

File Systems Basics

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Outline

- File System Overview
- File System Architecture
- File System Semantics
- File System Operations

File System Overview

- System that permanently stores data
- Usually layered on top of a lower-level physical storage medium
- Divided into logical units called "files"
 - Addressable by a filename ("foo.txt")
 - Usually supports hierarchical nesting (directories)
- A file path joins file & directory names into a relative or absolute address to identify a file ("/home/aaron/foo.txt")

File (an abstraction)

- A (potentially) large amount of information or data that lives a (potentially) very long time
 - Often much larger than the memory of the computer
 - Often much longer than any computation
 - Sometimes longer than life of machine itself
- (Usually) organized as a linear array of bytes or blocks
 - Internal structure is imposed by application
 - (Occasionally) blocks may be variable length
- (Often) requiring concurrent access by multiple processes
 - Even by processes on different machines!

<u>Directory – A Special Kind of File</u>

- A tool for users & applications to organize and find files
 - User-friendly names
 - Names that are meaningful over long periods of time

 The data structure for OS to locate files (i.e., containers) on disk

File System Basics

File:

- Named collection of logically related data
- Unix file: an uninterpreted sequence of bytes

File system:

- Provides a logical view of data and storage functions
- User-friendly interface
- Provides facility to create, modify, organize, and delete files
- Provides sharing among users in a controlled manner
- Provides protection

File Types and Attributes

File types:

- Regular files
- Directories
- Character special files: used for serial I/O
- Block special files: used to model disks [buffered I/O]
- File attributes: varies from OS to OS
 - Name, type, location, size, protection info, password, owner, creator, time and date of creation, last modification, access

File operations:

- Create, delete, open, close, read, write, append, get/set attributes
- File access:
 - Sequential, random

Attributes of Files

Name:

 Although the name is not always what you think it is!

Type:

 May be encoded in the name (e.g., .cpp, .txt)

Dates:

- Creation, updated, last accessed, etc.
- (Usually) associated with container
- Better if associated with content

Size:

 Length in number of bytes; occasionally rounded up

Protection:

- Owner, group, etc.
- Authority to read, update, extend, etc.

Locks:

For managing concurrent access

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Definition — File Metadata

- Information about a file
 - Maintained by the file system
 - Separate from file itself
 - Usually attached or connected to the file
 - E.g., in block # –1
 - Some information visible to user/application
 - Dates, permissions, type, name, etc.
 - Some information primarily for OS
 - Location on disk, locks, cached attributes

Observation – some attributes are not visible to user or program

- E.g., location
 - Location is stored in metadata
 - Location can change, even if file does not
 - Location is not visible to user or program

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File System

- Components: directory, authorization, file service and system service
 - Authorization service: between file and directory services
 - Directory service: used to keep track of the location of all resources in the system
 - File service provides a transparent way of accessing any file in the system in the same way
 - System service: file system's interface to hardware

File Systems

File service:

- Specification of what the file system offers to client
 - Actions
 - Client primitives
 - Parameters, application programming interface (API)
- Does not include how service is implemented

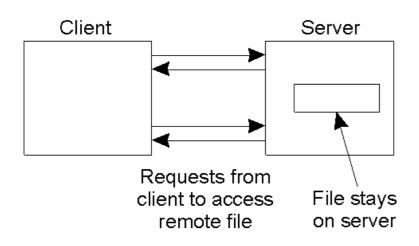
File server:

- Process that runs on a machine and implements file service
- Can have several servers on one machine (UNIX, DOS,...)
- ideally, clients do not know the distributed nature

Architectures

- How are FS generally organized?
 - Client-server architectures
 - Example: Sun Microsystem's NFS
 - File servers with a standardized view of its local file system; clients can access these files
 - Cluster-based distributed file systems
 - Example: file striping, partitioning the whole file system, GFS
 - Symmetric architectures
 - Fully symmetric organization based on p2p
 - Example: Ivy

Client-Server Architectures



- Client Server
 Old file
 New file
- 2. Accesses are done on client
- 3. When client is done, file is returned to server

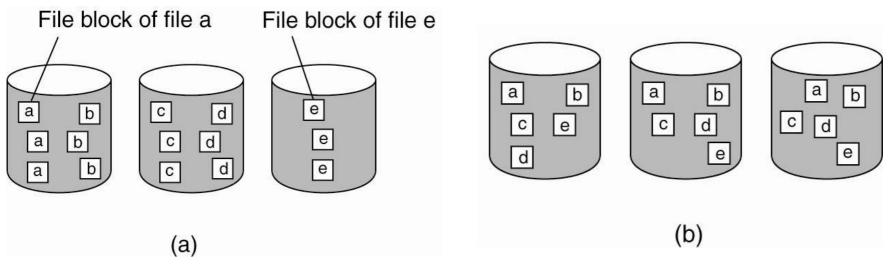
File moved to client

- Remote access model
 - Work done at the server
- Stateful server (e.g., databases)
- Pros & cons?

- Upload/download model
 - Work done at the client
- Stateless server
- Pros & cons?

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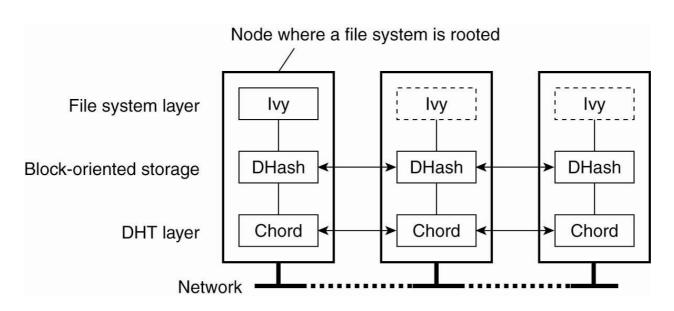
Cluster-based DFSs



- Figure (b): When server clusters are used for parallel applications
 - File-striping techniques, a single file is distributed across multiple servers
- Figure (a): For general-purpose applications, when file striping may not be effective
 - Partition the file system as a whole and simply store different files on different servers

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Symmetric Architectures



- Commonly used DHT-based systems for distributing data, combined with a key-based lookup mechanism
- Key difference: whether build a file system on top of a distributed storage layer
- Example: Ivy

System Structure

- How should the system be organized?
 - Are clients and server different?
 - Same process implements both functionality
 - Different processes, same machine
 - Different machines (a machine can either be client or server)
 - How are file and directory services organized-same server?
 - Different server processes: cleaner, more flexible, more overhead
 - Same server: just the opposite
 - Caching/no caching
 - server
 - client
 - How are updates handled?
 - File sharing semantics?
 - Server type: stateful vs. stateless

- Unix semantics: used in centralized systems
 - Read after write returns value written
 - System enforces absolute time ordering on all operations
 - Always returns most recent value
 - Changes immediately visible to all processes
- Issues in distributed systems
 - Single file server (no client caching): easy to implement UNIX semantics
 - Client file caching: improves performance by decreasing demand at the server; updates to the cached file are not seen by other clients
 - Conclusion:?

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- Session semantics (relaxed semantics)
 - Local changes only visible to process that opened file
 - File close => changes made visible to all processes
 - Allows local caching of file at client
- Problem:

– What if two or more clients are caching and modifying a file?

- No file update semantics (Immutable files):
 - Files are never updated/modified
 - Allowed file operations: CREATE and READ
 - Files are atomically replaced in the directory
 - Problems:
 - what if two clients want to replace a file at the same time?

Delete file in use by another process

Atomic transactions

- All file changes are delimited by a Begin and End transaction
- All files requests within the transaction are carried out in order
- The complete transaction is either carried out completely or not at all (atomicity)
- Serializable access
- Problem: ?

Operations on Files

Open, Close

- Gain or relinquish access to a file
- OS returns a file handle an internal data structure letting it cache internal information needed for efficient file access

Read, Write, Truncate

- Read: return a sequence of n bytes from file
- Write: replace n bytes in file, and/or append to end
- *Truncate:* throw away all but the first *n* bytes of file

Seek, Tell

- Seek: reposition file pointer for subsequent reads and writes
- Tell: get current file pointer

Create, Delete:

Conjure up a new file; or blow away an existing one

Methods for Accessing Files

Sequential access

Random access

Keyed (or indexed) access