

*Readings* Part III (Intro), Sections 12.1, 12.2, 12.3  
*Self test* Exercises 12.2-3, 12.3-1

## Lecture 05

### Dictionaries

keys: 5 7 2 0 20 4 9

**Objects:=** Sets  $s$  where each element  $x$  has field  $x.key$  - some totally ordered value. Keys are distinct.

#### Operations:

- $\text{search}(S, k)$ : return  $x \in S$  s.t.  $x.key = k$ , or NIL if no such  $x$
- $\text{delete}(S, x)$ : remove  $x$  from  $S$  (given element  $x \in S$  not just  $x.key$  or the values in the element)  
 Q: Why not  $\text{delete}(S, k)$ ?  
 A: Achieved with  $\text{delete}(S, \text{search}(S, k))$ , separates “searching phase” from “deletion” phase, making it possible to analyse each one separately  
 Notice that  $\text{delete}$  assumes that we not only know the element  $x$ ’s values but that we have a pointer to the actual element  $x$  in the data structure.
- $\text{insert}(S, x)$ : insert  $x$  in  $S$ ; if some  $y \in S$  has  $y.key = x.key$ , replace  $y$  by  $x$

5	7	2	8	20	4	9
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 size[7]

Complexity of search, insert for unsorted array  $\in \Theta(n)$ , delete  $\in \Theta(1)$

0	2	4	5	7	9	30
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 size[7]

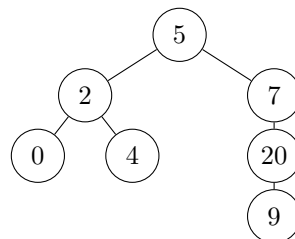
Complexity of search (binary)  $\in \Theta(\log n)$ , insert, delete  $\in \Theta(n)$

head  $\rightarrow 5 \leftrightarrow 7 \leftrightarrow 2 \leftrightarrow 0 \leftrightarrow 20 \leftrightarrow 4 \leftrightarrow 9$

Complexity of search, insert  $\in \Theta(n)$ , delete  $\in \Theta(1)$

head  $\rightarrow 0 \leftrightarrow 2 \leftrightarrow 4 \leftrightarrow 5 \leftrightarrow 7 \leftrightarrow 9 \leftrightarrow 20$

Complexity of search, insert  $\in \Theta(n)$ , delete  $\in \Theta(1)$



Complexity of search  $\in O(h)$

## Summary

Data Structure	search	insert	delete
unsorted array	$n$	$n$	1
sorted array	$\log n$	$n$	$n$
unsorted singly-linked list	$n$	$n$	$n$
unsorted doubly-linked list	$n$	$n$	1
sorted doubly-linked list	$n$	$n$	1
binary search tree	$n$	$n$	$n$
balanced search tree	$\log n$	$\log n$	$\log n$
direct-access table(+)	1	1	1
hash table	$n$	$n$	$n$

(+)  $\rightarrow$  all have space complexity  $O(n)$  except direct-access tables, which require space equal to size of universe, e.g., if keys are 32-bit integers, a direct-access table requires space  $\Omega(2^{32})$ .