Notice that, in the evaluation of the mean achievable rate, a factor $1/N$ is applied in order to evaluate the average rate per user, taking into consideration that in a cluster with more cells, there would be more users as well.

Thus, a decrease occurs in the mean achievable rate per user for large values of M , as it will be shown in Section 2.5.

2.4.3 Power allocation

Under the BD strategy, the transmission within each cluster is equivalent to a set of parallel non-interfering channels.

Therefore, the transmission power must be allocated in order to optimize some quality of service parameters, such as the sum rate or a weighted sum of the rates, for the users of each cluster.

This objective is subject to a maximum transmission power available at each BS, (1.45) and (1.48)

$$\sum_{i=1}^N \sum_{k=1}^{\ell} p_{ik} \|\bar{\mathbf{w}}_{j,ik}\|_2^2 \leq P_{\max} \quad (2.25)$$

for each of the M BS in the cluster.

The rate maximization problem is described in more detail in Section 1.4, and the solutions range from the simplest uniform power allocation, *cf.* Subsection 1.4.4, to an optimal allocation, *cf.* Subsection 1.4.1. In any case, the problem of power allocation is not the focus of this work because it can be solved separately and the actual powers could be inserted in the analytical expressions developed.

Hence, in the following a theoretical framework is derived for the uniform power allocation scheme, for the sake of simplicity, and an example with a different power allocation will be presented with the results.

With a uniform power allocation a common average power p_s is used for every stream of every user, as seen in (1.70). This value p_s varies according to the number of BS in the cluster, decreasing for a larger size of the cluster, since a fraction of the overall available power is spent in the coordination, to null the interference.

Substituting all the p_{ik} in (2.25) by the common value p_s it is easy to see that the condition in (2.25) is limited by the BS for which the following factor is maximum

$$\chi_j \triangleq \sum_{i=1}^N \sum_{k=1}^{\ell} \|\bar{\mathbf{w}}_{j,ik}\|_2^2 \quad (2.26)$$

Assuming that the coefficients of the precoding matrix, *i.e.* the elements of the vector $\bar{\mathbf{w}}_{j,ik}$, are Gaussian then χ_j is a Chi-squared random variables with $N' \triangleq N\ell$ degrees of freedom, and then the power p_s is related to the reciprocal value of the maximum of M random variables

$$p_s = \frac{P_{\max}}{\mathbb{E}\{\chi\}} \quad (2.27)$$

with

$$\chi = \max \{\chi_1, \dots, \chi_M\} \quad (2.28)$$

Then the probability distribution function of χ is given by

$$F_\chi(x) = P(N', x)^M \quad (2.29)$$

where $P(\cdot, \cdot)$ is the regularized Gamma function.

The mean value can be derived from the probability distribution function in (2.29) as

$$\mathbb{E}\{\chi\} = \int_0^\infty (1 - F_\chi(x)) dx \quad (2.30)$$

and it can be bounded using

$$(1 - e^{-\alpha x})^a \leq P(a, x) \leq (1 - e^{-\beta x})^a \quad (2.31)$$

with

$$\begin{aligned} \alpha &= \begin{cases} 1 & 0 < a < 1 \\ \Gamma(a+1)^{-\frac{1}{a}} & a > 1 \end{cases} \\ \beta &= \begin{cases} \Gamma(a+1)^{-\frac{1}{a}} & 0 < a < 1 \\ 1 & a > 1 \end{cases} \end{aligned} \quad (2.32)$$

where $\Gamma(\cdot)$ is the Gamma function.

The average value of χ is then bounded by

$$\frac{1}{\beta} [\psi(Mt+1) + \gamma_0] \leq \mathbb{E}\{\chi\} \leq \frac{1}{\alpha} [\psi(Mt+1) + \gamma_0] \quad (2.33)$$

with $\psi(\cdot)$ again the Euler's digamma function, and γ_0 the Euler-Mascheroni constant.

It is possible to rewrite (2.33) in terms of p_s as

$$P_{\max} \frac{\Gamma(N'+1)^{-\frac{1}{N'}}}{\psi(Mt+1) + \gamma_0} \leq p_s \leq P_{\max} \frac{1}{\psi(Mt+1) + \gamma_0} \quad (2.34)$$

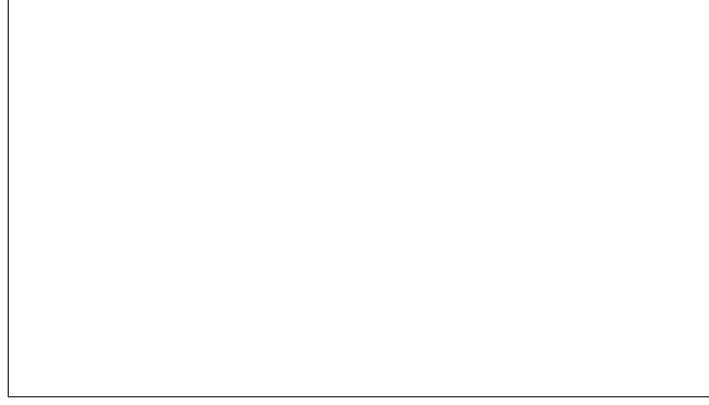


Figure 2.5: Normalized average power per stream p_s/P_{\max} for different antenna configurations: comparison between simulations and the bounds in (2.34).

In the evaluation of the rate, the lower bound in (2.34) will be considered, providing a lower bound to the average rate of each user.

The bounds for the power per stream p_s in (2.34) are compared in Figure 2.5 with the results obtained through simulations.

First thing that can be seen in Figure 2.5 is how the power p_s decreases with the size of the cluster, and this affects the mean achievable rate as it will be discussed in Section 2.5.

Secondly, Figure 2.5 shows a very good agreement between the analytical and simulation results for different antenna configurations. In particular, the upper bound is tight for small clusters while the lower bound becomes more accurate for bigger clusters.

2.4.4 Evaluation of the mean achievable rate

The performance of the coordination scheme will be measured by the mean achievable rate per user in the cluster

$$\bar{R}^{\text{BD}} = \frac{1}{N} \sum_{i=1}^N \bar{R}_i^{\text{BD}} \quad (2.35)$$

It is possible to derive a lower bound for each user's average rate in (2.11) by considering the inequality

$$\log(1+x) \geq \log(x) \quad (2.36)$$

so that the average rate for the i -th user (2.11) becomes

$$\begin{aligned}
\bar{R}_i^{\text{BD}} &\geq \frac{1}{\log(2)} \sum_{k=1}^{\ell} \int_0^{R_{\text{cell}}} \log \left(\frac{p_{ik} \kappa_{ik} u^{-\gamma}}{\sigma_i^2 + P_{\max} M_{\text{eq},i} (D_i - u)^{-\gamma}} \right) \frac{2u}{R_{\text{cell}}^2} du \\
&= \frac{1}{\log(2)} \sum_{k=1}^{\ell} \left\{ \log(\kappa_{ik}) + \log \left(\frac{p_{ik}}{P_{\max}} \right) + Z_i \right\}
\end{aligned} \tag{2.37}$$

where the interference model has been introduced, and Z_i is defined as

$$Z_i \triangleq \int_0^{R_{\text{cell}}} \log \left(\frac{u^{-\gamma}}{\frac{\sigma_i^2}{P_{\max}} + M_{\text{eq},i} (D_i - u)^{-\gamma}} \right) \frac{2u}{R_{\text{cell}}^2} du \tag{2.38}$$

The value of Z_i is derived in Appendix A, and it is

$$\begin{aligned}
Z_i &= \frac{\gamma}{2} + \log(\rho) + \bar{d}_i^2 \log \left(\frac{\bar{d}_i^\gamma}{M_{\text{eq},i} \rho + \bar{d}_i^\gamma} \right) \\
&\quad - 2\bar{d}_i^2 \gamma {}_2F_1 \left(1, \frac{1}{\gamma}; \frac{\gamma+1}{\gamma}; \frac{-\bar{d}_i^\gamma}{M_{\text{eq},i} \rho + \bar{d}_i^\gamma} \right) \\
&\quad + \frac{\bar{d}_i^2}{2} \gamma {}_2F_1 \left(1, \frac{2}{\gamma}; \frac{\gamma+2}{\gamma}; \frac{-\bar{d}_i^\gamma}{M_{\text{eq},i} \rho + \bar{d}_i^\gamma} \right) \\
&\quad - (\bar{d}_i^2 - 1) \log \left(\frac{(\bar{d}_i - 1)^\gamma}{M_{\text{eq},i} \rho + (\bar{d}_i - 1)^\gamma} \right) \\
&\quad + 2\bar{d}_i (\bar{d}_i - 1) \gamma {}_2F_1 \left(1, \frac{1}{\gamma}; \frac{\gamma+1}{\gamma}; \frac{-(\bar{d}_i - 1)^\gamma}{M_{\text{eq},i} \rho} \right) \\
&\quad - \frac{(\bar{d}_i - 1)^2}{2} \gamma {}_2F_1 \left(1, \frac{2}{\gamma}; \frac{\gamma+2}{\gamma}; \frac{-(\bar{d}_i - 1)^\gamma}{M_{\text{eq},i} \rho} \right)
\end{aligned} \tag{2.39}$$

where ${}_2F_1(\cdot)$ is the hypergeometric function, and ρ is the SNR as defined in (1.25).

Combining (2.19), (2.34), (2.35), and (2.37) the analytical expression for the mean achievable rate per user is

$$\begin{aligned}
\bar{R}^{\text{BD}} &\geq \frac{1}{N \log(2)} \left[\log(G) + \sum_{m=1}^{N\ell} \psi(N\ell - m) \right. \\
&\quad \left. + N\ell \log \left(\frac{\Gamma(N' + 1)^{-\frac{1}{N'}}}{\psi(Mt + 1) + \gamma_0} \right) + \ell \sum_{i=1}^N Z_i \right]
\end{aligned} \tag{2.40}$$

where G is the combination of all the G_i in (2.24).

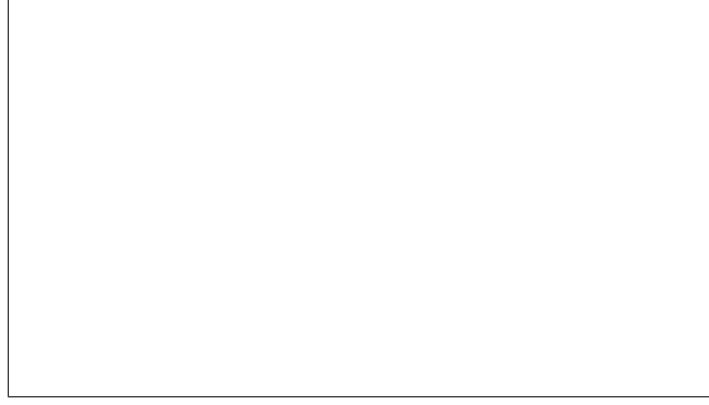


Figure 2.6: Scenario used for the simulations, with 7 hexagonal tiers.

2.5 Numerical Results

In this section, the results derived from the analytical expression in (2.40) are compared with the results obtained through simulations.

The scenario for the simulations is a network composed of 169 cells, laid out as 7 concentric tiers of hexagonal cells Figure 2.6.

The cell radius, unless stated otherwise, it is assumed to be $R_{\text{cell}} = 1.4$ km.

All the results are averaged over 5000 random trials. In each of these trials the position of the users was randomly set according to a uniform distribution inside each cell, Section 2.3. Also, for each of the trials, a random channel was generated according to the model described in Section 1.2, with a path loss exponent of $\gamma = 3.8$.

The parameter evaluated in the simulations is the achievable rate defined in (2.10), in which the different variables required (transmission power, interference power, $\hat{\lambda}_{ik}$, etc) were obtained by simulations.

As it has already been mentioned, the clusters considered are static, *i.e.* they are fixed and do not change for all the simulations. Despite this, not all the clusters must have the same shape for the same cluster size M . In fact, in the simulations, the clusters were generated following a heuristic approach that tries to group cells in a compact way, with a regular shape, by minimizing the sum of the inter-cell distances, thus to avoid long clusters. Note, however, that some values of M do not allow for regular clusters, *i.e.* hexagonal, to be formed. Figure 2.7 shows an example of this situation, where not all the clusters have the same shape.

2.5.1 Comparison with simulations and with Minimum Mean Squared Error (MMSE)

In order to validate the analytical expression (2.40), first it is compared with the mean achievable rate simulated using BD and, although the analytical derivations did not consider it, the



Figure 2.7: Irregular shaping of the clusters, due to the heuristic clustering algorithm used.



Figure 2.8: Mean achievable rate per user as a function of the number of cells in the cluster for $r = 2$, $t = 2$, variable values of SNR, $\gamma = 3.8$.



MMSE described in [20].

The antenna configuration used for this first comparison is $r = t = 2$, and also different values of the SNR, as defined in (1.25), are used so to observe the behaviour in different SNR regimes.

Figure 2.8 shows the mean achievable rate as a function of the cluster size M .

As expected, the MMSE approach outperforms the BD strategy at low SNR. On the other hand, for moderate values of the SNR BD is able to provide comparable, and even more favorable, results, thus showing that the interference dominates over the noise for regimes other than the low SNR regime.

A very good agreement between the theoretical result in (2.40) and the simulations is clear in Figure 2.8, where also some variations can be seen in the simulation results. This is mainly due to the variability of the cluster shape, as seen in Figure 2.7, for different cluster sizes, and not that much due to a low number of simulations that were averaged. That irregular



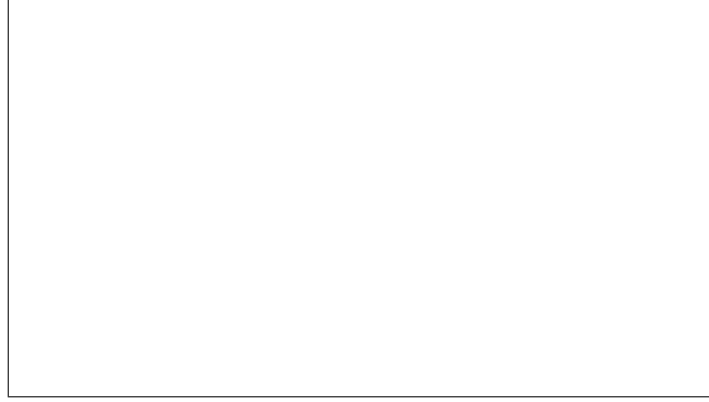


Figure 2.9: Mean achievable rate as a function of the SNR for different values of the path-loss coefficient γ .

shape of the clusters, despite being more or less controlled in the heuristic cluster selection algorithm, affects the simulation results in the form of the variability shown in the figures.

[21] points out at a fundamental limit of cooperation, and it is shown how the gains from cooperation cannot be unbounded, and that increasing the number of coordinated elements may saturate the performance achieved. This very same behaviour can be observed in Figure 2.8, both for BD and MMSE, where the mean achievable rate do not grow unboundedly with the cluster size, and an optimum value of the size M can be found.

Figure 2.9 also shows a similar behaviour. In this case, the mean achievable rate is represented as a function of the SNR, and there is an SNR at which the rate stops growing. This threshold SNR depends on the propagation path loss exponent γ because it directly determines the influence of the interference. In particular, the saturation SNR for BD is higher than for MMSE. For the former, it is always above 20 dB, for the scenarios considered, even for very small path loss exponents. This means that the saturation occurs for relatively large values of the SNR which is of practical importance because it would be possible to deliver good performance, using BD, within a practical range of SNR values.

Under the restriction that the same P_{\max} is transmitted for all values of γ , the saturation occurs at different levels for each γ , although the general conclusions do not change.

In order to complete the validation of the theoretical results with the simulations, a fixed value of SNR= 25 dB and different antenna configurations were considered in Figure 2.10, still using uniform power allocation. The discrepancies between the theoretical results and the simulations are not only due to the approximations, but also to the fact that in the simulation scenario not all the clusters have the same shape, despite having the same number of cells.

2.5.2 Effect of the power allocation

The cause of the decrease of the rate with respect to the cluster size, as seen in Figure 2.8, is two-fold:

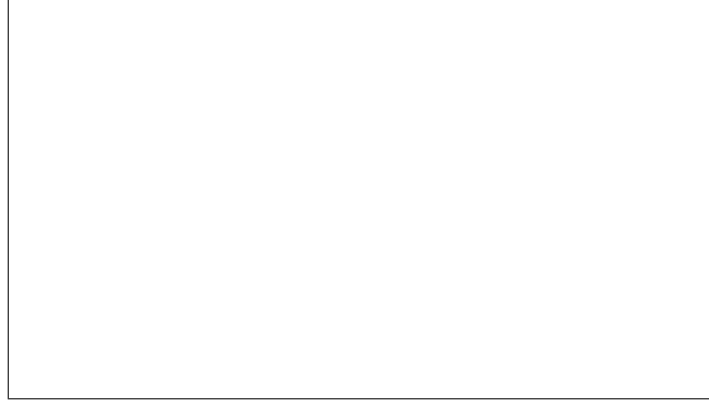


Figure 2.10: Mean achievable rate per user as a function of the cluster size for SNR= 25 dB, $\gamma = 3.8$ and different antenna configurations.

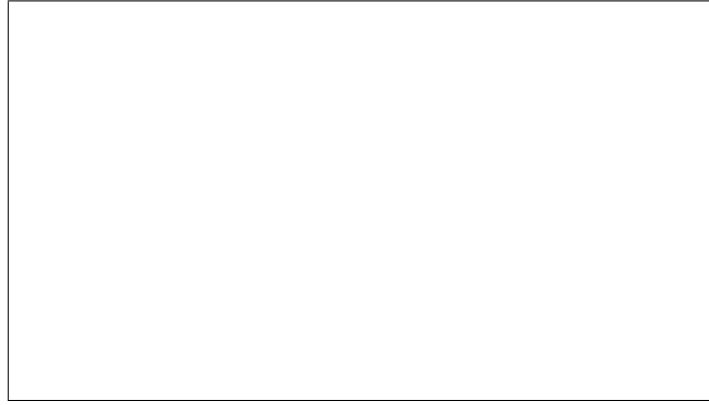


Figure 2.11: Comparison between different power allocation schemes, namely uniform and optimal. Mean achievable rate per user as a function of the cluster size for $r = t = 2$, $\gamma = 3.8$ and different values of SNR.

- First, the value of the attenuation experienced by each data stream, $\hat{\lambda}_{ik}$ decreases when the size of the cluster increases, as shown in Figure 2.4.
- Second, the power assigned to each data stream, the terms p_{ik} that for a uniform power allocation are all equal to p_s , also decreases as the the cluster grows, Figure 2.5, due to a coordination "loss".

In this section, a different power allocation scheme, other than the uniform, was used in order to verify if the behavior observed in Figure 2.8, where the rate decreases with the cluster size, is due to the power allocation.

The optimal power allocation used was obtained by means of numerical optimization, as described in [10], and those powers were plugged into (2.40) instead of the uniform power allocation.



Figure 2.12: Optimum cluster size as a function of the SNR, $\gamma = 3.8$ and different antenna configurations.

In Figure 2.11 the rates obtained with the uniform and the optimum power allocation are compared and represented versus the size of the cluster.

Although, as expected, the optimal outperforms the uniform power allocation, both curves show a similar trend, meaning that a reduced growth of the rate (and in some situations a reduction of the rate itself) is not due to the power allocation scheme, but it is the manifestation of a fundamental limitation of the coordination scheme, along the lines of the results in [21].

2.5.3 Optimum cluster size

Figure 2.8 through Figure 2.11 show a common trend which is the rate having a reduced growth, or even a reduction, with the cluster size.

Given this, it is possible and interesting to find the cluster size M that can maximize the mean achievable rate.

In the case of the rate actually decreasing with M , the optimum value can be readily obtained as the value of M for which the maximum rate is obtained.

When the rate does not decrease, but its growth with M is diminished, it is not possible to find a maximum for the rate. The optimum value of M is calculated by considering the relative change of the rate. The optimum is found when the marginal increase of the rate is below a given percentage. In the case under study, the threshold was set to a 10%.

Figure 2.12 represents the optimum value of M as a function of the SNR, for different antenna configurations, and for the power allocation schemes considered until now, uniform and optimum.

It can be seen that, for a wide range of SNR, the optimum value is limited to around 7–10 cells.

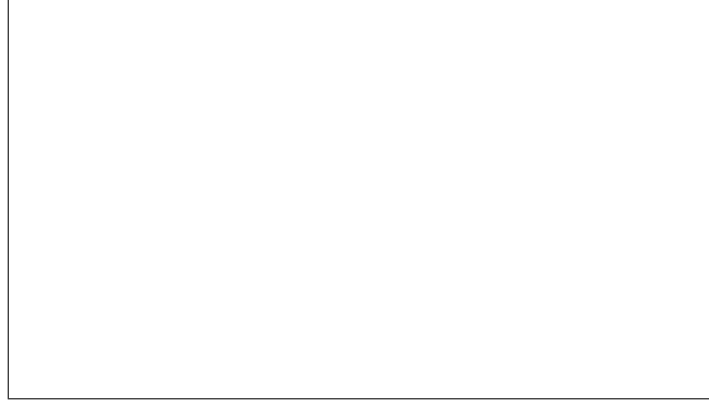


Figure 2.13: Mean achievable rate and payload rate as a function of the cluster size M with different SNRs.

Only for high SNR is it more convenient to increase the cluster size, since the reduction of the interference can compensate the decrease of the cluster gain due to the decrease of the factors $\hat{\lambda}_{i_k}$. This only happens for the case of considering the optimal power allocation, because in the case of the uniform power scheme there is the additional decrease of the power allocated to each stream, p_s .

2.5.4 Effect of signalling overhead

It is common in the literature to not take into account the effect of the signalling overhead, and the same has been done in all the previous results of this work.

However, if a certain percentage of the available resources are dedicated to channel estimation and control signalling, the effective SNR and the payload that can be delivered are reduced with respect to the global achievable rate.

In [21] the overhead, incurred by channel estimation requirements, is accounted for by a percentage α which should grow, at least, linearly with the cluster size, up to a maximum value, to prevent it from being greater than 100%. In their case, the SNR and the rate were effectively reduced by a factor of α , getting $\text{SNR}_{\text{eff}} = (1 - \alpha) \text{SNR}$ and $R_{\text{eff}} = (1 - \alpha) R$.

In order to show the effect of the overhead on the achievable rate, in Figure 2.13 a very conservative approach is adopted, in which the value of the reduction factor α scales linearly with the cluster size M , up to a maximum of 10% for a cluster size of 19.

Figure 2.13 shows the comparison of considering and not considering the overhead. It can be seen that even with a small amount of overhead, increasing values of M lead to a worse performance.

Moreover, it should be stressed that the actual definition of signalling overhead and its management is usually delegated to the operator implementation, and this is seldom defined in the standards. Thus, its quantitative effect can change considerably depending on how the overhead is defined.

Appendix A

Derivation of Z_i

Using the change of variable

$$v = \frac{u}{R_{\text{cell}}} \quad (\text{A.1})$$

in (2.38), it becomes

$$Z_i = \int_0^1 \log \left(\frac{v^{-\gamma}}{\frac{\sigma_i^2}{P_{\max} R_{\text{cell}}^{-\gamma}} + M_{\text{eq},i} (\bar{d}_i - v)^\gamma} \right) 2v dv \quad (\text{A.2})$$

Recall the definition of the SNR ρ in (1.25), then

$$\begin{aligned} Z_i &= \int_0^1 \log \left(\frac{v^{-\gamma}}{\frac{1}{\rho} + M_{\text{eq},i} (\bar{d}_i - v)^{-\gamma}} \right) 2v dv \\ &= \underbrace{2 \int_0^1 \log \left(\frac{1}{\frac{1}{\rho} + M_{\text{eq},i} (\bar{d}_i - v)^{-\gamma}} \right) v dv}_{\text{I}} - \underbrace{2\gamma \int_0^1 \log(v) v dv}_{\text{II}} \end{aligned} \quad (\text{A.3})$$

The term II in (A.3) is evaluated using (2.723) from [19], namely

$$\int v \log(v) dv = v^2 \left(\frac{\log(v)}{2} - \frac{1}{4} \right) \quad (\text{A.4})$$

Using the change of variable

$$w = \bar{d}_i - v \quad (\text{A.5})$$

in the term I of (A.3), it becomes

$$\begin{aligned} (\text{I}) &= - \int_{\bar{d}_i}^{\bar{d}_i-1} \log \left(\frac{\rho}{1 + M_{\text{eq},i} \rho w^{-\gamma}} \right) (\bar{d}_i - w) dw \\ &= \underbrace{\int_{\bar{d}_i-1}^{\bar{d}_i} \log(\rho) (\bar{d}_i - w) dw}_{\text{III}} + \underbrace{\int_{\bar{d}_i-1}^{\bar{d}_i} \log \left(\frac{1}{1 + M_{\text{eq},i} \rho w^{-\gamma}} \right) (\bar{d}_i - w) dw}_{\text{IV}} \end{aligned} \quad (\text{A.6})$$

The term III of (A.6) is straightforward to calculate

$$\int_{\bar{d}_i-1}^{\bar{d}_i} \log(\rho) (\bar{d}_i - w) dw = \log(\rho) \quad (\text{A.7})$$

On the other hand, the term IV in (A.4) can be computed using the following result from [22]

$$\begin{aligned} \int_{\bar{d}_i-1}^{\bar{d}_i} \log \left(\frac{1}{1 + M_{\text{eq},i} \rho w^{-\gamma}} \right) (\bar{d}_i - w) dw &= w \left[\frac{2\bar{d}_i - w}{2} \log \left(\frac{w^\gamma}{M_{\text{eq},i} \rho + w^\gamma} \right) \right. \\ &\quad \left. - \bar{d}_i \gamma {}_2F_1 \left(1, \frac{1}{\gamma}; \frac{\gamma+1}{\gamma}; \frac{-w^\gamma}{M_{\text{eq},i} \rho} \right) \right. \\ &\quad \left. + \frac{w^\gamma}{4} {}_2F_1 \left(1, \frac{2}{\gamma}; \frac{\gamma+2}{\gamma}; \frac{-w^\gamma}{M_{\text{eq},i} \rho} \right) \right] \end{aligned} \quad (\text{A.8})$$

where

$${}_2F_1(a, b; c; z) = \frac{\Gamma(c)}{\Gamma(b)\Gamma(c-b)} \int_0^1 \frac{x^{b-1} (1-x)^{c-b-1}}{(1-x)^a} dx \quad (\text{A.9})$$

is the hypergeometric function.

Combining the previous, Z_i can be expressed as

$$\begin{aligned}
Z_i = & \frac{\gamma}{2} + \log(\rho) + \bar{d}_i^2 \log \left(\frac{\bar{d}_i^\gamma}{M_{\text{eq},i}\rho + \bar{d}_i^\gamma} \right) \\
& - 2\bar{d}_i^2 \gamma {}_2F_1 \left(1, \frac{1}{\gamma}; \frac{\gamma+1}{\gamma}; \frac{-\bar{d}_i^\gamma}{M_{\text{eq},i}\rho + \bar{d}_i^\gamma} \right) \\
& + \frac{\bar{d}_i^2}{2} \gamma {}_2F_1 \left(1, \frac{2}{\gamma}; \frac{\gamma+2}{\gamma}; \frac{-\bar{d}_i^\gamma}{M_{\text{eq},i}\rho + \bar{d}_i^\gamma} \right) \\
& - (\bar{d}_i^2 - 1) \log \left(\frac{(\bar{d}_i - 1)^\gamma}{M_{\text{eq},i}\rho + (\bar{d}_i - 1)^\gamma} \right) \\
& + 2\bar{d}_i (\bar{d}_i - 1) \gamma {}_2F_1 \left(1, \frac{1}{\gamma}; \frac{\gamma+1}{\gamma}; \frac{-(\bar{d}_i - 1)^\gamma}{M_{\text{eq},i}\rho} \right) \\
& - \frac{(\bar{d}_i - 1)^2}{2} \gamma {}_2F_1 \left(1, \frac{2}{\gamma}; \frac{\gamma+2}{\gamma}; \frac{-(\bar{d}_i - 1)^\gamma}{M_{\text{eq},i}\rho} \right)
\end{aligned} \tag{A.10}$$

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