

AVL Tree

1. It is a BST

2. Height of left subtree - Height of right subtree
 - 9 - 1, 0, 1, 2

Balance factor

If the ~~is~~ balanced factor
 is greater than or less

than this, then (4) case arises

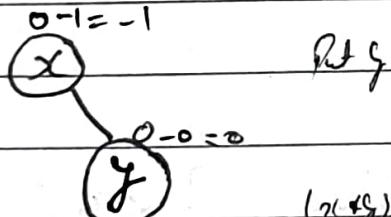
left Rotation

\downarrow x, y, z

~~Left Case I~~
 Left Rotation

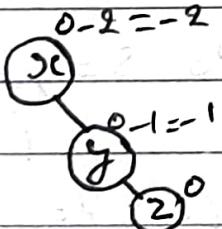
(2)

- no left and right height
 left-height of right = 0 (Balanced)



(x,y) both
balanced

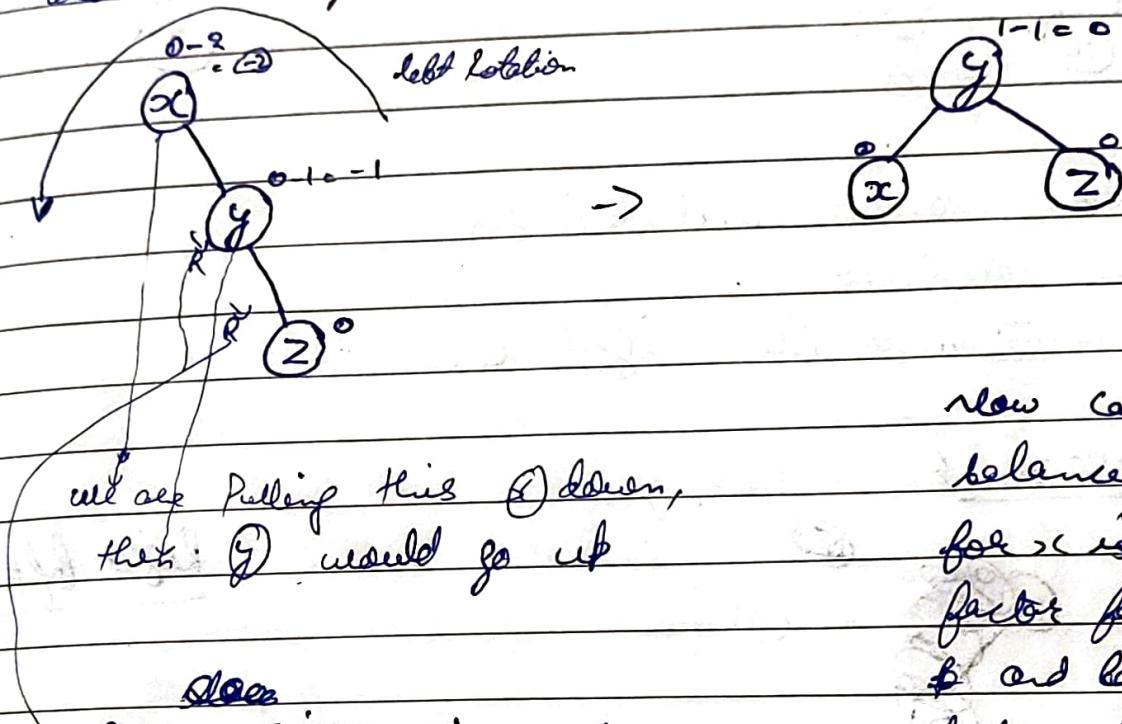
- now left and right subtree
 height of left subtree = $0 + 1$ so height of left subtree -
 height of right subtree = $0 - 1 = -1$ (Balanced)



- $y + z$ are balanced factor
 but x is not balanced because

height of left subtree is 0 & height of right subtree is of x is 2 $\Rightarrow 0-2=-2$ which not satisfies the Balanced factor Condition.

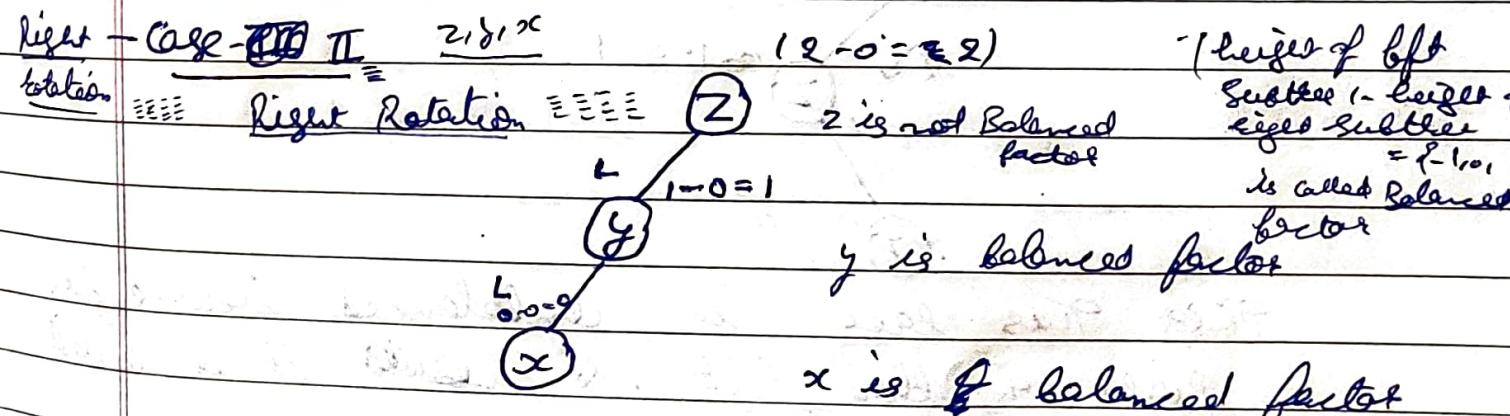
now we do left rotation, note that (left rotation) always the middle element go to root



we see putting this y down,
then y would go up

now calculate balanced factor
for x is 0, balance
factor for z is 0
& and balanced
factor for y is
 $1-1=0$

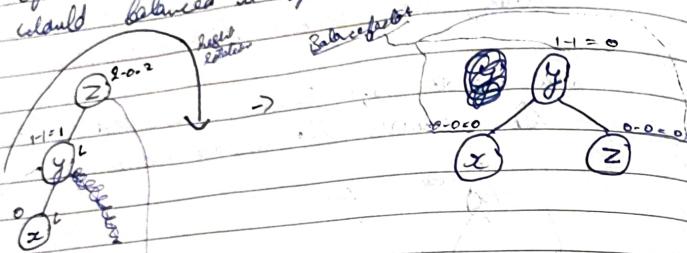
~~Right Right~~ Right Right ki wajah se
ha ~~is~~ unbalanced, isles
& then rotation should be
left rotation.



x is ~~not~~ balanced factor

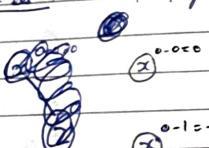
~~This is~~ This tree is unbalanced because of

left insertion and again left inserted, so we would balance it by left rotation.

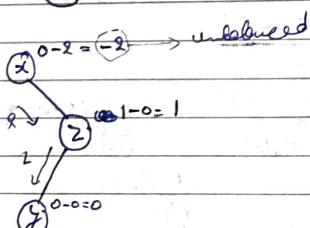


medium element will become root
pull \textcircled{z} into the right one

Case III : x, z, y



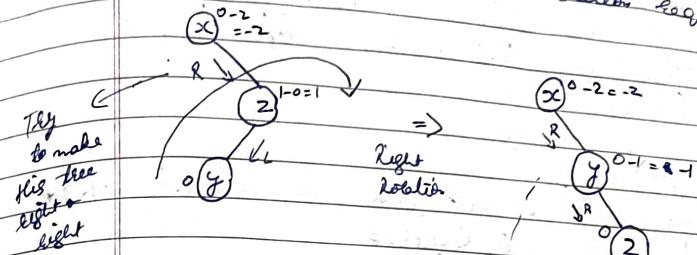
Right Left Rotation



This tree is unbalanced because first right + then left, so called it is called right-left rotation.

III

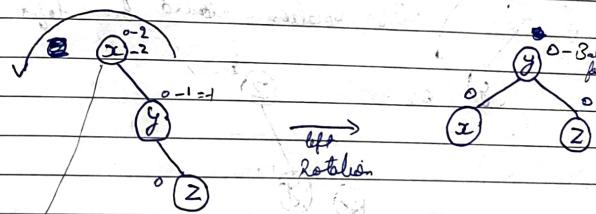
In Case I & Case II we done only 1 rotation left + right but in Case III, 2 rotations required



try to make his tree light + light

This tree is still unbalanced because \textcircled{x} do

This tree is unbalanced balanced because light + light. so we have to perform left rotation in this tree.



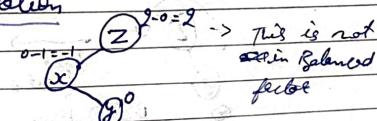
Pull \textcircled{x} down, then

\textcircled{y} would go up, medium element (\textcircled{y}) become root

Case IV

left light
right rotation.

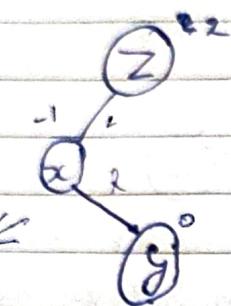
left Right Rotation



x, z, y

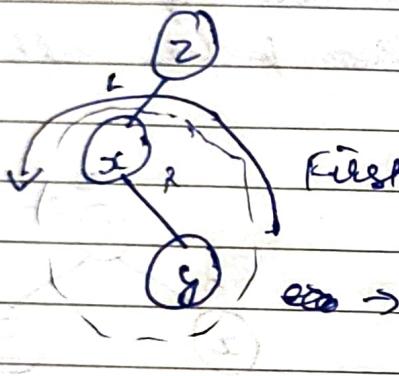
$z = 0 = 2 \rightarrow$ This is not
unbalanced

This tree is unbalanced because of first left insertion is there and then right insertion is there, left and then right this is like left-right rotation.

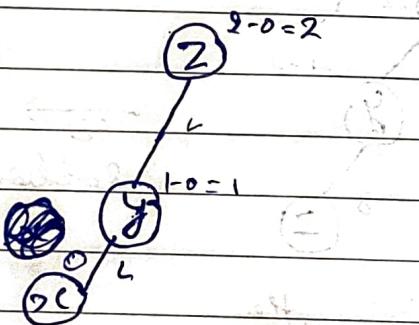


In this case, 2 rotations would be required

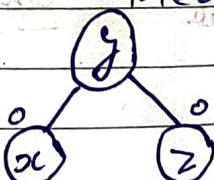
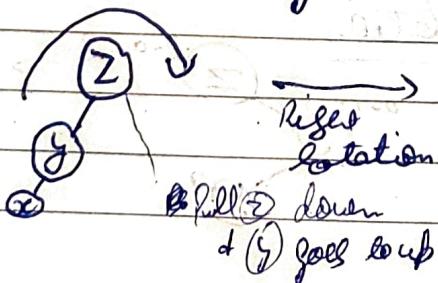
try to make this tree left-left balanced



First rotation could be left rotation



Again unbalanced, if tree is unbalanced because of left + left, so we could balance it by right rotation.



This 4 types of conditions could be used to balance out AVL tree

1, Left Rotation

There are 2 reasons when tree is unbalanced because of two right insertions

2, Right Rotation

There are 2 reasons when tree is unbalanced because of two left insertions

3, ~~Right~~ Right - Left Rotation

when tree is unbalanced due to first right and then left insertion

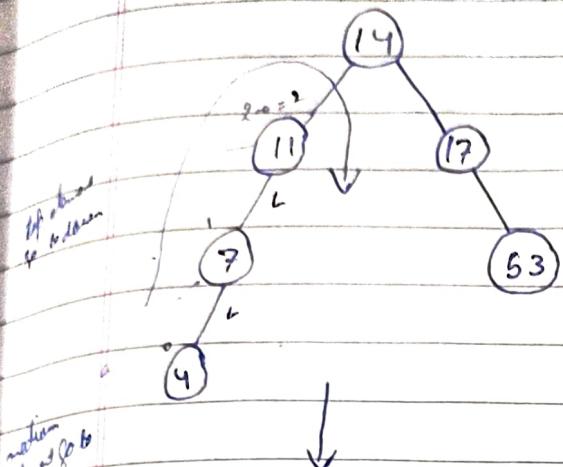
4, Left - ~~Left~~ Right Rotation

when tree is unbalanced due to first left + then right right insertion

68 larger - right larger
B.F or Balance factor = $-1, 0, +1$

Construct AVL tree by inserting the following data

14, 17, 11, 7, 53, 4, 13, 12, 8, 60, 19, 16, 20

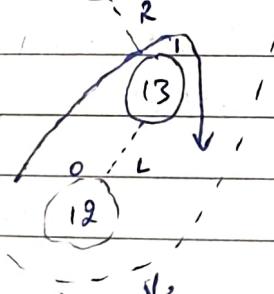
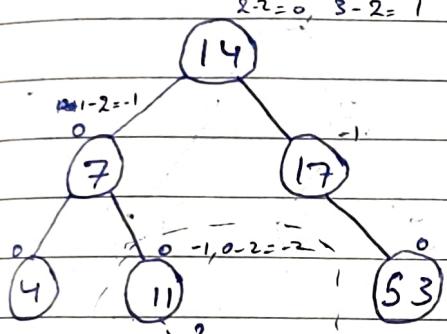


first element is root

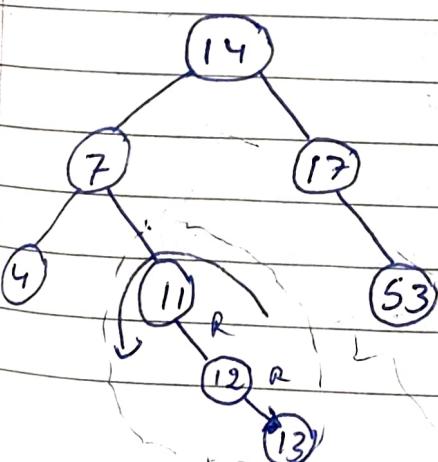
$17 > 14$ balance factor is $0 + 1 = 1$
also balance factor is 0
 $B.F \text{ of } 7 \text{ is } 0, \text{ now } B.F \text{ of } 11 \text{ is } 1 +$
 $B.F \text{ of } 14 \text{ is } 2 - 1 = 1$

$B.F \text{ of } 4 \text{ is } 0, B.F \text{ of } 7 \text{ is } 1, B.F \text{ of }$
 $11 \text{ is } 2 - 0 = 2$ so this is not balanced tree.

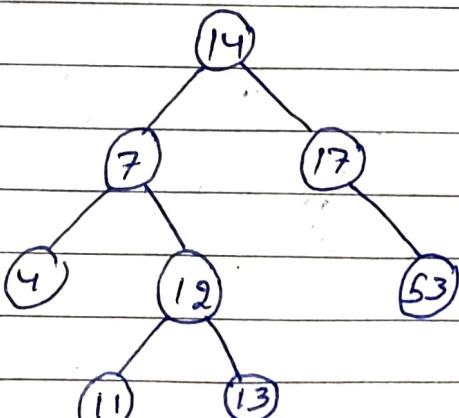
This tree is unbalanced due to left + left rotation or LL rotation, we need to balance it by right rotation.



This tree is unbalanced due to right left rotation, first we make it right right + then we perform left rotation.

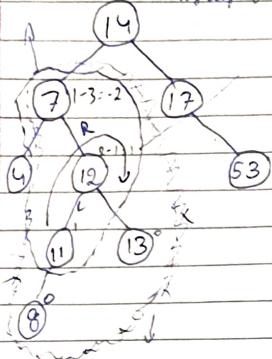


\rightarrow



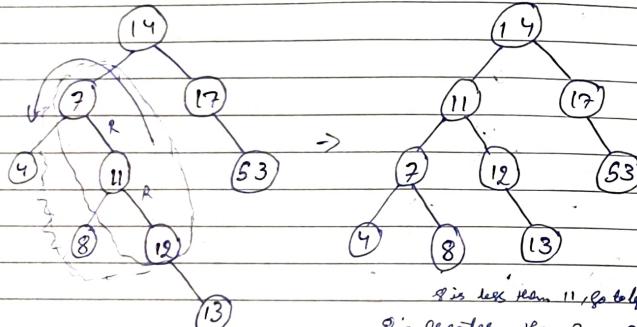
This is our critical node

initially S

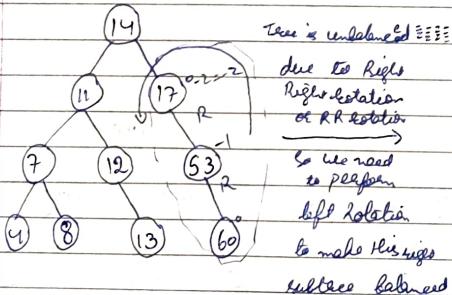


Check check the median of this node is 7, 11, 12 inserted order
11 go to the left of 17 & 18 are 13, 14

Tree is also unbalanced due to Right Left Rotation or RL rotation, we can balance it by first making it a right Right rotation + then perform left rotation



8 is less than 11, so go to left of 11 &
8 is greater than 7, so go to right of 7



Tree is unbalanced due to Right Right rotation or RR rotation

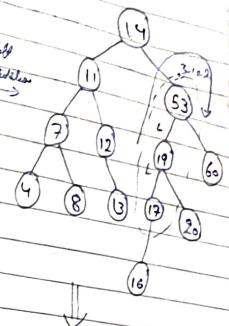
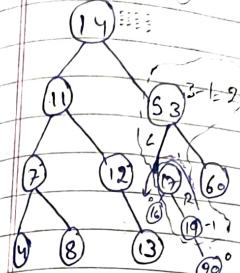
so we need to perform left rotation to make this right subtree balanced

11 go to the left of 17 & 18 are 13, 14

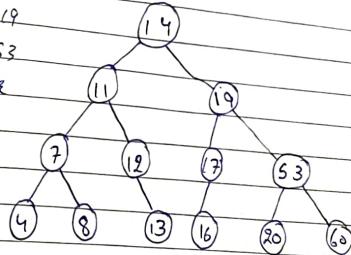


Tree is unbalanced

To left right rotation
try to make 14, 16
& then perform right's rotation



20 is present to the right of 19
but to the left of 19, 53
is present. 20 is lesser
than 53 so 20 go to
the left of 53



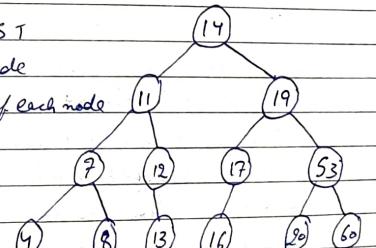
Deletion in A AVL Tree

1. Deletion same as in BST

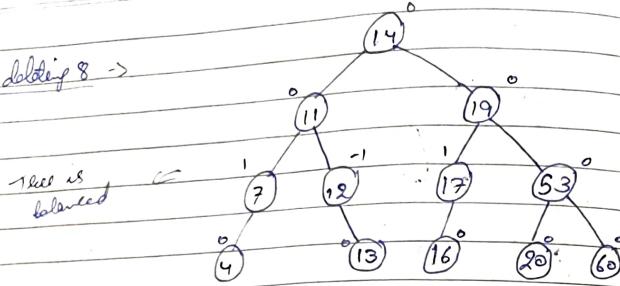
2. After deleting a node

check balance factor of each node
every time

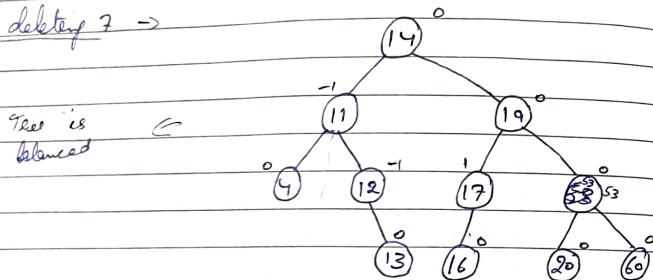
Delete - 8, 11, 14, 17



After deleting 8 →



After deleting 7 →



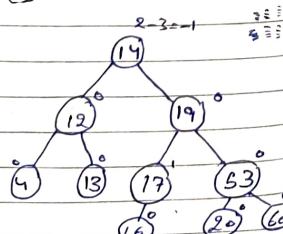
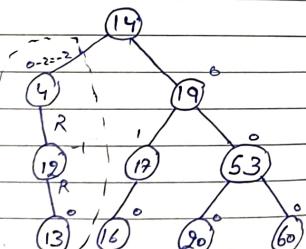
After deleting 11 →

Both left & right children of 11 is present, so we can delete 11 by inorder predecessor or inorder successor. We choose inorder predecessor.

Now check Balanced factor
Tree is not balanced.

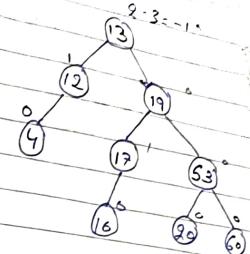
Tree is unbalanced due to right-left rotation or L-R rotation, so we can balance it by using left rotation.

Now again check balanced factor



After deleting 14 →

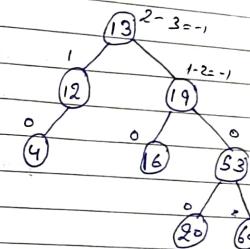
14 is having both left and right child, so we can delete 14 by inorder predecessor or inorder successor, so we choose inorder predecessor.



Now check balanced factor

After deleting 17 →

There is only one child present in 17, so we can easily replace 17 with 16.



Now check balanced factor

If tree is balanced we can delete new node, if tree is not balanced first we have to balance it & then we can delete a node.

Code for AVL Tree insertion

Binary Search Tree 14. java

```
int height (Node N) {
    if (N == null)
        return 0;
```

return N.height;

2

int max (int a, int b) {
 return (a > b) ? a : b;

3

Node rightRotate (Node y) {
 Node x = y.left;
 Node temp = x.right;

x.height = y;
y.left = temp;

// update heights

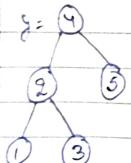
y.height = max (height (y.left), height (y.right)) + 1;

x.height = max (height (x.left), height (x.right)) + 1;

return x;

4

Subtract key's value

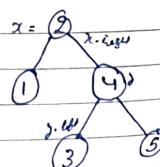


$$x = g(y.left)$$

$$\text{temp} = 3 \text{ (l1.right)}$$

$$x.l.height = 4 \text{ (y)}$$

$$y.left = 3 \text{ (temp)}$$



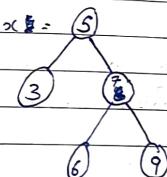
45 39
Node leftRotate (Node x) {
 Node y = x.right;
 Node temp = y.left;

y.left = x;
x.right = temp;
// update heights

x.height = max (height (x.left), height (x.right)) + 1;
y.height = max (height (y.left), height (y.right)) + 1;
// return new root
return y;

5

Explain left rotate



$$y = 7 \text{ (x.right)}$$

$$\text{temp} = 6 \text{ (7.left)}$$

$$y.left = 6 \text{ (temp)}$$

$$x.height = 6 \text{ (x)}$$

$$y = 7$$

$$x = 5$$

$$x = 5$$

$$x = 6$$

int getBalance (Node N) {

if (N == null)

return 0;

return height (N.left) - height (N.right);

6

Node insert (Node node, int key) {

if (!node == null)

return (new Node (key));

```

if (key < node.key)
    node.left = insert (node.left, key);
else if (key > node.key)
    node.right = insert (node.right, key);
else // Duplicate keys not allowed
    return node;
return node;

```

```

node.height = 1 + max (height (node.left),
                      height (node.right));

```

int balance = getBalance (node);
 // left left case

```

if (balance > 1 && key < node.left.key)
    return rightRotate (node);
else if (key > node.right.key)
    return leftRotate (node);

```

```

if (balance < -1 && key > node.right.key)
    return leftRotate (node);

```

// left right case

```

if (balance > 1 && key > node.left.key) {
    node.left = leftRotate (node.left);
    return rightRotate (node);
}

```

}

// Right left case

```

if (balance < -1 && key < node.right.key) {
    node.right = rightRotate (node.right);
    return leftRotate (node);
}

```

}

// return the (unchanged) node pointer
 return node;

}

AVL Tree - Deletion (with code)

Binary Search Tree 15. func

```

Node deleteNode (Node root, int key)
{

```

```

    if (root == null)
        return root;

```

```

    if (key < root.key)

```

```

        root.left = deleteNode (root.left, key);
    else if (key > root.key)

```

```

        root.right = deleteNode (root.right, key);
    else if (key == root.key)

```

```

        if (root.left == null)

```

```

            return root.right;
        else if (root.right == null)

```

```

            return root.left;
        else // in order successor (smaller
             // in right subtree)

```

```

            root.key = minValue (root.right);
            root.right = deleteNode (root.right, root.key);
        }
    }
}
```

```

    root.right = deleteNode (root.right, root.key);
}
```

}

```

    if (root == null)
        return root;

```

$\text{root.right} = \text{max}(\text{height}(\text{root.left}), \text{height}(\text{root.right})) + 1$

int balance = getBalance(root);

// left left case

if (balance > 1 && getBalance(root.left) >= 0)

return rightRotate(root);

// left right case

if (balance > 1 && getBalance(root.left) < 0) {

root.left = leftRotate(root.left);

return rightRotate(root);

3

// right right case

if (balance < -1 && getBalance(root.right) <= 0)

return leftRotate(root);

// right left case

if (balance < -1 && getBalance(root.right) > 0) {

root.right = rightRotate(rightRotate(root.right));

return leftRotate(root);

4

return root;

7