

Binary Trees

Trees

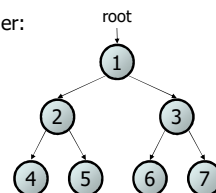
- **tree**: A directed, acyclic structure of linked nodes.
 - *directed*: Has one-way links between nodes.
 - *acyclic*: No path wraps back around to the same node twice.
 - *binary tree*: One where each node has at most two children.

- A **binary tree** can be defined as either:

- empty (`null`), or
- a **root** node that contains:

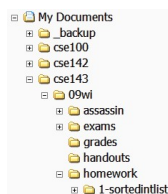
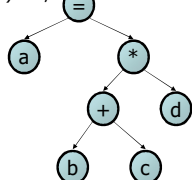
- **data**,
- a **left** subtree, and
- a **right** subtree.

– (The left and/or right subtree could be empty.)



Trees in computer science

- folders/files on a computer
- family genealogy; organizational charts
- AI: decision trees
- compilers: parse tree
 - $a = (b + c) * d;$



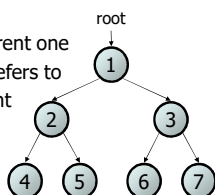
Programming with trees

- Trees are a mixture of linked lists and recursion
 - considered very elegant (perhaps beautiful!) by Ceng nerds
 - difficult for novices to master
- Common student remark #1:
 - "My code doesn't work, and I don't know why."
- Common student remark #2:
 - "My code works, and I don't know why."

Terminology

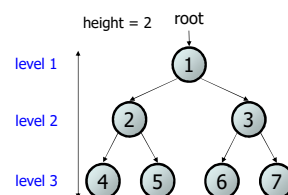
- **node**: an object containing a data value and left/right children
- **root**: topmost node of a tree
- **leaf**: a node that has no children
- **branch**: any internal node; neither the root nor a leaf

- **parent**: a node that refers to the current one
- **child**: a node that the current node refers to
- **sibling**: a node with a common parent



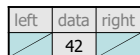
Terminology (cont.)

- **subtree**: the tree of nodes reachable to the left/right from the current node
- **height**: length of the longest path to a leaf from the given node
- **depth**: length of the path from the root to a given node
- **full tree**: one where every branch has 2 children

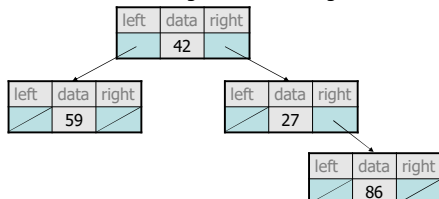


A tree node for integers

- A basic **tree node object** stores data and refers to left/right



- Multiple nodes can be linked together into a larger tree



IntTreeNode class

```
// An IntTreeNode object is one node in a binary tree of ints.
class IntTreeNode {
public:
    int data;           // data stored at this node
    IntTreeNode *left;  // reference to left subtree
    IntTreeNode *right; // reference to right subtree

    // Constructs a leaf node with the given data.
    IntTreeNode(int val) {
        data = val;
        left = nullptr;
        right = nullptr;
    }

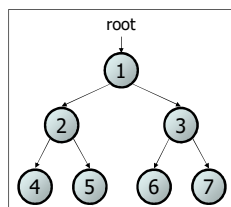
    // Constructs a branch node with the given data and links.
    IntTreeNode(int val, IntTreeNode *l, IntTreeNode *r) {
        data = val;
        left = l;
        right = r;
    }
}
```

IntTree class

```
// An IntTree object represents an entire binary tree of ints.
class IntTree {
private:
    IntTreeNode *root; // null for an empty tree

public:
    methods
}
```

- Client code talks to the `IntTree`, not to the node objects inside it
- Methods of the `IntTree` create and manipulate the nodes, their data and links between them

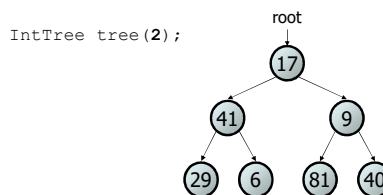


IntTree constructor

- Assume we have the following constructors:

```
IntTree(IntTreeNode *r)
IntTree(int height)
```

- The 2nd constructor creates a tree and fills it with nodes with random data values from 1..100 until it is full at the given height.

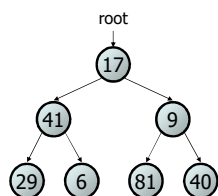


Printing a tree

- Add a method `print` to the `IntTree` class that prints the elements of the tree, separated by spaces.
- A node's left subtree should be printed before it, and its right subtree should be printed after it.

Example: `tree.print();`

29 41 6 17 81 9 40



Solution

```
// An IntTree object represents an entire binary tree of ints.
class IntTree {
public:
    void print() {
        print(root);
        cout << endl; // end the line of output
    }
    // other methods
private:
    IntTreeNode *root; // null for an empty tree

    void print(IntTreeNode *r) {
        // (base case is implicitly to do nothing on null)
        if (r != null) {
            // recursive case: print left, center, right
            print(r->left);
            cout << r->data << " ";
            print(r->right);
        }
    }
}
```

Style for tree methods

```
class IntTree {
public:
    type function_name(parameters) {
        function_name (root, parameters);
    }

private:
    IntTreeNode * root;
    type function_name (IntTreeNode *r, parameters) {
        ...
    }
};
```

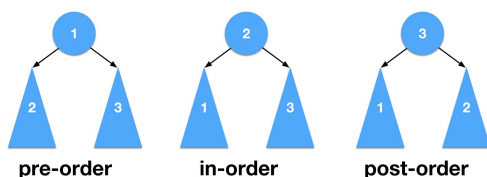
- Tree methods are often implemented recursively
 - with a public/private pair
 - the private version accepts the root node to process

Traversals

- traversal**: An examination of the elements of a tree.
 - A pattern used in many tree algorithms and methods
- Common orderings for traversals:
 - pre-order**: visit the *current* node, visit the left subtree, then visit the right subtree
 - in-order**: visit the left subtree, visit the *current* node, then visit the right subtree
 - post-order**: visit the left subtree, visit the right subtree, then visit the *current* node

Traversing a Binary Tree

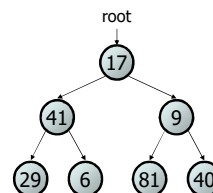
- Comparing the tree traversal methods:



(The numbers above refer to the order of traversal.)

- The subtrees are traversed **recursively**!

Traversal example

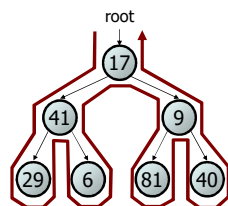


- pre-order: 17 41 29 6 9 81 40
- in-order: 29 41 6 17 81 9 40
- post-order: 29 6 41 81 40 9 17

Traversal trick

- To quickly generate a traversal:
 - Trace a path around the tree.
 - As you pass a node on the proper side, process it.

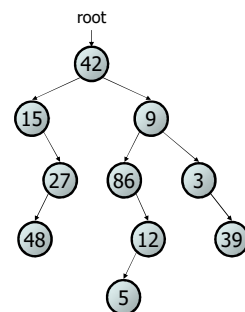
- pre-order: left side
- in-order: bottom
- post-order: right side



- pre-order: 17 41 29 6 9 81 40
- in-order: 29 41 6 17 81 9 40
- post-order: 29 6 41 81 40 9 17

Exercise 1

- Give pre-, in-, and post-order traversals for the following tree:



- Pre-order:
42 15 27 48 9 86 12 5 3 39
- In-order:
15 48 27 42 86 5 12 9 3 39
- Post-order:
48 27 15 5 12 86 39 3 42

Preorder traversal

```
void preorder(IntTreeNode *r) {
    if (r != nullptr) {
        cout << r->data << " ";
        preorder(r->left);
        preorder(r->right);
    }
}
```

Postorder Traversal

```
void postorder(IntTreeNode *r) {
    if (r != nullptr) {
        postorder(r->left);
        postorder(r->right);
        cout << r->data << " ";
    }
}
```

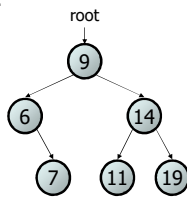
Exercise

- Add a method named `printSideways` to the `IntTree` class that prints the tree in a sideways indented format, with right nodes above roots above left nodes, with each level 4 spaces more indented than the one above it.

– Example: Output from the tree below:

```

      19
    14
  11
9
  7
  6
```



Exercise solution

```
// Prints the tree in a sideways indented format.
void printSideways() {
    printSideways(root, "");
}

void printSideways(IntTreeNode *r, string indent) {
    if (r != nullptr) {
        printSideways(r->right, indent + "    ");
        cout << indent << r->data << endl;
        printSideways(r->left, indent + "    ");
    }
}
```

Finding the maximum value in a binary tree

```
class IntTree {
public:
    ...
    int getMax () {
        return getMax (root);
    }
    ...
private:
    IntTreeNode * root;
    ...
    int getMax(IntTreeNode *r) {
        ...
    }
};
```

Finding the maximum value in a binary tree

```
int getMax(IntTreeNode *r) {
    int root_val, left, right, max;
    max = -1; // Assuming all values are positive integers
    if (r != nullptr) {
        root_val = r->data;
        left = getMax(r->left);
        right = getMax(r->right);
        // Find the largest of the three values.
        if (left > right)
            max = left;
        else
            max = right;
        if (root_val > max)
            max = root_val;
    }
    return max;
}
```

Adding up all values in a Binary Tree

```
int find_sum(IntTreeNode *r){
    if (r==nullptr)
        return 0;
    else
        return (r->data +
                find_sum(r->left) + find_sum(r->right) );
}
```

Exercise

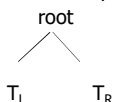
Add a method `count_leaves` to the `IntTree` class that counts the leaves of a binary tree.

```
public:
    int count_leaves () {
        return count_leaves (root);
    }
private:
    int count_leaves(IntTreeNode *r) {
        // TODO
    }
```

Height of Binary Tree

- The height of a binary tree T can be defined *recursively* as:

- If T is empty, its height is **-1**.
- If T is non-empty tree, then since T is of the form

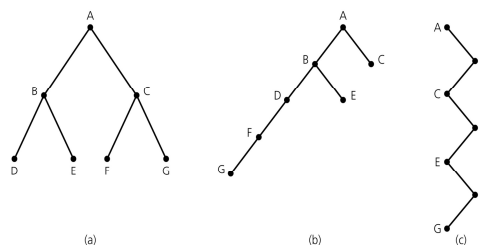


the height of T is 1 greater than the height of its root's taller subtree; i.e.

$$\text{height}(T) = 1 + \max\{\text{height}(T_L), \text{height}(T_R)\}$$

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Height of Binary Tree (cont.)



Binary trees with the same nodes but different heights

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Number of Binary trees with Same # of Nodes

$n=0 \rightarrow$ empty tree

$n=1 \rightarrow$ (1 tree)

$n=2 \rightarrow$ (2 trees)

$n=3 \rightarrow$ (5 trees)

In general:

$$\text{Catalan number } C(n) = (2n)! / ((n+1)!n!)$$

Different number of structurally different Binary trees is : Catalan(N)

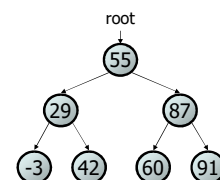
Different number of Binary Trees: $N!$ Catalan(N)

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Binary Search Trees

- binary search tree (BST)** is a binary tree that is either:

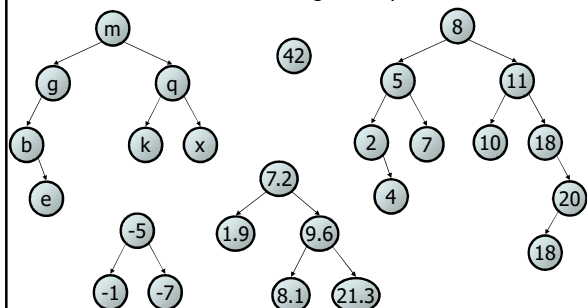
- empty (`null`), or
- a root node R such that:
 - every element of R 's left subtree contains data "less than" R 's data,
 - every element of R 's right subtree contains data "greater than" R 's,
 - R 's left and right subtrees are also binary search trees.



- BSTs store their elements in sorted order, which is helpful for searching/sorting tasks.

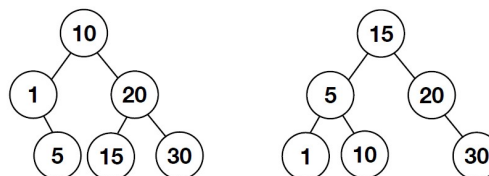
Exercise

- Which of the trees shown are legal binary search trees?



Inorder traversal of BST

- Let's work out the **in-order traversal** results of the following two valid BSTs.



- For both, in order traversal gives the same result: 1, 5, 10, 15, 20, 30. This is clearly sorted!

Hey! these are all different things

Please do not confuse them

- Binary Search:**
an algorithm on a sorted array.
- Binary Tree**
a tree where nodes have no more than 2 children.
- Binary Search Tree**
a binary tree with a special ordering property

Search in a BST

- However, **Binary Search** and **BST** are related, because the way you search in a BST is similar to performing a binary search in an ordered array.
- Find 31

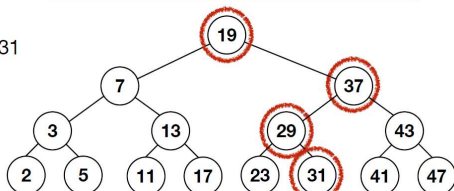
| | | | | | | | | | | | | | | |
|---|---|---|---|----|----|----|----|----|----|----|----|----|----|----|
| 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | 31 | 37 | 41 | 43 | 47 |
| 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | 31 | 37 | 41 | 43 | 47 |
| 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | 31 | 37 | 41 | 43 | 47 |
| 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | 31 | 37 | 41 | 43 | 47 |

Search in a BST

- Find 31

| | | | | | | | | | | | | | | |
|---|---|---|---|----|----|----|----|----|----|----|----|----|----|----|
| 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | 31 | 37 | 41 | 43 | 47 |
| 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | 31 | 37 | 41 | 43 | 47 |
| 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | 31 | 37 | 41 | 43 | 47 |
| 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | 31 | 37 | 41 | 43 | 47 |

- Find 31



What is the maximum number of nodes you would need to examine to perform any search?

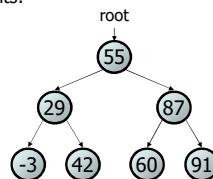
Search in a BST

- To summarize, you start from the root node, then choose to go left or right depending on the comparison result. The search ends when either you've found the target or you've reached a leaf.
- The maximum number of steps is the **tree height**.
- As in binary search, search in BST can achieve $O(\log N)$ time. However, this requires the BST to be balanced (i.e. the height should be small).
- If you have a poorly constructed BST (e.g. degenerated to a linked list), you won't get the $O(\log N)$ performance!

Binary Search Tree Class

- Convert the `IntTree` class into a `SearchTree` class.
 - The elements of the tree will constitute a legal binary search tree.
- Add a method `contains` to the `SearchTree` class that searches the tree for a given integer, returning `true` if found.
 - If a `SearchTree` variable `tree` referred to the tree below, the following calls would have these results:

- `tree.contains(29) → true`
- `tree.contains(55) → true`
- `tree.contains(63) → false`
- `tree.contains(35) → false`

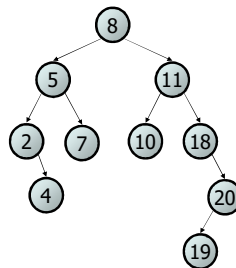


Method contains

```
// Returns whether this tree contains the given integer.
public:
    bool contains(int val){
        return contains(root, val);
    }
private:
    bool contains(IntTreeNode *r, int val){
        if (r == nullptr)
            return false;
        else {
            if (r->data == val)
                return true;
            else if (r->data > val)
                return contains(r->left, val);
            else return contains(r->right, val);
        }
    }
}
```

Adding to a BST

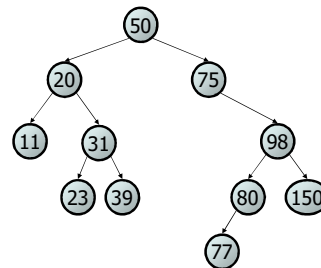
- Suppose we want to add the value 14 to the BST below.
 - Where should the new node be added?
- Where would we add the value 3?
- Where would we add 7?
- If the tree is empty, where should a new value be added?
- What is the general algorithm?



Adding exercise

- Draw what a binary search tree would look like if the following values were added to an initially empty tree in this order:

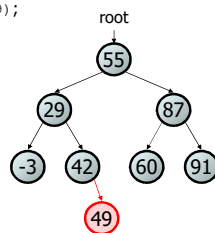
50
20
75
98
80
31
150
39
23
11
77



Implementing add

- Let's add a method `add` to the `SearchTree` class that adds a given integer value to the tree. Assume that the elements of the `SearchTree` constitute a legal binary search tree, and add the new value in the appropriate place to maintain ordering.

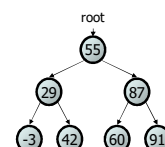
`tree.add(49);`



Code

```
// Adds the given value to this BST in sorted order.
public:
    void add(int value) {
        add(root, value);
    }
private:
    void add(IntTreeNode *&r, int value) {
        if (r == nullptr)
            r = new IntTreeNode(value);
        else if (r->data > value)
            add(r->left, value);
        else if (r->data < value)
            add(r->right, value);
        // else a duplicate
    }
}
```

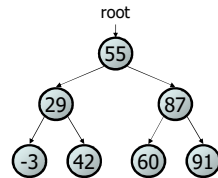
- Think about the case when `r` is a leaf..



Exercise

- Add a method `getMin` to the `IntTree` class that returns the minimum integer value from the tree. Assume that the elements of the `IntTree` constitute a legal binary search tree. Throw a `NoSuchElementException` if the tree is empty.

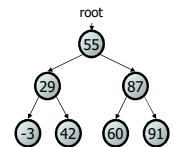
```
int min = tree.getMin(); // -3
```



Solution

```
// Returns the minimum value from this BST.
// Throws a NoSuchElementException if the tree is empty.
public:
int getMin() {
    if (root == nullptr)
        throw new NoSuchElementException();
    return getMin(root);
}

private:
int getMin(IntTreeNode* r) {
    if (r->left == nullptr)
        return r->data;
    else
        return getMin(r->left);
}
```



Find max: Iterative method

```
// Returns the largest value from this BST.
// Throws a NoSuchElementException if the tree is empty.
public:
int getMax() {
    return getMax(root);
}

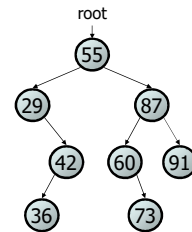
private:
int getMax(IntTreeNode *r){
    if (r == nullptr)
        throw new NoSuchElementException();
    while (r->right != nullptr)
        r = r->right;
    return r->data;
}
```

Removing from a BST

Possible cases for the node to be removed:

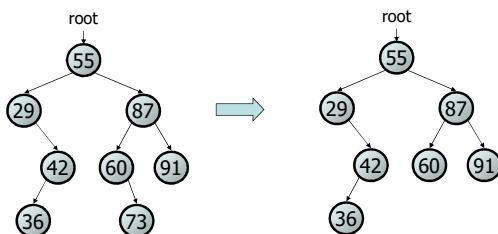
1. a leaf
2. a node with only one child (left or right child)
3. a node with both children

```
• tree.remove(73);
• tree.remove(29);
• tree.remove(42);
• tree.remove(55);
```



Case 1: Removing a Leaf Node

1. a leaf: replace with null

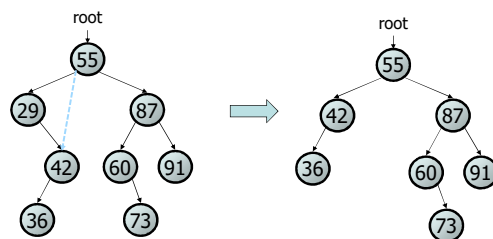


```
tree.remove(73);
```

Case 2: Remove a Node with one child

- 2.1 a node with a left child only:
- 2.2 a node with a right child only:

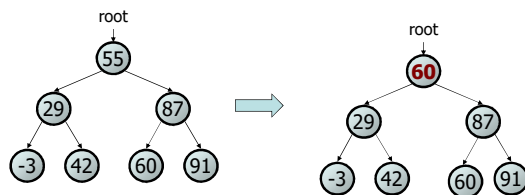
replace with left child
replace with right child



```
tree.remove(29);
```


Case 3: Remove a node with two children

3. a node with **both** children: replace with **min from right**



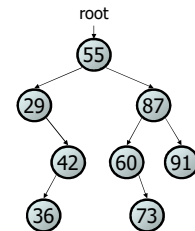
tree.remove(55);

remove method

- Add a method `remove` to the `IntTree` class that removes a given integer value from the tree, if present. Assume that the elements of the `IntTree` constitute a legal binary search tree, and remove the value in such a way as to maintain ordering.

```

• tree.remove(73);
• tree.remove(29);
• tree.remove(87);
• tree.remove(55);
  
```



remove method

```

// Removes the given value from this BST, if it exists.
public:
    void remove(int value) {
        remove(root, value);
    }
private:
    void remove(IntTreeNode *r, int value) {
        if (r == nullptr)
            return;
        else if (r->data > value)
            remove(r->left, value);
        else if (r->data < value)
            remove(r->right, value);
        else // r->data == value; remove this node
            if (r->left != nullptr && r->right != nullptr) {
                // case 3: both children; replace w/ min from R
                r->data = getMin(r->right); //copy value here
                remove(r->right, r->data);
            }
            else { // case 2: only child or case 1: leaf node
                IntTreeNode * oldNode = r;
                r = (r->left != nullptr)? r->left : r->right;
                delete oldNode;
            }
    }
}
  
```