

An Overview of Céu

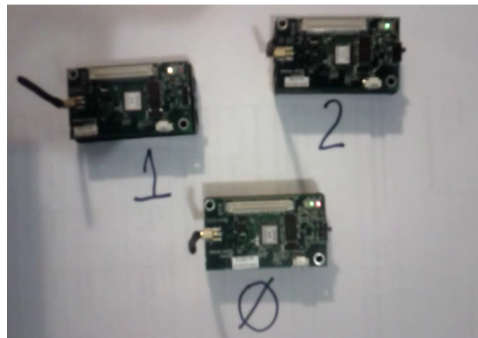
Francisco Sant'Anna

`francisco@ime.uerj.br`



Céu goals

- Front-end applications
 - GUIs, Games, interactive apps
 - Desktops → Arduino (*4K RAM*)
- Ruled by the environment
 - Immediate/real-time feedback
 - Global consensus
- Logical reasoning
 - Concurrency w/ Determinism



Céu non-goals

- Back-end infrastructure
 - High-performance servers
 - Clusters, Cloud computing
- Independent sessions/actors
 - Latency in communication
 - Distribution
- The C10M problem
 - Concurrency w/ Parallelism



“Hello world!” in Céu

- Blinking a LED
 1. *on* ↔ *off* every 500ms
 2. *stop* after “press”
 3. *restart* after 2s

- Compositions

- seq, loop, par (*trails*)
 - At any level of depth
- ~~state variables / communication~~

```
loop do
  par/or do
    loop do
      await 500ms;
      _leds_toggle();
    end
    with
      await PRESS;
    end
  end
  await 2s;
end
```

Lines of execution
=
Trails (in Céu)

From “Structured Programming” To “Structured Reactive Programming”



- Control Structures
 - Sequences, Loops, Conditionals
- Blocks, Scopes, Locals
 - Lexical memory management
- Subroutines
 - Abstraction mechanism
- What about reactivity?
 - Environment event → Short-lived callback
 - No more loops, scopes, etc.
 - Breaks structured programming
 - “Callbacks as our Generations' goto”
[Miguel de Icaza]
- The `await` statement
 - *Imperative-reactive* nature
- Compositions
 - Control structures + parallels
- Synchronous execution model
 - Time ~ Sequence of events
- Esterel did this back in the '80s
- What about abstractions and modularization?

Two Blinking LEDs

```
loop do
  par/
  1
  e
  with
  a
end
await 2s;
end
```

```
par do
```

```
  loop do
```

```
    par/or do
```

```
      every 500ms do
        _leds_toggle(1);
      end
```

```
    PRESS
    bt==1;
```

The need for abstractions!

```
  par/or do
```

```
    every 500ms do
      _leds_toggle(2);
    end
```

```
  with
```

```
    var int bt = await PRESS
    until bt==2;
```

```
  end
```

```
  await 2s;
```

```
end
```

```
end
```

Céu Abstractions

- code/await: a procedure that can await

```
code/await Led (var int led, var int but) -> FOREVER
do
```

```
  loop do
    par/or do
      every 500ms do
        _leds_toggle(led);
      end
    with
      var int bt = await PRESS
        until bt==but;
    end
    await 2s;
  end
```

```
end
```

A code/await
reacts directly to
the environment

Two Blinking LEDs

```
code/await Led (var int led, var int but) -> FOREVER  
do  
    <...>  
end
```

par do

await Led(1,1);

with

await Led(2,2);

end

LEDs are reacting
in parallel

await invokes a
code/await and waits for
its termination

6-Blinking LEDs

```
code/await Led (var int led, var int but) -> FOREVER
do
    <...>
end

par do
    await Led(1,1);
with
    await Led(2,2);
with
    await Led(3,3);
with
    await Led(4,4);
with
    await Led(5,5);
with
    await Led(6,6);
end
```


Lexical Scope

```
code/await Led (var int led, var int but) -> FOREVER  
do  
    <...>  
end
```

```
loop do  
    par/or do  
        var int bt = await PRESS until bt==0;  
        with  
            ← await Led(1,1);  
            with  
                ← await Led(2,2);  
                with  
                    ← await Led(3,3);  
                    with  
                        ← await Led(4,4);  
                        with  
                            ← await Led(5,5);  
                            with  
                                ← await Led(6,6);  
                            end  
                        end  
                    end  
                end  
            end  
        end  
    end  
end
```

code/await has
lexical scope

code/await out of scope:
data reclaimed **and body aborted**

code/await management is
as simple as local variables

**Structured, Reactive,
Abstract.**

***But we still have a static
semantics!***

6-Blinking LEDs with Lexical Scope

```
loop do
  par/or do
    loop do
      await 500ms;
      _leds_toggle(i);
    end
  with
    await i-PRESS;
  end
end
```

The diagram illustrates the transformation of a parallel loop into a nested structure with lexical scope. On the left, a code block shows a `loop do` containing a `par/or do` block. Inside the `par/or do` block, there is a `loop do` with an `await 500ms;` and a `_leds_toggle(i);` statement, followed by an `end` statement. To the right of this `end` statement, there is a `with` block containing `await i-PRESS;` and an `end` statement. Three arrows originate from the `end` statement of the inner `loop do` and point to the corresponding `end` statements in the transformed code on the right: the `end` of the innermost `loop do`, the `end` of the `with` block, and the `end` of the outermost `loop do`.

```
loop do
  par/or do
    par do
      <body-1>
    with
      ...
    with
      <body-n>
    end
  end
  with
    var int bt =
      await PRESS
    until bt==0;
  end
end
```

Another Example

```
code/await Bird (var int y, var int speed) -> FOREVER
do
  <...>
end
```

Reaction to the
environment is
abstracted inside
the body

```
par do
  await Bird(100,100);
with
  await Bird(300,200);
end
```

On instantiation,
only the interface
matters

```
code/await Bird (var int y, var int speed) -> FOREVER  
do  
    <...>  
end
```

```
loop do  
    par/or do  
        await Bird(100,100);  
    with  
        await Bird(100,100);  
    with  
        await MOUSE_CLICK;  
    end  
end
```

```
code/await Bird (var int y, var int speed) -> FOREVER
do
    <...>
end
```

A pool is a container for
code/await instances

```
pool[5] Bird birds;
```

```
loop i in [1 -> 5] do
```

```
    spawn Bird(70*i, 100+10*i) in birds;
```

```
end
```

```
await FOREVER;
```

spawn invokes a
code/await and continues

```
code/await Bird (var int y, var int speed) -> FOREVER
do
    <...>
end

pool[5] Bird birds;
    ↑
every 1s do
    spawn Bird(20+_rand()%HEIGHT, 100+_rand()%100) in birds;
end
```

```
code/await Bird (var int y, var int speed) -> FOREVER
do
    <...>
end

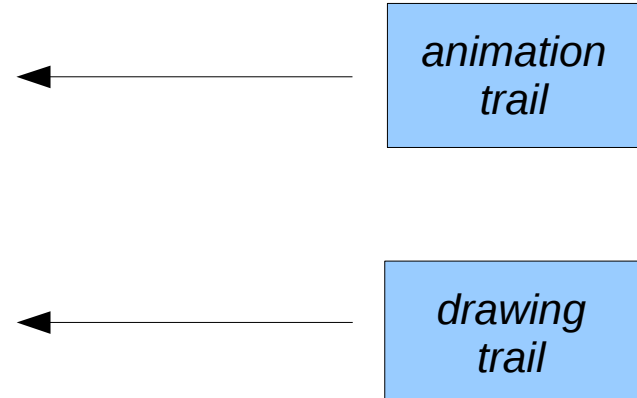
pool[] Bird birds;
every 1s do
    spawn Bird(...) in birds;
end
```



```
code/await Bird (var int y, var int speed) -> FOREVER
do
  par do
    every UPDATE do
      <animate-bird>
    end
  with
    every DRAW do
      <draw-bird>
    end
  end
end

pool[] Bird birds;

every 1s do
  spawn Bird(...) in birds;
end
```



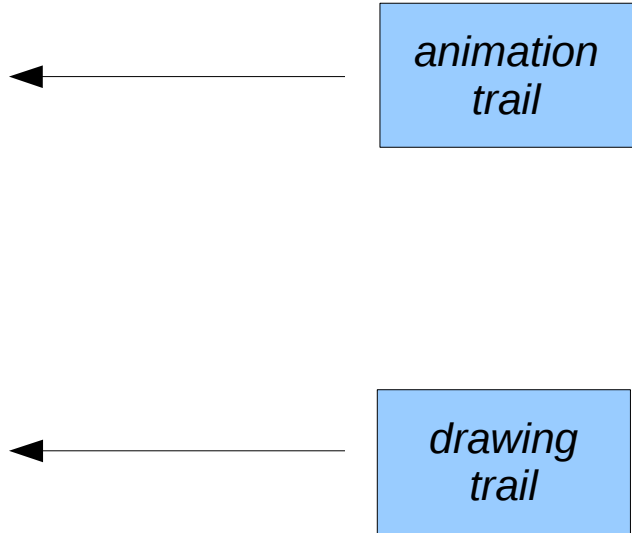
animation trail

drawing trail

```
code/await Bird (var int y, var int speed) -> FOREVER
do
  par do
    every UPDATE do
      if x >= WIDTH then
        break;
      end
      <animate-bird>
    end
  with
    every DRAW do
      <draw-bird>
    end
  end
end

pool[] Bird birds;

every 1s do
  spawn Bird(...) in birds;
end
```

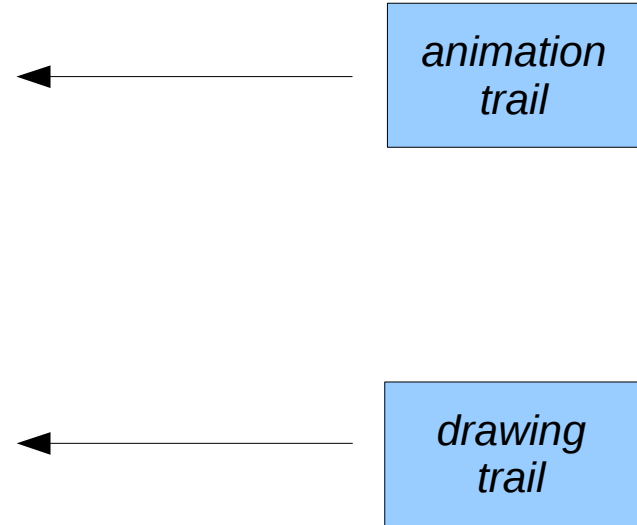


The diagram illustrates the execution flow of the Bird code. A blue box labeled "animation trail" has an arrow pointing to the `break;` line within the `every UPDATE do` block. Another blue box labeled "drawing trail" has an arrow pointing to the `<draw-bird>` line within the `every DRAW do` block.

```
code/await Bird (var int y, var int speed) -> FOREVER
do
  par do
    every UPDATE do
      if x >= WIDTH then
        break;
      end
      <animate-bird>
    end
  with
    every DRAW do
      <draw-bird>
    end
  end
end

pool[5] Bird birds;

every 1s do
  spawn Bird(...) in birds;
end
```



The diagram illustrates the execution flow of the provided code. It features two blue rectangular boxes on the right side. The top box is labeled "animation trail" and has an arrow pointing to the red "break;" line within the "if x >= WIDTH then" block. The bottom box is labeled "drawing trail" and has an arrow pointing to the "<draw-bird>" line within the "every DRAW do" block.

```
code/await Bird (var int y, var int speed) -> void  
do
```

```
  par/or do
```

```
    every UPDATE do  
      <animate-bird/break>  
    end
```

animation
trail

```
  with
```

```
    every DRAW do  
      <draw-bird>  
    end
```

drawing
trail

```
  end
```

the whole body terminates
(natural termination)

```
end
```

```
pool[5] Bird birds;
```

```
every 1s do
```

```
  spawn Bird(...) in birds;
```

```
end;
```

natural termination automatically
reclaims dynamic code/await
instances from memory pools

no need for a
free primitive or
garbage collection

spawned code/await
is anonymous

```
code/await Bird (var int y, var int speed) -> void  
do
```

```
<...>
```

```
end
```

pools also have
lexical scope

pool out of scope:
data and execution body of
all alive `code/await`
instances are reclaimed

```
loop do
```

```
  par/or do
```

```
    pool Birds[] birds;  
    every 1s do  
      spawn Bird(...) in birds;  
    end
```

```
  with
```

```
    await MOUSE_CLICK;
```

```
  end
```

```
end
```

heap allocation with
lexical memory management
(vs. garbage collection)

```
code/await Bird (var int y, var int speed) -> void
do
    <...>
end

loop do
    watching MOUSE_CLICK do
        pool Birds[] birds;
        every 1s do
            spawn Bird(...) in birds;
        end
    end
end
```

```
loop do
```

```
  watching MOUSE_CLICK do
```

```
    pool Birds[] birds;
```

```
    par do
```

```
      every 1s do
```

```
        spawn Bird(...) in birds;
```

```
      end
```

```
    with
```

```
      every UPDATE do
```

```
        var& Rect rct1;
```

```
        event& void coll1;
```

```
        loop (rct1,coll1) in birds do
```

```
          var& Rect rct2;
```

```
          event& void coll2;
```

```
          loop (rct2,col2) in birds do
```

```
            if (&rct1<&rct2) and Int(&rct1,&rct2) then
```

```
              emit coll1;
```

```
              emit coll2;
```

```
              break;
```

```
            end
```

```
          end
```

```
        end
```

```
      end
```

```
    end
```

```
  end
```

```
end
```

[birds-10]
(collision/iterator)

temporary alias
(valid only inside the loop)

iterate over the birds

access to internal state
of a code/await in the pool

notifies both code/await
about the collision

```
code/await Bird (var int y, var int speed)
    -> (var& Rect rct, event& void collide)
    -> FOREVER
```

do

```
var Rect my_rct = val Rect(20,y, 50,45);
rct = &my_rct;
```

```
event void my_collide;
collide = &my_collide;
```

a code/await may expose
and share its internal state

```
par/or do
    watching my_collide do
        every UPDATE do
            <animate/break> // accesses my_rct
        end
    end
    with
        every DRAW do
            <draw-bird> // accesses my_rct
        end
    end
end
```

end


```
code/await Bird (...) -> (...) -> void
do
  <...>
  par/or do
    watching my_collide do
      every UPDATE do
        <animate-bird/break>
      end
    end
    every UPDATE do
      <animate-bird-y>
      if my_rct.y >= HEIGHT then
        break;
      end
    end
  end
  with
    every DRAW do
      <draw-bird>
    end
  end
end
end
```

```

code/await Bird (...) -> (...) -> void
do
  <...>
  event bool hide;
  par/or do
    watching my_collide do
      every UPDATE do
        <animate-bird/break>
      end
    end
    every UPDATE do
      <animate-bird-fall>
    end
    watching ls do
      loop do
        emit hide(true);
        await 100ms;
        emit hide(false);
        await 100ms;
      end
    end
  with
    pause/if hide do
      every DRAW do
        <draw-bird>
      end
    end
  end
end
end

```

[birds-12]
(blink / pause/if)

```
event bool freeze;
```

```
par do
```

```
  pause/if freeze do
```

```
    <...> // bird pool, creation and collision
```

```
  end
```

```
with
```

```
  loop do
```

```
    var Key key1 = await KEY_PRESS until key1.sym=='p';  
    emit freeze(true);
```

```
  par/or do
```

```
    every DRAW do
```

```
      <draw-pause>
```

```
    end
```

```
  with
```

```
    var Key key2 = await KEY_PRESS until key2.sym=='p';
```

```
  end
```

```
  emit freeze(false);
```

```
end
```

```
end
```

```
pool Birds[] birds;  
par do  
  every 1s do  
    spawn Bird(...) in birds;  
  end  
with  
  every UPDATE do  
    <collision-detection>  
  end  
with  
  loop do  
    var Mouse mse = await MOUSE_CLICK;  
    var&? Rect found = do  
      var&? Rect rct;  
      loop (rct,_) in birds do  
        if Rect_vs_Mouse(&rct,mse) then  
          escape &rct;  
        end  
      end  
    end;  
    watching found do  
      every DRAW do  
        <draw-line>  
      end  
    end  
  end  
end
```

[birds-14]
(line)

checks if a bird
was clicked

watches the bird
while drawing
the line

Summary

- Structured Programming ->
Structured Synchronous Reactive Programming
 - sequences, loops, conditionals
 - await + parallels + abortion
 - procedures/abstractions
 - code/await
 - static/dynamic abstractions
 - lexical scope

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