Student Information

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Answer 1

a)
$$G_1 = \{V_1, \Sigma_1, R_1, S_1\}$$
, $V_1 = \{S_1, a, b\}$, $\Sigma_1 = \{a, b\}$ and
$$R = \{S_1 \rightarrow S_1 a S_1 b S_1 b S_1,$$

$$S_1 \rightarrow S_1 b S_1 a S_1 b S_1,$$

$$S_1 \rightarrow S_1 b S_1 b S_1 a S_1,$$

$$S_1 \rightarrow e\}$$

b)
$$G_2 = \{V_2, \Sigma_2, R_2, S_2\}$$
, $V_2 = \{S_2, K, L, a, b\}$, $\Sigma_2 = \{a, b\}$ and $R_2 = \{S_2 \rightarrow K \mid L, K \rightarrow KaKaKaK, K \rightarrow KaKbKaK, K \rightarrow KbKaKaK, K \rightarrow L, L \rightarrow LaLbL, L \rightarrow e\}$

c) PDA that accepts the L₁ is:

Let $G = \{V, \Sigma, R, S\}$ and we can construct a PDA M such that L(G) = L(M). Let M = $(K, \Sigma, \Gamma, \Delta, s, F)$ where $K = \{q\}$ $F = \{q\}$ $\Sigma = \{a, b\}$ $\Gamma = \{ K, a, b \}$ $\Delta = \{((q,a,e),(q,K)),$ a, ε\K ((q,a,b),(q,a)),((q,a,K),(q,Ka)),a, b∖a ((q,b,e),(q,b)),a, K\Ka ((q,b,K),(q,a)),((q,b,a),(q,e))b, K∖a b, a∖ε b, ε\b

d) We define grammar G_3 for $L_3 = L_1 \cup L_2$.

$$L(G_3) = L(G_1) \cup L(G_2).$$

Firstly, we should check L₁ and L₂ have disjoint sets of non-terminals.

$$V_1 - \Sigma_1 = \{ S_1 \}$$

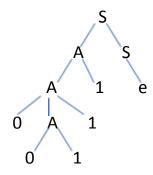
 $V_2 - \Sigma_2 = \{ S_2, K, L \}$

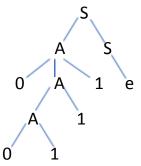
We can conclude their sets of non-terminals are disjoint.

According to the Union property of context-free languages (Theorem 3.5.1):

Answer 2

a) Grammars such as G', with strings that have two or more distinct parse trees, are called ambiguous. If we choose sample as "00111" we can draw 2 different parse trees. So, G_1 is ambiguous.





b) The unambiguous grammar for $L(G_1)$:

$$R = \{ S \rightarrow A$$

$$A \rightarrow AA \mid A1 \mid T$$

$$T \rightarrow 0T1 \mid 01 \}$$

c) The leftmost derivation of the string "00111":

$$\mathsf{S} \to \mathsf{A} \to \mathsf{A1} \to \mathsf{0T1} \to \mathsf{01}$$

Parse Tree:

