

---

---

---

---

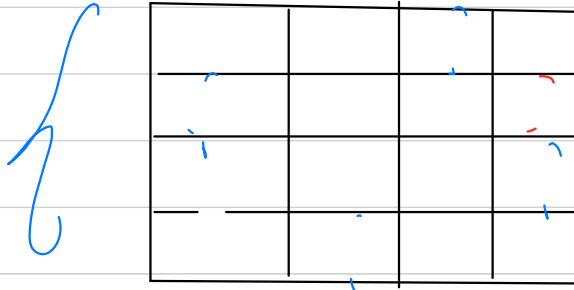
---



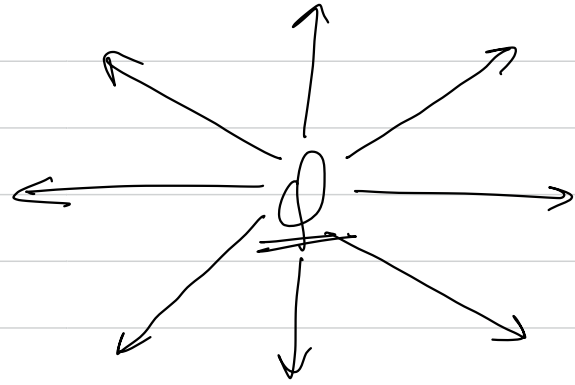
Q2 What is pruning??

↳ pruning is in the context with abandoning those sub problems which will never lead to a successful answer.

★ N Queen problem ==

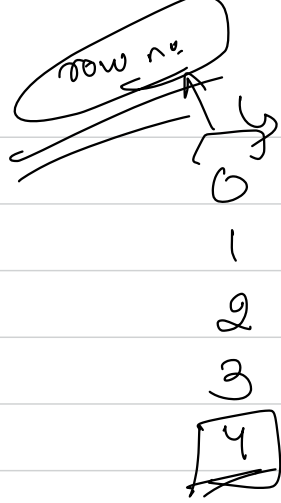


N x N



(N=4) → ans = 2

Given a  $N \times N$  board, find the no. of ways to place  $N$  queens, so that no queen attacks the other.



2

→ $Q_1$	...		
		$Q_2$	
	...		

$Q_1$   $Q_2$   $Q_3$   $Q_4$

$N \times N$

$N=4$

→ Before I place a queen, I will check that  
is it a safe spot to place.

## Backtracking

⇒

Let's try to write recursive solution.

# Base Case ⇒

# Recursion task → Recursively place  $N-1$  queen in my  $N-1$  rows

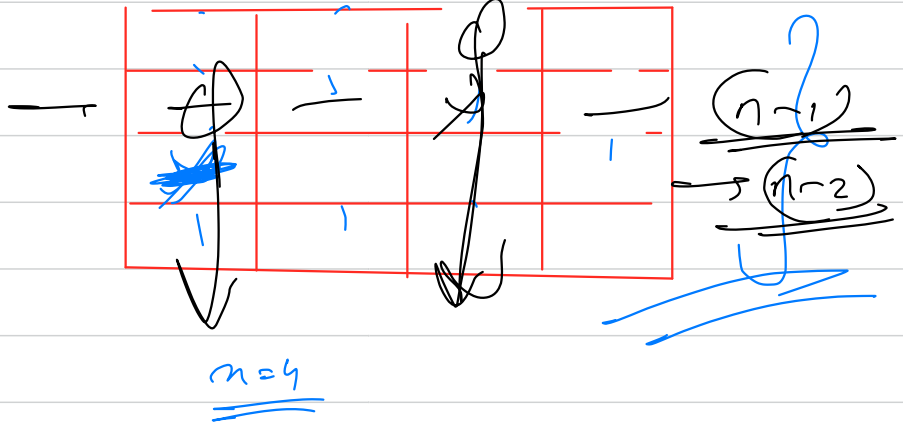
# Self work → check if the spot is safe  
if yes then place the queen

we will place our queen in current row

CallStack



→ 70 lines



count = 1

0,0

space complexity  $\rightarrow O(n^2)$

$$T(n) = n \times \underbrace{T(n-1)}^{\nwarrow} + O(n^2)$$

$$T(n-1) = (n-1) \times T(n-2) + O(n^2)$$

$$T(n-2) = (n-2) \times T(n-3) + O(n^2)$$

$$\vdots$$

$$T(2) = 2 \times T(1) + O(n^2)$$



$$T(n) = n \times \underbrace{T(n-1)} + O(n^2)$$

$$T(n) = n \left( \underbrace{n \times T(n-2)} + O(n^2) \right) + O(n^2)$$

$$T(n) = n \left( (n-1) \times \left( (n-2) \times T(n-3) \dots \right) \right)$$

$\hookrightarrow$   ~~$O(n!)$~~   $\rightarrow$   $O(n!)$

$\downarrow$   
Tighter bound

~~loose bound~~  
 $\rightarrow$   $n$

Q Given a  $N \times N$  board with a knight placed at  $(0, \underline{0})$ . Check if the knight can cover all the cells of the board, using its usual chess moves without visiting any cell twice.

7

Km (0-9)	0				17	8	
				9			16
	1.					18	7
					10	15	
		2.			23	6	19
				25	20	11	14
			3.	22	13	24	5
					4.	21	12

8x8