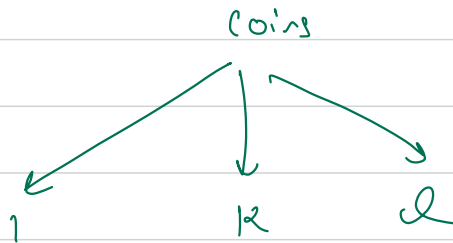
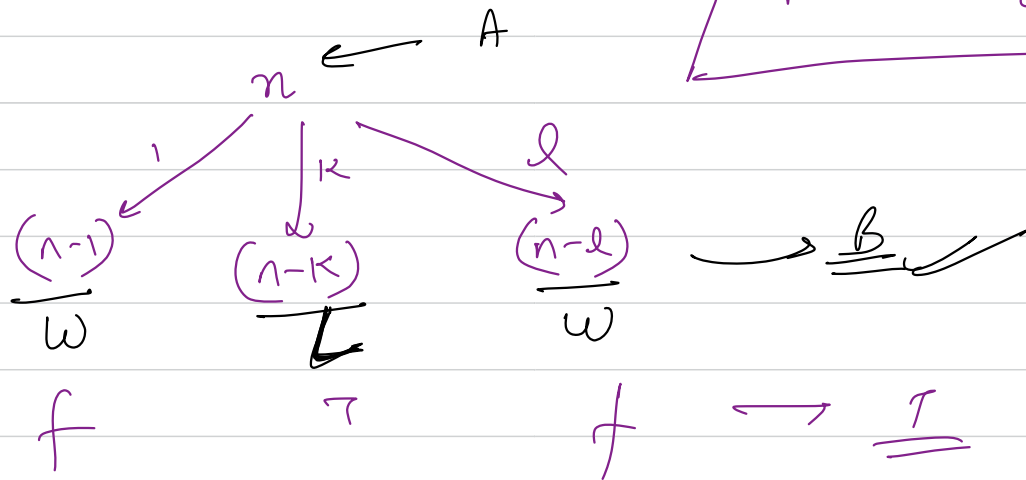


Q=) Mcoins SPOJ $\rightarrow n \rightarrow$ no. of coins



(A), B
starts first

if $(n=1) \rightarrow$ A wins always



$f(n)$ recursively give

for given n coins
who will win

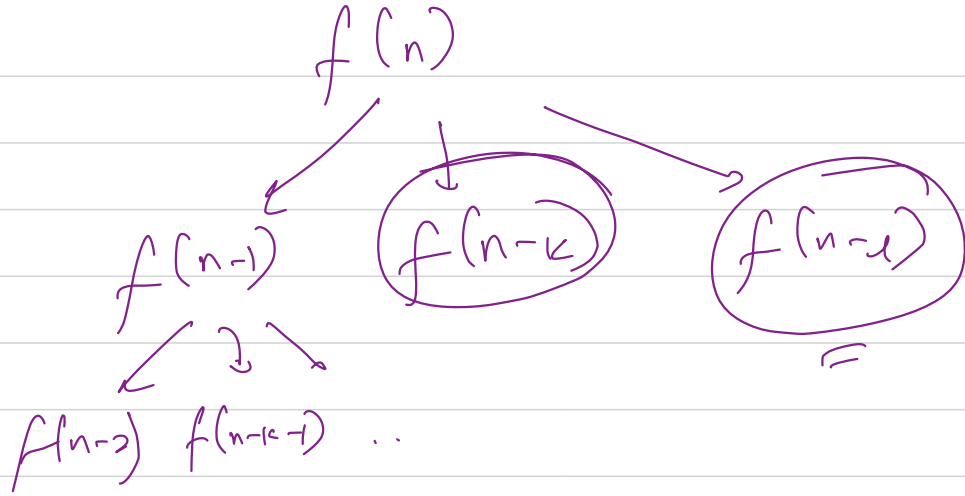
$f(1) \rightarrow \underline{\underline{A}}$

=

$$\left\{ \begin{array}{l} f(n-1) \rightarrow f \\ \text{or} \\ f(n-k) \rightarrow T \quad n-k \geq 1 \\ \text{or} \\ f(n-l) \rightarrow \underline{\underline{f}} \quad n-l \geq 1 \end{array} \right.$$

we say if A wins so that boolean True &
if B wins that's a boolean false

Overlapping
subproblem



$f(n-k)$

$f(n-1)$

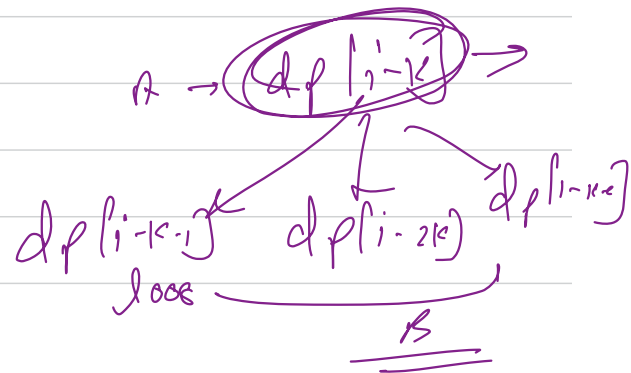
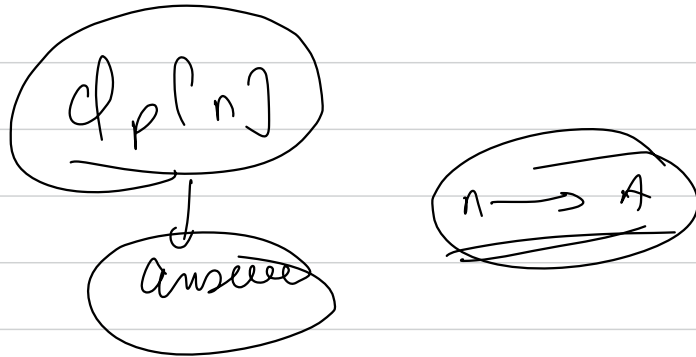
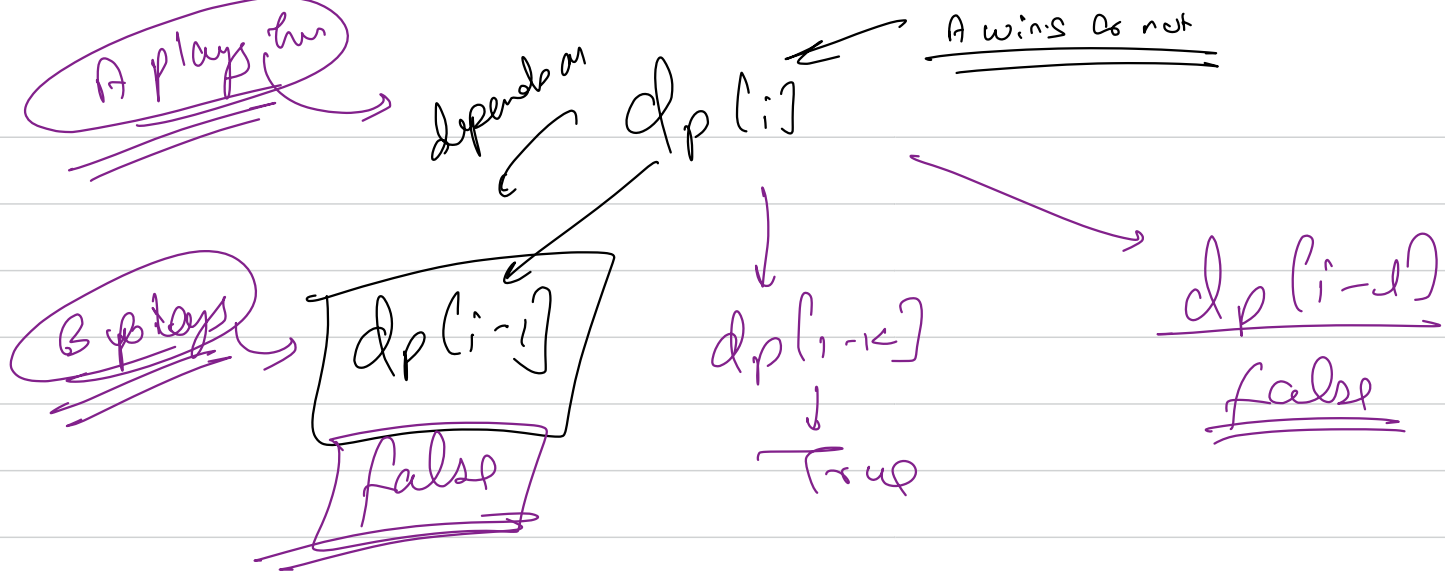
Bu \rightarrow $dp[i] \rightarrow 1$ \rightarrow A wins

$f(i) \rightarrow 1 \rightarrow$ True
 $0 \rightarrow$ false

$dp[i]$ \rightarrow $dp[i-1] \rightarrow$
 $dp[i-2] \rightarrow$
 $dp[i-3] \rightarrow$
A wins here
if anyone is false

State of dp

$dp[i]$ \rightarrow if we get to a state i ($i = \text{no. of coins}$) then
whether any player wins or not by coming to $n=0$



$dp[i] \leftarrow$

if any player
comes to this
state it will
never do not

Qn You are given a 2D grid of $n \times m$ dimension filled with integers. find a rectangle / subgrid with maximum sum of elements.

1	2	-1	-4	-20
-8	-3	4	2	1
3	8	10	1	3
-4	-1	1	2	-6

Imp Interview

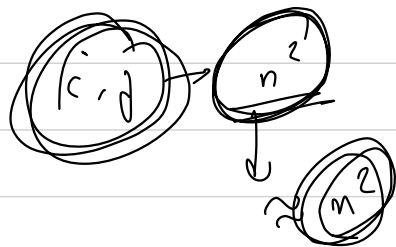
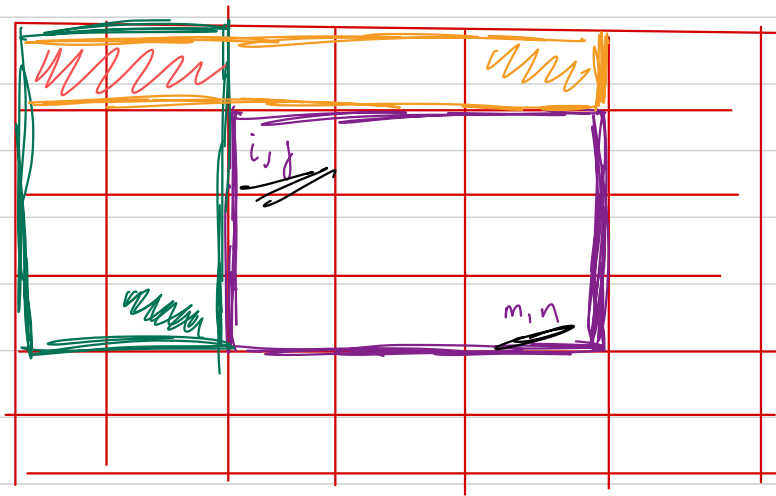
Kadane's algorithm

↓
max sum subarray
in a 1D array

→ ans

$$\underline{\underline{O(n^2)}}$$

Prefer
Sum



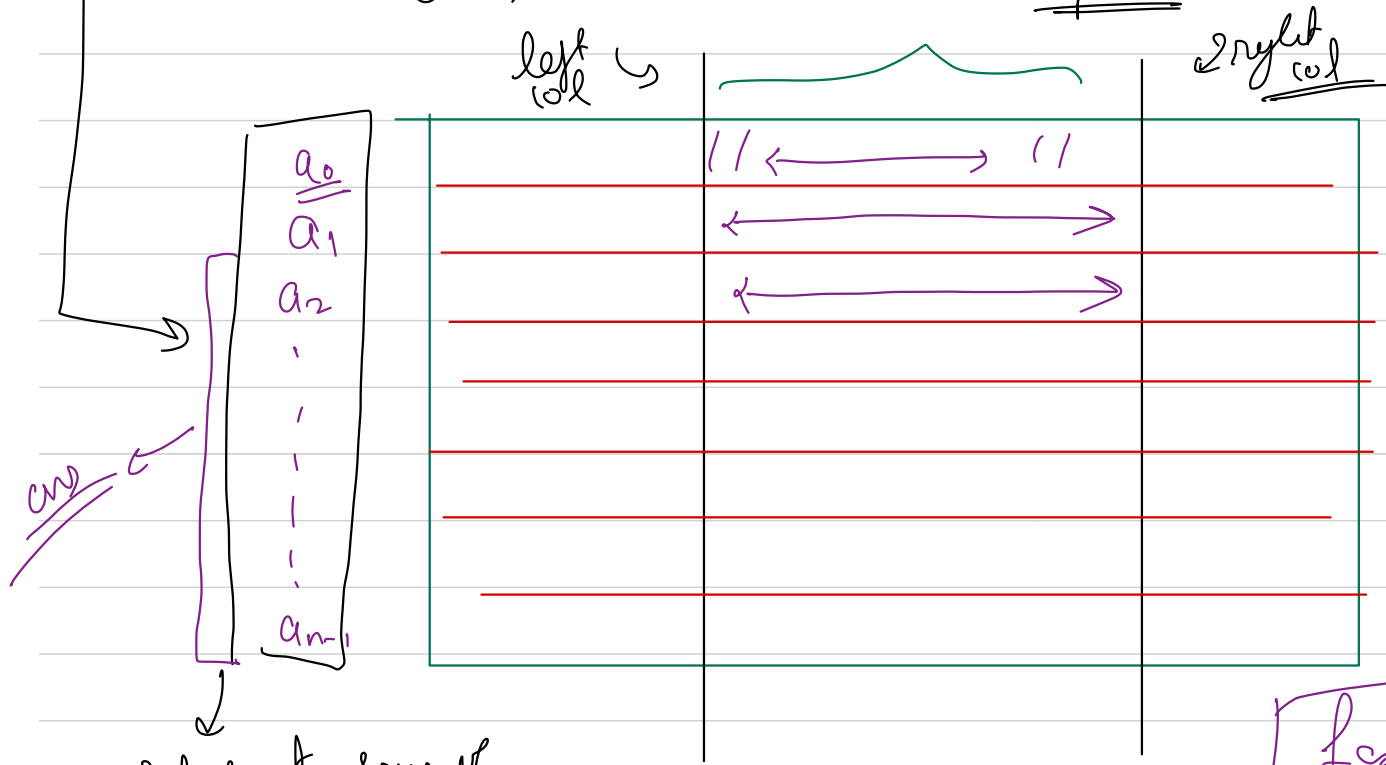
$$\boxed{O(n^4)}$$

Sum Purple - green - yellow + orange

prefer sum grid $\rightarrow \underline{\underline{O(n^2)}}$

Can we optimize \leq ?

Kadane's algo → maximum sum subarray in a 1D array
 $O(n)$ time $O(1)$ space



n^2
 pairs of
 columns

represent sum of
 row elements of i^{th} row
 between
 left & right col

for all pairs of
 left & right
 columns

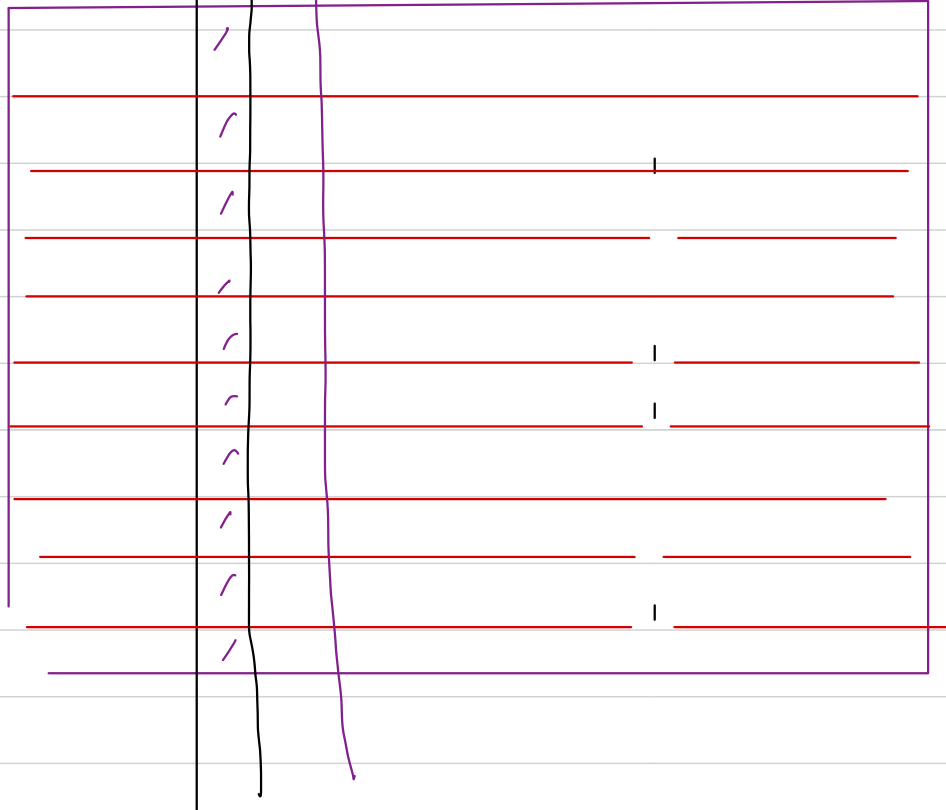
$l=1$

$r=2$

$r=3$

$r=4$

$$n^2 \times (n+n)$$



radius

C_n

2

$$n^6 \rightarrow n^4 \rightarrow n^3 \leftarrow$$

n^3

$$n^2(n+n)$$

↓
general
column
per

$O(n^3)$