

# SQUIP

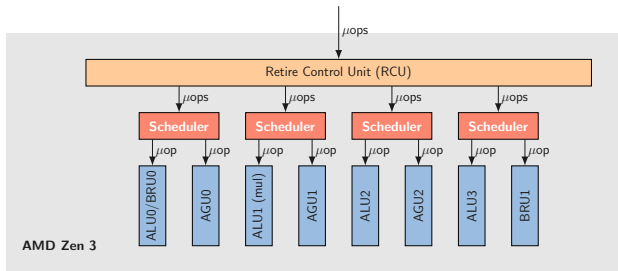
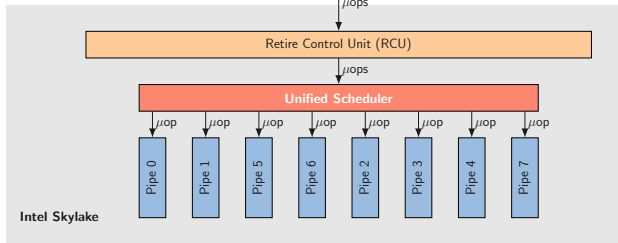
## Exploiting the Scheduler Queue Contention Side Channel

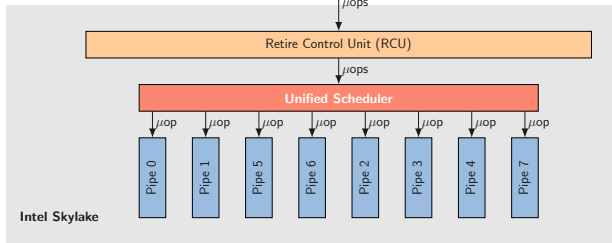
**Stefan Gast**<sup>1</sup>   **Jonas Juffinger**<sup>1</sup>   **Martin Schwarzl**<sup>1</sup>   **Gururaj Saileshwar**<sup>2</sup>   **Andreas Kogler**<sup>1</sup>  
**Simone Franza**<sup>1</sup>   **Markus Köstl**<sup>1</sup>   **Daniel Gruss**<sup>1</sup>

2023-05-23

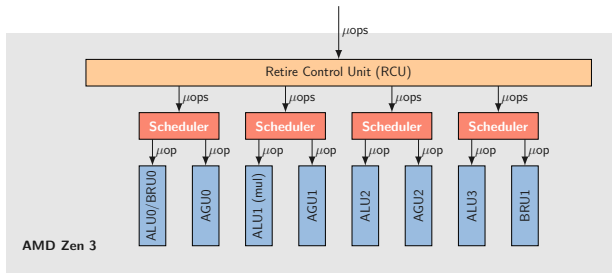
<sup>1</sup>Graz University of Technology

<sup>2</sup>Georgia Institute of Technology





Can we exploit this difference?



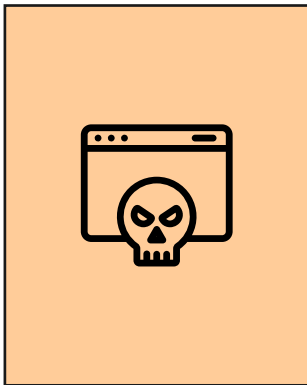
## Attacker



## Victim

```
mov (%rsi), %rax
mul %rbx
add $0x8, %rsi
add %rcx, %rax
mov %r8, %rcx
adc $0x0, %rdx
nop
```

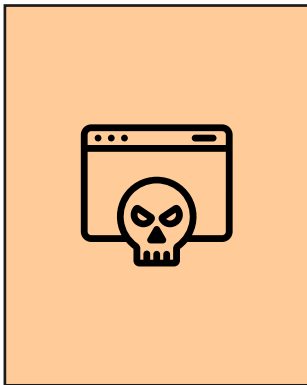
## Attacker



## Victim

```
mov (%rsi), %rax  
mul %rbx  
add $0x8, %rsi  
add %rcx, %rax  
mov %r8, %rcx  
adc $0x0, %rdx  
nop
```

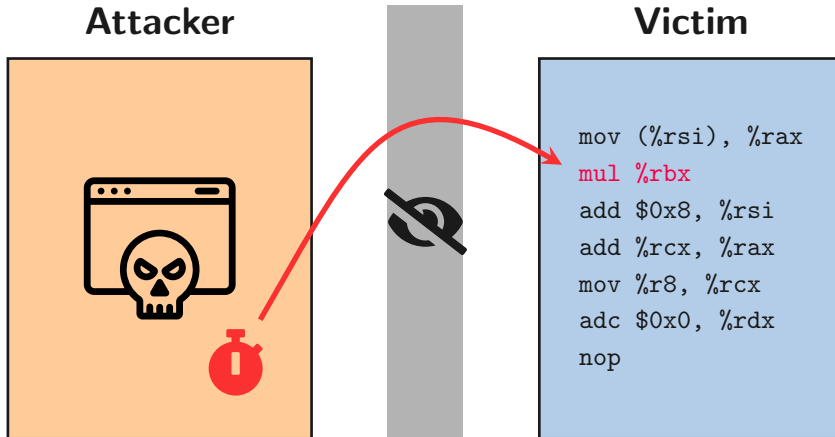
## Attacker

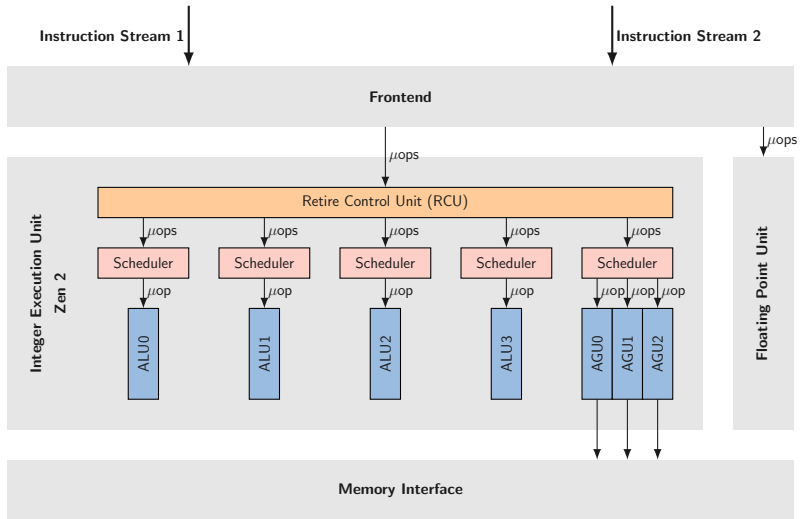


## Victim

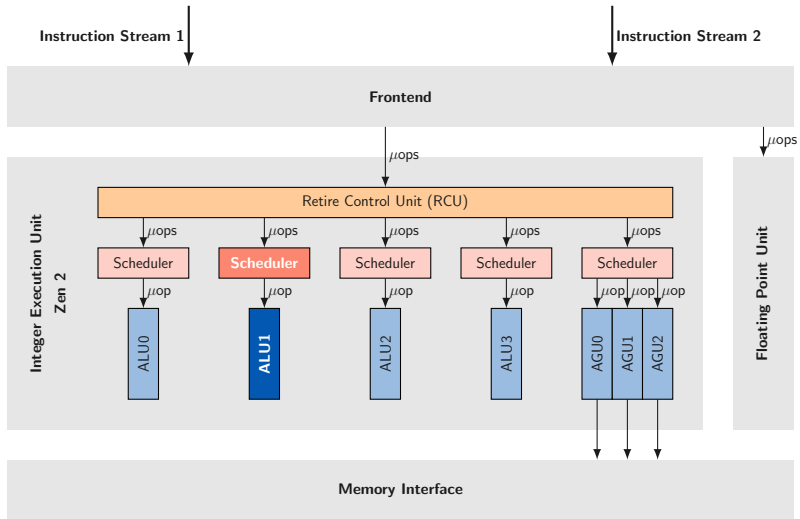
```
mov (%rsi), %rax  
mul %rbx  
add $0x8, %rsi  
add %rcx, %rax  
mov %r8, %rcx  
adc $0x0, %rdx  
nop
```

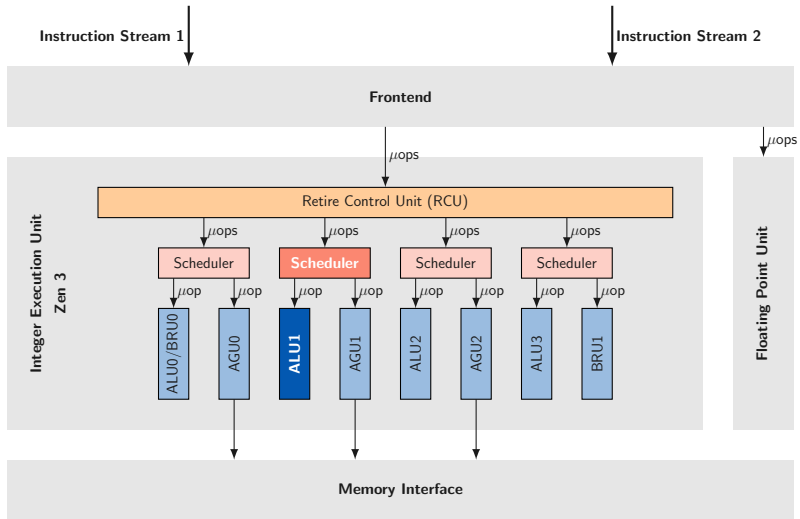


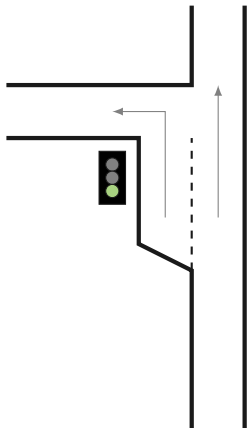


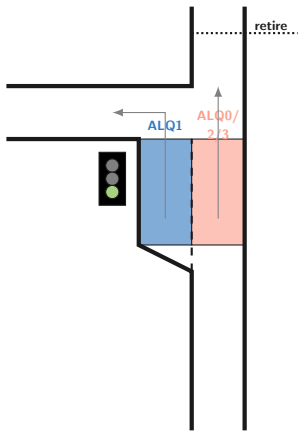


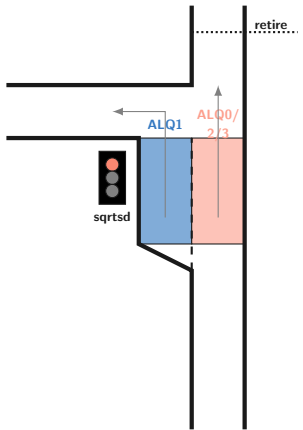


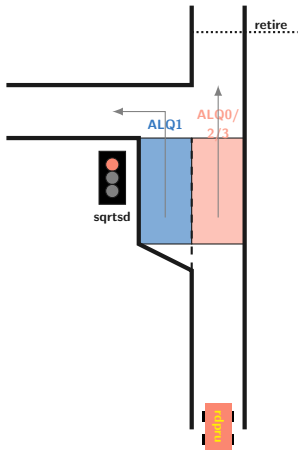


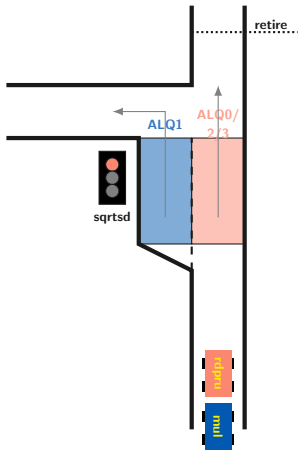


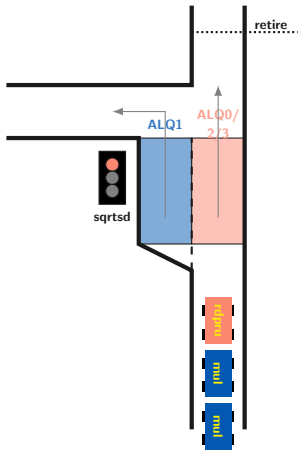




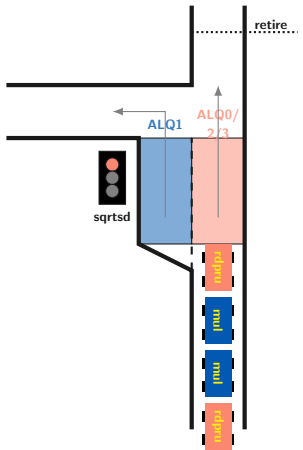


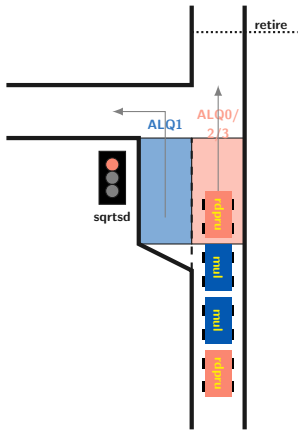


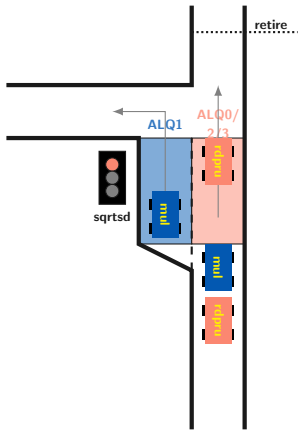


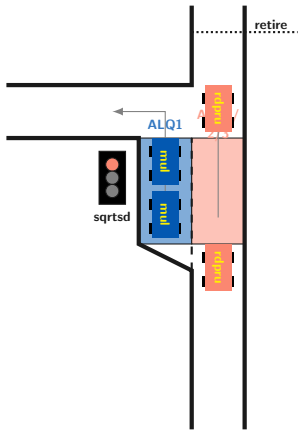


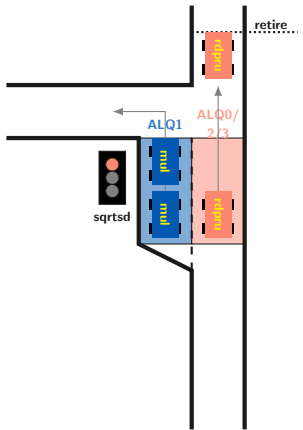


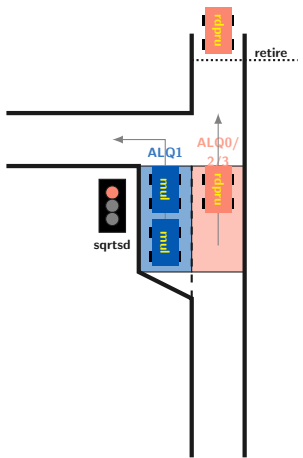


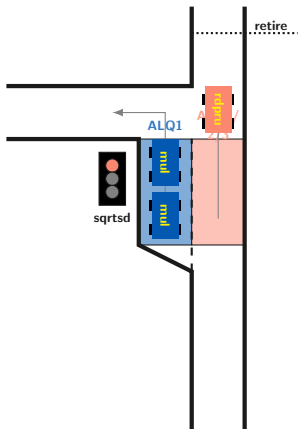


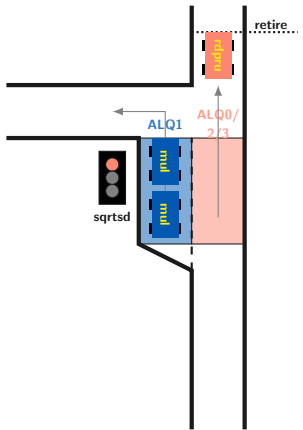




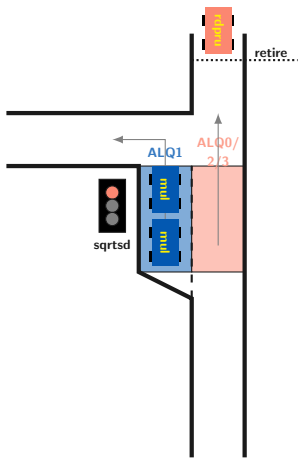


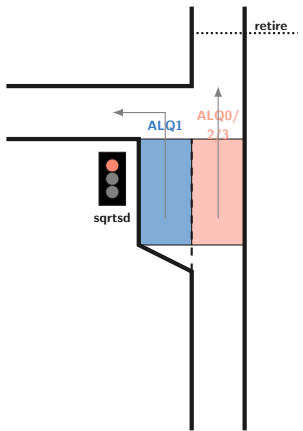


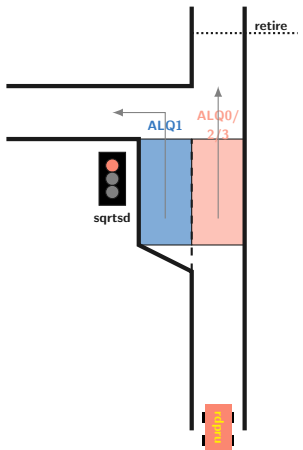


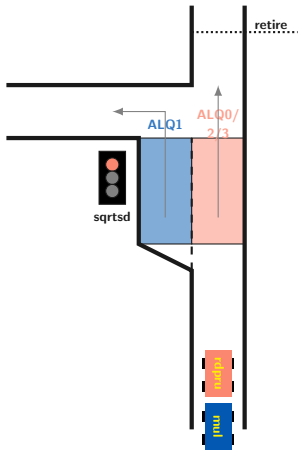


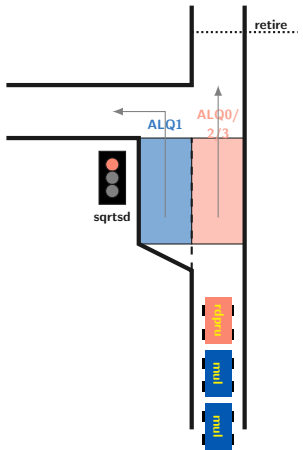


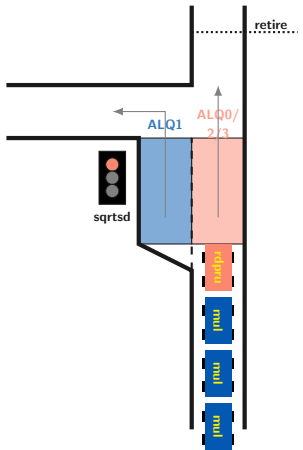


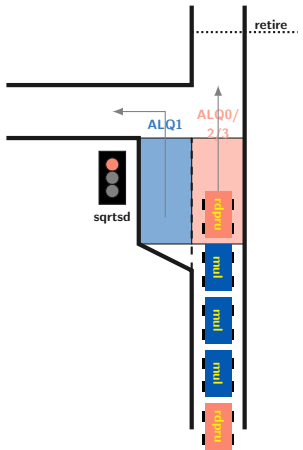


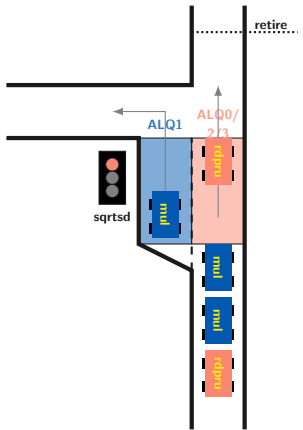




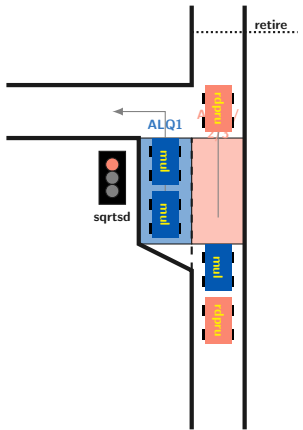


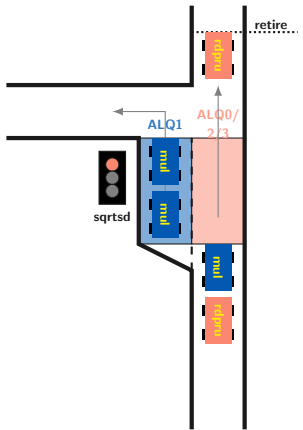


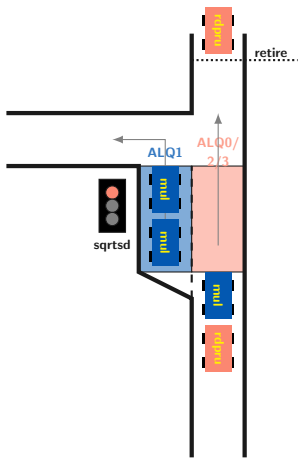


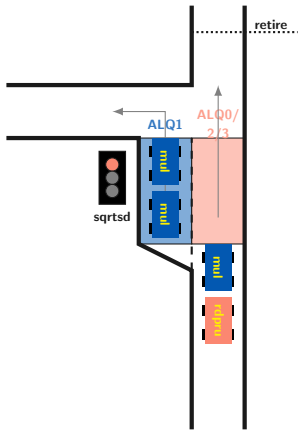




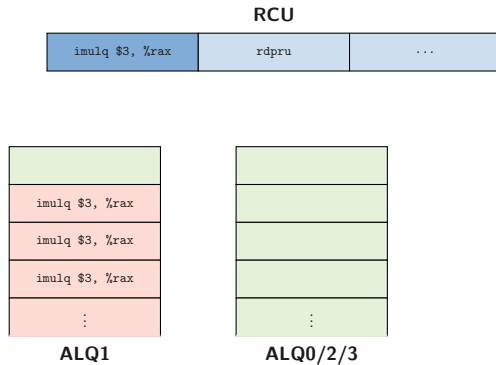




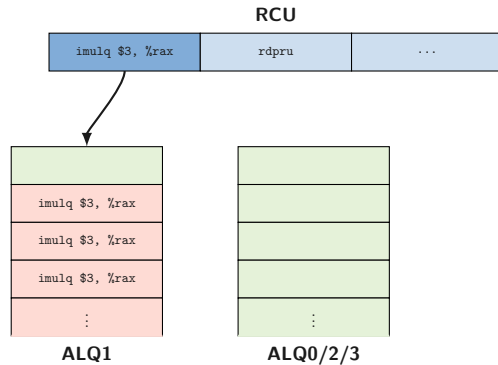




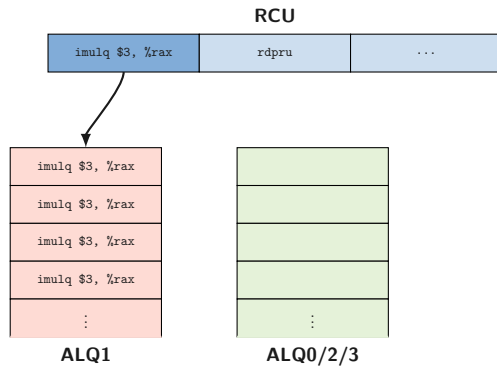
# What if a scheduler queue is full?



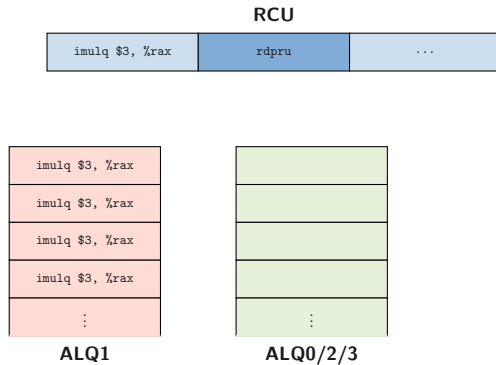
# What if a scheduler queue is full?



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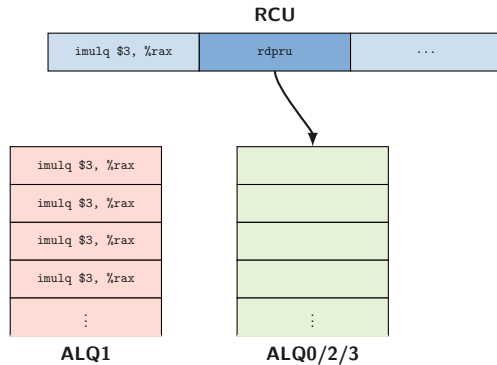


# What if a scheduler queue is full?

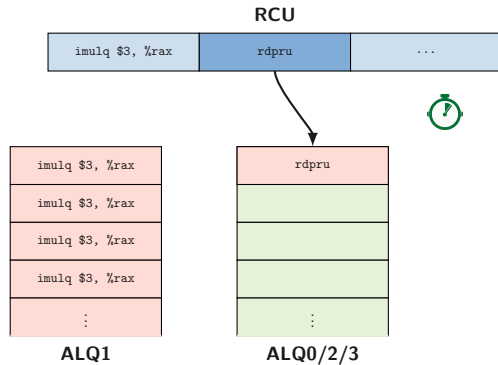




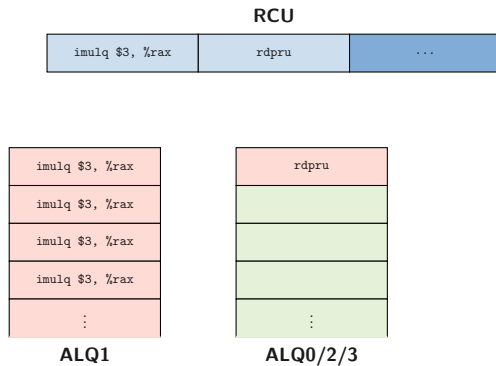
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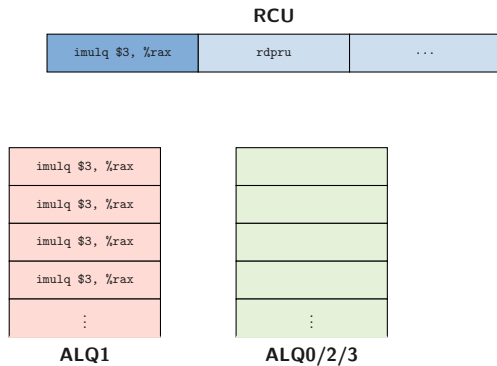
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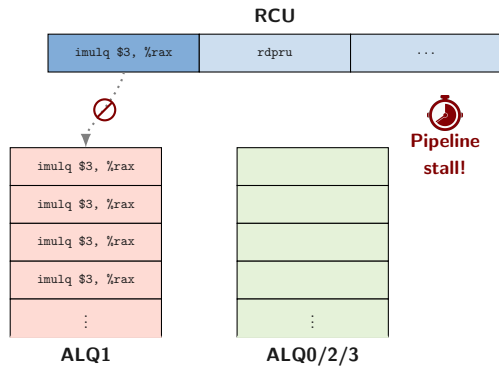
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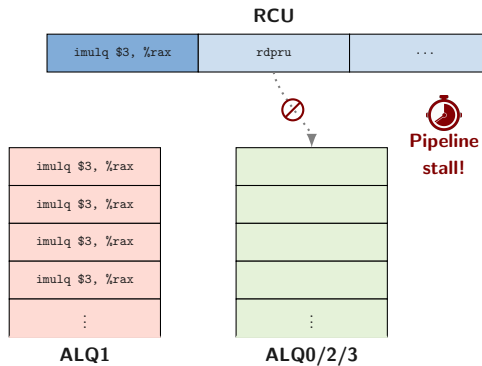
# What if a scheduler queue is full?



# What if a scheduler queue is full?



# What if a scheduler queue is full?



```
rdtsc / rdpru
```

...

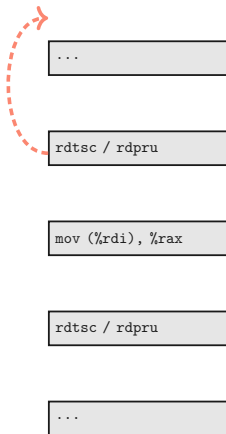
rdtsc / rdpru

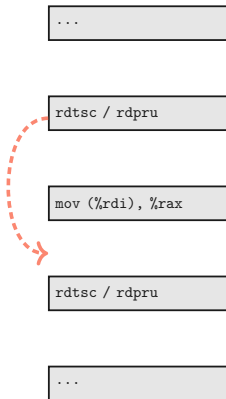
mov (%rdi), %rax

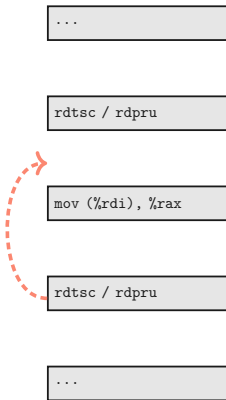
rdtsc / rdpru

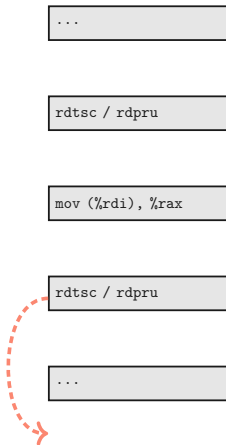
...











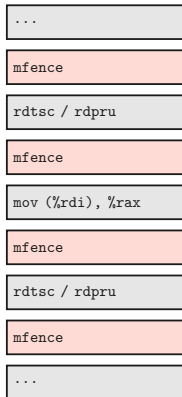
...

rdtsc / rdpru

mov (%rdi), %rax

rdtsc / rdpru

...



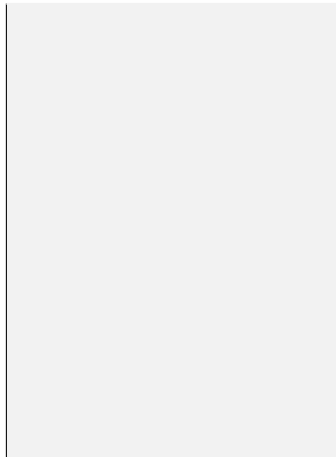
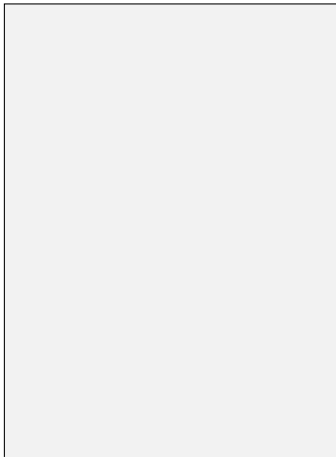


Adding  
serializing  
instructions  
around rdtsc



Exploiting  
out-of-order  
execution of rdtsc

# Let's code that!





# Let's code that!

```
movl $10000, %eax  
loop:  
subl $1, %eax  
jnz loop
```

drain  
queue



# Let's code that!

```
movl $10000, %eax  
loop:  
subl $1, %eax  
jnz loop
```

drain  
queue

```
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
# ...
```

fill  
queue



# Let's code that!

```
movl $10000, %eax  
loop:  
subl $1, %eax  
jnz loop
```

```
movq $12345678, %r15  
cvtsi2sd %r15, %xmm0  
sqrtsd %xmm0, %xmm0  
sqrtsd %xmm0, %xmm0  
sqrtsd %xmm0, %xmm0  
cvtsd2si %xmm0, %r15
```

drain  
queue

delay  
multipli-  
cations

```
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
# ...
```

fill  
queue



# Let's code that!

```
movl $1, %ecx
```

```
movl $10000, %eax  
loop:  
subl $1, %eax  
jnz loop
```

```
movq $12345678, %r15  
cvtsi2sd %r15, %xmm0  
sqrtsd %xmm0, %xmm0  
sqrtsd %xmm0, %xmm0  
sqrtsd %xmm0, %xmm0  
cvtsd2si %xmm0, %r15
```

drain  
queue

delay  
multipli-  
cations

```
rdpru  
movl %eax, %ebx
```

```
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
# ...
```

read  
time

fill  
queue



# Let's code that!

```
movl $1, %ecx
```

```
movl $10000, %eax  
loop:  
subl $1, %eax  
jnz loop
```

```
movq $12345678, %r15  
cvtsi2sd %r15, %xmm0  
sqrtsd %xmm0, %xmm0  
sqrtsd %xmm0, %xmm0  
sqrtsd %xmm0, %xmm0  
cvtsd2si %xmm0, %r15
```

drain  
queue

delay  
multipli-  
cations

```
rdpru  
movl %eax, %ebx
```

```
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
imulq $3, %r15  
# ...
```

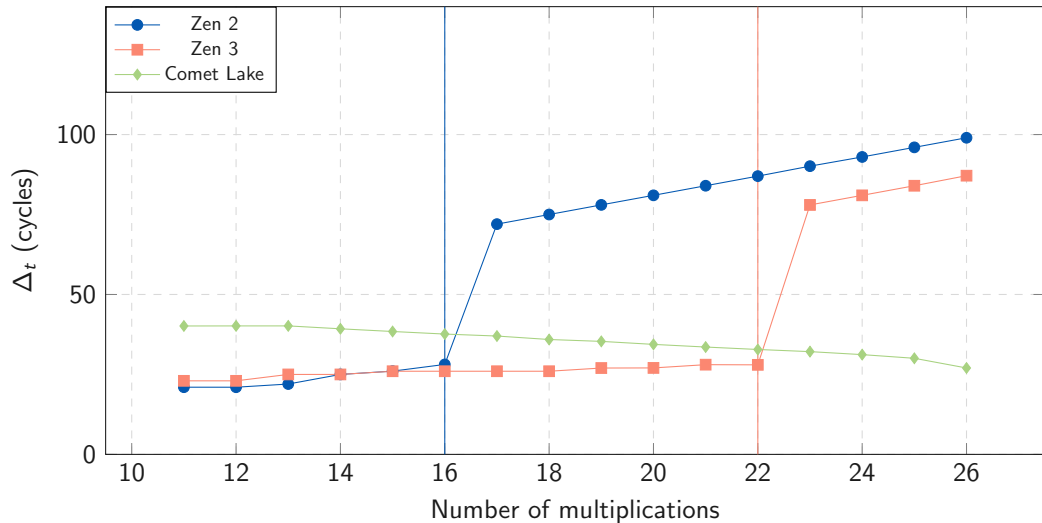
```
rdpru  
subl %ebx, %eax
```

read  
time

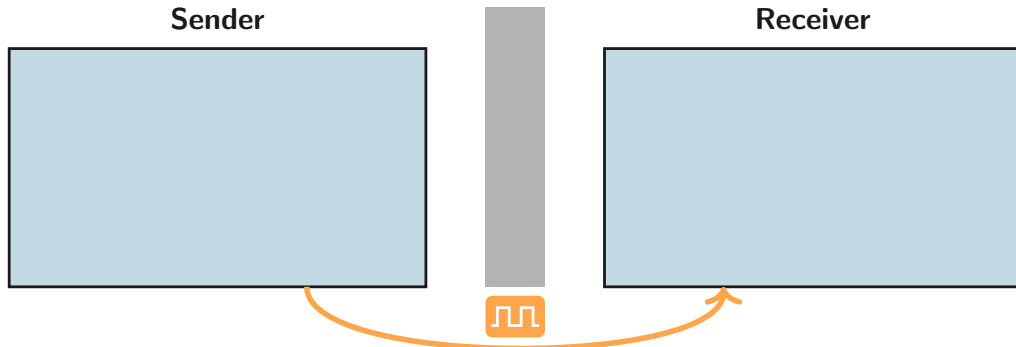
fill  
queue

read  
time

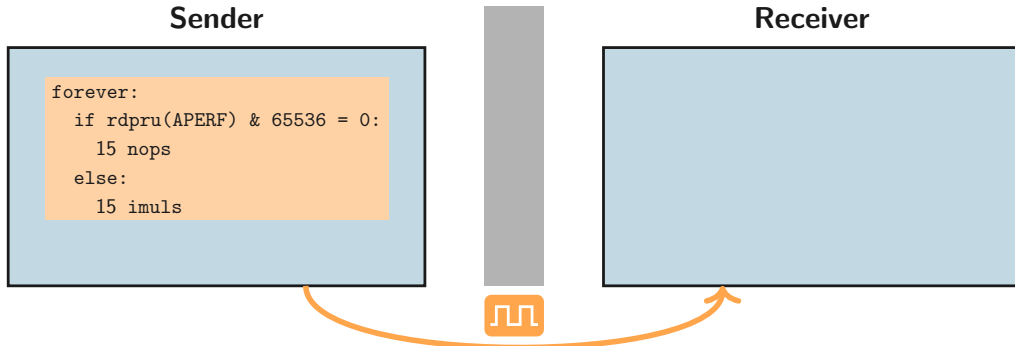
# Timing Differences on Zen 2, Zen 3 and Comet Lake



# Observing Multiplications of the Sibling Thread

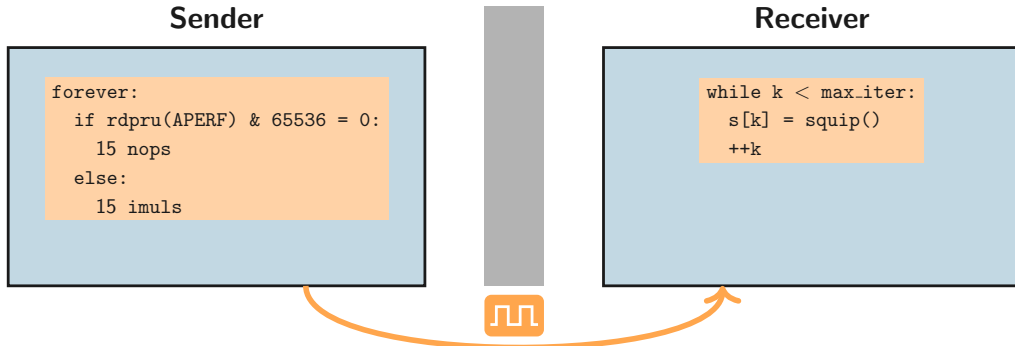


# Observing Multiplications of the Sibling Thread

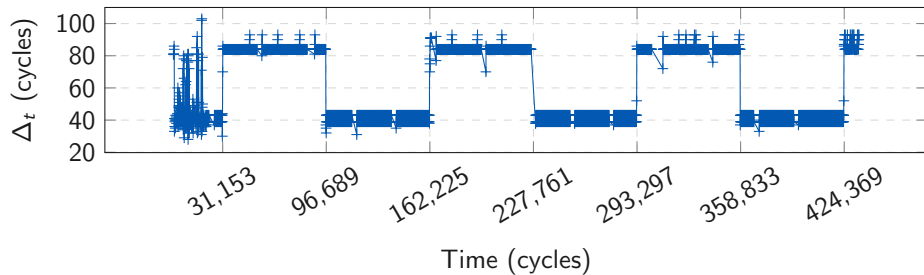




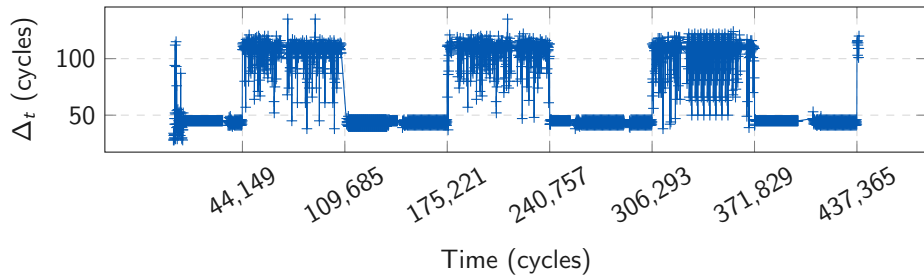
# Observing Multiplications of the Sibling Thread



# Observing Multiplications of the Sibling Thread (Zen 2)



# Observing Multiplications of the Sibling Thread (Zen 3)





Scenario	CPU	Raw Tx Rate	Error Rate

10 000 random messages, each with 32 kbit payload

Scenario	CPU	Raw Tx Rate	Error Rate
Cross-Process	Ryzen 7 3700X (Zen 2)	2.195 Mbit s <sup>-1</sup>	0.71 %
	Ryzen 7 5800X (Zen 3)	2.700 Mbit s <sup>-1</sup>	0.62 %

10 000 random messages, each with 32 kbit payload

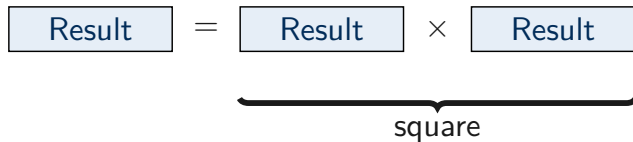
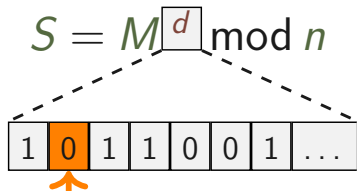
Scenario	CPU	Raw Tx Rate	Error Rate
Cross-Process	Ryzen 7 3700X (Zen 2)	2.195 Mbit s <sup>-1</sup>	0.71 %
	Ryzen 7 5800X (Zen 3)	2.700 Mbit s <sup>-1</sup>	0.62 %
Cross-VM	Ryzen 7 3700X (Zen 2)	0.873 Mbit s <sup>-1</sup>	3.18 %
	Ryzen 7 5800X (Zen 3)	0.892 Mbit s <sup>-1</sup>	0.75 %
	EPYC 7443 (Zen 3)	0.874 Mbit s <sup>-1</sup>	0.96 %

10 000 random messages, each with 32 kbit payload

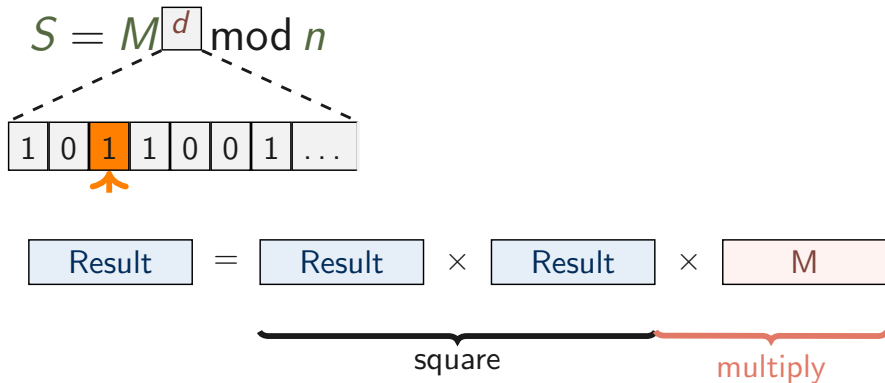
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	EPYC 7443 (Zen 3)	0.874 Mbit s <sup>-1</sup>	0.96 %
Cross-VM (SEV)	EPYC 7443 (Zen 3)	0.873 Mbit s <sup>-1</sup>	1.47 %

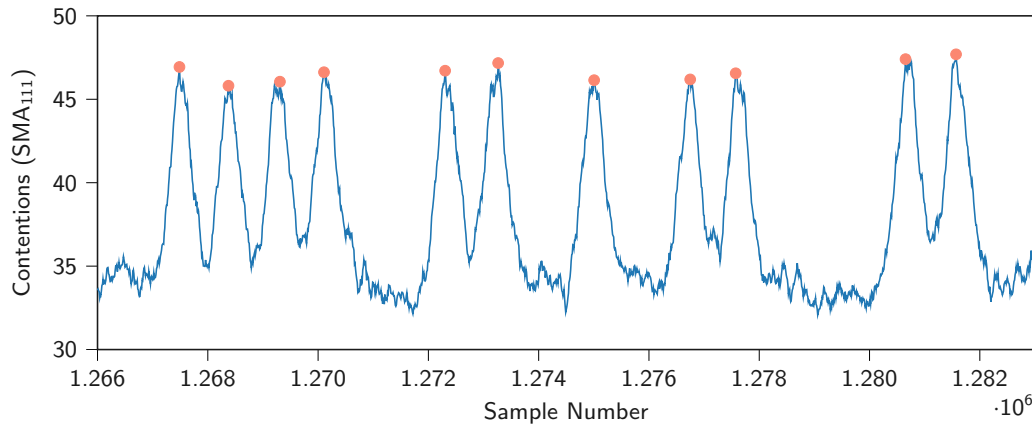
10 000 random messages, each with 32 kbit payload

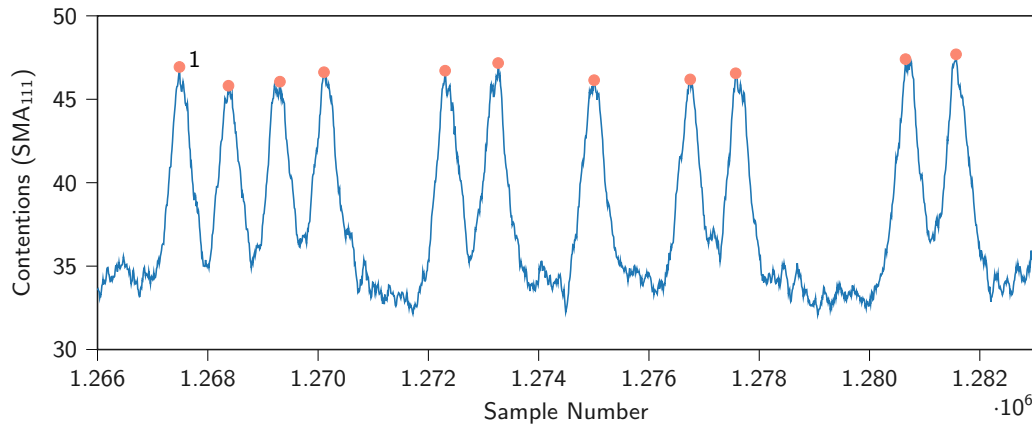


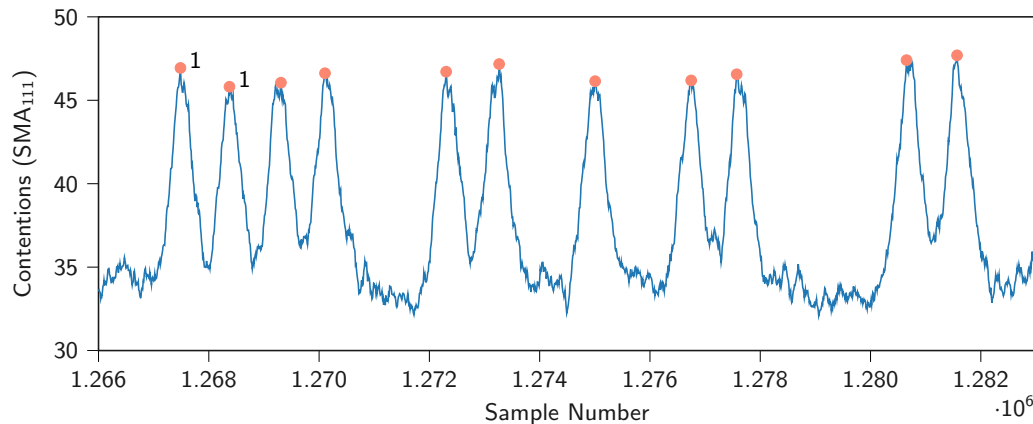


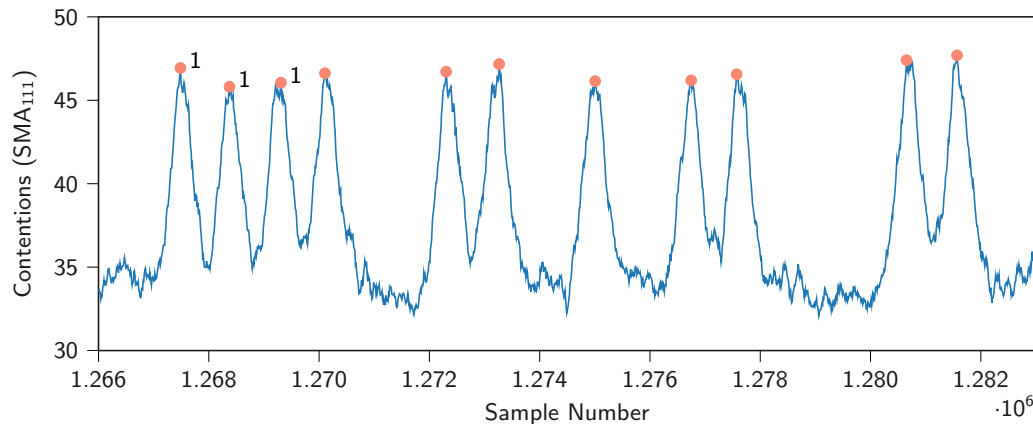
# Attacking RSA (Square+Multiply)

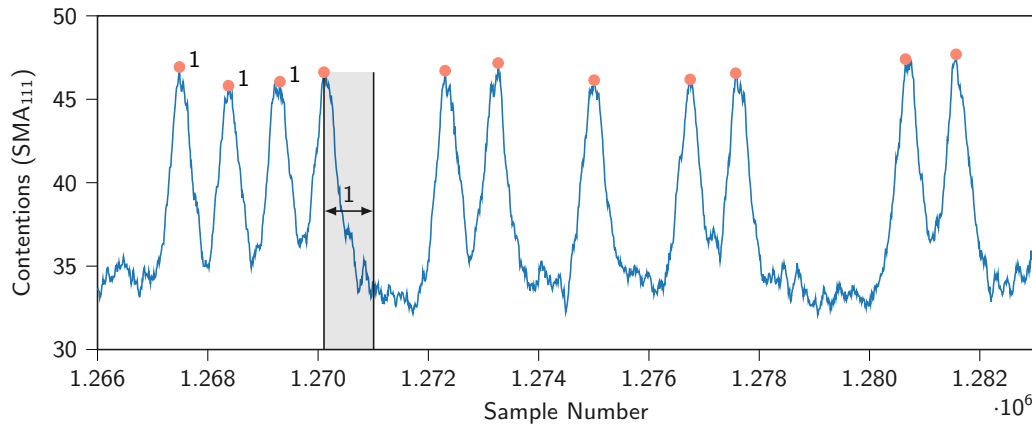


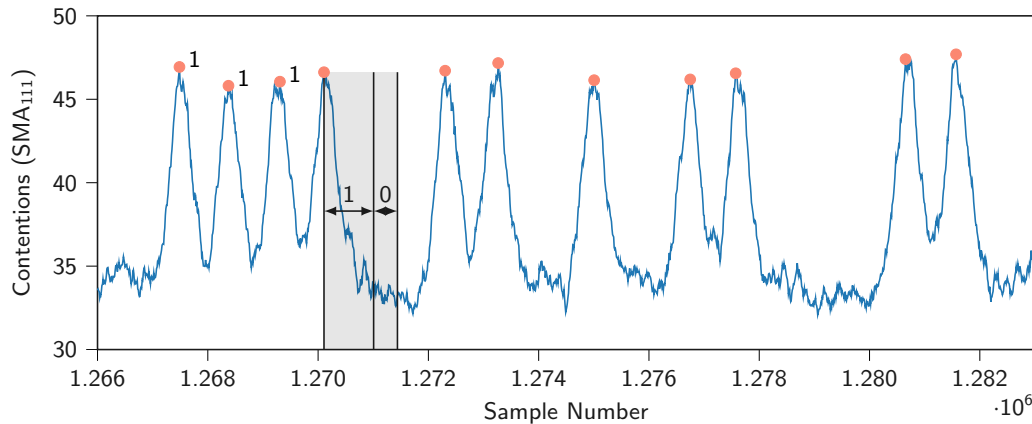




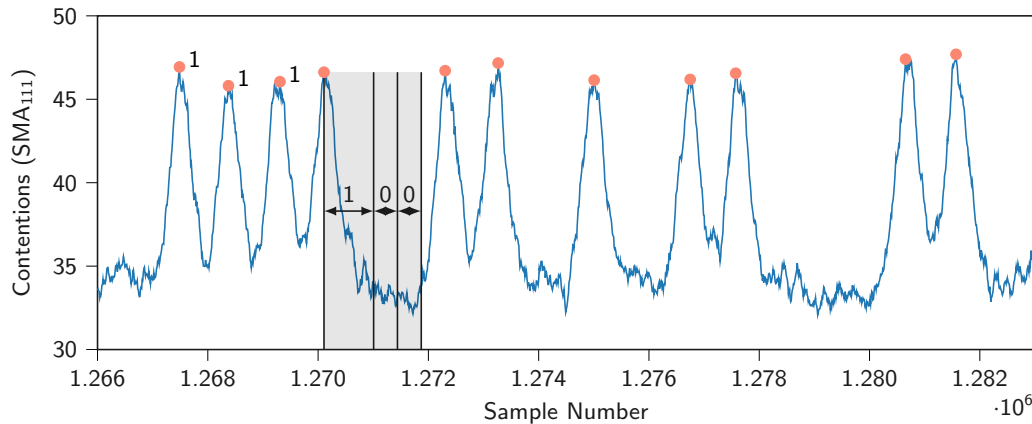


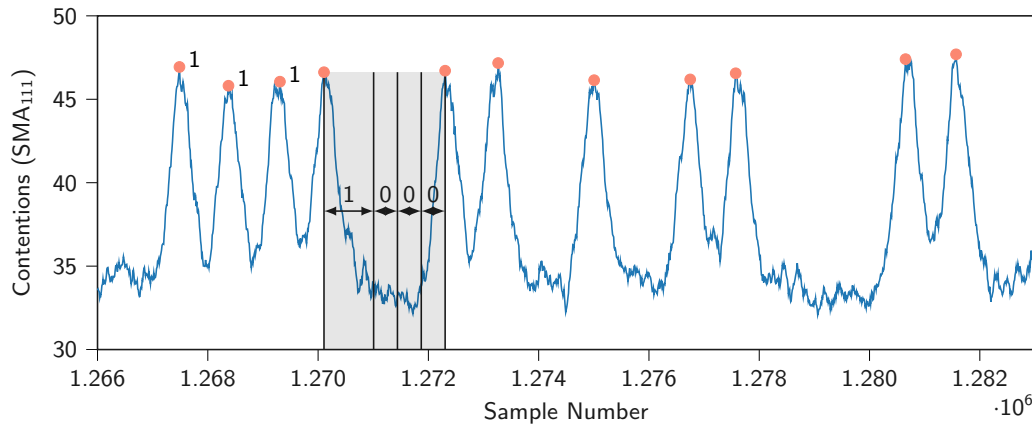


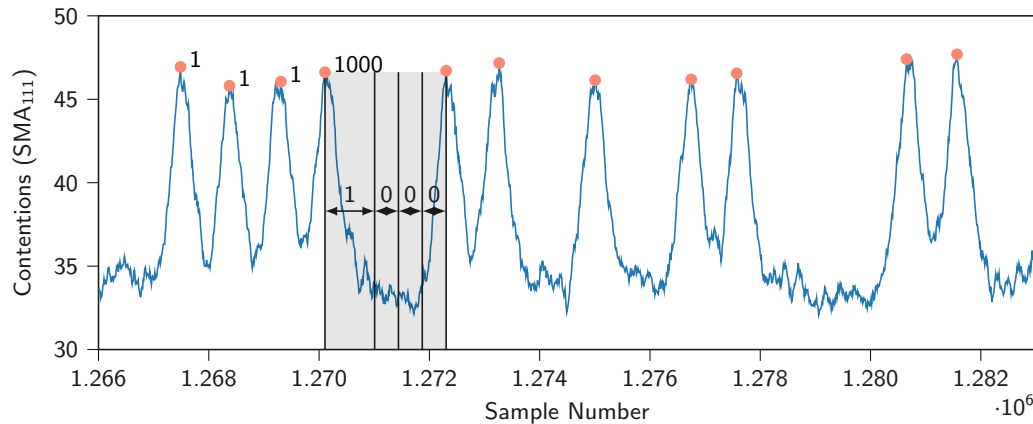


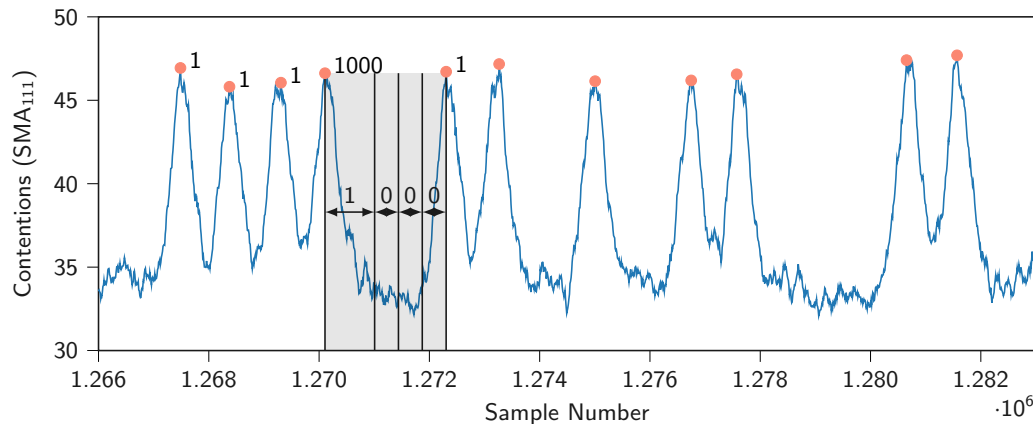


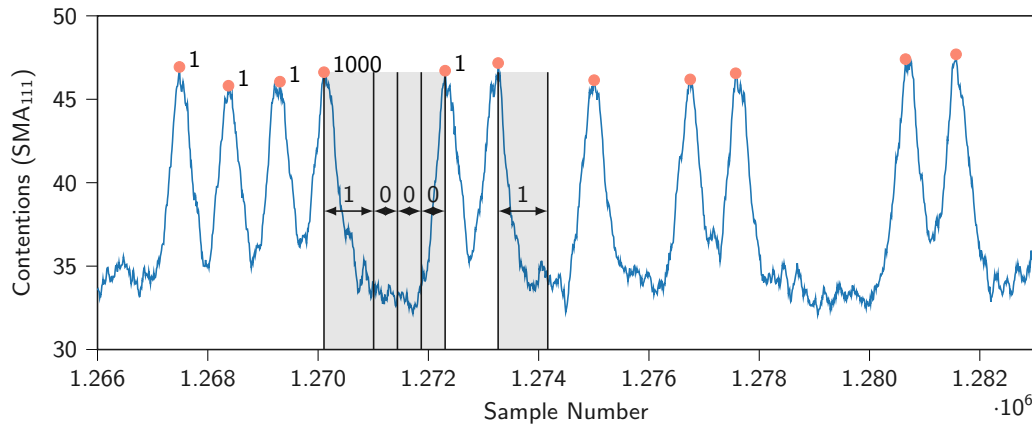


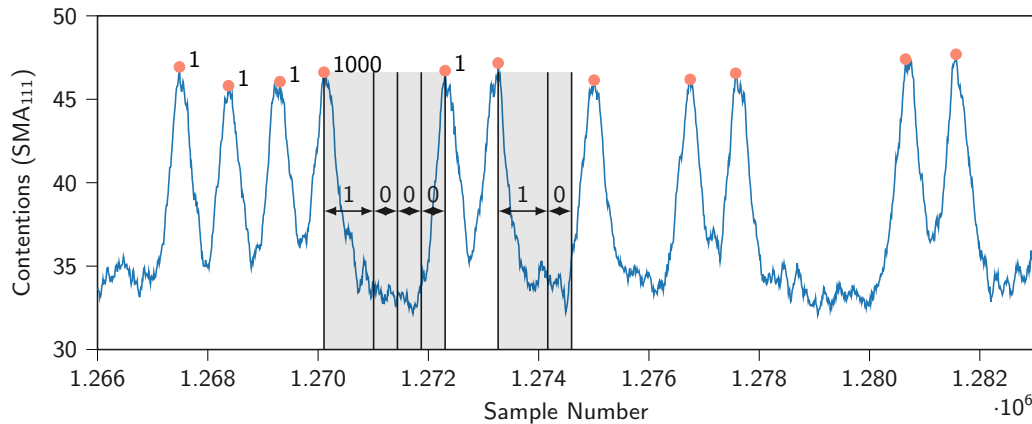




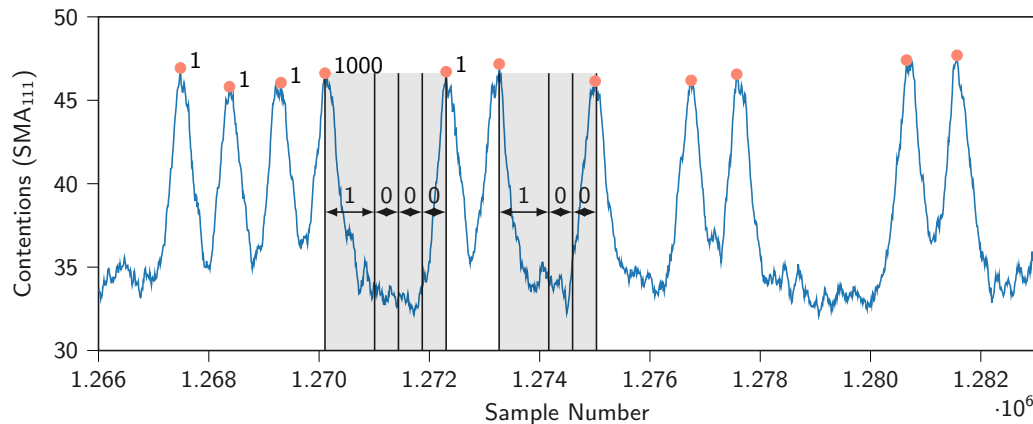


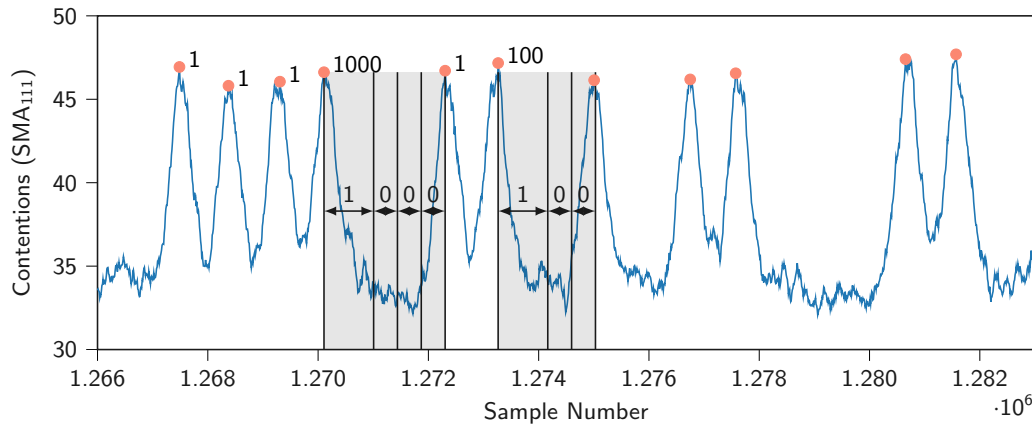






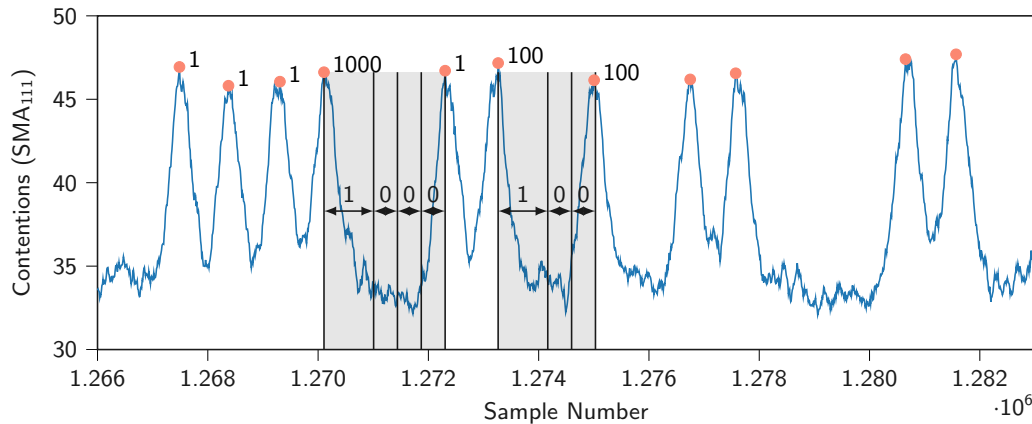
# Attacking RSA: Key Recovery

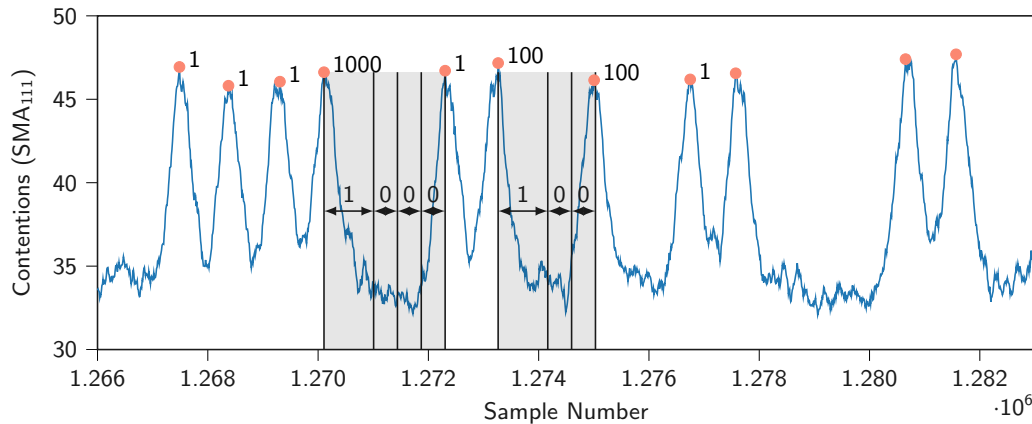




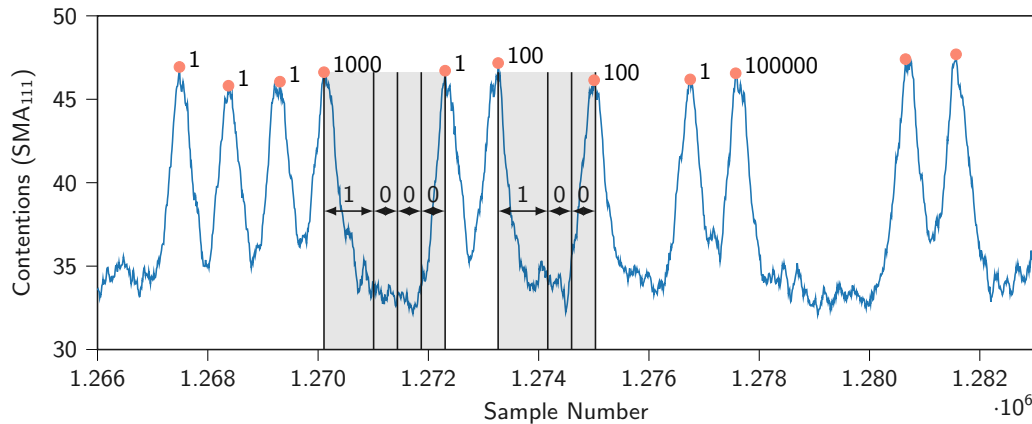


# Attacking RSA: Key Recovery

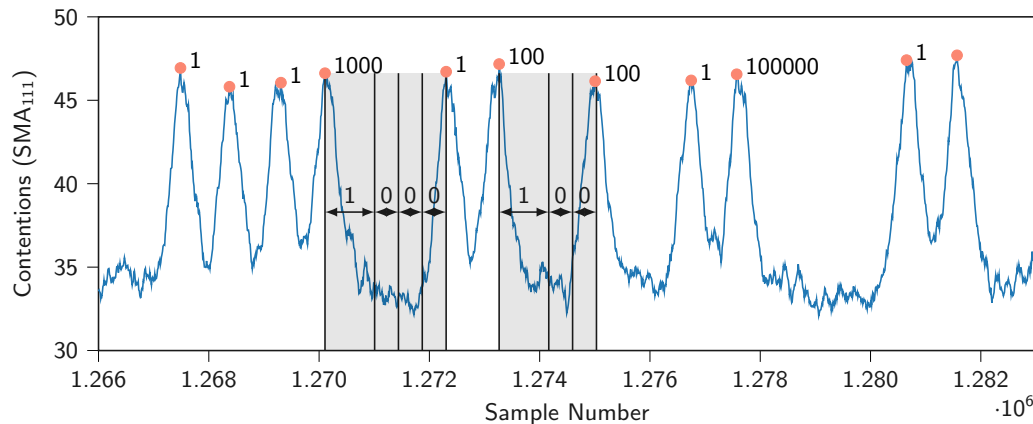




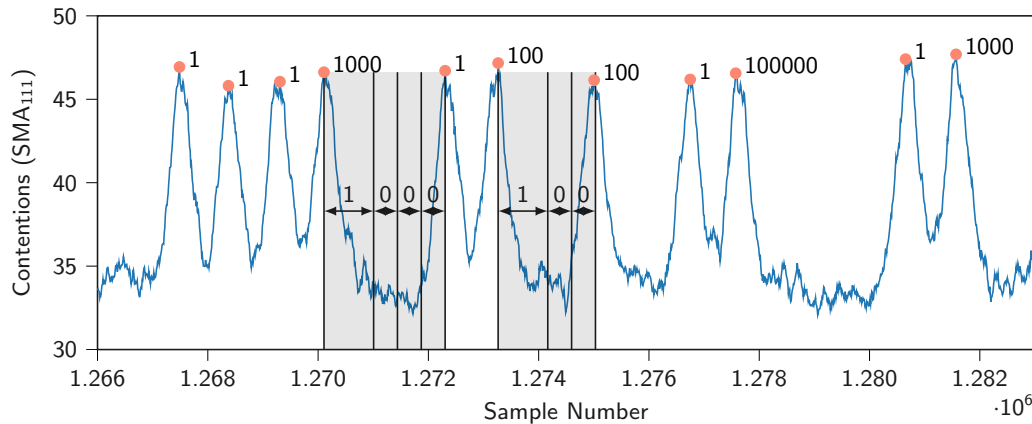
# Attacking RSA: Key Recovery



# Attacking RSA: Key Recovery

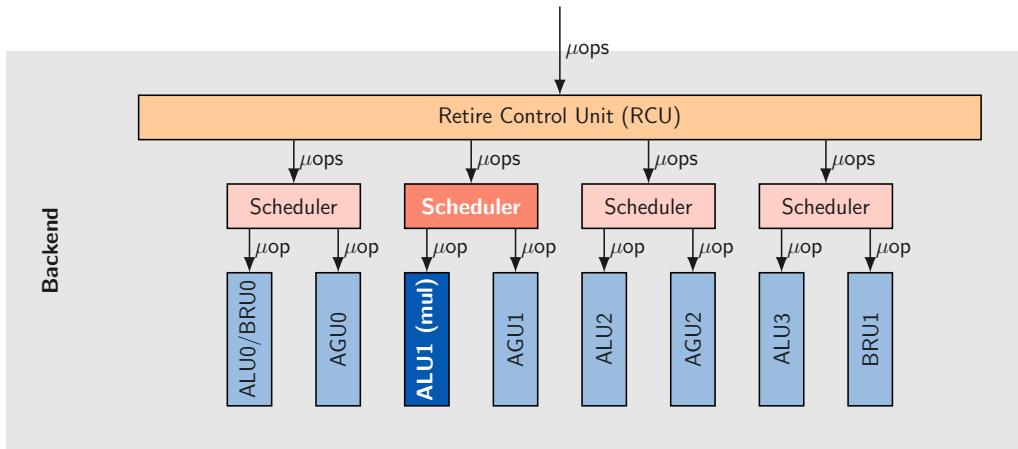


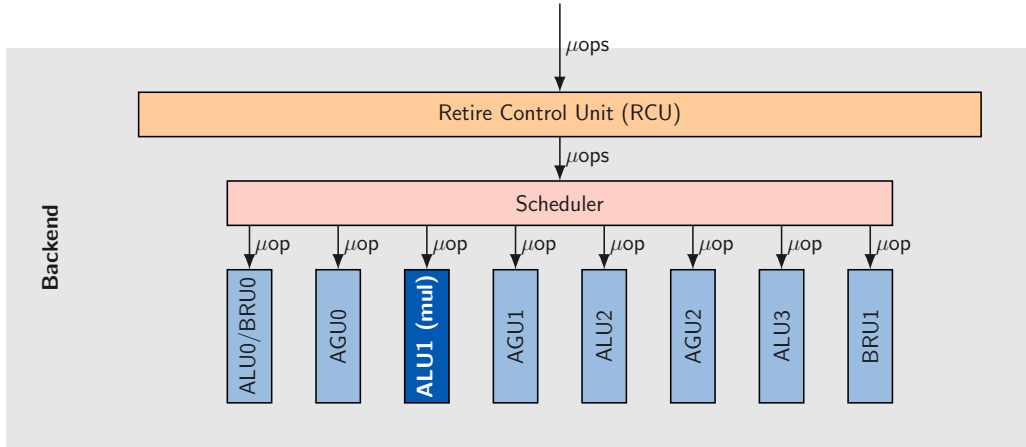
# Attacking RSA: Key Recovery



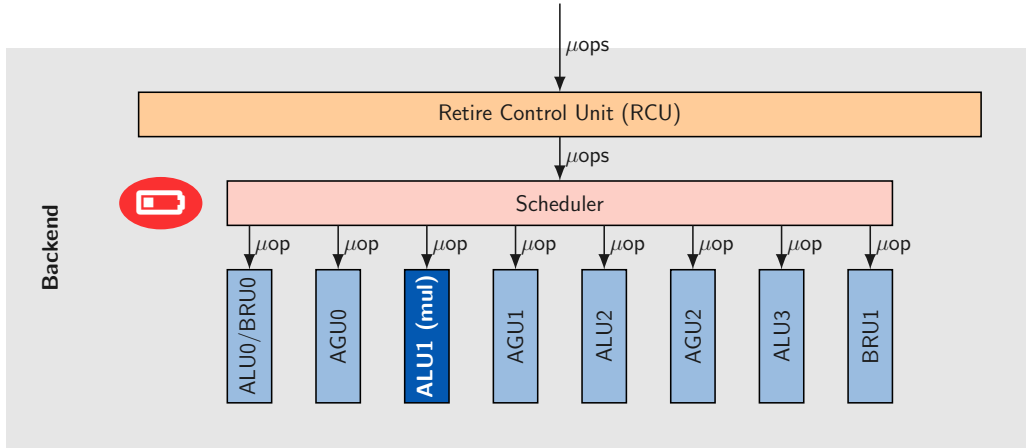
Scenario	Edit Distance	Error Rate	Recording Time
Cross-Process	4.9 bit	0.12 %	41 min
Cross-VM	17.8 bit	0.5 %	38 min

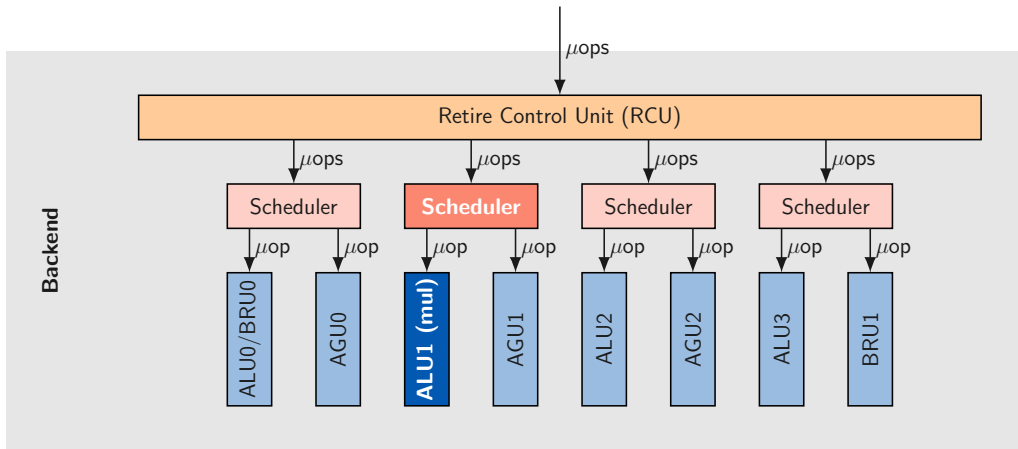
- 10 keys, generated with `openssl genrsa`
- CPU: AMD Ryzen 7 5800X (Zen 3)
- 50500 traces per key

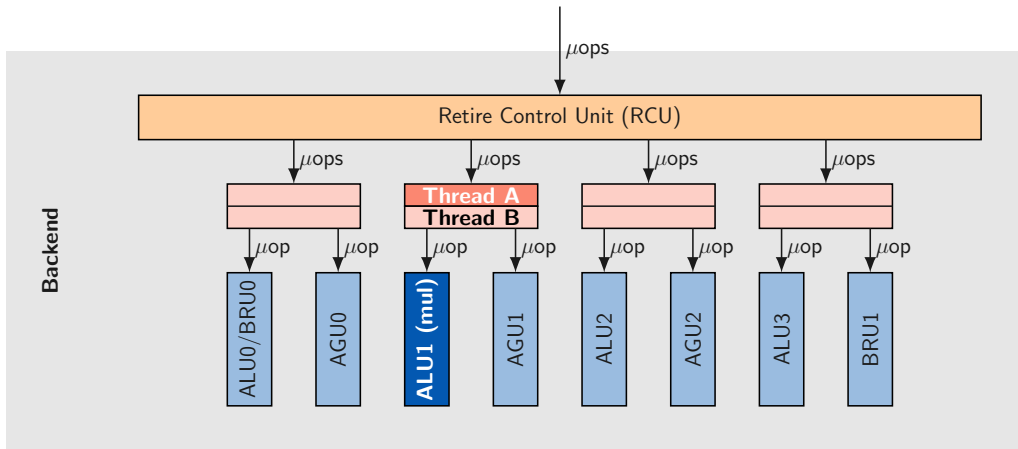


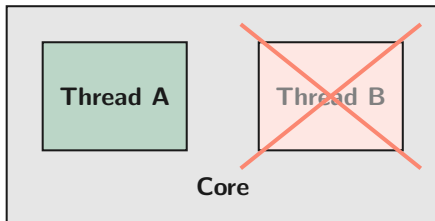












# SQUIP

## Exploiting the Scheduler Queue Contention Side Channel

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2023-05-23

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