

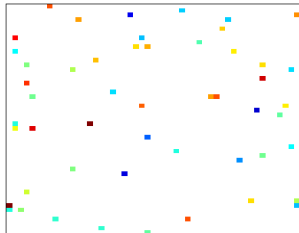
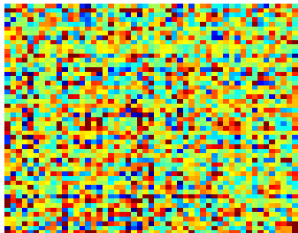
# Matrix Completion: Fundamental Limits and Efficient Algorithms

Sewoong Oh

PhD Defense  
Stanford University

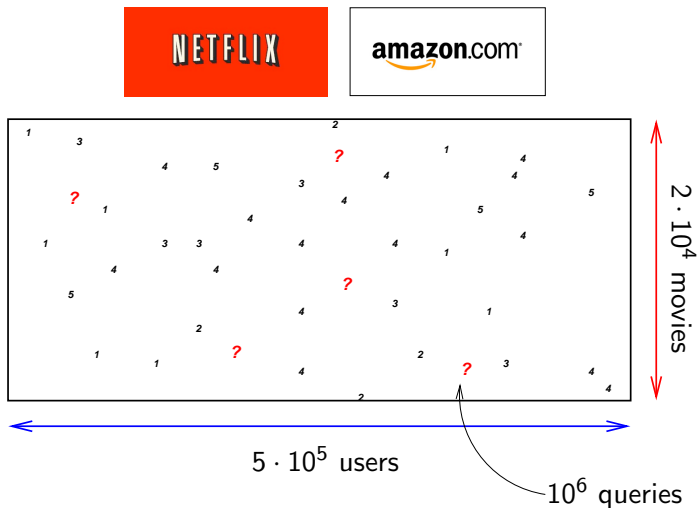
July 23, 2010

# Matrix completion



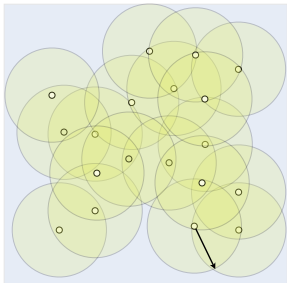
- Find the missing entries in a huge data matrix

## Example 1. Recommendation systems

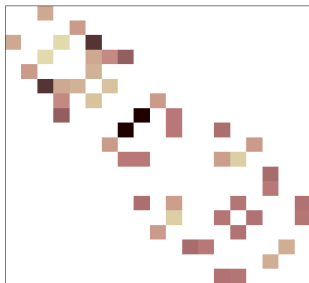


- Given less than 1% of the movie ratings
- Goal: Predict missing ratings

## Example 2. Positioning

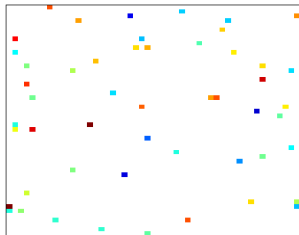
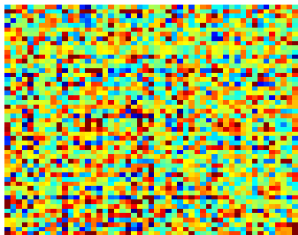


Distance Matrix



- Only distances between close-by sensors are measured
- Goal: Find the sensor positions up to a rigid motion

# Matrix completion



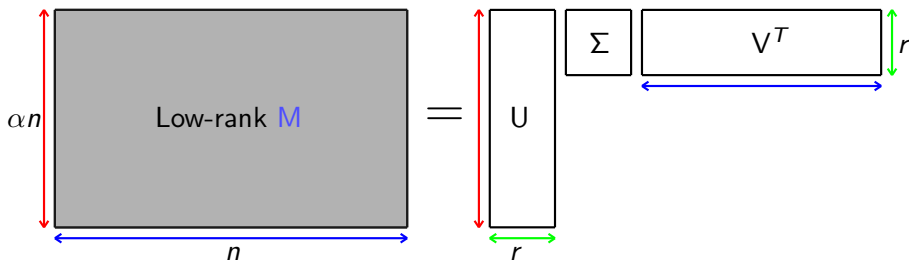
- More applications:
  - ▶ Computer vision: Structure-from-motion
  - ▶ Molecular biology: Microarray
  - ▶ Numerical linear algebra: Fast low-rank approximations
  - ▶ etc.

# Outline

- 1 Background
- 2 Algorithm and main results
- 3 Applications in positioning

## Background

# The model

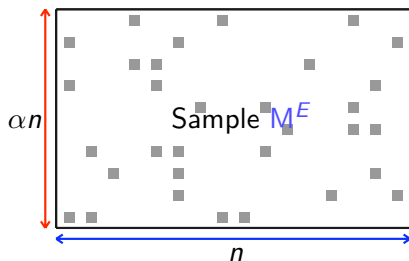


- Rank- $r$  matrix  $M$
- Random uniform sample set  $E$
- Sample matrix  $M^E$

$$M_{ij}^E = \begin{cases} M_{ij} & \text{if } (i,j) \in E \\ 0 & \text{otherwise} \end{cases}$$



# The model



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## Which matrices?

- Pathological example

$$M = \begin{bmatrix} 1 & 0 & \cdots & 0 \\ 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \cdot [1] \cdot [1 \quad 0 \quad \cdots \quad 0]$$

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- [Candès, Recht '08]  $M = U\Sigma V^T$  has coherence  $\mu$  if

$$A0. \quad \max_{1 \leq i \leq \alpha n} \sum_{k=1}^r U_{ik}^2 \leq \mu \frac{r}{n}, \quad \max_{1 \leq j \leq n} \sum_{k=1}^r V_{jk}^2 \leq \mu \frac{r}{n}$$

$$A1. \quad \max_{i,j} \left| \sum_{k=1}^r U_{ik} V_{jk} \right| \leq \mu \frac{\sqrt{r}}{n}$$

- Intuition

- ▶  $\mu$  is small if singular vectors are well balanced
- ▶ We need low-coherence for matrix completion

## Previous work

### Rank minimization

minimize     $\text{rank}(X)$   
subject to    $X_{ij} = M_{ij}, (i,j) \in E$

- NP-hard

# Previous work

## Rank minimization

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- NP-hard

## Heuristic [Fazel '02]

minimize  $\|X\|_*$   
subject to  $X_{ij} = M_{ij}, (i,j) \in E$

- Convex relaxation
- Nuclear norm

$$\|X\|_* = \sum_{i=1}^n \sigma_i(X)$$

- Can be solved using Semidefinite Programming(SDP)

## Previous work

- [Candès, Recht '08]

- ▶ Nuclear norm minimization reconstructs  $M$  **exactly** with high probability, if

$$|E| \geq C \mu r n^{6/5} \log n$$

- ▶ Surprise?

# Previous work

- [Candès, Recht '08]

- ▶ Nuclear norm minimization reconstructs  $M$  **exactly** with high probability, if

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- ▶ Degrees of freedom  $\simeq (1 + \alpha)rn$
- ▶ Open questions
  - ★ Optimality: Do we need  $n^{6/5} \log n$  samples?
  - ★ Complexity: SDP is computationally expensive
  - ★ Noise: Can not deal with noise

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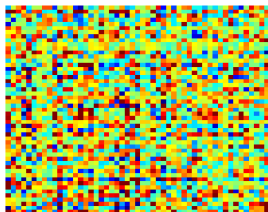
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A new approach to Matrix Completion: **OPTSPACE**

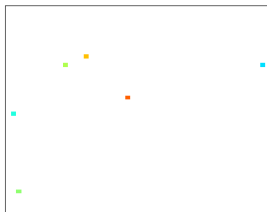


## Example: $2000 \times 2000$ rank-8 random matrix

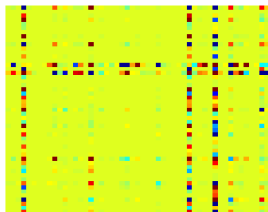
low-rank matrix  $M$



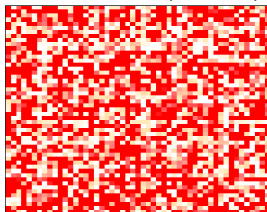
sampled matrix  $M^E$



OPTSPACE output  $\hat{M}$



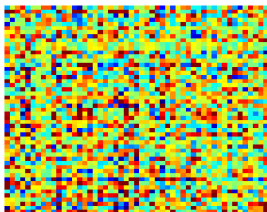
squared error  $(M - \hat{M})^2$



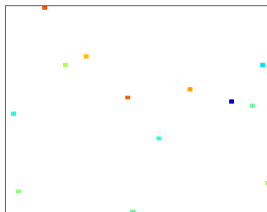
0.25% sampled

## Example: $2000 \times 2000$ rank-8 random matrix

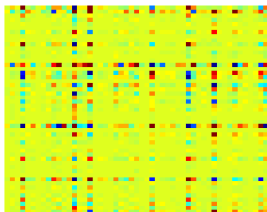
low-rank matrix  $M$



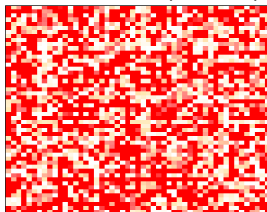
sampled matrix  $M^E$



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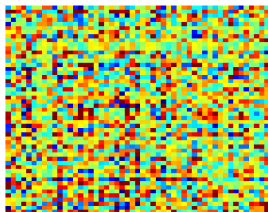
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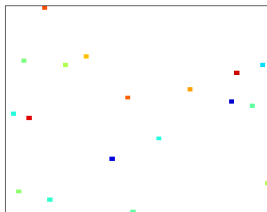
0.50% sampled

## Example: $2000 \times 2000$ rank-8 random matrix

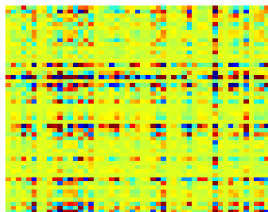
low-rank matrix  $M$



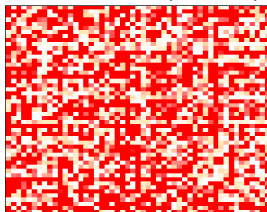
sampled matrix  $M^E$



OPTSPACE output  $\hat{M}$



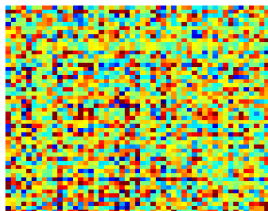
squared error  $(M - \hat{M})^2$



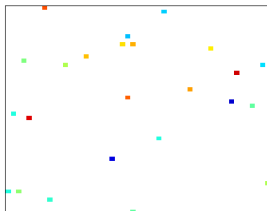
0.75% sampled

## Example: $2000 \times 2000$ rank-8 random matrix

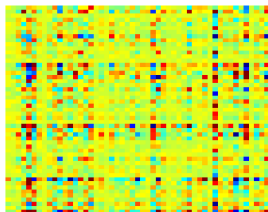
low-rank matrix  $M$



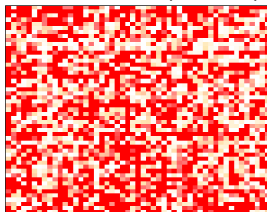
sampled matrix  $M^E$



OPTSPACE output  $\hat{M}$



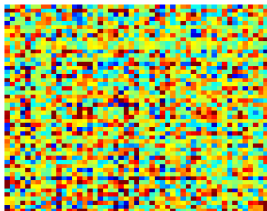
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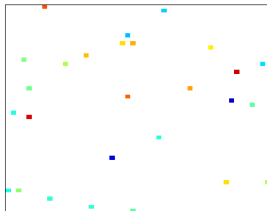
1.00% sampled

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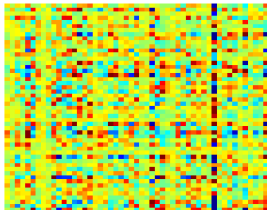
low-rank matrix  $M$



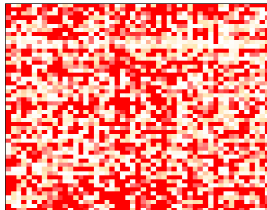
sampled matrix  $M^E$



OPTSPACE output  $\hat{M}$



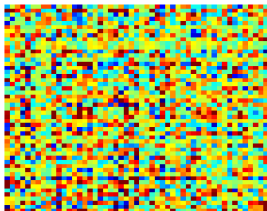
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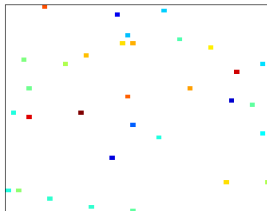
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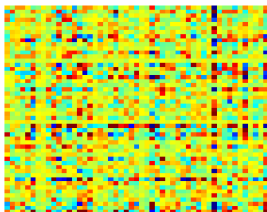
low-rank matrix  $M$



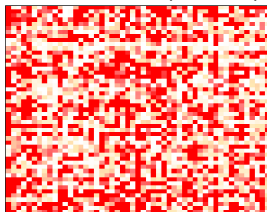
sampled matrix  $M^E$



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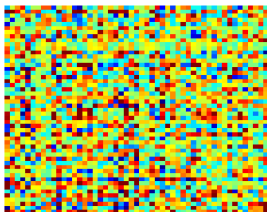
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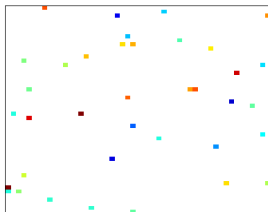
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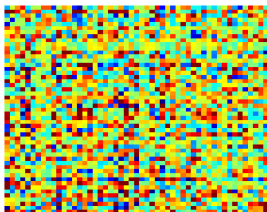
low-rank matrix  $M$



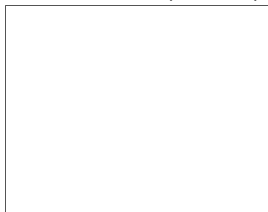
sampled matrix  $M^E$



OPTSPACE output  $\hat{M}$



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1.75% sampled

# Algorithm



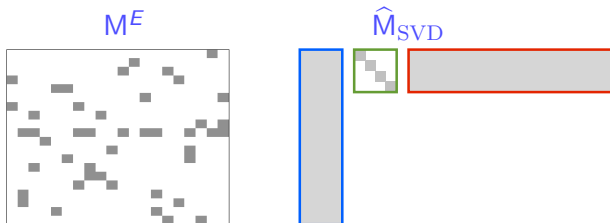
## Naïve approach

- Singular Value Decomposition (SVD)

$$M^E = \sum_{i=1}^n \sigma_i x_i y_i^T$$

- Compute rank- $r$  approximation  $\hat{M}_{\text{SVD}}$

$$\hat{M}_{\text{SVD}} \triangleq \frac{\alpha n^2}{|E|} \sum_{i=1}^r \sigma_i x_i y_i^T$$



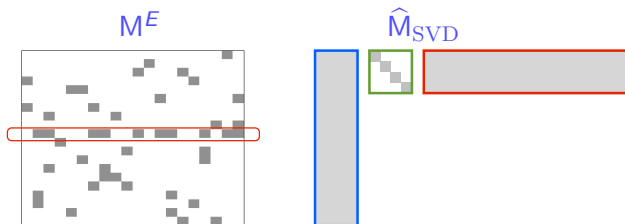
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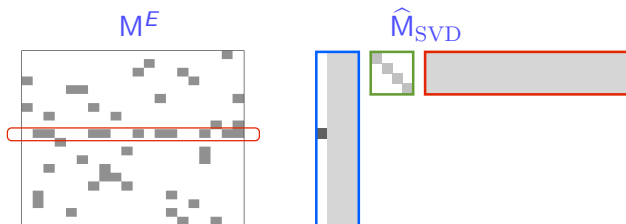
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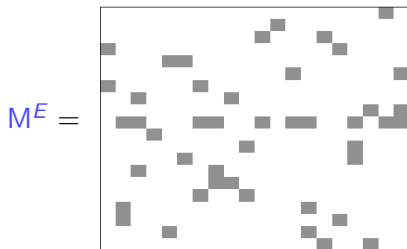
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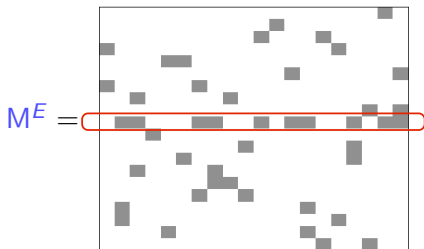
# Trimming



$$\tilde{M}_{ij}^E = \begin{cases} 0 & \text{if } \deg(\text{row}_i) > 2|E|/\alpha n \\ 0 & \text{if } \deg(\text{col}_j) > 2|E|/n \\ M_{ij}^E & \text{otherwise} \end{cases}$$

$\deg(\cdot)$  is the number of samples in that row/column

# Trimming



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# Algorithm

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## OPTSPACE

---

**Input :** sample indices  $E$ , sample values  $M^E$ , rank  $r$

**Output :** estimation  $\hat{M}$

- 1: Trimming
  - 2: Compute  $\hat{M}_{\text{SVD}}$  using SVD
  - 3: Greedy minimization of the residual error
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- $\hat{M}_{\text{SVD}}$  can be computed efficiently for sparse matrices



# Main results

## Theorem

For any  $|E|$ ,  $\hat{M}_{\text{SVD}}$  achieves, with high probability,

$$\text{RMSE} \leq CM_{\max} \sqrt{\frac{nr}{|E|}}$$

- $\text{RMSE} = \left( \frac{1}{\alpha n^2} \sum_{i,j} (M - \hat{M}_{\text{SVD}})_{ij}^2 \right)^{1/2}$
- $M_{\max} \triangleq \max_{i,j} |M_{ij}|$

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$$\text{RMSE} \leq CM_{\max} \sqrt{\frac{nr}{|E|}}$$

- [Achlioptas, McSherry '07]

If  $|E| \geq n(8 \log n)^4$ , with high probability,

$$\text{RMSE} \leq 4M_{\max} \sqrt{\frac{nr}{|E|}}$$

- For  $n = 10^5$ ,  $(8 \log n)^4 \simeq 7.2 \cdot 10^7$

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## Netflix dataset

A single user rated 17,000 movies.

“Miss Congeniality”: 200,000 ratings.

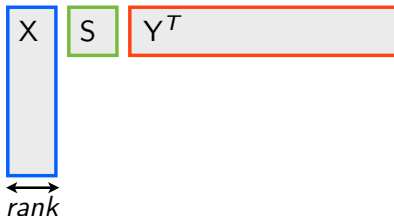
Can we do better?

## Greedy minimization of residual error

- Starting from  $(X_0, Y_0)$  for  $\hat{M}_{\text{SVD}} = X_0 S_0 Y_0^T$ , use gradient descent methods to solve

$$\begin{array}{ll} \text{minimize} & F(X, Y) \\ \text{subject to} & X^T X = \mathbb{I}, Y^T Y = \mathbb{I} \end{array}$$

$$F(X, Y) \triangleq \min_{S \in \mathbb{R}^{r \times r}} \sum_{(i,j) \in E} \left( M_{ij}^E - (XSY^T)_{ij} \right)^2$$



- Can be computed efficiently for sparse matrices

# Algorithm

---

## OPTSPACE

---

**Input** : sample indices  $E$ , sample values  $M^E$ , rank  $r$

**Output** : estimation  $\hat{M}$

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# Main results

## Theorem (Trimming+SVD)

$\hat{M}_{\text{SVD}}$  achieves, with high probability,

$$\text{RMSE} \leq CM_{\max} \sqrt{\frac{nr}{|E|}}$$

## Theorem (Trimming+SVD+Greedy minimization)

OPTSPACE reconstructs  $M$  exactly, with high probability, if

$$|E| \geq C \mu r n \max\{\mu r, \log n\}$$

# OPTSPACE is order-optimal

## Theorem

*If  $\mu$  and  $r$  are bounded, OPTSPACE reconstructs  $M$  exactly, with high probability, if*

$$|E| \geq C n \log n$$

- Lower bound (coupon collector's problem):  
If  $|E| \leq C' n \log n$ , then exact reconstruction is impossible



# OPTSPACE is order-optimal

## Theorem

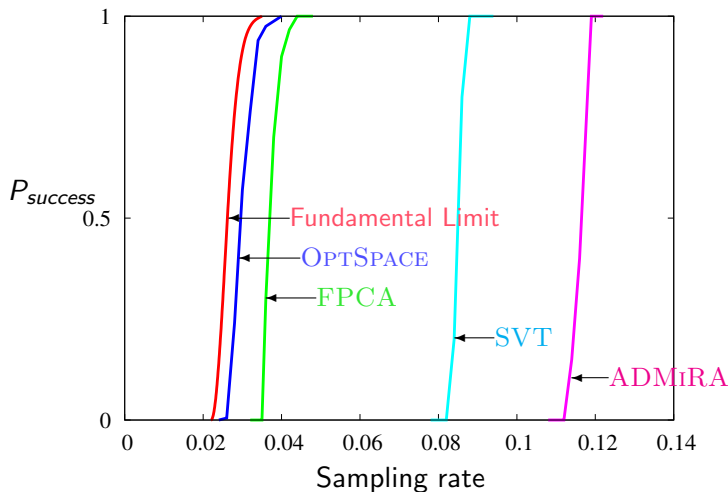
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- Lower bound (coupon collector's problem):  
If  $|E| \leq C' n \log n$ , then exact reconstruction is impossible
- Nuclear norm minimization:  
[Candès, Recht '08, Candès, Tao '09, Recht '09, Gross et al. '09]  
If  $|E| \geq C'' n (\log n)^2$ , then exact reconstruction by SDP

# Comparison

- $1000 \times 1000$  rank-10 matrix  $M$

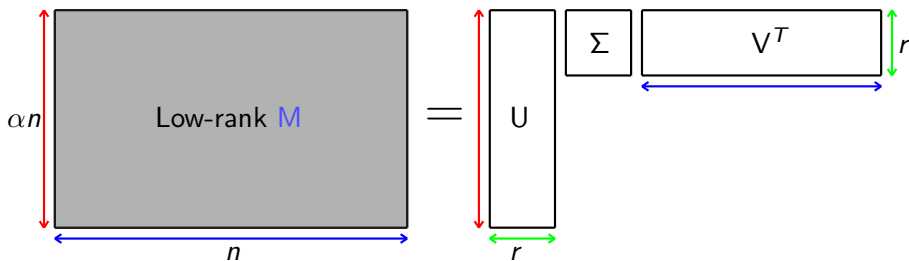


Fundamental Limit [Singer, Cucuringu '09], FPCA [Ma, Goldfarb, Chen '09],  
SVT [Cai, Candès, Shen '08], ADMiRA [Lee, Bresler '09]

## Story so far

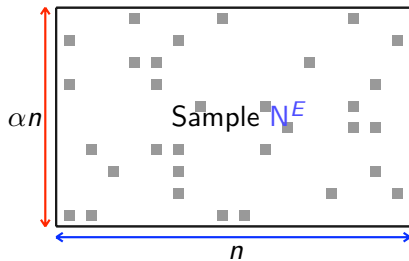
- OPTSPACE reconstructs  $M$  from a few sampled entries, when  $M$  is exactly low-rank and samples are exact
- In reality,
  - ▶  $M$  is only approximately low-rank
  - ▶ samples are corrupted by noise

# The model with noise



- Rank- $r$  matrix  $M$
- Random sample set  $E$
- Sample noise  $Z^E$
- Sample matrix  $N^E = M^E + Z^E$

# The model with noise



- Rank- $r$  matrix  $M$
- Random sample set  $E$
- Sample noise  $Z^E$
- Sample matrix  $N^E = M^E + Z^E$

# Main results

## Theorem

For  $|E| \geq C\mu r n \max\{\mu r, \log n\}$ , OPTSPACE achieves, with high probability,

$$\text{RMSE} \leq C' \frac{n\sqrt{r}}{|E|} \|Z^E\|_2,$$

provided that the RHS is smaller than  $\sigma_r(M)/n$ .

- $\|\cdot\|_2$  is the spectral norm

# OPTSPACE is order-optimal when noise is i.i.d. Gaussian

## Theorem

For  $|E| \geq C\mu rn \max\{\mu r, \log n\}$ , OPTSPACE achieves, with high probability,

$$\text{RMSE} \leq C' \sigma_z \sqrt{\frac{r n}{|E|}}$$

provided that the RHS is smaller than  $\sigma_r(M)/n$ .

- Lower bound: [Candès, Plan '09]

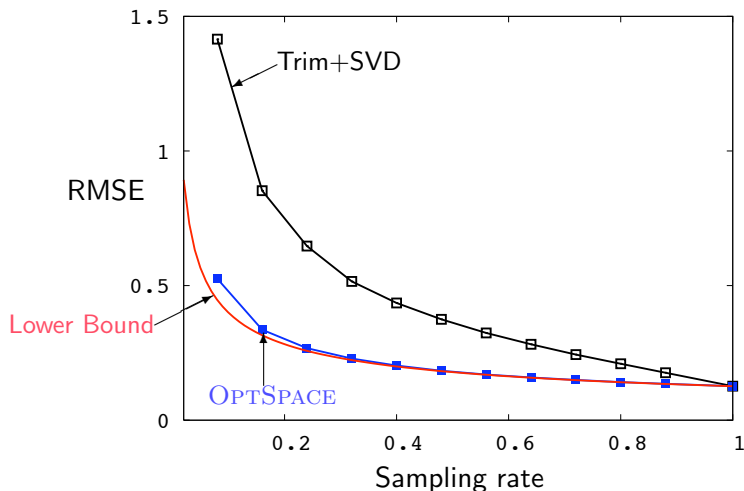
$$\text{RMSE} \geq \sigma_z \sqrt{\frac{2 r n}{|E|}}$$

- Trimming + SVD

$$\text{RMSE} \leq \underbrace{C M_{\max} \sqrt{\frac{r n}{|E|}}}_{\text{missing entries}} + \underbrace{C' \sigma_z \sqrt{\frac{r n}{|E|}}}_{\text{sample noise}}$$

## Comparison

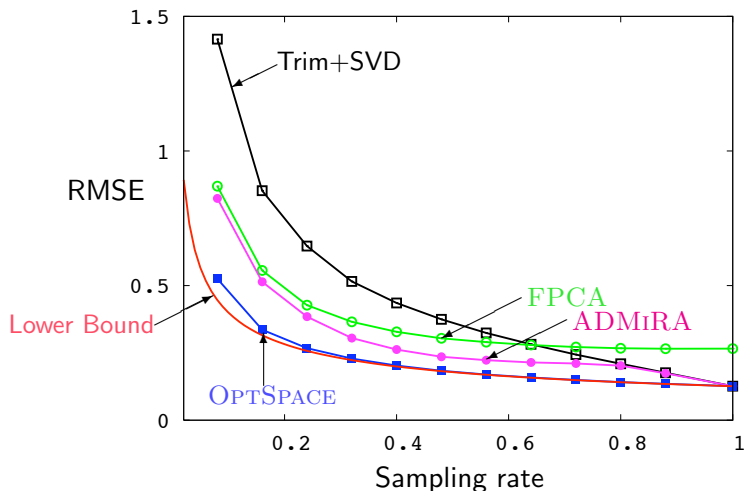
- $500 \times 500$  rank-4 matrix  $M$ , Gaussian noise with  $\sigma_z = 1$
- Example from [Candès, Plan '09]





## Comparison

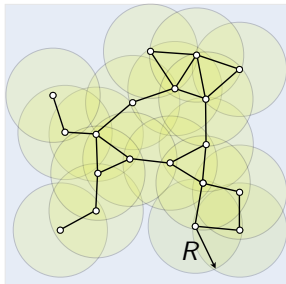
- $500 \times 500$  rank-4 matrix  $M$ , Gaussian noise with  $\sigma_z = 1$
- Example from [Candès, Plan '09]



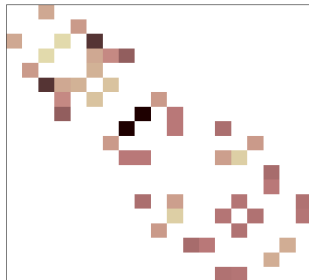
FPCA [Ma, Goldfarb, Chen '09], ADMiRA [Lee, Bresler '09]

## Positioning

# The model

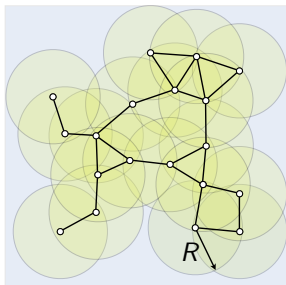


Distance Matrix  $D$

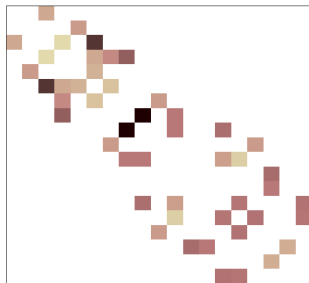


- $n$  wireless devices uniformly distributed in a bounded convex region
- Distance measurements between devices within radio range  $R$
- Find the locations up to a rigid motion

# The model

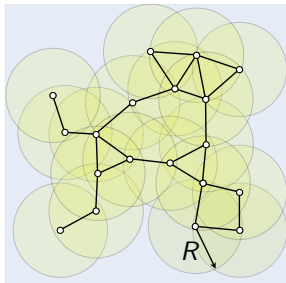


Distance Matrix  $D$

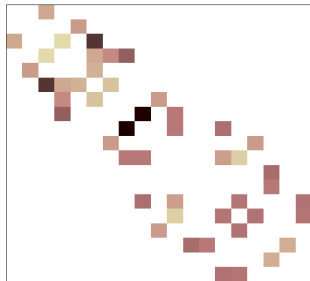


- How is it related to Matrix Completion?
  - ▶ Need to find the missing entries
  - ▶  $\text{rank}(D) = 4$
- How is it different?
  - ▶ Non-uniform sampling
  - ▶ Rich information not used in Matrix Completion

# The model



Distance Matrix  $D$



- MDS-MAP [Shang et al. '03]
  1. Fill in the missing entries with shortest paths
  2. Compute rank-4 approximation

# Main results

## Theorem

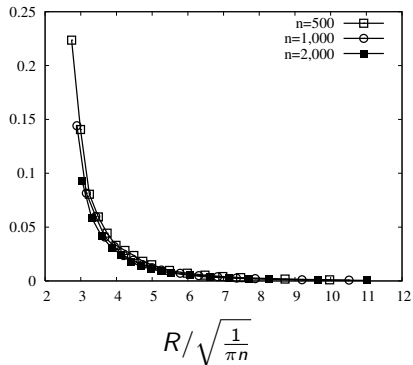
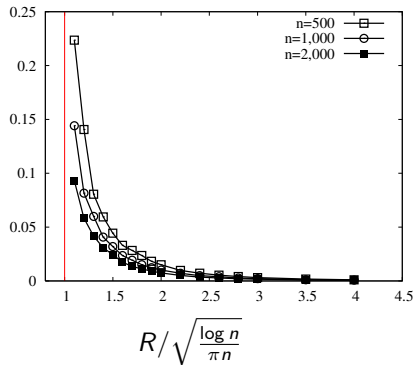
For  $R > C\sqrt{\frac{\log n}{n}}$ , with high probability,

$$\text{RMSE} \leq \frac{C}{R} \sqrt{\frac{\log n}{n}} + o(1) .$$

- $\text{RMSE} = \left( \frac{1}{n^2} \sum_{i,j} (D - \hat{D})_{ij}^2 \right)^{1/2}$
- Lower Bound:  
If  $R < \sqrt{\frac{\log n}{\pi n}}$ , then the graph is disconnected
- Generalized to quantized measurements and distributed algorithms
- We can add Greedy Minimization step

# Numerical simulation

## RMSE



# Conclusion

- Matrix Completion is an important problem with many practical applications
- OPTSPACE is an efficient algorithm for Matrix Completion
- OPTSPACE achieves performance close to the fundamental limit



Special thanks to:

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