

Context-Free Languages

Now refer to the description of **Walks** on page 22.

Let SAW be the language of strings over the alphabet $\{N, S, E, W\}$ that represent self-avoiding walks.

Problem 8. [8 marks]

Prove or disprove: The language SAW is context-free.

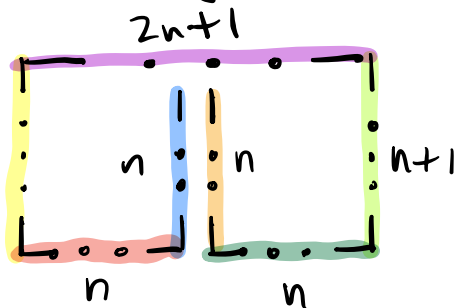
Your submission can be typed or hand-written, but it must be in PDF format and saved as a file prob8.pdf.

Suppose the language SAW is context-free
and let k be the number of non-terminal symbols
in a CFG for SAW.
now let $N = 2^k$ for simplicity.

Then by the pumping lemma for context free
languages every $w \in \text{SAW}$ such that $|w| > 2^k - 1$
can be written in the form $w = uvxyz$
where $|v| \geq 1$, $|vxy| \leq 2^k$ and all $i \geq 0$, such
that $uv^ixy^iz \in \text{SAW}$.

It's important to note that the rules on SAW
dictate that every substring of $w \in \text{SAW}$ must follow
that $|N| \neq |S| \vee |E| \neq |W|$

now: we choose a string $w = S^n W^n N^{n+1} E^{2n+1} S^{n+1} W^n N^n$
 which by map diagram we can represent as:

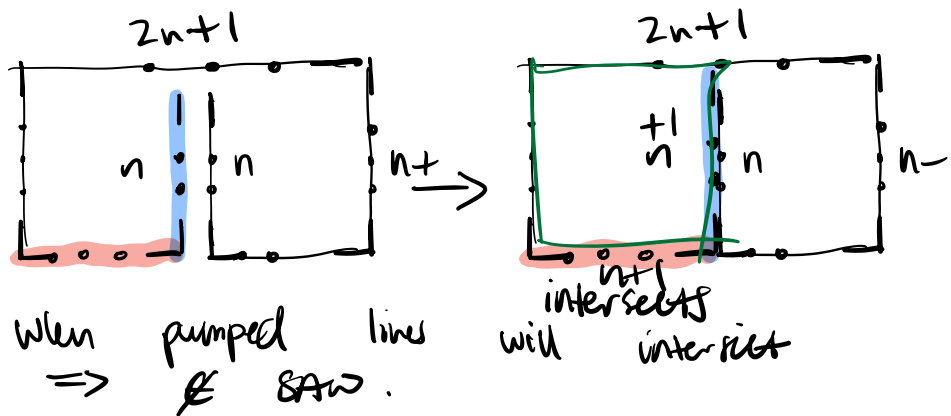


Case 1: $S^n W^n \dots$

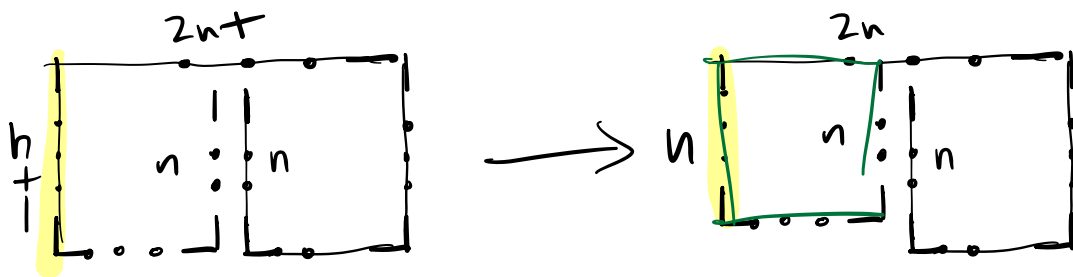
pumping S - if $i=1 \Rightarrow S^{n+1} W^n N^{n+1} E^n$,

Case 2: $|S| = |W|, |W| = |E| \Rightarrow \notin \text{SAW}$
 pumping W - if $i=1 \Rightarrow W^{n+1} N E^{2n+1} S^{n+1} W^n$
 equal parts $\Rightarrow \notin \text{SAW}$.

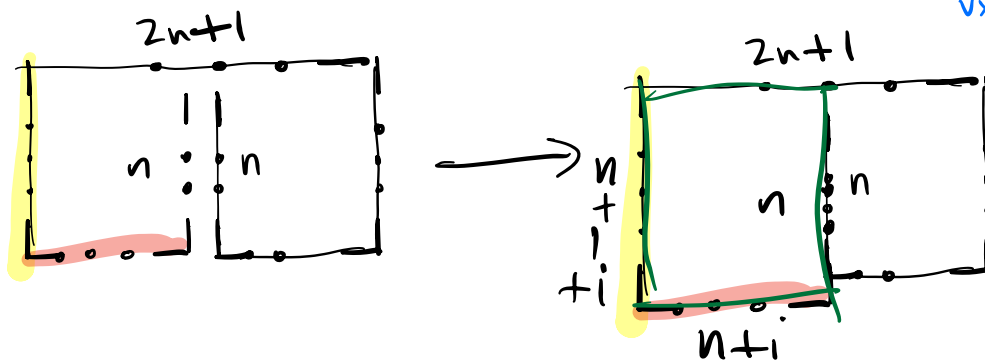
Case 3:
 pumping WNS -



Case 4: pumping N^{n+1} , if we let $i=0$ then we will reach $S^n W^n N^n E^n$ as a substring for which does not satisfy SAW.



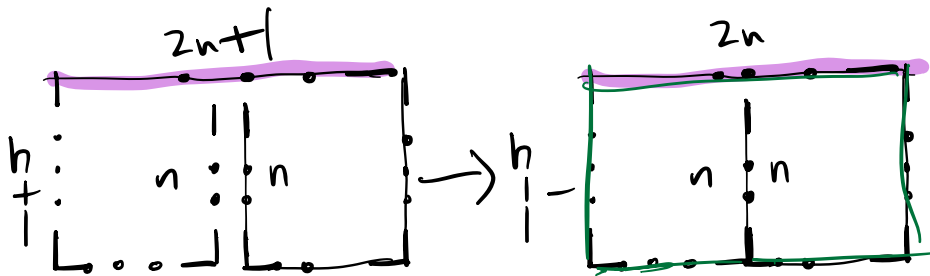
Case 5: pumping combination of $W^N N^{N+1}$
 vxy .



Any combination of W^N and N^N will result in an intersection occurring such that some substring of W

Case 6: pumping E , if we let $i=0$ then

$W^N N^{N+1} E^{2n} S^{N+1} W^N$, which forms a a substring that does not exist in a word that is an element of SAW



* I ran out of time, the rest of the cases are mostly similar to the ones provided except on the opposite side -

Apologies for the hand wariness, only came up with this late Friday night.