Homework 3

Please submit this work by Wednesday, April 23 at 11:59pm.

Problem 1: The Programming Language L2

Terms in the language L2 are as follows:

This is an enrichment of last week's language L with boolean constants, some new operations and a new IF form for conditional evaluation. As in L, whitespace is allowed between L2 tokens and the lexical scanner must tolerate it.

Your task this week is to implement L2, and furthermore to implement an optimization that removes unused variables as a separate pass during compilation.

Construct your L2 implementation by copying your L source code into a new directory hw3/L2. (Be careful not to copy *everything* from hw2/L into this directory, because you don't want to copy all the hidden syn files.)

Implementation of L2 will consist largely of making straightforward modifications to your implementation of L. You will need new tokens, an enhanced lexer and an enhanced parser, new AST forms, new ScopedAST forms. Furthermore, since evaluation of an L2 program can yield an integer or a boolean, your evaluator needs a new type. Whereas in L the function Eval.eval had typescopedAST.exp -> value where value was a synonym for int, in L2 the same function has the type scopedAST.exp -> value.value where the module value is as follows:

Furthermore, unlike L programs, evaluation of scannable, parseable, scopeable L2 programs is no longer guaranteed to proceed without error, because there exists the possibility of nonsensical programs like +(8,T) or IF(12,13,14). The complete set of nonsensical L2 programs is as follows:

logical negation of a number,

- testing whether a boolean constant is zero,
- performing a less-than comparison where one or both arguments are booleans,
- "adding" a boolean, or
- having a number in the test position in a conditional.

Each of these constitutes a new kind of error, aruntime error. Therefore we add RuntimeError to the errors we associated with the previous

phases:syntaxError, ParseError and UnboundVariable. (It is worth making clear that we do not loosely equate numbers and booleans as C does.)

For these reasons, the EVAL signature for L2 needs to look like this:

```
signature EVAL = sig
  exception RuntimeError of string
  val eval : ScopedAST.exp -> Value.value
end
```

Please do not impose a type discipline on L2; the language specifically lacks a static type system. RaiseruntimeErrors as they arise during evaluation.

Besides implementing L2 with its enriched set of values and operations, implement a useless-variable-elimination optimization on scoped ASTs that adheres to the following signature:

```
signature OPTIMIZE = sig
  val optimize : ScopedAST.exp -> ScopedAST.exp
end
```

The optimization should transform programs such that any variable that is bound but not used should be eliminated. For example, the program "LET(x<-1,0)" should be transformed simply to "0", and "LET(x<-1,LET(y<-2,+(x,1)))" to "LET(x<-1,+(x,1))". (Of course, all useless variables can be "eliminated" simply by evaluating the program, but here we are providing experience with program transformation and compiler structure.)

As such, the Interpret structure can be rewritten as follows:

In summary, problem 1 consists of doing the following:

- copying your hw2/L source code (but not the hidden svn baggage) into a new directoryhw3/L2,
- including the value structure defined above,

- modifying the EVAL signature as above,
- hacking on the lexer, parser, scoper and evaluator to support L2, raising RuntimeErrorsfor
 programs that are ultimately nonsensical but otherwise fine (with respect to the previous
 compilation phases),
- adding the OPTIMIZE signature as above and implementing an Optimize structure to match it,
- modifying Interpret to include the optimization pass.

You are encouraged to fix and/or improve your L source code as part of your work on L2.

Problem 2

TaPL p. 24 gives a grammar for a simple term language, not altogether unlike L2. As presented in lecture last week, the grammar generates terms that can't be evaluated, like <code>iszero false</code> (an unevaluatable term). In a programming language implementation, the impasses in evaluating terms like <code>iszero false</code> correspond to runtime errors.

Rewrite the grammar on p. 24 such that it produces only terms that will never lead to a nonsensical state (runtime error) at any point in their evaluation, even taking into consideration that for every conditional, either branch may be evaluated. For example, the grammar should generate succ 0 but not succ false, and all terms if true then ... but no termsif 0 then We expect only a simple plaintext presentation.