Lecture 20

Minimum Spanning Trees (cont.)

Kruskal's MST Algorithm

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kruskal(G=(V,E), w:E\to R): T\leftarrow \{\} sort edges so w(e_1)\leq w(e_2)\leq \ldots \leq w(e_n) for i\leftarrow 1 to n \text{ \# let } (u_i,v_i)=e_i if u_i,v_i in different connected components of T: T\leftarrow T\cup \{e_i\}
```

Disjoint Set ADT

Objects: Collection of non-empty sets

$$s = \{s_1, s_2, \dots, s_k\}$$

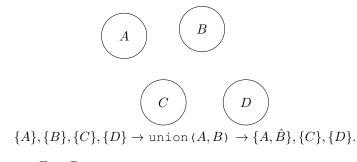
where each s_i is a non-empty set that has no element in common with any other s_i . That is, $s_i \cap s_j = \{\}$. Each set is identified by a unique element called its "representative".

Operations:

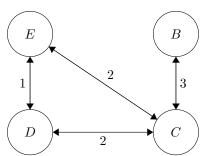
make-set (x): Given element x that does not already belong to one of the sets, create a new set $\{x\}$ that contains only x (and assign x as the representative).

find-set (x): Given element x, return the representative of the set that contains x (or NIL if x not in a set).

union (x,y): Given two distinct elements x and y, let s_x be the set that contains x and s_y be the set that contains y. Form a new set consisting of $s_x \cup s_y$, and remove s_x and s_y from the collection. Pick a representative for the new set.



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 \begin{aligned} & \text{kruskal} \ (G = (V, E), \ w : E \to R) : \\ & T \leftarrow \{\} \\ & \text{sort edges so} \ w(e_1) \leq w(e_2) \leq \ldots \leq w(e_n) \\ & \text{for} \ v \in V : \\ & \text{makeset} \ (v) \\ & \text{for} \ i \leftarrow 1 \ \text{to} \ |E| \\ & \text{\#} \ \text{let} \ (u_i, v_i) = e_i \\ & \text{if find-set} \ (u_i) \ ! = \ \text{find-set} \ (v_i) : \\ & T \leftarrow T \cup \{e_i\} \\ & \text{if} \ u_i, v_i \ \text{in different connected components of} \ T : \\ & \text{union} \ (u_i, v_i) \end{aligned}
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 $\{\hat{C}\}, \{\hat{B}\}, \{\hat{D}\}, \{\hat{E}\} \to \{\hat{C}\}, \{\hat{B}\}, \{E, \hat{D}\} \to \{B\}, \{\hat{D}, C, E\} \to \{\hat{B}, D, C, E\}$