

CS3343 Analysis of Algorithms Fall 2019

Homework 1

Due 9/5/19 before 11:59pm

Justify all of your answers with comments/text in order to receive full credit.
Completing the assignment in L^AT_EX will earn you extra credit on Midterm 1.

1. Longest Increasing Subsequence (8 points)

Consider the following problem:

Input: An array $A[1 \dots n]$ of integers

Output: The largest integer m such that the array $A[1 \dots n]$ has subsequence of length m which is strictly increasing.

The following pseudocode finds the length of the longest of the given array $A[1 \dots n]$ by considering all possible subsequences:

Algorithm 1 longestSubSeq(int $A[1 \dots n]$)

```
1:  $k = 0;$ 
2: while (true) do
3:   // (I) The longest increasing subsequence of  $A$  has length  $\leq n - k$ 
4:   do
5:     Let  $B = A$  with  $k$  of its elements removed
6:     if ( isIncreasing(B) ) then
7:       return  $n - k$ ;
8:     end if
9:   while ( there are other untested choices of  $k$  elements to remove )
10:   $k++;$ 
11: end while
```

The following code checks if an array is increasing (i.e., each number is smaller than the next in the array).

Algorithm 2 isIncreasing(int $C[1 \dots n]$)

```
1:  $i = 1;$ 
2: while  $i < n$  do
3:   if (  $C[i] \geq C[i + 1]$  ) then
4:     return false;
5:   end if
6:    $i++;$ 
7: end while
8: return true;
```

Example: `longestSubSeq([2, 4, 3, 8, 5, 5, 6, 7, 9])` returns 6

Justification: $[2, 4, 3, 8, 5, 5, 6, 7, 9] = [2, 4, 5, 6, 7, 9]$ which is a longest increasing subsequence of the original array.

- (1) (2 points) Consider running longestSubSeq on the array:

$[1\cancel{X}9, 100, 112, 1\cancel{X}, 113, 1\cancel{X}0, 1\cancel{X}, 115, 120] = \boxed{100 \ 112 \ 113 \ 115 \ 120}$

What does longestSubSeq return and what array B causes this return?

- (2) (4 points) Use induction to prove the loop invariant (I) is true and then use this to prove the correctness of the algorithm. Specifically complete the following:
- Base case
 - Inductive step
(Hint: you can assume that line 5 will successively try all possible choices of how to remove k elements from A .)
 - Termination step
(Hint: the outer loop never terminates but consider what can you say about the k value that causes us to return.)
- (3) (1 point) Give the best-case runtime of longestSubSeq in asymptotic (i.e., O) notation as well as a description of an array which would cause this behavior.
- (4) (1 point) Give the worst-case runtime of longestSubSeq in asymptotic (i.e., O) notation as well as a description of an array which would cause this behavior.

2. Asymptotic Notation (8 points)

Show the following using the definitions of O , Ω , and Θ .

- (1) (2 points) $2n^3 + n^2 + 4 \in \Theta(n^3)$
 (2) (2 points) $3n^4 - 9n^2 + 4n \in \Theta(n^4)$

(Hint: careful with the negative number)

- (3) (4 points) Suppose $f(n) \in O(g_1(n))$ and $f(n) \in O(g_2(n))$. Which of the following are true? Justify your answers using the definition of O . Give a counter example if it is false.
- $f(n) \in O(5 \cdot g_1(n) + 100)$
 - $f(n) \in O(g_1(n) + g_2(n))$
 - $f(n) \in O(\frac{g_1(n)}{g_2(n)})$
 - $f(n) \in O(\max(g_1(n), g_2(n)))$

3. Summations (4 points)

Find the order of growth of the following sums.

$$(1) \sum_{i=10}^n (5i + 3) = \sum_{i=0}^n 5i + \sum_{i=0}^n 3 = 5 \sum_{i=0}^n i + 3n - 27 = 5 \left(\frac{n(n+1)}{2} \right) + 3n - 27 = \frac{5n^2 + 5n}{2} + 3n - 27 \quad \text{so, } \sum_{i=0}^n 5i + 3 \in \Theta(n^2)$$

$$(2) \sum_{i=0}^{\log_2(n)} 2^i \quad (\text{for simplicity you can assume } n \text{ is a power of 2})$$

$$2^0 + 2^1 + 2^2 + \dots + 2^{\log_2 n} \text{ copies}$$

$$\text{so, } \sum_{i=0}^{\log_2 n} 2^i \in \Theta(n)$$

① Longest Increasing Subsequence.

Base case: $k=0$

(I) The longest increasing sequence of A has length $\leq n-k$
is true because each element is increasing
sequence.

inductive step: assume A has length $\leq n-k$ we need to show
prove $A[i] \leq n-k+1$ ↗
↳ LSA(A) $\neq n-k$
↳ decreasing

→ If $greatest$ was $n-k$ then one of the 3 arrays that we slice would have been increasing since there's tries on B arrays of length $n-k$.

Termination step: the program terminates when all elements are tested and finds the longest sequence until the last element.

best case: best case would be $O(n^2)$
in the given array is already sorted in increasing order.

worst case: $O(2^n)$ if the array is sorted in descending order.

② Asymptotic Notation

$$\textcircled{1} \quad 2n^3 + n^2 + 4 \in O(n^3)$$

$$f(n) \leq c \cdot g(n) \quad \forall n \geq n_0$$

$$2n^3 + n^2 + 4 \leq 2n^3 + n^3 + 4n^3$$

$$2n^3 + n^3 + 4 \leq 7n^3$$

$$c = 7 \quad n_0 = 1$$

$$f(n) \in O(n^3)$$

$$f(n) \geq c \cdot \Omega(g(n)) \quad \forall n \geq n_0$$

$$2n^3 + n^3 + 4n^3 \geq 2n^3 + 0 + 0 \\ \geq 2n^3$$

$$c = 2 \quad n_0 = 1$$

$$f(n) \in \Omega(n^3)$$

$$\text{so, } f(n) \in \Theta(n^3)$$

$$\textcircled{2} \quad 3n^4 - 9n^2 + 4n \in O(n^4)$$

$$f(n) \leq g(n) \quad \forall n \geq n_0$$

$$3n^4 - 9n^2 + 4n \leq 3n^4$$

$$c = 3 \quad n_0 = 1$$

$$f(n) \geq c \cdot \Omega(g(n)) \quad \forall n \geq n_0$$

$$3n^4 - 9n^2 + 4n \geq 3n$$

$$c = 3 \\ n_0 = 1$$

$$f(n) \in \Omega(n^4)$$

$$\text{so, } f(n) \in \Theta(n^4)$$

$$\textcircled{3a} \quad f(n) \in O(5g_1(n) + 100)$$

True

$$\text{we know } f(n) \in O(g_1(n))$$

$$\text{so } 0 \leq f(n) \leq c \cdot g_1(n) \quad \text{for } n \geq n_0$$

$$\textcircled{3b} \quad f(n) \in O(g_1(n) + g_2(n))$$

True:

$$f(n) \in O(g_1(n))$$

$$f(n) \in O(g_2(n))$$

$$f(n) \leq c_1 \cdot (g_1(n) + g_2(n))$$

$$f(n) \leq c_2 \cdot g_1(n) + g_2(n) \quad \text{for all } n \geq n_1$$

$$\textcircled{3c} \quad f(n) \in O\left(\frac{g_1(n)}{g_2(n)}\right)$$

False:

$$n \in O(n) \quad n^2 \in O(n^2) \text{ but } (n^2 \cdot n^2) \notin O(n^2)$$

$$\textcircled{3d} \quad f(n) \in O(\max(g_1(n), g_2(n)))$$

True

$$f(n) \in O(g_1(n)) \Rightarrow f(n) \leq c_1 \cdot g_1(n) \quad \forall n \geq n_1$$

$$f(n) \in O(g_2(n)) \Rightarrow f(n) \leq c_2 \cdot g_2(n)$$

$$f(n) \in O(\max(g_1(n), g_2(n))) \Rightarrow c \cdot \max(g_1(n), g_2(n))$$

for all $n \geq n_0$
and $c \geq 1$

