



Searching Algorithms

Sequential Search
Algo

Binary Search
Algo

⇒ Sequential Search Algorithm

⇒ Straight-forward algorithm.

⇒ Searches the given key element in the list of n elements by checking successive elements in the list until a match is found or list exhausted.

⇒ Also called Linear Search.

* Algorithm Sequential Search $[A(0 \dots n-1), K]$

// searched for a given value in a given array by sequential search

// input \rightarrow an array $A[0 \dots n-1]$ and a search key K .

// output \rightarrow The index of the first element in A that matches K or -1 if there are no matching elements.

$i \leftarrow 0$

While $i < n$ and $a[i] \neq K$ do

$i \leftarrow i + 1$

if $i < n$

return i

else

return -1 .

★

$a[0]$	$a[1]$	$a[2]$	$a[3]$	$a[4]$
15	12	7	10	8

$n = 5$

$K = 7 \rightarrow$ searching element.

iteration \rightarrow 1. $i = 0$ while $0 < 5$ and $15 \neq 7 \rightarrow$ true.

$i = 0 + 1 = 1$

iteration 2 \rightarrow

$i = 1$

$a[1]$

while $1 < 5$ and $12 \neq 7 \Rightarrow$ true.

$i = 1 + 1 = 2$

iteration 3 \rightarrow

$i = 2$

$a[2]$

while $2 < 5$ and $7 = 7 \Rightarrow$ False.

true

False

$2 < 5 \rightarrow$ true

return $i \rightarrow a[i] \rightarrow a[2]$

is, returned is

⇒ Time Complexity

↳ best case ⇒ Search key is 1st element
only one iteration is required →

$$C_{\text{best}}(n) = 1$$

↳ Average case efficiency ⇒

$$C_{\text{avg}}(n) = \frac{n+1}{2} = \frac{5+1}{2} = \frac{6}{2} = 3$$

↓
comparisons

↳ worst case ⇒

$$C_{\text{worst}}(n) = n$$

Size of array.

★

Binary Search

⇒ array should be sorted

A

2	6	13	21	36	47	63	81	97
0	1	2	3	4	5	6	7	8

$$\text{mid} = \frac{0 + 8}{2} = 4$$

↓
a[mid] = 36

Search key = 21

$$a[\text{mid}] = 36$$

$$a[3] = 21$$

$36 > 21$ ⇒ Search Left side

9	6	13	21	36
0	1	2	3	4

$$\text{mid} = \frac{0+4}{2} = 2$$

$$A[2] = 13$$

$13 < 21 \Rightarrow$ Search Right side

13	21	36
2	3	4

$$\text{mid} = \frac{2+4}{2} = \frac{6}{2} = 3$$

$$A[3] = 21$$

$21 = 21 \Rightarrow$ Search Successful.

Algorithm

Binary Search \Rightarrow

follows
divide &
conquer
Strategy

$$l \leftarrow 0; r \leftarrow n-1$$

While $l \leq r$ do

$$m \leftarrow \left\lfloor \frac{(l+r)}{2} \right\rfloor$$

if $K = A[m]$ return m

else if $K < A[m]$ $r \leftarrow m-1$

else $l \leftarrow m+1$

return -1

Handwritten title: Insertion Sort

↳ Best case $\Rightarrow O(1)$

Worst case $\Rightarrow O(\log n)$

avg case $\Rightarrow O(\log n)$

↳

$$\frac{n * (\log n)}{n+1}$$