

Principles of Programming

(4190.306)

Chung-Kil Hur (허충길)

Department of Computer Science and Engineering
Seoul National University

Syllabus

➤ Lecture & Practice Session

- Mon & Wed, 9:00 ~ 10:50 (302-208)
- <https://github.com/snu-sf-class/pp201802>
- Bring your laptop to lectures & practice sessions

➤ Instructor

- Chung-Kil Hur
- <http://sf.snu.ac.kr/gil.hur/>

➤ Teaching Assistant

- Yonghyun Kim, Sung-Hwan Lee, Minki Cho, Dongjoo Kim.
Contact TAs at pp2018fall@sf.snu.ac.kr

➤ Grading (tentative)

- Attendance: 5%
- Assignments: 25%
- Midterm exam: 30%
- Final exam: 40%

Free Laptop Rental

You can rent a laptop from the department for free.

Who am I?

➤ Prof. Chung-Kil (Gil) Hur [허충길]

- Education: KAIST (B.S.), Univ of Cambridge (Ph.D.)
- Software Foundations Lab
<http://sf.snu.ac.kr>
- Research Topics
 - Software Verification
 - Low-level Language Semantics (C/C++/LLVM/Rust)
 - Relaxed-Memory Concurrency
- Our collaborators
 - [UK] U of Cambridge, Microsoft Research Cambridge
 - [Germany] Max Planck Institute for Software Systems
 - [France] INRIA
 - [USA] Princeton, Yale, Upenn, Utah, Google, IBM, Mozilla.
- Publications
 - 10 top conference papers (last 5 years at SNU).
PLDI(5), POPL(2), OOPSLA(1), ICFP(1), AAAI(1)
- (High school) Bronze medal in IMO 1994.

Introduction

Overview

➤ Part 1

Functional Programming with Function Applications

➤ Part 2

Object-Oriented Programming

➤ Part 3

Type Classes for Specifications

➤ Part 4

Imperative Programming with Memory Updates

Imperative vs. Functional Programming

➤ Imperative Programming

- Computation by memory reads/writes
- Sequence of read/write operations
- Repetition by loop
- More procedural
- Easier to write efficient code

```
sum = 0;
i = n;
while (i > 0) {
    sum = sum + i;
    i = i - 1;
}
```

➤ Functional Programming

- Computation by function application
- Composition of function applications
- Repetition by recursion
- More declarative
- Easier to write safe code

```
def sum(n) =
    if (n <= 0)
        0
    else
        n + sum(n-1)
```

Both Imperative & Functional Style Supported

- Many languages support both imperative & functional style
 - More imperative: Java, Javascript, C++, Python, Rust, ...
 - More functional: OCaml, SML, Lisp, Scheme, ...
 - Middle: Scala
 - Purely functional: Haskell

Object-Oriented Programming

➤ Object-Oriented Programming

- Classes/Objects: data + methods
- Can separate Specification & Implementation

```
class Point(x: Int, y: Int) {  
  // data  
  val px : Int = x  
  val py : Int = y  
  // methods  
  def plus(q : Point) : Point =  
    new Point(px + q.px, py + q.py)  
}  
val p1 : Point = new Point(3,4)  
val p2 : Point = new Point(1,10)  
val p3 : Point = p1.plus(p2)
```

Why Scala?

➤ Why Scala?

- Equally well support both imperative & functional style
- Many advanced features (both OOP & Type class supported)
- Compatible with Java

PART 1

Functional Programming with Function Applications

Names, Functions and Evaluations

Values, Expressions, Names

➤ Types and Values

- A type is a set of values
- `Int`: $\{-2147483648, \dots, -1, 0, 1, \dots, 2147483647\}$ //32-bit integers
- `Double`: 64-bit floating point numbers // real numbers in practice
- `Boolean`: $\{\text{true}, \text{false}\}$
- ...

➤ Expressions

- Composition of
values, names, primitive operations, ...

➤ Name Binding

- Binding expressions to names

➤ Examples

```
def a = 1 + (2 + 3)
def b = 3 + a * 4
```

Evaluation

➤ Evaluation

- Reducing an expression into a value
- Strategy
 1. Take a name or an operator (outer to inner)
 2. (name) Replace the name with its associated expression
 3. (name) Evaluate the expression
 4. (operator) Evaluate its operands (left to right)
 5. (operator) Apply the operator to its operands

➤ Examples

$5+b \sim 5+(3+a*4) \sim 5+(3+(1+(2+3))*4) \sim \dots \sim 32$

Functions and Substitution

➤ Functions

- Expressions with Parameters
- Binding functions to names

```
def f(x: Int): Int = x + a
```

➤ Evaluation by substitution

- ...
- (function) Evaluate its operands (left to right)
- (function)
Replace the function application by the expression of the function
Replace its parameters with the operands

$$5 + f(f(3) + 1) \sim 5 + f((3 + a) + 1) \sim \dots \sim 5 + f(10) \sim 5 + (10 + a) \\ \sim \dots \sim 21$$

Simple Recursion

➤ Recursion

- Use X in the definition of X
- Powerful mechanism for repetition
- Nothing special but just rewriting

```
def sum(n: Int) : Int =  
  if (n <= 0)  
    0  
  else  
    n + sum(n-1)
```

```
sum(2) ~ if (2<=0) 0 else (2+sum(2-1)) ~  
2+sum(1) ~ 2+(if (1<=0) 0 else (1+sum(1-1))) ~  
2+(1+sum(0)) ~ 2+(1+(if (0<=0) 0 else (0+sum(0-1))))  
~ 2+(1+0) ~ 3
```


Termination/Divergence

Evaluation may not terminate

➤ Termination

- An expression may reduce to a value

➤ Divergence

- An expression may reduce forever

```
def loop: Int = loop
```

```
loop ~ loop ~ loop ~ ...
```

Evaluation strategy: Call-by-value, Call-by-name

$f(e1, e2)$

➤ Call-by-value

- Evaluate the arguments first, then apply the function to them

➤ Call-by-name

- Just apply the function to its arguments, without evaluating them.

```
def square (x: Int) = x * x
```

```
[cbv]square(1+1) ~ square(2) ~ 2*2 ~ 4
```

```
[cbn]square(1+1) ~ (1+1)*(1+1) ~ 2*(1+1) ~ 2*2 ~ 4
```

CBV, CBN: Differences

➤ Call-by-value

- Evaluates arguments once

➤ Call-by-name

- Do not evaluate unused arguments

➤ Question

- Do both always result in the same value?

Scala's evaluation strategy

➤ Call-by-value

- By default

➤ Call-by-name

- Use “=>”

```
def one(x: Int, y: =>Int) = 1
```

```
one(1+2, loop)
```

```
one(loop, 1+2)
```

Scala's name binding strategy

➤ Call-by-value

- Use “val” (also called “field”) e.g. `val x = e`
- Evaluate the expression first, then bind the name to it

➤ Call-by-name

- Use “def” (also called “method”) e.g. `def x = e`
- Just bind the name to the expression, without evaluating it
- Mostly used to define functions

```
def a = 1 + 2 + 3
```

```
val a = 1 + 2 + 3 // 6
```

```
def b = loop
```

```
val b = loop
```

```
def f(a: Int, b: Int): Int = a*b - 2
```

Conditional Expressions

➤ If-else

- `if (b) e1 else e2`
- `b` : Boolean expression
- `e1, e2`: expressions of the same type

➤ Rewrite rules:

- `if (true) e1 else e2 → e1`
- `if (false) e1 else e2 → e2`

```
def abs(x: Int) = if (x >= 0) x else -x
```

Boolean Expressions

➤ Boolean expression

- true, false
- !b
- b && b
- b || b
- e <= e, e >= e, e < e, e > e, e == e, e != e

➤ Rewrite rules:

- !true → false
- !false → true
- true && b → b
- false && b → false
- true || b → true
- false || b → b

true && (loop == 1) ~ loop == 1 ~ loop == 1

Exercise: and, or

➤ Write two functions

- `and(x,y) == x && y`
- `or(x,y) == x || y`
- Do not use `&&`, `||`

```
and(false,loop==1)
```

```
~ if (false) loop==1 else false
```

```
~ false
```

```
and(true,loop==1)
```

```
~ if (true) loop==1 else false
```

```
~ loop==1 ~ loop==1 ...
```


Solution

```
def and(x: Boolean, y: =>Boolean) =  
  if (x) y else false
```

```
def or(x: Boolean, y: =>Boolean) =  
  if (x) true else y
```

Exercise: square root calculation

➤ Calculate square roots with Newton's method

```
def isGoodEnough(guess: Double, x: Double) =  
    ??? // guess*guess is 99.9% close to x
```

```
def improve(guess: Double, x: Double) =  
    (guess + x/guess) / 2
```

```
def sqrtIter(guess: Double, x: Double): Double =  
    ??? // repeat improving guess until it is good enough
```

```
def sqrt(x: Double) =  
    sqrtIter(1, x)
```

```
sqrt(2)
```

Solution

➤ Calculate square roots with Newton's method

```
def isGoodEnough(guess: Double, x: Double) =  
    guess*guess/x > 0.999 && guess*guess/x < 1.001
```

```
def improve(guess: Double, x: Double) =  
    (guess + x/guess) / 2
```

```
def sqrtIter(guess: Double, x: Double): Double =  
    if (isGoodEnough(guess,x)) guess  
    else sqrtIter(improve(guess,x),x)
```

```
def sqrt(x: Double) =  
    sqrtIter(1, x)
```

```
sqrt(2)
```

Blocks and Name Scoping

Blocks in Scala

➤ Block

- ```
{ val x1 = e1
 def x2 = e2
 e
}
```
- Is an expression
- Allow nested name binding
- Allow arbitrary order of “**def**”s, but not “**val**”s (think about why)

# Scope of names

## ➤ Block

```
val t = 0
def f(x: Int) = t + g(x)
def g(x: Int) = x * x
val x = f(5)
val r = {
 val t = 10
 val s = f(5)
 s - t
}
t + r
```

- A definition inside a block is only accessible within the block
- A definition inside a block shadows definitions of the same name outside the block
- A definition is accessible inside a nested block unless it is shadowed
- A function is evaluated under the environment where it is defined, not the environment where it is invoked.

# Rewriting for Blocks

```
1: val t = 0
2: def f(x: Int) = t + g(x)
3: def g(x: Int) = x*x
4: val x = f(5)
```

```
5: val r = {
6: val t = 10
7: val s = f(5)
8: s - t
9: t + r
```

## ➤ Evaluation with Environment

$[], 1 \sim [t=0], 2 \sim [\dots, f=(x)t+g(x)], 3 \sim [\dots, g=(x)x*x], 4$   
 $\sim [\dots, x=25], 5 \sim [\dots]:[], 6 \sim [\dots]:[t=10], 7$   
 $\sim [\dots]:[\dots, s=25], 8 \sim [\dots, r=15], 9 \sim [\dots], 15$

4:  $[t=0, f=\dots, g=\dots]:[x=5], t+g(x) \sim 0+g(5) \sim 25$

$g(5): [t=0, f=\dots, g=\dots]:[x=5], x*x \sim 5*5 \sim 25$

7:  $[t=0, f=\dots, g=\dots, x=25]:[x=5], t+g(x) \sim 0+g(5) \sim 25$

$g(5): [t=0, f=\dots, g=\dots, x=25]:[x=5], x*x \sim 5*5 \sim 25$

Note: The same name cannot be defined twice in the same block

## Example: def with no arguments

```
1: val t = 0
2: def x = t+t // is treated as def x() = t
3: val r = {
4: val t = 10
5: x } // is treated as x()
```

### ➤ Evaluation with Environment

$[], 1 \sim [t=0], 2 \sim [\dots, x=()t+t], 3 \sim [\dots]:[], 4$   
 $\sim [\dots]:[t=10], 5 \sim [\dots, r=0], 6$

5:  $[t=0, x=()t+t]:[], t+t \sim 0$



# Semi-colons and Parenthesis

## ➤Block

- Can write two definitions/expressions in a single line using ;
- Can write one definition/expression in two lines using (), but can omit () when clear

*// ok*

```
val r = {
 val t = 10; val s = square(5); t +
 s }
```

*// Not ok*

```
val r = {
 val t = 10; val s = square(5); t
 + s }
```

*// ok*

```
val r = {
 val t = 10; val s = square(5); (t
 + s) }
```

# Exercise: Writing Better Code using Blocks

➤ Make the following code better

```
def isGoodEnough(guess: Double, x: Double) =
 guess*guess/x > 0.999 && guess*guess/x < 1.001
```

```
def improve(guess: Double, x: Double) =
 (guess + x/guess) / 2
```

```
def sqrtIter(guess: Double, x: Double): Double = {
 if (isGoodEnough(guess,x)) guess
 else sqrtIter(improve(guess,x),x)
}
```

```
def sqrt(x: Double) =
 sqrtIter(1, x)
```

```
sqrt(2)
```

# Solution

```
def sqrt(x: Double) = {
 def sqrtIter(guess: Double, x: Double): Double = {
 if (isGoodEnough(guess, x)) guess
 else sqrtIter(improve(guess, x), x)
 }
 def isGoodEnough(guess: Double, x: Double) = {
 val ratio = guess * guess / x
 ratio > 0.999 && ratio < 1.001
 }
 def improve(guess: Double, x: Double) =
 (guess + x/guess) / 2

 sqrtIter(1, x)
}

sqrt(2)
```

# Lazy Call-By-Value

# Lazy call-by-value

## ➤ Lazy call-by-value

- Use “lazy val” e.g. `lazy val x = e`
- Evaluate the expression **first time it is used**, then bind the name to it

```
def f(c: Boolean, i: =>Int): Int = {
 lazy val iv = i
 if (c) 0
 else iv * iv * iv
}
```

```
f(true, {println("ok"); 100+100+100+100})
f(false, {println("ok"); 100+100+100+100})
```

# Tail Recursion

# Recursion needs care

## ➤ Summation function

- Write a summation function `sum` such that  
 $\text{sum}(n) = 1+2+\dots+n$
- Test  
`sum(10), sum(100), sum(1000), sum(10000),`  
`sum(100000), sum(1000000)`
- What's wrong? (Think about evaluation)

# Recursion: Try

```
def sum(n: Int): Int =
 if (n <= 0) 0 else (n+sum(n-1))
```



# Recursion: Tail Recursion

```
import scala.annotation.tailrec
```

```
def sum(n: Int): Int = {
 @tailrec def sumItr(res: Int, m: Int): Int =
 if (m <= 0) res else sumItr(m+res,m-1)
 sumItr(0,n)
}
```

# Higher-Order Functions

# Functions as Values

## ➤ Functions

- Functions are normal values of function types  $(A_1, \dots, A_n \Rightarrow B)$ .
- They can be copied, passed and returned.
- Functions that take functions as arguments or return functions are called higher-order functions.
- Higher-order functions increase code reusability.

# Examples

```
def sumLinear(n: Int): Int =
 if (n <= 0) 0 else n + sumLinear(n-1)
```

```
def sumSquare(n: Int): Int =
 if (n <= 0) 0 else n*n + sumSquare(n-1)
```

```
def sumCubes(n: Int): Int =
 if (n <= 0) 0 else n*n*n + sumCubes(n-1)
```

Q: How to write reusable code?

# Examples

```
def sum(f: Int=>Int, n: Int): Int =
 if (n <= 0) 0 else f(n) + sum(f, n-1)
```

```
def linear(n: Int) = n
def square(n: Int) = n * n
def cube(n: Int) = n * n * n
```

```
def sumLinear(n: Int) = sum(linear, n)
def sumSquare(n: Int) = sum(square, n)
def sumCubes(n: Int) = sum(cube, n)
```

# Anonymous Functions

## ➤ Anonymous Functions

- Syntax

$(x_1: T_1, \dots, x_n: T_n) \Rightarrow e$

or

$(x_1, \dots, x_n) \Rightarrow e$

```
def sumLinear(n: Int) = sum((x: Int) => x, n)
def sumSquare(n: Int) = sum((x: Int) => x * x, n)
def sumCubes(n: Int) = sum((x: Int) => x * x * x, n)
```

Or simply

```
def sumLinear(n: Int) = sum((x) => x, n)
def sumSquare(n: Int) = sum((x) => x * x, n)
def sumCubes(n: Int) = sum((x) => x * x * x, n)
```

# Exercise

```
def sum(f: Int=>Int, a: Int, b: Int): Int =
 if (a <= b) f(a) + sum(f, a+1, b) else 0
```

```
def product(f: Int=>Int, a: Int, b: Int): Int =
 if (a <= b) f(a) * product(f, a+1, b) else 1
```

DRY (Do not Repeat Yourself) using a higher-order function, called “mapReduce”.

# Exercise

```
def mapReduce(reduce: (Int,Int)=>Int, inival: Int,
 f: Int=>Int, a: Int, b: Int): Int = {
 if (a <= b) reduce(f(a),mapReduce(reduce,inival,f,a+1,b))
 else inival
}
```

```
def sum(f: Int=>Int, a: Int, b: Int): Int =
 mapReduce((x,y)=>x+y,0,f,a,b)
```

```
def product(f: Int=>Int, a: Int, b: Int): Int =
 mapReduce((x,y)=>x*y,1,f,a,b)
```



# Closures for functional values

```
1: val t = 0
2: val f : Int=>Int = { // Note: What if using “def f” ?
3: val t = 10
4: def g(x: Int) : Int = x + t
5: g _ }
6: f(20)
```

\* Try: Evaluation without Closures

```
[], 1 ~ [t=0], 2 ~ [...]:[], 3 ~ [...]:[t=10], 4
~ [...]:[...,g=(x)x+t], 5 ~ [t=0,f=(x)x+t], 6 ~ [...], 20
6: [t=0,f=(x)x+t]:[x=20], x+t ~ 20+0 ~ 20
```

\* Evaluation with Closures

```
E1[], 1 ~ E1[t=0], 2 ~ E1[...] : E2[], 3 ~ E1[...] : E2[t=10], 4
~ E1[...] : E2[t=10,g=(x)x+t], 5
~ E1[t=0,f=((x)x+t,E2)], 6 ~ E1[...], 30
6: E1:E2:E3[x=20], x+t ~ 20+10 ~ 30
```

# Problem

// should be allowed

{

def f(x: Int) = g(x)

def g(x: Int) = 10

val x = f(10)

x

}

// should not be allowed

{

def f(x: Int) = g(x)

val x = f(10)

def g(x: Int) = 10

x

}

# Safety Checking

Safety checking for ruling out accessing undefined values is simple.

- For “val x = e”,  
all names in “e” should be defined before this definition.
- For “def x = e”,  
all names in “e” should be defined before the next “val” definition.

/\* The following code passes the safety checking \*/

```
{ def f(x: Int) = g(x)
 def g(x: Int) = 10
 val a = 10
 f(10) }
```

/\* The following code fails at the safety checking \*/

```
{ def f(x: Int) = g(x)
 val a = 10
 def g(x: Int) = 10
 f(10) }
```

# (Parameterized) code vs. values

- Functions defined using “def” are not values but parameterized code/expressions.
- Anonymous functions are values.
- Code  $f$  is converted to a value  $(f \ \_)$ .
- The compiler often infers and inserts missing conversions automatically.

- Anonymous functions can be seen as syntactic sugar:

$(x:T) \Rightarrow e$

is equivalent to

$\{ \text{def } \text{noname}(x:T) = e; (\text{noname } \_) \}$

where `noname` is not used in  $e$ .

- One can even write a recursive anonymous function in this way.
- Q: what's the difference between param. exps and function values?  
A: function values are “closures” (ie, param. exp. + env.)
- Q: how to implement call-by-name?  
A: The argument expression is converted to a closure.

# Example: call by name with closures

```
1: val t = 0
2: def f(x: =>Int) = t + x
3: val r = {
4: val t = 10
5: f(t*t) } // t*t is treated as ()=>t*t
```

## ➤ Evaluation with Closures

$E1[], 1 \sim E1[t=0], 2 \sim E1[t=0, f=(x:=>Int)t+x], 3$   
 $\sim E1[\dots]:E2[], 4 \sim E1[\dots]:E2[t=10], 5 \sim E1[\dots, r=100], 6$

5:  $E1:E3[x=(( )t*t, E2)], t+x$

$\sim 0 + x \sim 0 + 100 \sim 100$

x:  $E1:E2, t*t \sim 10*10 \sim 100$

# Currying

# Motivation

```
def sum(f: Int=>Int, a: Int, b: Int): Int =
 if (a <= b) f(a) + sum(f, a+1, b) else 0
def linear(n: Int) = n
def square(n: Int) = n * n
def cube(n: Int) = n * n * n
def sumLinear(a: Int, b: Int) = sum(linear, a, b)
def sumSquare(a: Int, b: Int) = sum(square, a, b)
def sumCubes(a: Int, b: Int) = sum(cube, a, b)
```

We want the following. How?

```
def sumLinear = sum(linear)
def sumSquare = sum(square)
def sumCubes = sum(cube)
```

# Solution

```
def sum(f: Int=>Int): (Int,Int)=>Int = {
 def sumF(a: Int, b: Int): Int =
 if (a <= b) f(a) + sumF(a+1, b) else 0
 sumF
}
```

```
def sumLinear = sum(linear)
def sumSquare = sum(square)
def sumCubes = sum(cube)
```



# Benefits

```
def sumLinear = sum(linear)
def sumSquare = sum(square)
def sumCubes = sum(cube)
```

`sumSquare(3, 10) + sumCubes(5, 20)`

We don't need to define the wrapper functions.

`sum(square)(3, 10) + sum(cube)(5, 20)`

# Multiple Parameter List

```
def sum(f: Int=>Int): (Int,Int)=>Int = {
 def sumF(a: Int, b: Int): Int =
 if (a <= b) f(a) + sumF(a+1, b) else 0
 sumF _
}
```

Or simply:

```
def sum(f: Int=>Int)(a: Int, b: Int): Int =
 if (a <= b) f(a) + sum(f)(a+1, b) else 0
```

Note that `sum(f)` is just code.

`(sum(f) _)` creates a closure like `(sumF _)`

The following code is inefficient. Think about why.

```
def sum(f: Int=>Int): (Int,Int)=>Int =
 (a,b) => if (a <= b) f(a) + sum(f)(a+1, b) else 0
```

# Currying and Uncurrying

- A function of type

$$(T_1, T_2, \dots, T_n) \Rightarrow T$$

can be turned into one of type

$$T_1 \Rightarrow (T_2 \Rightarrow (\dots \Rightarrow (T_n \Rightarrow T) \dots))$$

- This is called “currying” named after Haskell Brooks Curry.
- The opposite direction is called “uncurrying”.

# Currying using Anonymous Functions

```
def foo(x: Int, y: Int, z: Int)(a: Int, b: Int) =
 x + y + z + a + b
```

```
val f1 = (x: Int, z: Int, b: Int) => foo(x, 1, z)(2, b)
```

```
val f2 = foo(_: Int, 1, _: Int)(2, _: Int)
```

```
val f3 = (x: Int, z: Int) => ((b: Int) => foo(x, 1, z)(2, b))
```

```
f1(1, 2, 3)
```

```
f2(1, 2, 3)
```

```
f3(1, 2)(3)
```

# Exercise

Curry the mapReduce function.

```
def mapReduce(reduce: (Int, Int) => Int, inival: Int,
 f: Int => Int, a: Int, b: Int): Int = {
 if (a <= b) reduce(f(a), mapReduce(reduce, inival, f, a+1, b))
 else inival
}
```

```
def sum(f: Int => Int, a: Int, b: Int): Int =
 mapReduce((x, y) => x+y, 0, f, a, b)
```

```
def product(f: Int => Int, a: Int, b: Int): Int =
 mapReduce((x, y) => x*y, 1, f, a, b)
```

# Solution

```
def mapReduce(combine:(Int,Int)=>Int,inival: Int)
 (f: Int=>Int) (a: Int, b: Int): Int = {
 if (a <= b) combine(f(a),mapReduce(combine,inival)(f)(a+1,b))
 else inival
}
```

// need to make a closure since mapReduce is param. code.

```
def sum = mapReduce((x,y)=>x+y,0) _
```

// val is better than def. Think about why.

```
val product = mapReduce((x,y)=>x*y,1) _
```

# Exceptions

# Exception & Handling

```
class factRangeException(val arg: Int) extends Exception
```

```
def fact(n : Int): Int =
 if (n < 0) throw new factRangeException(n)
 else if (n == 0) 1
 else n * fact(n-1)
```

```
def foo(n: Int) = fact(n + 10)
```

```
try {
 println (fact(3))
 println (foo(-100))
} catch {
 case e : factRangeException => {
 println("fact range error: " + e.arg)
 }
}
```



# Datatypes

# Types so far

Types have introduction operations and elimination ones.

- Introduction: how to construct elements of the type
- Elimination: how to use elements of the type

## ➤ Primitive types

- Int, Boolean, Double, String, ...
- Intro for Int: ..., -2, -1, 0, 1, 2,
- Elim for Int: +, -, \*, /, <, <=, ...

## ➤ Function types

- $\text{Int} \Rightarrow \text{Int}$ ,  $(\text{Int} \Rightarrow \text{Int}) \Rightarrow (\text{Int} \Rightarrow \text{Int})$ , ...
- Intro:  $(x:T) \Rightarrow e$
- Elim:  $f(v)$

# Tuples

## ➤ Tuples

Intro:

- $(1,2,3) : (\text{Int}, \text{Int}, \text{Int})$
- $(1, \text{"a"}) : (\text{Int}, \text{String})$

Elim:

- $(1, \text{"a"}, 10)._1 = 1$
- $(1, \text{"a"}, 10)._2 = \text{"a"}$
- $(1, \text{"a"}, 10)._3 = 10$

Only up to length 22

# Structural Types (a.k.a. Record Types): Examples

```
def bar (x: Int) = x+1
val foo = new //or, object foo
{
 val a = 1+2
 def b = a + 1
 def f(x: Int) = b + x
 def f(x: String) = "hello" + x
 val g : Int => Int = bar _
}

foo.b
foo.f(3)
foo.f("gil")
```

```
def g(x: {val a: Int; def b: Int;
 def f(x: Int): Int; def f(x: String): String}) =
 x.f(3)
g(foo)
```

# Structural Types: Evaluation

```
1: def bar (x: Int) = x+1
2: val foo = new //or, object foo
3: { val a = 2 + 1
4: def f(x: Int) = a + x
5: val g : Int => Int = bar _
6: }
7: val b = foo.f(1)
8: foo.g(b)
```

## ➤ Evaluation with Closures

```
E1[], 1 ~ E1[bar=(x)x+1], 2 ~ E1[...]:E2[], 3 ~
E1[...]:E2[a=3], 4 ~
E1[...]:E2[a=3, f=(x)a+x], 5 ~
E1[...]:E2[a=3, f=(x)a+x, g=((x)x+1, E1)], 6 ~
E1[bar=(x)x+1, foo=(E2)], 7 ~
E1[bar=(x)x+1, foo=(E2), b=4], 8 ~ 5
7: E1:E2:E3[x=1], a+x ~ 3+1 ~ 4
8: E1:E4[x=4], x+1 ~ 4+1 ~ 5
```

# Type Alias

```
val gn = 0
object foo {
 val a = 3
 def b = a + 1
 def f(x: Int) = b + x + gn
}
```

foo.f(3)

```
type Foo = {val a: Int; def b: Int; def f(x: Int): Int}
```

```
def g(x: Foo) = {
 val gn = 10
 x.f(3)
}
```

g(foo)

# Algebraic Datatypes

## ➤ Ideas

- $T = C \text{ of } T * \dots * T$   
|  $C \text{ of } T * \dots * T$   
|  $\dots$   
|  $C \text{ of } T * \dots * T$

- E.g.

Attr = Name of String  
| Age of Int  
| DOB of Int \* Int \* Int  
| Height of Double

Intro:

Name(“Chulsoo Kim”), Name(“Younghee Lee”), Age(16),  
DOB(2000,3,10), Height(171.5), ...

# Algebraic Datatypes: Recursion

## ➤ Recursive ADT

- E.g.

$\text{IList} = \text{INil}$

|  $\text{ICons of Int} * \text{IList}$

Intro:

$\text{INil}()$ ,

$\text{ICons}(3, \text{INil}())$ ,

$\text{ICons}(2, \text{ICons}(3, \text{INil}()))$ ,

$\text{ICons}(1, \text{ICons}(2, \text{ICons}(3, \text{INil}())))$ ,

...



# Algebraic Datatypes In Scala

## ➤ Attr

```
sealed abstract class Attr
case class Name(name: String) extends Attr
case class Age(age: Int) extends Attr
case class DOB(year: Int, month: Int, day: Int) extends Attr
case class Height(height: Double) extends Attr
```

```
val a : Attr = Name("Chulsoo Kim")
val b : Attr = DOB(2000, 3, 10)
```

## ➤ IList

```
sealed abstract class IList
case class INil() extends IList
case class ICons(hd: Int, tl: IList) extends IList
```

```
val x : IList = ICons(2, ICons(1, INil()))
def gen(n: Int) : IList =
 if (n <= 0) INil()
 else ICons(n, gen(n-1))
```

# Exercise

IOption = INone  
          | ISome of Int

BTree = Leaf  
        | Node of Int \* BTree \* BTree

```
sealed abstract class IList
case class INil() extends IList
case class ICons(hd: Int, tl: IList) extends IList

def x : IList = ICons(2, ICons(1, INil()))
```

# Solution

```
sealed abstract class IOption
case class INone() extends IOption
case class ISome(some: Int) extends IOption
```

```
sealed abstract class BTree
case class Leaf() extends BTree
case class Node(value: Int, left: BTree, right: BTree)
 extends BTree
```

# Pattern Matching

## ➤ Pattern Matching

- A way to use algebraic datatypes

```
e match {
 case C1(...) => e1
 ...
 case Cn(...) => en
}
```

# Pattern Matching: An Example

```
def length(xs: IList) : Int =
 xs match {
 case /Ni/() => 0
 case /Cons(x, tl) => 1 + length(tl)
 }
```

length(x)

# Advanced Pattern Matching

## ➤ Advanced Pattern Matching

```
e match {
 case P1 => e1

 ...

 case Pn => en
}
```

- One can combine constructors and use `_` and `|` in a pattern.  
(E.g) `case ICons(x, INil()) | ICons(x, ICons(_, INil())) => ...`
- The given value `e` is matched against the first pattern `P1`.  
If succeeds, evaluate `e1`.  
If fails, `e` is matched against `P2`.  
If succeeds, evaluate `e2`.  
If fails, ...
- The compiler checks exhaustiveness (ie, whether there is a missing case).

# Advanced Pattern Matching: An Example

```
def secondElmt(xs: IList) : IOption =
 xs match {
 case /Nil() | /Cons(_, /Nil()) => /None()
 case /Cons(_, /Cons(x, _)) => /Some(x)
 }
```

Vs.

```
def secondElmt2(xs: IList) : IOption =
 xs match {
 case /Nil() | /Cons(_, /Nil()) => /None()
 case /Cons(_, /Cons(x, /Nil())) => /Some(x)
 case _ => /None()
 }
```

Vs.

```
def secondElmt2(xs: IList) : IOption =
 xs match {
 case /Cons(_, /Cons(x, /Nil())) => /Some(x)
 case _ => /None() }
```

# Pattern Matching on Int

```
def factorial(n: Int) : Int =
 n match {
 case 0 => 1
 case _ => n * factorial(n-1)
 }
```

```
def fib(n: Int) : Int =
 n match {
 case 0 | 1 => 1
 case _ => fib(n-1) + fib(n-2)
 }
```



# Pattern Matching with If

```
def f(n: Int) : Int =
 n match {
 case 0 | 1 => 1
 case _ if (n <= 5) => 2
 case _ => 3
 }
```

```
def f(t: BTree) : Int =
 t match {
 case Leaf() => 0
 case Node(n,_,_) if (n <= 10) => 1
 case Node(_,_,_) => 2
 }
```

# Exercise

Write a function “find(t: BTree, x: Int) : Boolean” that checks whether x is in t.

```
sealed abstract class BTree
case class Leaf() extends BTree
case class Node(value: Int, left: BTree, right: BTree)
 extends BTree
```

# Solution

```
def find(t: BTree, i: Int) : Boolean =
 t match {
 case Leaf() => false
 case Node(n, lt, rt) =>
 if (i == n) true
 else find(lt, i) || find(rt, i)
 }
```

```
def t: BTree = Node(5, Node(4, Node(2, Leaf(), Leaf()), Leaf()),
 Node(7, Node(6, Leaf(), Leaf()), Leaf()))
```

```
find(t, 7)
```

```
find(t, 1)
```

# Type Checking & Inference (Concept)

# What Are Types For?

## ➤ Typed Programming

```
def id1(x: Int): Int = x
def id2(x: Double): Double = x
```

- At run time, type information is erased (ie, `id1 = id2`)

## ➤ Untyped Programming

```
def id(x) = x
```

- Do not care about types at compile time.
  - But, many such languages check types at run time paying cost.
  - Without run-time type check, errors can be badly propagated.
- ## ➤ What is compile-time type checking for?
- Can detect type errors at compile time.
  - Increase Readability (Give a good abstraction).
  - Soundness: Well-typed programs raise no type errors at run time.

# Type Checking and Inference

## ➤ Type Checking

$$x_1:T_1, x_2:T_2, \dots, x_n:T_n \vdash e : T$$

- `def f(x: Boolean): Boolean = x > 3`  
=> Type error
- `def f(x: Int): Boolean = x > 3`  
=> OK. `f: (x: Int)Boolean`

## ➤ Type Inference

$$x_1:T_1, x_2:T_2, \dots, x_n:T_n \vdash e : ?$$

- `def f(x: Int) = x > 3`  
=> OK by type inference. `f: (x: Int)Boolean`
- Too much type inference is not good. Why?

You can learn how type checking & inference work in  
**4190.310 Programming Languages**

# Parametric Polymorphism

# Parametric Polymorphism: Functions

## ➤ Problem

```
def id1(x: Int): Int = x
def id2(x: Double): Double = x
```

- Can we avoid DRY?
- Polymorphism to the rescue!

## ➤ Parametric Polymorphism (a.k.a. For-all Types)

```
def id[A](x: A) : A = x
```

- The type of `id` is `[A](val x:A)=>A`
- `id` is a parametric expression.
- `id[T] _` is a value of type `T=>T` for any type `T`.

[We will learn other kinds of polymorphism later.]



# Examples

```
def id[A](x:A) = x
id(3)
id("abc")
```

```
def applyn[A](f: A => A, n: Int, x: A): A =
 n match {
 case 0 => x
 case _ => f(applyn(f, n - 1, x))
 }
```

```
applyn((x:Int)=>x+1, 100, 3)
applyn((x:String)=>x+"!", 10, "gil")
applyn(id[String], 10, "hur")
```

```
def foo[A,B](f: A=>A, x: (A,B)) : (A,B) =
 (applyn[A](f, 10, x._1), x._2)
```

```
foo[String, Int]((x:String)=>x+"!", ("abc", 10))
```

# Full Polymorphism using Scala's trick

```
type Applyn = {def apply[A](f: A=>A, n: Int, x: A): A}
object applyn { // val applyn = new {
 def apply[A](f: A=>A, n: Int, x: A): A =
 n match {
 case 0 => x
 case _ => f(applyn(f, n-1, x))
 }
}
applyn((x:String)=>x+"!", 10, "gil")
// applyn.apply[String]((x:String)=>x+"!", 10, "gil")

def foo(f: Applyn): String = {
 val a:String = f[String]((x:String)=> x + "!", 10, "gil")
 val b:Int = f[Int]((x:Int)=> x + 2, 10, 5)
 a + b.toString()
}
foo(applyn)
```

# Summary

## ➤ Main points

- Definitions allow type parameters at the beginning.  
(E.g.) `def id[A](x: A) : A = x`
- Values does not allow type parameters.  
(E.g.) `[A](A=>A)` is not a valid value type.
- Record types allow type parameters in their “def” fields.  
(E.g.) `{def apply[A](x: A): A}` is a valid value type.

## ➤ Minor points

- The compiler automatically inserts “`.apply`” when a record is applied to arguments.
- The function types are encoded as “class types”, which we will learn in Part II.

# Parametric Polymorphism: Datatypes

```
sealed abstract class MyOption[A]
case class MyNone[A]() extends MyOption[A]
case class MySome[A](some: A) extends MyOption[A]
```

```
sealed abstract class MyList[A]
case class MyNil[A]() extends MyList[A]
case class MyCons[A](hd: A, tl: MyList[A]) extends MyList[A]
```

```
sealed abstract class BTree[A]
case class Leaf[A]() extends BTree[A]
case class Node[A](value: A, left: BTree[A], right: BTree[A])
extends BTree[A]
```

```
def x: MyList[Int] = MyCons(3, MyNil())
def y: MyList[String] = MyCons("abc", MyNil())
```

# Exercise

BSTree[A] = Leaf

| Node of Int \* A \* BSTree[A] \* BSTree[A]

```
def lookup[A](t: BSTree[A], k: Int) : MyOption[A] =
 ???
```

```
def t : BSTree[String] =
 Node(5, "My5",
 Node(4, "My4", Node(2, "My2", Leaf(), Leaf()), Leaf()),
 Node(7, "My7", Node(6, "My6", Leaf(), Leaf()), Leaf()))
```

lookup(t, 7)

lookup(t, 3)

# Solution

```
sealed abstract class BSTree[A]
case class Leaf[A]() extends BSTree[A]
case class Node[A](key: Int, value: A, left: BSTree[A], right:
BSTree[A]) extends BSTree[A]
def lookup[A](t: BSTree[A], key: Int) : MyOption[A] =
 t match {
 case Leaf() => MyNone()
 case Node(k,v,l t,r t) =>
 k match {
 case _ if key == k => MySome(v)
 case _ if key < k => lookup(l t,key)
 case _ => lookup(r t, key)
 }
 }
def t : BSTree[String] =
 Node(5, "My5",
 Node(4, "My4", Node(2, "My2", Leaf(), Leaf()), Leaf()),
 Node(7, "My7", Node(6, "My6", Leaf(), Leaf()), Leaf()))
lookup(t, 7)
lookup(t, 3)
```

# A Better Way

```
sealed abstract class BTree[A]
case class Leaf[A]() extends BTree[A]
case class Node[A](value: A, left: BTree[A], right: BTree[A])
 extends BTree[A]
```

```
type BSTree[A] = BTree[(Int, A)]
```

```
def lookup[A](t: BSTree[A], k: Int) : MyOption[A] =
 ???
```

```
def t : BSTree[String] =
 Node((5, "My5"),
 Node((4, "My4"), Node((2, "My2"), Leaf(), Leaf()), Leaf()),
 Node((7, "My7"), Node((6, "My6"), Leaf(), Leaf()), Leaf()))
```

```
lookup(t, 7)
```

# Solution

```
type BSTree[A] = BTree[(Int,A)]
```

```
def lookup[A](t: BTree[A], key: Int) : MyOption[A] =
 t match {
 case Leaf() => MyNone()
 case Node((k,v),lt,rt) =>
 k match {
 case _ if key == k => MySome(v)
 case _ if key < k => lookup(lt,key)
 case _ => lookup(rt, key)
 }
 }
```

```
def t : BTree[String] =
 Node((5, "My5"),
 Node((4, "My4"), Node((2, "My2"), Leaf(), Leaf()), Leaf()),
 Node((7, "My7"), Node((6, "My6"), Leaf(), Leaf()), Leaf()))
```

```
lookup(t, 7)
lookup(t, 3)
```



# Polymorphic Option (Library)

## ➤ Option[T]

Intro:

- None
- Some(x)
- Library functions

Elim:

- Pattern matching
- Library functions

Some(3) : Option[Int]

Some("abc"): Option[String]

None: Option[Int]

None: Option[String]

# Polymorphic List (Library)

## ➤ List[T]

Intro:

- Nil
- $x :: L$
- Library functions

Elim:

- Pattern matching
- Library functions

`“abc”::Nil : List[String]`

`List(1,3,4,2,5) = 1::3::4::2::5::Nil : List[Int]`

# **PART 2**

## **Object-Oriented Programming**

# **Sub Type Polymorphism (Concept)**

# Motivation

We want:

```
object tom {
 val name = "Tom"
 val home = "02-880-1234"
}
```

```
object bob {
 val name = "Bob"
 val mobile = "010-1111-2222"
}
```

```
def greeting(r: ???) = "Hi " + r.name + ", How are you?"
greeting(tom)
greeting(bob)
```

Note that we have

```
tom: {val name: String; val home: String}
bob: {val name: String; val mobile: String}
```

# Sub Types to the Rescue!

```
type NameHome = { val name: String; val home: String }
type NameMobile = { val name: String; val mobile: String }
type Name = { val name: String }
```

NameHome <: Name (NameHome is a sub type of Name)

NameMobile <: Name (NameMobile is a sub type of Name)

```
def greeting(r: Name) = "Hi " + r.name + ", How are you?"
greeting(tom)
greeting(bob)
```

# Sub Types

- The sub type relation is kind of the subset relation.
- But they are **NOT** the same.
- $T <: S$   
Every element of T **can be used as** that of S.
- *Cf.* T is a subset of S.  
Every element of T **is** that of S.
- Why polymorphism?  
A function of type  $S \Rightarrow R$  can be used as  $T \Rightarrow R$  for many sub types T of S.  
Note that  $S \Rightarrow R <: T \Rightarrow R$  when  $T <: S$ .

# Two Kinds of Sub Types

## ➤ Structural Sub Types

- The system implicitly determines the sub type relation by the structures of data types.
- Structurally equivalent types are the same.

## ➤ Nominal Sub Types

- The user explicitly specify the sub type relation using the names of data types.
- Structurally equivalent types with different names may be different.



# Structural Sub Types

# General Sub Type Rules

- Reflexivity:

For any type T, we have:

$$T <: T$$

- Transitivity:

For any types T, S, R, we have:

$$T <: R \quad R <: S$$

=====

$$T <: S$$

# Sub Types for Special Types

- Nothing: The empty set
- Any: The set of all values

- For any type T, we have:

$\text{Nothing} \leq T \leq \text{Any}$

- Example

```
val a : Int = 3
val b : Any = a
def f(a: Nothing) : Int = a
```

# Sub Types for Records

- Permutation

$$\{\dots; x: T1; y: T2; \dots\} <: \{\dots; y: T2, x: T1; \dots\}$$

- Width

$$\{\dots; x: T; \dots\} <: \{\dots; \dots\}$$

- Depth

$$T <: S$$

$$\{\dots; x: T; \dots\} <: \{\dots; x: S; \dots\}$$

# Sub Types for Records

- Example

{val x: { val y: Int; val z: String}, val w: Int}

<: (by permutation)

{val w: Int; val x: { val y: Int; val z: String} }

<: (by depth & width)

{val w: Int; val x: {val z: String} }

# Sub Types for Tuples

- Depth

$$T <: S$$

=====

$$(\dots; x: T; \dots) <: (\dots; x: S; \dots)$$

# Sub Types for Functions

- Function Sub Type

$$T <: T' \quad S <: S'$$

=====

$$(T' \Rightarrow S) <: (T \Rightarrow S')$$

- Example

```
def foo(s: {val a: Int; val b: Int}) :
 {val x: Int; val y: Int} = {
 object tmp {
 val x = s.b
 val y = s.a
 }
 tmp
 }
val gee:
 {val a: Int; val b: Int; val c: Int} =>
 {val x: Int} =
 foo _
```

# Classes



# Class: Parameterized Record

```
type gee_type = {val a: Int; def b: Int; def f(z: Int): Int}
def gee_fun(x: Int, y: Int) = new {
 val a : Int = x
 def b : Int = a + y
 def f(z: Int) : Int = b + y + z
}
val gee : gee_type = gee_fun(10, 20)
```

```
class foo_type(x: Int, y: Int) {
 val a : Int = x
 def b : Int = a + y
 def f(z: Int) : Int = b + y + z
}
val foo : foo_type = new foo_type(10, 20)
```

- use: foo.a foo.b foo.f
- foo is a value of foo\_type
- gee is a value of gee\_type

# Class: No Structural Sub Typing

## ➤ Records: Structural sub-typing

`foo_type <: gee_type`

## ➤ Classes: Nominal sub-typing

`gee_type <: foo_type`

`val v1 : gee_type = foo`

`val v2 : foo_type = gee // type error`

# Structural Types vs. Nominal Types

## ➤ Structural Types

- Includes arbitrary values with the required structures as elements
- Allows arbitrary types with the required structures as sub types
- Cannot assume any properties on their elements

## ➤ Nominal Types

- Includes only specific values as elements
- Allows only specific types as sub types
- Can assume specific properties on their elements

# Class: Can be Recursive!

```
class MyList[A](v: A, nxt: Option[MyList[A]]) {
 val value : A = v
 val next : Option[MyList[A]] = nxt
}
type YourList[A] = Option[MyList[A]]

val t : YourList[Int] =
 Some(new MyList(3, Some (new MyList(4, None))))
```

# Note on Null value

- `null`: The special element of every class & structural type
- This value is needed to construct disjoint union types using classes in Java, which, however, is not as elegant and type safe as algebraic data types (ADTs):
  - Such disjoint union types can contain junk values (not elegant).
  - Null-point exception can be raised at run time (not type safe).
- For this reason, it is discouraged to use `null` in Scala although Scala supports `null` for compatibility with Java.
- Instead, it is encouraged to use ADTs, which themselves are classes and thus take advantages of both ADT and class.

# Simplification using Argument Members

```
class MyList[A](v: A, nxt: Option[MyList[A]]) {
 val value = v
 val next = nxt
}
```

```
class MyList[A](val value:A, val next:Option[MyList[A]]) {
}
```

```
class MyList[A](val value:A, val next:Option[MyList[A]])
```

# Simplification using Companion Object

```
class MyList[A](v: A, nxt: Option[MyList[A]]) {
 val value = v
 val next = nxt
}
```

```
object MyList // val MyList = new
{ def apply[A](v: A, nxt: Option[MyList[A]]) =
 new MyList(v,nxt)
}
```

```
type YourList[A] = Option[MyList[A]]
```

```
val t0 = None
```

```
val t1 = Some(new MyList(3, Some (new MyList(4, None))))
```

```
val t2 = Some(MyList(3, Some (MyList(4, None))))
```

# Exercise

Define a class “MyTree[A]” for binary trees:

```
MyTree[A] =
 (value: A) *
 (left: Option[MyTree[A]]) *
 (right: Option[MyTree[A]])
```



# Solution

```
class MyTree[A](v: A,
 lt: Option[MyTree[A]],
 rt: Option[MyTree[A]]) {
 val value = v
 val left = lt
 val right = rt
}
```

```
type YourTree[A] = Option[MyTree[A]]
```

```
val t0 : YourTree[Int] = None
```

```
val t1 : YourTree[Int] = Some(new MyTree(3, None, None))
```

```
val t2 : YourTree[Int] =
 Some(new MyTree(3, Some (new MyTree(4, None, None)), None))
```

# Nominal Sub Typing for Classes

# Nominal Sub Typing, a.k.a. Inheritance

```
class foo_type(x: Int, y: Int) {
 val a : Int = x
 def b : Int = a + y
 def f(z: Int) : Int = b + y + z
}
```

```
class gee_type(x: Int) extends foo_type(x+1,x+2) {
 val c : Int = f(x) + b
}
```

```
gee_type <: foo_type
```

```
(new gee_type(30)).c
def test(f: foo_type) = f.a + f.b
test(new foo_type(10,20))
test(new gee_type(30))
```

# Overriding 1

```
class foo_type(x: Int, y: Int) {
 val a : Int = x
 def b : Int = a + y
 def f(z: Int) : Int = b + y + z
}
class gee_type(x: Int) extends foo_type(x+1,x+2) {
 override def f(z: Int) = b + z
 // or, override def f(z: Int) = super.f(z) * 2
 val c : Int = f(x) + b
}
(new gee_type(30)).c
```

**Q:** Can we override with a different type?

```
override def f(z: String): Int = 77 //No, arg: diff type
def f(z: String): Int = 77 //Yes, arg: diff type
override def f(z: Int): Nothing = ??? //Yes, ret: sub type
```

## Overriding 2

```
class foo_type(x: Int, y: Int) {
 val a : Int = x
 def b : Int = a + y
 def f(z: Int) : Int = b + y + z
}
```

```
class gee_type(x: Int) extends foo_type(x+1,x+2) {
 override def b = 10
}
```

```
(new gee_type(30)).f(0)
```

## Example: My List

```
class MyList[A]()
class MyNil[A]() extends MyList[A]
class MyCons[A](val hd: A, val tl: MyList[A])
 extends MyList[A]
val t: MyList[Int] =
 new MyCons(3, (new MyCons(4, new MyNil()))))
```

## Example: MyList

```
class MyList[A]
```

```
class MyNil[A]() extends MyList[A]
```

```
object MyNil { def apply[A]() = new MyNil[A]() }
```

```
class MyCons[A](val hd: A, val tl: MyList[A])
 extends MyList[A]
```

```
object MyCons {
 def apply[A](hd:A, tl:MyList[A]) = new MyCons[A](hd, tl)}
```

```
val t: MyList[Int] = MyCons(3, MyNil())
```

```
def length(x: MyList[Int]) = ???
```

# Case Class

```
class MyList[A]() { ... }
```

```
case class MyNil[A]() extends MyList[A] { ... }
```

```
object MyNil { def apply[A]() = new MyNil[A]() }
```

```
case class MyCons[A](val hd: A, val tl: MyList[A])
 extends MyList[A] { ... }
```

```
object MyCons {
```

```
 def apply[A](hd:A, tl:MyList[A]) = new MyCons[A](hd, tl)}
```

```
val t: MyList[Int] = MyCons(3, MyNil())
```

+ Pattern Matching

*Cf.* sealed abstract class MyList[A]



# Exercise

Define “MyTree[A]” using sub class.

```
class MyTree[A](v: A,
 lt: Option[MyTree[A]],
 rt: Option[MyTree[A]]) {
 val value = v
 val left = lt
 val right = rt
}
```

```
type YourTree[A] = Option[MyTree[A]]
```

# Solution

```
sealed abstract class MyTree[A]
case class Empty[A]() extends MyTree[A]
case class Node[A](value: A,
 left: MyTree[A],
 right: MyTree[A])
 extends MyTree[A]

val t : MyTree[Int] =
 Node(3, Node(4, Empty(), Empty()), Empty())

t match {
 case Empty() => 0
 case Node(v, l, r) => v
}
```

# **Abstract Classes for Specification**

# Abstract Class: Specification

## ➤ Abstract Classes

- Can be used to abstract away the implementation details.

Abstract classes for Specification  
Concrete sub-classes for Implementation

# Abstract Class: Specification

## ➤ Example Specification

```
abstract class Iter[A] {
 def getValue: Option[A]
 def getNext: Iter[A]
}
```

```
def sumElements[A](f: A=>Int)(xs: Iter[A]) : Int =
 xs.getValue match {
 case None => 0
 case Some(n) => f(n) + sumElements(f)(xs.getNext)
 }
```

```
def sumElementsId(xs: Iter[Int]) =
 sumElements((x: Int)=>x)(xs)
```

# Concrete Class: Implementation

```
sealed abstract class MyList[A] extends Iter[A]
case class MyNil[A]() extends MyList[A] {
 def getValue = None
 def getNext = throw new Exception("...")
}
```

```
case class MyCons[A](hd: A, tl: MyList[A])
 extends MyList[A]
{
 def getValue = Some(hd)
 def getNext = tl
}
```

```
val t1 = MyCons(3, MyCons(5, MyCons(7, MyNil())))
```

```
sumElementsId(t1)
```

## Exercise

Define `IntCounter(n)` that implements the specification `Iter[A]`.

```
class IntCounter(n: Int) extends Iter[Int] {
 def getValue = ???
 def getNext = ???
}
```

```
sumElementsId(new IntCounter(100))
```

# Solution

Define `IntCounter(n)` that implements the specification `Iter[A]`.

```
class IntCounter(n: Int) extends Iter[Int] {
 def getValue = if (n >= 0) Some(n) else None
 def getNext = new IntCounter(n-1)
}
```

```
sumElementsId(new IntCounter(100))
```



## **More on Abstract Classes**

## Problem: Iter for MyTree

```
abstract class Iter[A] {
 def getValue: Option[A]
 def getNext: Iter[A]
}
```

```
sealed abstract class MyTree[A]
case class Empty[A]() extends MyTree[A]
case class Node[A](value: A,
 left: MyTree[A],
 right: MyTree[A]) extends MyTree[A]
```

Q: Can MyTree[A] implement Iter[A]?

Try it, but it is not easy.

# Solution: Better Specification

```
abstract class Iter[A] {
 def getValue: Option[A]
 def getNext: Iter[A]
}
```

```
abstract class Iterable[A] {
 def iter : Iter[A]
}
```

```
def sumElementsGen[A](f: A=>Int)(xs: Iterable[A]) : Int =
 sumElements(f)(xs.iter)
```

# Let's Use MyList

```
sealed abstract class MyList[A] extends Iter[A]
case class MyNil[A]() extends MyList[A] {
 def getValue = None
 def getNext = throw new Exception("...")
}
case class MyCons[A](val hd: A, val tl: MyList[A])
 extends MyList[A] {
 def getValue = Some(hd)
 def getNext = tl
}
```

# MyTree <: Iterable (Try)

```
sealed abstract class MyTree[A] extends Iterable[A]
```

```
case class Empty[A]() extends MyTree[A] {
 val iter = MyNil()
}
```

```
case class Node[A](value: A,
 left: MyTree[A],
 right: MyTree[A]) extends MyTree[A] {
 // "val iter" is more specific than "def iter",
 // so it can be used in a sub type.
 // In this example, "val iter" is also
 // more efficient than "def iter".
 val iter = MyCons(value, ???(left.iter, right.iter))
}
```

# Extend MyList with append

```
sealed abstract class MyList[A] extends Iter[A] {
 def append(lst: MyList[A]) : MyList[A]
}

case class MyNil[A]() extends MyList[A] {
 def getValue = None
 def getNext = throw new Exception("...")
 def append(lst: MyList[A]) = lst
}

case class MyCons[A](val hd: A, val tl: MyList[A])
 extends MyList[A]
{
 def getValue = Some(hd)
 def getNext = tl
 def append(lst: MyList[A]) = MyCons(hd, tl.append(lst))
}
```

# MyTree <: Iterable

```
sealed abstract class MyTree[A] extends Iterable[A] {
 def iter : MyList[A]
 /* Note:
 def iter : Int // Type Error (no bug in Scala)
 // because not (Int <: Iter[A])
 */
}

case class Empty[A]() extends MyTree[A] {
 val iter = MyNil()
}

case class Node[A](value: A,
 left: MyTree[A],
 right: MyTree[A]) extends MyTree[A] {
 val iter = MyCons(value, left.iter.append(right.iter))
}
```

# Test

```
val t : MyTree[Int] =
 Node(3, Node(4, Node(2, Empty(), Empty()),
 Node(3, Empty(), Empty())),
 Node(5, Empty(), Empty()))

sumElementsGen((x: Int) => x)(t)
```



# Problem: Inefficiency & Stack Overflow

```
def sumN[A](f: A=>Int)(n: Int, xs: Iterable[A]) : Int = {
 def sumIter(res : Int, n: Int, xs: Iter[A]) : Int =
 if (n <= 0) res
 else xs.getValue match {
 case None => res
 case Some(v) => sumIter(f(v) + res, n-1, xs.getNext)
 }
 sumIter(0,n,xs.iter)
}
```

// Problem: Stack Overflow

```
{ val t: MyTree[Int] = generateTree(100)
 sumN((x: Int) => x)(100, t)
}
```

# Incremental Computation for Iteration: Ver 1

```
class MyTreeIter[A](val lst: MyList[MyTree[A]]) extends Iter[A] {
 val valAndNext : (Option[A], MyList[MyTree[A]]) = lst match {
 case MyNil() => (None, MyNil())
 case MyCons(hd, tl) => hd match {
 case Empty() => throw new Exception("...")
 case Node(v, Empty(), Empty()) => (Some(v), tl)
 case Node(v, lt, Empty()) => (Some(v), MyCons(lt, tl))
 case Node(v, Empty(), rt) => (Some(v), MyCons(rt, tl))
 case Node(v, lt, rt) => (Some(v), MyCons(lt, MyCons(rt, tl)))
 }
 }
 val getValue = valAndNext._1
 def getNext = new MyTreeIter(valAndNext._2)
}
```

# Incremental Computation for Iteration: Ver 1

```
sealed abstract class MyTree[A] extends Iterable[A]
case class Empty[A]() extends MyTree[A] {
 val iter = new MyTreeIter(MyNil())
}
case class Node[A](value: A,
 left: MyTree[A],
 right: MyTree[A]) extends MyTree[A]
{
 val iter = new MyTreeIter(MyCons(this, MyNil()))
}

{ val t: MyTree[Int] = generateTree(100)
 sumN((x: Int) => x)(100, t)
 // compare it with sumN((x: Int) => x)(100000000, t)
}
```

# Incremental Computation for Iteration: Ver 2

```
class MyTreeIter[A](v : Option[A], cont : => MyTreeIter[A])
 extends Iter[A]
{
 def append(ti: MyTreeIter[A]) : MyTreeIter[A] =
 getValue match {
 case None => ti
 case Some(v) => new MyTreeIter(getValue, getNext.append(ti))
 }
 val getValue = v
 def getNext : MyTreeIter[A] = cont
}
```

# Incremental Computation for Iteration: Ver 2

```
sealed abstract class MyTree[A] extends Iterable[A] {
 def iter : MyTreeIter[A]
}
case class Empty[A]() extends MyTree[A] {
 val iter = new MyTreeIter(None, throw new Exception("..."))
}
case class Node[A](value: A,
 left: MyTree[A],
 right: MyTree[A]) extends MyTree[A]
{
 val iter =
 new MyTreeIter(Some(value), left.iter.append(right.iter))
}

{ val t: MyTree[Int] = generateTree(100)
 sumN((x: Int) => x)(100, t)
 // compare it with sumN((x: Int) => x)(100000000, t)
}
```

# Iter <: Iterable

```
abstract class Iterable[A] {
 def iter : Iter[A]
}
```

```
abstract class Iter[A] extends Iterable[A] {
 def getValue: Option[A]
 def getNext: Iter[A]
 def iter = this
}
```

```
val lst : MyList[Int] =
 MyCons(3, MyCons(4, MyCons(2, MyNil())))
```

```
sumElementsGen ((x:Int)=>x)(lst)
```

# Wrapper for Inheritance

# Using a Wrapper Class

```
abstract class Iter[A] {
 def getValue: Option[A]
 def getNext: Iter[A]
}
```

```
class ListIter[A](val list: List[A]) extends Iter[A] {
 def getValue = list.headOption
 def getNext = new ListIter(list.tail)
}
```



# MyTree Using ListIter

```
abstract class Iterable[A] {
 def iter : Iter[A]
}

sealed abstract class MyTree[A] extends Iterable[A] {
 override def iter : ListIter[A]
}

case class Empty[A]() extends MyTree[A] {
 val iter : ListIter[A] = new ListIter(Ni)
}

case class Node[A](value: A,
 left: MyTree[A],
 right: MyTree[A])
 extends MyTree[A] {
 val iter : ListIter[A] = new ListIter(
 value::(left.iter.list ++ right.iter.list))
}
```

# Test

```
val t : MyTree[Int] =
 Node(3, Node(4, Node(2, Empty(), Empty()),
 Node(3, Empty(), Empty()))),
 Node(5, Empty(), Empty()))

sumElementsGen((x: Int) => x)(t)
```

# **Abstract Class With Abstract Types**

# Using an Abstract Type

```
abstract class Iterable[A] {
 type iter_t
 def iter: iter_t
 def getValue(i: iter_t) : Option[A]
 def getNext(i: iter_t) : iter_t
}

def sumElements[A](f:A=>Int)(xs: Iterable[A]) : Int = {
 def sumElementsIter(i: xs.iter_t) : Int =
 xs.getValue(i) match {
 case None => 0
 case Some(n) => f(n) + sumElementsIter(xs.getNext(i))
 }
 sumElementsIter(xs.iter)
}
```

# MyTree Using ListIter

```
sealed abstract class MyTree[A] extends Iterable[A] {
 type iter_t = List[A]
 def getValue(i: List[A]): Option[A] = i.headOption
 def getNext(i: List[A]): List[A] = i.tail
}

case class Empty[A]() extends MyTree[A] {
 val iter: List[A] = Nil
}

case class Node[A](value: A,
 left: MyTree[A], right: MyTree[A])
 extends MyTree[A] {
 val iter = value :: (left.iter ++ right.iter) //Pre-order
 //val iter = left.iter ++ (value :: right.iter) // In-order
 //val iter = left.iter ++ (right.iter ++ List(value))
 //Post-order
}
```

# Test

```
val t : MyTree[Int] =
 Node(3, Node(4, Node(2, Empty(), Empty()),
 Node(3, Empty(), Empty()))),
 Node(5, Empty(), Empty()))

sumElements((x: Int) => x)(t)
```

# Abstract Classes with Arguments

# Abstract Class with Arguments

```
abstract class IterableHE[A](eq: (A,A) => Boolean)
 extends Iterable[A]
{
 def hasElement(a: A) : Boolean = {
 def hasElementIter(i: iter_t) : Boolean =
 getValue(i) match {
 case None => false
 case Some(n) =>
 if (eq(a,n)) true
 else hasElementIter(getNext(i))
 }
 hasElementIter(iter)
 }
}
```



# MyTree

```
sealed abstract class MyTree[A](eq: (A, A) => Boolean)
 extends Iterable[A](eq) {
 type iter_t = List[A]
 def getValue(i: List[A]) : Option[A] = i.headOption
 def getNext(i: List[A]) : List[A] = i.tail
 }

case class Empty[A](eq: (A, A) => Boolean)
 extends MyTree[A](eq) {
 val iter : List[A] = Nil
 }

case class Node[A](eq: (A, A) => Boolean,
 value: A, left: MyTree[A], right: MyTree[A])
 extends MyTree[A](eq) {
 val iter : List[A] = value :: (left.iter ++ right.iter)
 }
```

# Test

```
val leq = (x: Int, y: Int) => x == y
```

```
val lEmpty = Empty(leq)
```

```
def lNode(n: Int, t1: MyTree[Int], t2: MyTree[Int]) =
 Node(leq, n, t1, t2)
```

```
val t : MyTree[Int] =
 lNode(3, lNode(4, lNode(2, lEmpty, lEmpty),
 lNode(3, lEmpty, lEmpty)),
 lNode(5, lEmpty, lEmpty))
```

```
sumElements((x: Int) => x)(t)
```

```
t.hasElement(5)
```

```
t.hasElement(10)
```

# Alternatively, Argument Elimination

```
abstract class IterableHE[A]
 extends Iterable[A]
{
 def eq(a:A, b:A) : Boolean
 def hasElement(a: A) : Boolean = {
 def hasElementIter(i: iter_t) : Boolean =
 getValue(i) match {
 case None => false
 case Some(n) =>
 if (eq(a,n)) true
 else hasElementIter(getNext(i))
 }
 hasElementIter(iter)
 }
}
```

# MyTree

```
sealed abstract class MyTree[A] extends Iterable[A] {
 type iter_t = List[A]
 def getValue(i : List[A]) : Option[A] = i.headOption
 def getNext(i: List[A]) : List[A] = i.tail
}

case class Empty[A](_eq: (A,A)=>Boolean) extends MyTree[A] {
 def eq(a:A, b:A) = _eq(a,b)
 val iter : List[A] = Nil
}

case class Node[A](_eq: (A,A)=>Boolean,
 value: A, left: MyTree[A], right: MyTree[A])
 extends MyTree[A] {
 def eq(a:A, b:A) = _eq(a,b)
 val iter : List[A] = value :: (left.iter ++ right.iter)
}
```

# Test

```
val leq = (x: Int, y: Int) => x == y
```

```
val lEmpty = Empty(leq)
```

```
def lNode(n: Int, t1: MyTree[Int], t2: MyTree[Int]) =
 Node(leq, n, t1, t2)
```

```
val t : MyTree[Int] =
 lNode(3, lNode(4, lNode(2, lEmpty, lEmpty),
 lNode(3, lEmpty, lEmpty)),
 lNode(5, lEmpty, lEmpty))
```

```
sumElements((x: Int) => x)(t)
```

```
t.hasElement(5)
```

```
t.hasElement(10)
```

## **More on Classes**

# Motivating Example

```
class Primes(val prime: Int, val primes: List[Int]) {
 def getNext: Primes = {
 val p = computeNextPrime(prime + 2)
 new Primes(p, primes ++ (p :: Nil))
 }
 def computeNextPrime(n: Int) : Int =
 if (primes.forall((p: Int) => n%p != 0)) n
 else computeNextPrime(n+2)
}
```

```
def nthPrime(n: Int): Int = {
 def go(primes: Primes, k: Int): Int =
 if (k <= 1) primes.prime
 else go(primes.getNext, k - 1)
 if (n == 0) 2 else go(new Primes(3, List(3)), n)
}
nthPrime(10000)
```

# Multiple Constructors

```
class Primes(val prime: Int, val primes: List[Int]) {
 def this() = this(3, List(3))
 def getNext: Primes = {
 val p = computeNextPrime(prime + 2)
 new Primes(p, primes ++ (p :: Nil))
 }
 def computeNextPrime(n: Int) : Int =
 if (primes.forall((p: Int) => n%p != 0)) n
 else computeNextPrime(n+2)
}
```

```
def nthPrime(n: Int): Int = {
 def go(primes: Primes, k: Int): Int =
 if (k <= 1) primes.prime
 else go(primes.getNext, k - 1)
 if (n == 0) 2 else go(new Primes, n)
}
nthPrime(10000)
```



# Access Modifiers

## ➤ Access Modifiers

- Private: Only the class can access the member.
- Protected: Only the class and its sub classes can access the member.

# Using Access Modifiers

```
class Primes private (val prime: Int, protected val primes: List[Int])
{ def this() = this(3, List(3))
 def getNext: Primes = {
 val p = computeNextPrime(prime + 2)
 new Primes(p, primes ++ (p :: Nil))
 }
 private def computeNextPrime(n: Int) : Int =
 if (primes.forall((p: Int) => n%p != 0)) n
 else computeNextPrime(n+2)
}
```

```
def nthPrime(n: Int): Int = {
 def go(primes: Primes, k: Int): Int =
 if (k <= 1) primes.prime
 else go(primes.getNext, k - 1)
 if (n == 0) 2 else go(new Primes, n)
}
nthPrime(10000)
```

# Traits for Multiple Inheritance

# Multiple Inheritance Problem

## ➤ Multiple Inheritance

- The famous “diamond problem”

```
class A(val a: Int)
class B extends A(10)
class C extends A(20)
class D extends B, C.
```

Problem 1: What is the value of (new D).a ?

Problem 2: The constructor of A must be executed once because A may contain side effects such as sending messages over the network.

# Java's Solution: Interface

## ➤ Interface

- An interface cannot contain any implementation but only types of its methods.
- A class can inherit implementations from only one parent class but implement multiple interfaces.

# Scala's Solution: Trait

## ➤ Traits

- A trait can implement any of its methods, but should have only one constructor with no arguments.
- A trait can
  - “extend” only one trait or (abstract) class with no arguments
  - “with” multiple traits.
- An (abstract) class can
  - “extend” only one trait or (abstract) class with any arguments
  - “with” multiple traits.

# Example

```
class A(val a : Int) {
 def this () = this(0)
}
trait B {
 def f(x: Int): Int = x
}
trait C extends A with B {
 def g(x: Int): Int = x + a
}
trait D extends B {
 def h(x: Int): Int = f(x + 50)
}
class E extends A(10) with C with D {
 override def f(x: Int) = x * a
}

val e = new E
```

# Algorithm for Multiple Inheritance

## ➤ Algorithm

- Give a linear order among all ancestors by “post-order” traversing without revisiting the same node.
- Invoke the constructors once in that order.

N.B. Post-order traversal of a class C means

- Recursively post-order traverse C’s first parent;
- ...
- Recursively post-order traverse C’s last parent; and
- Visit C.

By post-order traversing from “E” in the previous example, we have the order: A(10) → B → C → D → E

```
val e = new E
e.a // 10
e.f(100) // 100*10
e.g(100) // 100 + 10
e.h(100) // (100 + 50) * 10
```



# A Simple Example With Traits

# Motivation

```
abstract class Iter[A] {
 def getValue: Option[A]
 def getNext: Iter[A]
}
```

```
class ListIter[A](val list: List[A]) extends Iter[A] {
 def getValue = list.headOption
 def getNext = new ListIter(list.tail)
}
```

```
abstract class Dict[K,V] {
 def add(k: K, v: V): Dict[K,V]
 def find(k: K): Option[V]
}
```

Q: How can we extend ListIter and implement Dict?

# Specification using Traits

```
// abstract class Dict[K,V] {
// def add(k: K, v: V): Dict[K,V]
// def find(k: K): Option[V] }
```

```
trait Dict[K,V] {
 def add(k: K, v: V): Dict[K,V]
 def find(k: K): Option[V]
}
```

# Implementing Traits

```
class ListIterDict[K,V]
 (eq: (K,K)=>Boolean, list: List[(K,V)])
 extends ListIter[(K,V)](list)
 with Dict[K,V]
{
 def add(k:K,v:V) = new ListIterDict(eq,(k,v)::list)
 def find(k: K) : Option[V] = {
 def go(l: List[(K, V)]): Option[V] = l match {
 case Nil => None
 case (k1, v1) :: tl =>
 if (eq(k, k1)) Some(v1) else go(tl) }
 go(list) }
}
```

# Test

```
def sumElements[A](f: A=>Int)(xs: Iter[A]) : Int =
 xs.getValue match {
 case None => 0
 case Some(n) => f(n) + sumElements(f)(xs.getNext)
 }
def find3(d: Dict[Int,String]) = {
 d.find(3)
}

val d0 = new ListIterDict[Int,String]((x,y)=>x==y, Nil)
val d = d0.add(4, "four").add(3, "three")

sumElements[(Int,String)](x=>x._1)(d)
find3(d)
```

# **Mixin with Traits**

# Motivation: Mixin Functionality

```
abstract class Iter[A] {
 def getValue: Option[A]
 def getNext: Iter[A]
}
```

```
class ListIter[A](list: List[A]) extends Iter[A]
{
 def getValue = list.headOption
 def getNext: Iter[A] = new ListIter(list.tail)
}
```

```
trait MRIter[A] extends Iter[A] {
 def mapReduce[B,C](combine: (B,C)=>C, ival: C, f: A=>B): C = ???
}
```

# Mixin Composition

```
trait MRIter[A] extends Iter[A] {
 override def getNext: MRIter[A]
 def mapReduce[B,C](combine: (B,C)=>C, ival: C, f: A=>B): C =
 getValue match {
 case None => ival
 case Some(v) =>
 combine(f(v), getNext.mapReduce(combine, ival, f))
 }
}
```

```
class MRListIter[A](list: List[A])
 extends ListIter (list) with MRIter[A]
{
 override def getNext: MRIter[A] = new MRListIter(list.tail)
}
```

```
val mr = new MRListIter[Int](List(3,4,5))
mr.mapReduce[Int,Int]((b,c)=>b+c,0,(a)=>a*a)
```



# Mixin Composition: A Better Way

```
trait MRIter[A] extends Iter[A] {
 def mapReduce[B,C](combine: (B,C)=>C, ival: C, f: A=>B): C = {
 def go(c: Iter[A]): C = c.getValue match {
 case None => ival
 case Some(v) => combine(f(v), go(c.getNext))
 }
 go(this)
 }
}
```

```
class MRListIter[A](list: List[A])
 extends ListIter(list) with MRIter[A]
```

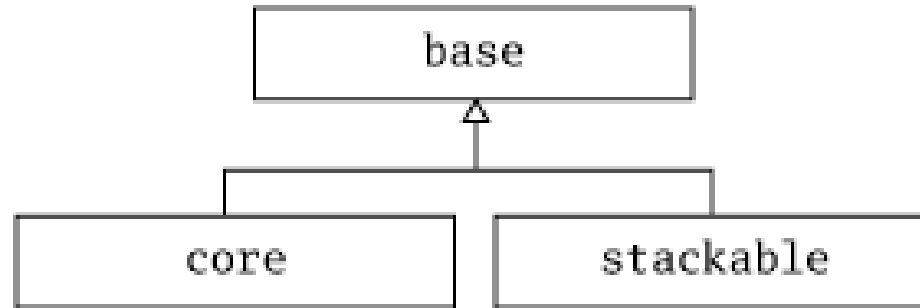
```
val mr = new MRListIter[Int](List(3,4,5))
```

```
// or, val mr = new ListIter(List(3,4,5)) with MRIter[Int]
```

```
mr.mapReduce[Int,Int]((b,c)=>b+c,0,(a)=>a*a)
```

# Stacking with Traits

# Typical Hierarchy in Scala



- **BASE**  
Interface (trait or abstract class)
- **CORE**  
Functionality (trait or concrete class)
- **CUSTOM**  
Modifications (each in a separate, composable trait)

# IntStack: Base

## ➤BASE

```
trait IntStack {
 def get(): (Int, IntStack)
 def put(x: Int): IntStack
}
```

# IntStack: Core

## ➤CORE

```
class BasicIntStack protected (xs: List[Int]) extends IntStack
{
 override val toString = "Stack:" + xs.toString
 def this() = this(Nil)
 protected def mkStack(xs: List[Int]): IntStack =
 new BasicIntStack(xs)
 def get(): (Int, IntStack) = (xs.head, mkStack(xs.tail))
 def put(x: Int): IntStack = mkStack(x :: xs)
}
```

```
val s0 = new BasicIntStack
val s1 = s0.put(3)
val s2 = s1.put(-2)
val s3 = s2.put(4)
val (v1,s4) = s3.get()
val (v2,s5) = s4.get()
```

# IntStack: Custom Modifications

## ➤CUSOM

```
trait Doubling extends IntStack {
 abstract override def put(x: Int): IntStack = super.put(2 * x)
}
```

```
trait Incrementing extends IntStack {
 abstract override def put(x: Int): IntStack = super.put(x + 1)
}
```

```
trait Filtering extends IntStack {
 abstract override def put(x: Int): IntStack =
 if (x >= 0) super.put(x) else this
}
```

# IntStack: Stacking

## ➤ Stacking

```
class DIFIntStack protected (xs: List[Int])
 extends BasicIntStack(xs)
 with Doubling with Incrementing with Filtering
{
 def this() = this(//)
 override def mkStack(xs: List[Int]): IntStack =
 new DIFIntStack(xs)
}
```

```
val s0 = new DIFIntStack
val s1 = s0.put(3)
val s2 = s1.put(-2)
val s3 = s2.put(4)
val (v1,s4) = s3.get()
val (v2,s5) = s4.get()
```

# Intersection Types



# Intersection Types

## ➤ Typing Rule

$$\frac{t : T1 \quad t : T2}{t : T1 \text{ with } T2}$$

## ➤ Example

```
trait A { val a: Int = 0 }
trait B { val b: Int = 0 }
class C extends A with B {
 override val a = 10
 override val b = 20
 val c = 30
}
```

```
val x = new C
val y: A with B = x
y.a // 10
y.b // 20
y.c // type error
```

# Subtype Relation for “with”

The subtype relation for “with” is structural.

- Permutation

=====

$$\dots \text{ with } T1 \text{ with } T2 \dots <: \dots \text{ with } T2 \text{ with } T1 \dots$$

- Width

=====

$$\dots \text{ with } T \dots <: \dots \dots$$

- Depth

=====

$$T <: S$$

=====

$$\dots \text{ with } T \dots <: \dots \text{ with } S \dots$$

# Additional Resources

## ➤ Traits

- <http://www.scala-lang.org/old/node/126>

## ➤ Mixin Composition

- <http://www.scala-lang.org/old/node/117>

## ➤ Stackable Trait Pattern

- [http://www.artima.com/scalazine/articles/stackable\\_trait\\_pattern.html](http://www.artima.com/scalazine/articles/stackable_trait_pattern.html)

## ➤ Multiple Inheritance via Traits

- <https://www.safaribooksonline.com/blog/2013/05/30/traits-how-scala-tames-multiple-inheritance/>

## ➤ UCSD CSE 130

- <http://cseweb.ucsd.edu/classes/wi14/cse130-a/lectures/scala/02-classes.html>

# **PART 3**

## **Type Classes for Specifications**

# Specifications over Parameter Types

# Subtype Polymorphism

```
trait Ord {
 // this cmp that < 0 iff this < that
 // this cmp that > 0 iff this > that
 // this cmp that == 0 iff this == that
 def cmp(that: Ord): Int

 def ==(that: Ord): Boolean = (this.cmp(that)) == 0
 def < (that: Ord): Boolean = (this.cmp(that) < 0
 def > (that: Ord): Boolean = (this.cmp(that) > 0
 def <= (that: Ord): Boolean = (this.cmp(that) <= 0
 def >= (that: Ord): Boolean = (this.cmp(that) >= 0
}

def max3(a: Ord, b: Ord, c: Ord) : Ord =
 if (a <= b) { if (b <= c) c else b }
 else { if (a <= c) c else a }
```

\* Problem: hard (almost impossible) to define `OrdInt <: Ord`

# Specification over Parameter Types

```
trait Ord[A] {
 def cmp(that: Ord[A]): Int
 def getValue : A

 def ==(that: Ord[A]): Boolean = (this.cmp(that)) == 0
 def < (that: Ord[A]): Boolean = (this.cmp that) < 0
 def > (that: Ord[A]): Boolean = (this.cmp that) > 0
 def <= (that: Ord[A]): Boolean = (this.cmp that) <= 0
 def >= (that: Ord[A]): Boolean = (this.cmp that) >= 0
}

def max3[A](a: Ord[A], b: Ord[A], c: Ord[A]) : Ord[A] =
 if (a <= b) {if (b <= c) c else b }
 else {if (a <= c) c else a }

class OInt(val getInt : Int) extends Ord[OInt] {
 def cmp(that: Ord[OInt]) = getInt.compare(that.getValue.getInt)
 def getValue = this
}

max3(new OInt(3), new OInt(2), new OInt(10)).getValue.getInt
```

## Further example: Ordered Bag

```
class Bag[A <: Ord[A]] protected (val toList: List[A]) {
 def this() = this(//)
 def add(x: A) : Bag[A] = {
 def go(elmts: List[A]): List[A] =
 elmts match {
 case // => x :: //
 case e :: _ if (x < e) => x :: elmts
 case e :: _ if (x == e) => elmts
 case e :: rest => e :: go(rest)
 }
 new Bag(go(toList))
 }
}

val emp = new Bag[0Int]()
val b = emp.add(new 0Int(3)).add(new 0Int(2)).add(new 0Int(10))
b.toList.map((x)=>x.getInt)
```

Works, but Very Awkward!!!



# Type Classes

# Completely Separating Ord from Int

```
abstract class Ord[A] {
 def cmp(me: A, you: A): Int

 def ==(me: A, you: A): Boolean = cmp(me, you) == 0
 def < (me: A, you: A): Boolean = cmp(me, you) < 0
 def > (me: A, you: A): Boolean = cmp(me, you) > 0
 def <= (me: A, you: A): Boolean = cmp(me, you) <= 0
 def >= (me: A, you: A): Boolean = cmp(me, you) >= 0
}

def max3[A](a: A, b: A, c: A)(implicit ord: Ord[A]) : A =
 if (ord.<=(a, b)) {if (ord.<=(b, c)) c else b }
 else {if (ord.<=(a, c)) c else a }

implicit val intOrd : Ord[Int] = new Ord[Int] {
 def cmp(me: Int, you: Int) = me - you }
```

max3(3, 2, 10) // 10

# Syntactic Sugar: new A with B with C { ... }

```
new A with B with C {
 code
}
```

is equivalent to

```
{
 class _tmp_ extends A with B with C {
 code
 }
 new _tmp_
}
```

# Implicit

## ➤ Implicit

- An argument is given “implicitly”

```
def foo(s: String)(implicit t: String) = s + t
```

```
implicit val exclamation : String = "!!!!!!"
```

```
foo("Hi")
```

```
foo("Hi")("???) // can give it explicitly
```

# Bag Example

```
class Bag[A] protected (val toList: List[A])(implicit ord: Ord[A])
{ def this()(implicit ord: Ord[A]) = this(Nil)(ord)
```

```
 def add(x: A) : Bag[A] = {
 def go(elmts: List[A]) : List[A] =
 elmts match {
 case Nil => x :: Nil
 case e :: _ if (ord.<(x,e)) => x :: elmts
 case e :: _ if (ord.==(x,e)) => elmts
 case e :: rest => e :: go(rest)
 }
 new Bag(go(toList))
 }
}
```

```
implicit val intOrd : Ord[Int] = new Ord[Int] {
 def cmp(me: Int, you: Int) = me - you }
```

```
(new Bag[Int]()).add(3).add(2).add(10).toList
```

# Bootstrapping Implicits

*// lexicographic order*

```
implicit def tup2Ord[A, B](implicit ordA: Ord[A], ordB: Ord[B]) = {
 new Ord[(A, B)] {
 def cmp(me: (A, B), you: (A, B)) : Int = {
 val c1 = ordA.cmp(me._1, you._1)
 if (c1 != 0) c1
 else { ordB.cmp(me._2, you._2) }
 }
 }
}
```

```
val b = new Bag[(Int, (Int, Int))]
b.add((3, (3, 4))).add((3, (2, 7))).add((4, (0, 0))).toList
```

# With Different Orders

```
val intOrdRev : Ord[Int] =
 new Ord[Int] { def cmp(me: Int, you: Int) = you - me }

(new Bag[Int])(intOrdRev).add(3).add(2).add(10).toList
```

# Type Classes With Multiple Parameters



# Iter

```
// trait Iter[A] {
// def getValue: Option[A]
// def getNext: Iter[A]
// }
```

```
abstract class Iter[I, A] {
 def getValue(i: I): Option[A]
 def getNext(i: I): I
}
```

```
def sumElements[I](xs: I)(implicit proxy: Iter[I, Int]) : Int =
 proxy.getValue(xs) match {
 case None => 0
 case Some(n) => n + sumElements(proxy.getNext(xs)) }
}
```

```
def printElements[I, A](xs: I)(implicit proxy: Iter[I, A]) : Unit =
 proxy.getValue(xs) match {
 case None =>
 case Some(n) => {print/n(n); printElements(proxy.getNext(xs))}}
```

# List

```
implicit def listIter[A] : Iter[List[A], A] =
 new Iter[List[A],A] {
 def getValue(a: List[A]) = a.headOption
 def getNext(a: List[A]) = a.tail
 }
```

```
val l = List(3,5,2,1)
```

```
sumElements(l) //sumElements(l)(listIter[Int])
```

```
printElements(l) //printElements(l)(listIter[Int])
```

# Iterable

```
// trait Iterable[A] {
// def iter : Iter[A]
// }
```

```
abstract class Iterable[R,I,A] {
 def iter(a: R): I
 def iterProxy: Iter[I, A]
}
```

```
def sumElements2[R,I](xs: R)(implicit proxy: Iterable[R,I,Int]) =
 sumElements(proxy.iter(xs))(proxy.iterProxy)
//sumElements[1](proxy.iter(xs))(proxy.iterProxy)
```

```
def printElements2[R,I,A](xs: R)(implicit proxy: Iterable[R,I,A]) =
 printElements(proxy.iter(xs))(proxy.iterProxy)
//printElements[1,A](proxy.iter(xs))(proxy.iterProxy)
```

# MyTree

```
sealed abstract class MyTree[A]
case class Empty[A]() extends MyTree[A]
case class Node[A](value: A, left: MyTree[A], right: MyTree[A])
 extends MyTree[A]

implicit def treeIterable[A](implicit proxy: Iter[List[A], A]) :
 Iterable[MyTree[A], List[A], A] =
 new Iterable[MyTree[A], List[A], A] {
 def iter(a: MyTree[A]): List[A] = a match {
 case Empty() => Nil
 case Node(v, left, right) => v :: (iter(left) ++ iter(right))
 }
 val iterProxy = proxy }

val t : MyTree[Int] =
 Node(3, Node(4, Empty(), Empty()), Node(2, Empty(), Empty()))
sumElements2(t) //sumElements2(t)(treeIterable[Int])
printElements2(t) //printElements2(t)(treeIterable[Int])
```

# Iter being Iterable

```
implicit def iterIterable[I,A](implicit proxy: Iter[I,A]) :
 Iterable[I,A] =
 new Iterable[I,A] {
 def iter(a: I) = a
 val iterProxy = proxy
 }
```

```
// val l = List(3,5,2,1)
```

```
sumElements2(l) //sumElements2(iterIterable(listIter[Int]))
```

```
printElements2(l) //printElements2(iterIterable(listIter[Int]))
```

# Higher-kind Type Classes

# Iter

```
import scala.language.higherKinds
//trait Iter[I,A] {
// def getValue(a: I): Option[A]
// def getNext(a: I): I }
abstract class Iter[I[_]] {
 def getValue[A](a: I[A]) : Option[A]
 def getNext[A](a: I[A]) : I[A]
}
def sumElements[I[_]](xs: I[Int])(implicit itr: Iter[I]): Int = {
 itr.getValue(xs) match {
 case None => 0
 case Some(n) => n + sumElements(itr.getNext(xs)) }
}
def printElements[I[_],A](xs: I[A])(implicit itr: Iter[I]): Unit = {
 itr.getValue(xs) match {
 case None =>
 case Some(n) => {print/n(n); printElements(itr.getNext(xs))}}
}
```

# List

```
implicit val listIter : Iter[List] =
 new Iter[List] {
 def getValue[A](a: List[A]) = a.headOption
 def getNext[A](a: List[A]) = a.tail
 }
```

```
val l = List(3,5,2,1)
```

```
sumElements(l) //sumElements(l)(listIter)
```

```
printElements(l) //printElements(l)(listIter)
```



# Iterable

```
//trait Iterable[R, I, A] {
 // def iter(a: R): I
 // def iterProxy: Iter[I, A]
 //}
```

```
abstract class Iterable[R[_], I[_]] {
 def iter[A](a: R[A]): I[A]
 def iterProxy: Iter[I]
}
```

```
def sumElements2[R[_], I[_]](xs: R[Int])
 (implicit proxy: Iterable[R, I]) =
 sumElements(proxy.iter(xs))(proxy.iterProxy)
//sumElements[I](proxy.iter(xs))(proxy.iterProxy)
```

```
def printElements2[R[_], I[_], A](xs: R[A])
 (implicit proxy: Iterable[R, I]) =
 printElements(proxy.iter(xs))(proxy.iterProxy)
//printElements[I, A](proxy.iter(xs))(proxy.iterProxy)
```

# MyTree

```
sealed abstract class MyTree[A]
case class Empty[A]() extends MyTree[A]
case class Node[A](value: A, left: MyTree[A], right: MyTree[A])
 extends MyTree[A]

implicit val treeIterable : Iterable[MyTree, List] =
 new Iterable[MyTree, List] {
 def iter[A](a: MyTree[A]) : List[A] = a match {
 case Empty() => Nil
 case Node(v, left, right) => v :: (iter(left) ++ iter(right)) }
 val iterProxy = implicitly[Iter[List]]
 }

val t : MyTree[Int] =
 Node(3, Node(4, Empty(), Empty()), Node(2, Empty(), Empty()))
sumElements2(t) //sumElements2(t)(treeIterable)
printElements2(t) //printElements2(t)(treeIterable)
```

# Implicitly

## ➤ Definition

```
def implicitly[A](implicit proxy: A) : A = proxy
```

## ➤ Example

```
implicit val treeIterable :
```

```
Iterable[MyTree, List] = new Iterable[MyTree, List] {
 def iter[A](a: MyTree[A]) : List[A] = a match {
 case Empty() => Nil
 case Node(v, left, right) => v :: (iter(left) ++ iter(right))
 }
 val iterProxy = implicitly[Iter[List]]
}
```

```
implicit def treeIterable(implicit proxy: Iter[List]) :
```

```
Iterable[MyTree, List] = new Iterable[MyTree, List] {
 def iter[A](a: MyTree[A]) : List[A] = a match {
 case Empty() => Nil
 case Node(v, left, right) => v :: (iter(left) ++ iter(right))
 }
 val iterProxy = proxy
}
```

# Iter being Iterable

```
implicit def iterIterable[I[_]](implicit proxy: Iter[I]) :
 Iterable[I, I] = new Iterable[I, I]
{
 def iter[A](a: I[A]) = a
 def iterProxy = proxy
}
```

```
// val l = List(3, 5, 2, 1)
sumElements2(l) //sumElements2(l)(iterIterable(listIter))
printElements2(l) //printElements2(l)(iterIterable(listIter))
```

# Example: Functor Specification

```
trait Functor[F[_]] {
 def map[A,B](f: A=>B)(x: F[A]) : F[B]
}
```

```
def compose[F[_],A,B,C](g: B=>C)(f: A=>B)(a: F[A])
 (implicit proxy: Functor[F]): F[C] =
 proxy.map(g)(proxy.map(f)(a))
```

# Example: Functor Implementation

```
sealed abstract class MyTree[A]
case class Empty[A]() extends MyTree[A]
case class Node[A](value: A, left: MyTree[A], right: MyTree[A])
 extends MyTree[A]

implicit val ListFunctor : Functor[List] = new Functor[List] {
 def map[A,B](f: A=>B)(x: List[A]) = x.map(f)
}

implicit val MyTreeFunctor : Functor[MyTree] = new Functor[MyTree] {
 def map[A,B](f: A=>B)(x: MyTree[A]) : MyTree[B] = x match {
 case Empty() => Empty()
 case Node(v,l,r) => Node(f(v),map(f)(l),map(f)(r)) }
}

compose((x: Int)=>x*x)((x: Int)=>x+x)(List(1,2,3))
val t : MyTree[Int] =
 Node(3, Node(4, Empty(), Empty()), Node(2, Empty(), Empty()))
compose((x: Int)=>x*x)((x: Int)=>x+x)(t)
```

# Even Higher Kinds

```
// Iter: (* -> *) -> *
```

```
abstract class Iter[I[_]] {
 def getValue[A](a: I[A]) : Option[A]
 def getNext[A](a: I[A]) : I[A]
}
```

```
// Foo: ((* -> *) -> *) -> *
```

```
abstract class Foo[I[_[_]]] {
 def get : I[List]
}
```

```
def f(x: Foo[Iter]) : Iter[List] = x.get
```

# Turning Type Classes into OO Classes



# Iter

```
// trait Iter[A] {
// def getValue: Option[A]
// def getNext: Iter[A]
// }
```

```
abstract class Iter[I, A] {
 def getValue(i: I): Option[A]
 def getNext(i: I): I
}
```

```
def sumElements[I](xs: I)(implicit proxy: Iter[I, Int]) : Int =
 proxy.getValue(xs) match {
 case None => 0
 case Some(n) => n + sumElements(proxy.getNext(xs)) }
}
```

```
def printElements[I, A](xs: I)(implicit proxy: Iter[I, A]) : Unit =
 proxy.getValue(xs) match {
 case None =>
 case Some(n) => {print/n(n); printElements(proxy.getNext(xs))}}
```

# List

```
implicit def listIter[A] : Iter[List[A], A] =
 new Iter[List[A],A] {
 def getValue(a: List[A]) = a.headOption
 def getNext(a: List[A]) = a.tail
 }
```

```
val l = List(3,5,2,1)
```

```
sumElements(l) //sumElements(l)(listIter[Int])
```

```
printElements(l) //printElements(l)(listIter[Int])
```

# How to return data satisfying a specification?

```
def inclter(max: Int) : Iter[Int,Int] = new Iter[Int,Int] {
 def getValue(i: Int) = if (i <= max) Some(i) else None
 def getNext(i: Int) = i + 1
}
```

```
def getMyIter(isInc: Boolean) : ? = {
 if (isInc) ? // want to return 0 with inclter(10)
 else ? // want to return List(3,1,4) with listIter
}
```

```
val i1 = getMyIter(true)
printElements(i1)
val i2 = getMyIter(false)
printElements(i2)
```

What can we do?

# Turning Type Classes into OO Classes

```
import scala.language.higherKinds
import scala.language.implicitConversions

abstract class Box2[S[_],A] {
 type Data
 val d: Data
 val i: S[Data,A]
}

object Box2 {
 implicit def apply[S[_],D,A](dd: D)(implicit ii: S[D,A]):
 Box2[S,A] = new Box2[S,A] {
 type Data = D
 val d = dd
 val i = ii
 }
 implicit def methods[S[_],A](d: Box2[S,A]): S[d.Data,A] = d.i
}
```

# Test

```
def inclter(max: Int) : Iter[Int,Int] = new Iter[Int,Int] {
 def getValue(i: Int) = if (i <= max) Some(i) else None
 def getNext(i: Int) = i + 1
}
```

```
def getMyIter(isInc: Boolean) : Box2[Iter,Int] = {
 if (isInc) Box2(0)(inclter(10)) // Box2.apply(0)(inclter(10))
 else List(3,1,4) // Box2(List(3,1,4))(listIter[Int])
}
```

```
val i1 = getMyIter(true)
printElements(i1.d)(i1.i)
i1.getValue(i1.getNext(i1.d)) // i1.i.getValue(i1.i.getNext(i1.d))
```

```
val i2 = getMyIter(false)
printElements(i2.d)(i2.i)
i2.getValue(i2.getNext(i2.d))
```

# Can define “Box” for different kinds of specs

```
abstract class Box[S[_]] {
 type Data
 val d: Data
 val i: S[Data]
}
```

```
object Box {
 implicit def apply[D,S[_]](dd: D)(implicit ii: S[D]):
 Box[S] = new Box[S] {
 type Data = D
 val d = dd
 val i = ii
 }
 implicit def methods[S[_]](d: Box[S]): S[d.Data] = d.i
}
```

# Iterable

```
// abstract class Iterable[R,I,A] {
// def iter(a: R): I
// def iterProxy: Iter[I, A]
// }
```

```
abstract class Iterable[R,A] {
 def iter(a: R): Box2[Iter,A]
}
```

```
def sumElements2[R](xs: R)(implicit proxy: Iterable[R,Int]) = {
 val cs = proxy.iter(xs)
 sumElements(cs.d)(cs.i)
}
```

```
def printElements2[R,A](xs: R)(implicit proxy: Iterable[R,A]) = {
 val cs = proxy.iter(xs)
 printElements(cs.d)(cs.i)
}
```

# MyTree

```
sealed abstract class MyTree[A]
case class Empty[A]() extends MyTree[A]
case class Node[A](value: A, left: MyTree[A], right: MyTree[A])
 extends MyTree[A]

implicit def treeIterable[A] :
 Iterable[MyTree[A], A] = new Iterable[MyTree[A], A] {
 def iter(a: MyTree[A]) = {
 def go(l: MyTree[A]) : List[A] =
 l match {
 case Empty() => Nil
 case Node(v, left, right) => v :: (go(left) ++ go(right))
 }
 go(a) // Box2(go(a))(listIter[A])
 }
}

val t : MyTree[Int] =
 Node(3, Node(4, Empty(), Empty()), Node(2, Empty(), Empty()))
sumElements2(t) //sumElements2(t)(treeIterable[Int])
printElements2(t) //printElements2(t)(treeIterable[Int])
```



# Iter being Iterable

```
implicit def iterIterable[I,A](implicit proxy: Iter[I,A]) :
 Iterable[I,A] = new Iterable[I,A] {
 def iter(a: I) = a // Box2(a)(proxy)
 }
```

```
// val l = List(3,5,2,1)
```

```
sumElements2(l)
```

```
printElements2(l)
```

# Stacking with Type Classes

# IntStack Spec & Modifying Traits

```
trait IntStack[A] {
 def empty : A
 def get(s: A): (Int, A)
 def put(s: A)(x: Int): A
}
```

```
trait Doubling[A] extends IntStack[A] {
 abstract override def put(s: A)(x: Int): A = super.put(s)(2 * x)
}
```

```
trait Incrementing[A] extends IntStack[A] {
 abstract override def put(s: A)(x: Int): A = super.put(s)(x + 1)
}
```

```
trait Filtering[A] extends IntStack[A] {
 abstract override def put(s: A)(x: Int): A =
 if (x >= 0) super.put(s)(x) else s
}
```

# Implementation using List

```
trait ListStackImpl extends IntStack[List[Int]] {
 val empty = List()
 def get(s: List[Int]) = (s.head, s.tail)
 def put(s: List[Int])(x: Int) = x :: s
}
```

```
val stkDIF : IntStack[List[Int]] =
 new ListStackImpl
 with Doubling[List[Int]]
 with Incrementing[List[Int]]
 with Filtering[List[Int]]
```

```
val s0 = stkDIF.empty
val s1 = stkDIF.put(s0)(3)
val s2 = stkDIF.put(s1)(-2)
val s3 = stkDIF.put(s2)(4)
val (v1,s4) = stkDIF.get(s3)
val (v2,s5) = stkDIF.get(s4)
```

# Implementation using SortedIntStack

```
class SortedStack protected (xs: List[Int]) {
 override val toString = "Stack:" + xs.toString
 def this() = this(Nil)
 def get : (Int, SortedStack) = (xs.head, new SortedStack(xs.tail))
 def put(x: Int) : SortedStack = {
 def go(l: List[Int]) : List[Int] = l match {
 case Nil => x :: Nil
 case hd :: tl => if (x <= hd) x :: l else hd :: go(tl)
 }
 new SortedStack(go(xs))
 }
}
```

```
trait SortedStackImpl extends IntStack[SortedStack] {
 val empty = new SortedStack()
 def get(s: SortedStack) : (Int, SortedStack) = s.get
 def put(s: SortedStack)(x: Int) : SortedStack = s.put(x)
}
```

# Implementation using SortedIntStack: Test

```
val sortedDIF : IntStack[SortedStack] =
 new SortedStackImpl
 with Doubling[SortedStack]
 with Incrementing[SortedStack]
 with Filtering[SortedStack]
```

```
val s0 = sortedDIF.empty
val s1 = sortedDIF.put(s0)(3)
val s2 = sortedDIF.put(s1)(-2)
val s3 = sortedDIF.put(s2)(4)
val (v1,s4) = sortedDIF.get(s3)
val (v2,s5) = sortedDIF.get(s4)
```

# Separating methods from SortedIntStack

```
class SortedStack protected (private val xs: List[Int]) {
 override val toString = "Stack:" + xs.toString
 def this() = this(Nil)
}

object SortedStack {
 trait Impl extends IntStack[SortedStack] {
 val empty : SortedStack = new SortedStack
 def get(s: SortedStack) : (Int, SortedStack) =
 (s.xs.head, new SortedStack(s.xs.tail))
 def put(s: SortedStack)(x: Int) : SortedStack = {
 def go(l: List[Int]): List[Int] = l match {
 case Nil => x :: Nil
 case hd :: tl => if (x <= hd) x :: l else hd :: go(tl)
 }
 new SortedStack(go(s.xs))
 }
 }
}
```

# Separating methods from SortedIntStack: Test

```
val sortedDIF : IntStack[SortedStack] =
 new SortedStack.Impl
 with Doubling[SortedStack]
 with Incrementing[SortedStack]
 with Filtering[SortedStack]
```

```
val s0 = sortedDIF.empty
val s1 = sortedDIF.put(s0)(3)
val s2 = sortedDIF.put(s1)(-2)
val s3 = sortedDIF.put(s2)(4)
val (v1,s4) = sortedDIF.get(s3)
val (v2,s5) = sortedDIF.get(s4)
```



# **PART 4**

## **Imperative Programming with Memory Updates**

# Mutable Variables

## ➤ Mutable Variables

- Use “var” instead of “val” and “def”
- We can update the value stored in a variable.

```
class Main(i: Int) {
 var a = i
}
```

```
val m = new Main(10)
m.a // 10
m.a = 20
m.a // 20
m.a += 5 // m.a = m.a + 5
m.a // 25
```

# While loop

## ➤ While loop

- Syntax: `while (cond) body`  
Executes *body* while *cond* holds.
- It is equivalent to:

```
def mywhile(cond: =>Boolean)(body: =>Unit) : Unit =
 if (cond) { body; mywhile(cond)(body) } else ()
```

## ➤ Example

```
var i = 0
var sum = 0
while (i <= 100) { // mywhile (i <= 100) {
 sum += i
 i += 2
}
sum // 2550
```

# For loop

## ➤ For loop

- Syntax: `for (i <- collection) body`  
Executes *body* for each *i* in *collection*.
- It is equivalent to:

```
def myfor[A](xs: Traversable[A])(f: A => Unit) : Unit =
 xs.foreach(f)
```

## ➤ Example

```
var sum = 0
for (i <- 0 to 100 by 2) { // myfor (0 to 100 by 2) { i =>
 sum += i
}
sum // 2550
```

# Additional Resources

## ➤ UCSD CSE 130

- <http://cseweb.ucsd.edu/classes/wi14/cse130-a/lectures/scala/00-crash.html>
- <http://cseweb.ucsd.edu/classes/wi14/cse130-a/lectures/scala/01-iterators.html>

Thanks for your hard work!