## 图之 习题选讲

浙江大学 陈 越



# 旅游规划



### 题意理解

- ■城市为结点
- 公路为边
  - □ 权重1: 距离
  - □ 权重2: 收费
- ■単源最短路
  - □ Dijkstra 距离
  - □ 等距离时按收费更新

#### Sample Input:

2 3 1 20

1/20 0 2/20 4/10 2 2/30 3 1/20



### 核心算法

```
void Dijkstra( Vertex s )
{ while (1) {
                                                   V = \lambda \psi  \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi \lambda \psi 
                                                       if (这样的v不存在)
                                                                               break:
                                                      collected[V] = true;
                                                       for (V的每个邻接点W)
                                                                                  if ( collected[W] == false )
                                                                                                                       if ( dist[V]+E<sub><V,W></sub> < dist[W] ) {</pre>
                                                                                                                                                      dist[W] = dist[V] + E_{\langle V,W \rangle};
                                                                                                                                                   path[W] = V;
                                                                                                                                                     cost[W] = cost[V] + C<sub><V,W></sub>;
                                                                                                                     else if ( (dist[V]+E<sub><V,W></sub> == dist[W])
                                                                                                                                                                                                                                                                                    && (cost[V]+C_{v,W} < cost[W]))
                                                                                                                                                cost[W] = cost[V] + C_{\langle V,W \rangle};
```



### 其他类似问题

- ■要求数最短路径有多少条
  - $\square$  count[s] = 1;
  - □ 如果找到更短路: count[W]=count[V];
  - □ 如果找到等长路: count[W]+=count[V];
- ■要求边数最少的最短路
  - $\square$  count[s] = 0;
  - □ 如果找到更短路: count[W]=count[V]+1;
  - □ 如果找到等长路: count[W]=count[V]+1;

