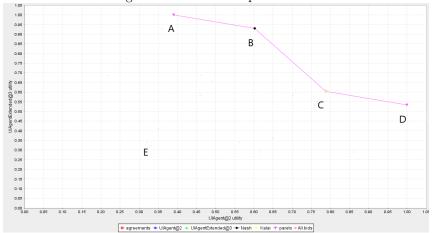
Multi-Agent Systems Lab - Report 1

Group 7

March 9, 2022

1 1st Question





(a) When defining the fairest arrangement, there are multiple possibilities. One could argue that the solution that minimizes the difference between the two parties is the fairest. In that case, point E would be the answer. This might be a fair outcome, but it is not a realistic outcome, as it is literally the outcome with the lowest combined utility. Any useful fair outcome therefore should be Pareto-optimal (its front is delimited by the pink line above).

We propose a way to get equal expected utility for both parties, but still get a Pareto-optimal outcome. For this, we need to allow for probability in the outcome of the negotiation. That is, having a certain outcome with 70% probability and another with 30% is a valid outcome. Keep in mind that we first throw the weighted coin and then choose an entire outcome, not throw a separate coin for each issue. The outcome space is no longer a set of points P, but the convex hull of P. This is because every linear combination of points of P can be achieved by simply having the outcome with the corresponding points and probability.

Note that not all points that were on the Pareto front are still on the Pareto front. In our example, C is no longer on the Pareto front, as taking B with probability 0.4 and D with probability 0.6 will have a higher expected utility for both parties.

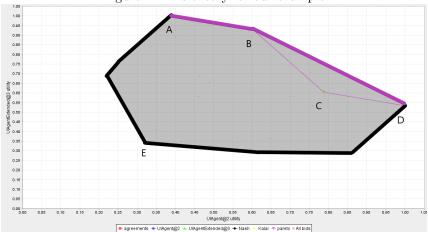
Now, the Pareto front consists of only outcomes. This is because any linear combination of two previous outcomes is now an outcome as well (and all vertices of CH(P) are in P as well).

Then, we will simply determine a Nash solution for this new problem. This Nash solution is on the Pareto front, so we can simply find the two vertices of the Pareto front that surround the Nash solution, and the fairest solution is a linear combination of those two.

In our example, the only vertices of the Pareto front are A, B, and D. A is the outcome where agent B gets everything, D is the outcome where agent A gets everything, and B is the outcome (Milan, 2 weeks, Hostel). This gives the following values for the agents:

	Agent A	Agent B
A	3.9	10
В	6	9.3
D	10	5.3

Figure 2: This theory for our example



Now, it is clear that the Nash solution does not lie on the line segment AB, so it lies on the line segment BD. Some elementary math tells us that the utility values are (7.65, 7.65) and the solution is 0.4125B + 0.58875D, so if we will take B with probability 0.4125 and D with probability 0.5887, we get the highest product of expected utilities.

(b) When the Random Bidders negotiate against themselves, we concluded that the Pareto Optimal Outcome is not consistently reached, since the acceptance strategy is also random; even if a bid is made on the Pareto Frontier, there is no guarantee that it will be accepted. The following figures (3 through 6) graph the aforementioned negotiations. Even with high minimum bidding targets, the bidders do not seem to reach their optimal outcomes (it does happen, but as mentioned before, it is extremely inconsistent).

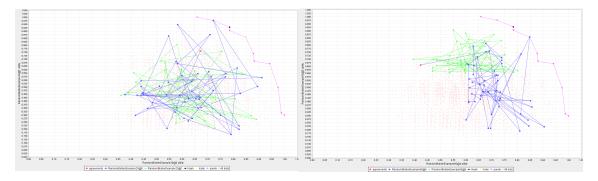


Figure 3: Minimum Bidding Target = 0.20

Figure 4: Minimum Bidding Target = 0.60

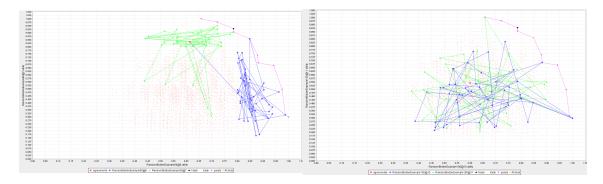


Figure 5: Minimum Bidding Target = 0.80

Figure 6: Minimum Bidding Target = 1.00

Notably, the graphs are extremely similar for the minimum targets of 0.2 and 1.0, since the agent attempts 100 times to find a bid with a fitting utility, and it rarely ever finds one with 100% utility, it is as if it were completely random.

Lastly, when doing the same negotiations, but using the Boulware and Conceder agents, the Pareto Optimal Outcome is always reached, primarily benefiting Boulware. When generating



Figure 7: Boulware vs Conceder

bids they generate a random bid with utility around their target, which decreases with time throughout the negotiation. In this case, the Boulware Party gets the better deal, as their target utility goes down much more slowly than the Conceder's.