

Q=) flowers  $\rightarrow$  K the group occurs for white

there is constraint of grouping on white but not  
on red

4 flowers

K=2

R R R R — 1

W W R R — 1

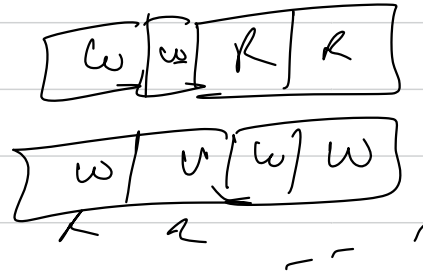
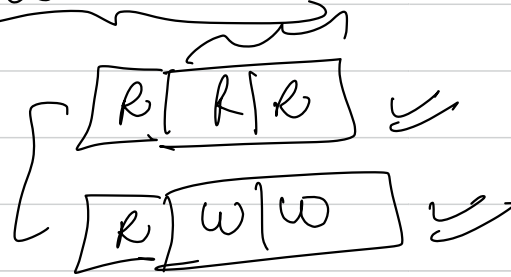
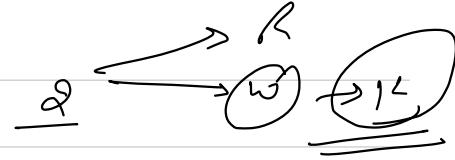
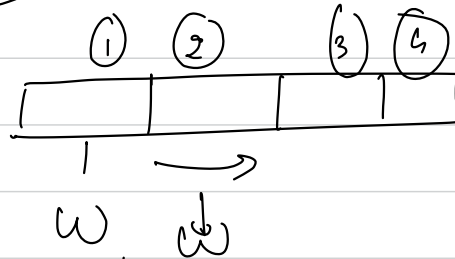
R W W R — 1

W W W W — 1

R R W W — 1

5

is DP even used 22



1 query  $\rightarrow$   $O(1)$  or  $(\log_2 n)$

How can we efficiently calc no. of ways for any value  $i$   $\rightarrow$   $i$  flowers will be calcs

$i=1$   $\rightarrow$   $k$

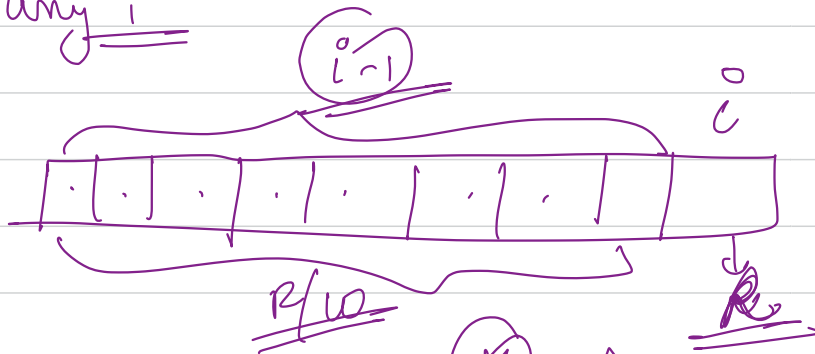
$k=1$

$k \geq 1$

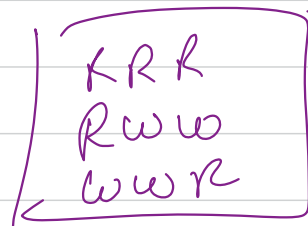
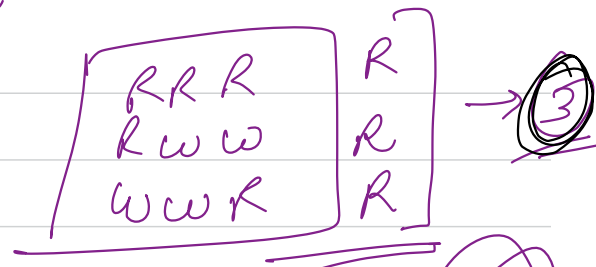
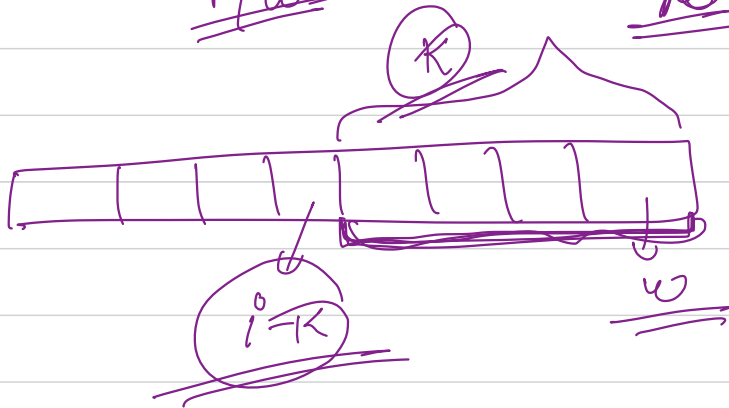
$i=0$   $\rightarrow$  one valid way  
 $1$

for any  $i$

~~$k=8$~~

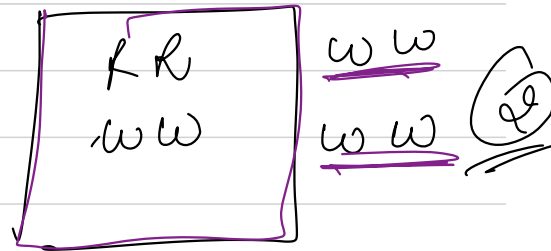


~~$i=10$~~



$i=4$

$i=4-2$



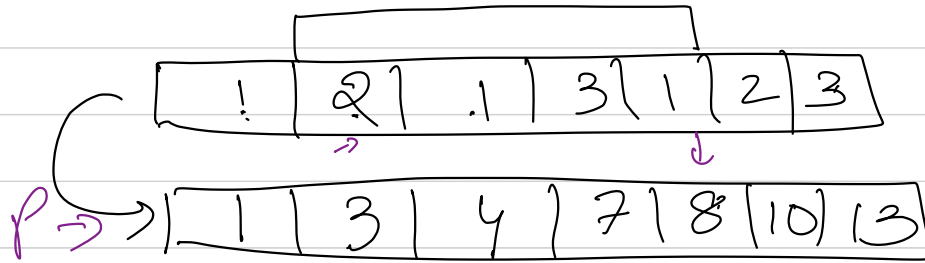
$$\hookrightarrow \underline{f(i, k)} = \begin{cases} f(i-1, k) & \text{if } \underline{i < k} \\ f(i-1, k) + f(i-k, k) & \text{if } i \geq k \end{cases}$$

$f(i, k)$  which  
 counts the no of  
 ways to take  
 $i$  from

dp array  $\rightarrow$  precompute for all value of i

$\hookrightarrow$  prefix sum

[a, b]

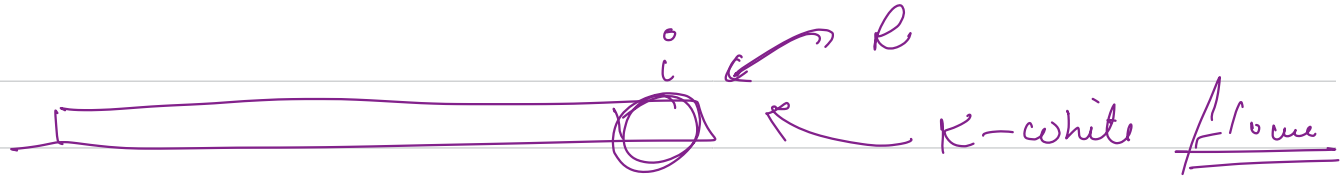


$$[8 - 3 + 2]$$

$$8 - 1 \Rightarrow 7$$

$$p[b] - p[a] + \overbrace{arr[a]} \quad \checkmark$$

$$p[b] - p[a-1] \quad \underline{\underline{\quad}}$$



$k=2$

$\begin{matrix} R & R & R & R \\ R & W & W & R \\ W & W & R & R \end{matrix}$

$i=3$   
 $i=4$

$k-2=2$

3  
 $i=4$

$i=5$

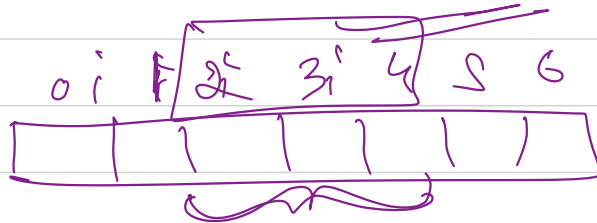
$\begin{matrix} R & R \\ W & W \end{matrix}$

$\begin{matrix} W & W \\ W & W \end{matrix}$

2

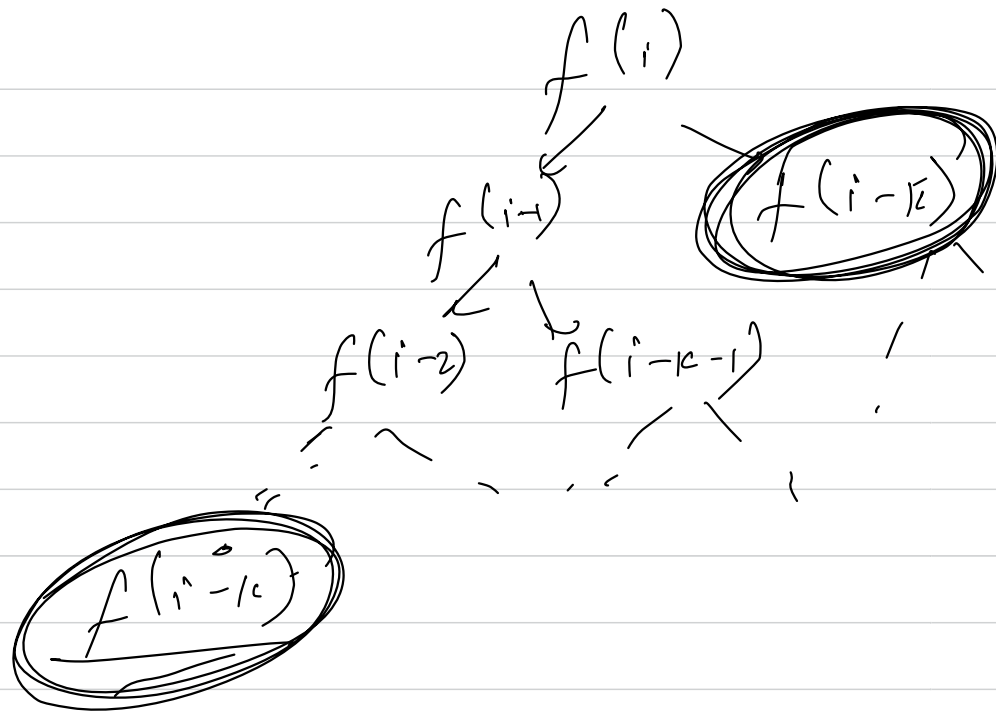
$1 \leq 10^9$   
 $n \leq 10^5$

TL  
 $O(n)$



dp

poj



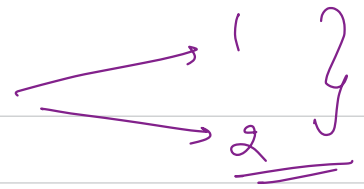
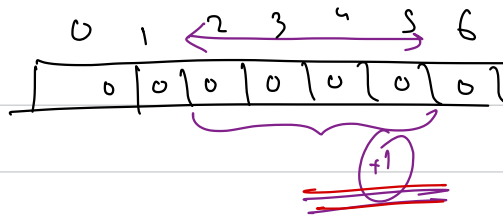


$a_1, a_2, a_3, \dots, a_n$

$\rightarrow p_1, p_2, p_3, \dots, p_n$

$$p_k = \sum_{i=1}^k a[i] \quad \leftarrow \text{prefix sum}$$

$$\text{sum}(a, b) = \underline{\underline{p[b] - p[a-1]}}$$



$[L, R]$   $(+1)$

$[L, R]$   $\rightarrow i \rightarrow [L, R]$

$+1 \rightarrow$  how many times I was added

0, 2

1, 4

0, 1

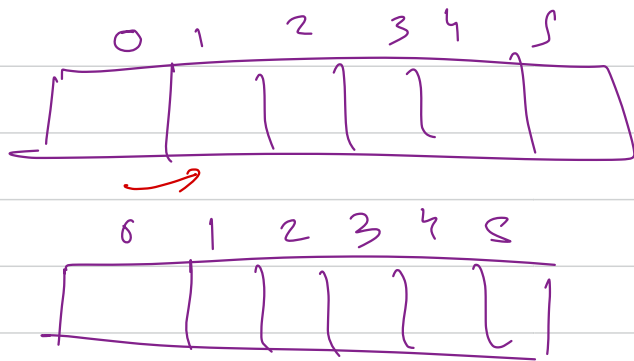
4, 5

under  $\rightarrow \underline{0} \rightarrow$  only the queue starts from 0

under  $\rightarrow \underline{1} \rightarrow$  all queue starts from 0 + starts from 1 - ending at 0

$O(n)$  start

end



$l, r$

$start[l]++$

$end[r]++$

on

$idx \rightarrow 0 \rightarrow S_0$

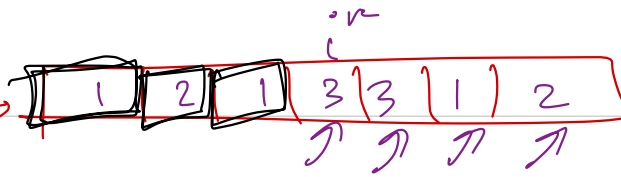
$idx \rightarrow 1 \rightarrow S_0 + S_1 - e_0$  ✓

$idx \rightarrow 2$   $\rightarrow S_0 + S_1 + S_2 - e_0 - e_1$  ✓

$S_0 + S_1 - e_0$  +  $S_2 - e_1$

$idx \rightarrow 3 \rightarrow S_0 + S_1 + S_2 + S_3 - e_2 - e_1 - e_0$  ✓

result of range addition



$dp[i]$

coins in  $n$ th box

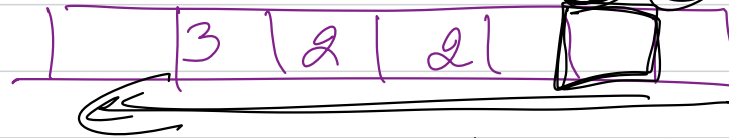
Super sum

X

$X, X+1, X+2, \dots$

at least X coins

coins



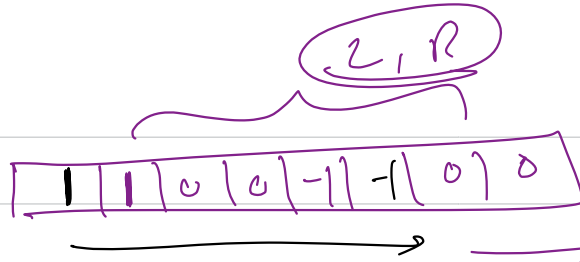
Super sum

stores that at the  $i$ th index how many

boxes have exactly  $i$  coins

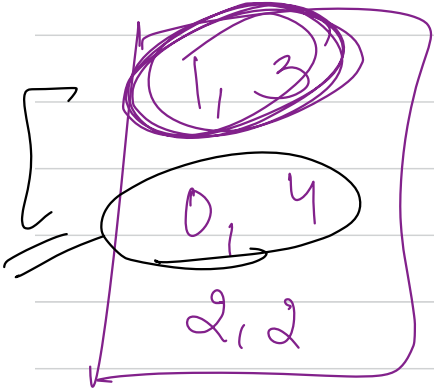
$\rightarrow$  X at coins[X]

Diff array



Prefix sum

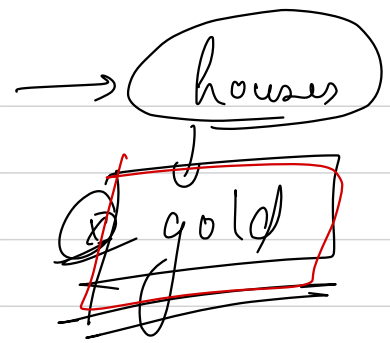
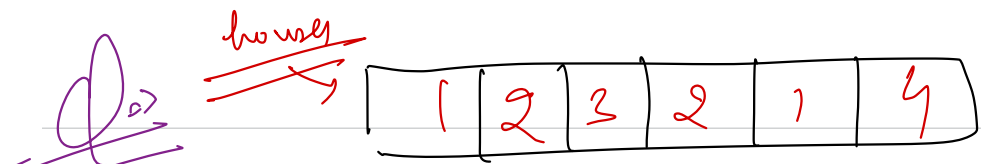
1, 2, 2, 2, 1, 0, 0, 0



$1, 3$

$+1$

$arr[l] += 1$   
 $arr[r+1] -= 1$



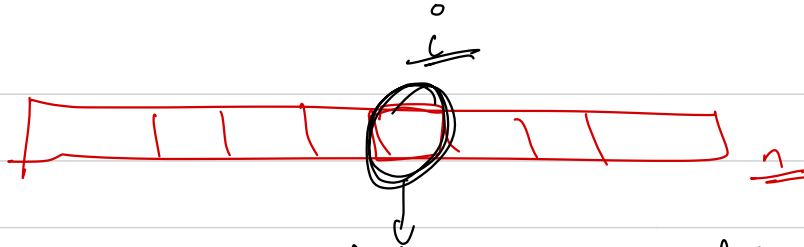
Robber

cannot rob  
of consecutive  
houses

max amount

[1, 2, 3, 1]

Q



whether the robber should rob  $i$ <sup>th</sup>  
house or not

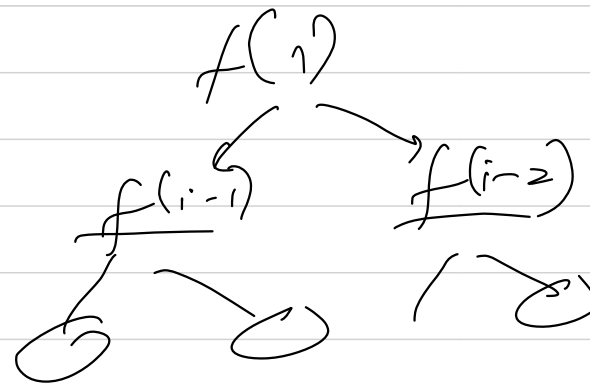
if he is  
robber in  $(i-1)$ <sup>th</sup> house can't be robbed

2 DP

$$f(i) = \max(a[i] + f(i-2), f(i-1))$$

↓  
 change ore  
 rokey n  
 in hour

man profit  
 tell me in  
 hour





$$dp[0] = a[0]$$

$$dp[i] = \underline{\max(a[0], a[i])}$$

$$dp[i] = \max(a[i] + dp[i-2], \underline{dp[i-1]})$$

$$\underline{\underline{dp[n]}}$$