Master Thesis – Introductory seminar

Topic: Design and Prototypical Implementation of a High-Level Synchronization Component for Dynamic Updates of Task Run Queues in L4 Fiasco.OC/Genode

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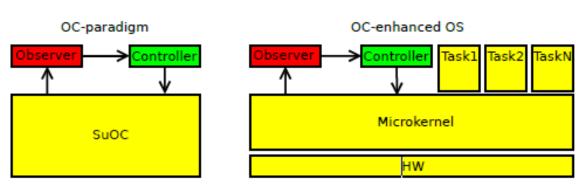
Advisors: Daniel Krefft, M.Sc.; Sebastian Eckl, M.Sc.

Outline

- Motivation
- Problem statement
- System constraints
- Related work
- Comparison/Evaluation of existing methods
- Design
- Implementation/Evaluation

Motivation

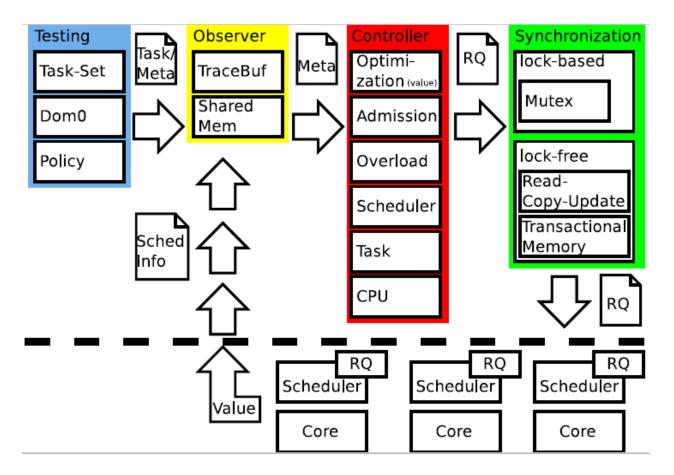
- KIA4SM Cooperative Integration Architecture for Future Smart Mobility Solutions
 - Communication between Intelligent Transport Systems(ITS)
 - Universally applicable ECUs
 - Organic computing
 - Requirements of KIA4SM
 - Flexible scheduling of tasks per ECU
 - Observer controller architecture: Decentralised OC
 - Observers collect data
 - Controllers analyse data



Problem Statement

- Controller produces a new Run Queue(RQ)
 - Update the run queue
 - Start executing new run queue
- Producer-Consumer problem

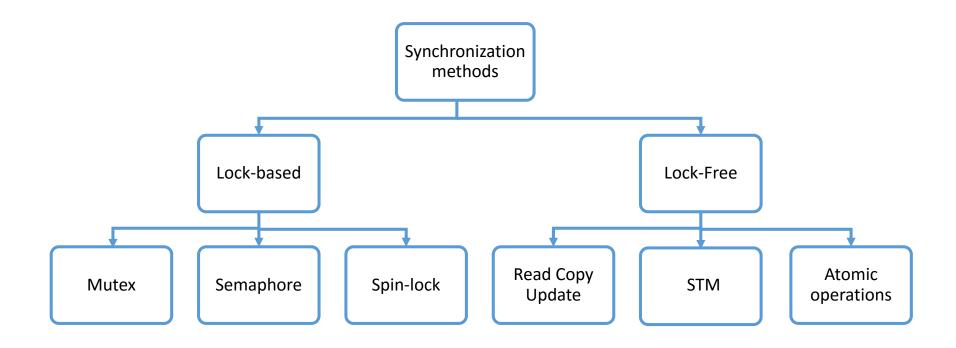
Synchronization of Run Queue



System constraints

- Real time system requirements
 - Safety Hard real time system
 - Efficiency Embedded system
 - Utilization Threads should reach their deadline
 - Predictable thread synchronization mechanism
- Organic-Computing-Compatibility

Related work: Synchronization Methods



Lock-based: Mutex

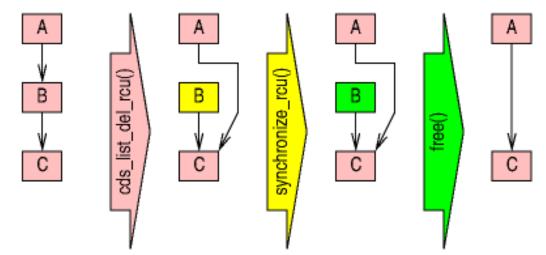
- Simple
- Easy to implement
- Only one thread can access/Modify at a time
- Deadlocks
- Priority inversion
- Starvation

Lock-free: Read Copy Update

- Concurrent read and write operations to the same data
- Three fundamental mechanisms
 - Publish-Subscribe Mechanism (for insertion)
 - Wait For Pre-Existing RCU Readers to Complete (for deletion)
 - Maintain Multiple Versions of Recently Updated Objects (for readers)
- Publish-Subscribe Mechanism
 - Publish: Execute pointer update in right order
 - Subscribe: Fetch the data in right order
 - Memory barrier and compiler directives

Lock-free: Read Copy Update(2)

- Wait For Pre-Existing RCU Readers to Complete
- Deletion of element from linked list



Grace period

Lock-free: Read Copy Update(3)

Advantages

- Concurrent access to shared data structure
- Low computation and storage overhead
- Deterministic completion time

Drawbacks

- Not a great tool when there are multiple writers
- Inconsistent data at times
- Difficult to implement

Lock-Free: Software Transactional Memory

- Concurrency control paradigm that provides atomic and isolated execution for regions of code
- Atomicity, Isolation
- Language extensions are required
 - Data versioning and conflict detection
- Global state of the program is consistent
- GCC 4.7
 - Eg:

```
void testfunc(int *x, int *y) {
    __transaction_atomic {
          *x += *y;
     }
}
```

Lock-Free: Software Transactional Memory(2)

Advantages

- Lock-free code is hard to implement
- More reliable performance
- Easy to Code (Uses language extensions)
- No-deadlocks
- Nested transactions are fine
- Compiler/Runtime figures out conflicts

Limitations

- Running expensive operations multiple times
- Overhead of transactional monitoring
- Difficult to debug
- On fewer cores(<4) performance is not increased

Lock-Free: Atomic operations

- Compare and Swap(CAS)
 - Basic primitive for many lock-free algorithms
 - Handled by CPU instructions
 - Atomically update by comparing memory contents to a given value, if value is same then update

```
int CAS(int *ptr,int oldvalue,int newvalue)
```

- Perform this operation repeatedly in a loop until it succeeds
- Limitations:
 - ABA problem
 - Considerable overhead (both time and space)

Evaluation of Existing methods

| | Mutex | RCU | STM |
|----------------------|-------|-----|-----|
| Implementation | + | | ++ |
| Read-speed | | ++ | + |
| Write-speed | | + | + |
| Deadlocks | | ++ | + |
| Overhead | + | ++ | |
| Security/Consistency | ++ | - | + |

Design

- Finding the right point for synchronization
 - The current running thread should execute
 - System should not go into unsafe state
- Empty Run Queue (RQ)



• Static point in time

Variable point in time



Implementation/Evaluation

- Includes getting the RQ from Controller and Switching it successfully with existing RQ
- Successful execution of new task set/RQ
- Finding the Ideal time for switching the run queue
- Benchmarking different synchronization methods for
 - Speed
 - Safety

Timeline

| Apr 15 – May 15 | May 15 – Jun 15 | Jun 15 – Jul 15 | July 15 – Aug 15 | Aug 15- Sept 15 | Sept15 – Oct 15 |
|-----------------|-----------------|-----------------|------------------|-----------------|-----------------|
| Research and | Design | Implementation | Implementation | Testing | Write up |
| Related work | | | Testing | Benchmarking | |

Conclusion

- KIA4SM project
 - Flexible scheduling of tasks per ECU
 - Dynamic thread scheduling necessary to get fault tolerant system
- Observer-Controller Architecture
 - Produce-Consumer problem
 - Run queue(RQ) synchronization
 - Finding right method and time for synchronization
- State of the Art algorithms
 - Lock-based: Mutex
 - Lock-Free: RCU, STM
- Theoretically RCU gives us the best choice in Lock-Free algorithms

References

- KIA4SM Sebastian Eckl, Daniel Krefft, Uwe Baumgarten
- Software transactional memory in gcc Dave boutcher
 - GCC benchmarking with STM, Locks
- What Is RCU? Paul E. McKenney, IBM Distinguished Engineer
 - RCU fundamentals
- Software Transactional Memory, Siyuan Zhou
 - STM basics, comparison to locks
- A Framework for Implementing Objects and Scheduling Tasks in Lock-Free Real-Time Systems - James H. Anderson and Srikanth Ramamurthy
 - RCU usage in real time systems
- Extending RCU for Realtime and Embedded Workloads: Paul E McKenney